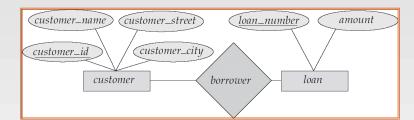
Chapter 6: Entity-Relationship Model

- A database can be modeled as:
 - a collection of entities,
 - relationship among entities.
- An entity is an object that exists and is distinguishable from other objects.
 - Example: specific person, company, event, plant
- Entities have attributes
 - Example: people have names and addresses
- An **entity set** is a set of entities of the same type that share the same properties.
 - Example: set of all persons, companies, trees, holidays

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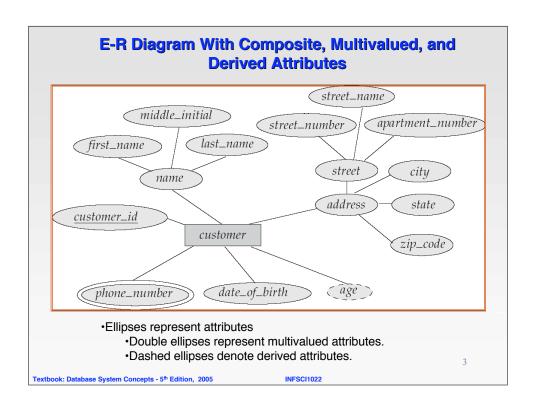
E-R Diagrams

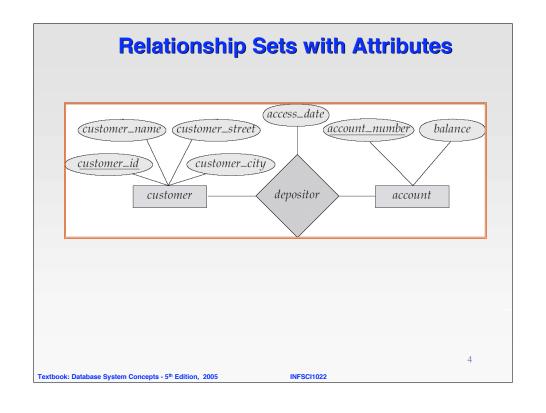


- Rectangles represent entity sets.
- Diamonds represent relationship sets.
- Lines link attributes to entity sets and entity sets to relationship sets.
- Underline indicates primary key attributes

2

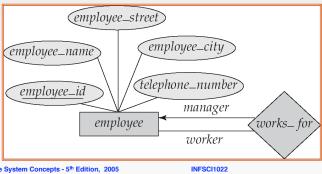
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Roles

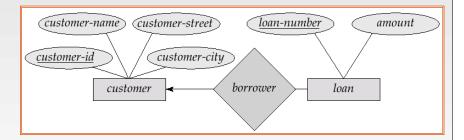
- Entity sets of a relationship need not be distinct
- The labels "manager" and "worker" are called **roles**; they specify how employee entities interact via the works_for relationship set.
- Roles are indicated in E-R diagrams by labeling the lines that connect diamonds to rectangles.
- Role labels are optional, and are used to clarify semantics of the relationship



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One-To-Many Relationship

■ In the one-to-many relationship a loan is associated with at most one customer via borrower, a customer is associated with several (including 0) loans via borrower

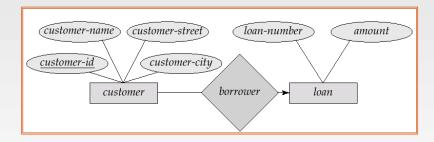


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Many-To-One Relationships

In a many-to-one relationship a loan is associated with several (including 0) customers via borrower, a customer is associated with at most one loan via borrower

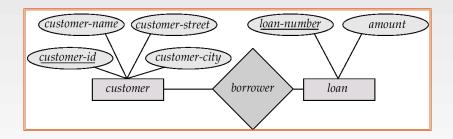


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Many-To-Many Relationship

- A customer is associated with several (possibly 0) loans via borrower
- A loan is associated with several (possibly 0) customers via borrower

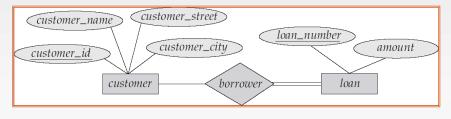


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Participation of an Entity Set in a Relationship Set

