

Künstliche Intelligenz - Übung 4

Julian Dobmann, Kai Kruschel

Aufgabe 2

"To prove logical consequence take the axioms and negated conjecture in CNF, and check for unsatisfiability."

Vermutung/Conjecture: I always avoid a kangaroo.

Negierte Vermutung:

```
cnf(avoid_kangaroo,negated_conjecture,  
    ( ~ avoided(the_kangaroo) ) ).
```

Von dieser (negierten) Vermutung ist die Unerfüllbarkeit zu zeigen.

DPLL-Algo

Quelle: <http://www.cs.miami.edu/~geoff/Courses/CSC648-12S/Content/DPLL.shtml>

Mit römischen Ziffern sind "Breakpoints" im Algorithmus markiert.

```

def dpll(S)
  # I
  if (S is empty)
    return "satisfiable"
  end

  # II
  while (there is some *unit clause* {p} or some pure literal p)
    # III
    if ({p} and {~p} are in S)
      return "unsatisfiable"
    else
      S = simplify(S,p)
    end
  end

  # IV
  if (S is empty)
    return "satisfiable"
  end

  # V
  # pick some literal q from a shortest clause in S
  q = pick(S)

  # VI
  if (dpll(simplify(S,q)) == "satisfiable")
    return "satisfiable"
  else
    return dpll(simplify(S,~q))
  end
end

def simplify(S, p)
  delete every clause in S containing p
  delete every occurrence in S of ~p
  return S
end

```

Anwendung des DPLL-Algorithmus: `dpll(S)` .

```

S = {
  ~ in_house(Cat) | cat(Cat),
  ~ gazer(Gazer) | suitable_pet(Gazer),
  ~ detested(Detested) | avoided(Detested),
  ~ carnivore(Carnivore) | prowler(Carnivore),

```

```

~ cat(Cat) | mouse_killer(Cat),
~ take_to_me(Taken_animal) | in_house(Taken_animal),
~ kangaroo(Kangaroo) | ~ suitable_pet(Kangaroo),
~ mouse_killer(Killer) | carnivore(Killer),
takes_to_me(Animal) | detested(Animal),
~ prowler(Prowler) | gazer(Prowler),
kangaroo(the_kangaroo),
~ avoided(the_kangaroo)
}
@I: S is not empty
@II: S contains a unit clause: kangaroo(the_kangaroo)
@III: ~ kangaroo(the_kangaroo) is not in S
-> S = simplify(S, kangaroo(the_kangaroo))
    S = {
        ~ in_house(Cat) | cat(Cat),
        ~ gazer(Gazer) | suitable_pet(Gazer),
        ~ detested(Detested) | avoided(Detested),
        ~ carnivore(Carnivore) | prowler(Carnivore),
        ~ cat(Cat) | mouse_killer(Cat),
        ~ take_to_me(Taken_animal) | in_house(Taken_animal),
        ~ suitable_pet(Kangaroo),
        ~ mouse_killer(Killer) | carnivore(Killer),
        takes_to_me(Animal) | detested(Animal),
        ~ prowler(Prowler) | gazer(Prowler),
        ~ avoided(the_kangaroo)
    }
@II: S contains a unit clause: ~ avoided(the_kangaroo)
@III: avoided(the_kangaroo) is not in S
-> S = simplify(S, ~ avoided(the_kangaroo))
    S = {
        ~ in_house(Cat) | cat(Cat),
        ~ gazer(Gazer) | suitable_pet(Gazer),
        ~ detested(Detested),
        ~ carnivore(Carnivore) | prowler(Carnivore),
        ~ cat(Cat) | mouse_killer(Cat),
        ~ take_to_me(Taken_animal) | in_house(Taken_animal),
        ~ suitable_pet(Kangaroo),
        ~ mouse_killer(Killer) | carnivore(Killer),
        takes_to_me(Animal) | detested(Animal),
        ~ prowler(Prowler) | gazer(Prowler)
    }
@II: S contains a unit clause: ~ suitable_pet(Kangaroo)
@III: suitable_pet(Kangaroo) is not in S
-> S = simplify(S, ~ suitable_pet(Kangaroo))
    S = {
        ~ in_house(Cat) | cat(Cat),
        ~ gazer(Gazer),
        ~ detested(Detested),

