

CS323 Operating Systems Deadlock III & Memory Management

Yuanyuan Zhou
Lecture 13
2/17/2003

Content of this lecture

- Administrative announcements
- Deadlock recovery
- Summary

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Administrative

- None

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Review

- Resource: preemptible or non-preemptible
- Deadlock avoidance
 - Unsafe vs. safe state
 - Banker algorithm
- Deadlock detection
 - Wait graph

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Recovery From Deadlock

OPTIONS:

- Kill deadlocked processes and release resources
- Kill one deadlocked process at a time and release its resources
- Rollback all or one of the processes to a checkpoint that occurred before they requested any resources

Note: with rollback, difficult to prevent indefinite postponement

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Deadlock Summary

- In general, deadlock detection or avoidance is expensive
- Must evaluate cost of deadlock against detection or avoidance costs
- Deadlock avoidance and recovery may cause indefinite postponement

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Memory Management

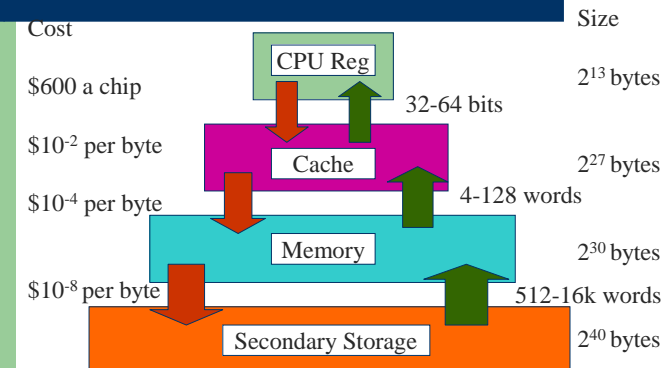
- Storage Hierarchy
- Resident Monitor
- Fixed Partitions for Spooling and Multiprogramming
- Variable Sized Partitions
- Storage Placement Strategies
- Compaction

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Storage Hierarchy



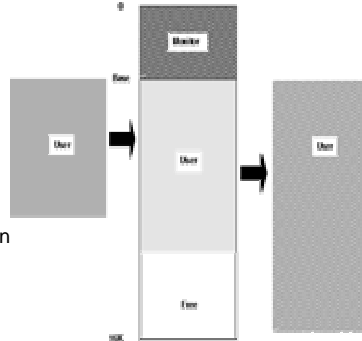
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Memory Manager

- Manage memory hierarchy
 - Monitor used and free memory
 - Allocate memory to processes
 - Reclaim (De-allocate) memory
 - Swapping between main memory and disk



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Mono-programming without Swapping

- Run one process at a time
 - simplest possible memory management scheme
- Memory is shared only between OS and the process.
- Three different ways to organize memory:
 - OS at the bottom of memory in RAM (Random Access Memory); was used in mainframes and minicomputers; not used anymore
 - OS is in ROM (Read-Only Memory) at the top of memory; is used in some palm-tops and embedded systems
 - Device drivers are in ROM at the top of memory and the rest is in RAM; was used by early PCs (e.g., running MS-DOS), where the portion in the ROM is called BIOS (Basic Input Output System).

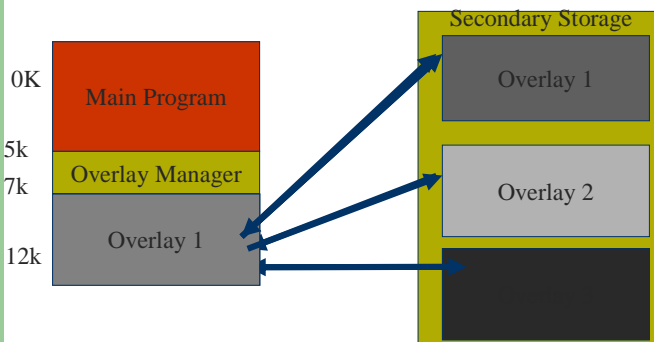
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Overlaying

Used when process memory requirement exceeds the physical memory space



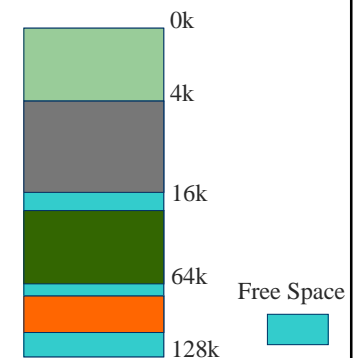
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Multiprogramming with Fixed Partitions

- Divide memory into n (possible unequal) partitions.
- Problem:
 - Fragmentation



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Fixed Partition Allocation

- Separate input queue for each partition
 - Requires sorting the incoming jobs and putting them into separate queues
 - Inefficient utilization of memory
 - when the queue for a large partition is empty but the queue for a small partition is full. Small jobs have to wait to get into memory even though plenty of memory is free.
- One single input queue for all partitions.
 - Allocate a partition where the job fits in.
 - Best Fit
 - Available Fit

