

# Divide-and-conquer algorithms

CS 146 - Spring 2017

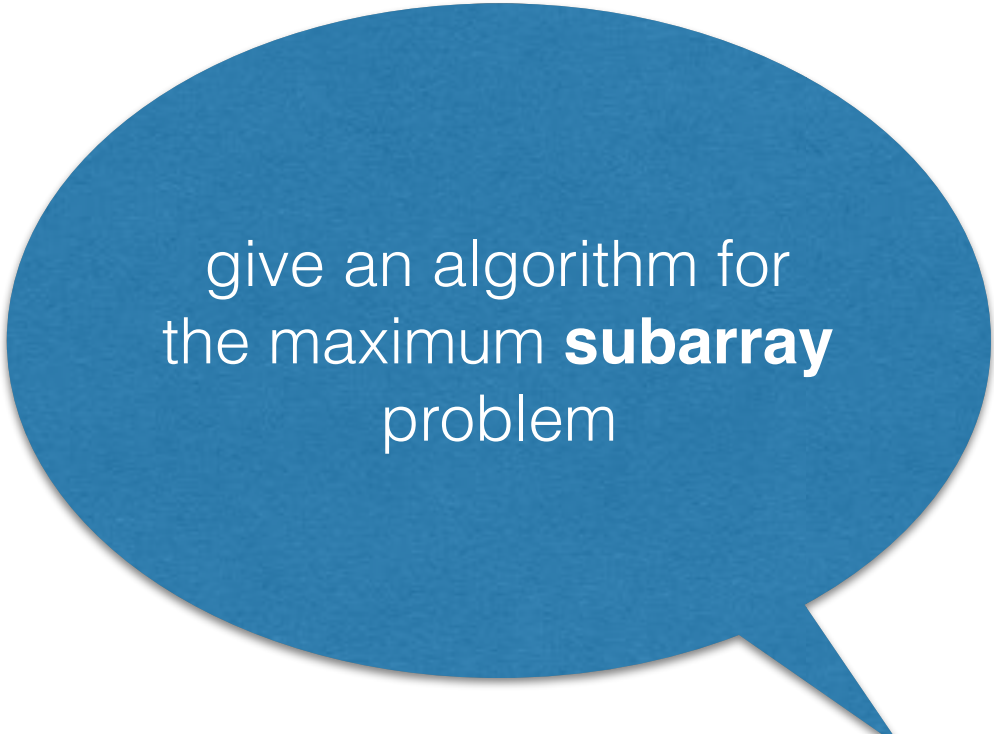
# Today

- Review: mergesort
- Anatomy of a divide-and-conquer algorithm
- Example: maximum subarray
- Example: Karatsuba's multiplication

# Question

what is the maximum sum of values  
among all **contiguous subsequences**  
of this sequence ?

