

A 2D TQFT approach to topological graph polynomials and graphs in thickened surfaces

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Overview

The Yamada polynomial [Yamada '89] is a quantum invariant of spatial graphs.

It satisfies a contraction-deletion relation like classic graph polynomials.

$$\left[\begin{array}{c} \nearrow \\ \searrow \end{array} \right] = \left[\begin{array}{c} \nearrow \\ \searrow \\ \nearrow \\ \searrow \end{array} \right] - \left[\begin{array}{c} \nearrow \\ \searrow \end{array} \right]$$

- chromatic polynomial
- flow polynomial
- Tutte polynomial

For planar graphs, it "is" the flow polynomial.

- Bollobás-Riordan polynomial

Extensions of flow polynomial to nonplanar graphs yield invariants of virtual spatial graphs:

[Fleming and Mellor '07], [McPhail-Snyder and M. '18], [Deng, Jin, and Kauffman '18]

Overview

These are relatively easy-to-compute invariants of virtual spatial graphs

- Can we find all extensions of the Yamada polynomial to virtual spatial graph invariants?

The arrow polynomial [Dye and Kauffman '09] is an invariant of virtual knots that incorporates more topological information.

- Can similar information be incorporated into a virtual Yamada polynomial?

Virtual spatial graphs

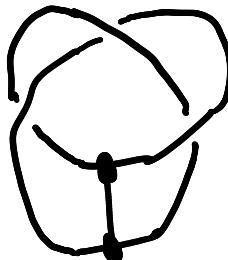
A **ribbon graph** is a "thickened" graph.

vertices \leadsto disks

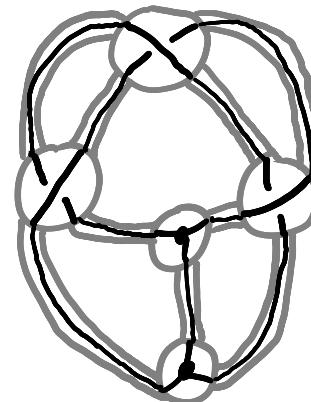
edges \leadsto thin rectangular strips

A **spatial graph** is a ribbon graph embedded in S^3 (equivalently, in $S^2 \times [0,1]$)

A **spatial graph diagram** is a planar graph with special degree-four vertices for crossings:



as a planar ribbon graph:



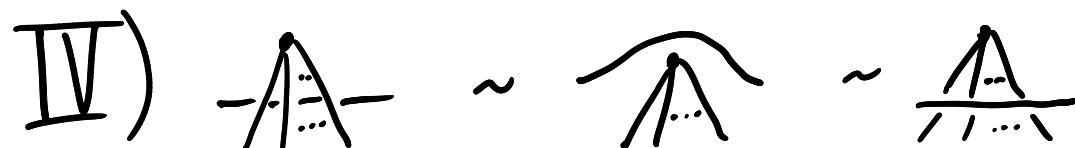
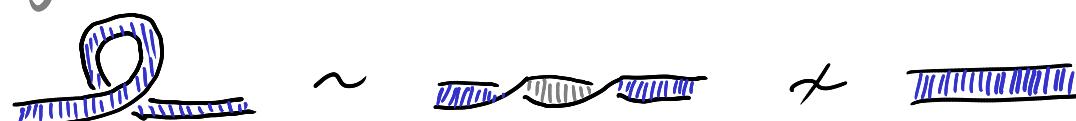
Virtual spatial graphs

Spatial graphs / isotopy \approx spatial graph diagrams / "Reidemeister moves"

Moves:



\leftarrow Regular Isotopy due to strips:

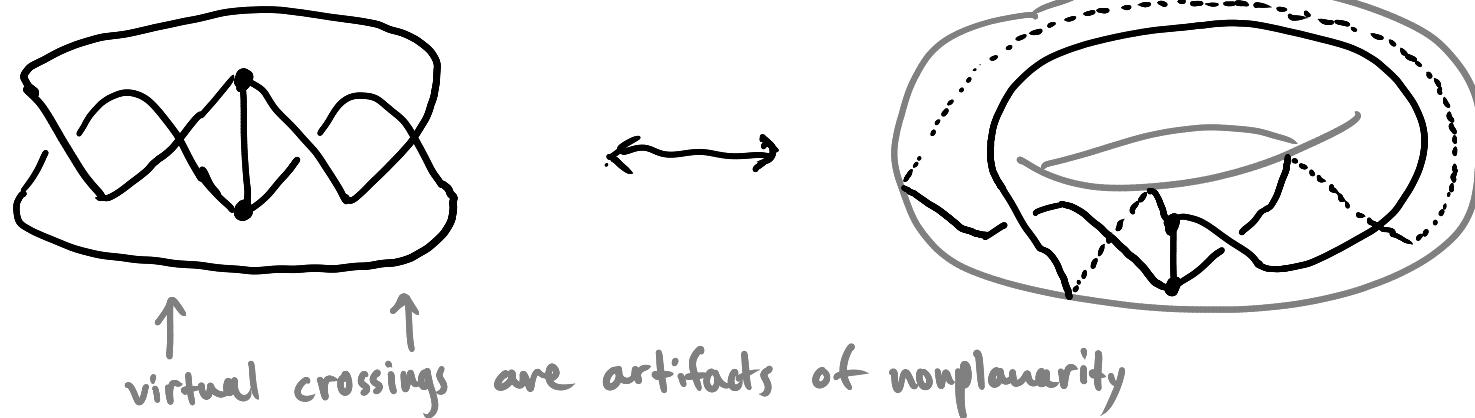


Virtual spatial graphs

A virtual spatial graph is a spatial graph diagram without regard for planarity!

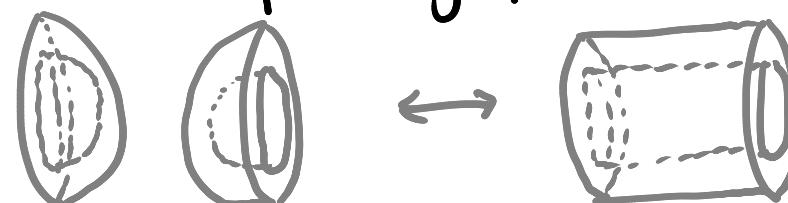
Equivalence is still via Reidemeister moves.

Ex



virtual crossings are artifacts of nonplanarity

[Carter, Kamada, and Saito '02], [Kuperberg '03] Virtual spatial graphs are ribbon graphs in thickened surfaces, modulo stabilization :



(away from graph)

Thm Can always maximally destabilize then isotope.

The Yamada polynomial

G - spatial graph

$R(G; A) \in \mathbb{Z}[A^{\pm 1}]$ is Yamada polynomial, defined by the skein relations

$$1) [G \amalg \bullet] = [G]$$

$$2) \begin{bmatrix} \nearrow & \searrow \\ \bullet & \cdot \\ \searrow & \nearrow \end{bmatrix} = \begin{bmatrix} \nearrow & \searrow \\ \cdot & \bullet \\ \searrow & \nearrow \end{bmatrix} - \begin{bmatrix} \rightarrow & \leftarrow \\ \cdot & \cdot \\ \cdot & \cdot \end{bmatrix}$$

contraction - deletion relation

$$3) \begin{bmatrix} \nearrow \circ \\ \cdot \end{bmatrix} = (A + 1 + A^{-1}) \begin{bmatrix} \nearrow \\ \cdot \end{bmatrix}$$

$$4) \begin{bmatrix} \times \\ \diagup \quad \diagdown \end{bmatrix} = A \begin{bmatrix}) \\ (\end{bmatrix} + A^{-1} \begin{bmatrix} \vee \\ \diagdown \quad \diagup \end{bmatrix} - \begin{bmatrix} \times \\ \bullet \end{bmatrix}$$