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Relocation

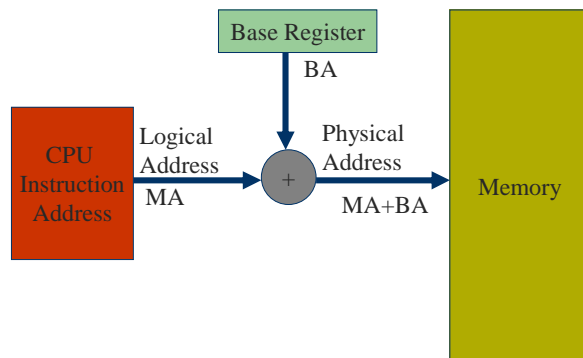
- Correct starting address when a program should start in the memory
- Different jobs will run at different addresses
 - When a program is linked, the linker must know at what address the program will begin in memory.
- Logical addresses, Virtual addresses
 - Logical address space, range (0 to max)
- Physical addresses, Physical address space
 - range (R+0 to R+max) for base value R.
- User program never sees the real physical addresses
- Memory-management unit (MMU)
 - map virtual to physical addresses.
- relocation register
 - Mapping requires hardware (MMU) with the base register

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Relocation Register



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Protection

- Problem:
 - How to prevent a malicious process to write or jump into other user's or OS partitions
- Solution:
 - Base bounds registers

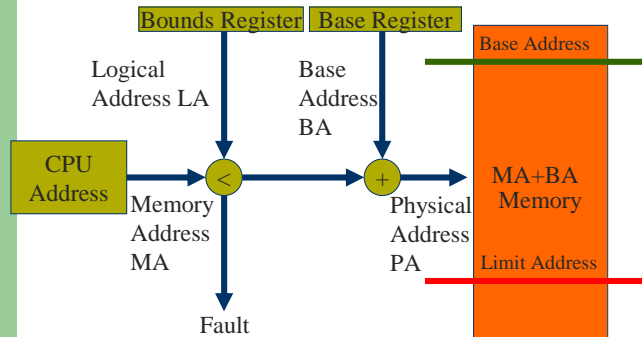


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Base Bounds Registers



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Swapping

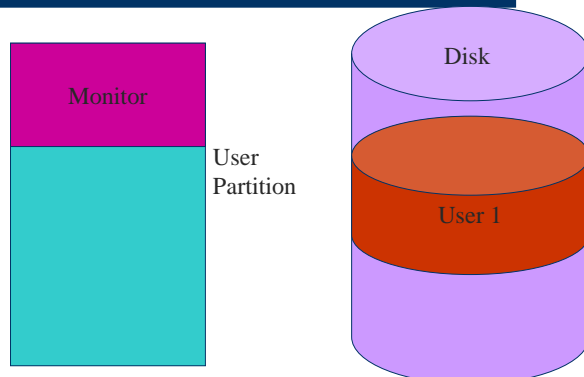
- Move a part of or the whole process to disk
- Allows several processes to share a fixed partition
- Processes that grow can be swapped out and swapped back in a bigger partition

