

7. Consider the following Grammar:

$$\begin{aligned} \langle S \rangle &::= wc\langle S \rangle \\ \langle S \rangle &::= \{\langle L \rangle\} \\ \langle S \rangle &::= s; \\ \langle L \rangle &::= \langle L \rangle \langle S \rangle \\ \langle L \rangle &::= \epsilon \end{aligned}$$

Part 1:

Letting $\langle S \rangle$ be the start symbol, list all possible strings consisting only of terminals for productions that can be reached by applying:

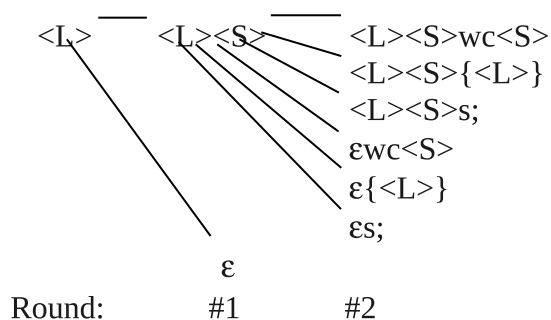
- a) 1 round of productions
- b) 2 rounds of productions
- c) 3 rounds of productions.

Part 2:

Do the same for $\langle L \rangle$, assuming that it is the start symbol.

A “round of productions” applied to a string s of terminal and non-terminal symbols is a set of strings of terminal and non-terminal symbols that can be reached by applying productions to all non-terminals in s . In subsequent rounds, start with all strings of terminals and non-terminals generated in the previous round.

To get you started, suppose we start with $\langle L \rangle$ and do two rounds.



The answer for 1 round of productions for $\langle L \rangle$ is ϵ , and the answer for 2 rounds of productions is $s;$. All other generated strings have one or more non-terminals in them and are therefore not “strings consisting only of terminals.” Note that when we write strings of all terminals in our answer, we drop ϵ unless it stands alone.