13, -3, -25, 20, -3, -16, -23, 18, 20, -7, 12, -5, -22, 15, -4, 7



give an algorithm for  
the maximum **subarray**  
problem

# Recall: mergesort

initial sequence

5 2 4 6 1 3 2 6

split

5 2 4 6

1 3 2 6

split

split

5 2

4 6

1 3

2 6

split

split

split

split

5

2

4

6

1

3

2

6

sorted sequence

1 2 2 3 4 5 6 6

merge

2 4 5 6

1 2 3 6

merge

merge

2 5

4 6

1 3

2 6

merge

merge

merge

merge

5

2

4

6

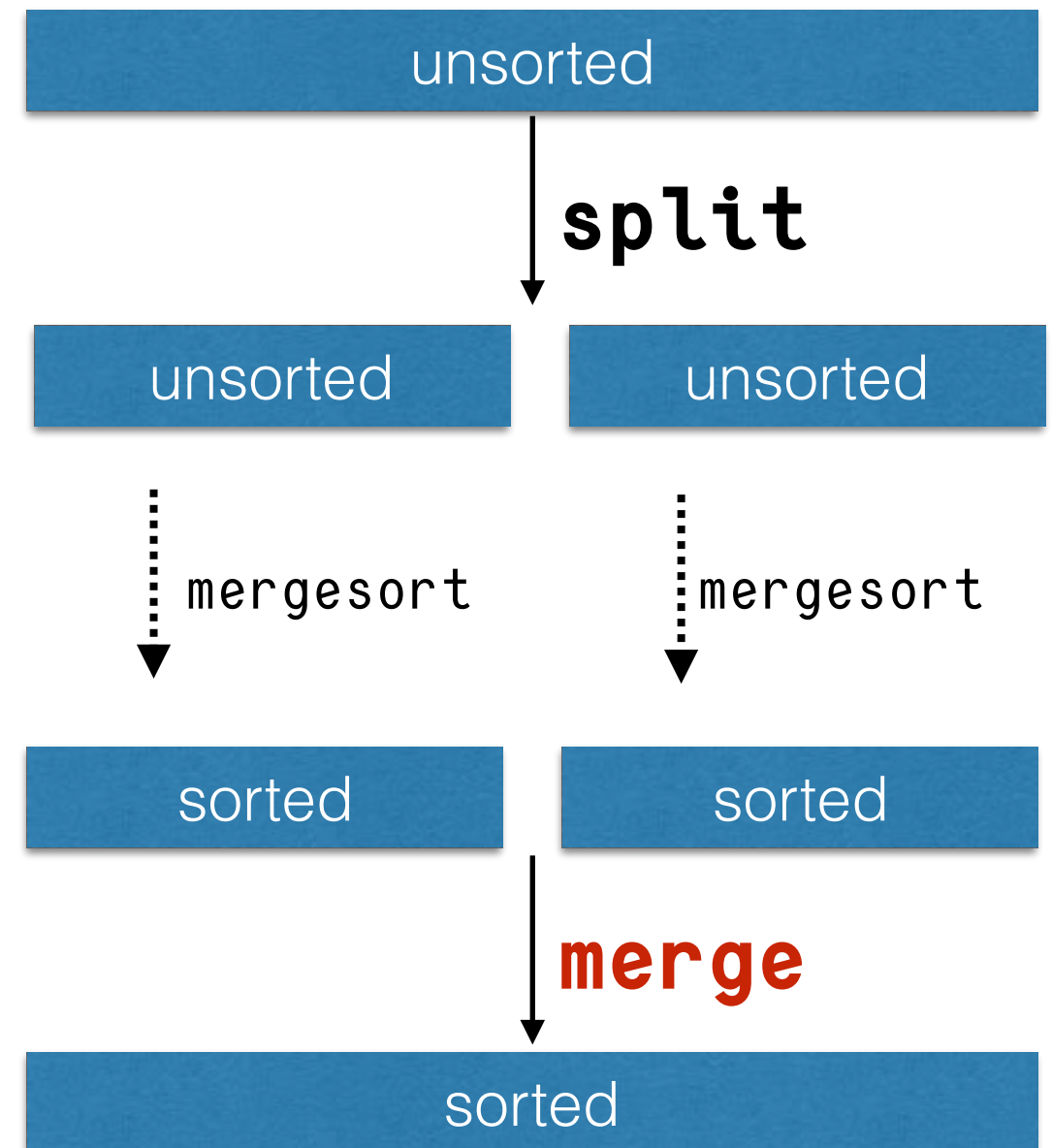
1

3

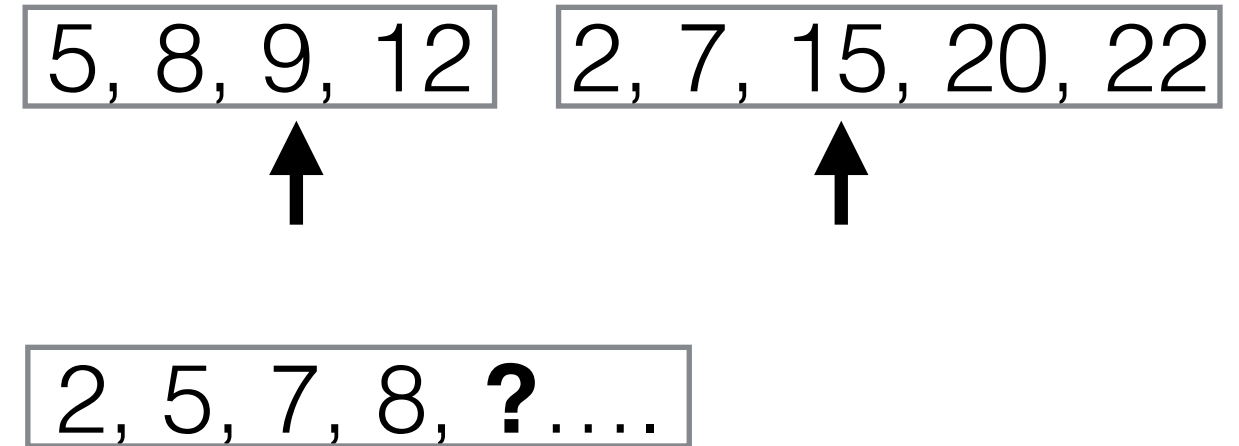
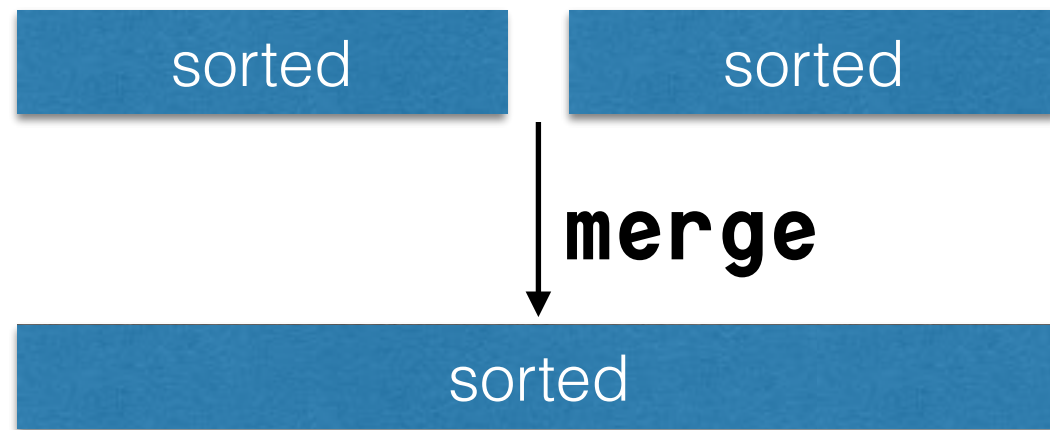
2

6

```
void mergesort(list) {  
    if (length(list) <= 1) return;  
    split list into left and  
    .           right sublists  
  
    mergesort(left);  
    mergesort(right);  
    merge left and right;  
}
```

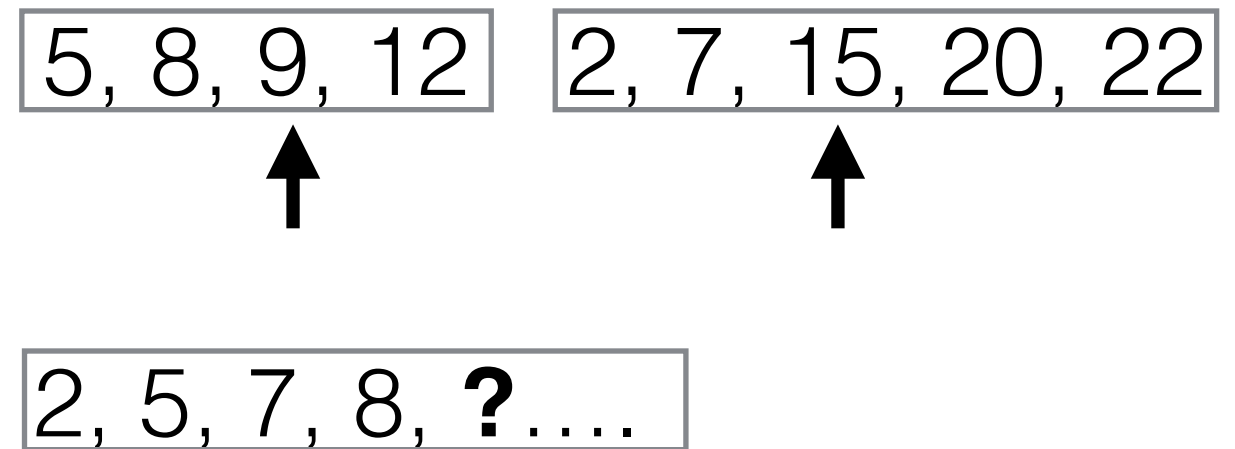
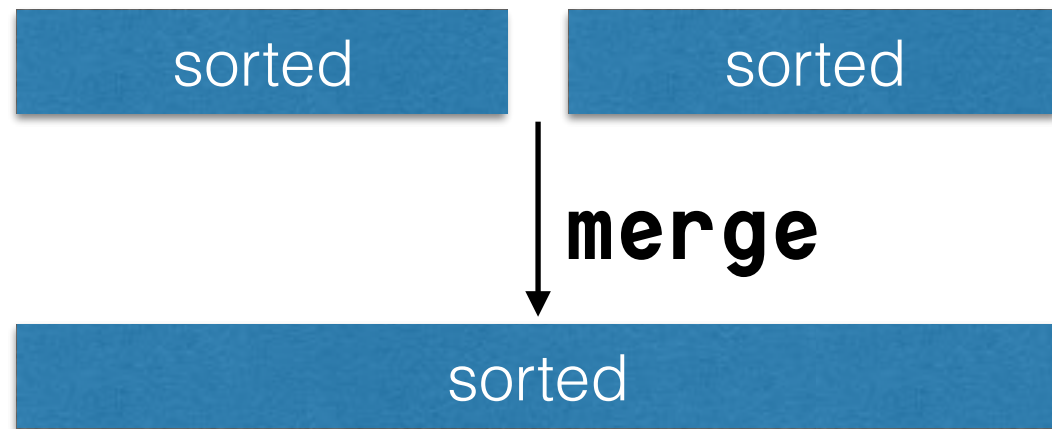


