

# 1 A proof of focalisation

Let  $\vdash \Gamma; \Pi$  denote a sequent of  $\text{LL}_{\text{foc}}$  as defined in [Laurent04].

We define the focalised syntactic phase model as  $(M, \perp, \varphi)$  where  $M$  is the free commutative monoid over formulas of MALL,  $\perp = \{\Gamma \in M \mid \vdash \Gamma\}$ , and  $\varphi(X) = \{X\}^\perp$  for positive atoms  $X$ . Let  $\llbracket A \rrbracket$  be the interpretation of a formula  $A$  in this model.

For a formula  $A$ , let  $|A|$  denote the number of main negative subformulas in  $A$ . Define  $\Psi_A$  as an  $|A|$ -ary monotonous operator on  $\mathcal{P}(M)$  by induction:

- $\Psi_N(N_1) = N_1$  if  $N$  is negative
- $\Psi_X() = \{X^\perp\}$
- $\Psi_{B \otimes C}(B_1, \dots, B_{|B|}, C_1, \dots, C_{|C|}) = \Psi_B(B_1, \dots, B_{|B|}) \cdot \Psi_C(C_1, \dots, C_{|C|})$
- $\Psi_{B \oplus C}(B_1, \dots, B_{|B|}, C_1, \dots, C_{|C|}) = \Psi_B(B_1, \dots, B_{|B|}) \cup \Psi_C(C_1, \dots, C_{|C|})$
- $\Psi_1() = \{\emptyset\}$
- $\Psi_0() = \emptyset$

**Lemma 1.** *For any formula  $A$  with main negative subformulas  $A_1, \dots, A_{|A|}$ ,  $\Psi_A(\{A_1\}^\perp, \dots, \{A_{|A|}\}^\perp)^{\perp\perp} \subseteq \{A\}^\perp$ .*

*Proof.* To simplify the notations, let  $\vdash \Gamma; N$  mean  $\vdash \Gamma, N$ ; when  $N$  is a negative formula. Let  $\text{Foc}(A) = \{\Gamma \in M \mid \vdash \Gamma; A\}$ . Clearly  $\text{Foc}(A) \subseteq \{A\}^\perp$  by the *foc* rule.

We prove by induction on  $A$  that  $\Psi_A(\text{Foc}(A_1), \dots, \text{Foc}(A_{|A|})) \subseteq \text{Foc}(A)$ :

- If  $A$  is negative, the result is trivial.
- If  $A = X$ , then  $\Psi_X() = \{X^\perp\} \subseteq \text{Foc}(X)$  by the *ax* rule.
- If  $A = B \otimes C$ , then we have

$$\begin{aligned} \Psi_A(\text{Foc}(A_1), \dots, \text{Foc}(A_{|A|})) &= \Psi_B(\text{Foc}(B_1), \dots, \text{Foc}(B_{|B|})) \cdot \Psi_C(\text{Foc}(C_1), \dots, \text{Foc}(C_{|C|})) \\ &\subseteq \text{Foc}(B) \cdot \text{Foc}(C) \end{aligned}$$

by the induction hypothesis; moreover,

$$\frac{\vdash \Gamma; B \quad \vdash \Delta; C}{\vdash \Gamma, \Delta; B \otimes C} \otimes$$

hence  $\text{Foc}(B) \cdot \text{Foc}(C) \subseteq \text{Foc}(B \otimes C)$ , from which the result follows.

- If  $A = B \oplus C$ , then we have

$$\begin{aligned} \Psi_A(\text{Foc}(A_1), \dots, \text{Foc}(A_{|A|})) &= \Psi_B(\text{Foc}(B_1), \dots, \text{Foc}(B_{|B|})) \cup \Psi_C(\text{Foc}(C_1), \dots, \text{Foc}(C_{|C|})) \\ &\subseteq \text{Foc}(B) \cup \text{Foc}(C) \end{aligned}$$

by the induction hypothesis; moreover,

$$\frac{\vdash \Gamma; B}{\vdash \Gamma; B \oplus C} \oplus_1 \quad \frac{\vdash \Delta; C}{\vdash \Delta; B \oplus C} \oplus_2$$

hence  $\text{Foc}(B) \cup \text{Foc}(C) \subseteq \text{Foc}(B \oplus C)$ , from which the result follows.

- If  $A = 1$ , clearly  $\Psi_1() = \{\emptyset\} \subseteq \text{Foc}(1)$  by the 1 rule.
- If  $A = 0$ , clearly  $\Psi_0() = \emptyset \subseteq \text{Foc}(0)$ .

