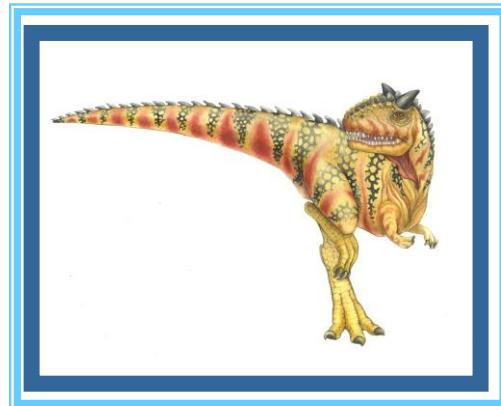


Chapter 6: Deadlocks





Chapter 6: Deadlocks

- The Deadlock Problem
- System Model
- Deadlock Characterization
- Methods for Handling Deadlocks
- Deadlock Prevention
- Deadlock Avoidance
- Deadlock Detection
- Recovery from Deadlock

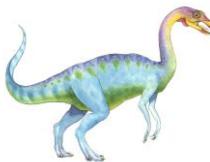




Chapter Objectives

- To develop a description of deadlocks, which prevent sets of concurrent processes from completing their tasks
- To present a number of different methods for preventing or avoiding deadlocks in a computer system





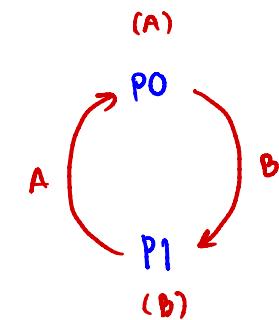
The Deadlock Problem

- A set of blocked processes each holding a resource and waiting to acquire a resource held by another process in the set
- Example
 - System has 2 disk drives
 - P_1 and P_2 each hold one disk drive and each needs another one
- Example
 - semaphores A and B , initialized to 1

P_0
 $A = 1 - 1 = 0$
wait(A);
wait(B);

ନୀତିକାଳୀଙ୍କ କାର୍ଯ୍ୟ

P_1
 $B = 1 - 1 = 0$
wait(B)
wait(A)

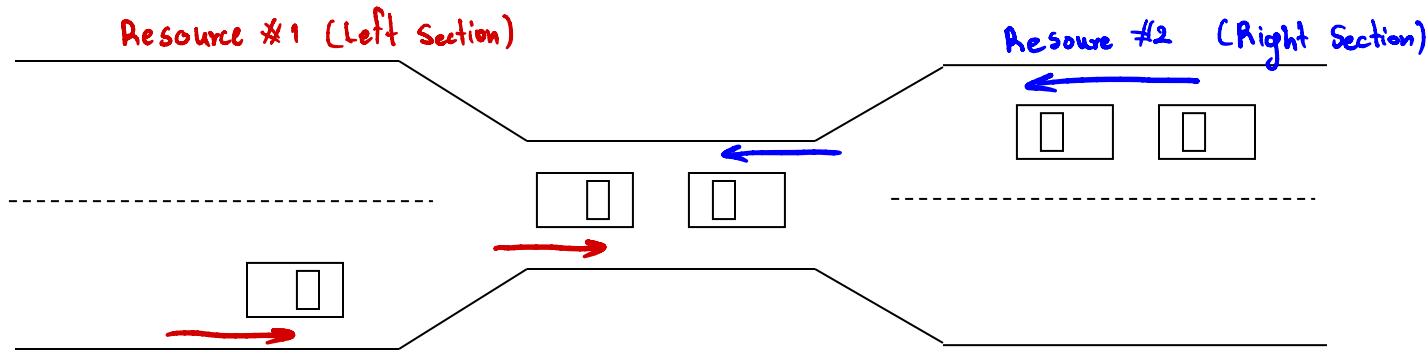


acquire: ଅଧିକରଣକରଣ





Bridge Crossing Example



- Traffic only in one direction
- Each section of a bridge can be viewed as a resource
- If a deadlock occurs, it can be resolved if one car backs up (preempt resources and rollback)
- Several cars may have to be backed up if a deadlock occurs
- Starvation is possible *Ex. Some car may always be preempted and never cross the bridge*
- Note – Most OSes do not prevent or deal with deadlocks





System Model

- ນີ້ກັນຫາກ ເປັນ ທົນດ
ຮົມຕະໜີ 1
ຮົມຕະໜີ 2.
- Resource types R_1, R_2, \dots, R_m
CPU cycles, memory space, I/O devices
 - Each resource type R_i has W_i instances.
 - Each process utilizes a resource as follows:
 - **request**
 - **use**
 - **release**





Deadlock Characterization

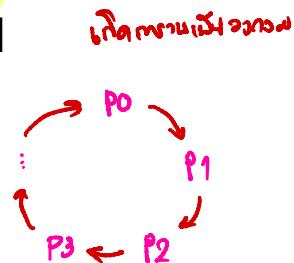
ຕົວລີ້ນ ຂໍເກືອບໄຂພຮອມກັນ ໄສດັບຈະໄລຍ້ເລືດ Deadlock

Deadlock can arise if four conditions hold simultaneously.

- **Mutual exclusion:** only one process at a time can use a resource
- **Hold and wait:** a process holding at least one resource is waiting to acquire additional resources held by other processes
- **No preemption:** a resource can be released only voluntarily by the process holding it, after that process has completed its task
- **Circular wait:** there exists a set $\{P_0, P_1, \dots, P_n\}$ of waiting processes such that P_0 is waiting for a resource that is held by P_1 , P_1 is waiting for a resource that is held by P_2 , ..., P_{n-1} is waiting for a resource that is held by P_n , and P_n is waiting for a resource that is held by P_0 .

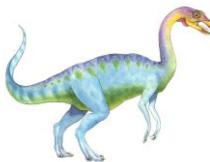
Process wait for
resource

ໃຫຍວະນາ
ຕົວລີ້ນ
Deadlock



simultaneously : ເກີດປຶ້ນໃນເວລາເດືອກັນ

voluntarily : ກະທຳເສົ່ງແລ້ວ = willingly without being forced



Resource-Allocation Graph

ການຝ່າຍອດສ້າງໃນພາສາ

A set of vertices V and a set of edges E .

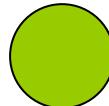
- V is partitioned into two types:
 - $P = \{P_1, P_2, \dots, P_n\}$, the set consisting of all the processes in the system
 - $R = \{R_1, R_2, \dots, R_m\}$, the set consisting of all resource types in the system
- request edge – directed edge $P_i \rightarrow R_j$ P_i requests R_j
- assignment edge – directed edge $R_j \rightarrow P_i$ R_j is assigned (or allocated) to P_i





Resource-Allocation Graph (Cont.)