# Recall: merge step



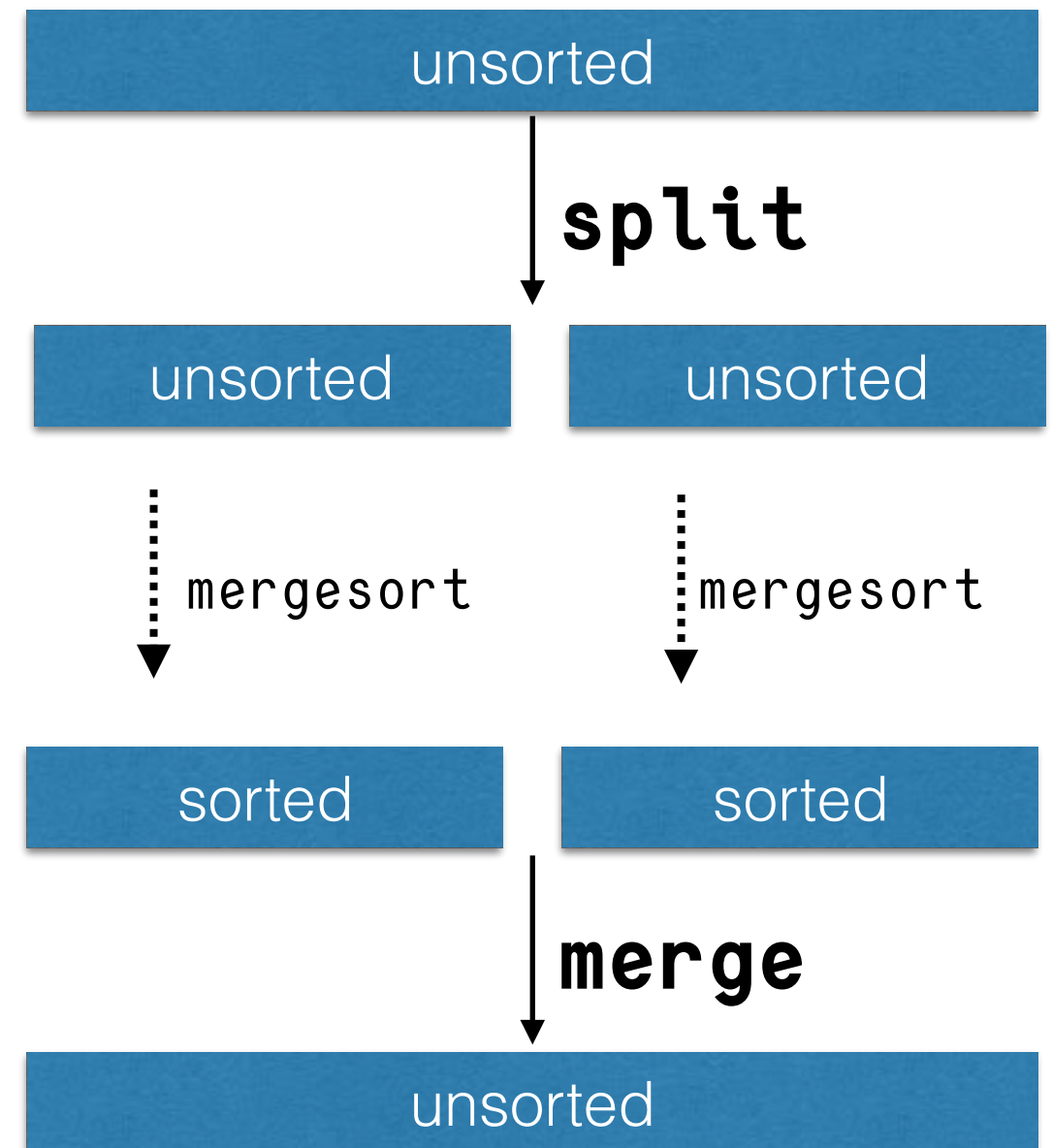
```
merge(a, b) {  
    initialize empty result list,      set pointers to beginning of the 2 lists  
    while there are 2 elements left to compare  
        pick the smaller of the 2 being pointed at  
        append it to the result list,   increment its pointer  
    append rest of remaining input list to result and return that  
}
```

# #comparisons in merge?



How do you analyze this??

```
void mergesort(list) {  
    if (length(list) <= 1) return;  
    split list into left and  
    right sublists  
  
    mergesort(left);  
    mergesort(right);  
    merge left and right;  
}
```

