

Morse theory and persistent homology of geometric complexes

Ulrich Bauer

TUM

February 9, 2023

Funded by



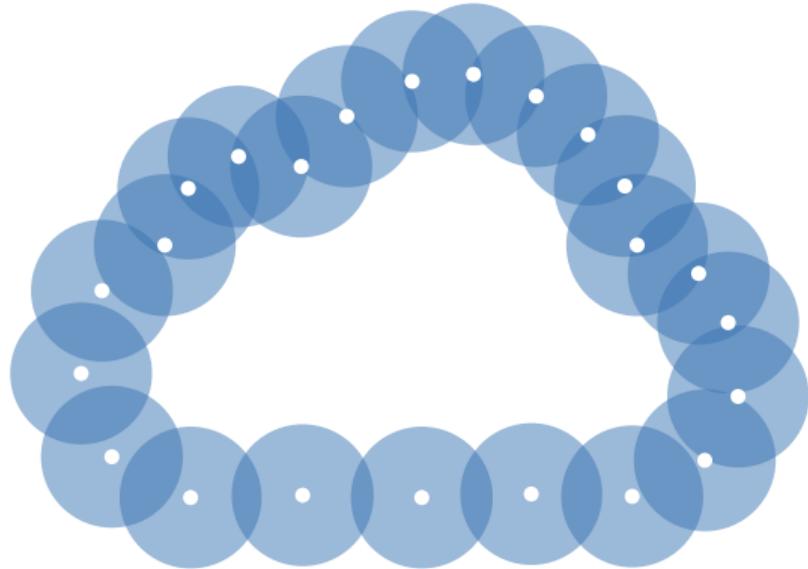
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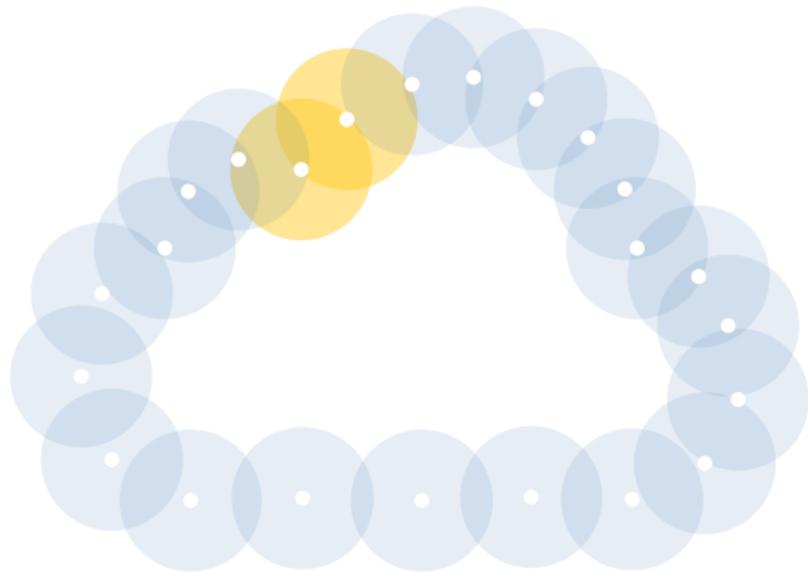
SFB
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Discretization
in Geometry
and Dynamics



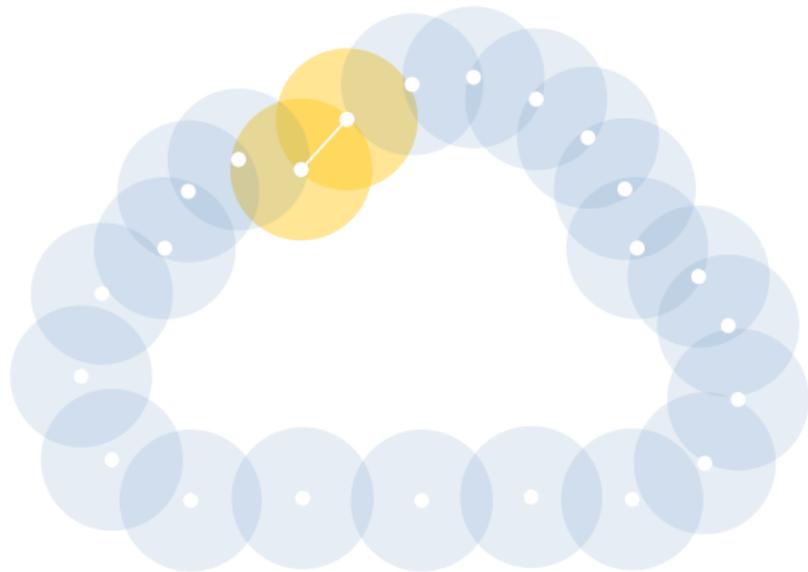
Čech complexes



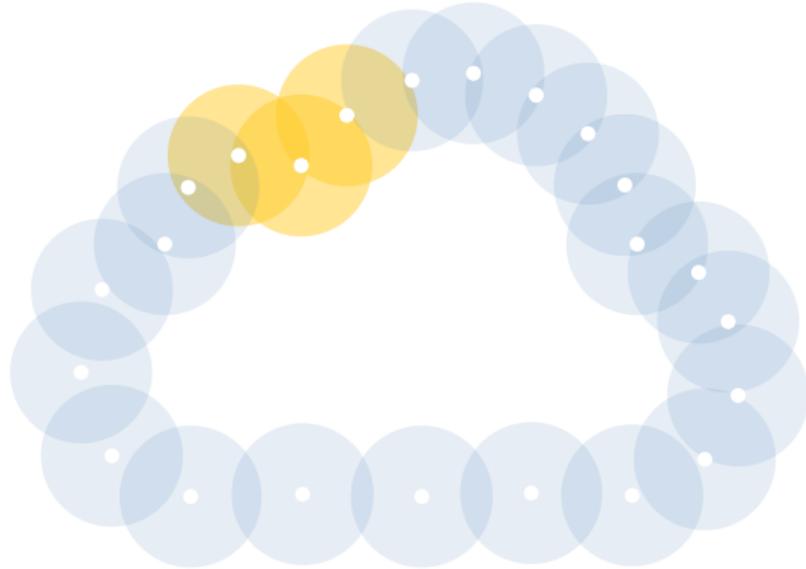
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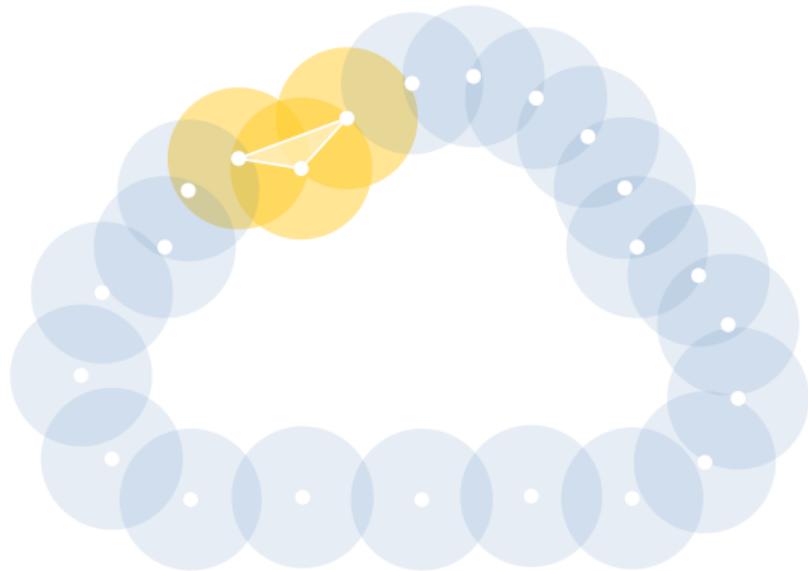
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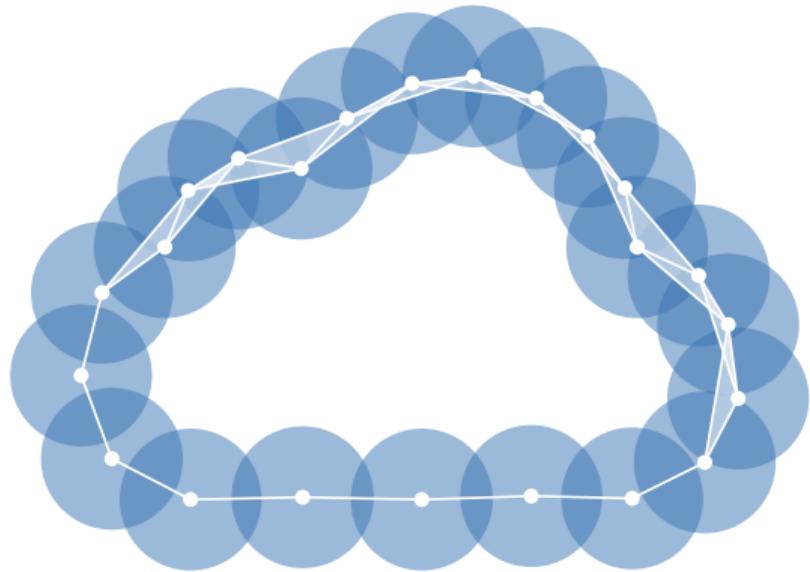
Čech complexes



Čech complexes



Čech complexes



Nerves

Definition (Alexandrov 1928)

Let $\mathcal{U} = (U_i)_{i \in I}$ be a cover of a space X . The *nerve* of \mathcal{U} is the simplicial complex

$$\text{Nrv}(\mathcal{U}) = \{J \subseteq I \mid |J| < \infty \text{ and } \bigcap_{i \in J} U_i \neq \emptyset\}.$$

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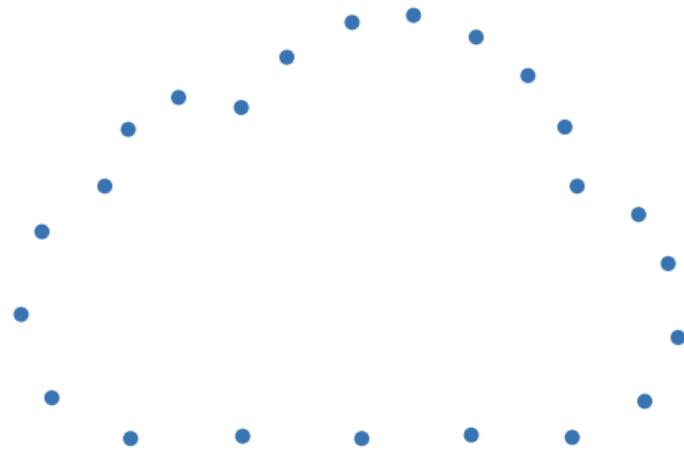


U. Bauer, M. Kerber, F. Roll, and A. Rolle

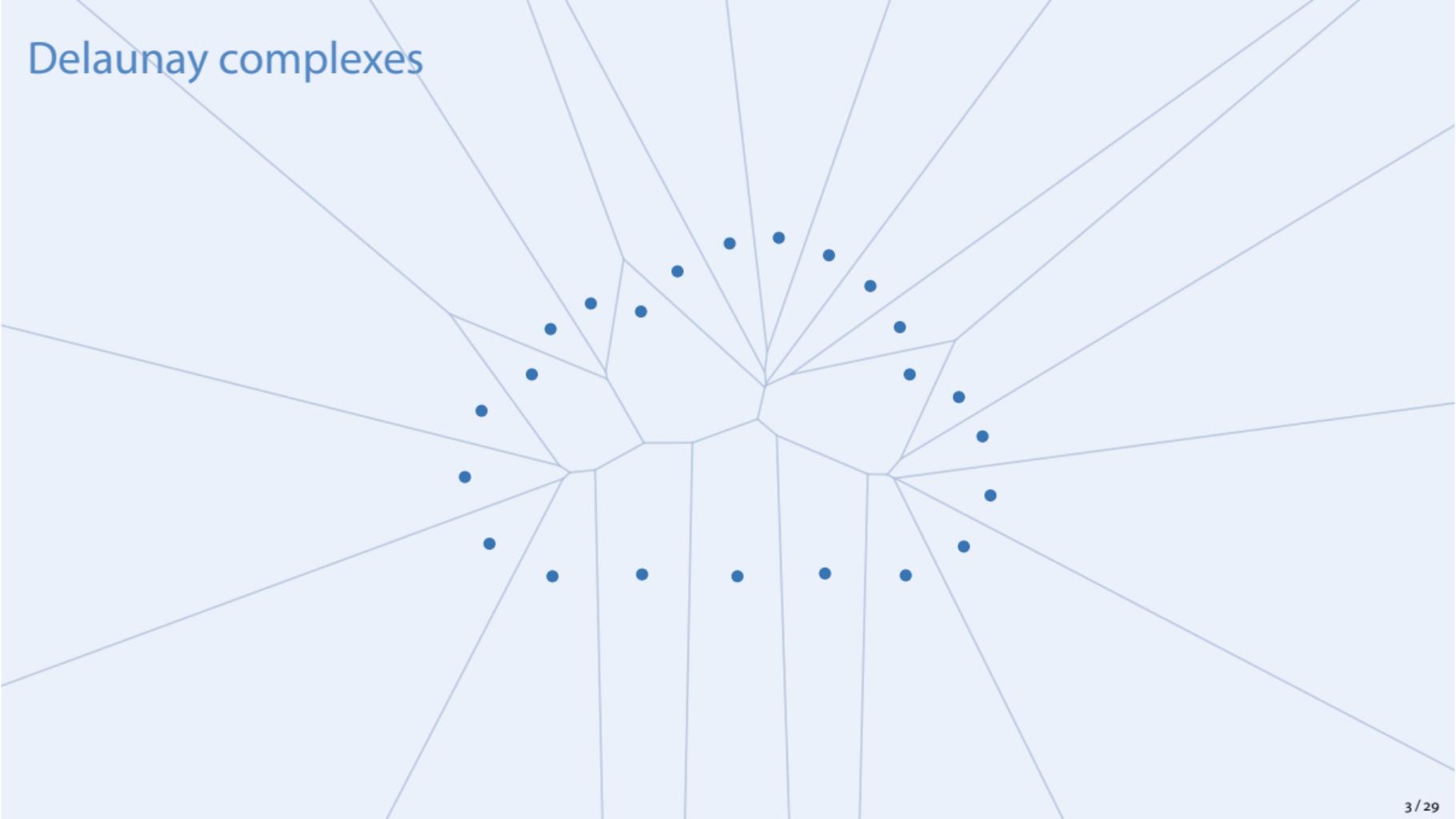
A Unified View on the Functorial Nerve Theorem and its Variations