```

```

    ~ carnivore(Carnivore) | prowler(Carnivore),
    ~ cat(Cat) | mouse_killer(Cat),
    ~ take_to_me(Taken_animal) | in_house(Taken_animal),
    ~ mouse_killer(Killer) | carnivore(Killer),
    takes_to_me(Animal) | detested(Animal),
    ~ prowler(Prowler) | gazer(Prowler)
}
@II: S contains a unit clause: ~ gazer(Gazer)
@III: gazer(Gazer) is not in S
-> S = simplify(S, ~ gazer(Gazer))
    S = {
        ~ in_house(Cat) | cat(Cat),
        ~ detested(Detested),
        ~ carnivore(Carnivore) | prowler(Carnivore),
        ~ cat(Cat) | mouse_killer(Cat),
        ~ take_to_me(Taken_animal) | in_house(Taken_animal),
        ~ mouse_killer(Killer) | carnivore(Killer),
        takes_to_me(Animal) | detested(Animal),
        ~ prowler(Prowler)
    }
@II: S contains a unit clause: ~ detested(Detested)
@III: detested(Detested) is not in S
-> S = simplify(S, ~ detested(Detested))
    S = {
        ~ in_house(Cat) | cat(Cat),
        ~ carnivore(Carnivore) | prowler(Carnivore),
        ~ cat(Cat) | mouse_killer(Cat),
        ~ take_to_me(Taken_animal) | in_house(Taken_animal),
        ~ mouse_killer(Killer) | carnivore(Killer),
        takes_to_me(Animal),
        ~ prowler(Prowler)
    }
@II: S contains a unit clause: ~ prowler(Prowler)
@III: prowler(Prowler) is not in S
-> S = simplify(S, ~ prowler(Prowler))
    S = {
        ~ in_house(Cat) | cat(Cat),
        ~ carnivore(Carnivore),
        ~ cat(Cat) | mouse_killer(Cat),
        ~ take_to_me(Taken_animal) | in_house(Taken_animal),
        ~ mouse_killer(Killer) | carnivore(Killer),
        takes_to_me(Animal)
    }
@II: S contains a unit clause: ~ carnivore(Carnivore)
@III: carnivore(Carnivore) is not in S
-> S = simplify(S, ~ carnivore(Carnivore))
    S = {
        ~ in_house(Cat) | cat(Cat),

```

```

    ~ cat(Cat) | mouse_killer(Cat),
    ~ take_to_me(Taken_animal) | in_house(Taken_animal),
    ~ mouse_killer(Killer),
    takes_to_me(Animal)
}
@II: S contains a unit clause: ~ mouse_killer(Killer)
@III: mouse_killer(Killer) is not in S
-> S = simplify(S, ~ mouse_killer(Killer))
    S = {
        ~ in_house(Cat) | cat(Cat),
        ~ cat(Cat),
        ~ take_to_me(Taken_animal) | in_house(Taken_animal),
        takes_to_me(Animal)
    }
@II: S contains a unit clause: ~ cat(Cat)
@III: cat(Cat) is not in S
-> S = simplify(S, ~ cat(Cat))
    S = {
        ~ in_house(Cat),
        ~ take_to_me(Taken_animal) | in_house(Taken_animal),
        takes_to_me(Animal)
    }
@II: S contains a unit clause: ~ in_house(Cat)
@III: in_house(Cat) is not in S
-> S = simplify(S, ~ in_house(Cat))
    S = {
        ~ take_to_me(Taken_animal),
        takes_to_me(Animal)
    }
@II: S contains a unit clause: ~ take_to_me(Taken_animal)
@III: take_to_me(Taken_animal) is in S
-> return "unsatisfiable"

```

Das Ergebnis "unsatisfiable" bedeutet, dass die ursprüngliche (nicht negierte) Vermutung erfüllbar ist, also eine logische Konsequenz aus den Axiomen ist.

Aufgabe 1

a) $\forall Z \exists Y \forall X (f(X, Y)) \Leftrightarrow (f(X, Z) \wedge \neg f(X, X))$

Äquivalenz in zwei Implikationen umformen:

$$\begin{aligned} \forall Z \exists Y \forall X ((f(X, Y) \Rightarrow (f(X, Z) \wedge \neg f(X, X))) \\ \wedge (f(X, Z) \wedge \neg f(X, X)) \Rightarrow f(X, Y)) \end{aligned}$$

Implikationen in Disjunktionen umformen:

$$\begin{aligned} \forall Z \exists Y \forall X (\neg f(X, Y) \vee (f(X, Z) \wedge \neg f(X, X))) \\ \wedge \neg (f(X, Z) \wedge \neg f(X, X)) \vee f(X, Y) \end{aligned}$$

