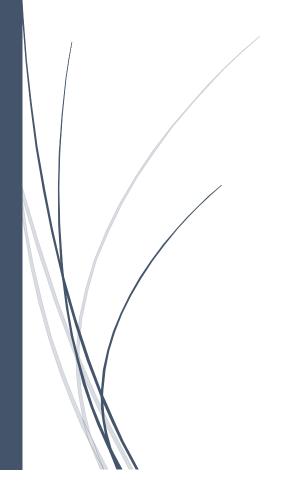
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# Hash Function and Public Key Cryptography

CSCE 465 Homework 5



Nicholas Charchenko

### **Book Problems**

- 1. 5.3: It should take around 2^33 attempts to find a collision. We know that the probability that a given Type 2 message has the same message digest as one of the 2^32 Type 1 messages is roughly 2^32/2^64, equivalent to 1/2^32, so it is likely that one of the 2^32 Type 2 messages matches one of the 2^32 Type 1 messages by the birthday problem.
- 2. 5.4: The digest size will remain 64 bits. Therefore, the iterations will remain the same, and if a palindrome exists within the message space, then the message will be the same forward and backward, leading to a collision when the two messages are hashed.
- 3. 5.14: For ~x, if x is random, then the output will be random. For x XOR y, if either x or y is random, then the output will be random. For x OR y, a fixed 1 bit will lead to non-random output, while a fixed 0 bit or both x and y being random will lead to a random output. For x AND y, the output will be random if there is either a fixed 1 bit or if both x and y are random. For the selection function, if any two of x,y, and z are random, then the output will be random. For the majority function, if any two are random, then the output will be random. For x XOR y XOR z, at least two of x,y, and z must be random. For the last function, if either y, or both x and z are random, then the output will be random. Note: It is assumed that if any fixed variable can tamper with the results, then the result is not truly random.
- 4. 6.2: The attacker will not be able to decrypt the Diffie-Hellman values sent to him and so will not be able to compute the shared secrets.
- 5. 6.8: We can compute  $m_1^j \mod n$  for any positive integer j by this formula:  $\left(m_1^d\right)^j \mod n = \left(m_1^j\right)^d \mod n$ . We apply this principle and then compute the signature of  $m_1^x$  using the fact that the signature of  $m_1^{-1}$  is  $(m_1^{-1})^d \mod n = (m_1^{-d}) \mod n = (m_1^d)^{-1} \mod n$ . We can similarly compute the signature of  $m_2^k$  for any k. Now we know how to compute the signatures of  $m_1^j$  and  $m_2^k$ . We can compute the signature of the product with the formula  $m_1*m_2=(m_1*m_2)^d \mod n = ((m_1)^d \mod n)*((m_2)^d \mod n) \mod n$ . Another property to note is that if j=0, then the signature of  $m_1^j=m_1^{-j}$ .

# Task 1: Generating Message Digest and MAC

For this task, I created a file called example.txt with the sentence "This is an example." Then, I used OpenSSL to implement MD5, SHA-1, and SHA-256 algorithms.

```
[04/11/2018 15:18] seed@ubuntu:~/Desktop$ openssl dgst -md5 example.txt
MD5(example.txt)= 46edc6541babd006bb52223c664b29a3
[04/11/2018 15:18] seed@ubuntu:~/Desktop$ openssl dgst -sha1 example.txt
SHA1(example.txt)= a6f153801c9303d73ca2b43d3be62f44c6b66476
[04/11/2018 15:19] seed@ubuntu:~/Desktop$ openssl dgst -sha256 example.txt
SHA256(example.txt)= c80a97041f15ba166b9a3e8fc2b09726d778bc3bd9338d4befe34b46707
ebeec
```

Figure 1: MD5, SHA-1, and SHA-256 hash outputs

I can observe that MD5 gives a 32-bit hash length, SHA-1 gives a 40-bit hash length, and SHA-256 gives a 64-bit hash length. This proves that MD5 has a 128-bit hash value, SHA-1 has a 160-bit hash value, and SHA-256 has a 256-bit hash value.