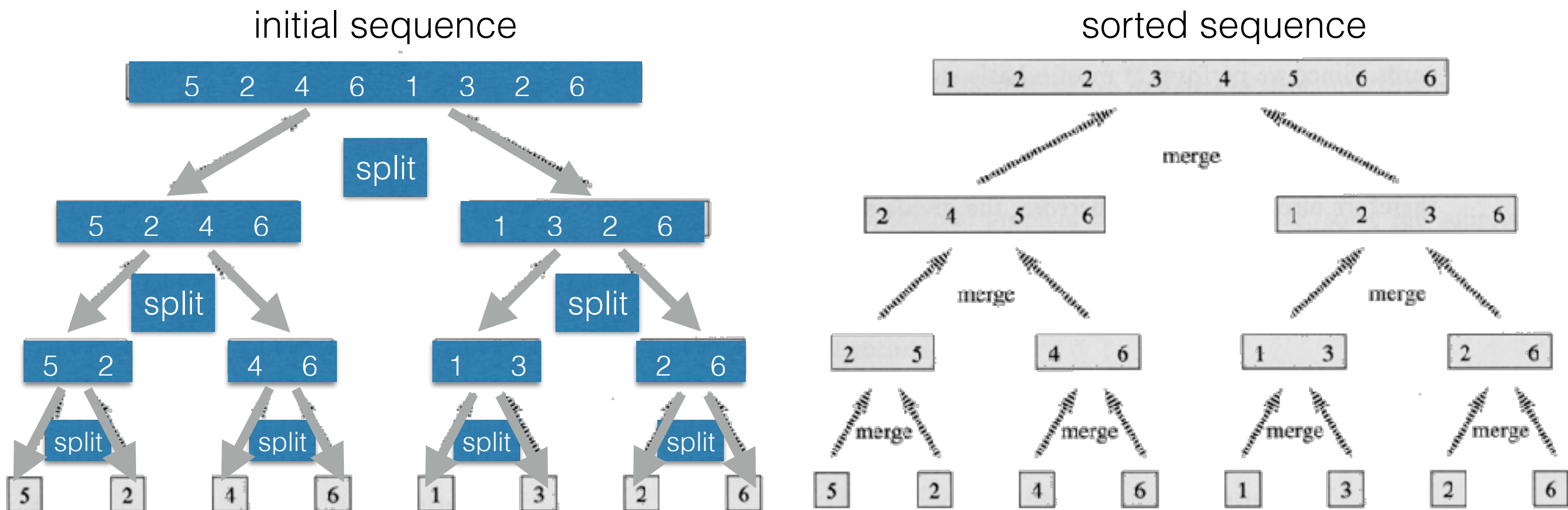


Recall the main idea: sum up work done in initial call AND all subsequent calls



Recall the main idea: sum up work done in initial call AND all subsequent calls

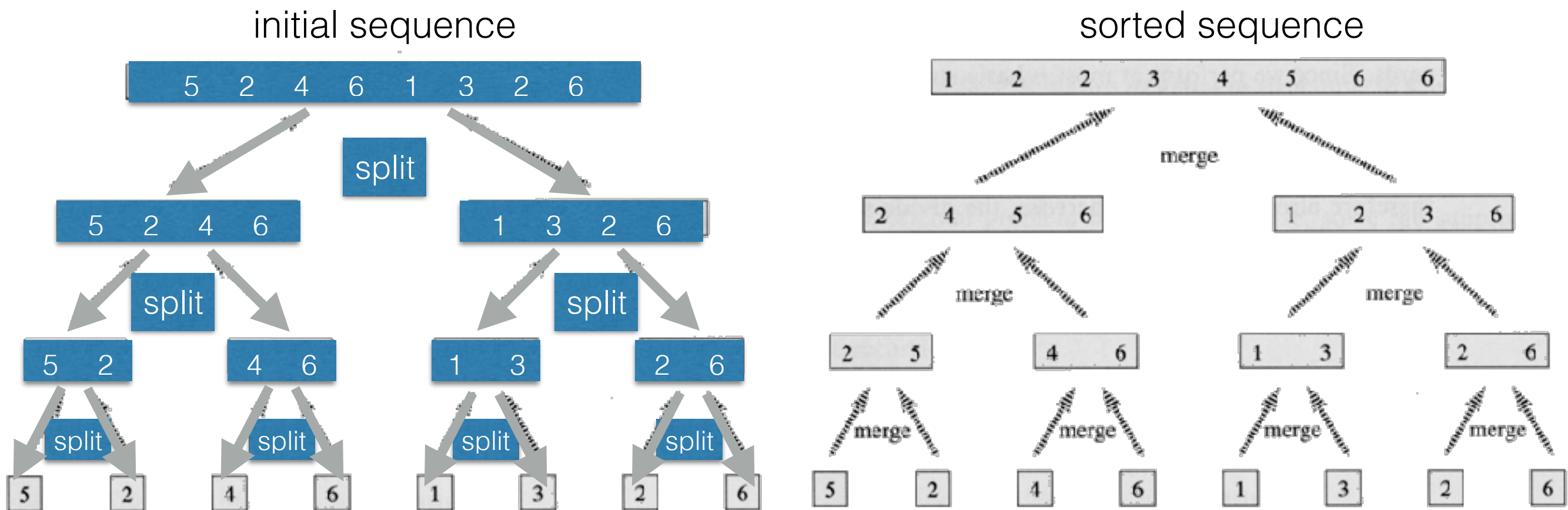
How do you analyze this??



Obs 1: each call does 1 split and 1 merge and merge dominates

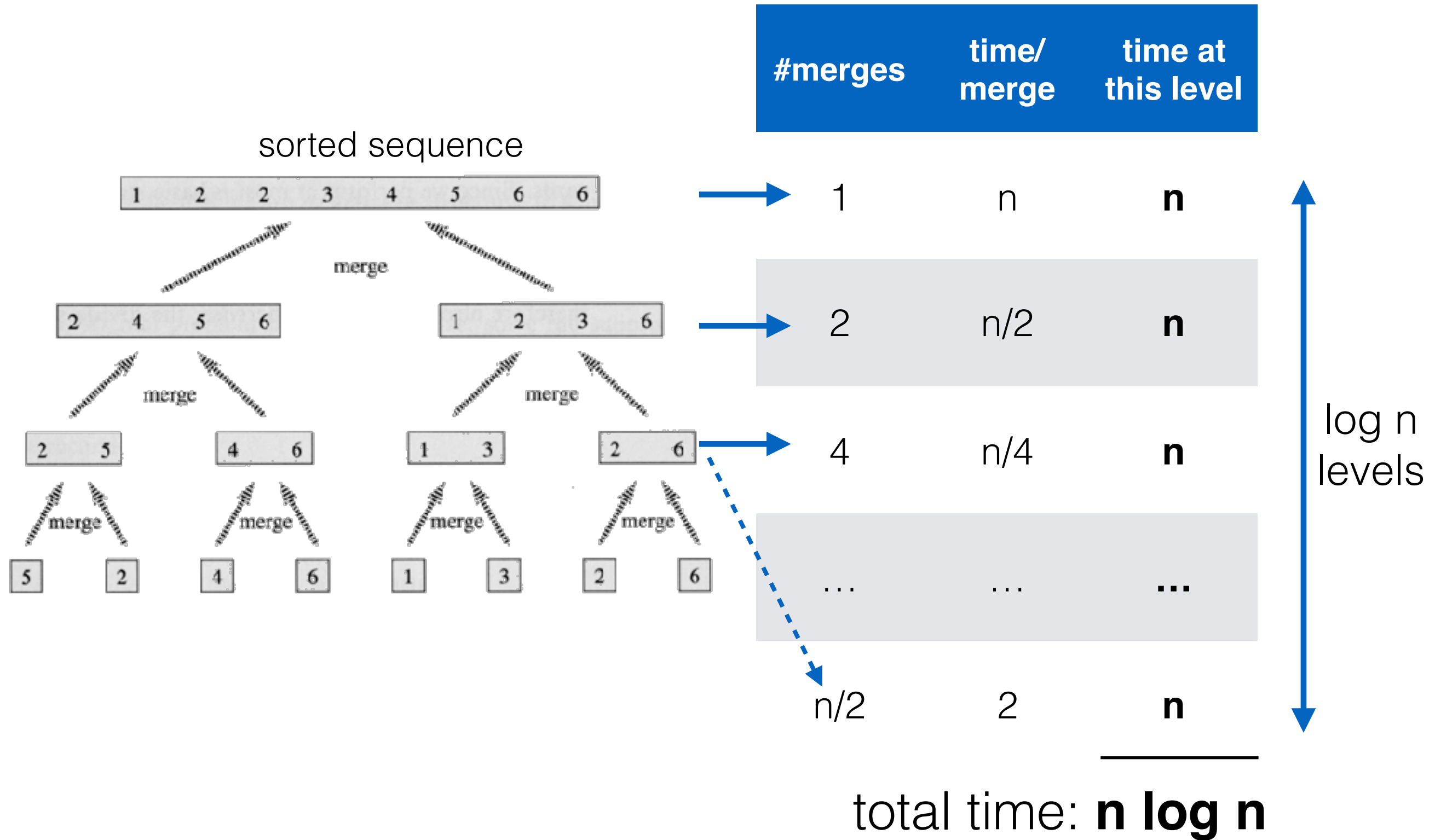
Recall the main idea: sum up work done in initial call AND all subsequent calls

