



Managed by Jivan Jyot Trust, Amroli

**PROF. V. B. SHAH INSTITUTE OF MANAGEMENT | R. V. PATEL COLLEGE OF COMMERCE (ENG. MED.),
V. L. SHAH COLLEGE OF COMMERCE (GUJ. MED.) | SUTEX BANK COLLEGE OF COMPUTER APPLICATIONS & SCIENCE, AMROLI.**

Accredited by National Assessment and Accreditation Council with 'B' Grade
All colleges are Affiliated to Veer Narmad South Gujarat University, Surat

B.Sc. (I.T.) / M.Sc. (I.T.) 1s Semester

Course: 104: Fundamentals of Programming Using C-1

Unit 4: Array

- **4.1 One dimensional Array**
- **4.2 Declaration & Initialization of Array**
- **4.3 Two dimensional array**
 - **4.3.1 Declaration**
 - **4.3.2 Accessing Matrix Elements**
 - **4.3.3 Operations on matrix elements and entire matrices**
- **4.4 Array manipulation**
 - **4.4.1 Searching**
 - **4.4.2 Insertion**
 - **4.4.3 Deletion**
 - **4.4.4 Modification**
 - **4.4.5 Sorting**
- **4.5 Multidimensional Array**

4.1 One-Dimensional Array

A **one-dimensional array (1D array)** stores a list of elements **of the same data type** in **continuous memory locations**.

This makes searching, sorting, and processing extremely fast.

Why 1D Arrays Are Needed

Suppose we want to store marks of 50 students.

Using 50 variables is impossible:

```
int m1, m2, m3, ..., m50;
```

Instead:

```
int marks[50];
```

Advantages:

- ✓ Memory-efficient
- ✓ Easy to process using loops
- ✓ Fast access using index
- ✓ Predictable memory layout



Managed by Jivan Jyot Trust, Amroli

**PROF. V. B. SHAH INSTITUTE OF MANAGEMENT | R. V. PATEL COLLEGE OF COMMERCE (ENG. MED.),
V. L. SHAH COLLEGE OF COMMERCE (GUJ. MED.) | SUTEX BANK COLLEGE OF COMPUTER APPLICATIONS & SCIENCE, AMROLI.**

Accredited by National Assessment and Accreditation Council with 'B' Grade

All colleges are Affiliated to Veer Narmad South Gujarat University, Surat

4.2 Declaration & Initialization of Array

Declaration

`data_type array_name[size];`

Examples:

`int marks[5];`

`float prices[10];`

`char vowels[5];`

Logic Behind Indexing

Indexes start from 0 because:

Memory offset calculated as:

`address = base_address + (index * size_of_data_type)`

So index 0 naturally refers to the base address.

Initialization

1. Complete Initialization

`int a[5] = {10, 20, 30, 40, 50};`

2. Partial Initialization (rest become 0)

`int a[5] = {1, 2}; // → {1,2,0,0,0}`

3. Automatic Size

`int a[] = {5, 10, 15};`

4. Runtime Input

`for(i=0;i<5;i++)`

`scanf("%d",&a[i]);`

Logic During Runtime Input

Each scanf reads individual elements and stores them in consecutive memory blocks.



4.3 Two-Dimensional Array (Matrix)

A 2D array is like a table of rows and columns.

Example:

```
int a[3][3];
```

Represents:

```
[ a00 a01 a02 ]
```

```
[ a10 a11 a12 ]
```

```
[ a20 a21 a22 ]
```

4.3.1 Declaration

```
int matrix[rows][columns];
```

Example:

```
int a[4][3]; // 4 rows, 3 columns
```

4.3.2 Accessing Matrix Elements

```
a[row][column]
```

Logic Behind Accessing Element

For element $a[i][j]$:

address = base + $((i * \text{total_columns}) + j) * \text{size}$

This is called **row-major order** (C language default).

Example

Input for 2×2 matrix:

```
for(i=0;i<2;i++)
```

```
    for(j=0;j<2;j++)
```

```
        scanf("%d",&a[i][j]);
```

Here:

- Outer loop changes row
 - Inner loop fills each column of that row
- This ensures predictable tabular input.



4.3.3 Operations on Matrix

1. Matrix Addition

$$c[i][j] = a[i][j] + b[i][j];$$

Logic

Each element in the resulting matrix is sum of elements at **same position**.

Example:

$$A = \begin{bmatrix} 1 & 2 \\ 3 & 4 \end{bmatrix}$$

$$B = \begin{bmatrix} 5 & 6 \\ 7 & 8 \end{bmatrix}$$

$$C = \begin{bmatrix} 1+5 & 2+6 \\ 3+7 & 4+8 \end{bmatrix}$$

2. Matrix Subtraction

$$c[i][j] = a[i][j] - b[i][j];$$

Logic

Position-wise subtraction.

