

ACM 104 Homework 1

① By the definition of matrix multiplication, we have that $AB = C$, where $c_{ij} = \sum_{k=1}^p a_{ik} b_{kj}$

We observe that $c_j = \sum_{k=1}^p a_k \cdot b^{kj}$, so $C = [c_1 \ c_2 \ \dots \ c_n]$

Which can be written as $C = \left[\sum_{k=1}^p a_k \cdot b^k \ \sum_{k=1}^p a_k b^{k_2} \ \dots \ \sum_{k=1}^p a_k b^{k_n} \right]$

C can also be written as a sum of matrices
 $C = \sum_{k=1}^p [a_k b^{k_1} \ \dots \ a_k b^{k_n}]$

which is $\sum_{k=1}^p a_k b^k$. Thus, $AB = \sum_{k=1}^p a_k b^k$

② Let $A \in \mathbb{M}_{m \times n}$ be a strictly upper triangular matrix.
 By definition, it must be true that:

$$a_{ij} = \begin{cases} 0 & \text{if } i \geq j \\ R & \text{if } i < j \end{cases}$$

(claim) If A is strictly upper triangular, then for $A^k, k \geq 1$:

$$a_{ij}^k = \begin{cases} 0 & \text{if } i \geq j - k + 1 \\ R & \text{if } i < j - k + 1 \end{cases}$$

For $k=n$, the condition $i \geq j - k + 1$ will always be true.
 Thus, A is nilpotent.

To prove this, we want to show that for $A^k, i \geq j - k + 1 \Rightarrow a_{ij} = 0$

Base case: $k=1$. This case is trivial. The claim holds by definition if $k=1$.

Inductive step: Need to show that claim holds for all $k > 1$ as well.

Let p be a positive integer and let $A^p = B$.
 Then, $A^{p+1} = A \cdot A^p = A \cdot B$

We let $A \cdot B = C$, where $c_{ij} = \sum_{r=1}^n a_{ir} \cdot b_{rj}$

We see that $a_{ir} = 0$ for $r \leq i$ (by definition of A) and that $b_{rj} = 0$ for $r \geq j-p+1$.

So for $i \geq j-p$ we have that for any $r \in \{1 \dots n\}$, $c_{ij} = 0$ because it will be true that either $a_{ir} = 0$ for $r \leq i$ or $b_{rj} = 0$ for $r \geq j-p+1$.

Thus, any strictly upper triangular matrix is nilpotent.

(3) We want to use the following definitions for P_n, L_n , and U_n . (Code attached)

$$P_n = I_n, \quad L_n = \begin{bmatrix} 1 & & & & 0 \\ -\frac{1}{2} & \frac{1}{2} & 1 & & \\ & -\frac{1}{3} & \frac{1}{3} & 1 & \\ & & -\frac{1}{4} & \frac{1}{4} & \ddots \\ 0 & & & \ddots & -\frac{1}{n} \end{bmatrix} \quad \begin{array}{l} \text{For all } i=j+1, \quad l_{ij} = -\left(\frac{i-1}{i}\right) \\ \text{For all } i=j, \quad l_{ij} = 1 \\ \text{All other elements are 0} \end{array}$$

$$U_n = \begin{bmatrix} \frac{2}{1} & -1 & & & 0 \\ \frac{3}{2} & -1 & & & \\ & \frac{4}{3} & -1 & \ddots & \\ & & \ddots & \ddots & -1 \\ 0 & & & & \frac{n+1}{n} \end{bmatrix} \quad \begin{array}{l} \text{For all } i=j, \quad u_{ij} = \frac{j+1}{j} \\ \text{For all } i=j-1, \quad u_{ij} = -1 \\ \text{All other elements are 0} \end{array}$$

$$P_n A_n = L_n U_n \rightarrow \text{If } P_n = I_n \rightarrow A_n = L_n U_n \quad \text{let } C_n = L_n U_n$$

$$\text{Then } c_{ij} = \sum_{k=1}^n l_{ik} \cdot u_{kj}$$

$$\text{If } i=j=1: \quad c_{11} = \sum_{k=1}^n l_{1k} \cdot u_{1j} = 2$$

$$\text{If } i=j-1: \quad c_{ij} = -1 \cdot 1 = -1 \quad (\text{Gives } -1 \text{ values on super diagonal})$$

$$\text{If } i=j \neq 1: \quad c_{ij} = -\left(\frac{i-1}{i}\right) \cdot (-1) + (1) \cdot \left(\frac{i+1}{j}\right) = 2 \quad (\text{Gives 2's on diagonal})$$

$$\text{If } i=j+1: \quad c_{ij} = -\left(\frac{i-1}{i}\right)\left(\frac{i+1}{j}\right) = -\left(\frac{i+1}{j+1}\right)\left(\frac{i+1}{j}\right) = -1 \quad (\text{Gives } -1 \text{'s on sub-diagonal})$$

If $i \geq j+1$ or $i \leq j-1$: No overlap; $c_{ij} = 0$. Thus, the definitions of P_n, L_n , and U_n work.

