

ETCH

Define **Once**, Etch **Forever**.



A language that *keeps you safe* before it even runs,
yet feels like *you're just scripting* - **until it blows your mind!**

Why Etch? 🤔

The Origin Story Nobody Asked For (But You're Getting Anyway)

The Endless Quest: 🔍

- Spent *years* hunting for the "perfect" scripting engine.
- Lua? *Too bare*. Python? *Too slow*. JavaScript? *Too... JavaScript*.
- Tried them all. They all lacked *something*.
- Performance ⚡ vs Safety 🛡 vs Simplicity 📝. You can only **pick two**.

Why Etc? 🤔

The Origin Story Nobody Asked For (But You're Getting Anyway)

The "Why Not?" Moment: 💡

- **Plot twist:** Compilers are actually **FUN** to build! 🎉
- Living in the age of AI: "Make an app!" 🤖
- Me: "...What if I make a LANGUAGE instead?" 🎯
- **YOLO Compiler Theory!**

Why Etch? 🤔

The Origin Story Nobody Asked For (But You're Getting Anyway)

The Mad Experiment: 🚀

- Built Etch as a *playground* for compiler tech
- Prove correctness? *Why not!*
- Safety without complexity? *Challenge accepted!*
- **Etch: Because life's too short** 🔥

The Bugs That Haunt Your Dreams

Those sneaky runtime gremlins - the ones that compile just fine, only to explode right when you hit play.

Numerical Safety Issues:

- ✗ Division by zero crashes
- ✗ Integer overflow vulnerabilities

Memory Safety Issues:

- ✗ Nullptr dereferencing
- ✗ Uninitialized variable bugs






Array and Bounds Issues:

- ✗ Array bounds errors
- ✗ Buffer overflows
- ✗ Platform-dependent behavior




Compile-Time Safety Verification

Etch proves safety properties at compile-time through static analysis.

Safety Guarantees:

-  No division by zero
-  No integer overflow
-  No nullptr dereferences
-  No uninitialized variables
-  No array out of bounds errors

Additional Benefits:

-  Dead code automatically eliminated
-  Redundant checks removed
-  Errors caught before deployment

nuqneH, Etch!

```
1 fn main() -> void {  
2     print("Hello, World!");  
3 }
```

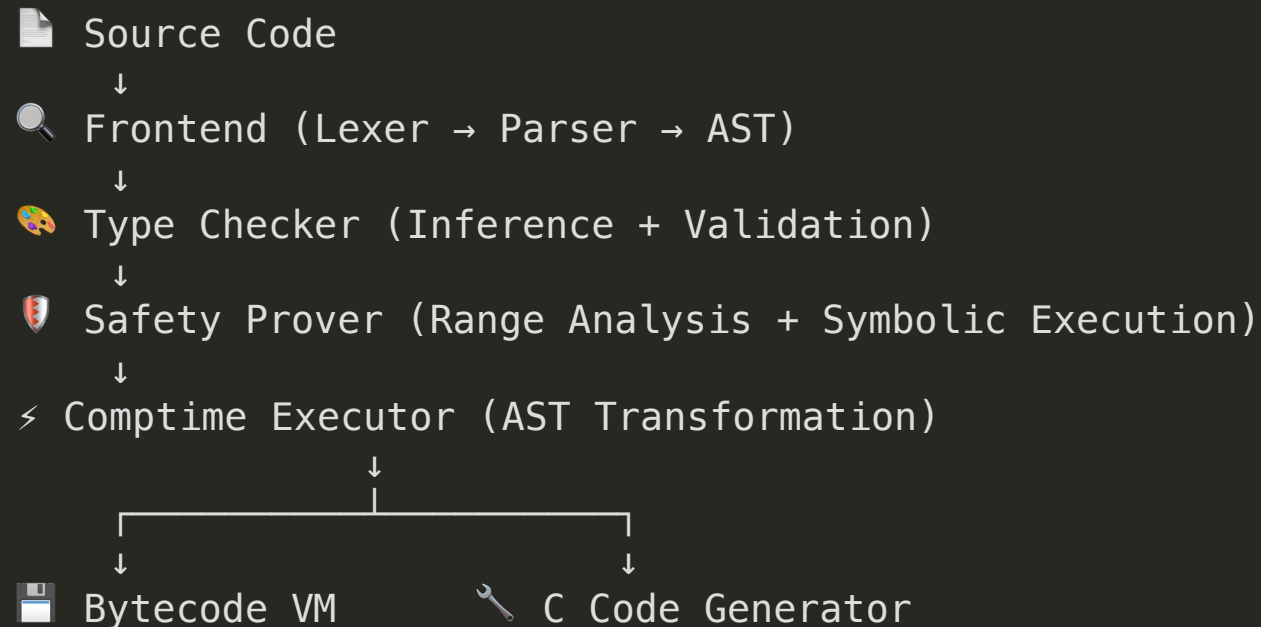
C-like syntax you already know-*simple, clean, familiar.*