Preprint, arXiv:2203.03571, 2022

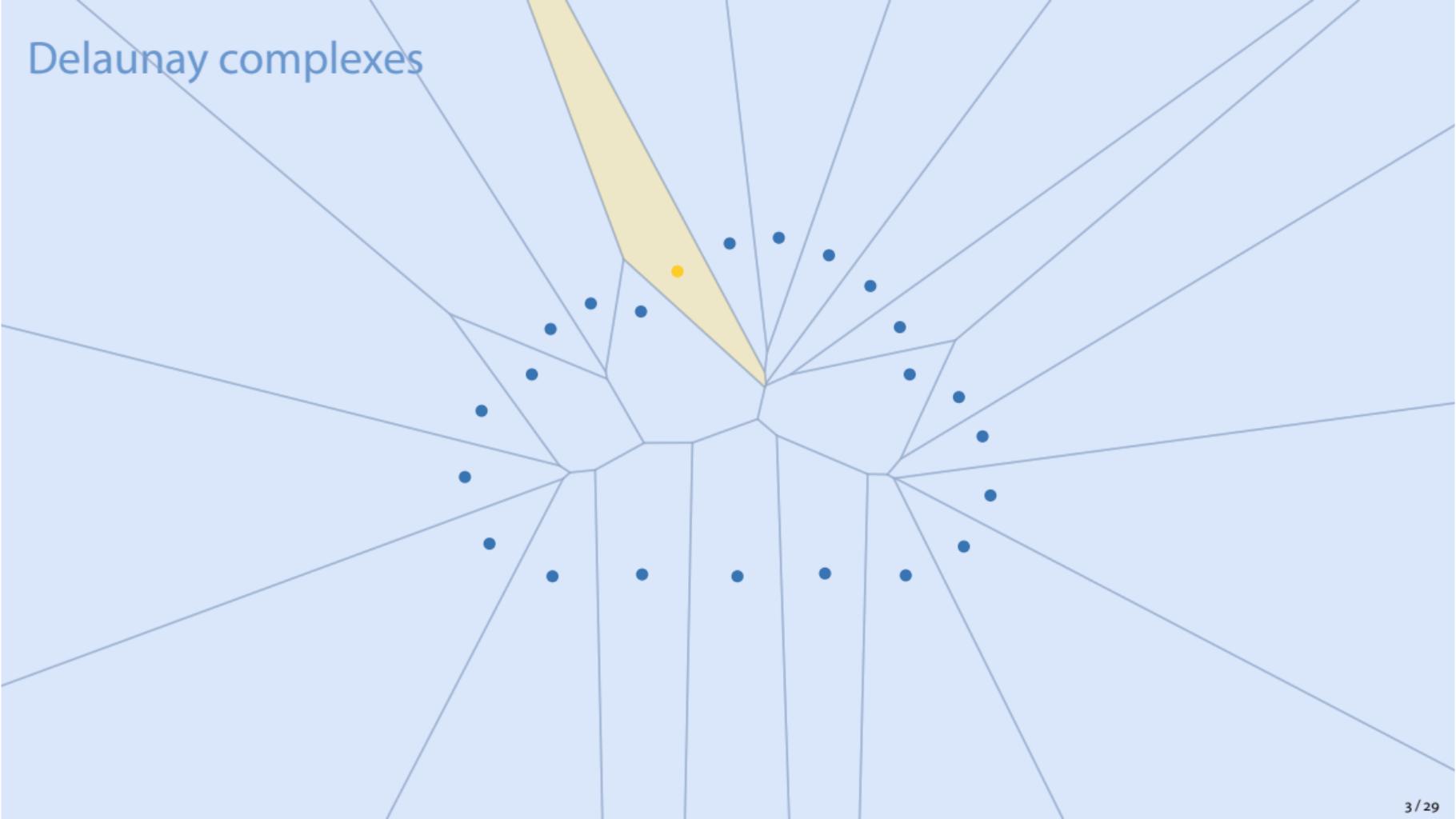
Delaunay complexes



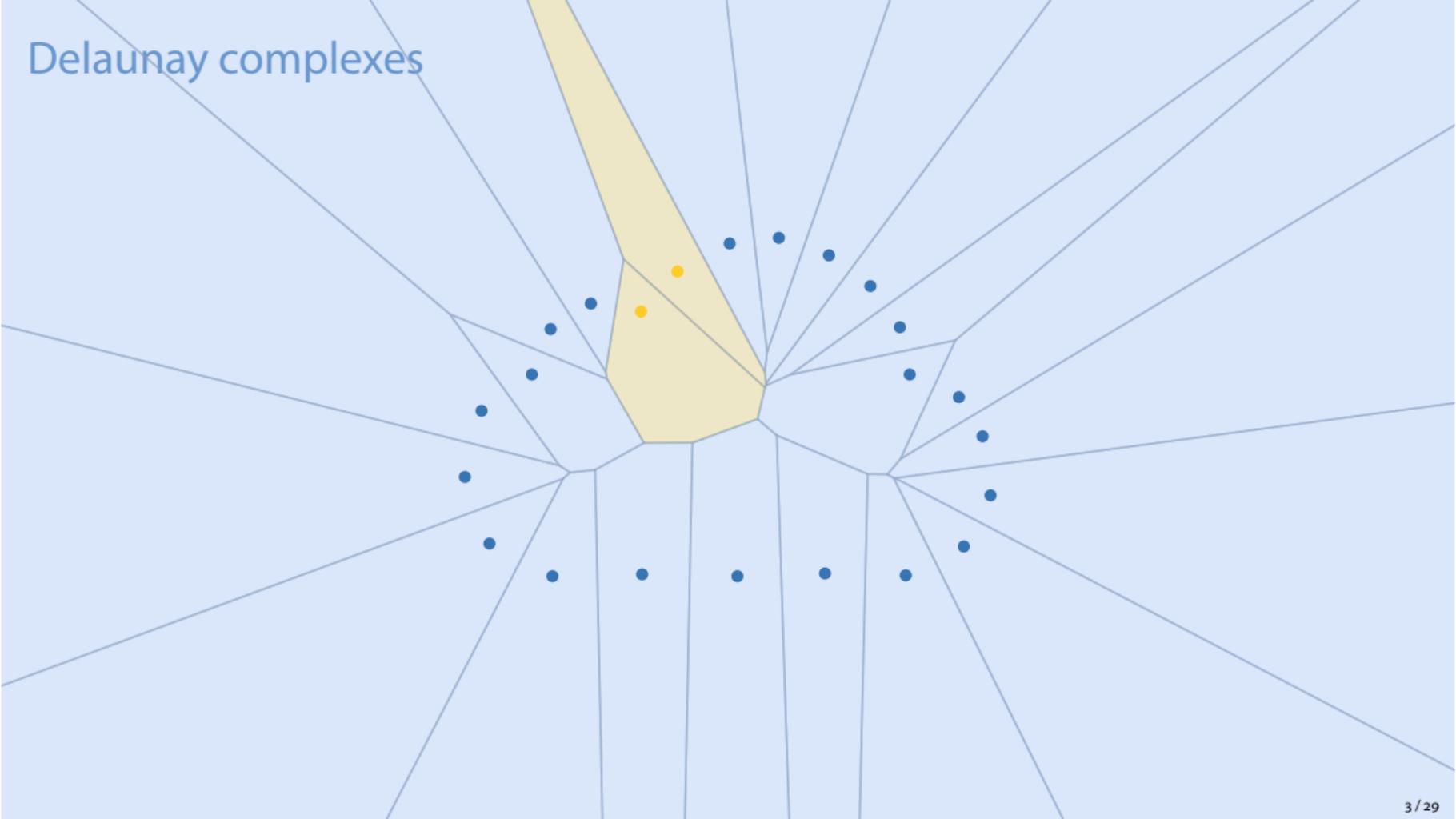
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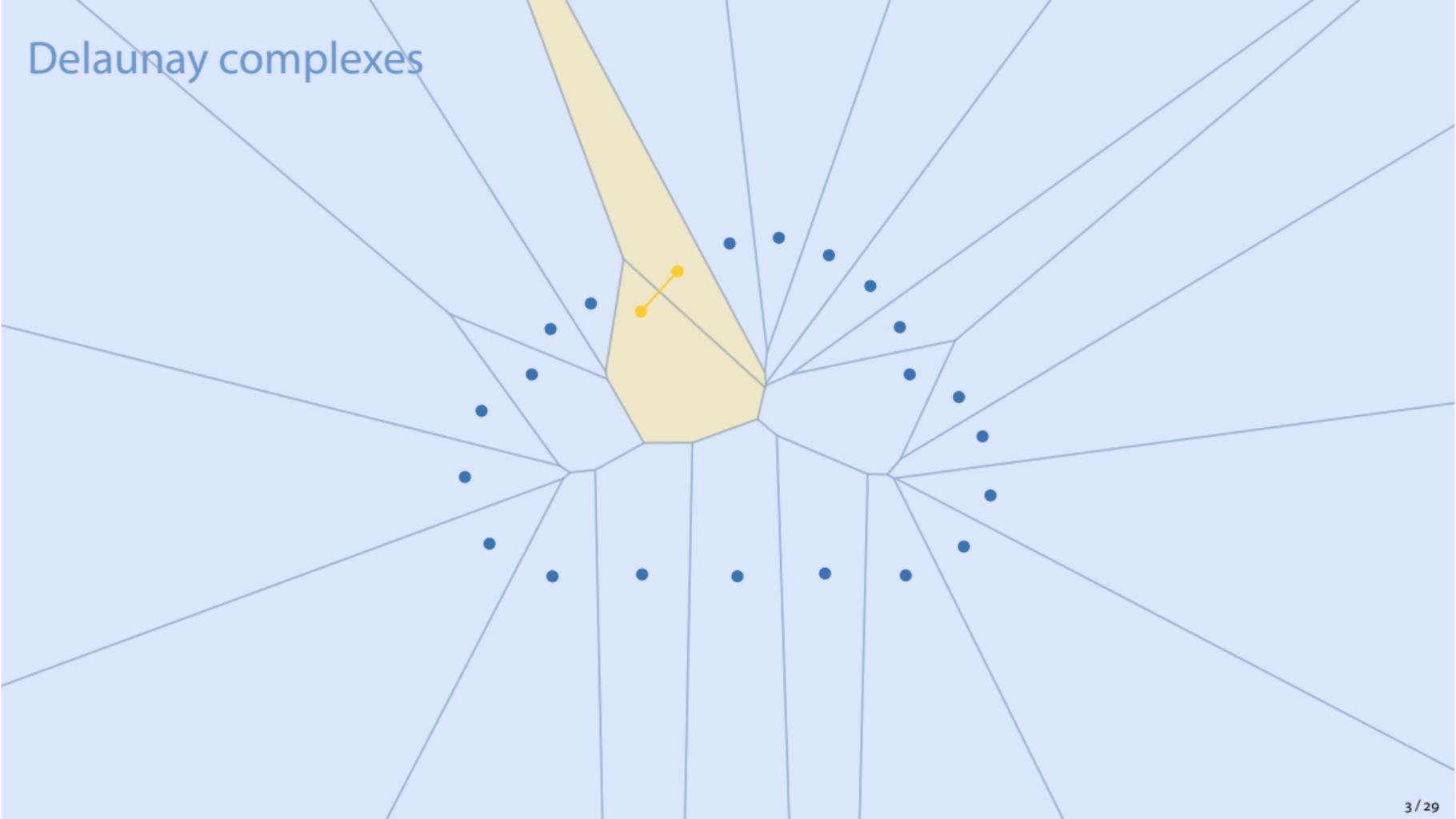
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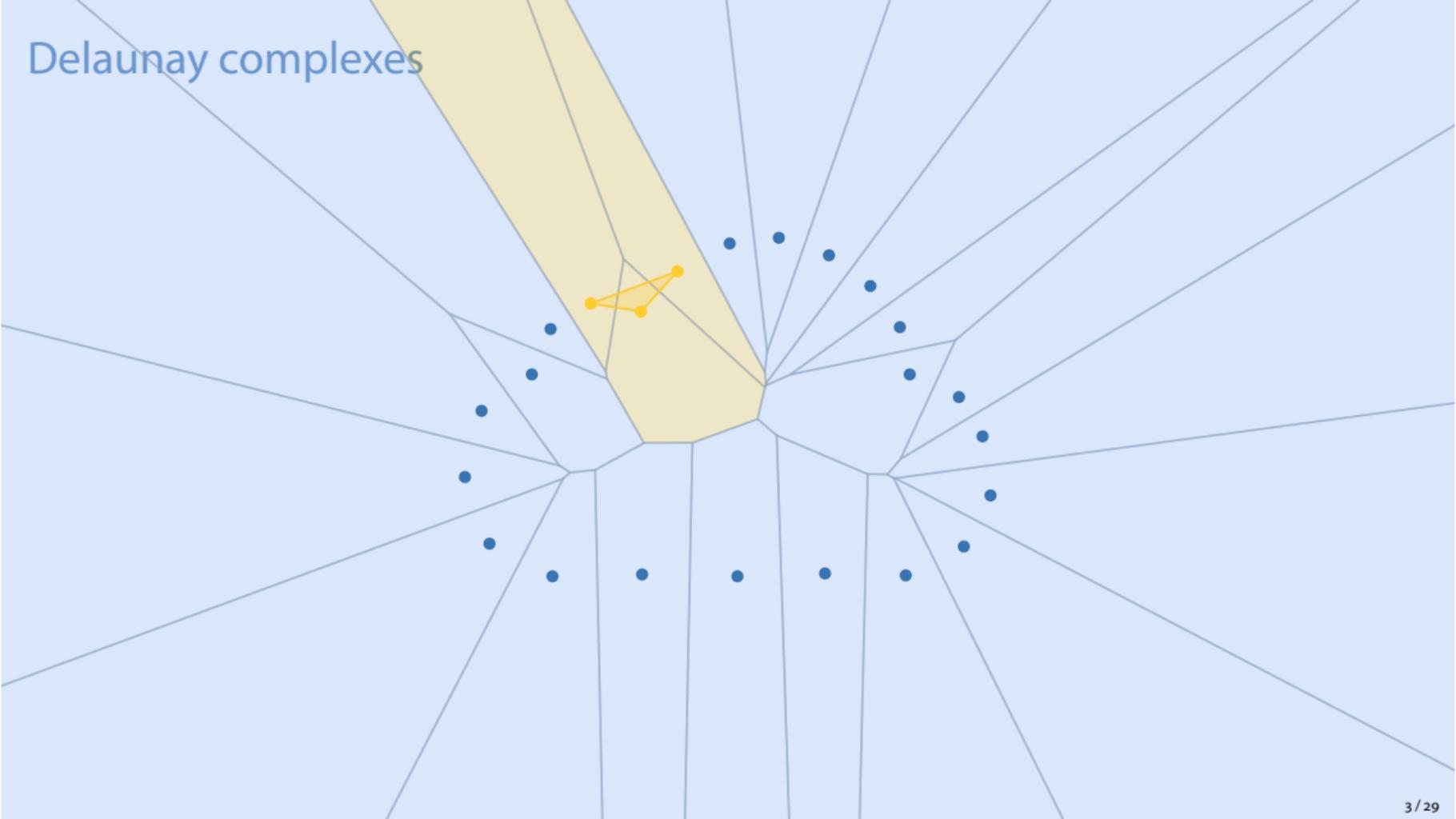
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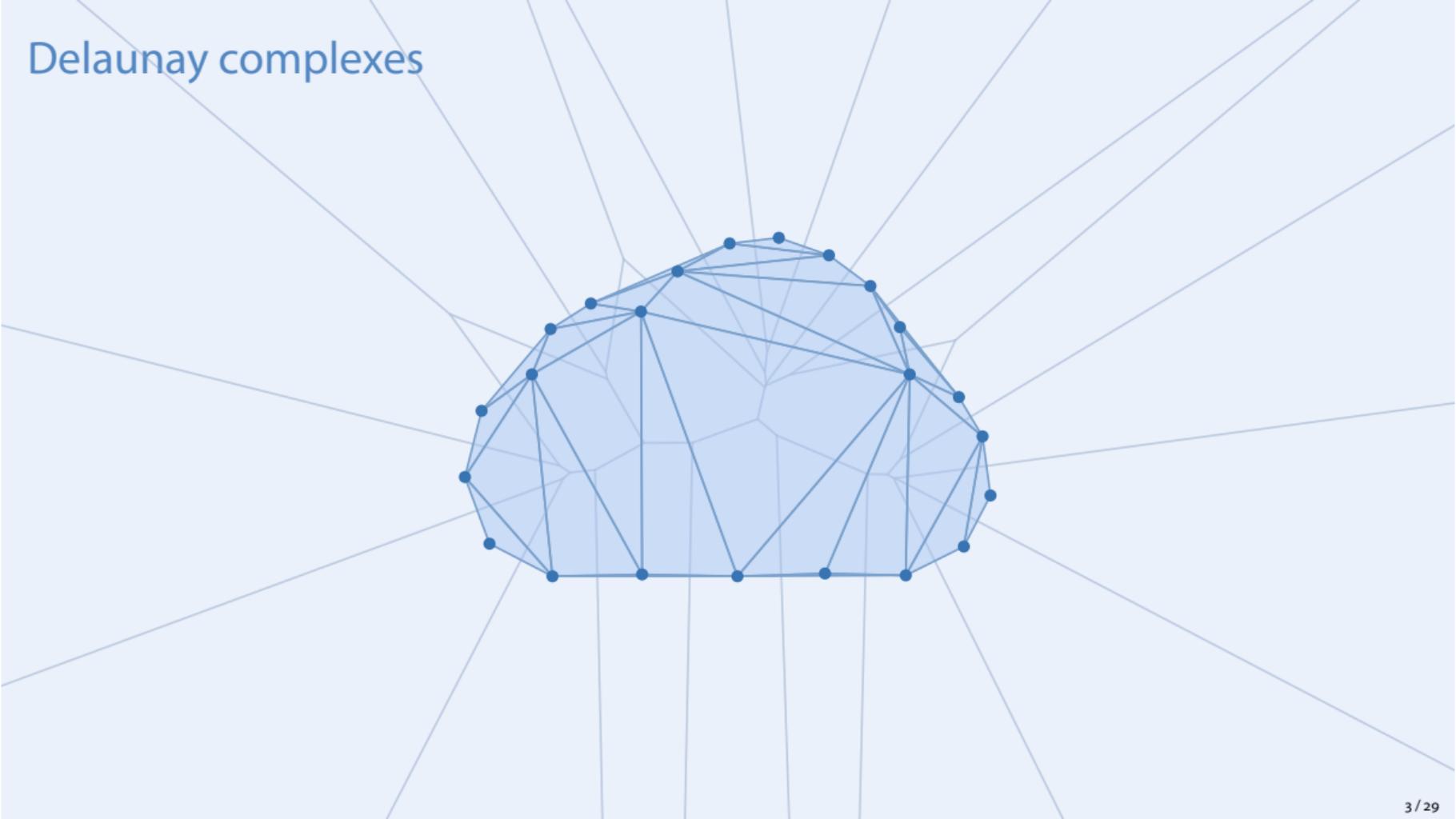
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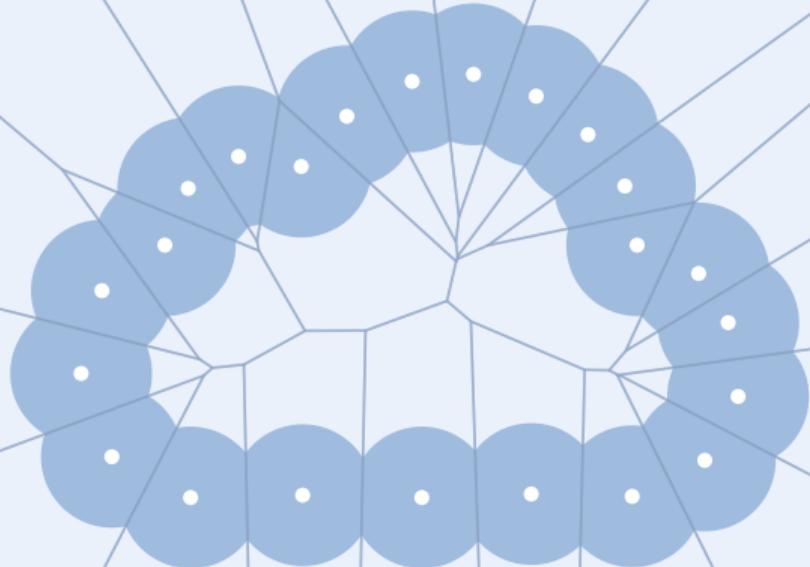
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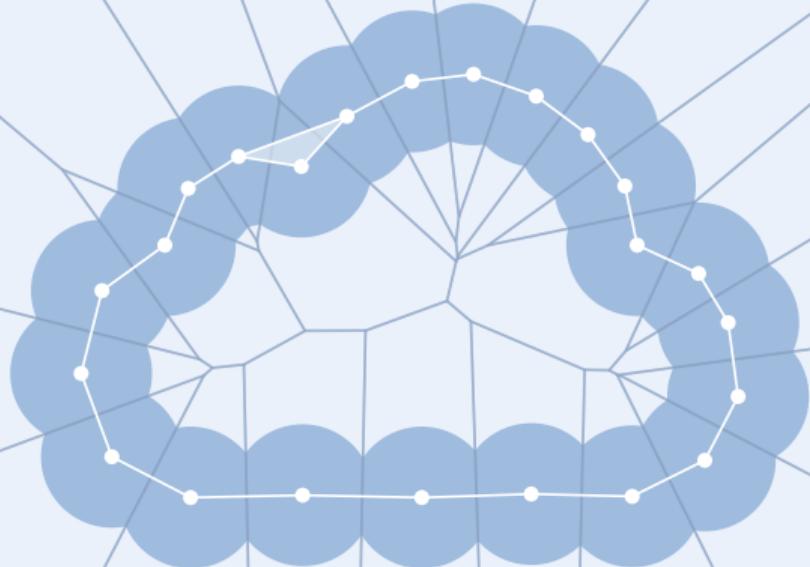
