

1 A proof of focalisation

Let $\vdash \Gamma; \Pi$ denote a sequent of LL_{foc} as defined in [Laurent04].

We define the focalised syntactic phase model as (M, I, \perp, φ) where M is the free commutative monoid over formulas of LL with $?A$ and $?A, ?A$ identified, $I = \{?\Gamma \mid \Gamma \in M\}$, $\perp = \{\Gamma \in M \mid \vdash \Gamma; \}$, and $\varphi(X) = \{X\}^\perp$ for positive atoms X . Let $\llbracket A \rrbracket$ be the interpretation of a formula A in this model.

Note that $I \subseteq 1$ because of the $?w$ rule, so $\llbracket !A \rrbracket = (\llbracket A \rrbracket \cap I)^{\perp\perp}$.

To simplify the notation, let $\vdash \Gamma; N$ mean $\vdash \Gamma, N$; when N is a negative formula. Let $\text{Foc}(A) = \{\Gamma \in M \mid \vdash \Gamma; A\}$. Clearly $\text{Foc}(A) \subseteq \{A\}^\perp$ by the *foc* rule, and in particular $\text{Foc}(N) = \{N\}^\perp$ for N negative.

We use the decomposition of exponential connectives alluded to in [Laurent04, section 4.1]:

$$!A = \downarrow \sharp A \qquad ?A = \uparrow \flat A$$

We extend our definitions to this decomposition with $\llbracket \sharp A \rrbracket = \llbracket A \rrbracket$ and $\text{Foc}(\sharp A) = \{A\}^\perp$, which is all that is needed here; this ensures the important property that $\text{Foc}(N) = \text{Foc}(N)^{\perp\perp}$ for N negative.

For a formula A , let $|A|$ denote the number of main negative subformulas in A (where $\sharp B$ is the main negative subformula in $!B$). Define Ψ_A as an $|A|$ -ary monotonous operator on $\mathcal{P}(M)$ by induction:

- $\Psi_N(N_1) = N_1$ if N is negative
- $\Psi_X() = \{X^\perp\}$
- $\Psi_{B \otimes C}(B_1, \dots, B_{|B|}, C_1, \dots, C_{|C|}) = \Psi_B(B_1, \dots, B_{|B|}) \cdot \Psi_C(C_1, \dots, C_{|C|})$
- $\Psi_{B \oplus C}(B_1, \dots, B_{|B|}, C_1, \dots, C_{|C|}) = \Psi_B(B_1, \dots, B_{|B|}) \cup \Psi_C(C_1, \dots, C_{|C|})$
- $\Psi_1() = \{\emptyset\}$
- $\Psi_0() = \emptyset$
- $\Psi_{!B}(B_1) = B_1 \cap I$

Lemma 1. For any formula A with main negative subformulas $A_1, \dots, A_{|A|}$,

$$\llbracket A \rrbracket \subseteq \Psi_A(\llbracket A_1 \rrbracket, \dots, \llbracket A_{|A|} \rrbracket)^{\perp\perp}$$

Proof. By induction, using positivity results from [Girard99, appendix F]: $(X^{\perp\perp} \cdot Y^{\perp\perp})^{\perp\perp} \subseteq (X \cdot Y)^{\perp\perp}$ and $(X^{\perp\perp} \cup Y^{\perp\perp})^{\perp\perp} \subseteq (X \cup Y)^{\perp\perp}$.

- If A is negative, then we have $\llbracket A \rrbracket = \llbracket A \rrbracket^{\perp\perp}$ because $\llbracket A \rrbracket$ is a fact.
- If $A = X$, let $\Gamma \in \{X\}^\perp$ and $\Delta \in \{X^\perp\}^\perp$. We have

$$\frac{\vdash \Gamma, X; \quad \vdash X^\perp, \Delta;}{\vdash \Gamma, \Delta;} \quad n\text{-cut}$$

from which $\{X\}^\perp \subseteq \{X^\perp\}^{\perp\perp}$ follows; therefore $\llbracket X \rrbracket = \{X\}^\perp \subseteq \{X^\perp\}^{\perp\perp} = \Psi_X()^{\perp\perp}$.