No need to learn *Klingon* to say hello! 



Language Architecture

Multi-stage compilation pipeline:



✓ Each stage adds safety guarantees



The Compilation Pipeline

Staged transformation with safety guarantees at each step:

Stage	Process	Output
1	Lexer	Token stream (frontend/lexer)
2	Parser	AST (frontend/ast - frontend/parser)
3	Type Checker	Typed + Inferred AST (typechecker/)
4	Safety Prover	Proven AST (prover/)
5	Comptime Executor	Transformed AST (comptime/)
6	Bytecode Generator	Register VM bytecode (interpreter/)
7	Cache	Disk (___etch___/ directory)
8	VM / C Backend	Execution or native binary (backend/)



Frontend: Lexer → Parser

Token-based lexical analysis:






- 📍 Position tracking for error reporting
- ⚡ Efficient single-pass scanning
- 🎯 Source location preservation

Recursive descent parser:




- 🌳 Produces strongly-typed AST
- 📝 Expression kinds: Binary ops, Calls, Arrays, Pattern matching
- 🔁 Statement kinds: Declarations, Control flow, Assignments
- 🛡️ Syntax validation during parsing

AST Node Types

Expression variants (20+ types):

-  **Literals:** ekInt, ekFloat, ekString, ekBool, ekChar
-  **Operations:** ekBin, ekUn, ekCall
-  **Collections:** ekArray, ekIndex, ekSlice, ekArrayLen
-  **Pattern matching:** ekMatch, ekOptionSome, ekResultOk
-  **Objects:** ekObjectLiteral, ekFieldAccess, ekNew

Statement variants (15+ types):

-  **Declarations:** skLet, skVar, skFun, skType
-  **Control flow:** skIf, skWhile, skFor, skReturn
-  **Others:** skAssign, skDefer, skComptime

Type System: Three-Phase Process

Phase 1: Type Collection

- Gather all type definitions and global variables
- Build forward reference table
- Create type environment

Type System: Three-Phase Process

Phase 2: Type Inference

- Hindley-Milner style inference for generics
- Constraint generation and unification
- Return type inference from body
- Generic function instantiation on demand







Type System: Three-Phase Process

Phase 3: Type Validation




- Expression type checking (expressions.nim)
- Statement type checking (statements.nim)
- Full program validation

Type Kinds

Primitive Types:





-  tkBool (1 byte)
-  tkChar (1 byte)
-  tkInt (8 bytes, signed)
-  tkFloat (8 bytes, double)
-  tkString
-  tkVoid

Composite Types:

-  tkArray - Fixed and dynamic arrays
-  tkRef - References with generational tracking
-  tkObject - Structured data with fields

Type Kinds

Advanced Types:

-  `tkOption` - Optional values (some/none)
-  `tkResult` - Result types (ok/error)
-  `tkGeneric` - Generic type parameters
-  `tkUnion` - Sum types



Prover: Safety Through Analysis

Safety analysis through symbolic execution

Core Components:

- 🔍 **expression_analysis**: Range propagation, expression evaluation
- ➕ **binary_operations**: Arithmetic operation range inference
- 🔁 **symbolic_execution**: Control flow analysis
- ⚡ **function_evaluation**: Pure function compile-time execution

Key Concept is tracking values

- known/unknown
- min/max range
- non-zero/non-null
- initialized/used/last use
- is array/string
- array/string size

Every variable has a proven state at every program point!