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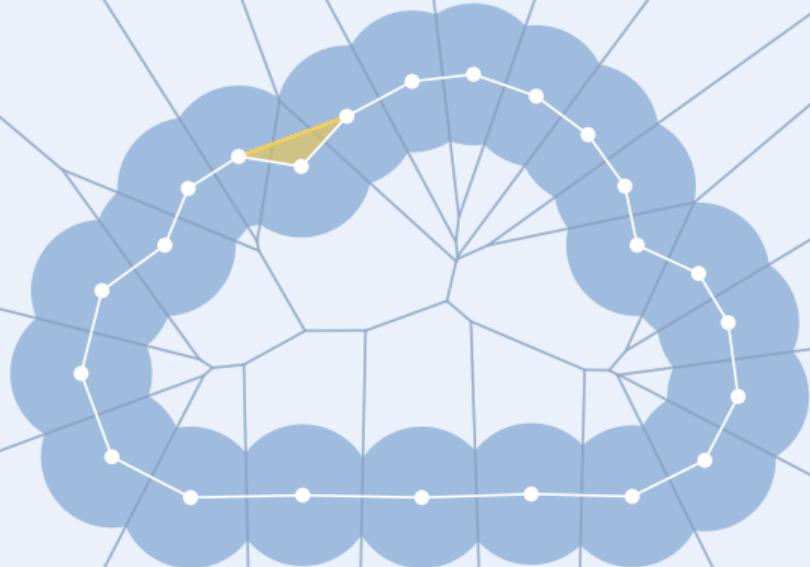
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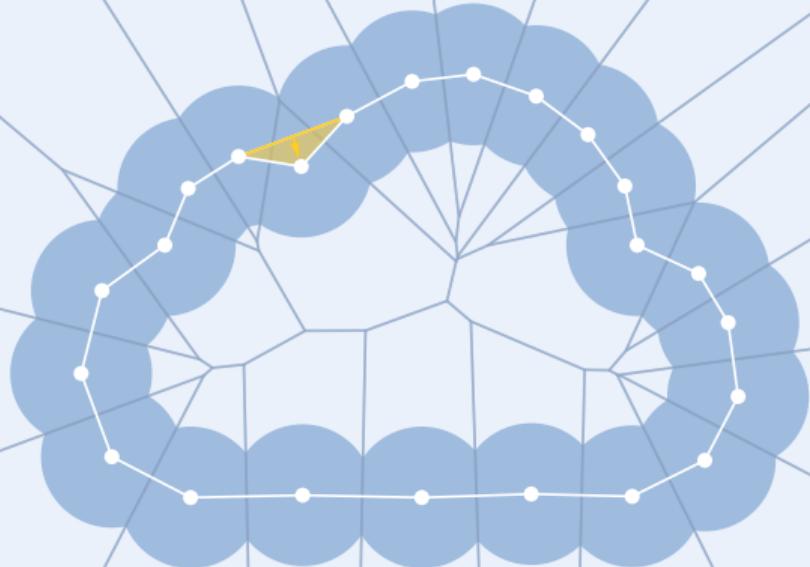
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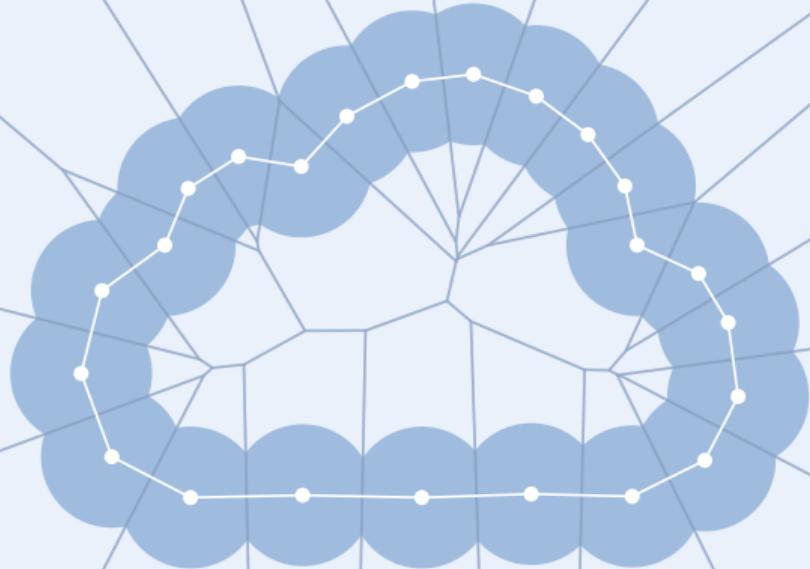
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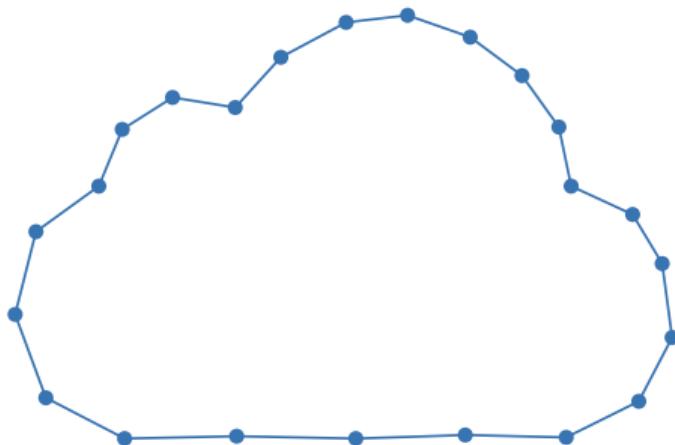
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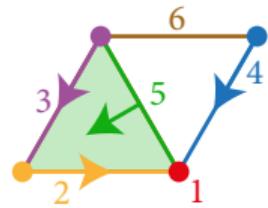
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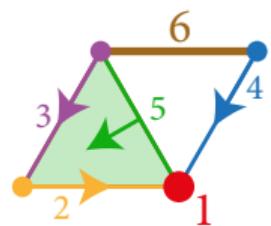
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Discrete Morse theory



