

Induced matchings and the algebraic stability of persistence barcodes

Ulrich Bauer

TUM

July 12, 2015

Geometry Workshop Seggau

Joint work with Michael Lesnick (IMA)



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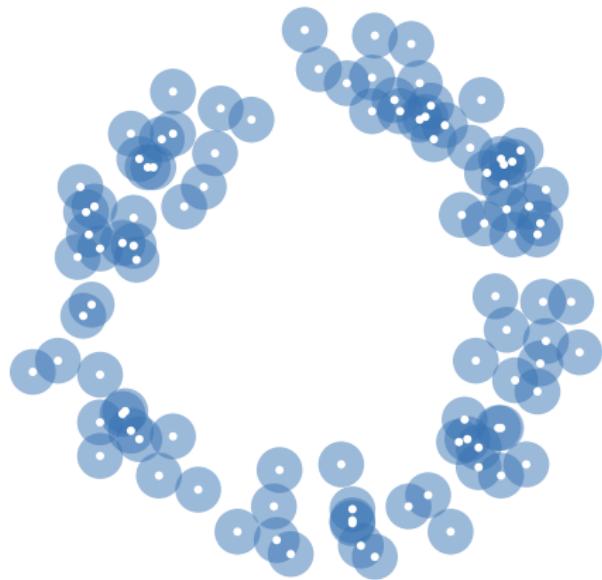