Virtual Yamada polynomials

An extension of R to virtual spatial graphs corresponds to a ribbon graph invariant f satisfying

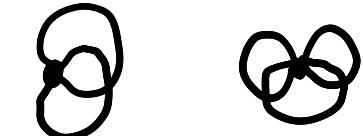
$$1) f(G \amalg \bullet) = f(G)$$

$$2) f(\text{graph with a crossing}) = f(\text{graph without crossing}) - f(\Rightarrow \Leftarrow)$$

$$3) f(\text{graph with a loop}) = (Q-1) f(\text{graph without loop})$$

Reduces virtual spatial graphs to interleaved bouquets

ex



Then: define $f(X) = Af(\text{graph}) + A'f(\text{graph with a loop}) - f(\text{graph with a crossing})$ with $Q = A + 2 + A'$

Classification

A virtual Yamada polynomial factors through the ring \mathcal{R} of formal $\mathbb{Z}[A^{\pm 1}]$ -linear combinations of virtual spatial graphs modulo the previous relations; disjoint union is multiplication.

Ihm $\mathcal{R} \cong \mathbb{Z}[A^{\pm 1}][\text{ } \otimes \text{ , } \otimes \text{ , } \dots, \text{ vertex sum of } n \text{ copies of } \text{ } \otimes \text{ , } \dots]$ (polynomial ring)

Or A virtual Yamada polynomial is determined by its value on these n -fold vertex sums.
 (No constraints.)

Calculation note

It is easier to consider a variant:

let $\phi := | - \cdot |$ be symbol for unexpanded edges, and

$$\text{i)} \quad \begin{array}{c} \text{---} \\ | \end{array} = \times$$

$$\text{ii)} \quad \begin{array}{c} \text{---} \\ | \end{array} = Q \begin{array}{c} \text{---} \\ | \end{array}$$

ex

$$\begin{aligned} &= \text{genus-2 surface} - \text{genus-1 surface} - \text{genus-1 surface} + \text{genus-0 surface} = \text{genus-2 surface} - Q - Q + 1 \\ &\quad - \text{genus-1 surface} + \text{genus-0 surface} + \text{genus-0 surface} - \dots - Q + 1 + 1 - 1 \end{aligned}$$

$$= 2 - 3Q + X_1 \in \mathbb{Z}[A^{\pm 1}, X_1, X_2, X_3, \dots]$$

$(Q = A + 1 + A^{-1})$

$(X_g$ for component
with genus g .)

Cor If X_g appears in the polynomial for a virtual spatial graph,
then its **virtual genus** is $\geq g$ (every diagram has genus $\geq g$).

Special cases

[Fleming and Mellor] $X_g = Q^{2g}$

[McPhail-Snyder and M.] $X_g = Q^g$

[Deng, Jin, and Kauffman] essentially $X_g = x^g$, giving an invariant in $\mathbb{Z}[A^{\pm 1}, x]$

Hence it interpolates the above two

A symmetric monoidal category $2\text{Cob}^{\text{open}}$ [Lauda and Pfeiffer '08]

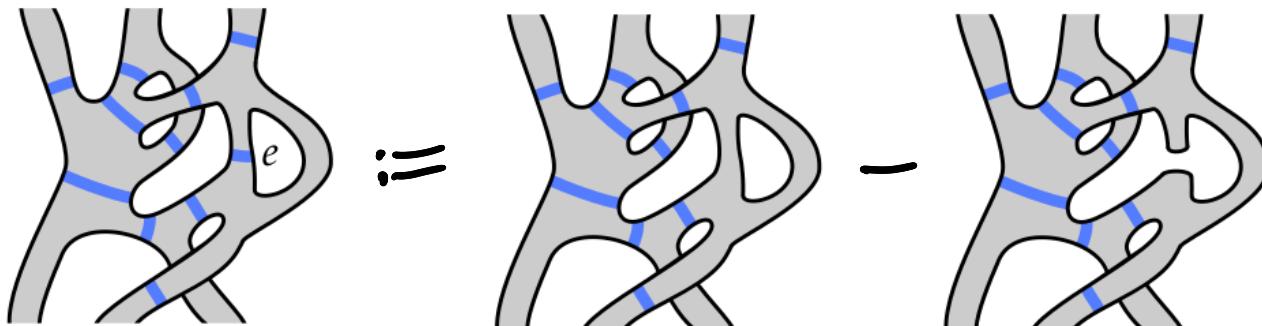
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Objects: disjoint unions of oriented intervals.