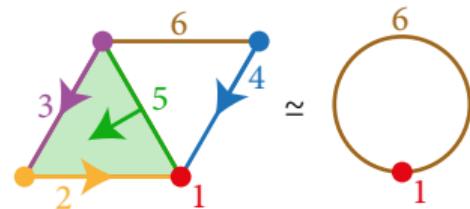
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Theorem (Forman 1998)

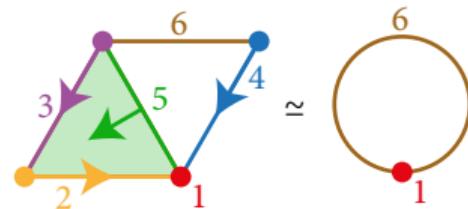
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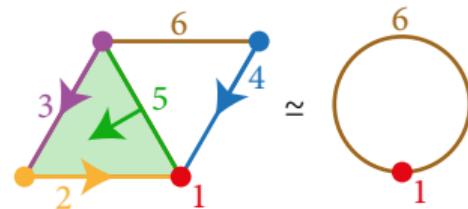
Discrete Morse functions – and their gradients – encode collapses of sublevel sets:



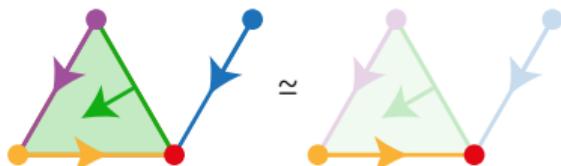
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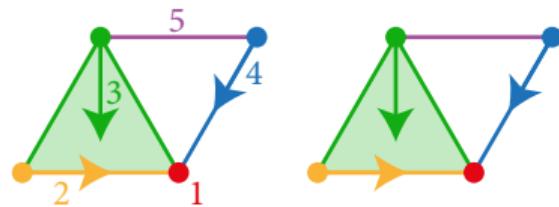


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Generalizing discrete Morse theory

Generalized discrete Morse functions/gradients:



Morse theory for Čech and Delaunay complexes

Proposition (B, Edelsbrunner 2014)

The Čech complexes and the Delaunay complexes (alpha shapes) are sublevel sets of (generalized) discrete Morse functions.

Morse theory for Čech and Delaunay complexes

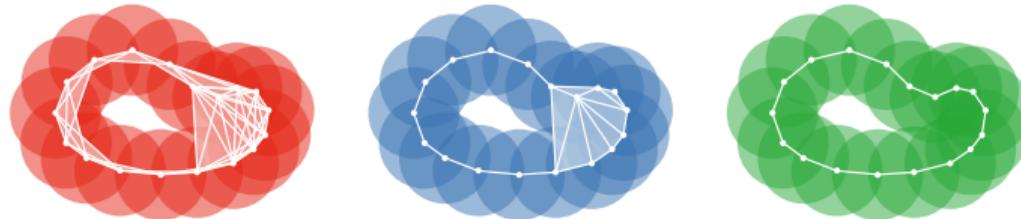
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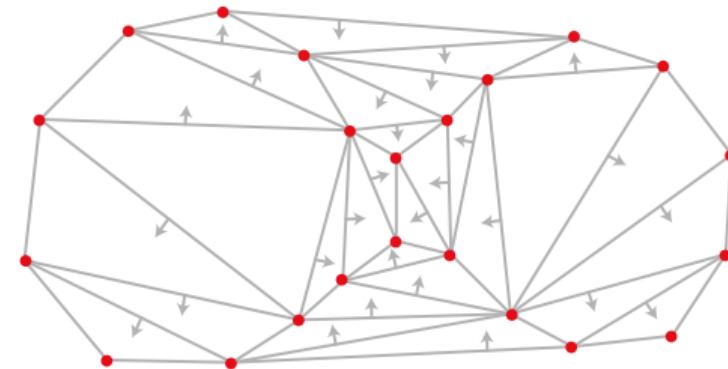
Čech, Delaunay, and Wrap complexes (at any scale r) are related by collapses encoded by a single discrete gradient field:

$$\text{Cech}_r X \searrow \text{Del}_r X \searrow \text{Wrap}_r X.$$



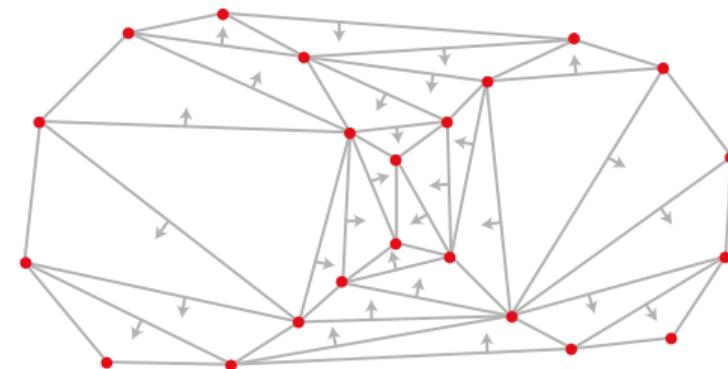
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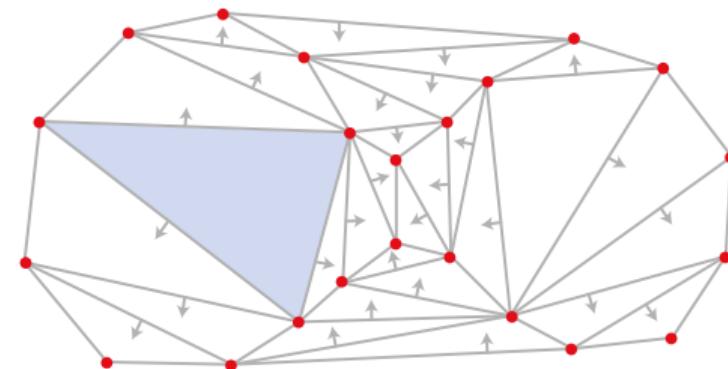
Definition (Edelsbrunner 1995; B, Edelsbrunner 2017)

Define $\text{Wrap}_r(X)$ as the smallest subcomplex of $\text{Del } X$ that

- contains all critical simplices with circumradius $\leq r$ and
- is a union of face intervals of V .

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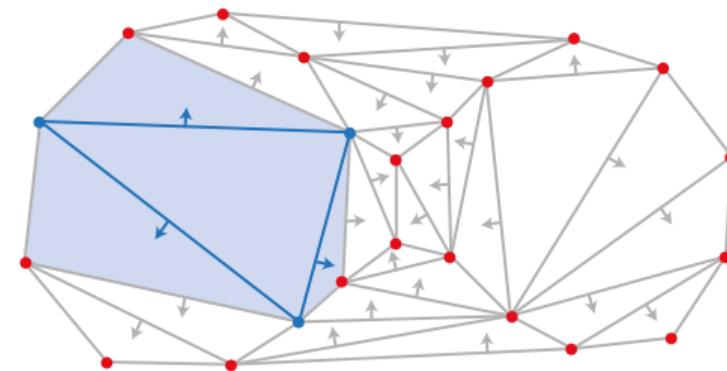
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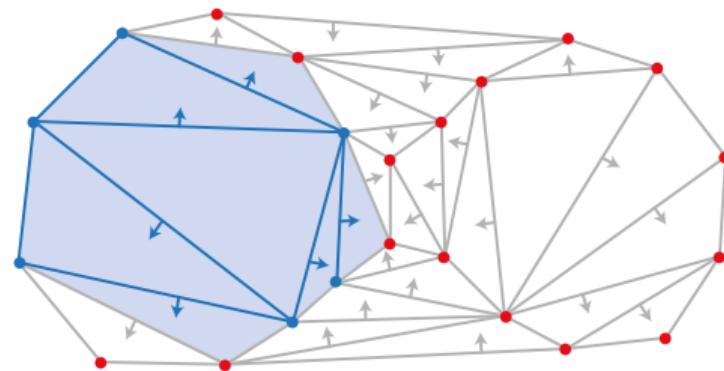
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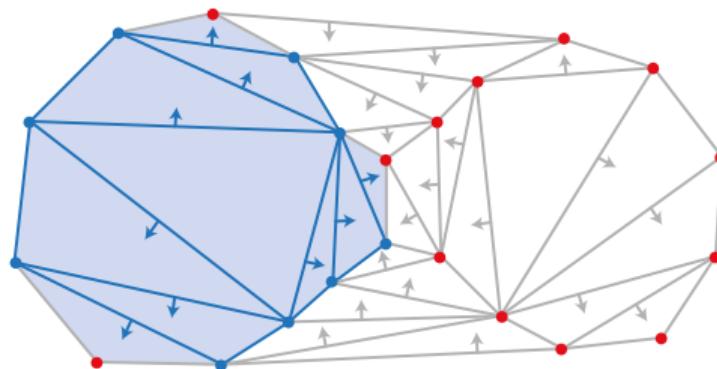
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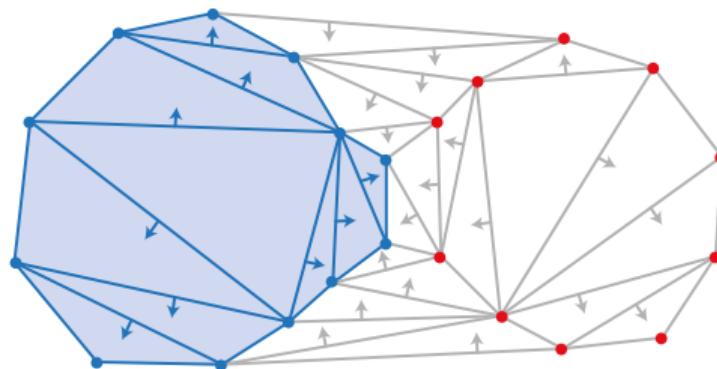
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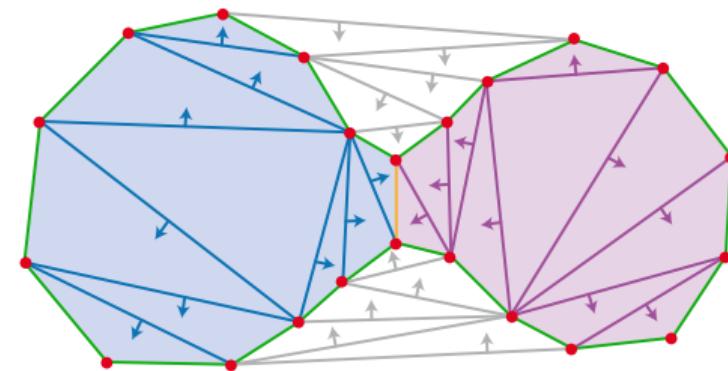
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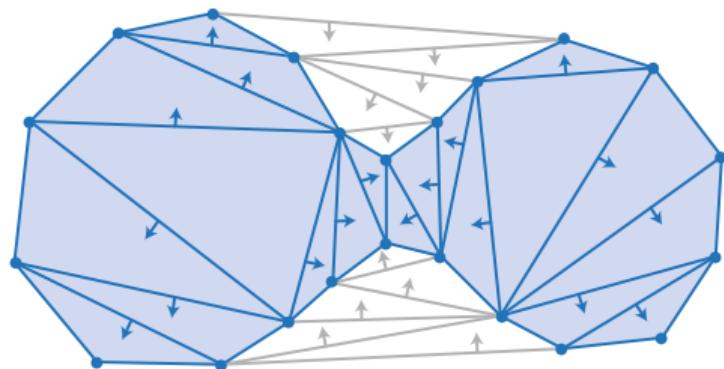
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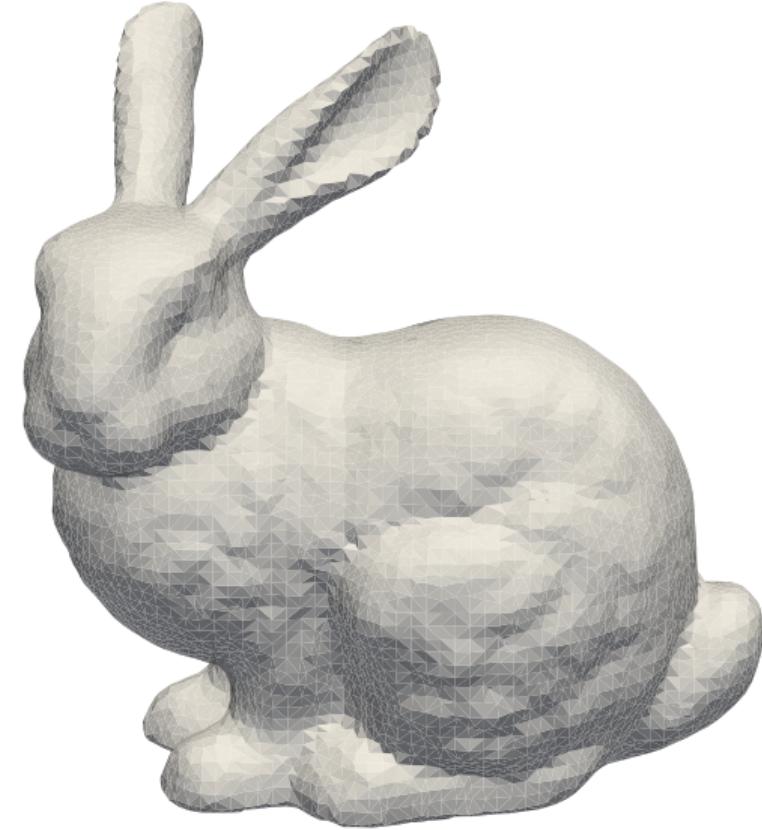
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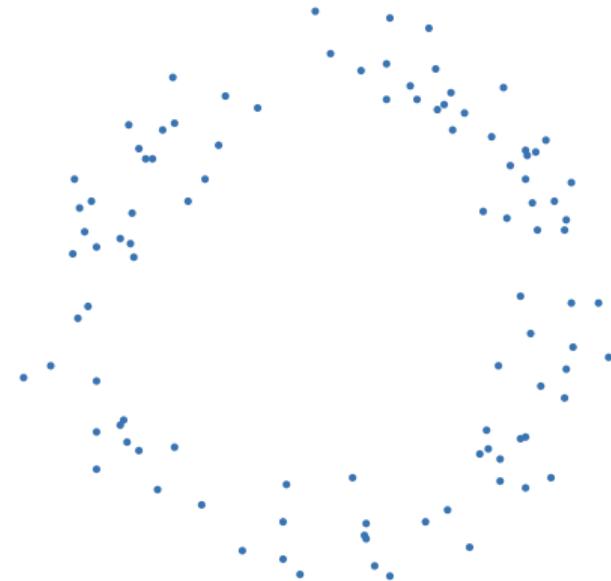
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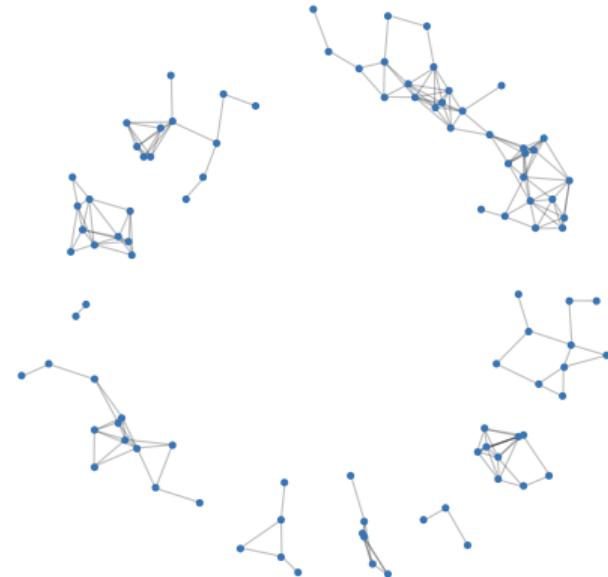
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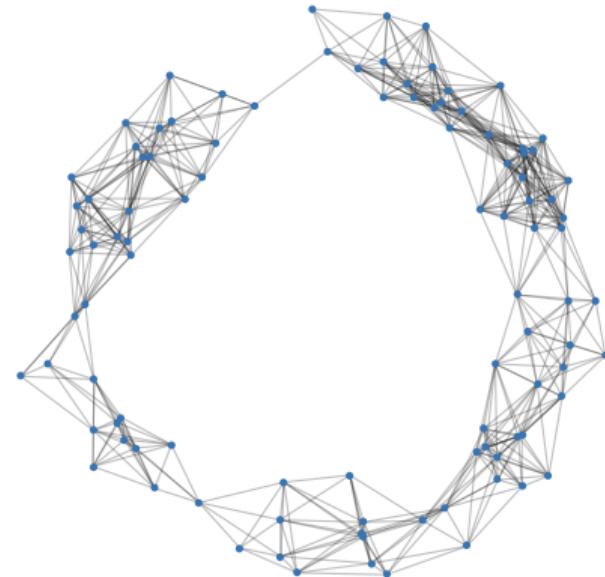
Vietoris–Rips complexes