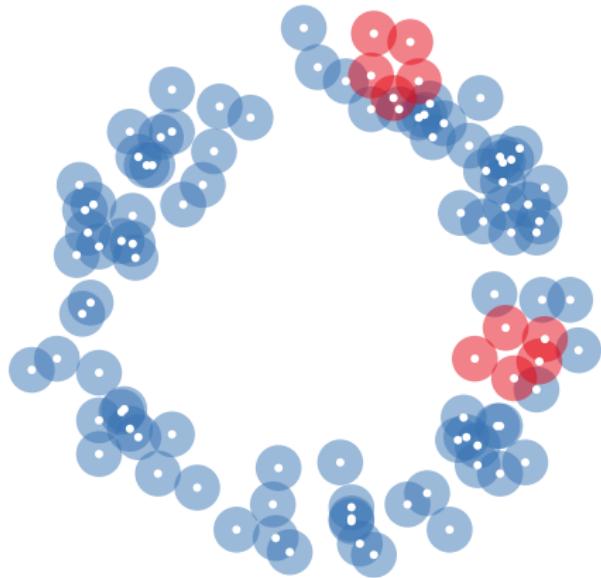
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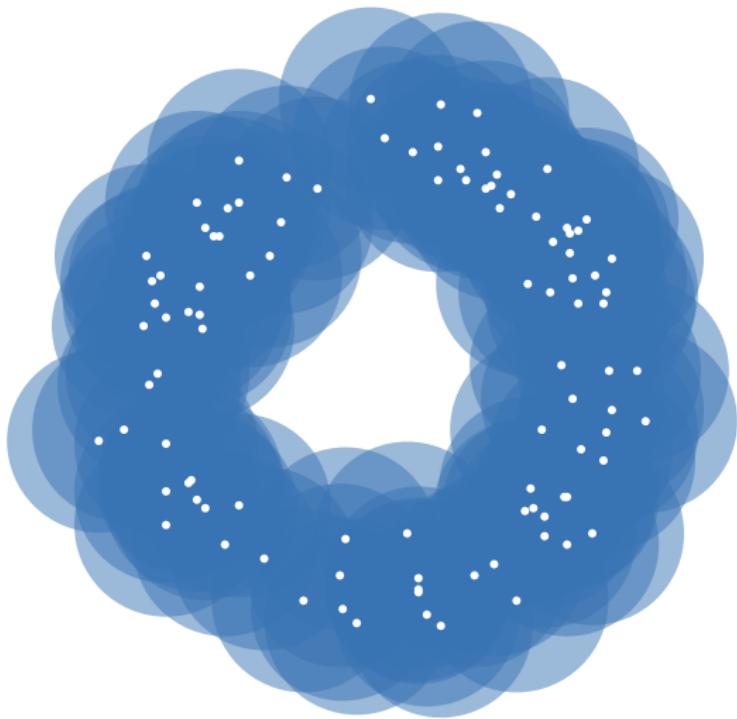
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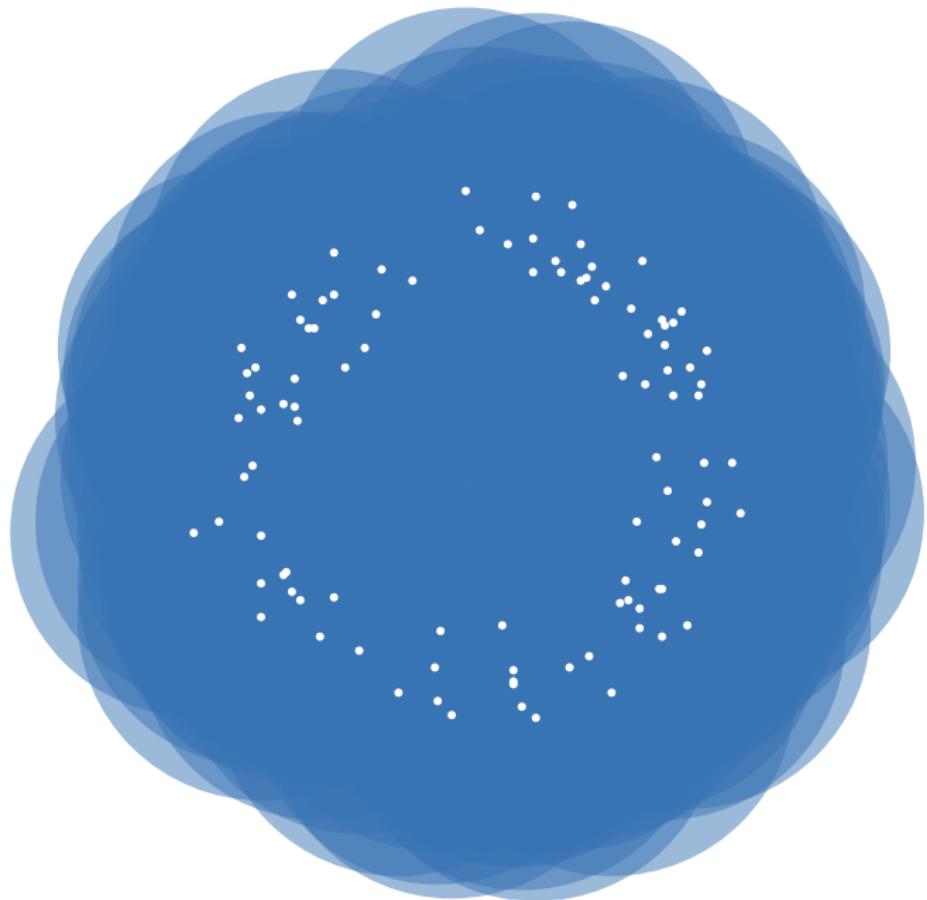