- If $A = B \otimes C$, then

$$\begin{aligned} \llbracket B \otimes C \rrbracket &= (\llbracket B \rrbracket \cdot \llbracket C \rrbracket)^{\perp\perp} \\ &\subseteq (\Psi_B(\llbracket B_1 \rrbracket, \dots, \llbracket B_{|B|} \rrbracket)^{\perp\perp} \cdot \Psi_C(\llbracket C_1 \rrbracket, \dots, \llbracket C_{|C|} \rrbracket)^{\perp\perp})^{\perp\perp} && \text{by the induction hypothesis} \\ &\subseteq (\Psi_B(\llbracket B_1 \rrbracket, \dots, \llbracket B_{|B|} \rrbracket) \cdot \Psi_C(\llbracket C_1 \rrbracket, \dots, \llbracket C_{|C|} \rrbracket))^{\perp\perp} && \text{by positivity} \\ &= \Psi_{B \otimes C}(\llbracket B_1 \rrbracket, \dots, \llbracket B_{|B|} \rrbracket, \llbracket C_1 \rrbracket, \dots, \llbracket C_{|C|} \rrbracket)^{\perp\perp} \end{aligned}$$

- If $A = B \oplus C$, then

$$\begin{aligned}
\llbracket B \oplus C \rrbracket &= (\llbracket B \rrbracket \cup \llbracket C \rrbracket)^{\perp\perp} \\
&\subseteq (\Psi_B(\llbracket B_1 \rrbracket, \dots, \llbracket B_{|B|} \rrbracket)^{\perp\perp} \cup \Psi_C(\llbracket C_1 \rrbracket, \dots, \llbracket C_{|C|} \rrbracket)^{\perp\perp})^{\perp\perp} && \text{by the induction hypothesis} \\
&\subseteq (\Psi_B(\llbracket B_1 \rrbracket, \dots, \llbracket B_{|B|} \rrbracket) \cup \Psi_C(\llbracket C_1 \rrbracket, \dots, \llbracket C_{|C|} \rrbracket))^{\perp\perp} && \text{by positivity} \\
&= \Psi_{B \oplus C}(\llbracket B_1 \rrbracket, \dots, \llbracket B_{|B|} \rrbracket, \llbracket C_1 \rrbracket, \dots, \llbracket C_{|C|} \rrbracket)^{\perp\perp}
\end{aligned}$$

- If $A = 1$ then $\llbracket 1 \rrbracket = \{\emptyset\}^{\perp\perp}$ by definition.
- If $A = 0$ then $\llbracket 0 \rrbracket = \emptyset^{\perp\perp}$ by definition.
- If $A = !B$ then $\llbracket !B \rrbracket = (\llbracket B \rrbracket \cap I)^{\perp\perp} = \Psi_{!B}(\llbracket \sharp B \rrbracket)^{\perp\perp}$.

□

Lemma 2. For any formula A with main negative subformulas $A_1, \dots, A_{|A|}$,

$$\Psi_A(\text{Foc}(A_1), \dots, \text{Foc}(A_{|A|})) \subseteq \text{Foc}(A)$$

Proof. By induction:

- If A is negative, $\Psi_A(\text{Foc}(A)) = \text{Foc}(A)$ by definition.
- If $A = X$, then $\Psi_X() = \{X^\perp\} \subseteq \text{Foc}(X)$ by the *ax* rule.
- If $A = B \otimes C$, then we have

$$\begin{aligned}
\Psi_A(\text{Foc}(A_1), \dots, \text{Foc}(A_{|A|})) &= \Psi_B(\text{Foc}(B_1), \dots, \text{Foc}(B_{|B|})) \cdot \Psi_C(\text{Foc}(C_1), \dots, \text{Foc}(C_{|C|})) \\
&\subseteq \text{Foc}(B) \cdot \text{Foc}(C)
\end{aligned}$$

by the induction hypothesis; moreover,

$$\frac{\vdash \Gamma; B \quad \vdash \Delta; C}{\vdash \Gamma, \Delta; B \otimes C} \otimes$$

hence $\text{Foc}(B) \cdot \text{Foc}(C) \subseteq \text{Foc}(B \otimes C)$, from which the result follows.