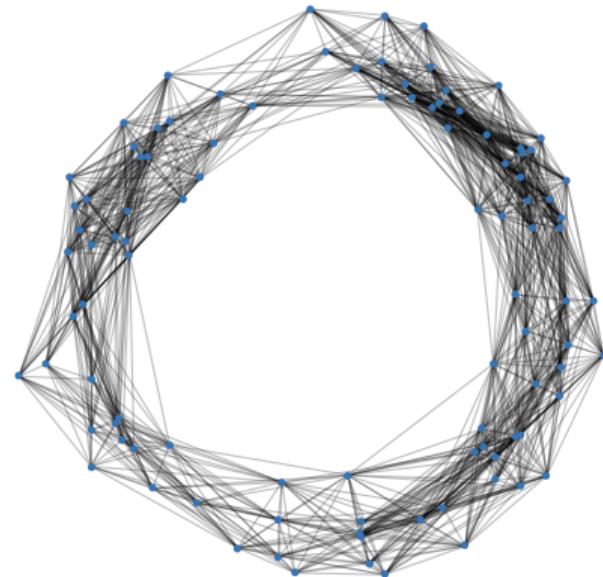
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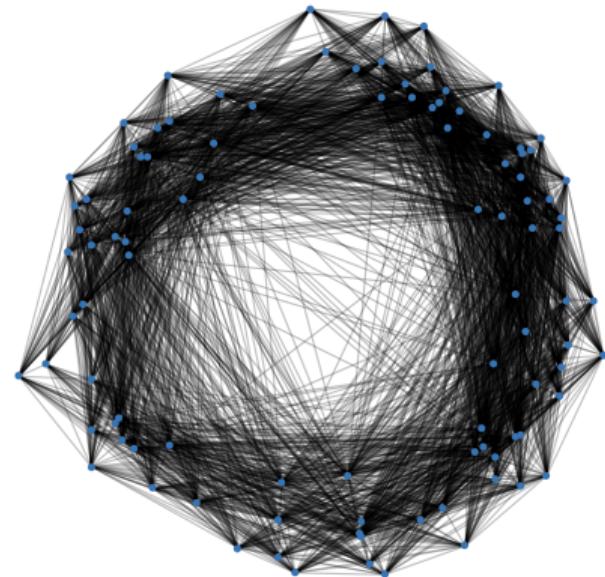
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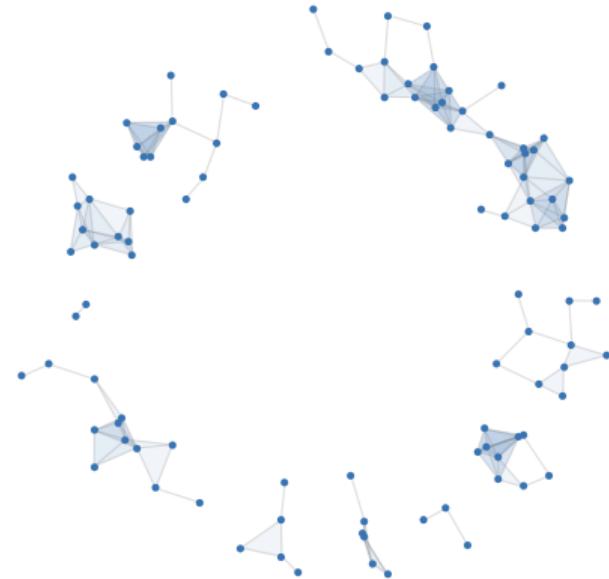
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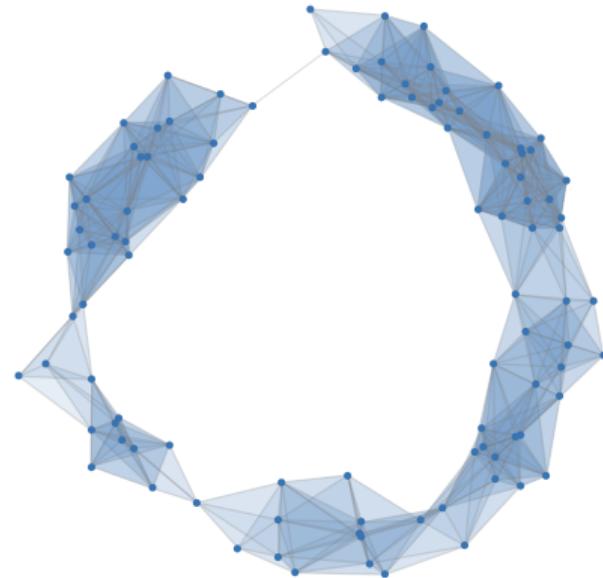
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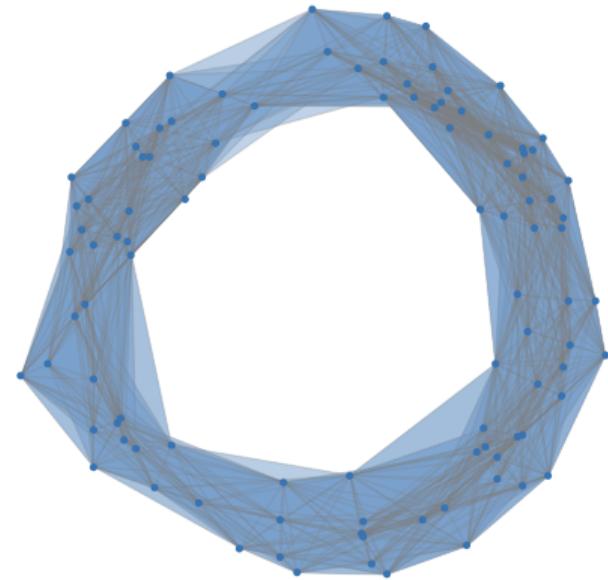
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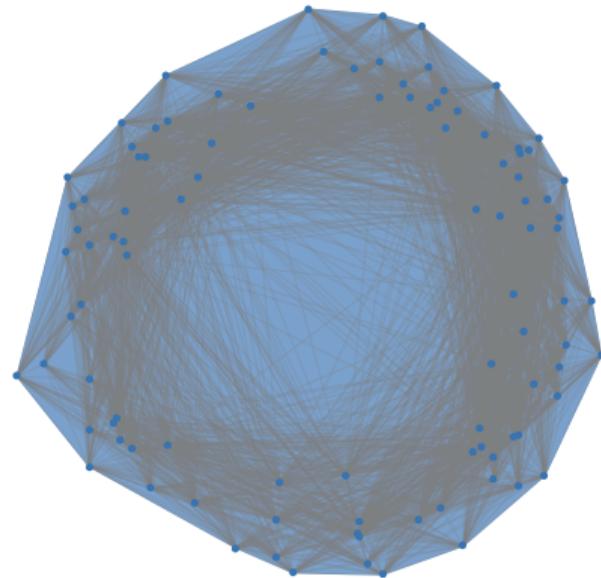
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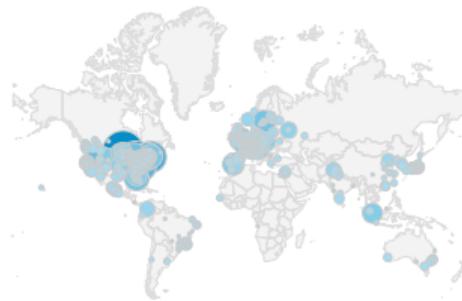
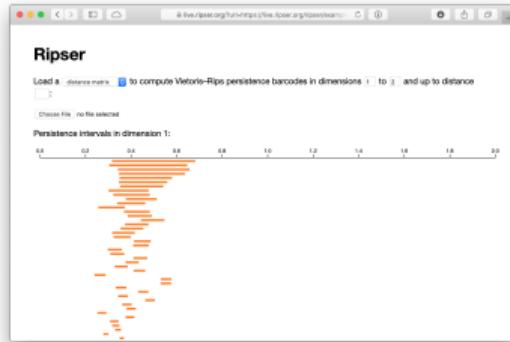
For a metric space X , the *Vietoris–Rips complex* at $t > 0$ is the simplicial complex

$$\text{Rips}_t(X) = \{S \subseteq X \mid S \neq \emptyset \text{ finite}, \text{diam } S \leq t\}.$$

Ripser: software for computing Vietoris–Rips persistence barcodes

Open source software (ripser.org)

- significantly faster / more memory efficient than previous codes
- de facto standard for topological data analysis applications
- most popular persistent homology project on GitHub



Ripser users from 935 different locations

Computational improvements based on

- *implicit matrix representations*
- connecting persistence to *discrete Morse theory*

Apparent pairs

Ripser uses the following construction for a computational shortcut:

Definition

In a simplexwise filtration $(K_i = \{\sigma_1, \dots, \sigma_i\})_i$, two simplices (σ_i, σ_j) are an *apparent pair* if

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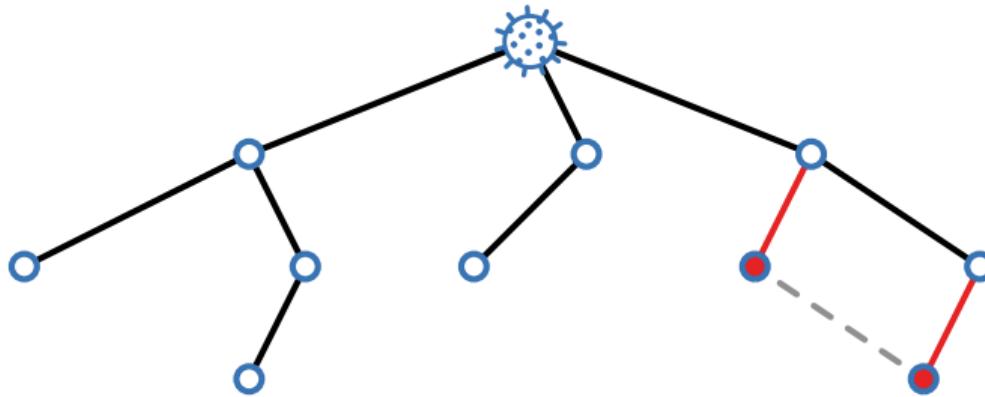
Proposition

The apparent pairs form a discrete gradient.