3. Matrix Multiplication

Formula:

$$C[i][j] = \text{sum} (A[i][k] * B[k][j])$$

Logic

- Take **row** from A
- Take **column** from B
- Multiply corresponding elements
- Add all products

Example Step-by-Step:

A (2×3) and B (3×2)

$$\begin{array}{cc} A & B \\ \begin{bmatrix} 1 & 2 & 3 \\ 4 & 5 & 6 \end{bmatrix} & \begin{bmatrix} 1 & 2 \\ 3 & 4 \\ 5 & 6 \end{bmatrix} \end{array}$$

Compute C[0][0]:

$$\begin{aligned} & (1 \times 1) + (2 \times 3) + (3 \times 5) \\ & = 1 + 6 + 15 \\ & = 22 \end{aligned}$$

This logic applies for all elements.



Managed by Jivan Jyot Trust, Amroli

**PROF. V. B. SHAH INSTITUTE OF MANAGEMENT | R. V. PATEL COLLEGE OF COMMERCE (ENG. MED.),
V. L. SHAH COLLEGE OF COMMERCE (GUJ. MED.) | SUTEX BANK COLLEGE OF COMPUTER APPLICATIONS & SCIENCE, AMROLI.**

Accredited by National Assessment and Accreditation Council with 'B' Grade

All colleges are Affiliated to Veer Narmad South Gujarat University, Surat

4. Diagonal Elements

Main diagonal

$a[i][i]$

Secondary diagonal

$a[i][n-i-1]$

Why?

Because secondary diagonal moves backward across columns.

5. Row & Column Sum

Row Sum Logic

Fix row

Vary column

Column Sum Logic

Fix column

Vary row

6. Transpose

$t[j][i] = a[i][j];$

Logic

Row becomes column → column becomes row.



4.4 Array Manipulation (With Logic)

These operations help manage data inside arrays.

4.4.1 Searching

1. Linear Search

Algorithm:

1. Start from index 0
2. Compare each element with key
3. Stop if found

Logic

Sequential check until match found.

Example:

```
for(i=0;i<n;i++)
```

```
    if(a[i] == key)
```

```
        found = i;
```

Time Complexity: **$O(n)$**

2. Binary Search

Works only on sorted arrays.

Logic:

1. Find middle element
2. Compare key with mid
3. If key < mid → search left
4. Else → search right

This halves the array each time.

Time Complexity: **$O(\log n)$**

4.4.2 Insertion

To insert element at position pos:

1. Start from last element
2. Shift all elements **right**
3. Insert value at pos

Example:

```
for(i=n-1;i>=pos;i--)
```

```
    a[i+1] = a[i];
```

```
a[pos] = value;
```

Logic

Shifting is necessary because arrays have fixed contiguous storage, so you cannot "create space" without shifting.



4.4.3 Deletion

Algorithm:

1. Start from position pos
2. Shift elements **left**
3. Overwrite the element to be deleted

Example:

```
for(i=pos; i<n; i++)  
    a[i] = a[i+1];
```

Logic

Left shifting closes the gap created after deleting.

4.4.4 Modification

Changing a specific element.

```
a[pos] = new_value;
```

Logic

Direct access by index makes modification **O(1)** constant time.

4.4.5 Sorting

A. Bubble Sort

Compares adjacent elements and swaps if needed.

Logic:

- Bubbles larger element to end in each pass.

Example (Step-by-step for [5,3,1]):

Pass 1:

5 > 3 → swap → [3,5,1]

5 > 1 → swap → [3,1,5]

Pass 2:

3 > 1 → swap → [1,3,5]

Final: [1,3,5]

B. Selection Sort

- Find smallest element
- Put at correct position

Logic:

- Select minimum
- Place at front
- Continue for rest



Managed by Jivan Jyot Trust, Amroli

**PROF. V. B. SHAH INSTITUTE OF MANAGEMENT | R. V. PATEL COLLEGE OF COMMERCE (ENG. MED.),
V. L. SHAH COLLEGE OF COMMERCE (GUJ. MED.) | SUTEX BANK COLLEGE OF COMPUTER APPLICATIONS & SCIENCE, AMROLI.**

Accredited by National Assessment and Accreditation Council with 'B' Grade

All colleges are Affiliated to Veer Narmad South Gujarat University, Surat

C. Insertion Sort

- Consider first element sorted
- Insert next element in sorted part

Logic:

Similar to playing cards sorting in a hand.

4.5 Multidimensional Array

Arrays with more than 2 dimensions.

Example:

3D array:

```
int box[3][4][2];
```

Logic

Represents data as:

- 3 layers
- Each layer has 4 rows
- Each row has 2 elements

Used in:

- Scientific simulations
- 3D graphics
- Multi-level data tables

Example access:

```
box[i][j][k]
```