- Process

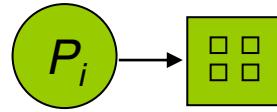


A vertex representing a process

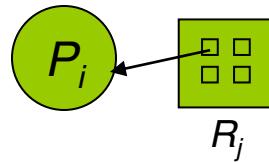
- Resource Type with 4 instances

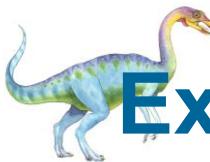


- P_i requests instance of R_j



- P_i is holding an instance of R_j

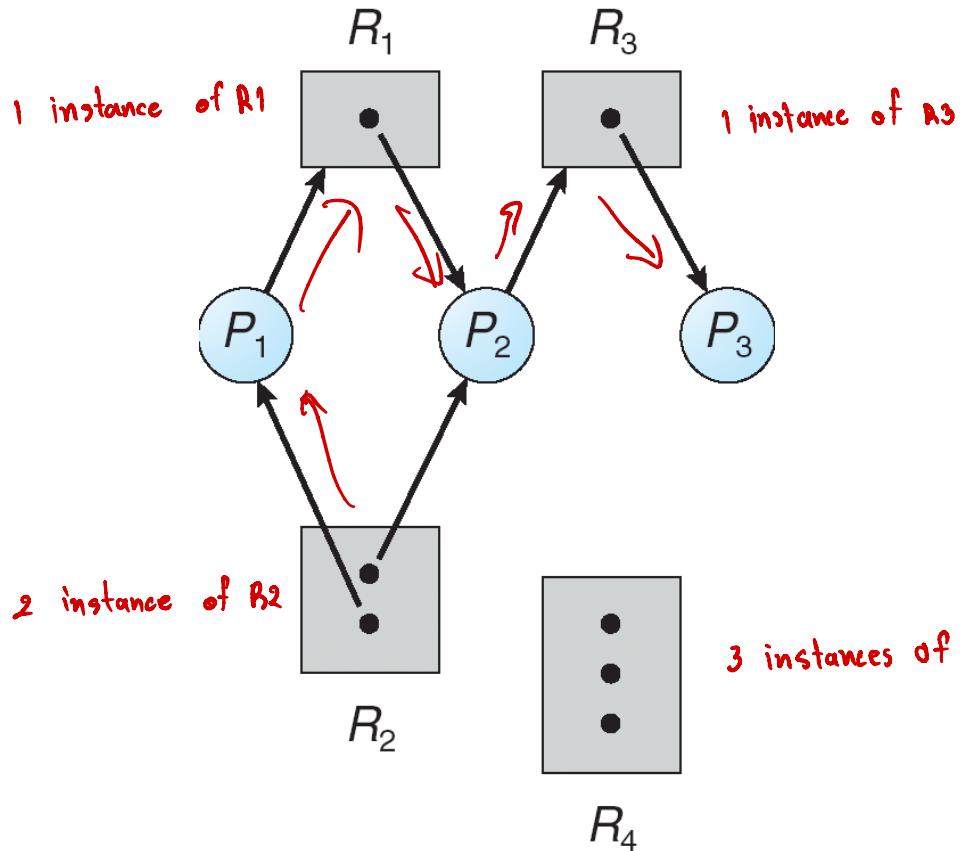




Example of a Resource Allocation Graph

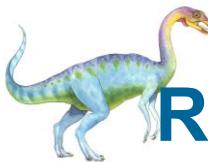
There are 6 edges in the graph

$E = \{ \text{set request \& assignmet} \}$

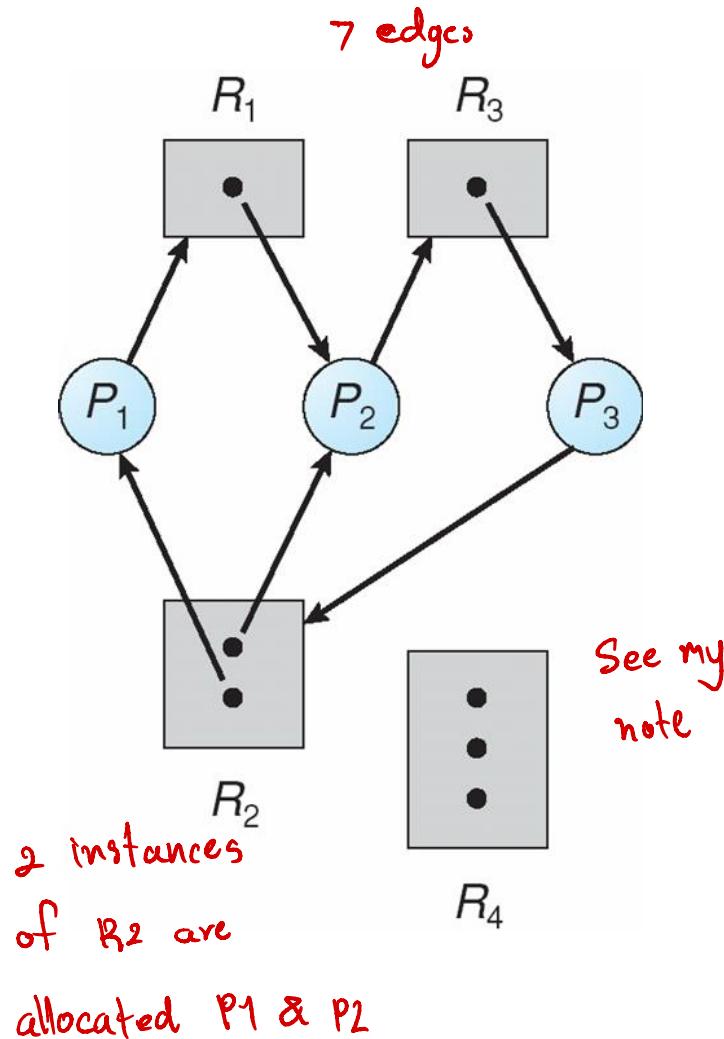


ได้กราฟ $E = \{P_1 \rightarrow R_1,$
 $P_2 \rightarrow R_3,$
 $R_1 \rightarrow P_2,$
 $R_2 \rightarrow P_1,$
 $R_2 \rightarrow P_2,$
 $R_3 \rightarrow P_3\}$





Resource Allocation Graph With A Deadlock



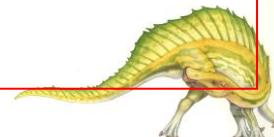
ได้กราฟ **E** คือ ?

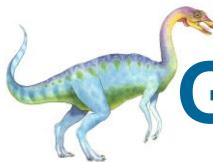
= { $P_1 \rightarrow R_1, P_2 \rightarrow R_3, P_3 \rightarrow R_2,$
 $R_1 \rightarrow P_2, R_2 \rightarrow P_1, R_2 \rightarrow P_3,$
 $R_3 \rightarrow P_3$ }