How do you analyze this??



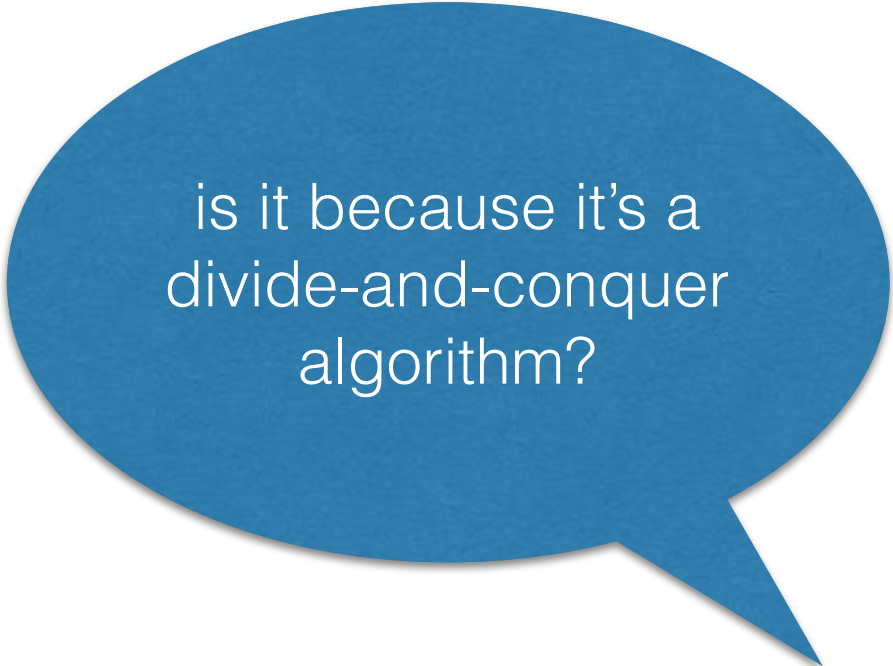
Obs 2: each call at the same level of the recursion tree does the same amount of work

# Mergesort analysis



# What makes mergesort fast?

- Insertion sort:  $O(n^2)$
- Selection sort:  $O(n^2)$
- Mergesort:  $O(n \log n)$



is it because it's a  
divide-and-conquer  
algorithm?

what is the maximum sum of values  
among all **contiguous subsequences**  
of this sequence ?

13, -3, -25, 20, -3, -16, -23, 18, 20, -7, 12, -5, -22, 15, -4, 7

divide-and-conquer idea?

13, -3, -25, 20, -3, -16, -23, 18, 20, -7, 12, -5, -22, 15, -4, 7

given the max sum in each half,  
can we get the best in the whole array?

13, -3, -25, 20, -3, -16, -23, 18, 20, -7, 12, -5, -22, 15, -4, 7

.....

the best contiguous subsequence can be

- entirely contained on the left of split
- entirely contained on the right of split
- cross the split