- If $A = B \oplus C$, then we have

$$\begin{aligned}
\Psi_A(\text{Foc}(A_1), \dots, \text{Foc}(A_{|A|})) &= \Psi_B(\text{Foc}(B_1), \dots, \text{Foc}(B_{|B|})) \cup \Psi_C(\text{Foc}(C_1), \dots, \text{Foc}(C_{|C|})) \\
&\subseteq \text{Foc}(B) \cup \text{Foc}(C)
\end{aligned}$$

by the induction hypothesis; moreover,

$$\frac{\vdash \Gamma; B}{\vdash \Gamma; B \oplus C} \oplus_1 \quad \frac{\vdash \Delta; C}{\vdash \Delta; B \oplus C} \oplus_2$$

hence $\text{Foc}(B) \cup \text{Foc}(C) \subseteq \text{Foc}(B \oplus C)$, from which the result follows.

- If $A = 1$, clearly $\Psi_1() = \{\emptyset\} \subseteq \text{Foc}(1)$ by the 1 rule.
- If $A = 0$, clearly $\Psi_0() = \emptyset \subseteq \text{Foc}(0)$.

- If $A = !B$, then $\Psi_{!B}(\text{Foc}(\sharp B)) = \{B\}^\perp \cap I$, and

$$\frac{\vdash ?\Gamma, B;}{\vdash ?\Gamma; !B} !$$

hence $\{B\}^\perp \cap I \subseteq \text{Foc}(!B)$, from which the result follows.

□

Lemma 3. For any formula A , $\llbracket A \rrbracket \subseteq \text{Foc}(A)^{\perp\perp}$.

Proof. By induction:

- If A is a positive formula with main negative subformulas $A_1, \dots, A_{|A|}$, then

$$\begin{aligned} \llbracket A \rrbracket &\subseteq \Psi_A(\llbracket A_1 \rrbracket, \dots, \llbracket A_{|A|} \rrbracket)^{\perp\perp} && \text{by lemma 1} \\ &\subseteq \Psi_A(\text{Foc}(A_1)^{\perp\perp}, \dots, \text{Foc}(A_{|A|})^{\perp\perp})^{\perp\perp} && \text{by the induction hypothesis} \\ &= \Psi_A(\text{Foc}(A_1), \dots, \text{Foc}(A_{|A|}))^{\perp\perp} && \text{because } A_1, \dots, A_{|A|} \text{ are negative} \\ &\subseteq \text{Foc}(A)^{\perp\perp} && \text{by lemma 2} \end{aligned}$$

- If $A = \sharp B$, then by the induction hypothesis $\llbracket \sharp B \rrbracket = \llbracket B \rrbracket \subseteq \text{Foc}(B)^{\perp\perp} \subseteq \{B\}^\perp = \text{Foc}(\sharp B)^{\perp\perp}$.

Otherwise, it is enough to prove $\llbracket A \rrbracket \subseteq \{A\}^\perp$.

- If $A = X^\perp$, we have $\{X\} \subseteq \{X^\perp\}^\perp$, therefore $\llbracket X^\perp \rrbracket = \{X\}^{\perp\perp} \subseteq \{X^\perp\}^{\perp\perp\perp} = \{X^\perp\}^\perp$.
- If $A = B \& C$, we have $\llbracket B \& C \rrbracket = \llbracket B \rrbracket \cap \llbracket C \rrbracket \subseteq \{B\}^\perp \cap \{C\}^\perp$ by the induction hypothesis; moreover,

$$\frac{\vdash \Gamma, B; \quad \vdash \Gamma, C;}{\vdash \Gamma, B \& C;} \&$$

hence $\{B\}^\perp \cap \{C\}^\perp \subseteq \{B \& C\}^\perp$, from which the result follows.