กราฟนี้มีการเกิดเป็น **Cycle** (วงรอบ) 2 วง ได้แก่
วงที่ 1

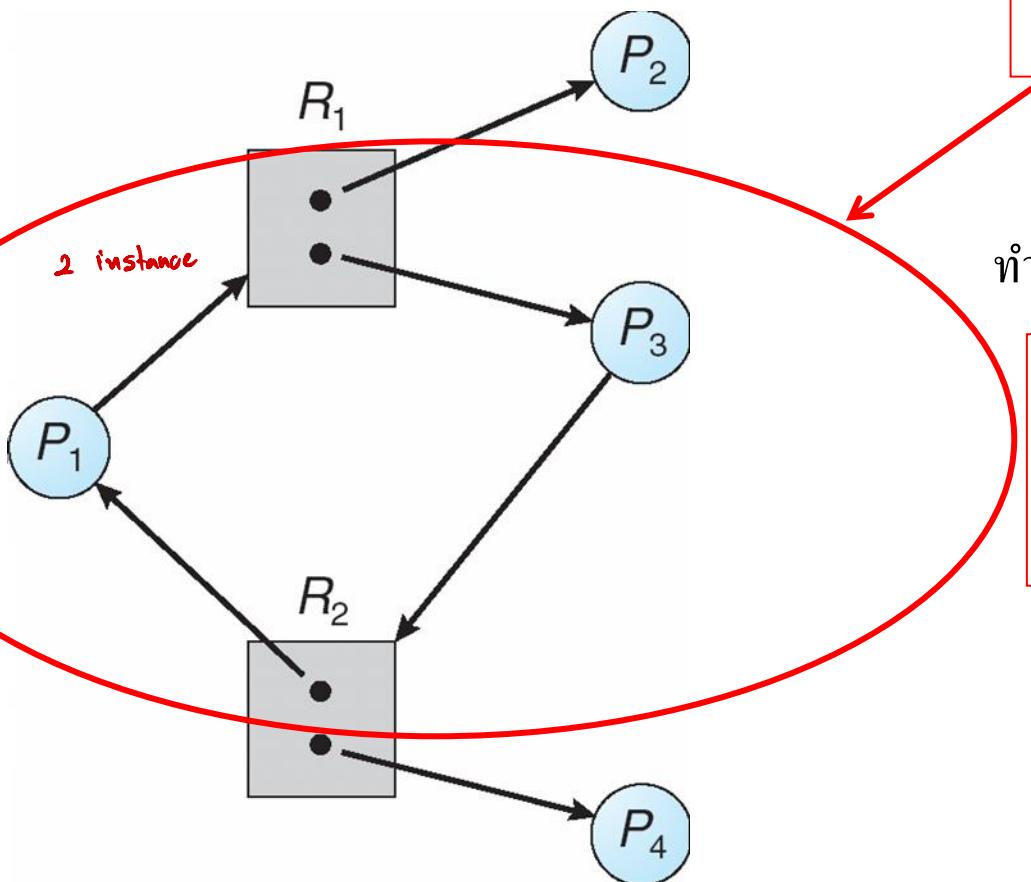
$P_1 \rightarrow R_1 \rightarrow P_2 \rightarrow R_3 \rightarrow P_3 \rightarrow R_2 \rightarrow P_1$

วงที่ 2 = $P_2 \rightarrow R_3 \rightarrow P_3 \rightarrow R_2 \rightarrow P_2$





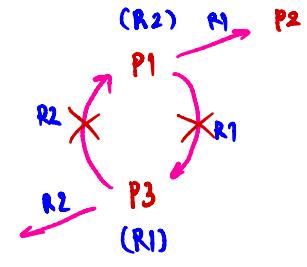
Graph With A Cycle But No Deadlock

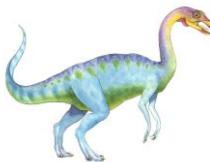


A Cycle

ทำไมถึงไม่เกิด Deadlock ???

หาก P_4 ปล่อย R_2 ให้กับระบบ
 P_3 จะได้ครอบครองทรัพยากร
และไม่เกิด Cycle





Basic Facts

- If graph contains no cycles \Rightarrow no deadlock
- If graph contains a cycle
 - if only one instance per resource type, then deadlock
 - if several instances per resource type, possibility of deadlock





Methods for Handling Deadlocks

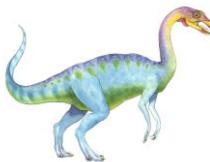
ป้องกัน
หลีกเลี่ยง

The system can use
- a deadlock - prevention schema
- or a deadlock - avoidance schema

- Ensure that the system will *never* enter a deadlock state
(กำหนดเงื่อนไขในการใช้ Resource)
- Allow the system to enter a **deadlock** state and then recover
(เมื่อเกิดปัญหา ตามแก้ไขทีหลัง)
- Ignore the problem and pretend that deadlocks never occur in the system; used by most operating systems, including UNIX

(มองข้ามปัญหา ทำเสมือนไม่มีการเกิด **Deadlock** ในระบบ
**** วิธีนี้เป็น 1 วิธีการที่ใช้ใน **OS** ส่วนใหญ่****)





Deadlock Prevention (ป้องกัน)

พิจารณาถึง การเกิด Deadlock ต้องมีเงื่อนไข ทั้ง 4 กรณีเกิดขึ้นพร้อมกัน

Restrain the ways request can be made

- **Mutual Exclusion** – not required for sharable resources; must hold for nonshareable resources

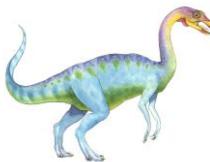
- **Hold and Wait** – must guarantee that whenever a process requests a resource, it does not hold any other resources
 - Require process to request and be allocated all its resources before it begins execution, or allow process to request resources only when the process has none
 - Low resource utilization; starvation possible

All or
nothing

(หาก Resource ว่างแต่ไม่สามารถให้ Process ถือครองได้เวลานานๆ สามารถนำ Resource มาใช้ประโยชน์เมื่อใช้เสร็จต้องรีบคืน เมื่อจะใช้ใหม่ต้อง Request ใหม่)

หากมี Process ต้องการใช้ Resource ที่ได้รับความนิยมมากๆ จะเกิด Starvation)





Deadlock Prevention (Cont.)

□ **No Preemption –**

- If a process that is holding some resources requests another resource that cannot be immediately allocated to it, then all resources currently being held are released
- Preempted resources are added to the list of resources for which the process is waiting
- Process will be restarted only when it can regain its old resources, as well as the new ones that it is requesting

□ **Circular Wait –** impose a total ordering of all resource types, and require that each process requests resources in an increasing order of enumeration

Impose : กำหนด





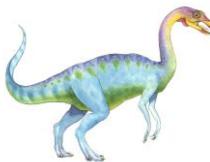
Deadlock Avoidance

Requires that the system has some additional *a priori* information available.
using fact that are then

- Simplest and most useful model requires that each process declare the *maximum number of resources* of each type that it may need.
- The deadlock-avoidance algorithm dynamically examines the resource-allocation state to ensure that there can never be a circular-wait condition.
- Resource-allocation state is defined by the number of available and allocated resources, and the *maximum demands* of the processes.

priori : บางส่วนก่อนหน้า





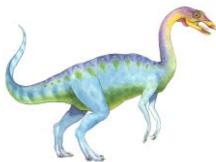
Safe State

QUESTION

- When a process requests an available resource, system must decide if immediate allocation leaves the system in a *safe state*.
- System is in *safe state* if there exists a *safe sequence* of all processes.
- Sequence $\langle P_1, P_2, \dots, P_n \rangle$ is safe if for each P_i , the resources that P_i can still request can be satisfied by currently available resources + resources held by all the P_j , with $j < i$.
 - If P_i resource needs are not immediately available, then P_i can wait until all P_j have finished.
 - When P_j is finished, P_i can obtain needed resources, execute, return allocated resources, and terminate.
 - When P_i terminates, P_{i+1} can obtain its needed resources, and so on.

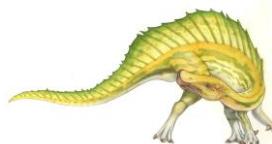
P_j will release its hold resources soon thus P_i gain them after that





Basic Facts

- If a system is in safe state \Rightarrow no deadlocks.
- If a system is in unsafe state \Rightarrow possibility of deadlock.
- Avoidance \Rightarrow ensure that a system will never enter an unsafe state.