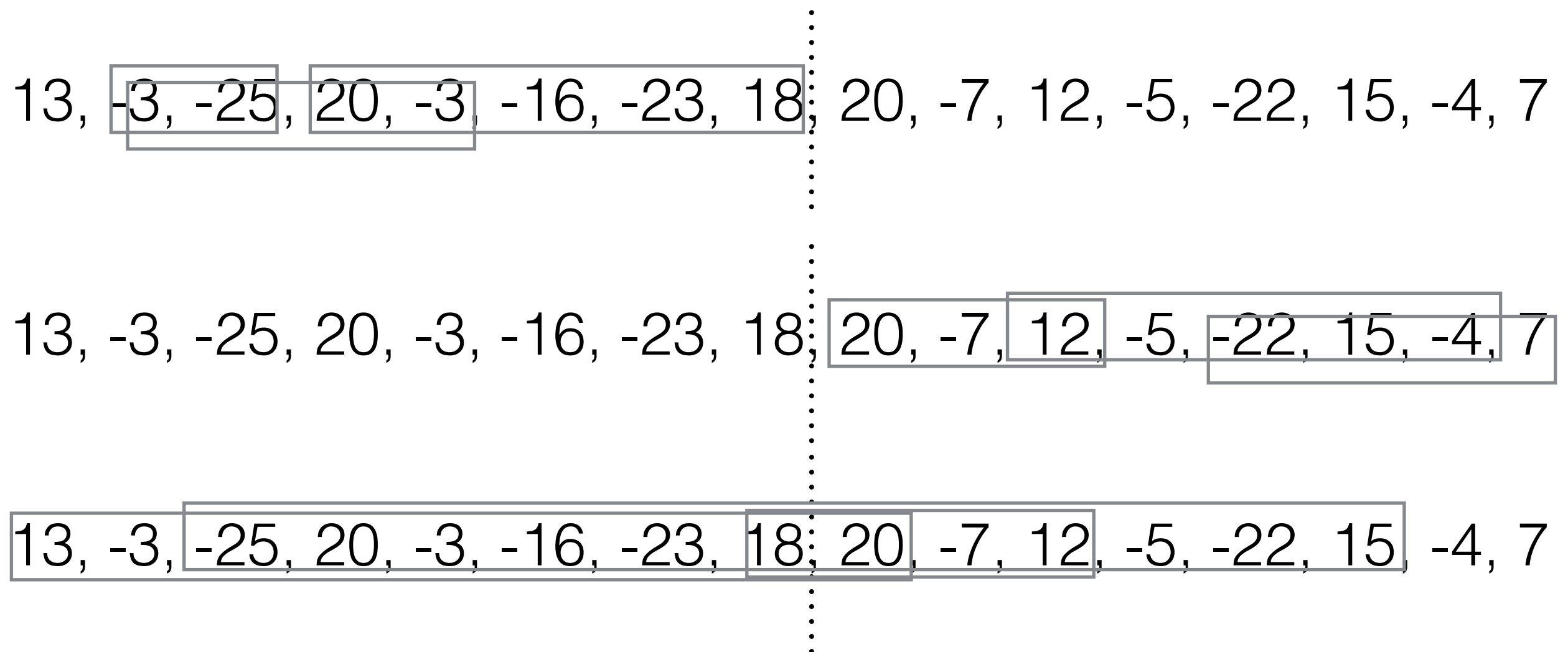
that's a job  
for recursion

let's do this  
ourselves

```

int maxSubarraySum(list) {
    leftMax = maxSubarraySum(left half of list)
    rightMax = maxSubarraySum(right half of list)
    crossMax = crossMax(list)
    return max(leftMax, rightMax, crossMax)
}

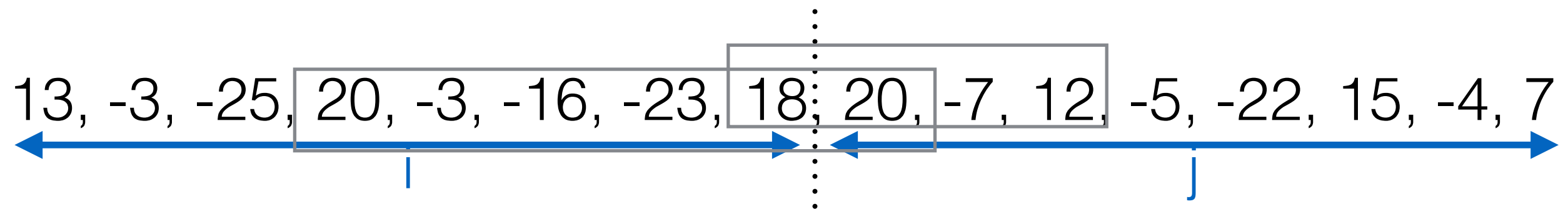
```



```

int maxSubarraySum(list) {
    leftMax = maxSubarraySum(left half of list)
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    return max(leftMax, rightMax, crossMax)
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```



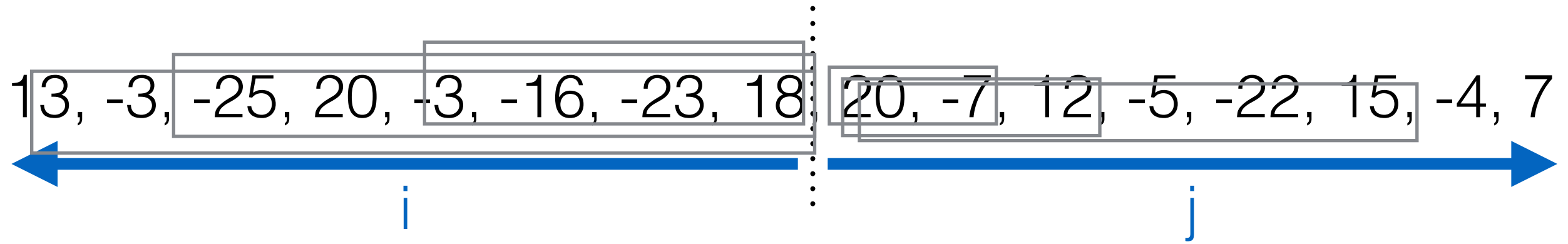
```

int crossMax(list) {
    int crossMax = -inf (or sum all neg vals in list)
    for (int i = 0; i < n/2-1; i++)
        for (int j = n/2+1; j < n; j++)
            sum = sum values in list from i to j
            crossMax = max(crossMax, sum)
    return crossMax
}

```

time?





```
int crossMax2(list) {  
    int rightCrossMax = -inf  
    int rightRunningSum = 0  
    for (i = n/2; i < n; i++) {  
        rightRunningSum += list(i)  
        rightCrossMax = max(rightCrossMax, rightRunningSum)  
    }  
    int leftCrossMax = ... (same idea, but going towards left)  
    return leftCrossMax + rightCrossMax;  
}
```

time?

# Time?