Morphisms: "open cobordisms" – surfaces with nonempty boundary

$$\left\{ \text{Symmetric monoidal functors } 2\text{D}\text{Cob}^{\text{open}} \rightarrow \text{Vect} \right\} \cong \left\{ \text{Symmetric Frobenius algebras} \right\}$$

Extend to have 1D defects defined by



Then get contraction-deletion ribbon graph invariants from functors.

Invariants from symmetric monoidal functors

Thm On ribbon graphs, every such invariant is from a symmetric Frobenius algebra that is a direct sum of some number of the following:

$$\begin{array}{ll}
 1) \text{ Mat}_n(\mathbb{C}) & \varepsilon: \text{Mat}_n(\mathbb{C}) \rightarrow \mathbb{C} \\
 & A \mapsto \text{ctr}A \\
 2) \mathbb{C}[x]/(x^n) & \varepsilon: \mathbb{C}[x]/(x^n) \rightarrow \mathbb{C} \\
 & x^k \mapsto \begin{cases} 0 & \text{if } k < n-1 \\ c & \text{if } k = n-1 \end{cases}
 \end{array}
 \quad \left. \begin{array}{l} n \geq 0 \text{ in } \mathbb{N} \\ c \neq 0 \text{ in } \mathbb{C} \end{array} \right\} \quad \left. \begin{array}{l} n \geq 2 \text{ in } \mathbb{N} \\ c \neq 0 \text{ in } \mathbb{C} \end{array} \right\}$$

Thm If the algebra is semisimple (i.e., $A \cong \text{Mat}_{n_1}(\mathbb{C}) \oplus \cdots \oplus \text{Mat}_{n_m}(\mathbb{C})$) then

$$G \mapsto \sum_{\substack{H \subseteq G \\ \text{spanning subgraph}}} (-1)^{E(G)-E(H)} \prod_F \sum_{i=1}^m (c_i^{-1} n_i)^{b_i(F)-1} n_i^{2-2g(F)}$$

Ribbon graph $H \subseteq G$ F $i=1$ $c_i^{-1} n_i$ $b_i(F)$ $n_i^{2-2g(F)}$
spanning subgraph com. component of H

Suitability for virtual Yamada polynomials

Need $\nabla \circ = (Q-1) \nabla$. Setting $c_i = Q^{-1} n_i$ works; then

$$G \mapsto \sum_{H \subseteq G} (-1)^{E(G)-E(H)} Q^{b_1(H)-b_0(H)} \prod_F Y_g \quad \text{with} \quad Y_g = \sum_{i=1}^m n_i^{2-2g}$$

For m and n_1, \dots, n_m generic, Y_0, Y_1, Y_2, \dots are algebraically independent.

Thm The collection of all symmetric monoidal functors $2D\text{Cob}^{\text{open}} \rightarrow \text{Vect}$ that extend to virtual Yamada polynomials embed \mathcal{R} in $(\mathbb{C}[A^{\pm 1}]\text{-Mod}$, ie, if G, G' are virtual spatial graphs with different images in \mathcal{R} , then there is such a functor that can distinguish them.

