Lecture 4: Integer arithmetic

September 9, 2019

601.229 Computer System Fundamentals



Integer arithmetic

- ► Integer representations based on fixed-size machine words are finite
- ▶ I.e., only a finite number of possible values can be represented
 - ightharpoonup For word with w bits, can represent 2^w possible values
- ➤ So, we should expect some (potentially) strange results when doing arithmetic using machine words
- ► These strange results can lead to surprising program behavior, including security vulnerabilities

Addition of unsigned values

Addition of unsigned values

- Same idea as what you learned in grade school
 - Start with least significant digit
 - As needed, carry excess into next-most-significant digit

```
\begin{array}{c} & 0 \\ 0110 \\ + & 0111 \end{array}
```

Example:
$$0110 + 0111$$

$$\begin{array}{c|c}
 & 100 \\
 & 0110 \\
 + & 0111 \\
\hline
 & 01 & carry 1
\end{array}$$

$$\begin{array}{c|c} & \textbf{1}000 \\ & 0\textbf{1}10 \\ + & 0\textbf{1}11 \\ \hline & \textbf{1}01 & \mathsf{carry} \ 1 \end{array}$$

$$egin{array}{cccc} {f 0}1000 & & & & \\ {f 0}110 & & & & \\ {f +} & {f 0}111 & & & \\ \hline {f 1}101 & {\it no carry} & & & \end{array}$$

$$\begin{array}{c} & 0110 \\ + & 0111 \\ \hline & 1101 & \text{done} \end{array}$$

Overflow

- ► If the sum of w-bit (unsigned) integer values is too large to represent using a w-bit word, overflow occurs
- ightharpoonup Effective sum of w bit integers a and b is

$$(a+b) \mod 2^w$$

```
\begin{array}{c} & 0 \\ 1110 \\ + & 0111 \end{array}
```

Example:
$$1110 + 0111$$

$$\begin{array}{c|cccc}
 & 100 \\
 & 1110 \\
 & & 0111 \\
\hline
 & 01 & carry 1
\end{array}$$

$$\begin{array}{c|c} & \mathbf{1}100 \\ & 1\mathbf{1}10 \\ + & 0\mathbf{1}11 \\ \hline & \mathbf{1}01 & \mathsf{carry} \ 1 \end{array}$$

$$\begin{array}{c|cccc}
 & \mathbf{1}1100 \\
 & \mathbf{1}110 \\
 & \mathbf{0}111 \\
\hline
 & \mathbf{0}101 & \mathsf{carry} \ 1
\end{array}$$

Example: 1110 + 0111

$$\begin{array}{r}
 11100 \\
 1110 \\
 + 0111 \\
\hline
 10101
\end{array}$$

True sum is 10101 (21), effective sum is 101 (5) (note 21 mod 16 = 5)

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Addition of signed values

Useful property of two's complement: addition is carried out exactly the same way for signed values as for unsigned values

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Useful property of two's complement: addition is carried out exactly the same way for signed values as for unsigned values

Signed addition example

Example:
$$0101 (5) + 1110 (-2)$$

Signed addition example

Example:
$$0101 (5) + 1110 (-2)$$

$$\begin{array}{r} & 0101 \\ + & 1110 \\ \hline & 10011 \end{array}$$

After truncating (discarding high bit of sum), effective sum is 0011 (3)



Signed overflow

What happens when sum of signed w-bit values can't be represented?

- ▶ If sum exceeds $2^{w-1} 1$, it becomes negative (overflow)
- ▶ If sum is less than -2^{w-1} , it becomes positive (negative overflow)

Signed addition example (overflow)

Example:
$$0100 (4) + 0101 (5)$$

$$\begin{array}{c} 0100 \\ + 0101 \end{array}$$

Signed addition example (overflow)

Example:
$$0100 (4) + 0101 (5)$$

$$\begin{array}{r} & 0100 \\ + & 0101 \\ \hline & 1001 \end{array}$$

Result is
$$-7 (-8 + 1)$$

Signed addition example (negative overflow)

Signed addition example (negative overflow)

Result (after truncating) is 7

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Two's complement negation and subtraction

- Negation: if x is a two-complement integer value, -x can be computed by inverting bits of x, then adding 1
 - ► Why?
- Subtraction:

$$a - b = a + -b$$

I.e., to compute a - b, compute -b, then add -b to a



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Integer arithmetic in C

Yeah