Negation in der zweiten Zeile auflösen:

$$\begin{aligned} \forall Z \exists Y \forall X (\neg f(X, Y) \vee (f(X, Z) \wedge \neg f(X, X)) \\ \wedge (\neg f(X, Z) \vee f(X, X) \vee f(X, Y))) \end{aligned}$$

erste Zeile aufteilen in Konjunktion von zwei Disjunktionen:

$$\begin{aligned} \forall Z \exists Y \forall X ((\neg f(X, Y) \vee (f(X, Z) \\ \wedge (\neg f(X, Y) \vee \neg f(X, X))) \\ \wedge (\neg f(X, Z) \vee f(X, X) \vee f(X, Y))) \end{aligned}$$

Skolemisieren, so dass $\exists Y \equiv skY(Z)$:

$$\begin{aligned} (\neg f(X, skY(Z)) \vee f(X, Z)) \\ \wedge (\neg f(X, skY(Z)) \vee \neg f(X, X)) \\ \wedge (\neg f(X, Z) \vee f(X, X) \vee f(X, skY(Z))) \end{aligned}$$

schließlich bringen wir das ganze in die Mengenschreibweise der KNF:

$$\{\neg f(X, skY(Z)) \vee f(X, Z), \neg f(X, skY(Z)) \vee \neg f(X, X), \neg f(X, Z) \vee f(X, X) \vee f(X, skY(Z))\}$$

b) $\forall X \forall Y (q(X, Y)) \Leftrightarrow \forall Z (f(Z, X) \Leftrightarrow f(Z, Y))$

aufteilen:

$$\begin{aligned} \forall X \forall Y (q(X, Y)) \Leftrightarrow \forall Z (f(Z, X)) \\ \forall Z (f(Z, X)) \Leftrightarrow f(Z, Y) \end{aligned}$$

c) $\forall X \exists Y ((p(X, Y) \Leftarrow \forall X \exists T q(Y, X, T)) \Rightarrow r(Y)$

aufteilen:

$$\begin{aligned} \forall X \exists Y (\forall X \exists T q(Y, X, T)) \Rightarrow r(Y) \\ \wedge (\forall X \exists T q(Y, X, T) \Rightarrow p(X, Y)) \end{aligned}$$

in Disjunktionen umwandeln:

$$\begin{aligned} \forall X \exists Y (\neg (\forall X \exists T q(Y, X, T))) \vee r(Y) \\ \wedge (\neg (\forall X \exists T q(Y, X, T)) \vee p(X, Y)) \end{aligned}$$

Quantoren nach außen ziehen:

$$\forall X \exists Y ((\exists X \forall T (\neg q(Y, X, T) \vee r(Y)) \\ \wedge (\exists X \forall T (\neg q(Y, X, T) \vee p(X, Y))))$$

d)

Aufgabe 3

gegeben

$$V = \{X, Y\} \\ F = \{vater_von/1, mutter_von/, max/0\} \\ P = \{verheiratet/2\}$$

gesucht

Herbrand Universum

$$\begin{aligned} &max, \\ &vater_von(max), \\ &vater_von(vater_von(max)) \\ &\dots \\ &mutter_von(max), \\ &mutter_von(mutter_von(max)) \\ &\dots \\ &vater_von(mutter_von(max)), \\ &vater_von(vater_von(mutter_von(max))), \\ &mutter_von(vater_von(max)), \\ &mutter_von(vater_von(vater_von(max))), \\ &\dots \\ &mutter_von(vater_von(mutter_von(vater_von(mutter_von(...)))))) \end{aligned}$$

Herbrand Basis

$$\begin{aligned} &verheiratet(max, vater_von(max)), \\ &verheiratet(max, mutter_von(max)), \\ &verheiratet(vater_von(max), mutter_von(max)), \\ &verheiratet(vater_von(vater_von(max)), mutter_von(vater_von(max))), \\ &\dots \end{aligned}$$

Herbrand Interpretation

D

D ist das Herbrand Universum

F

Die Identitätsfunktion:

$$\begin{aligned} &max \rightarrow max, \\ &vater_von(max) \rightarrow vater_von(max), \\ &\dots \end{aligned}$$

R

weist den Elementen der Herbrand Basis Wahrheitswerte zu:

$$\begin{aligned} &verheiratet(max, vater_von(max)) \rightarrow FALSE, \\ &verheiratet(vater_von(max), mutter_von(max)) \rightarrow TRUE, \\ &verheiratet(vater_von(vater_von(max)), mutter_von(vater_von(max))) \rightarrow TRUE, \\ &\dots \end{aligned}$$