Prover: Analysis Phases

Phase 1: Environment Setup

- Initialize environment with variable declarations
- Add all globals to environment



Prover: Analysis Phases

Phase 2: Global Analysis 🌐

- Analyze global variable initializations
- Track global value ranges



Prover: Analysis Phases

Phase 3: Function Analysis

- Prove main function first
- Analyze all reachable functions
- Validate safety at every operation



Prover: Analysis Phases

Phase 4: Property Validation ✓

- Check non-zero divisors (division/modulo)
- Verify array bounds safety
- Prevent integer overflow
- Ensure variable initialization

The Safety Prover in Action

```
1 // ✓ Safe!  
2 fn main() -> void {  
3     let divisor: int = rand(5, 10); // Range: [5, 10]  
4     let calculation: int = 100 / divisor; // ✓ Safe!  
5     print(calculation);  
6 }
```

Safe Example: divisor ranges [5, 10] → division succeeds ✓

The Safety Prover in Action

```
1 // × COMPILER ERROR
2 fn main() -> void {
3     let divisor: int = rand(5);           // Range: [0, 5]
4     let calculation: int = 100 / divisor; // × COMPILER ERROR
5     print(calculation);
6 }
7
8 /*
9  Compiling: xyz.etch
10 xyz.etch:4:32: error: cannot prove divisor is non-zero in main
11 3 |     let divisor: int = rand(5);           // Range: [0, 5]
12 4 |     let calculation: int = 100 / divisor; // × COMPILER ERROR
13   |                                     ^
14 5 |     print(calculation);
15 */
```




Unsafe Example: divisor ranges [0, 5] → *COMPILER ERROR* **×**

Compiler tracks value ranges and proves divisor should be non-zero.



Arrays with Safety Guarantees

```
1 fn main() -> void {  
2     let numbers: array[int] = [10, 20, 30, 40, 50];  
3  
4     let count: int = #numbers;           // Length operator  
5     let middle: int = numbers[count / 2]; // Bounds checked  
6     let slice = numbers[1:4];           // Safe slicing, inferred as array[int]  
7 }
```

-  Compile-time bounds checking when possible
-  Compiler enforces insertion of runtime checks when necessary
-  Clear error messages

Type System & Inference

```
1 fn main() -> void {  
2     let x: int = 42;           // Explicit type  
3     let y = 3.14;             // Inferred as float  
4     let name = "Etch";        // Inferred as string  
5     let numbers = [1, 2, 3];  // Inferred as array[int]  
6 }
```

- ✨ Strong static typing
- 🔮 Smart type inference
- ✅ No surprises






Types and Objects

Type Aliases, Unions and Object Definitions:

```
1 // Type aliases for clarity
2 type UserId = int;
3 type Email = distinct string;
4
5 // Union types for sum types
6 type IntOrString = int | string;
7
8 // Object types with fields
9 type Point = object {
10   x: int;
11   y: int;
12 };
13
14 fn main() -> void {
15   let p: Point = { x: 10, y: 20 };
16   print(p.x + p.y); // 30
17 }
```

Safety Guarantees:

-  All object fields must be initialized before use
-  Prover tracks initialization of each field
-  Compile-time error if accessing uninitialized field



Uniform Function Call Syntax

Call functions as methods using dot notation:

```
1 fn add(a: int, b: int) {  
2   return a + b;  
3 }  
4  
5 fn main() {  
6   var x: int = 10;  
7   var y: int = 20;  
8  
9   // Traditional call  
10  discard add(x, y);  
11  
12  // UFCS – first argument becomes receiver  
13  var result2: int = x.add(y);  
14  print(result2); // 30  
15  
16  // But also!  
17  print(5.add(15)); // 20  
18 }
```

Clean, readable method-style calls without OOP overhead! 🚀



Pattern Matching: Option & Result

Safe value extraction with exhaustive matching:

```
1 fn divideInts(a: int, b: int) -> result[int] {
2   if b == 0 {
3     return error("Division by zero");
4   } else {
5     return ok(a / b);
6   }
7 }
8
9 fn main() -> void {
10  let divResult: result[int] = divideInts(42, 6);
11  let message: string = match divResult {
12    ok(value) => {
13      "Success: " + toString(value);
14    }
15    error(err) => {
16      "Failed: " + err;
17    }
18  };
19  print(message); // "Success: 7"
20 }
```

Pattern Matching: Advanced

Option types and nested patterns:

```
1 fn tryGetElement(arr: array[int], index: int) -> option[int] {
2   if index >= 0 and index < #arr {
3     return some(arr[index]);
4   } else {
5     return none;
6   }
7 }
8
9 fn main() -> void {
10  let numbers: array[int] = [10, 20, 30];
11  let maybeValue: option[int] = tryGetElement(numbers, 1);
12
13  let result: string = match maybeValue {
14    some(value) => "Found: " + toString(value);
15    none => "Not found";
16  };
17  print(result); // "Found: 20"
18 }
```

Compiler enforces exhaustive pattern coverage! 