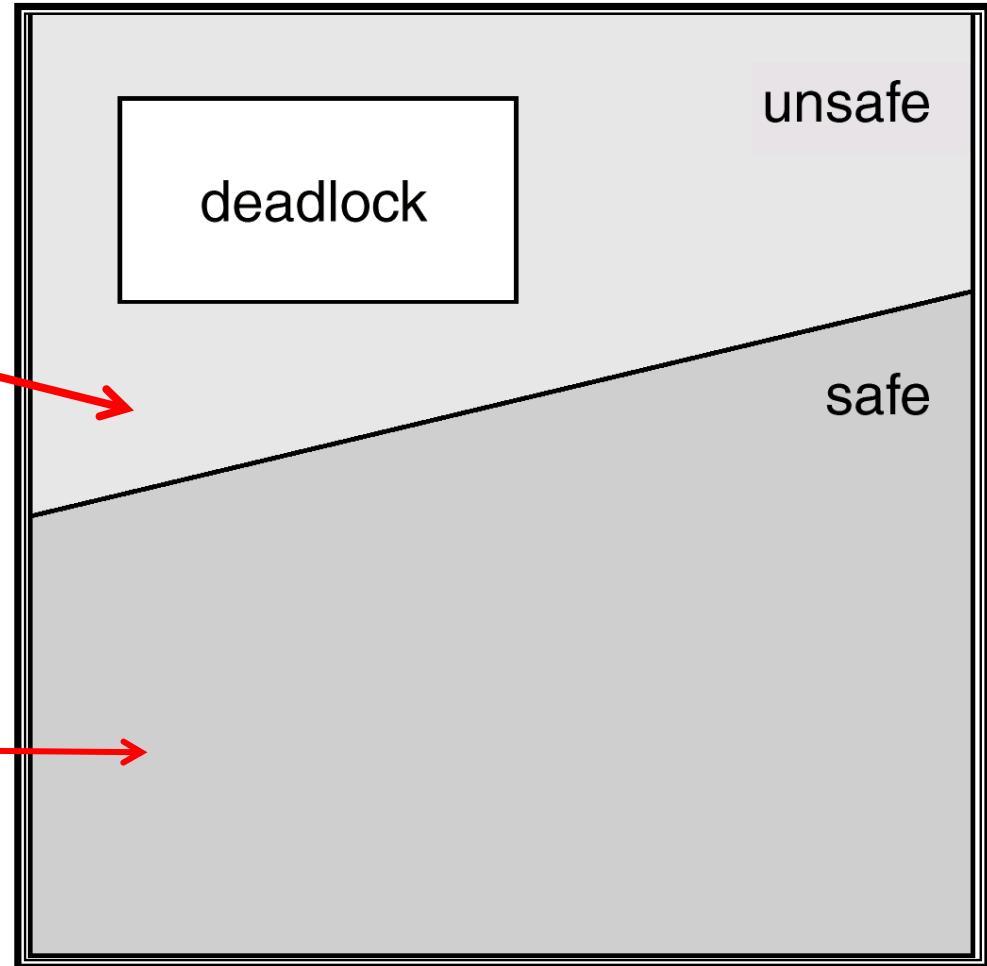




Safe, Unsafe , Deadlock State

อยู่ในสถานะ **unsafe** อาจจะไม่เกิด
Deadlock ได้

อยู่ในสถานะ **safe** ไม่เกิด
Deadlock แน่นอน





Avoidance algorithms

- Single instance of a resource type
 - Use a resource-allocation graph *nwnfnwmsr*

- Multiple instances of a resource type
 - Use the banker's algorithm





Resource-Allocation Graph Scheme

For deadlock avoidance with ONE instance of each resource type.

P_j

- Claim edge $P_i \rightarrow R_j$ indicated that process P_j may request resource R_i ; represented by a dashed line
- Claim edge converts to request edge when a process requests a resource
- Request edge converted to an assignment edge when the resource is allocated to the process
- When a resource is released by a process, assignment edge reconverts to a claim edge
- Resources must be claimed *a priori* in the system

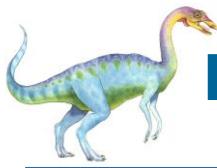
$P_i \rightarrow R_j$

$P_i \rightarrow R_i$

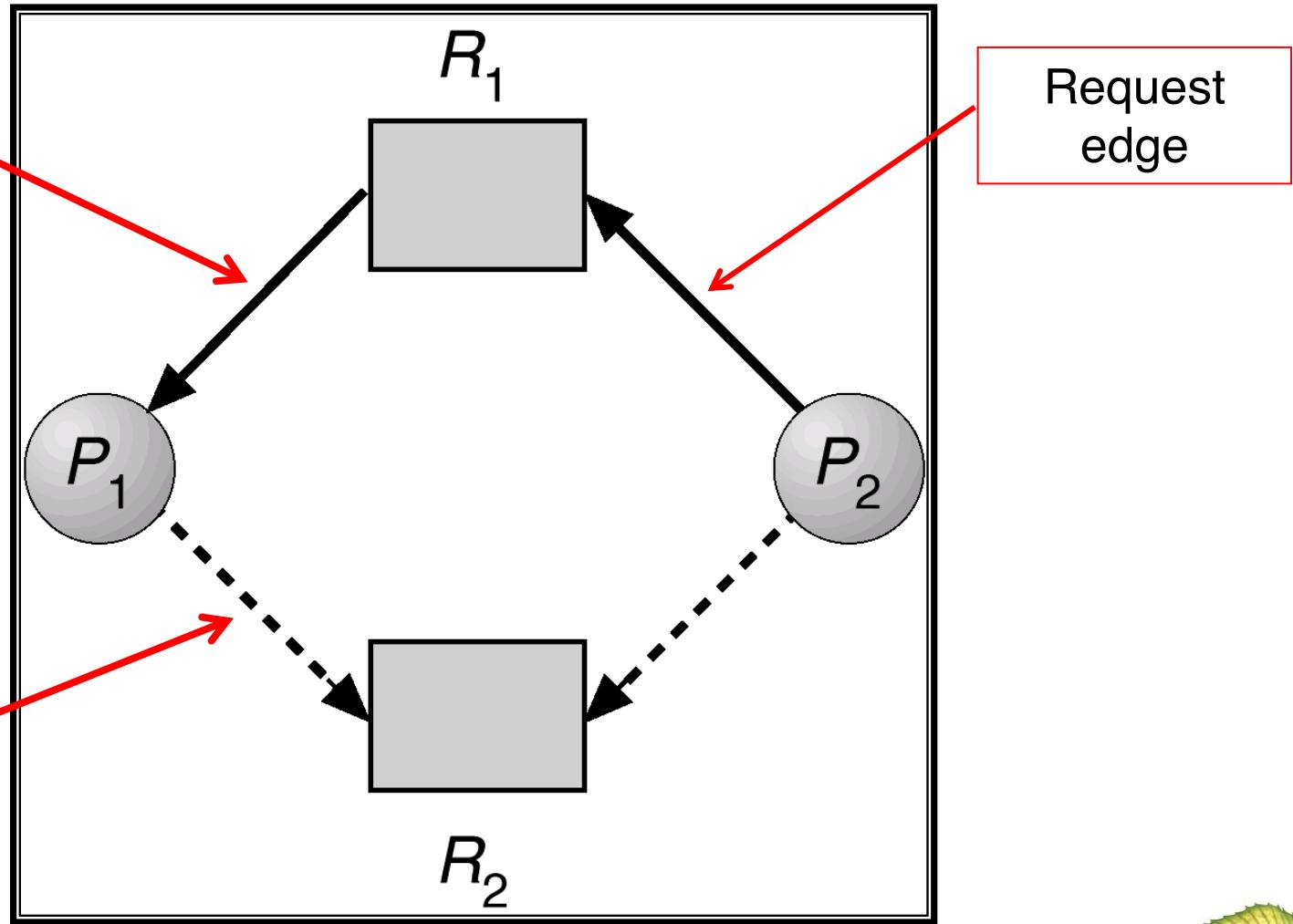
$R_i \rightarrow P_i$

Claim edge : เส้นความต้องการ (แสดงโดยใช้เส้นประ)