Since  $A_1, \dots, A_{|A|}$  are negative, we have

$$\begin{aligned}\Psi_A(\{A_1\}^\perp, \dots, \{A_{|A|}\}^\perp)^{\perp\perp} &= \Psi_A(\text{Foc}(A_1), \dots, \text{Foc}(A_{|A|}))^{\perp\perp} \\ &\subseteq \text{Foc}(A)^{\perp\perp} \\ &\subseteq \{A\}^{\perp\perp\perp} = \{A\}^\perp\end{aligned}$$

□

**Lemma 2.** For any formula  $A$  with main negative subformulas  $A_1, \dots, A_{|A|}$ ,  $\llbracket A \rrbracket = \Psi_A(\llbracket A_1 \rrbracket, \dots, \llbracket A_{|A|} \rrbracket)^{\perp\perp}$ .

*Proof.* By induction, using positivity results from [Girard99, appendix F]:  $(X^{\perp\perp} \cdot Y^{\perp\perp})^{\perp\perp} = (X \cdot Y)^{\perp\perp}$  and  $(X^{\perp\perp} \cup Y^{\perp\perp})^{\perp\perp} = (X \cup Y)^{\perp\perp}$ .

- If  $A$  is negative, then  $\llbracket A \rrbracket = \llbracket A \rrbracket^{\perp\perp}$  because  $\llbracket A \rrbracket$  is a fact.
- If  $A = X$ , let  $\Gamma \in \{X\}^\perp$  and  $\Delta \in \{X^\perp\}^\perp$ . We have

$$\frac{\vdash \Gamma, X; \quad \vdash X^\perp, \Delta;}{\vdash \Gamma, \Delta;} \text{ } n\text{-cut}$$

from which  $\{X\}^\perp \subseteq \{X^\perp\}^{\perp\perp}$  follows; moreover  $\{X^\perp\}^{\perp\perp} = \Psi_X()^{\perp\perp} \subseteq \{X\}^\perp$  by lemma 1, therefore  $\llbracket X \rrbracket = \{X\}^\perp = \{X^\perp\}^{\perp\perp}$ .

- If  $A = B \otimes C$ , then

$$\begin{aligned}\llbracket B \otimes C \rrbracket &= (\llbracket B \rrbracket \cdot \llbracket C \rrbracket)^{\perp\perp} \\ &= (\Psi_B(\llbracket B_1 \rrbracket, \dots, \llbracket B_{|B|} \rrbracket)^{\perp\perp} \cdot \Psi_C(\llbracket C_1 \rrbracket, \dots, \llbracket C_{|C|} \rrbracket)^{\perp\perp})^{\perp\perp} \\ &= (\Psi_B(\llbracket B_1 \rrbracket, \dots, \llbracket B_{|B|} \rrbracket) \cdot \Psi_C(\llbracket C_1 \rrbracket, \dots, \llbracket C_{|C|} \rrbracket))^{\perp\perp} \\ &= \Psi_{B \otimes C}(\llbracket B_1 \rrbracket, \dots, \llbracket B_{|B|} \rrbracket, \llbracket C_1 \rrbracket, \dots, \llbracket C_{|C|} \rrbracket)^{\perp\perp}\end{aligned}$$

- If  $A = B \oplus C$ , then

$$\begin{aligned}\llbracket B \oplus C \rrbracket &= (\llbracket B \rrbracket \cup \llbracket C \rrbracket)^{\perp\perp} \\ &= (\Psi_B(\llbracket B_1 \rrbracket, \dots, \llbracket B_{|B|} \rrbracket)^{\perp\perp} \cup \Psi_C(\llbracket C_1 \rrbracket, \dots, \llbracket C_{|C|} \rrbracket)^{\perp\perp})^{\perp\perp} \\ &= (\Psi_B(\llbracket B_1 \rrbracket, \dots, \llbracket B_{|B|} \rrbracket) \cup \Psi_C(\llbracket C_1 \rrbracket, \dots, \llbracket C_{|C|} \rrbracket))^{\perp\perp} \\ &= \Psi_{B \oplus C}(\llbracket B_1 \rrbracket, \dots, \llbracket B_{|B|} \rrbracket, \llbracket C_1 \rrbracket, \dots, \llbracket C_{|C|} \rrbracket)^{\perp\perp}\end{aligned}$$

- If  $A = 1$  then  $\llbracket 1 \rrbracket = \{\emptyset\}^{\perp\perp}$  by definition.
- If  $A = 0$  then  $\llbracket 0 \rrbracket = \emptyset^{\perp\perp}$  by definition.

□

**Lemma 3.** For any formula  $A$ ,  $\llbracket A \rrbracket \subseteq \{A\}^\perp$ .

*Proof.* By induction:

- If  $A = X^\perp$ , we have  $\{X^\perp\} \subseteq \text{Foc}(X) \subseteq \{X\}^\perp$ , therefore  $\llbracket X^\perp \rrbracket = \llbracket X \rrbracket^\perp = \{X\}^{\perp\perp} \subseteq \{X^\perp\}^\perp$ .
- If  $A = B \& C$ , we have  $\llbracket B \& C \rrbracket = \llbracket B \rrbracket \cap \llbracket C \rrbracket \subseteq \{B\}^\perp \cap \{C\}^\perp$  by the induction hypothesis; moreover,

$$\frac{\vdash \Gamma, B; \quad \vdash \Gamma, C;}{\vdash \Gamma, B \& C;} \&$$

hence  $\{B\}^\perp \cap \{C\}^\perp \subseteq \{B \& C\}^\perp$ , from which the result follows.