Defer: Guaranteed Cleanup

Execute cleanup code when scope exits:

```
1 fn main() {  
2     print("Start");  
3  
4     defer { print("Cleanup - runs last!"); }  
5  
6     print("Middle");  
7  
8     defer { print("Second cleanup"); }  
9  
10    print("End");  
11 }
```

Output:

```
Start  
Middle  
End  
Second cleanup  
Cleanup - runs last!
```

Defers execute in reverse order (LIFO) - Perfect for resource cleanup! 🧹



Import System: Modules & CFFI

Module Imports (Etch code):

```
1 import math
2 import math { sqrt, pow }
```

C FFI Imports (Native libraries):

```
1 import ffi cmath {
2   fn sin(x: float) -> float;
3   fn cos(x: float) -> float;
4   fn sqrt(x: float) -> float;
5 }
6
7 fn main() -> void {
8   var pi: float = 3.14159;
9   var sine: float = sin(pi / 2.0);
10  print(sine); // ~1.0
11 }
```

Zero-cost abstractions:







- 🚀 Direct C function calls (no overhead)
- 🔗 Dynamic library loading (runtime)
- ✅ Type-safe FFI boundaries

Compile-Time Execution

Comptime evaluation during compilation:

Not macros or templates-actual code execution in the compiler.

Comptime Use Cases

-  **Build-time configuration:** Different builds from same source
-  **Lookup tables:** Compute once at compile-time, use at runtime
-  **Feature flags:** Conditional compilation based on environment
-  **Resource embedding:** Templates, shaders, assets in binary
-  **Version information:** Embed git commit hash, build date
-  **Platform-specific code:** Single codebase for multiple targets

Comptime Basics

```
1 fn square(x: int) -> int {  
2     return x * x;  
3 }  
4  
5 fn main() -> void {  
6     // Prints 64 during compilation  
7     comptime{ print(square(8)); }  
8  
9     // Evaluates to constant at compile-time  
10    let sq: int = comptime(square(8));  
11    print(sq);  
12 }
```

Zero runtime overhead

- Function calls happen during compilation
- Results embedded as constants in bytecode
- No function call overhead at runtime

Comptime Blocks

```
1 fn main() -> void {  
2     comptime {  
3         print("Hello from the compiler!");  
4  
5         var i: int = 0;  
6         while i < 5 {  
7             i = i + 1;  
8         }  
9  
10        print(i);  
11    }  
12  
13    print("Hello from runtime!");  
14 }
```

Compile-time output:

```
Hello from the compiler!  
5
```

Runtime output:

```
Hello from runtime!
```

File Embedding

```
1 fn main() -> void {  
2     // File read at COMPILE-TIME  
3     let config: string = comptime(readFile("config.txt"));  
4     print(config); // File embedded in binary!  
5 }
```

Embed files directly into your binary

- No runtime I/O
- No missing file errors
- Single executable deployment

Code Injection

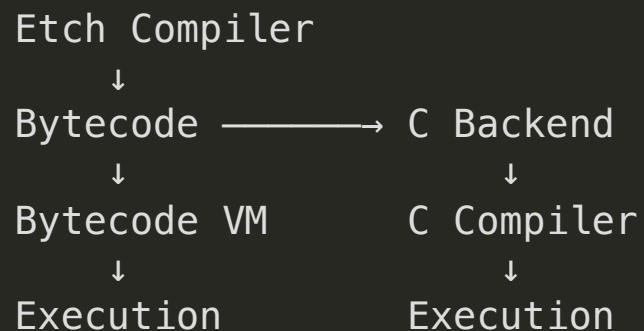
```
1 fn main() -> void {
2     comptime {
3         let env: string = readFile(".env");
4
5         if env == "production" {
6             inject("LOG_LEVEL", "int", 0);
7             inject("DEBUG", "bool", false);
8         } else {
9             inject("LOG_LEVEL", "int", 2);
10            inject("DEBUG", "bool", true);
11        }
12    }
13
14    if DEBUG { // Variable was injected!
15        print("Debug mode enabled");
16    }
17 }
```