The Bollobás - Riordan polynomial

Using $A = (\text{Mat}_n \mathbb{C})^m$ with $c_1 = \dots = c_n = c$, then the functor computes (a version of) the Bollobás - Riordan polynomial.

$$A \cong \mathbb{C}^m \otimes \text{Mat}_n \mathbb{C} \cong \mathbb{C}^m \otimes (\mathbb{C}^n)^* \otimes \mathbb{C}^n$$

Graphical calculus :
for functor

$$\begin{array}{c} \text{---} \\ | \quad | \\ \text{---} \end{array} \mapsto c' \quad \begin{array}{c} \text{---} \\ | \quad | \\ \text{---} \end{array} - \begin{array}{c} \text{---} \\ | \quad | \\ \text{---} \end{array}$$

$$\begin{array}{c} \text{---} \\ | \quad | \\ \text{---} \end{array} \mapsto c \quad \begin{array}{c} \text{---} \\ | \quad | \\ \text{---} \end{array}$$

with $\begin{array}{c} \circlearrowright \\ \text{---} \end{array} = n$ $\bullet = m$ $\begin{array}{c} \text{---} \\ | \quad | \\ \text{---} \end{array} = \begin{array}{c} \text{---} \\ | \quad | \\ \text{---} \end{array}$ $\begin{array}{c} \text{---} \\ | \quad | \\ \text{---} \end{array} = \begin{array}{c} \text{---} \\ | \quad | \\ \text{---} \end{array}$

If $c_{mn} = 1$ and $mn^2 = A + 2 + A^{-1}$, get a 2D locus of virtual Yamada polynomials.

Encompasses [Fleming and Mellor '07], [McPhail-Snyder and M. '18], [Deng, Jin, and Kauffman '18]

Note: with $m=1$, $c = -A^2$, $n = -A^2 - A^{-2}$, reduces to rescaled Kauffman bracket [Dasbach, Futer, Kalfagianni '08]

Arrow polynomial [Dye and Kauffman '09]

Kauffman bracket for oriented virtual links.

$$\left[\begin{array}{c} \nearrow \\ \times \\ \searrow \end{array} \right] = A \left[\begin{array}{c} \nearrow \\ \uparrow \\ \searrow \end{array} \right] + A^{-1} \left[\begin{array}{c} \nearrow \\ \curvearrowleft \\ \searrow \end{array} \right] \quad \rightarrow \text{polynomial in } \mathbb{Z}[A^{\pm 1}, X_1, X_2, \dots] \\ \text{using a rule for net \# of cusps.}$$

[M. '20 unpublished] Can compute generalization using a functional on skein module.

$L \subset \Sigma \times [0,1]$ oriented link, there is map $H_1(L) \rightarrow H_1(\Sigma)$

given simple closed curve $C \subset \Sigma$, there is bilinear map

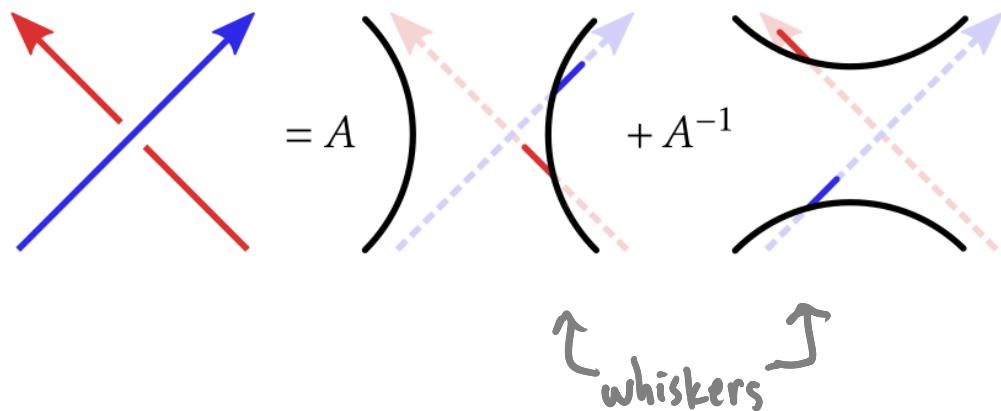
$$H_1(L) \otimes H_1(C) \xrightarrow{\cap} H_0(\Sigma) \xrightarrow{\epsilon} \mathbb{Z}$$

Fixing an $\alpha \in H_1(C)$, by U.C.T. get element of $H^*(L) \cong H_0(L)$.