- If  $A = B \wp C$ , let  $\Gamma \in \llbracket B \wp C \rrbracket = (\llbracket B \rrbracket^\perp \cdot \llbracket C \rrbracket^\perp)^\perp$ . By the induction hypothesis,  $\llbracket B \rrbracket \subseteq \{B\}^\perp$ , hence  $B \in \{B\}^{\perp\perp} \subseteq \llbracket B \rrbracket^\perp$ , and similarly  $C \in \llbracket C \rrbracket^\perp$ , therefore  $\vdash B, C, \Gamma; .$  Moreover,

$$\frac{\vdash \Gamma, B, C;}{\vdash \Gamma, B \wp C; } \wp$$

hence  $\Gamma \in \{B \wp C\}^\perp$ , therefore  $\llbracket B \wp C \rrbracket \subseteq \{B \wp C\}^\perp$ .

- If  $A = \top$ , we have  $\llbracket \top \rrbracket = M = \{\top\}^\perp$  by the  $\top$  rule.
- If  $A = \perp$ , we have  $\llbracket \perp \rrbracket = \perp \subseteq \{\perp\}^\perp$  by the  $\perp$  rule.
- Otherwise,  $A$  is a positive formula with main negative subformulas  $A_1, \dots, A_{|A|}$ . Then,

$$\begin{aligned} \llbracket A \rrbracket &= \Psi_A(\llbracket A_1 \rrbracket, \dots, \llbracket A_{|A|} \rrbracket)^{\perp\perp} && \text{by lemma 2} \\ &\subseteq \Psi_A(\{A_1\}^\perp, \dots, \{A_{|A|}\}^\perp)^{\perp\perp} && \text{by the induction hypothesis and monotonicity of } \Psi_A \\ &\subseteq \{A\}^\perp && \text{by lemma 1} \end{aligned}$$

□

**Corollary 3.1.** *For any multiset of formulas  $\Gamma = A_1, \dots, A_n$ ,  $\llbracket \Gamma \rrbracket \subseteq \{\Gamma\}^\perp$ .*

*Proof.* By lemma 3, we have  $\llbracket A_i \rrbracket \subseteq \{A_i\}^\perp$  for all  $1 \leq i \leq n$ , hence  $\{A_i\} \subseteq \{A_i\}^{\perp\perp} \subseteq \llbracket A_i \rrbracket^\perp$ , therefore  $\{\Gamma\} = \{A_1\} \cdots \{A_n\} \subseteq \llbracket A_1 \rrbracket^\perp \cdots \llbracket A_n \rrbracket^\perp$ .

Thus,  $\llbracket \Gamma \rrbracket = \llbracket A_1 \rrbracket \wp \cdots \wp \llbracket A_n \rrbracket = (\llbracket A_1 \rrbracket^\perp \cdots \llbracket A_n \rrbracket^\perp)^\perp \subseteq \{\Gamma\}^\perp$ . □

**Theorem 4** (Focalised completeness). *If a sequent  $\vdash \Gamma$  of MALL is valid in all phase models, then  $\vdash \Gamma$  has a focalised proof.*

*Proof.* In particular  $\emptyset \in \llbracket \Gamma \rrbracket$ , hence  $\emptyset \in \{\Gamma\}^\perp$  by corollary 3.1, therefore there is a proof  $\pi$  of  $\vdash \Gamma$ ; in  $\mathbf{LL}_{\text{foc}}$ . Then, using [Laurent04, section 3.2], we get a proof  $\pi'$  of  $\vdash \Gamma$ ; in  $\mathbf{LL}_{\text{Foc}}$ . Finally, by [Laurent04, proposition 2],  $\pi'^{\circ}$  is a cut-free, focalised proof of  $\vdash \Gamma$  in  $\mathbf{LL}$ . □

Combining this with the soundness theorem for phase models, we get:

**Corollary 4.1** (Focalisation). *Every provable sequent  $\vdash \Gamma$  of MALL has a focalised proof.*

## References

- [Laurent04] Olivier Laurent. ‘A proof of the focalization property of linear logic’. Apr. 2004. URL: <https://web.archive.org/web/20210225023814/https://perso.ens-lyon.fr/olivier.laurent/llfoc.pdf>.
- [Girard99] Jean-Yves Girard. ‘On the Meaning of Logical Rules I: Syntax Versus Semantics’. In: *Computational Logic*. Springer Berlin Heidelberg, 1999, pp. 215–272. ISBN: 978-3-642-58622-4.