```
int maxSubarraySum(list) {  
    leftMax = maxSubarraySum(left half of list)  
    rightMax = maxSubarraySum(right half of list)  
    crossMax = crossMax2(list)  
    return max(leftMax, rightMax, crossMax)  
}
```



time?

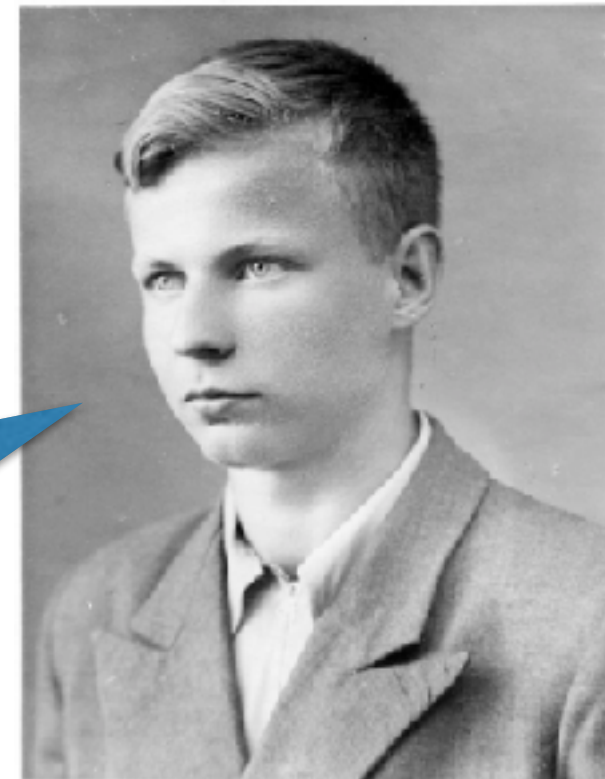
# The 1st divide-and-conquer algorithm



A. N. Kolmogorov

I conjecture that multiplication cannot be solved in faster than  $\Theta(n^2)$ .

Uh, ... Prof?  
I think you can...  
Will you check my work?



A. A. Karatsuba



Recall....

$$\begin{array}{r} \phantom{0}756 \\ \times \phantom{0}32 \\ \hline 1512 \\ 80 \end{array}$$

The grade-school multiplication algorithm is  $O(n^2)$ .

# Karatsuba fast multiplication

explained on the board

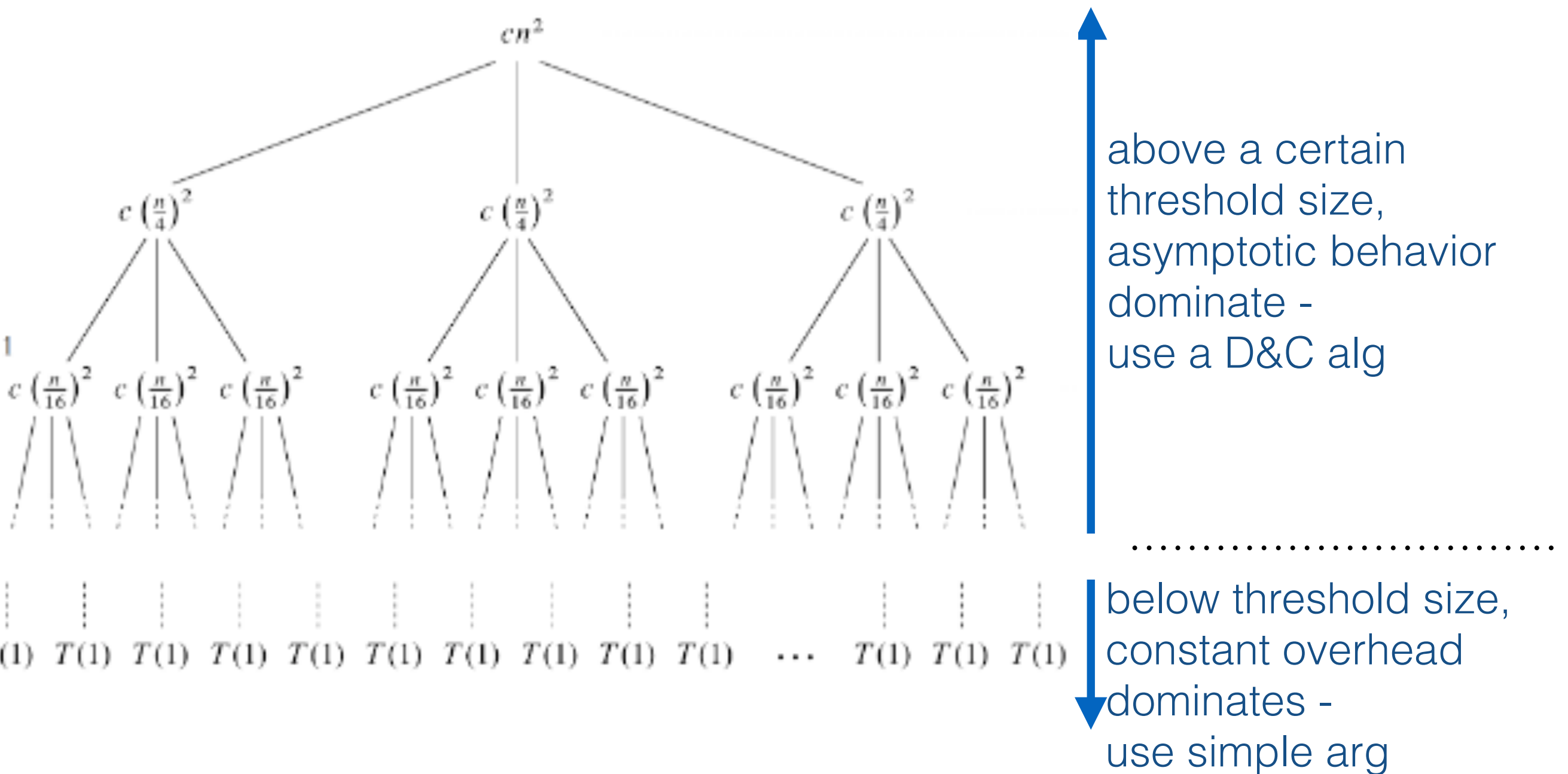
# Recap: divide-and-conquer algorithms

- anatomy of a divide-and-conquer algorithm
  - divide: split problem into smaller subproblems