ACM 104 Homework 1 Part 2

- (4)(A) A permutation matrix P represents a set of elementary row operations. The operations perform row swaps of the matrix that P is being applied to, so its operations can be undone by reapplying P . Because of this, $P = P^{-1}$.

An example of a permutation matrix could be: $P = \begin{bmatrix} 1 & & & \\ & 0 & 1 & \\ & 1 & 0 & \\ & & & 1 \end{bmatrix}$

row i
row j
col i
col j

$P = P^T$, P swaps rows i and j of the matrix if it is applied to. $P = P^{-1} = P^T$. Thus, any permutation matrix P is orthogonal.

- (4)(B) It is not true that an orthogonal matrix is necessarily a permutation matrix. A counterexample is: Let $M = \begin{bmatrix} 0 & -1 \\ -1 & 0 \end{bmatrix}$. $M = M^T$

$$M^{-1} = \frac{1}{-1} \begin{bmatrix} 0 & 1 \\ 1 & 0 \end{bmatrix} = \begin{bmatrix} 0 & -1 \\ -1 & 0 \end{bmatrix}$$

In this case, $M = M^{-1} = M^T$, but M is not a permutation matrix; it has values other than 1 or 0.

- (5) Let $A \in M_{n \times n}$ where A is some matrix. Let constant S be defined as: $S_{ij} = \begin{cases} a_{ii} & \text{if } i=j \\ \frac{a_{ij} + a_{ji}}{2} & \text{if } i \neq j \end{cases}$ (symmetric)

And let J be defined as: $J_{ij} = \begin{cases} 0 & \text{if } i=j \\ \frac{a_{ij} - a_{ji}}{2} & \text{if } i \neq j \end{cases}$ (skew symmetric)

Let $C = S + J$, then we have

$$C_{ij} = \begin{cases} a_{ij} & \text{if } i=j \\ \frac{a_{ij} + a_{ji}}{2} + \frac{a_{ij} - a_{ji}}{2} = \frac{2a_{ij}}{2} = a_{ij} & \text{if } i \neq j \end{cases}$$

Thus every matrix A can be written as $A = S + J$, $S^T = S$, $J^T = -J$

⑥ Attached

⑦(A) The difference between each adjacent column is $[1 \ 1 \ 1 \ \dots \ 1]$

Because of this, any 2 columns from A can be used to represent any of the other columns. Thus, the rank of A is 2.

⑧ Code attached. $x = [0 \ 0 \ 0 \ \dots \ 0 \ 0.01]$
 $x_{n-1} = 0$ and $x_n = 0.01$ (size n)

$x_n = 0.01$ is the only non-zero component of x .

```
% code for problem 3
n = 10;
A = zeros(n);

for i = 1:n
    for j = 1:n
        if i == j
            A(i, j) = 2;
        end
        if i == j + 1 || i == j - 1
            A(i, j) = -1;
        end
    end
end

% get L, U, and P
[L, U, P] = lu(A)
```

L =

Columns 1 through 7

1.0000	0	0	0	0	0	0
-0.5000	1.0000	0	0	0	0	0
0	-0.6667	1.0000	0	0	0	0
0	0	-0.7500	1.0000	0	0	0
0	0	0	-0.8000	1.0000	0	0
0	0	0	0	-0.8333	1.0000	0
0	0	0	0	0	-0.8571	1.0000
0	0	0	0	0	0	-0.8750
0	0	0	0	0	0	0
0	0	0	0	0	0	0

Columns 8 through 10

0	0	0
0	0	0
0	0	0
0	0	0
0	0	0
0	0	0
0	0	0
1.0000	0	0
-0.8889	1.0000	0
0	-0.9000	1.0000