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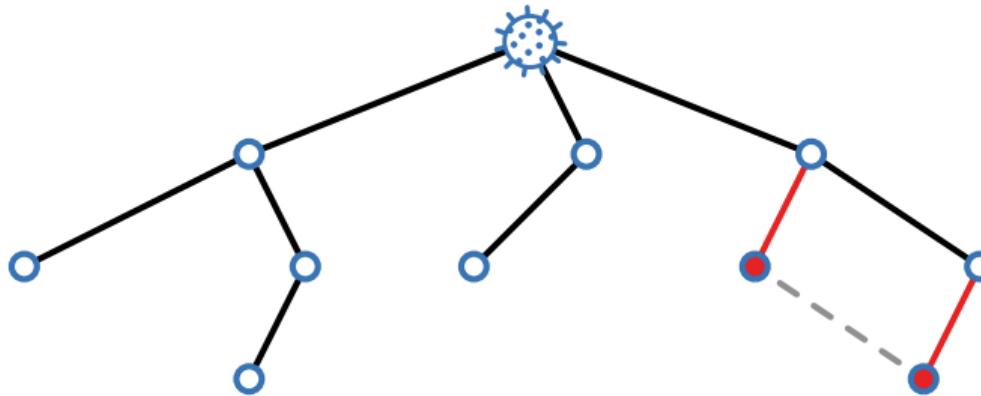
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Topology of viral evolution



Joint work with: A. Ott, M. Bleher, L. Hahn (Heidelberg), R. Rabadan, J. Patiño-Galindo (Columbia), M. Carrière (INRIA)

Topology of viral evolution



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Observation: Ripser runs unusually fast on genetic distance data

- SARS-CoV2 RNA sequences (spike protein)
- 25556 data points (2.8×10^{12} simplices in 2-skeleton)
- 120 s computation time (with data points ordered appropriately)

The Rips Contractibility Lemma

Theorem (Rips; Gromov 1988)

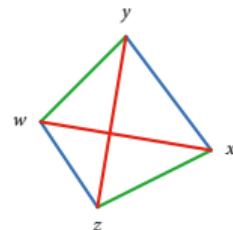
Let X be a δ -hyperbolic geodesic metric space. Then $\text{Rips}_t(X)$ is contractible for all $t \geq 4\delta$.

Gromov-hyperbolicity

Definition (Gromov 1988)

A metric space X is δ -hyperbolic if for all $w, x, y, z \in X$ we have

$$d(w, x) + d(y, z) \leq \max\{d(w, y) + d(x, z), d(w, z) + d(x, y)\} + 2\delta.$$

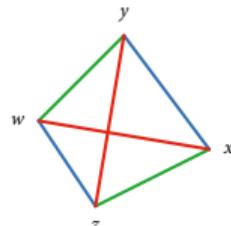


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- The hyperbolic plane is $\ln 2$ -hyperbolic

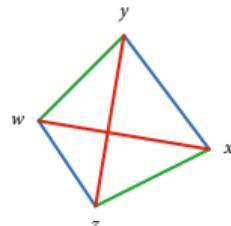


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- The hyperbolic plane is $\ln 2$ -hyperbolic



- 0-hyperbolic spaces are subspaces of trees

Rips contractibility for non-geodesic spaces

- What about non-geodesic spaces? In particular, finite metric spaces?
- Connection to apparent pairs/Ripser?

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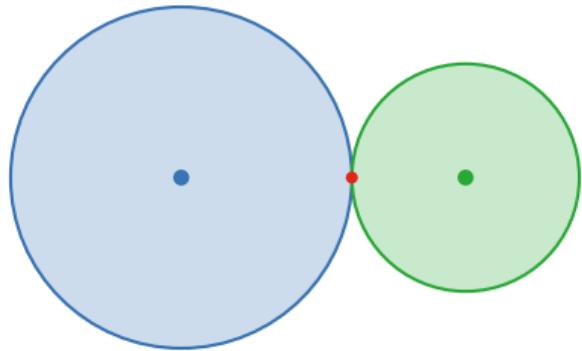
Theorem (B, Roll 2022)

Let X be a finite δ -hyperbolic space. Then there is a single discrete gradient encoding the collapses

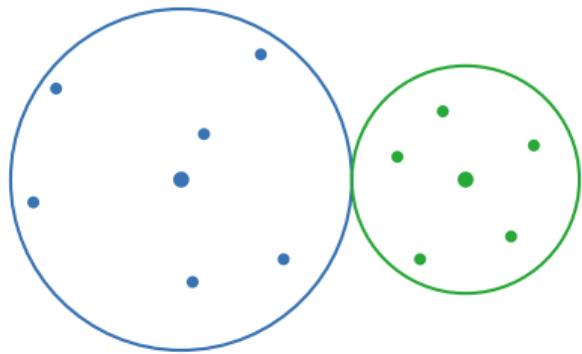
$$\text{Rips}_u(X) \searrow \text{Rips}_t(X) \searrow \{\ast\}$$

for all $u > t \geq 4\delta + 2\nu$, where ν is the geodesic defect of X .

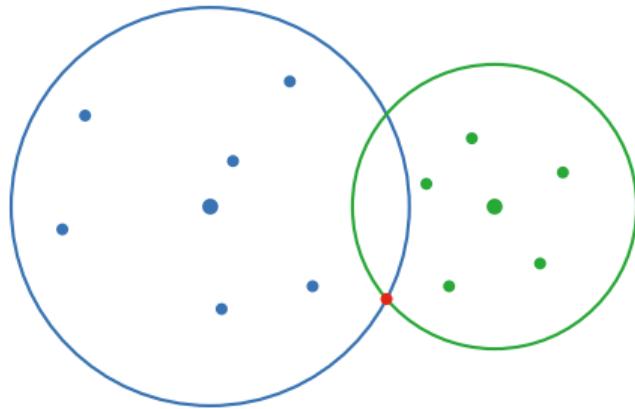
Geodesic defect



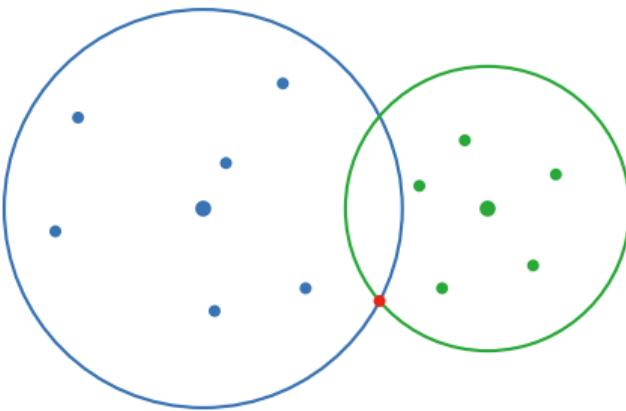
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Geodesic defect



Definition (Bonk, Schramm 2000)

A metric space X is v -geodesic if for all points $x, y \in X$ and all $r, s \geq 0$ with $r + s = d(x, y)$ we have

$$B_{r+v}(x) \cap B_{r+v}(y) \neq \emptyset.$$

The infimum of all such v is the *geodesic defect* of X .

Collapsing Vietoris–Rips complexes of generic trees

Consider a *generic* finite metric tree $T = (V, E)$ (distinct distances).

- $\text{diam}: 2^V \rightarrow \mathbb{R}$ is a generalized discrete Morse function.

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Theorem (B, Roll 2022)

This discrete gradient induces the collapses

$$\text{Rips}_t(X) \searrow T_t \quad \text{for all } t \in \mathbb{R}, \text{ with}$$

$$\text{Rips}_t(X) \searrow T \searrow \{\ast\} \quad \text{for all } t \geq \max l(E), \text{ and}$$

$$\text{Rips}_u(X) \searrow \text{Rips}_t(X) \quad \text{for all } (t, u] \cap l(E) = \emptyset.$$

In particular, the persistent homology is trivial in degrees > 0 .

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Example: phylogenetic trees

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However, the canonical gradient is not compatible with the apparent pairs gradient.

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Tie breaking for non-distinct pairwise distances:

- Choose total order on vertices
- Order edges lexicographically
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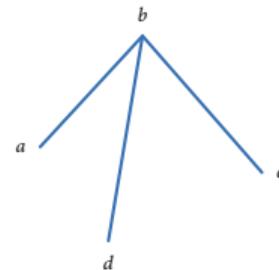
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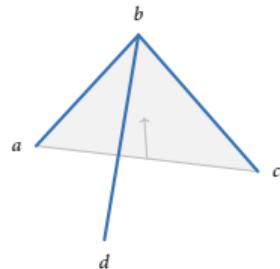


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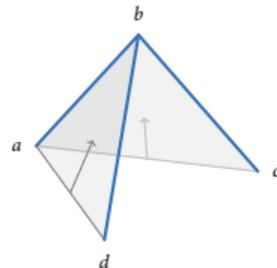


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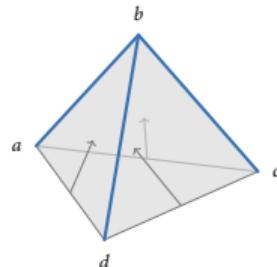


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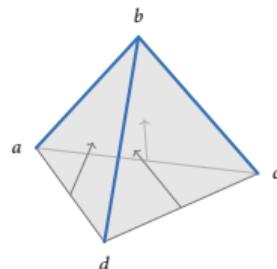


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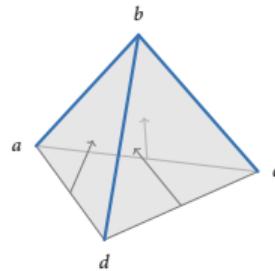
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Collapsing Rips complexes of trees with apparent pairs

Let X be the path length metric space for a weighted tree $T = (X, E)$.