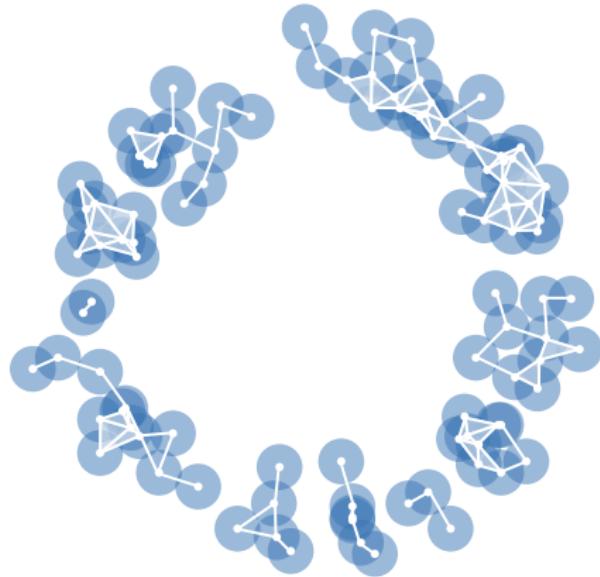


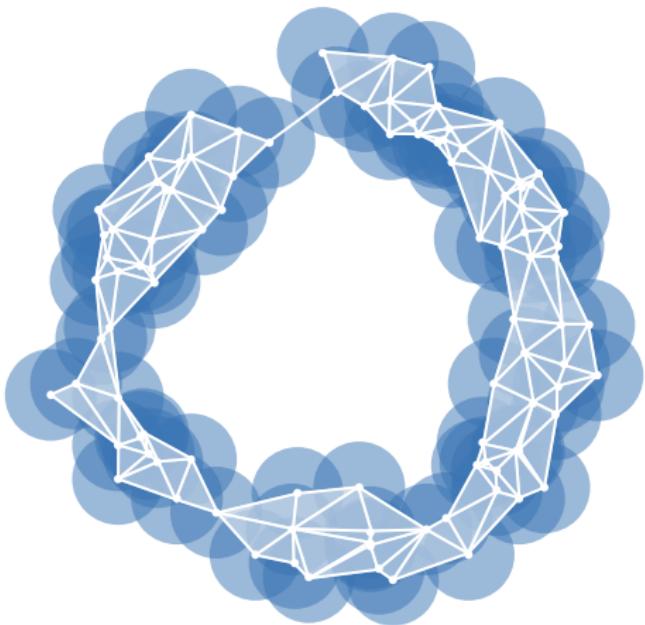


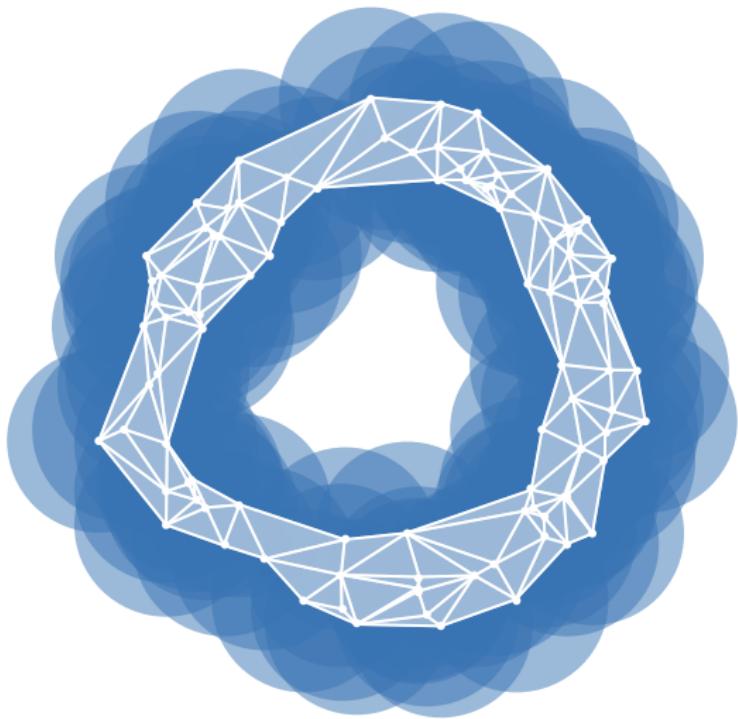