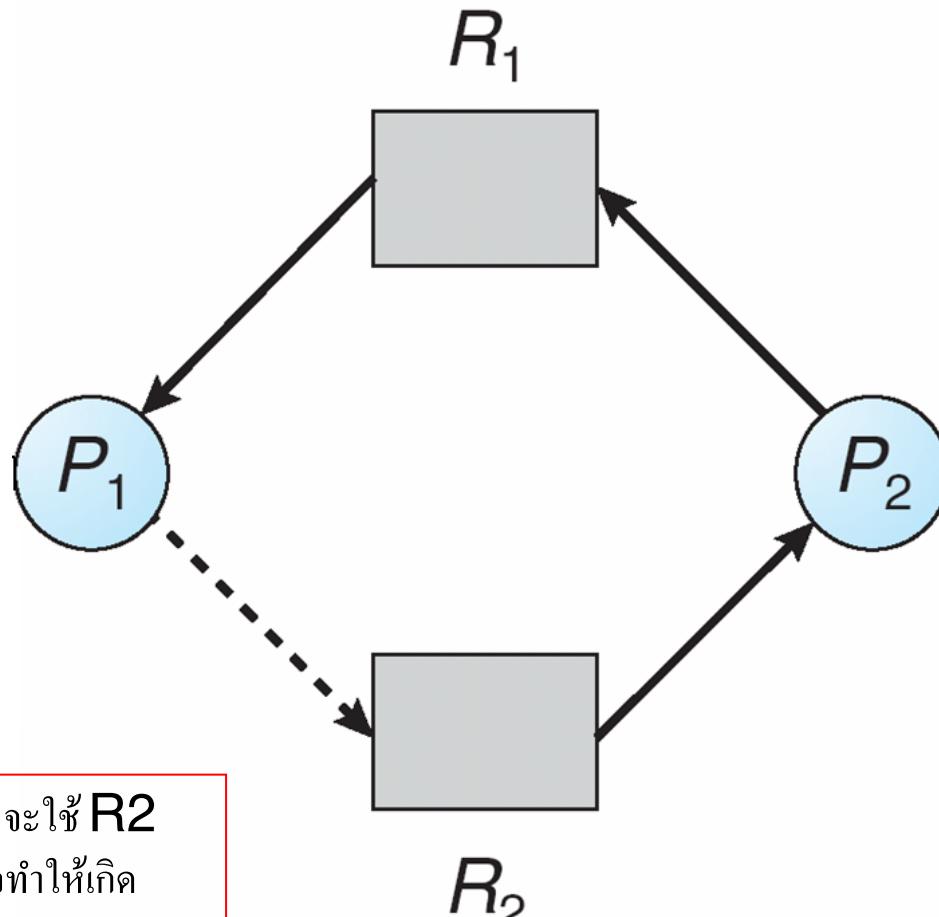


Resource-Allocation Graph For Deadlock Avoidance



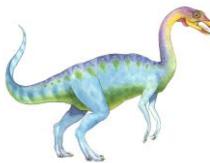


Unsafe State In Resource-Allocation Graph



หากมีความต้องการของ P_1 จะใช้ R_2
ระบบจะไม่อนุญาต เพราะอาจทำให้เกิด
Deadlock (ดูจากลักษณะอาจเกิด
วงรอบ ขึ้นได้)





Resource-Allocation Graph Algorithm

- Suppose that process P_i requests a resource R_j
- The request can be granted **only if** converting the request edge to an assignment edge does not result in the formation of a cycle in the resource allocation graph

ဒါမ္မာများကို ပြန်လည်
အသေးစိတ်ပေး.





Banker's Algorithm

For deadlock avoidance with MULTIPLE instance of each resource type

- Multiple instances.
- Each process must a priori claim maximum use.
- When a process requests a resource it may have to wait.
- When a process gets all its resources it must return them in a finite amount of time.





Data Structures for the Banker's Algorithm

Let n = number of processes, and m = number of resources types.

- **Available:** Vector of length m . If $\text{available}[j] = k$, there are k instances of resource type R_j available. (เก็บ Resource ที่ว่าง)
- **Max:** $n \times m$ matrix. If $\text{Max}[i,j] = k$, then process P_i may request at most k instances of resource type R_j . (เก็บจน. สูงสุดของ Resource ที่กระบวนการแต่ละตัว ต้องการใช้)
- **Allocation:** $n \times m$ matrix. If $\text{Allocation}[i,j] = k$ then P_i is currently allocated k instances of R_j . (เก็บจน. Resource ที่แต่ละกระบวนการครอบครองอยู่)
- **Need:** $n \times m$ matrix. If $\text{Need}[i,j] = k$, then P_i may need k more instances of R_j to complete its task. (เก็บจน. Resource ที่เหลืออยู่ ที่แต่ละกระบวนการยังคงต้องการใช้ เพื่อ ทำงานให้เสร็จสมบูรณ์)

$$\text{Need}[i,j] = \text{Max}[i,j] - \text{Allocation}[i,j].$$

(ที่ขอใช้ Resource สูงสุด – ที่ได้ Resource ครอบครองแล้ว)





Safety Algorithm

1. Let $Work$ and $Finish$ be vectors of length m and n , respectively.
Initialize:

$Work = Available$

the number of resource available for P_i to use of its work

$Finish[i] = false$ for $i = 1, 2, \dots, n$. for all processes P_i , initialize

the finish status to FALSE

2. Find and i such that both:

(a) $Finish[i] = false$

(b) $Need_i \leq Work$

ตรวจสอบเงื่อนไข

If ($Finish[i] == false$ AND $Need_i \leq work$)

If no such i exists, go to step 4.

3. $Work = Work + Allocation_i$ After P_i finished its task, all allocated resource must be released
 $Finish[i] = true$ to the system
go to step 2.
4. If $Finish[i] == true$ for all i , then the system is in a safe state.

Resource มีสถานะที่ available เพื่อรอในการทำงานครั้งต่อไป





Resource-Request Algorithm for Process P_i

Request = request vector for process P_i . If $Request_i[j] = k$ then process P_i wants k instances of resource type R_j .

1. If $Request_i \leq Need_i$, go to step 2. Otherwise, raise error condition, since process has exceeded its maximum claim.
Need_i = the number of resources that P_i needs more
2. If $Request_i \leq Available$, go to step 3. Otherwise P_i must wait, since resources are not available.
Available = the number of ready available resources in the system
3. Pretend to allocate requested resources to P_i by modifying the state as follows:

$Available = Available - Request_i$; Fewer resources are left in the system

$Allocation_i = Allocation_i + Request_i$; More resources allocated to P_i

$Need_i = Need_i - Request_i$; P_i need fewer of additional resources
(fewer $Need_i[j]$ left).

- If $safe \Rightarrow$ the resources are allocated to P_i .
- If $unsafe \Rightarrow P_i$ must wait, and the old resource-allocation state is restored

All allocated resources of P_i are released.

exceed: มากเกิน

pretend: การอ้างถึง





Example of Banker's Algorithm

- ◻ 5 processes P_0 through P_4 ; 3 resource types A (10 instances), B (5 instances), and C (7 instances).
 - ◻ Snapshot at time T_0 :

System-j of Rj is composed of

	<u>Allocation</u>	<u>Max</u>	<u>Available</u>
P_0	A B C 0 1 0	A B C 7 5 3	A B C 3 3 2
P_1	A B C 2 0 0	A B C 3 2 2	
P_2	A B C 3 0 2	A B C 9 0 2	
P_3	A B C 2 1 1	A B C 2 2 2	
P_4	A B C 0 0 2	A B C 4 3 3	





Example (Cont.)