# Task 2: Keyed Hash and HMAC

I created a separate file called hmacexample.txt and ran MD5, SHA-1, and SHA-256 hashing algorithms with three different keys: "a", "abcd", and "abcdefghijklmnopqrstuvwxyz" to show keys of varying lengths. The key length does not need to be fixed, and the HMAC lengths are the same as the lengths of the hash codes from Task 1.

```
[04/14/2018 14:35] seed@ubuntu:~/Desktop$ openssl dgst -md5 -hmac "a" hmacexampl
e.txt
HMAC-MD5(hmacexample.txt)= 22e9f0b7c3c0deb895c67ae34ea9dae6
[04/14/2018 14:59] seed@ubuntu:~/Desktop$ openssl dgst -md5 -hmac "abcd" hmacexa
mple.txt
HMAC-MD5(hmacexample.txt)= 404295b1a68aeb9efc083c917ce802df
[04/14/2018 14:59] seed@ubuntu:~/Desktop$ openssl dgst -md5 -hmac "abcdefghijklm
nopqrstuvwxyz" hmacexample.txt
HMAC-MD5(hmacexample.txt)= 4eff411ef91f3a801df0ab4dd48fbb13
[04/14/2018 15:00] seed@ubuntu:~/Desktop$ openssl dgst -sha1 -hmac "a" hmacexamp
le.txt
HMAC-SHA1(hmacexample.txt)= 7b2455572759eeb983a73a6656063fc8c60574b1
[04/14/2018 15:02] seed@ubuntu:~/Desktop$ openssl dgst -sha1 -hmac "abcd" hmacex
HMAC-SHA1(hmacexample.txt)= 14872a0bb8b463c10a6b3dddf102cc3e6aee68c7
[04/14/2018 15:04] seed@ubuntu:~/Desktop$ openssl dgst -sha1 -hmac "abcdefghijkl
mnopgrstuvwxyz" hmacexample.txt
HMAC-SHA1(hmacexample.txt)= 75d247f29199d754dc6bc16970e572246a610cbb
[04/14/2018 15:05] seed@ubuntu:~/Desktop$ openssl dgst -sha256 -hmac "a" hmacexa
mple.txt
HMAC-SHA256(hmacexample.txt)= e64fa62d78708340525cd981815fe26e7d5004804f8ef5796c
89f956e70dd486
[04/14/2018 15:05] seed@ubuntu:~/Desktop$ openssl dgst -sha256 -hmac "abcd" hmac
example.txt
HMAC-SHA256(hmacexample.txt)= eb1803f7a886db56394e62c08875f4d6c973b7fd22b2810ca6
19b4dcc525f3cf
[04/14/2018 15:05] seed@ubuntu:~/Desktop$ openssl dgst -sha256 -hmac "abcdefghij
klmnopqrstuvwxyz" hmacexample.txt
HMAC-SHA256(hmacexample.txt)= 9061dd4106f2e3a76fdd3e0ecef91d06ae21dfee8fdc658b71
b680ab6c194130
[04/14/2018 15:05] seed@ubuntu:~/Desktop$
```

Figure 2: Output for Task 2: Keyed Hash and HMAC

# Task 3: The Randomness of One-Way Hash

For this task, I created a file called task3.txt with the sentence, "This is an example for Task 3." (no quotes). I then did the hashing with MD5 and SHA-256 for H1. For H2, I changed "Task 3" to "task 3" using Ghex and ran the algorithms again. I observe that flipping this one bit causes the entire hash to change for both MD5 and SHA-256 algorithms. I compared the outputs for MD5 and only 1 bit was shared while SHA-256 had 6 shared bits between the two outputs. This shows that there is a strong avalanche effect with hash algorithms, similar to secret key encryption algorithms.

```
[04/14/2018 15:39] seed@ubuntu:~/Desktop$ openssl dgst -md5 task3.txt
MD5(task3.txt)= e92b190b4959481ad1421c158f26f5aa
[04/14/2018 15:39] seed@ubuntu:~/Desktop$ openssl dgst -sha256 task3.txt
SHA256(task3.txt)= 62b78c62465534f1f7e2817c4ce3c0fb032e35aaa4a32aba034d34f71b5cb46a
[04/14/2018 15:39] seed@ubuntu:~/Desktop$ openssl dgst -md5 task3.txt
MD5(task3.txt)= 5c6fa2ad50ccf4d98bf05907edb444a4
[04/14/2018 16:06] seed@ubuntu:~/Desktop$ openssl dgst -sha256 task3.txt
SHA256(task3.txt)= 45489a677d87472c6f7fe9f7b683f0bb73abe0a477e6542b7ae92026881026c9
[04/14/2018 16:06] seed@ubuntu:~/Desktop$ ■
```

Figure 3: Changing one bit of task3.txt makes for a significant change in the hash.

# Task 4: One-Way vs Collision-Free Properties

For this task, I designed a program using C to create experiments to perform attacks against both the one-way hash property and the collision-free hash property. To make it easier, I set it up so that we only had 24-bit values, since otherwise it could take a very long time to brute-force. I used random number generation to create messages to be hashed and then repeatedly made hashes using loops to run 15 trials for breaking the collision property, and 5 trials for breaking the one-way property because breaking the one-way property proved to take much longer, outputting the average number of attempts. I used both SHA-256 and MD5 algorithms in this task. From my observations, it is clear that the one-way property is more difficult to break.