  - recurse: solve each subproblem recursively
  - conquer: combine the results of the subproblems
- D&C not inherently more efficient. Need
  - evenly split subproblems
  - efficient divide and conquer steps

# Recap: divide-and-conquer algorithms

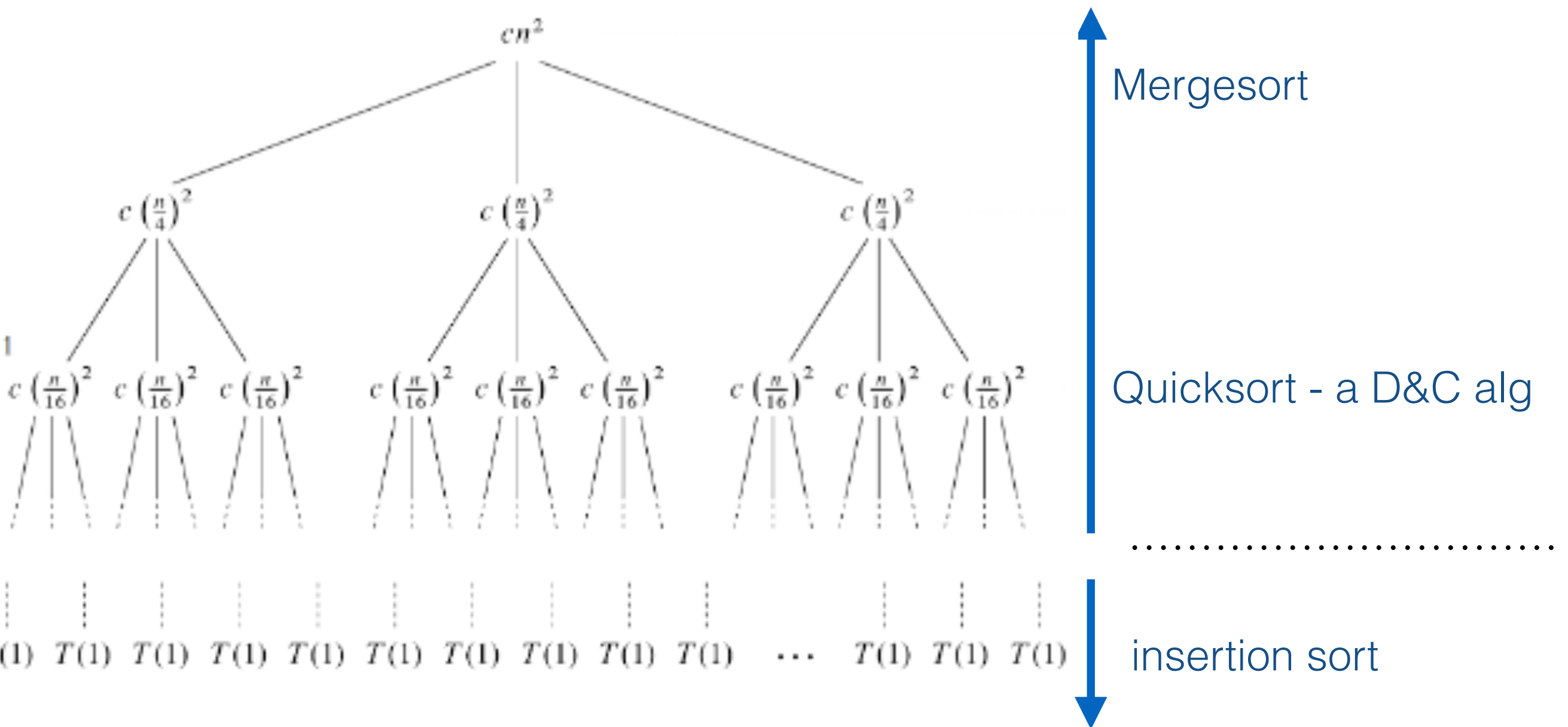
Problem	Non-D&C	D&C
sorting	insertion sort selection sort $O(n^2)$	mergesort quick sort $O(n \log n)$
max-subarray-sum	brute force $O(n^3)$	$O(n \log n)$
multiplication	grade school $O(n^2)$	Karatsuba $O(n^{\log_2 3})$

# Divide-and-conquer algorithms in practice



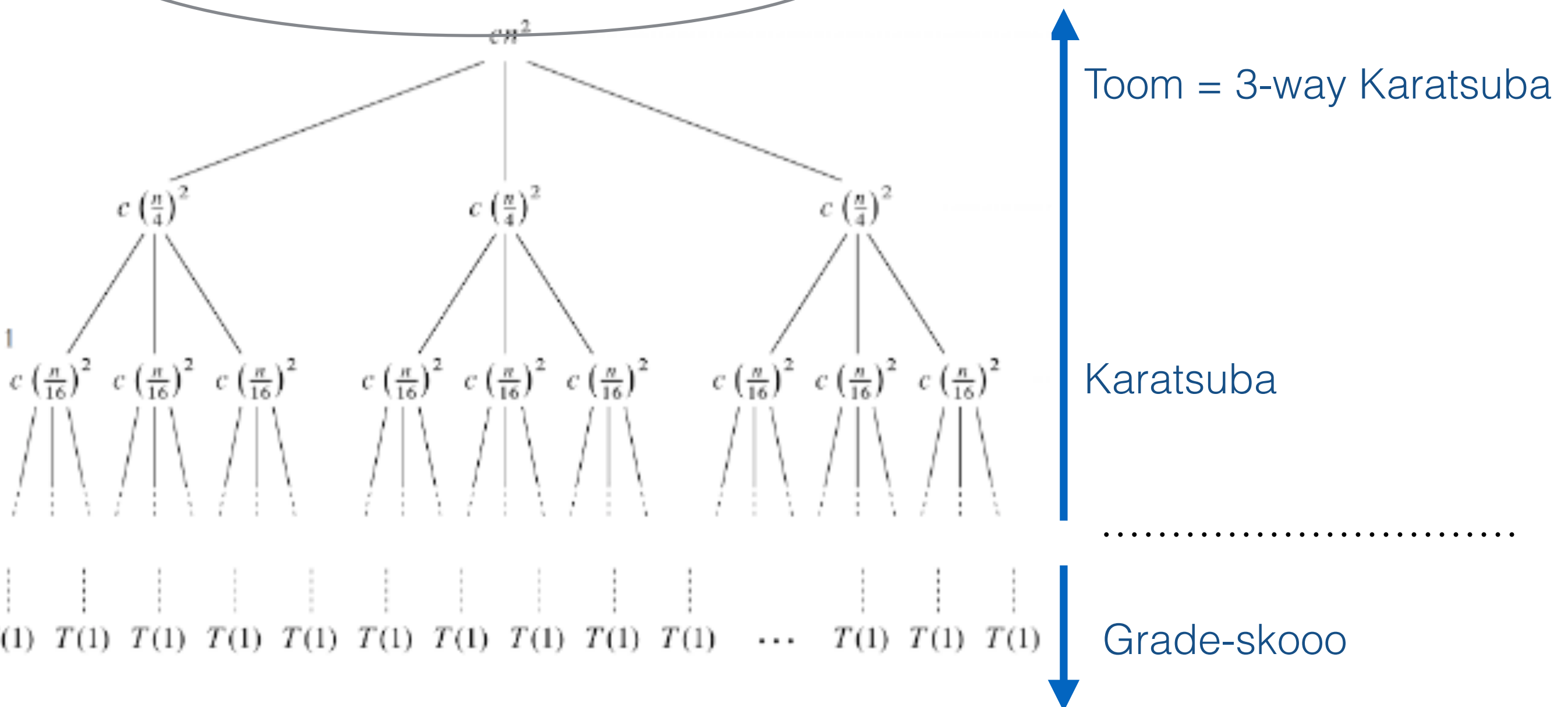


# Java implementation of Arrays.sort



# GMP library multiplication

GNU multi-precision  
(scientific computing)



# Next time...

- Why Karatsuba's running time is  $O(n^{\log_2 3})$
- How to analyze D&C algorithms
- Exam study questions, bring them...