- need more storage
□ The content of the matrix. **Need** is defined to be **Max – Allocation**.

	<u>Need</u>		
	A	B	C
P_0	7	4	3
P_1	1	2	2
P_2	6	0	0
P_3	0	1	1
P_4	4	3	1

7 5 3 - 0 1 0

4 3 3 - 0 0 2

Found at least 1 safe sequence*

- The system is in a **safe state** since the sequence $\langle P_1, P_3, P_4, P_2, P_0 \rangle$ satisfies safety criteria.



P6.30 i - Process P_i ; j - Resource Type R_j ;

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Home Insert Draw View Class Notebook

P.6.30 i - Process P_i ; j - Resource Type R_j

System_j = (10, 5, 7). *No more use.*

Step 1: Compute initial Available_j

Allocation_i:

	A	B	C	
③ P0	(0, 1, 0)	(7, 5, 3)		
① P1	(2, 0, 0)	(3, 2, 2)		
④ P2	(3, 0, 2)	(9, 0, 2)		
② P3	(2, 1, 1)	(2, 2, 2)		
⑤ P4	(0, 0, 2)	(4, 3, 3)		

$\sum_j = (7, 2, 5)$

Max_i:

	A	B	C	
③ P0	(0, 1, 0)	(7, 5, 3)		
① P1	(2, 0, 0)	(3, 2, 2)		
④ P2	(3, 0, 2)	(9, 0, 2)		
② P3	(2, 1, 1)	(2, 2, 2)		
⑤ P4	(0, 0, 2)	(4, 3, 3)		

Step 2: Compute & Check

Max_i - Allocation_i

Need_i:

	A	B	C	
③ P0	(6, 4, 3)	x	x	-
① P1	(1, 2, 2)	✓	-	-
④ P2	(6, 0, 0)	x	✓	-
② P3	(0, 1, 1)	✓	-	-
⑤ P4	(4, 3, 1)	✓	-	✓

Step 3: Update Available_j + Allocation_i

Available_j:

	A	B	C	
③ P0	(3, 3, 2)	+ (2, 0, 0)	① P1	
① P1	(5, 3, 2)	+ (2, 1, 1)	② P3	
④ P2	(7, 4, 3)	+ (0, 1, 0)	③ P0	
② P3	(7, 5, 3)	+ (3, 0, 2)	④ P2	
⑤ P4	(10, 5, 5)	+ (0, 0, 2)	⑤ P4	

[Checking]

IF (Need_i ≤ Available_j) THEN $\boxed{(10, 5, 7)} = \text{System}_j \therefore \text{O.K.}$

Available_j = Available_j + Allocation_i

// Add P_i to the safe sequence.

Available_j = System_j - \sum_j (Allocation_i)

= (10, 5, 7) - (7, 2, 5)

= (3, 3, 2).

Step A : Conclude

∴ A 'safe sequence' is found

The system is in the 'safe state'



Example P_1 Request (1,0,2)

- Check that **Request** \leq **Available** (that is $(1,0,2) \leq (3,3,2)$) \Rightarrow true.

	<u>Allocation</u>	<u>Need</u>	<u>Available</u>
	A B C	A B C	A B C
P_0	2 0 0 + 1 0 2	0 1 0 ↘	2 3 0 ←
P_1	3 0 2	0 2 0 ↙	3 3 2 - 1 0 2
P_2	3 0 2	6 0 0	??
P_3	2 1 1	0 1 1	1 2 2 - 1 0 2
P_4	0 0 2	4 3 1	

- Executing safety algorithm shows that sequence $\langle P_1, P_3, P_4, P_0, P_2 \rangle$ satisfies safety requirement.
 - Can request for (3,3,0) by P_4 be granted?
 - Can request for (0,2,0) by P_0 be granted?



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Semester 6501

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Soln.

$\therefore \text{Allocation}_j = \text{Allocation}_{j-1} + \text{Request}_j$

$= (2, 0, 0) + (1, 0, 2)$

$= (3, 0, 2)$.

No more use.

Step 2: Compute & Check

Max_j - Allocation_j

Available_j = Available_{j-1}

Allocation _j			Max _j			Need _j			Available _j		
A	B	C	A	B	C	P ₁	P ₂	P ₃	P ₄	P ₅	P ₆
④ P ₀	(0, 1, 0)		(7, 5, 3)			(2, 3, 0)					
① P ₁	(3, 0, 2)		(3, 2, 2)			(5, 3, 2)					
⑤ P ₂	(3, 0, 2)		(9, 0, 2)			(7, 4, 3)					
② P ₃	(2, 1, 1)		(2, 2, 2)			(0, 2, 0)					
③ P ₄	(0, 0, 2)		(4, 3, 3)			(6, 0, 0)					
	\sum_j					(0, 1, 1)					
	$= (8, 2, 7)$					(4, 3, 1)					

[Checking]

IF ($\text{Need}_j \leq \text{Available}_j$) THEN

$(2, 3, 0) + (3, 0, 1)$
 $(5, 3, 2) + (2, 1, 1)$
 $(7, 4, 3) + (0, 0, 0)$
 $(7, 4, 5) + (0, 1, 0)$
 $(7, 5, 5) + (3, 0, 2)$
 $(10, 5, 7) \Rightarrow \text{System Safe}$



Practice: Banker's Algorithm

- Snapshot at time T_0 :

	<u>Allocation</u>			<u>Max</u>			<u>Available</u>		
	<i>A</i>	<i>B</i>	<i>C</i>	<i>A</i>	<i>B</i>	<i>C</i>	<i>A</i>	<i>B</i>	<i>C</i>
P_0	0	1	0	7	5	3	3	3	2
P_1	2	0	0	3	2	2			
P_2	3	0	2	9	0	2			
P_3	2	1	1	2	2	2			
P_4	0	0	2	4	3	3			

- จาก Snapshot T_0 จงแสดงวิธีทำเพื่อที่จะตรวจสอบว่า ระบบจะอยู่ในสถานะที่ปลอดภัย (safe state) หรือไม่ ?? เมื่อมีลำดับของโปรเซส คือ P_0, P_1, P_2, P_3, P_4





Practice: Banker's Algorithm (Solution)

- คำนวณหา Need ของแต่ละ Process ก่อน แล้วจึงมาดำเนินการตาม Banker Algorithm

Snapshot at time T_0 :

	<u>Allocation</u>	<u>Max</u>	<u>Need</u>	<u>Available</u>
	A B C	A B C	A B C	A B C
P_0	0 1 0	7 5 3	7 4 3	3 3 2
P_1	2 0 0	3 2 2	1 2 2	
P_2	3 0 2	9 0 2	6 0 0	
P_3	2 1 1	2 2 2	0 1 1	
P_4	0 0 2	4 3 3	4 3 1	

ระบบจะอยู่ในสถานะที่ปลอดภัย (safe state) หรือไม่ ?? เมื่อมีลำดับของโปรเซส คือ P_3, P_2, P_1, P_4, P_0





Deadlock Detection

หากไม่มีการ Protection และ Avoidance

- Allow system to enter deadlock state
- Detection algorithm ตรวจสอบ
- Recovery scheme ฟื้นฟู





Single Instance of Each Resource Type

- Maintain **wait-for graph**
 - Nodes are processes.
 - $P_i \rightarrow P_j$ if P_i is waiting for P_j .
กระบวนการจับเป็นจั่ง
- Periodically invoke an algorithm that searches for a **cycle in the graph**.
- An algorithm to detect a cycle in a graph requires an order of n^2 operations, where n is the number of vertices in the graph.
 $n = \text{จำนวนตัวแปร} / \text{ค่าคงที่}$