- Total participation (indicated by double line): every entity in the entity set participates in at least one relationship in the relationship set
 - participation of loan in borrower is total
 - > every loan must have a customer associated to it via borrower
- Partial participation: some entities may not participate in any relationship in the relationship set
 - participation of customer in borrower is partial

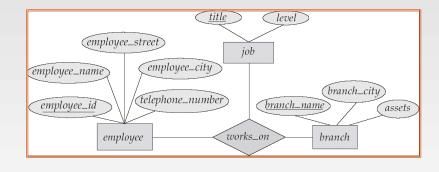


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E-R Diagram with a Ternary Relationship

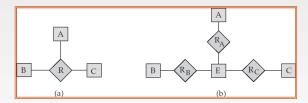


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Converting Non-Binary Relationships to Binary Form

- In general, any non-binary relationship can be represented using binary relationships by creating an artificial entity set.
 - Replace *R* between entity sets A, B and C by an entity set *E*, and three relationship sets:
 - 1. R_A , relating E and A
- $2.R_B$, relating E and B
- 3. R_C , relating E and C



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Weak Entity Sets

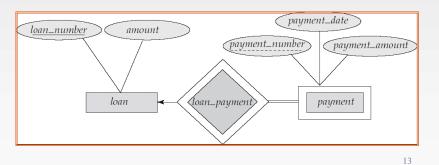
- An entity set that does not have a primary key is referred to as a weak entity set.
- The existence of a weak entity set depends on the existence of a identifying entity set
 - it must relate to the identifying entity set via a total, one-to-many relationship set from the identifying to the weak entity set
 - Identifying relationship depicted using a double diamond
- The **discriminator** (*or partial key*) of a weak entity set is the set of attributes that distinguishes among all the entities of a weak entity set.
- The primary key of a weak entity set is formed by the primary key of the strong entity set on which the weak entity set is existence dependent, plus the weak entity set's discriminator.

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Weak Entity Sets (Cont.)

- We depict a weak entity set by double rectangles.
- We underline the discriminator of a weak entity set with a dashed line.
- payment_number discriminator of the *payment* entity set
- Primary key for *payment* (*loan_number*, *payment_number*)



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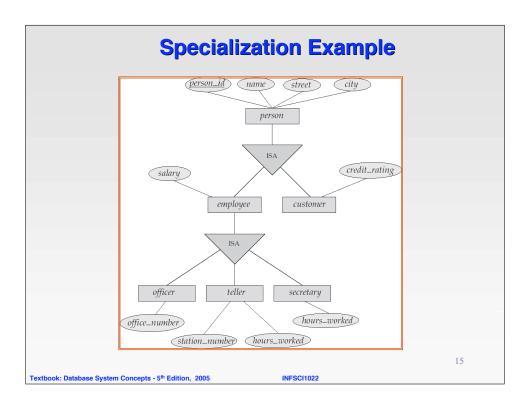
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Extended E-R Features: Specialization

- Top-down design process; we designate subgroupings within an entity set that are distinctive from other entities in the set.
- These subgroupings become lower-level entity sets that have attributes or participate in relationships that do not apply to the higher-level entity set.
- Depicted by a triangle component labeled ISA (E.g. customer "is a" person).
- Attribute inheritance a lower-level entity set inherits all the attributes and relationship participation of the higher-level entity set to which it is linked.

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Extended ER Features: Generalization

- A bottom-up design process combine a number of entity sets that share the same features into a higher-level entity set.
- Specialization and generalization are simple inversions of each other; they are represented in an E-R diagram in the same way.
- The terms specialization and generalization are used interchangeably.

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Design Constraints on a Specialization/Generalization

- Constraint on which entities can be members of a given lower-level entity set.
 - condition-defined
 - Example: all customers over 65 years are members of senior-citizen entity set; senior-citizen ISA person.
 - user-defined
- Constraint on whether or not entities may belong to more than one lower-level entity set within a single generalization.
 - Disjoint
 - > an entity can belong to only one lower-level entity set
 - Noted in E-R diagram by writing disjoint next to the ISA triangle
 - Overlapping
 - > an entity can belong to more than one lower-level entity set

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Design Constraints on a Specialization/Generalization (Cont.)