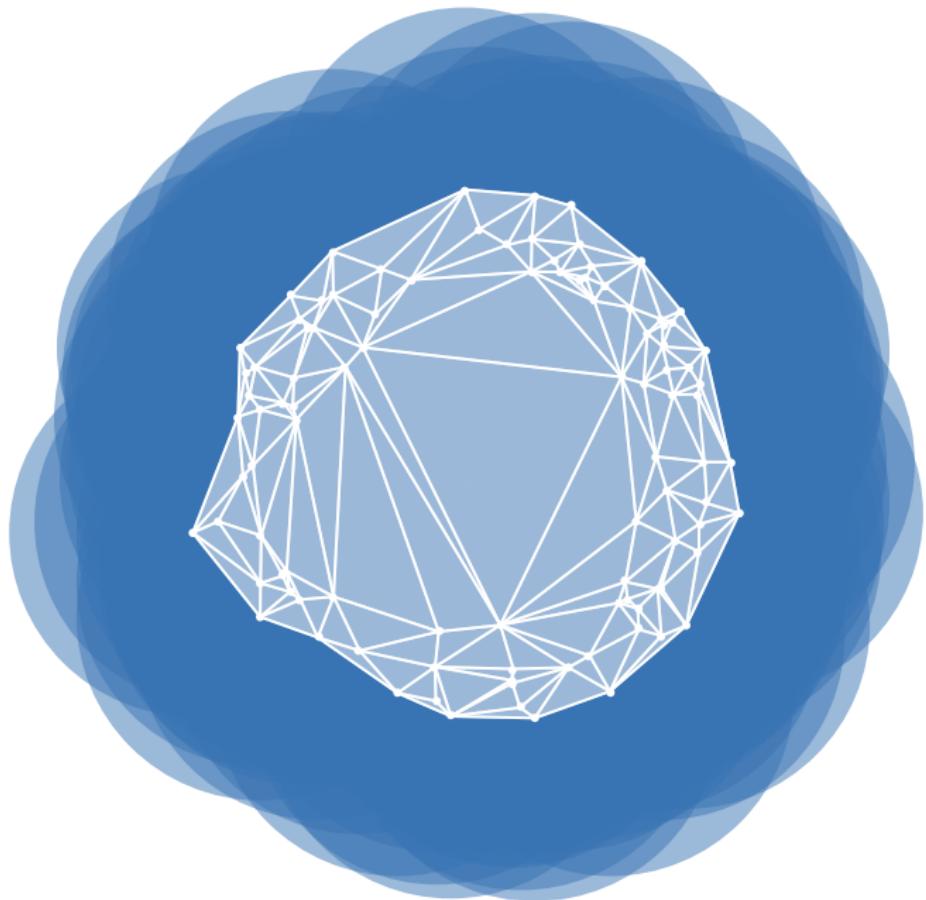




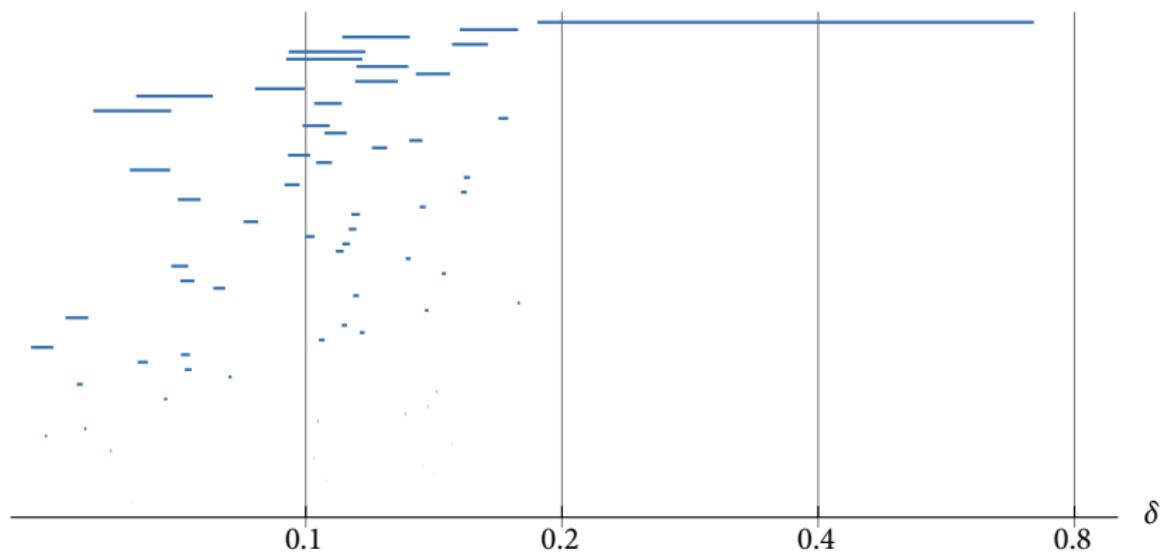
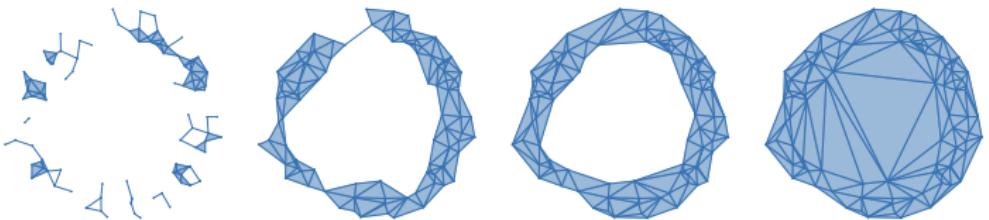