Weight C in Kauffman bracket expansion by $X_{\pm[C]}$ with $\pm[C] \in H_0(L)$, $X_0 := 1$

Arrow polynomial

Practically,



ex

$$= (-A^2 - A^{-2}) X_{\pm(1, -1)}$$

ex

$$= A + A^{-1}$$

$$= A(-A^2 - A^{-2}) X_{1,-1} + A^{-1}(-A^2 - A^{-2}) X_{1,1} \xrightarrow[\text{usual}{ } \text{Arrow polynomial}]{} A(-A^2 - A^{-2}) + A^{-1}(-A^2 - A^{-2}) X_1$$

$$X_{(a,b)} \mapsto X_{(a+b)/2}$$

Arrow Yamada polynomial

Consider G a virtual spatial graph, represented as a ribbon graph in $\Sigma \times [0,1]$.

Can expand G in Yamada Skein Module:

$$\text{Diagram} = A \text{ }) \text{ } (+ A^{-1} \text{ } \text{Diagram} - \text{Diagram} \quad (\text{along with contraction-deletion, etc.})$$

Also assume we used $\phi := | - \frac{1}{2}$ convention.

Result: linear combination of t-complexes in Σ .

How to put a functional on Yamada Skein Module so that one obtains a virtual spatial graph invariant?

Arrow Yamada polynomial

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Rather than just 1-complex $\Gamma \leadsto X_g$ with $g = \text{genus } \Gamma$ like before, use bilinear form

$$H_1(G) \otimes H_1(\Gamma) \xrightarrow{\Delta} H_0(\Sigma) \xrightarrow{\epsilon} \mathbb{Z}$$

For each connected Γ , take matrix M_Γ of this.

Define $[\Gamma] = M_\Gamma \cdot GL(H_1(\Gamma))$, the coset of action on matrices.
(so it doesn't depend on basis of $H_1(\Gamma)$)

Then can weight component Γ of expansion by $X_g Y_{M_\Gamma \cdot GL(H_1(\Gamma))}$

If we do not have a preferred basis for $H_1(G)$, can do double cosets

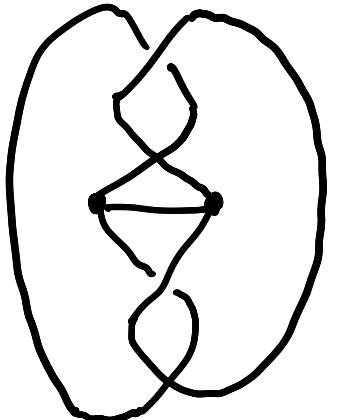
$GL(H_1(G)) \cdot M_\Gamma \cdot GL(H_1(\Gamma)) \rightsquigarrow$ Smith normal form

get sequence d_1, d_2, d_3, \dots with $d_i | d_{i+1}$

hence get variables $Y_{d_1, d_2, \dots}$

Arrow Yamada polynomial

ex



$$\begin{aligned}
 & -\frac{(-1+A)^3 (1+A)^2}{A^3} + \frac{(-1+A) (1+A)^2 (4+3A+2A^2+3A^3+4A^4) Y[1]}{A^4} - \\
 & \frac{(-1+A) (1+A)^4 (1+A+A^2+A^3+A^4) Y[1]^2}{A^5} - \frac{2 (-1+A) (1+A)^2 Y[2]}{A^2} - \frac{(-1+A)^3 (1+A)^4 X[1] \times Y[1, 2]}{A^4}
 \end{aligned}$$

Both Y and X terms indicate positive virtual genus.

Thanks!