# Network assignment 4

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#### **P2**

Sum of 16-bit integers:

```
\begin{array}{c} 00000000 & 00000001 \\ 00000010 & 00000011 \\ 00000100 & 00000101 \\ 00000110 & 00000111 \\ + & 00001000 & 00001001 \\ \hline 00010100 & 00011001 \\ \end{array}
```

Take 1's complement of the sum:  $checksum = 11101011 \ 11100110$ 

#### **P7**

(a) 
$$y = Np(1-p)^{N-1}$$
  
 $\frac{dy}{dp} = N(1-p)^{N-1} - Np \cdot (N-1)(1-p)^{N-2}$   
Let  $\frac{dy}{dp} = 0$   
 $N(1-p)^{N-1} - Np \cdot (N-1)(1-p)^{N-2} = 0$   
 $\implies p = 1 \text{ or } p = \frac{1}{N}$   
Since  $p < 1$ ,  $p = \frac{1}{N}$ 

(b) Plug in 
$$p = \frac{1}{N}$$
  
 $y = (1 - \frac{1}{N})^{N-1} = (1 - \frac{1}{N})^N \cdot (1 - \frac{1}{N})^{-1}$   
 $\lim_{N \to \infty} y = e^{-1} \cdot 1 = e^{-1}$ 

#### **P8**

$$\begin{split} y &= Np(1-p)^{2(N-1)} \\ \frac{\mathrm{d}y}{\mathrm{d}p} &= N(1-p)^{2(N-1)} - Np \cdot 2(N-1)(1-p)^{2N-3} \\ \mathrm{Let} \ \frac{\mathrm{d}y}{\mathrm{d}p} &= 0 \\ N(1-p)^{2(N-1)} - Np \cdot 2(N-1)(1-p)^{2N-3} &= 0 \\ \Longrightarrow p &= 1 \ \mathrm{or} \ p = \frac{1}{2N-1} \\ \mathrm{Since} \ p &< 1, \ p = \frac{1}{2N-1} \ \mathrm{Plug} \ \mathrm{in} \ p = \frac{1}{2N-1} \\ y &= \frac{N}{2N-1}(1-\frac{1}{2N-1})^{2(N-1)} \\ \mathrm{lim}_{N\to\infty} \ y &= e^{-1} \cdot \frac{1}{2} = \frac{1}{2e} \end{split}$$

### P14

For a 1Mbps Ethernet:  $100 \times 512$ bit times = 51.2msFor a 10Mbps Ethernet:  $100 \times 512$ bit times = 5.12ms

## P15

Last bit of B's jam signal reaches A: 225 + 273 = 498bit time A start retransmission: 498 + 96 = 594bit time A's retransmission signal reaches B: 594 + 225 = 819bit time

B returns to Step 2: 273 + 512 = 785bit time

Since 819 - 785 < 96, the retransmission won't collide and the retransmission of B will delay.