U =

Columns 1 through 7

2.0000	-1.0000	0	0	0	0	0
0	1.5000	-1.0000	0	0	0	0
0	0	1.3333	-1.0000	0	0	0
0	0	0	1.2500	-1.0000	0	0
0	0	0	0	1.2000	-1.0000	0
0	0	0	0	0	1.1667	-1.0000
0	0	0	0	0	0	1.1429
0	0	0	0	0	0	0
0	0	0	0	0	0	0
0	0	0	0	0	0	0

Columns 8 through 10

0	0	0
0	0	0
0	0	0
0	0	0
0	0	0
0	0	0
0	0	0
-1.0000	0	0
1.1250	-1.0000	0
0	1.1111	-1.0000
0	0	1.1000

$P =$

1	0	0	0	0	0	0	0	0	0
0	1	0	0	0	0	0	0	0	0
0	0	1	0	0	0	0	0	0	0
0	0	0	1	0	0	0	0	0	0
0	0	0	0	1	0	0	0	0	0
0	0	0	0	0	1	0	0	0	0
0	0	0	0	0	0	1	0	0	0
0	0	0	0	0	0	0	1	0	0
0	0	0	0	0	0	0	0	1	0
0	0	0	0	0	0	0	0	0	1

Published with MATLAB® R2016a

```
% code for problem 6

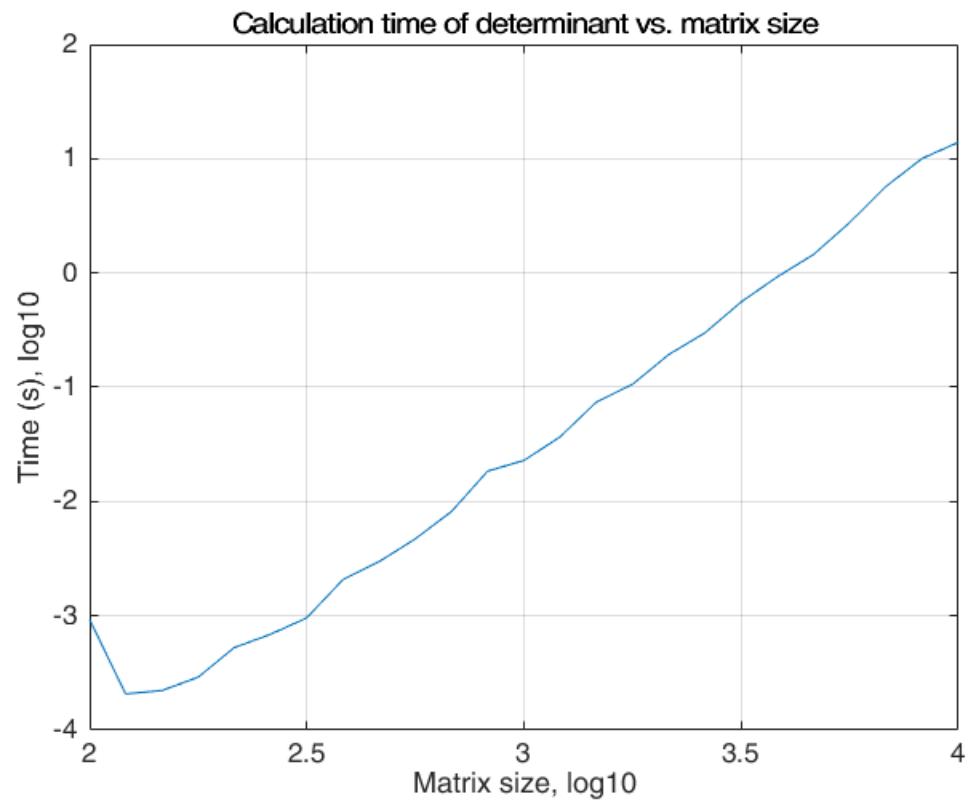
numvals = 25;
ns = logspace(2, 4, numvals); % part a
nLogs = zeros(1, numvals);

% part b
arrays = cell(1, numvals);
for i = 1 : numvals
    nInt = round(ns(i));
    arrays{i} = randn(nInt);
    nLogs(i) = log10(nInt);
end

%parts c and d
timeTaken = zeros(1, numvals);

for i = 1 : numvals
    tic
    det(arrays{i});
    timeTaken(i) = log10(toc);
end

% set up the plot
plot(nLogs, timeTaken)
title('Calculation time of determinant vs. matrix size')
xlabel('Matrix size, log10')
ylabel('Time (s), log10')
grid on
```



Published with MATLAB® R2016a

Published with MATLAB® R2016a