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Swapping

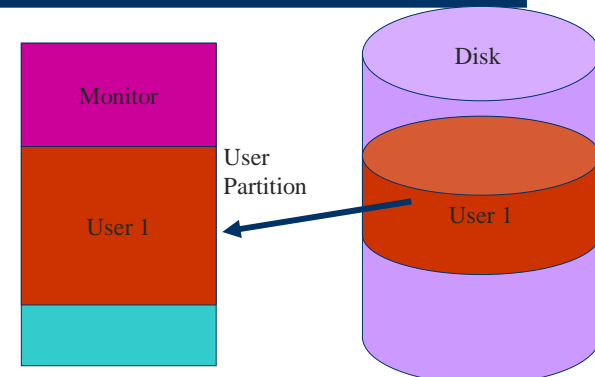


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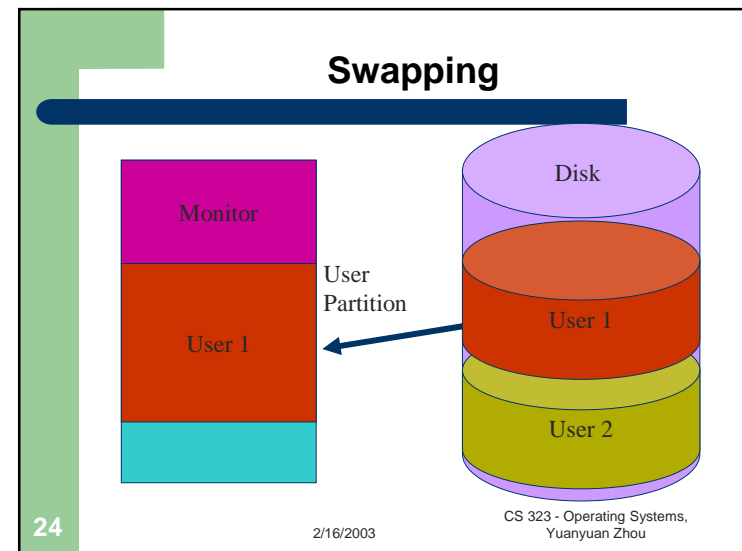
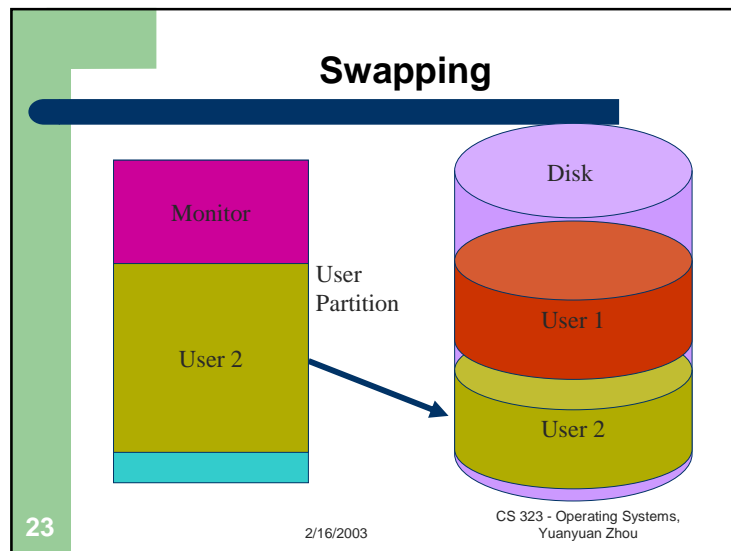
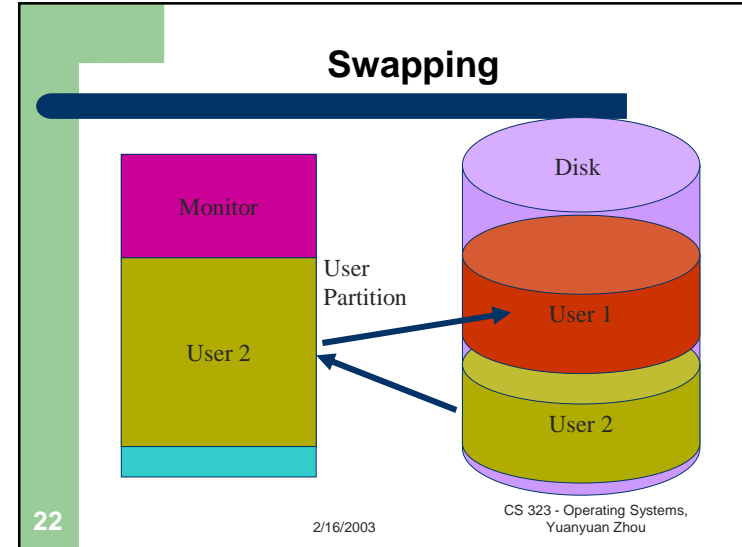
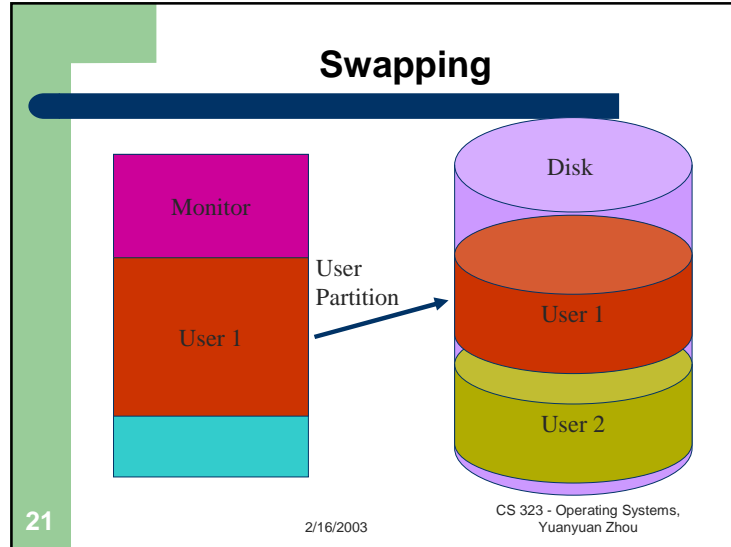
Swapping



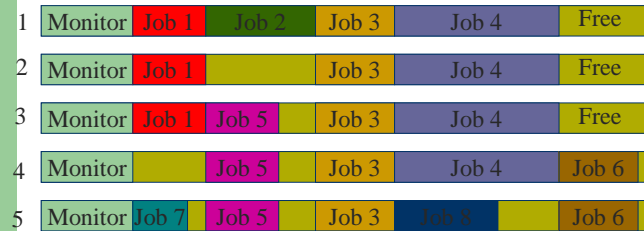
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Variable Partitions and Fragmentation



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Reminder

- Next lecture: Memory Management
- MP2

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