

Overview

- Basic I/O hardware
 - ports, buses, devices and controllers
- I/O Software
 - Interrupt Handlers, Device Driver, Device-Independent Software, User-Space I/O Software
- Real I/O devices
 - Disks, Character-Oriented Terminals, Graphical User Interfaces, Network Terminals,
- Power Management

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Device-Computer/Device-Device Communication

- Physically: via signals over a cable or through air
- Logically: via a connection point port (e.G., Serial port)
- Multiple devices are connected via a bus
 - A common set of wires and a rigidly defined protocol that specifies a set of messages that can be sent on the wires

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Devices

- Devices
 - Storage devices (disk, tapes)
 - Transmission devices (network card, modem)
 - Human interface devices (screen, keyboard, mouse)
 - Specialized device (joystick)

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Device Controller

- I/O units typically consist of a mechanical component and an electronic component. The electronic component is called the device controller or adapter. Example is a circuit board. The mechanical component is the device itself.
- Interface between controller and device is a very low level interface.
- Example:
 - Disk's controller converts the serial bit stream, coming off the drive into a block of bytes, and performs error correction. The block of bytes is first assembled in a buffer inside the controller. After its checksum has been verified, the error-free block is copied to main memory.
 - Built-in controllers

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I/O Controller

- Disk controller implements the disk side of the protocol that does: bad error mapping, prefetching, buffering, caching
- Controller has registers for data and control
- CPU and controllers communicate via I/O instructions and registers
- Memory-mapped I/O

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Host-controller interface: Interrupts

- CPU hardware has the interrupt report line that the CPU senses after executing every instruction
 - device raises an interrupt
 - CPU catches the interrupt and saves the state (e.g., Instruction pointer)
 - CPU dispatches the interrupt handler
 - interrupt handler determines cause, services the device and clears the interrupt
- Why interrupts?
- Real life analogy?

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I/O Ports

- 4 registers status, control, data-in, data-out
 - Status states whether the current command is completed, byte is available, device has an error, etc
 - Control host determines to start a command or change the mode of a device
 - Data-in host reads to get input
 - Data-out host writes to send output
- Size of registers 1 to 4 bytes

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Support for Interrupts

- Need the ability to defer interrupt handling during critical processing
- CPUs have two interrupt request lines
 - non-maskable interrupt (reserved for unrecoverable memory errors)
 - maskable interrupt (can be turned off by cpus before the execution of critical instructions
- Need efficient way to dispatch the proper device
 - Interrupt comes with an address (offset in interrupt vector) that selects a specific interrupt handling
- Need multilevel interrupts interrupt priority level

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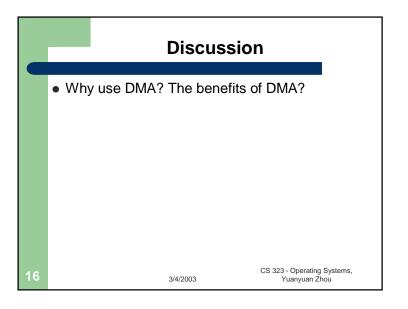
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At boot time, OS probes the hardware buses to determine what devices are present and installs corresponding interrupt handlers into the interrupt vector During I/O interrupt, controller signals that device is ready CS 323 - Operating Systems, Yuanyuan Zhou CS 323 - Operating Systems, Yuanyuan Zhou

Other Types of Interrupts Interrupt mechanisms are used to handle wide variety of exceptions: Division by zero, wrong address Virtual memory paging System calls (software interrupts, trap) Multi-threaded systems CS 323 - Operating Systems, Yuanyuan Zhou

Programmed I/O (PIO) use CPU to watch status bits and feed data into a controller register 1 byte at a time - EXPENSIVE for large transfers Direct memory access (DMA) use a special purpose processor, called a DMA controller CS 323 - Operating Systems, Yuanyuan Zhou



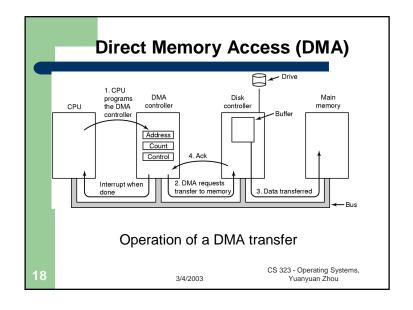
DMA-CPU Protocol Host writes a DMA command block into - pointer to source, pointer to destination, count of bytes to be transferred • CPU writes the address of this command block to the DMA controller and goes on with other work • DMA controller proceeds to operate the memory bus directly without help of main

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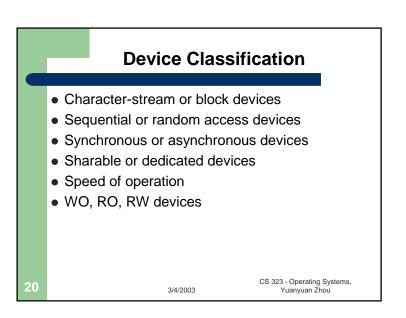
memory

CPU



DMA Issues Handshaking between DMA controller and the device controller Cycle stealing - DMA controller takes away CPU cycles when it uses CPU memory bus, hence blocks the CPU from accessing the memory • In general DMA controller improves the total system performance CS 323 - Operating Systems. 3/4/2003 Yuanyuan Zhou

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Block Devices Block device - transfers blocks of data (e.G., Disk device) Commands: read, write, seek (if random access device) Memory-mapped files access can be layered on top of block-device drivers

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Network Devices Network socket interface Commands; create socket, connect local socket to remote address, select information from a socket Other interfaces: half-duplex pipes, full-duplex fifos, message queues, sockets (UNIX) CS 323 - Operating Systems, Yuanyuan Zhou

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Character Devices Character device - transfers byte by byte (e.G., Keyboard) Commands: get, put one character Libraries can built on top of this device which provide line-at-a time access with buffering and editing services, etc Question: can block device be implemented using character device?