Metaprogramming without macros, generate code based on compile-time conditions



Two Execution Modes

Single source, dual execution paths:



🔧 **Bytecode VM:** Fast iteration, debugger support, portable caching

🚀 **C Backend:** Native performance, compiler optimizations, standalone binary

Development workflow: 🖋️ VM for rapid → 📦 C for deployment






Bytecode VM - Fast Development

Perfect for development iteration:

- ⚡ Instant execution (cached)
- 🐛 Full debugger support
- 📦 Portable bytecode
- 🚫 No C compiler needed
- 🔗 FFI to C libraries (runtime)
- 📝 Rich runtime errors

C Backend - Maximum Performance

Compiles to clean, readable C code:

-  Native machine code via gcc/clang
-  Platform optimizations (-O3)
-  System calling conventions
-  FFI to C libraries (linktime)
-  Standalone executable



Register-Based VM Architecture

⚡ RegVM: Lua-inspired register machine ⚡

Stack VM	Register VM
PUSH 5	LoadK r0, 5
PUSH 3	LoadK r1, 3
ADD	Add r2, r0, r1
POP r0	







Architecture Details:

- 🖨️ 256 registers per function frame (8-bit addressing)
- 📁 65536 constants per function (16-bit index)
- 🎯 3-address instruction format ($A = B \text{ op } C$)
- 🛠️ Multiple instruction encodings: ABC, ABx, AsBx, Ax
- ⚡ Fused instructions for common patterns



Bytecode Instructions

Instruction Categories:

Category	Instructions
 Load/Store	LoadK, Move, GetGlobal, SetGlobal
 Arithmetic	Add, Sub, Mul, Div, Mod, Neg
 Comparison	Eq, Lt, Le, Gt, Ge, Ne
 Control Flow	Jump, JumpIf, JumpIfNot, TestJump
 Function Calls	Call, Return
 Fused Instructions	AddAdd, MulAdd, LoadAddStore, EqStore



Bytecode Caching

Performance improvement: 🚀 faster subsequent runs

First run:

Source → Parse → Typecheck → Prove → Compile → Cache

Subsequent runs:

Source Hash Check → Load Cached Bytecode → Run

Cache invalidation:

- ✓ Source file changed → recompile
- ✓ Source hash mismatch → recompile
- ✓ Bytecode version changed → recompile
- ✓ Compiler binary changed → recompile
- ✓ Compiler flags changed → recompile

Etch compiler is written in Nim

Why Nim for compiler development:

- ✨ Compiles to C → Portable, fast execution
- ✨ Python-like syntax → Readable codebase
- ✨ Zero-cost abstractions → Efficient compilation
- ✨ Memory safe with no GC → No crashes and leaks
- ✨ Strong standard library → Less boilerplate
- ✨ Metaprogramming → DSL capabilities

Benefits for Etch:

- 🌳 Clean AST representation with algebraic types
- 🔄 Pattern matching for compiler passes
- 🔗 Easy C FFI for library integration
- ⌚ Fast compilation of the compiler itself (< 6 seconds)



VSCode Debugger Integration

Full DAP (Debug Adapter Protocol) Support ✨

Features:

- 🛑 Set breakpoints in .etch files
- ⏮ Step through execution (step in/out/over)
- 📊 View call stack
- 🔍 Inspect variables
- 👁 Watch expressions (in progress)
- 📌 Conditional breakpoints (in progress)

Debug Server:

- 🔌 DAP protocol implementation (console based)
- 🌐 TCP/IP communication (in progress)
- 🖨 Integrated with RegVM











Benchmark Results

Real benchmark data from **hyperfine** on an M3 (generated 2025-10-22)



Benchmark Results

Real benchmark data from **hyperfine** on an M3 (generated 2025-10-22)

Benchmark	C	VM	Python 3	C vs Py	VM vs Py
 Arithmetic ops	6.5ms	115.8ms	103.5ms	15.9x	0.9x
 Array ops	6.9ms	32.9ms	42.8ms	6.2x	1.3x
 For loops	10.7ms	18.1ms	39.4ms	3.7x	2.2x
 Function calls	14.2ms	56.8ms	32.9ms	2.3x	0.6x
 Math intensive	5.0ms	26.8ms	31.2ms	6.2x	1.2x
 Memory alloc	2.6ms	11.3ms	23.6ms	9.0x	2.1x
 Nested loops	6.6ms	66.0ms	40.3ms	6.1x	0.6x
 String ops	13.9ms	10.3ms	25.2ms	1.8x	2.4x

Key Takeaway: C backend ~1.8-15.9x faster than Python, VM competitive for many workloads



More Optimization Opportunities

Current Optimizations:

- Constant folding (partial)
- Dead code elimination
- Range-based check elimination
- Fused instructions
- Bytecode caching

Future Optimizations:

- Loop hoisting optimizations
- Common subexpression elimination
- Type-specialized instructions
- Function inlining
- Register coalescing

The compiler keeps getting faster!