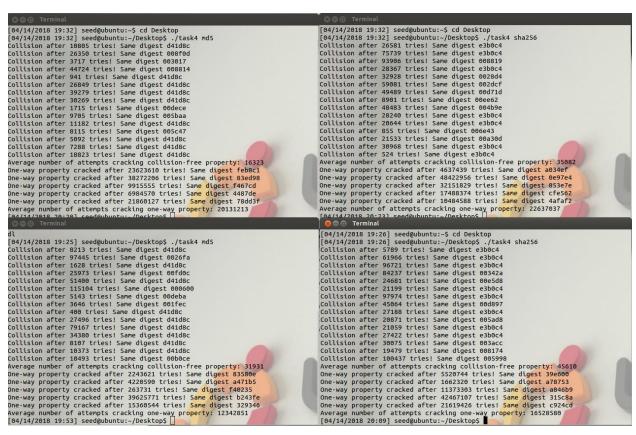


Figure 4: Brute-Force attacks on Collision and One-Way properties of both MD5 and SHA-256

We can think about this in context of the birthday problem. The birthday problem states that the probability that at least one student was born on a specific day is  $1-\left(\frac{364}{365}\right)^n$  for a class of n students. For n=30, this probability is around 7.9%. This is the same as breaking the one-way hash property. This is because breaking the one-way hash property requires that, given a message m and a hash H(m), we find a message m' such that H(m')=H(m), like finding a student born on a specific day. However, it is much easier to find a student that shares a birthday with any other student. This probability is equivalent to  $1-\frac{365!}{(365-n)!*365^n}$ , which is equal to around 70% for n=30. This is the same as breaking the collision hash property.

### Task 5: RSA vs AES

I used the following commands for Task 5:

- openssl genrsa -out privatekey.pem 1024
  - Generate private key
- openssl rsa -in mykey.pem -pubout -out mypubkey.pem
  - Generate public key
- time openssl rsautl -encrypt -in message.txt -pubin -inkey mypubkey.pem -out message.enc
  - Timing RSA encryption
- time openssl rsautl -decrypt -in message.enc -inkey mykey.pem -out message\_decrypted.txt
  - Timing RSA decryption
- time openssl enc -aes-128-cbc -in message.txt -out message.aes \-K a \-iv a
  - Timing AES-128-CBC encryption

The RSA decrypt averaged around 0.004 seconds, the RSA encrypt and AES-128-CBC encryption averaged around 0.003 seconds. Running the speed benchmark on RSA for 1024-bit keys yielded that signing on average took 0.001135 seconds and verifying took on average 0.000053 seconds. For AES speed benchmarks, I got the following results:

```
The 'numbers' are in 1000s of bytes per second processed.

type 16 bytes 64 bytes 256 bytes 1024 bytes 8192 bytes
aes-128 cbc 79778.77k 89718.42k 90975.36k 176637.59k 168611.77k
```

Figure 5: Speed tests on AES-128 CBC encryption

# Task 6: Create a Digital Signature

The following screenshot shows the commands and steps used for Task 6:

```
[04/14/2018 20:14] seed@ubuntu:~$ cd Desktop
[04/14/2018 20:14] seed@ubuntu:~/Desktop$ openssl genrsa -des3 -out myrsaCA.pem
1024
Generating RSA private key, 1024 bit long modulus
unable to write 'random state'
e is 65537 (0x10001)
Enter pass phrase for myrsaCA.pem:
Verifying - Enter pass phrase for myrsaCA.pem:
[04/14/2018 20:14] seed@ubuntu:~/Desktop$ openssl rsa -in myrsaCA.pem -pubout -o
ut myrsapubkev.pem
Enter pass phrase for myrsaCA.pem:
writing RSA key
[04/14/2018 20:17] seed@ubuntu:~/Desktop$ openssl dgst -sha256 -out example.sha2
56 -sign myrsaCA.pem example.txt
Enter pass phrase for myrsaCA.pem:
[04/14/2018 20:17] seed@ubuntu:~/Desktop$ openssl genrsa -des3 -out myrsaCA.pem
1024
Generating RSA private key, 1024 bit long modulus
.....+++++
unable to write 'random state'
e is 65537 (0x10001)
Enter pass phrase for myrsaCA.pem:
Verifying - Enter pass phrase for myrsaCA.pem:
[04/14/2018 20:17] seed@ubuntu:~/Desktop$ openssl dgst -sha256 -signature exampl
e.sha256 -verify myrsapubkey.pem example.txt
Verified OK
[04/14/2018 20:20] seed@ubuntu:~/Desktop$ openssl dgst -sha256 -signature exampl
e.sha256 -verify myrsapubkey.pem example.txt
Verification Failure
[04/14/2018 20:21] seed@ubuntu:~/Desktop$ openssl dgst -sha256 -signature exampl
e.sha256 -verify myrsapubkey.pem example.txt
Verified OK
[04/14/2018 20:33] seed@ubuntu:~/Desktop$ openssl dgst -sha256 -signature exampl
e.sha256 -verify myrsapubkey.pem example.txt
Verification Failure
[04/14/2018 20:33] seed@ubuntu:~/Desktop$
```

Figure 6: Task 6-Creating Digital Signature

I noticed that even the slightest change in example.txt would cause the subsequent verification to fail. This is because I first made a significant change (adding a single word), and got a failure. Then, I added a whitespace character and got a verification failure. To create the RSA public/private key pair, I followed the same process as in Task 5 for generating the key pair.

Digital signatures are important in security because they ensure integrity and non-repudiation. From performing Task 6, the tampered document fails to verify. Using this, we can tell if a document was tampered, and we know not to use it. If the verification succeeds, then we know that the document is legitimate and trustworthy, improving our security. Furthermore, digital signatures help ensure delivery and helps track who interacts with the information.