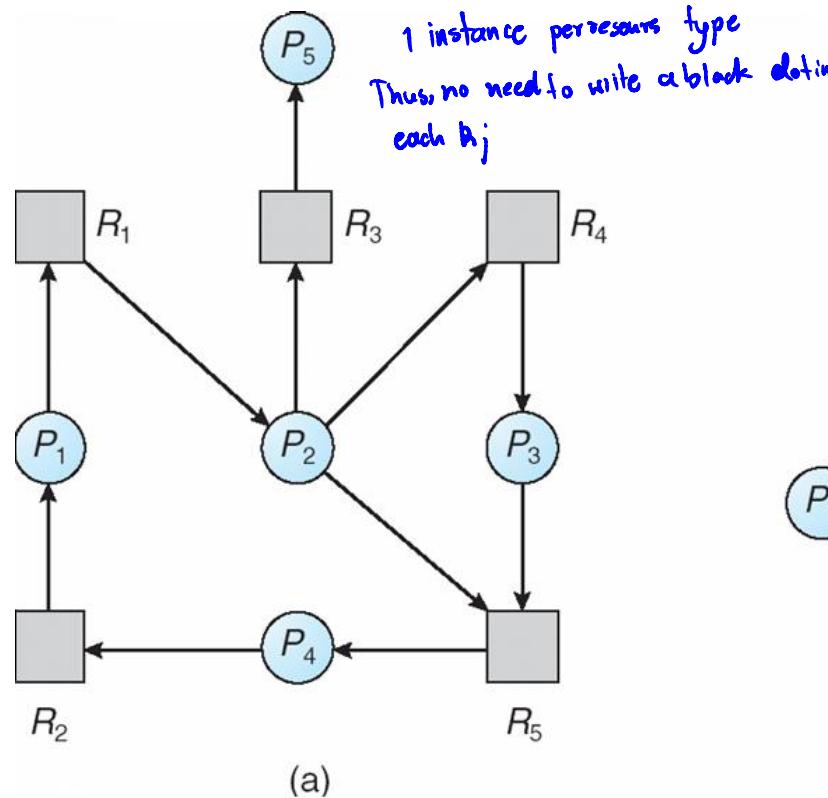
แปลงเป็น กราฟการรอคิวยทรัพยากร (wait-for graph)

จะไม่มีโหนดของ Resource แล้ว

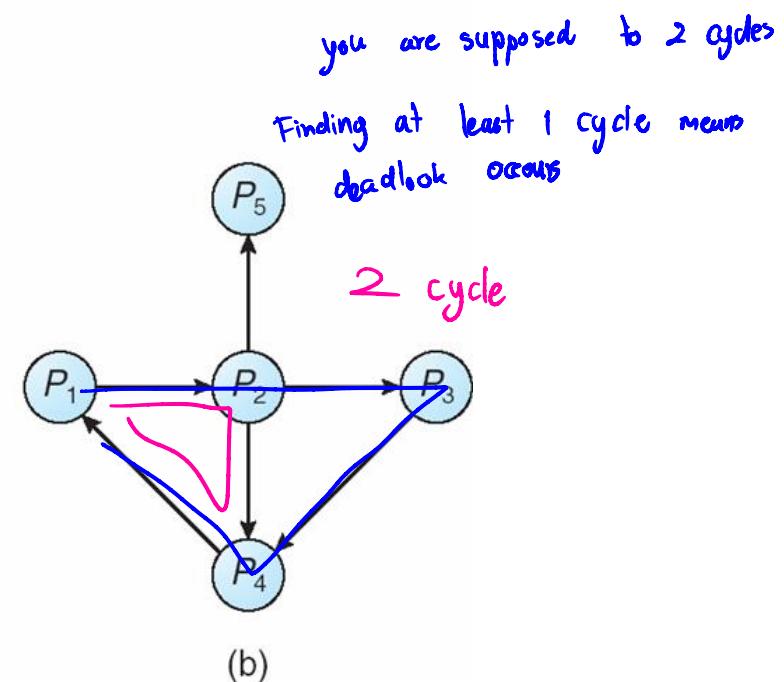




Resource-Allocation Graph and Wait-for Graph



Resource-Allocation Graph



Corresponding wait-for graph





Several Instances of a Resource Type

m — # of resource types

n — # of processes

- **Available**: A vector of length m indicates the number of **available resources** of each type.
- **Allocation**: An $n \times m$ matrix defines the number of **resources** of each type **currently allocated** to each process.
- **Request**: An $n \times m$ matrix indicates the **current request** of each process. If $\text{Request}[i,j] = k$, then process P_i is requesting k more instances of resource type R_j .





Detection Algorithm

1. Let Work and Finish be vectors of length m and n , respectively Initialize:

(a) Work = Available

(b) For $i = 1, 2, \dots, n$, if Allocation $_i \neq 0$, then
Finish[i] = false; otherwise, Finish[i] = true.

ของ Banker's Algo
For i=1,2 ..n then
Finish[i]=false;

2. Find an index i such that both:

(a) Finish[i] == false

(b) Request $_i \leq \text{Work}$

ตรวจสอบเงื่อนไข Request $_i$
If (Finish[i] == false AND Need $_i \leq \text{work}$)

P i : does not finish yet

P i 's additional
request \leq available

If no such i exists, go to step 4.





Detection Algorithm (Cont.)

3. $Work = Work + Allocation_i$
 $Finish[i] = true$
go to step 2
4. If $Finish[i] == \text{false}$, for some i , $1 \leq i \leq n$, then the system is in deadlock state. Moreover, if $Finish[i] == \text{false}$, then P_i is deadlocked

After P_i finished, all allocated resource must be returned to the system

Algorithm requires an order of $O(m \times n^2)$ operations to detect whether the system is in deadlocked state





Example of Detection Algorithm

- Five processes P_0 through P_4 ; three resource types A (7 instances), B (2 instances), and C (6 instances)
- Snapshot at time T_0 :

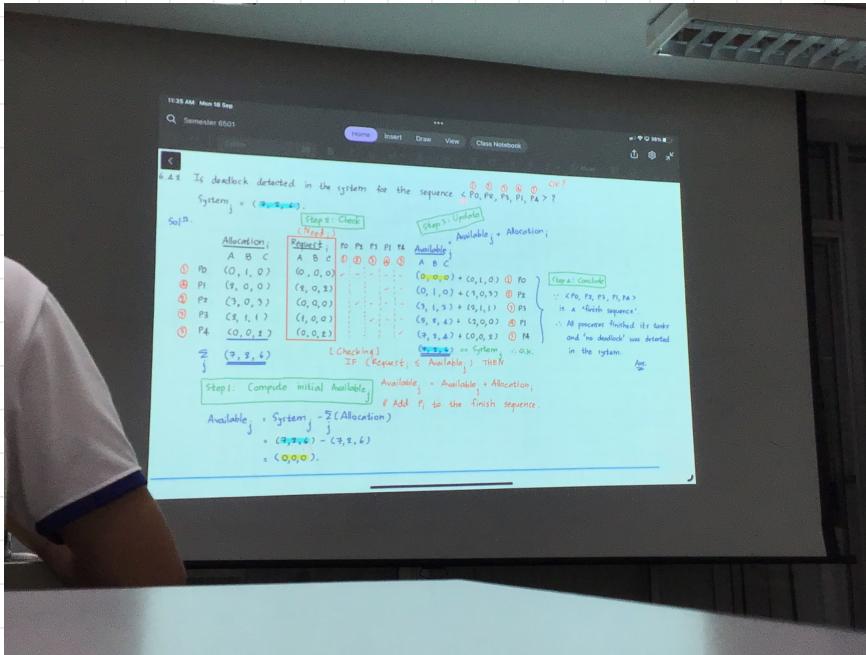
Process	Allocation			Request A B C	Available A B C
	A	B	C		
P_0	0	1	0	0 0 0	0 0 0
P_1	2	0	0	2 0 2	
P_2	3	0	3	0 0 0	
P_3	2	1	1	1 0 0	
P_4	0	0	2	0 0 2	

Similar to Need[i] in Banker.