Development Experience

Etch provides multiple tools for exploration:

Experimentation:

- ⚡ Test compile time evaluation limits
- 🔬 Understand prover range analysis
- 📊 Profile VM vs C backend performance
- 🎨 Generate code through metaprogramming

Tooling:





- 🎨 VSCode extension with syntax highlighting
- 🐛 DAP debugger integration
- 📝 Verbose logging for compiler internals
- 📈 Performance benchmarking with `just perf`

Learning compiler technology:

- 🔍 Inspect bytecode with verbose mode
- 🛡️ Study prover analysis output
- 📄 Compare generated C code
- 🚀 Understand optimization passes

Who Is Etch For?

Perfect for:

-  Compiler enthusiasts exploring PL design
-  Programmers trying safety without complexity
-  Game developers needing a fun scripting runtime
-  Discovering program verification



Current Status

Language Status: Active Development 🚧

What works:

- ✓ Core language features
- ✓ Safety prover with range analysis
- ✓ Compile-time execution
- ✓ Bytecode VM with caching
- ✓ C code generation backend
- ✓ VSCode debugger integration
- ✓ Test framework
- ✓ Performance benchmarking




Production ready? Not yet! ⚠️

Great for: 🧪 Experiments · 📖 Learning · 🔬 Research






Optimization Roadmap

Phase 1: Bytecode optimization

-  Re-enable optimizer
-  Enhanced constant folding
-  Integrate prover data into compiler

Phase 2: Instruction improvements

-  Jump target tables
-  ARG instructions
-  Reversed operations



Roadmap

Phase 3: Advanced optimizations 🚀

- 🔍 Peephole optimization
- 🧮 Common subexpression elimination
- 🔄 Loop optimizations

Phase 4: Type-aware optimization 🎨

- 🎯 Static type specialization
- 📦 Function inlining

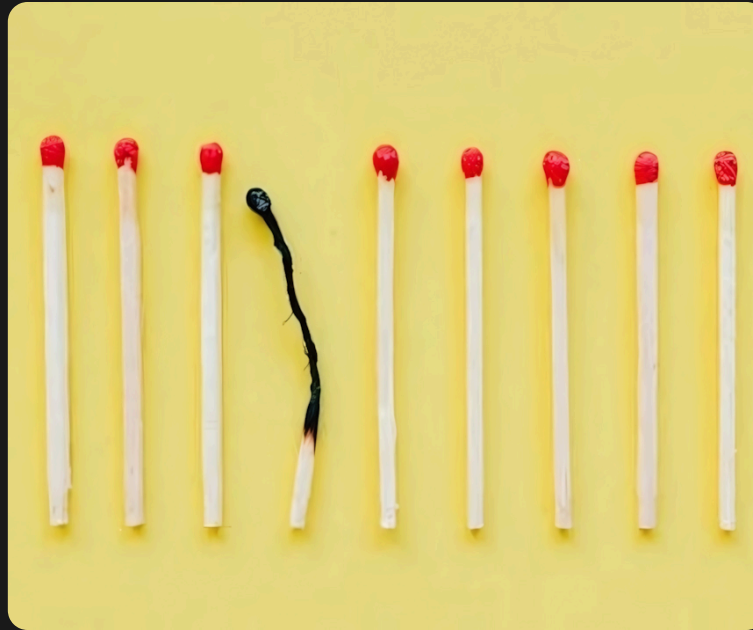
Let's see Etch in action! 🎬

Demos:

1. 🛡️ Safety proofs catching bugs
2. ⚡ Comptime execution
3. 🐛 Debugger in VSCode
4. 📊 Performance comparison
5. 🔧 C backend code generation

Questions?

"Define once, Etch forever."



Thank you! 🚀