- Completeness constraint -- specifies whether or not an entity in the higher-level entity set must belong to at least one of the lower-level entity sets within a generalization.
 - total : an entity must belong to one of the lower-level entity sets
 - partial: an entity need not belong to one of the lower-level entity sets

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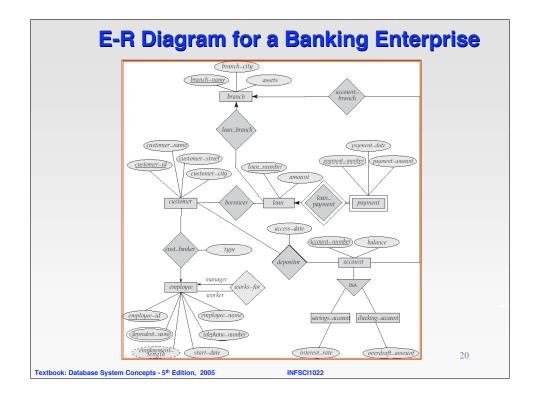
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E-R Design Decisions

- The use of an attribute or entity set to represent an object.
- Whether a real-world concept is best expressed by an entity set or a relationship set.
- The use of a strong or weak entity set.
- The use of specialization/generalization contributes to modularity in the design.

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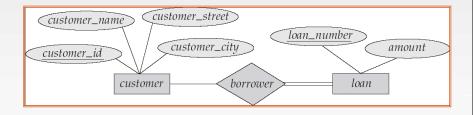
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Representing Entity Sets as Relation Schemas

- A strong entity set reduces to a schema with the same attributes.
- A weak entity set becomes a table that includes a column for the primary key of the identifying strong entity set

loan = (loan_number, amount)

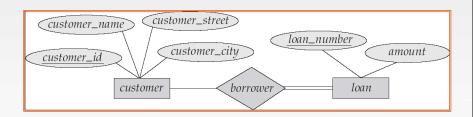


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Representing Relationship Sets as Relation Schemas

- A many-to-many relationship set is represented as a schema with attributes for the primary keys of the two participating entity sets, and any descriptive attributes of the relationship set.
- Example: schema for relationship set borrower borrower = (customer_id, loan_number)



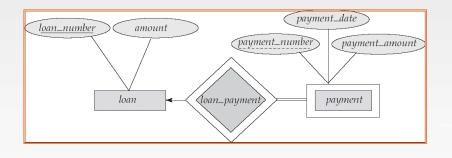
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Representing Weak Entity Sets as Relational Schemas

payment =

(<u>loan_number</u>, <u>payment_number</u>, payment_date, payment_amount)

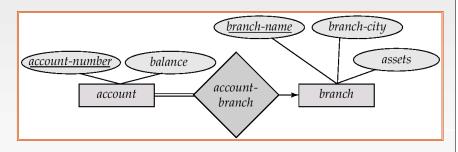


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Redundancy of Schemas

- Many-to-one and one-to-many relationship sets that are total on the many-side can be represented by adding an extra attribute to the "many" side, containing the primary key of the "one" side
- Example: Instead of creating a schema for relationship set account_branch, add an attribute branch_name to the schema arising from entity set account

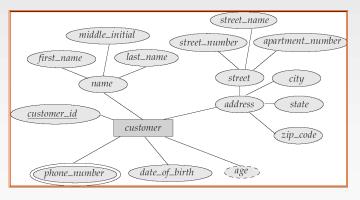


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Composite and Multivalued Attributes

- Composite attributes are flattened out by creating a separate attribute for each component attribute
- A multivalued attribute *M* of an entity *E* is represented by a separate schema *EM*
- Each value of the multivalued attribute maps to a separate tuple of the relation on schema EM



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Representing Specialization via Schemas

■ Method 1:

- Form a schema for the higher-level entity
- Form a schema for each lower-level entity set, include primary key of higher-level entity set and local attributes

schema	attributes	
person	name, street, city	
customer	name, credit_rating	
employee	name, salary	

 Drawback: getting information about, an employee requires accessing two relations, the one corresponding to the low-level schema and the one corresponding to the high-level schema

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Representing Specialization as Schemas (Cont.)

■ Method 2:

• Form a schema for each entity set with all local and inherited attributes

schema	attributes
person	name, street, city
customer	name, street, city, credit_rating
employee	name, street, city, salary

- If specialization is total, the schema for the generalized entity set (person) not required to store information
- Drawback: street and city may be stored redundantly for people who are both customers and employees

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