See my note

- Sequence $\langle P_0, P_2, P_3, P_1, P_4 \rangle$ will result in $Finish[i] = \text{true}$ for all i

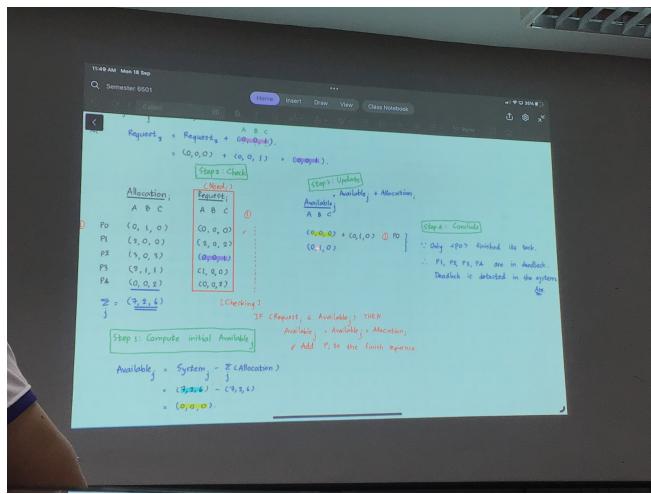
Thus, the system detects detectable deadlock.





Example (Cont.)

- P_2 requests an additional instance of type C



Request

	A	B	C
P_0	0	0	0
P_1	2	0	1
P_2	0	0	1
P_3	1	0	0
P_4	0	0	2

ឧបរកម្រាវសាស្ត្រ.

$$(7, 2, 6) - (7, 2, 6) \\ (0, 0, 0)$$

- State of system?

- Can reclaim resources held by process P_0 , but insufficient resources to fulfill other processes; requests
- Deadlock exists**, consisting of processes P_1, P_2, P_3 , and P_4

insufficient : ไม่เพียงพอ





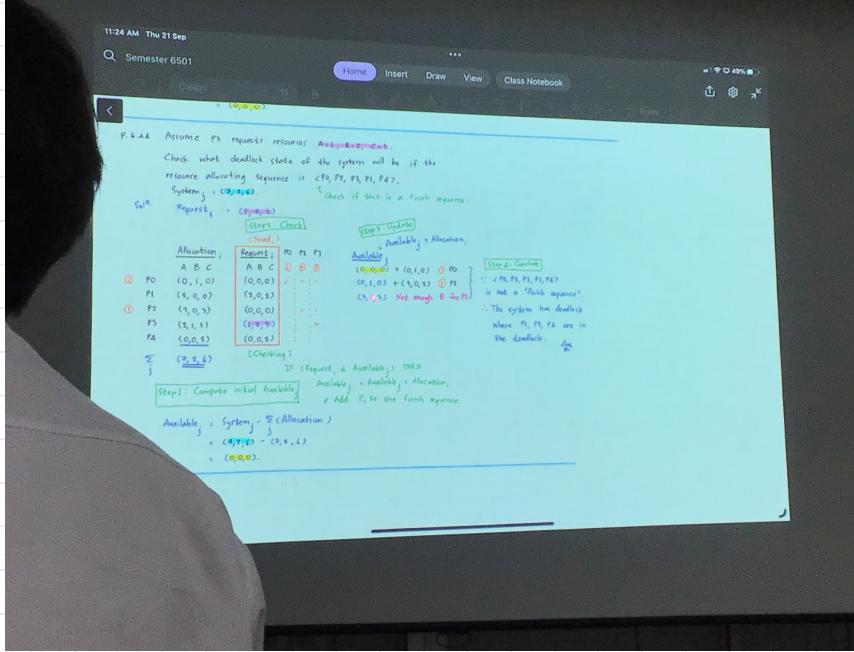
Practice: Deadlock Detection

- Snapshot at time T_0 :

	<u>Allocation</u>			<u>Request</u>			<u>Available</u>		
	A	B	C	A	B	C	A	B	C
P_0	0	1	0	0	0	0	0	0	1
P_1	2	0	0	2	0	2	0	1	0
P_2	3	0	3	0	1	0	3	1	3
P_3	2	1	1	1	2	3	5	2	4
P_4	0	0	2	0	0	2	$(1,0,0) \rightarrow (1,2,3)$ Not addition request but just change P_3 's request		
	7 2 6								

- จาก Snapshot T_0 เมื่อ P_3 ต้องการ (Request) ทรัพยากรชนิดต่างๆ (Resource type) A = 1, B=2, และ C =3 เมื่อลำดับการทำงานของโปรแกรม คือ P_0, P_2, P_3, P_1 และ P_4 ให้นักศึกษาแสดงวิธีทำในการตรวจสอบสถานะของระบบว่าจะมีสถานะเป็นอย่างไร ??







Detection-Algorithm Usage

- When, and how often, to invoke depends on:
 - How often a deadlock is likely to occur?
 - How many processes will need to be rolled back?
 - ▶ one for each disjoint cycle
- Roll back processes one-by-one until each cycle is broken.
- If detection algorithm is invoked arbitrarily, there may be many cycles in the resource graph and so we would not be able to tell which of the many deadlocked processes “caused” the deadlock

arbitrarily : ไม่มีกฎเกณฑ์



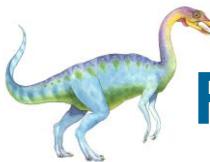


Recovery from Deadlock: Process Termination

(ยกเลิกกระบวนการที่เกิด **deadlock**)

- Abort all deadlocked processes
- Abort one process at a time until the deadlock cycle is eliminated
สิ่ง nào mr abort?
- In which order should we choose to abort?
 - Priority of the process *ความสำคัญของมัน* process *อัตราในการทำงานของมัน*
 - How long process has computed, and how much longer to completion
 - Resources the process has used *ทรัพยากร่วมที่มันใช้*
 - Resources process needs to complete *ต้องการ*
 - How many processes will need to be terminated *จำนวน processes*
 - Is process interactive or batch? *Batch processes are likely to be terminated first*
มีการตอบกลับ กับ ไม่มี





Recovery from Deadlock: Resource Preemption

ພວ.ຖົງກົມພາກຕູກ ແກ້ວການ

ກາຮເລືອກໃໝ່ ຈະຕ້ອງພິຈາລະນາຜລທີ່ຈະເກີດຂຶ້ນ 3 ຊ້ອ ດັ່ງນີ້

- Selecting a victim – minimize cost
- Rollback – return to some safe state, restart process for that state
- Starvation – same process may always be picked as victim, include number of rollback in cost factor

We must ensure that a process can be picked as a victim only a (small) finite number of times.
The most common SOLUTION is to include
the number of rollback in the factor.

Victim : ຜູ້ຮັບເຄራະໜ້າ (process)



End of Chapter 6