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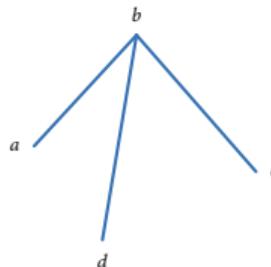
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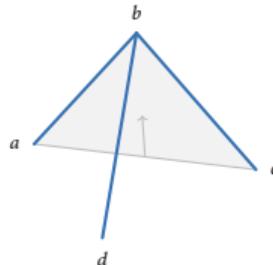
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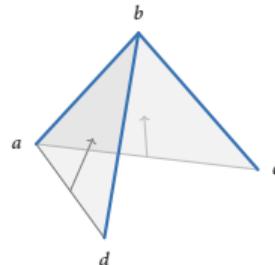
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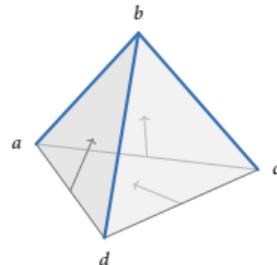
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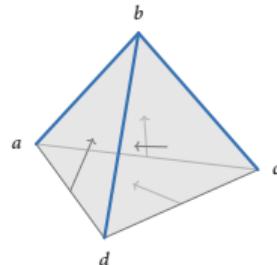
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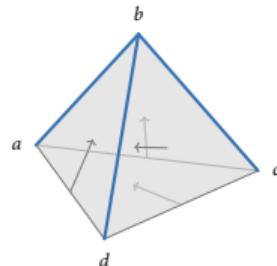
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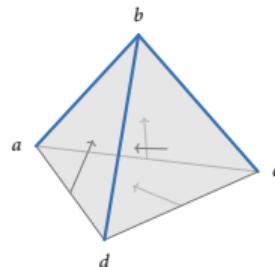
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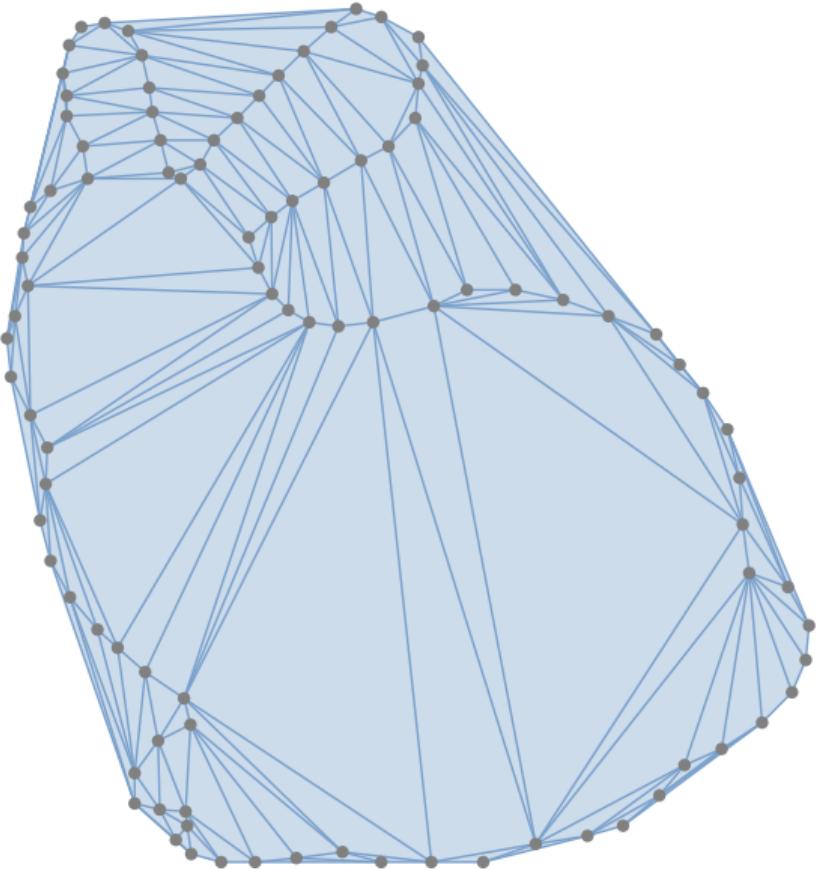
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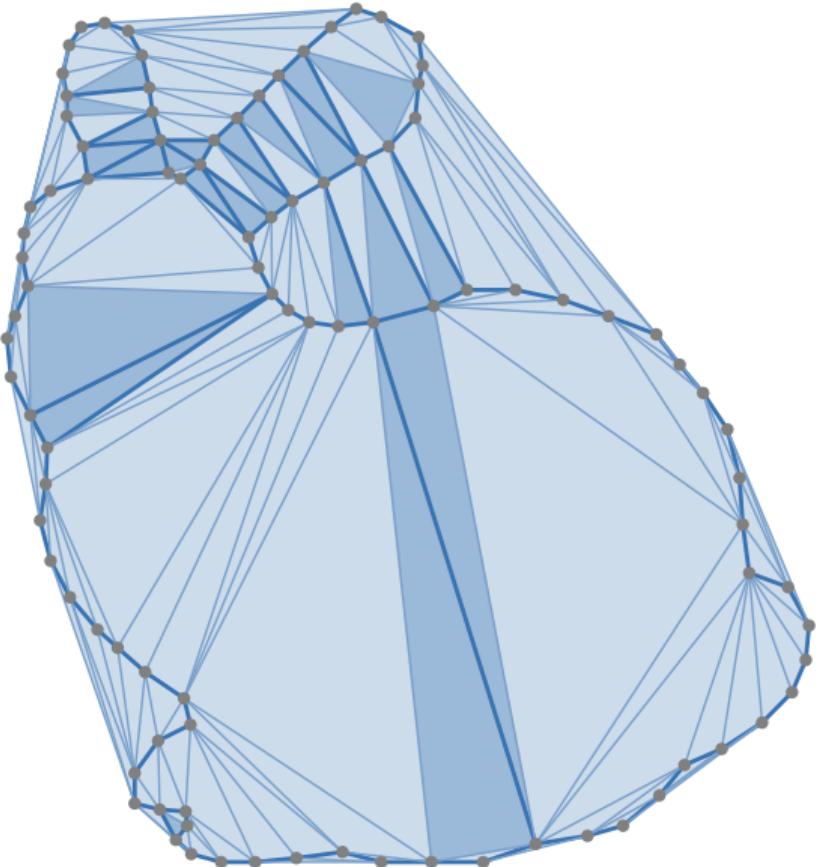
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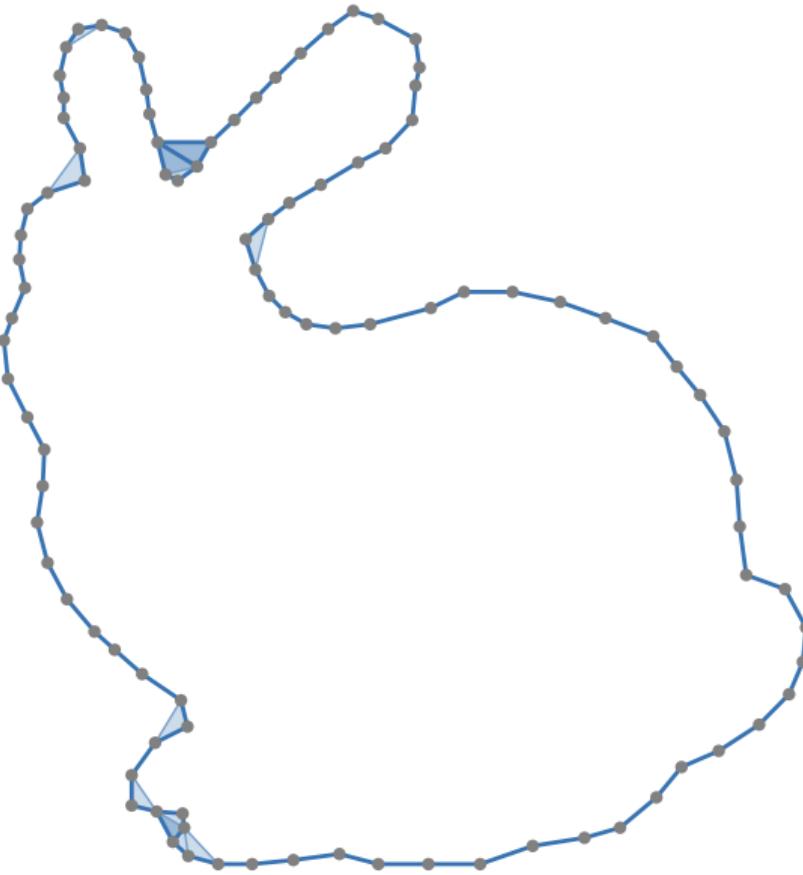
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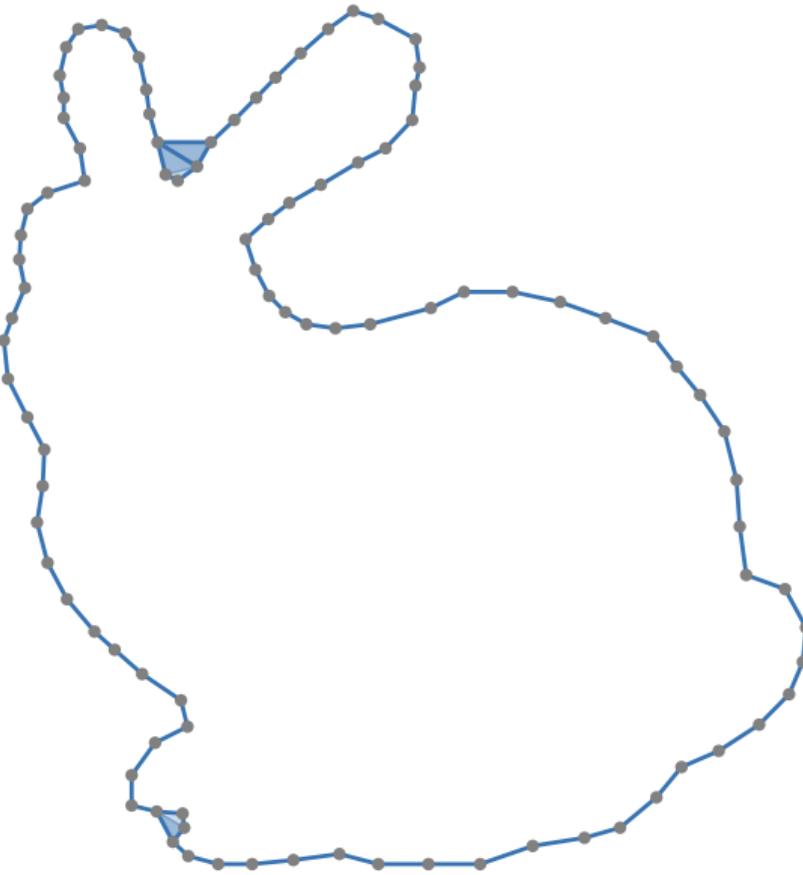


- Explains why Ripser is extraordinarily fast on genetic distances (tree-like)









Algebraic gradient flows and persistent homology

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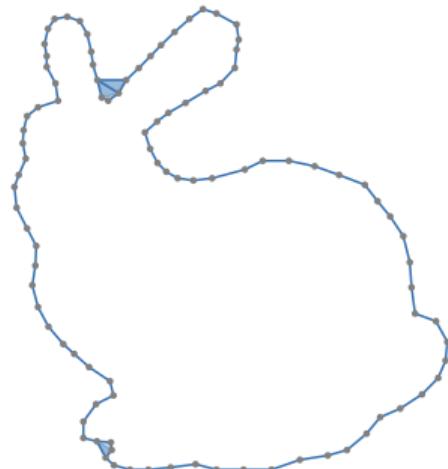
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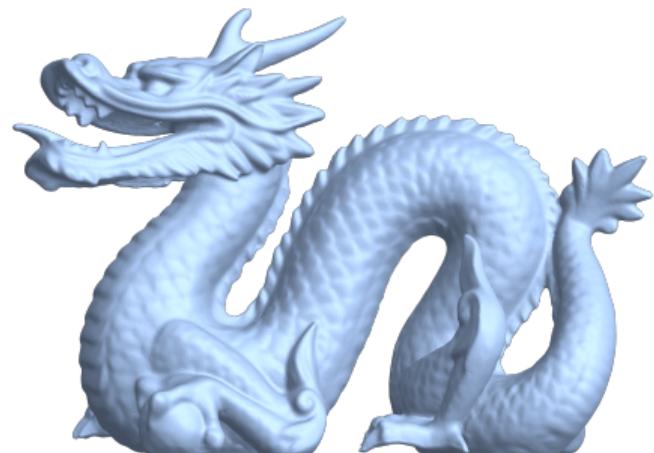
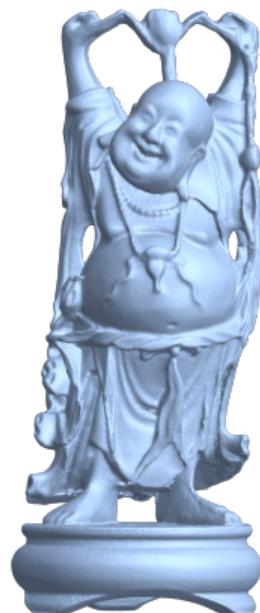
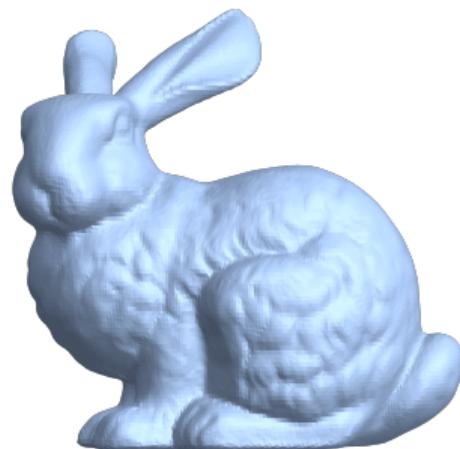
- *Persistence pairs correspond to algebraic gradient pairs.*
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- *The resulting cycle representatives are lexicographically optimal.*

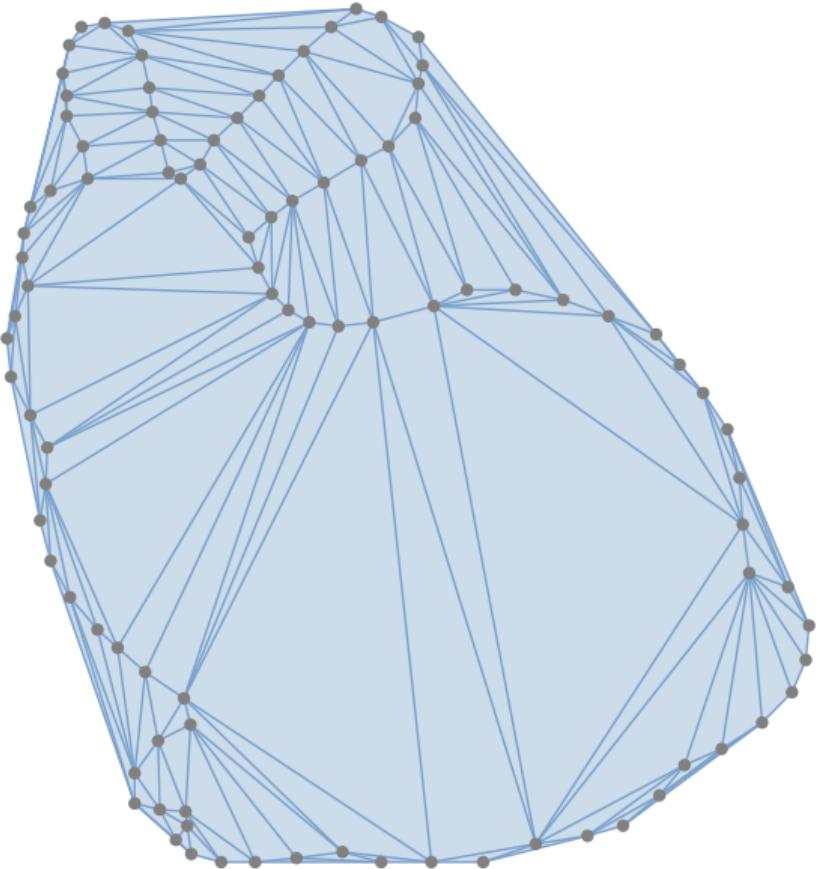
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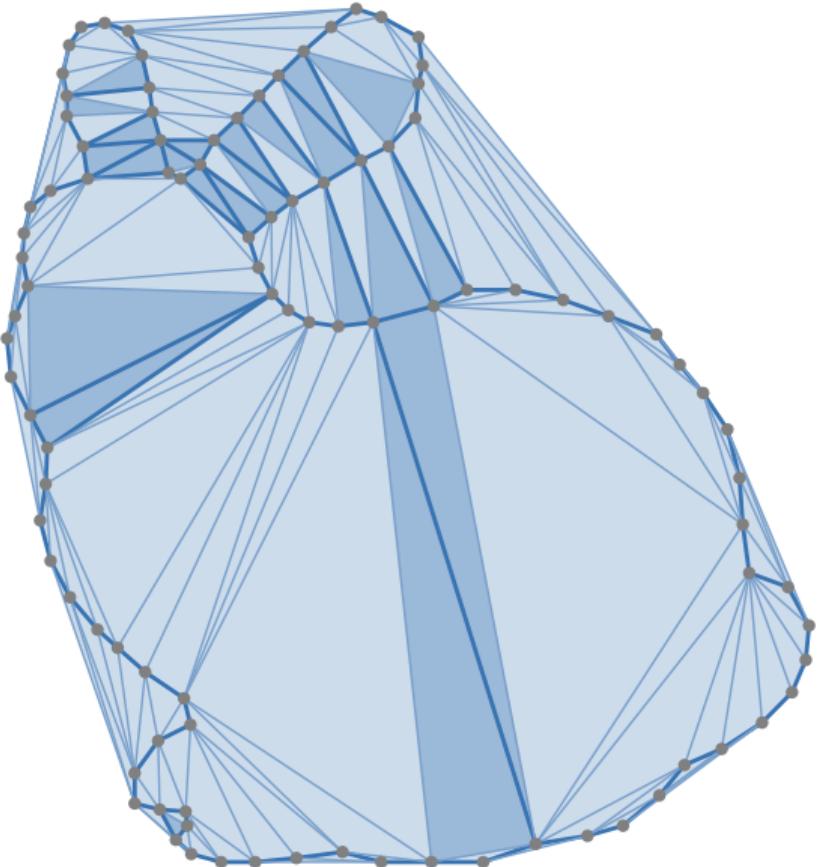
Let $X \subset \mathbb{R}$ be a finite subset in general position, and let $r \in \mathbb{R}$. Then the lexicographic optimal cycles of $\text{Del}_r(X)$ are supported on $\text{Wrap}_r(X)$.