- If $A = B \wp C$, let $\Gamma \in \llbracket B \wp C \rrbracket = (\llbracket B \rrbracket^\perp \cdot \llbracket C \rrbracket^\perp)^\perp$. By the induction hypothesis, $\llbracket B \rrbracket \subseteq \{B\}^\perp$, hence $B \in \{B\}^{\perp\perp} \subseteq \llbracket B \rrbracket^\perp$, and similarly $C \in \llbracket C \rrbracket^\perp$, therefore $\vdash B, C, \Gamma; .$ Moreover,

$$\frac{\vdash \Gamma, B, C;}{\vdash \Gamma, B \wp C;} \wp$$

hence $\Gamma \in \{B \wp C\}^\perp$, from which the result follows.

- If $A = \top$, we have $\llbracket \top \rrbracket = M = \{\top\}^\perp$ by the \top rule.
- If $A = \perp$, we have $\llbracket \perp \rrbracket = \perp \subseteq \{\perp\}^\perp$ by the \perp rule.
- If $A = ?B$, then $\llbracket ?B \rrbracket = (\llbracket B \rrbracket^\perp \cap I)^\perp$. By the induction hypothesis, $\llbracket B \rrbracket \subseteq \text{Foc}(B)^{\perp\perp}$, hence $\text{Foc}(B)^\perp \subseteq \llbracket B \rrbracket^\perp$ and $(\llbracket B \rrbracket^\perp \cap I)^\perp \subseteq (\text{Foc}(B)^\perp \cap I)^\perp$. Moreover, $?B \in \text{Foc}(B)^\perp \cap I$ because of the $?d$ rule, therefore $(\text{Foc}(B)^\perp \cap I)^\perp \subseteq \{?B\}^\perp$, from which the result follows.

□

Corollary 3.1. For any multiset of formulas $\Gamma = A_1, \dots, A_n$, $\llbracket \Gamma \rrbracket \subseteq \{\Gamma\}^\perp$.

Proof. By lemma 3, we have $\llbracket A_i \rrbracket \subseteq \text{Foc}(A_i)^{\perp\perp} \subseteq \{A_i\}^\perp$ for all $1 \leq i \leq n$, hence $\{A_i\} \subseteq \{A_i\}^{\perp\perp} \subseteq \llbracket A_i \rrbracket^\perp$, therefore $\{\Gamma\} = \{A_1\} \cdots \{A_n\} \subseteq \llbracket A_1 \rrbracket^\perp \cdots \llbracket A_n \rrbracket^\perp$.

Thus, $\llbracket \Gamma \rrbracket = \llbracket A_1 \rrbracket \wp \cdots \wp \llbracket A_n \rrbracket = (\llbracket A_1 \rrbracket^\perp \cdots \llbracket A_n \rrbracket^\perp)^\perp \subseteq \{\Gamma\}^\perp$. \square

Theorem 4 (Focalised completeness). *If a sequent $\vdash \Gamma$ of LL is valid in all phase models, then $\vdash \Gamma$ has a focalised proof.*

Proof. In particular $\emptyset \in \llbracket \Gamma \rrbracket$, hence $\emptyset \in \{\Gamma\}^\perp$ by corollary 3.1, therefore there is a proof π of $\vdash \Gamma$; in LL_{foc} . Then, using [Laurent04, section 3.2], we get a proof π' of $\vdash \Gamma$; in LL_{Foc} . Finally, by [Laurent04, proposition 2], π'^{lo} is a cut-free, focalised proof of $\vdash \Gamma$ in LL. \square

Combining this with the soundness theorem for phase models, we get:

Corollary 4.1 (Focalisation). *Every provable sequent $\vdash \Gamma$ of LL has a focalised proof.*

References

- [Laurent04] Olivier Laurent. ‘A proof of the focalization property of linear logic’. Apr. 2004. URL: <https://web.archive.org/web/20210225023814/https://perso.ens-lyon.fr/olivier.laurent/llfoc.pdf>.
- [Girard99] Jean-Yves Girard. ‘On the Meaning of Logical Rules I: Syntax Versus Semantics’. In: *Computational Logic*. Springer Berlin Heidelberg, 1999, pp. 215–272. ISBN: 978-3-642-58622-4.