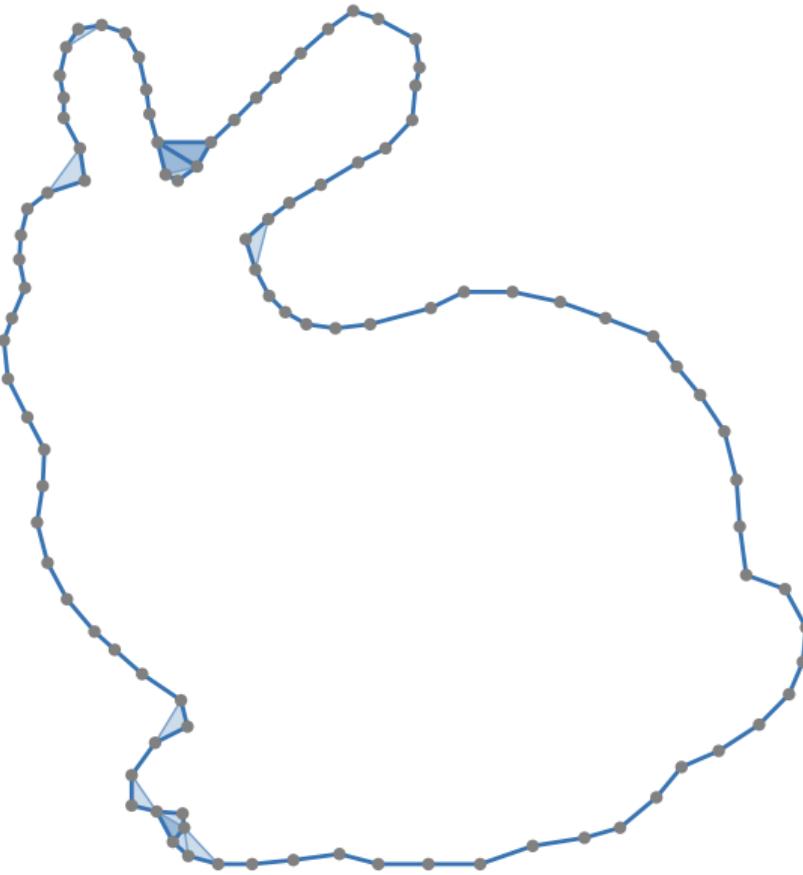


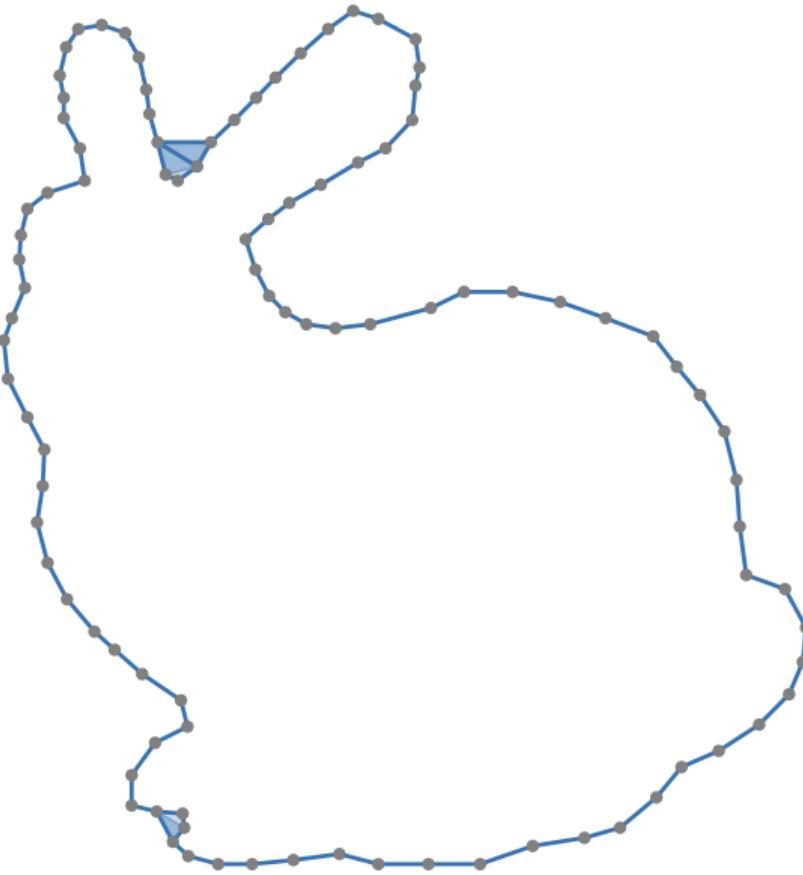
Point cloud reconstruction with lexicographic optimal cycles

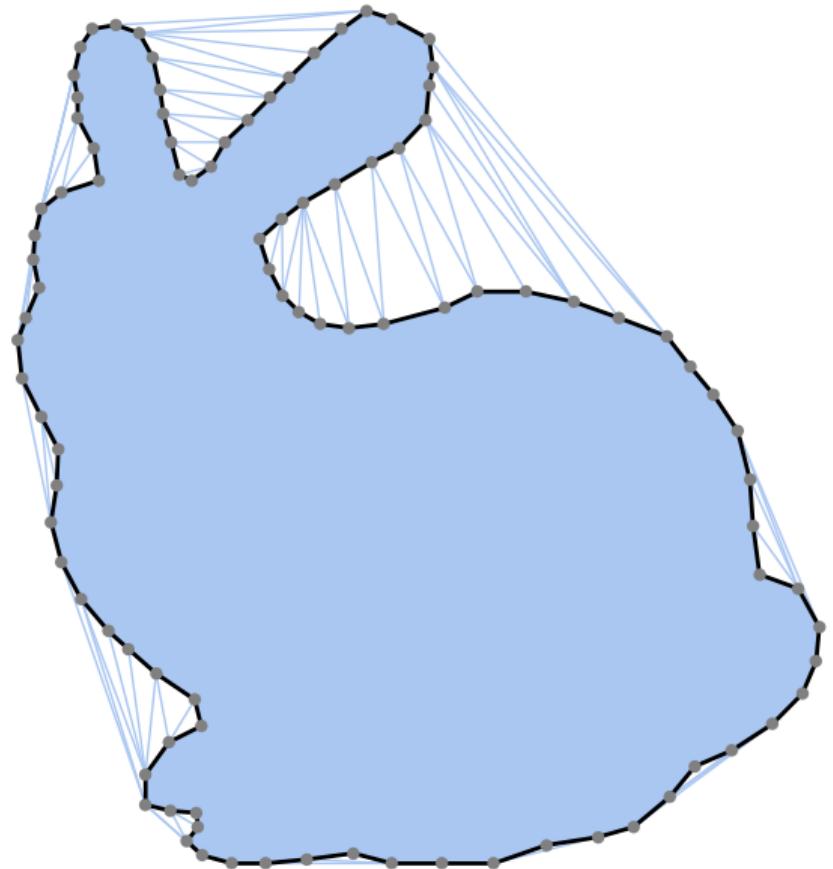


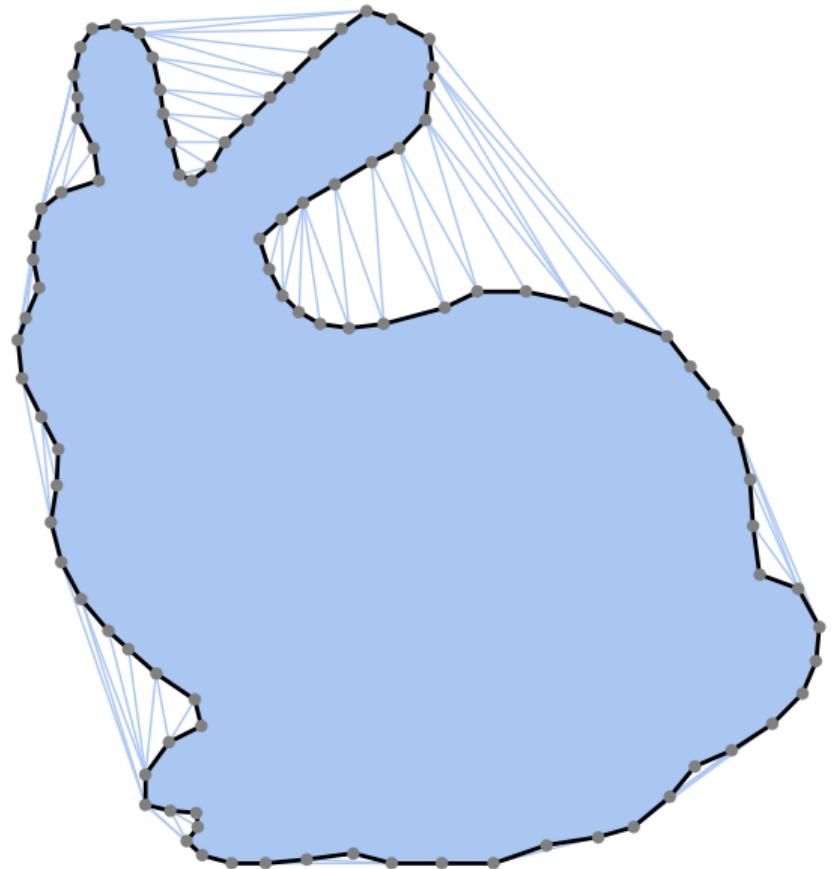


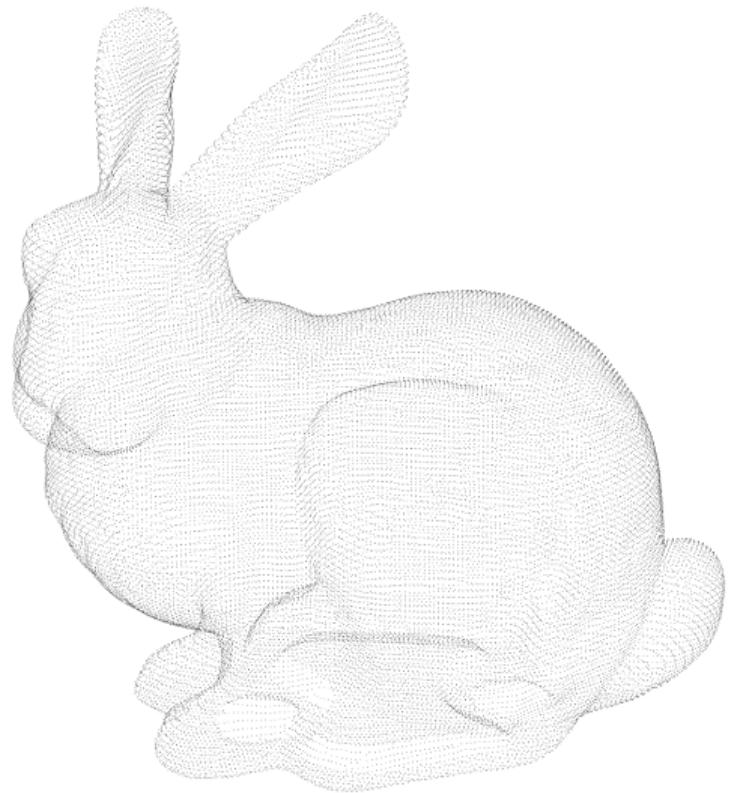


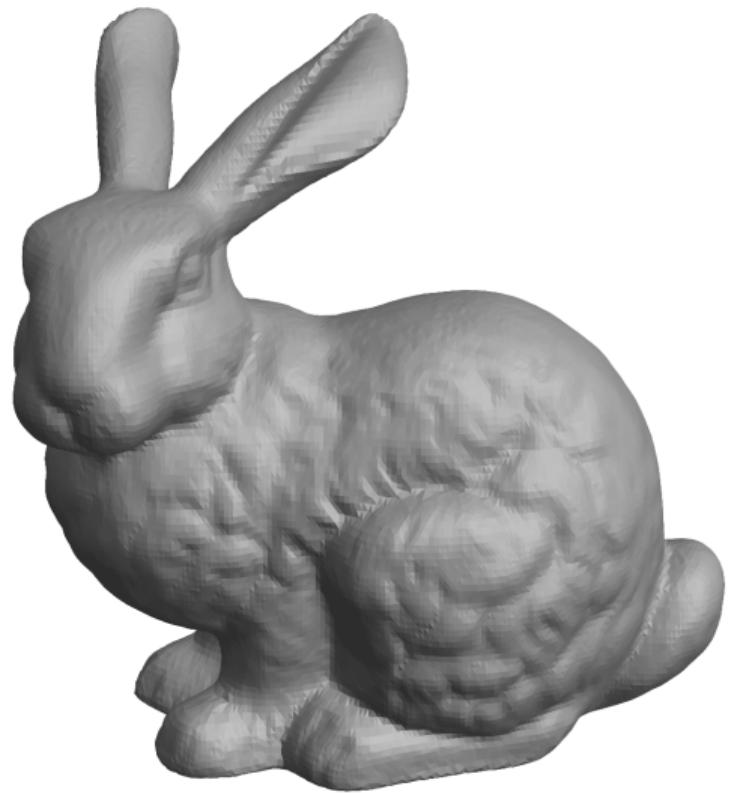












Thanks for your attention!

Further reading

-  [U. Bauer, M. Kerber, F. Roll, and A. Rolle](#)
A Unified View on the Functorial Nerve Theorem and its Variations
[Preprint, arXiv:2203.03571, 2022](#)
-  [U. Bauer, H. Edelsbrunner](#)
The Morse Theory of Čech and Delaunay Complexes
[Transactions of the AMS, 2017.](#)
-  [U. Bauer, F. Roll](#)
Gromov hyperbolicity, geodesic defect, and apparent pairs in Vietoris-Rips filtrations
[Symposium on Computational Geometry, 2022](#)
-  [U. Bauer, F. Roll](#)
Connecting Discrete Morse Theory and Persistence: Wrap Complexes and Lexicographic Optimal Cycles
[Preprint, arXiv:2212.02345, 2022](#)