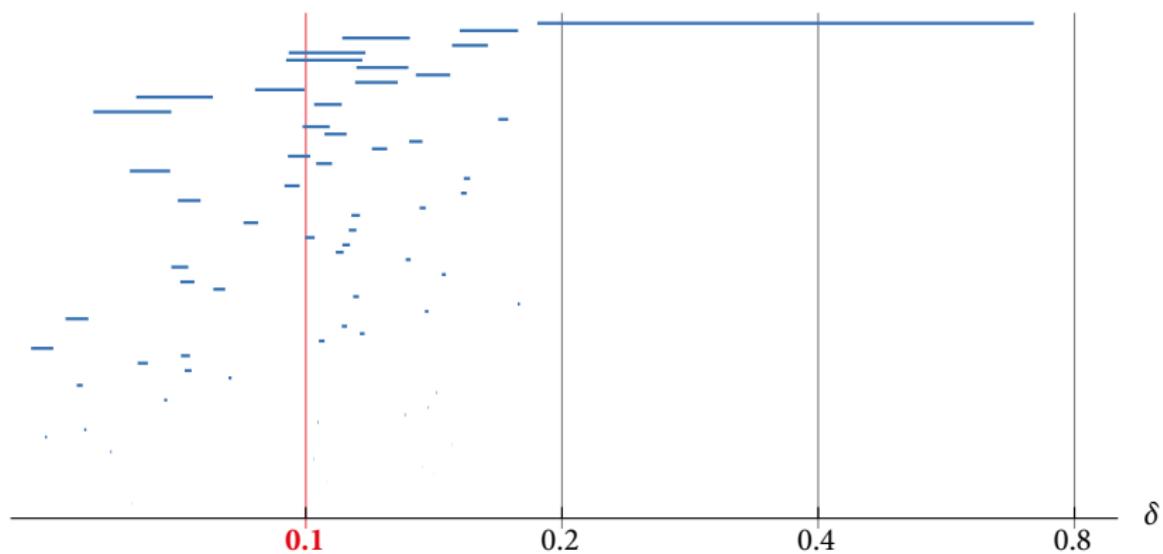
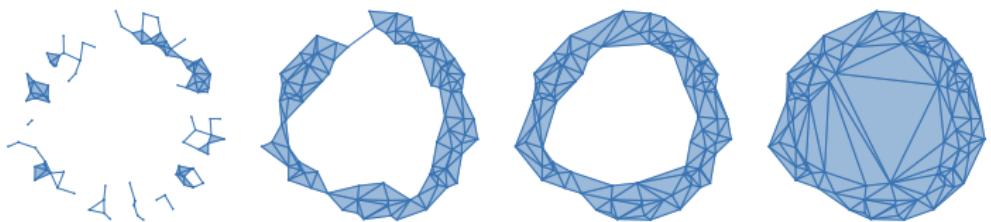




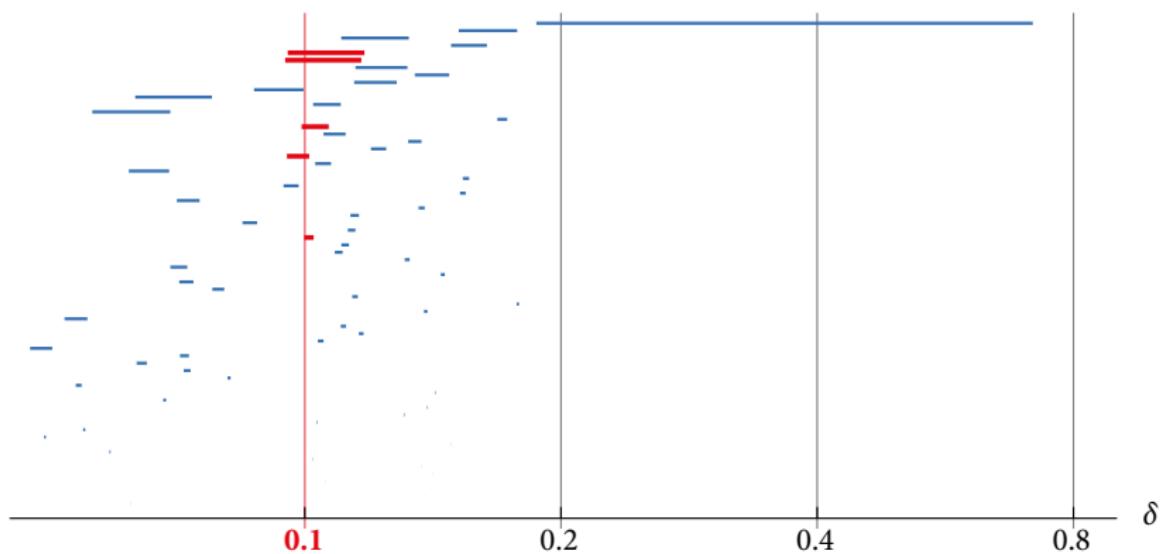
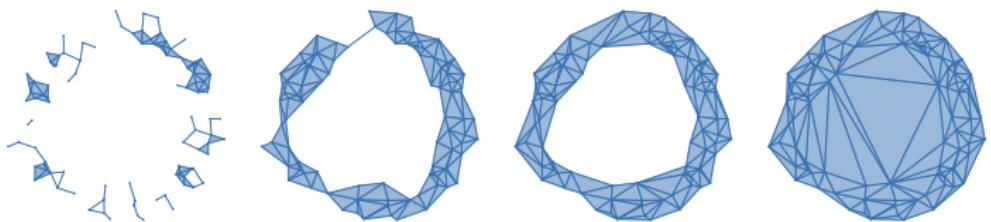
What is persistent homology?



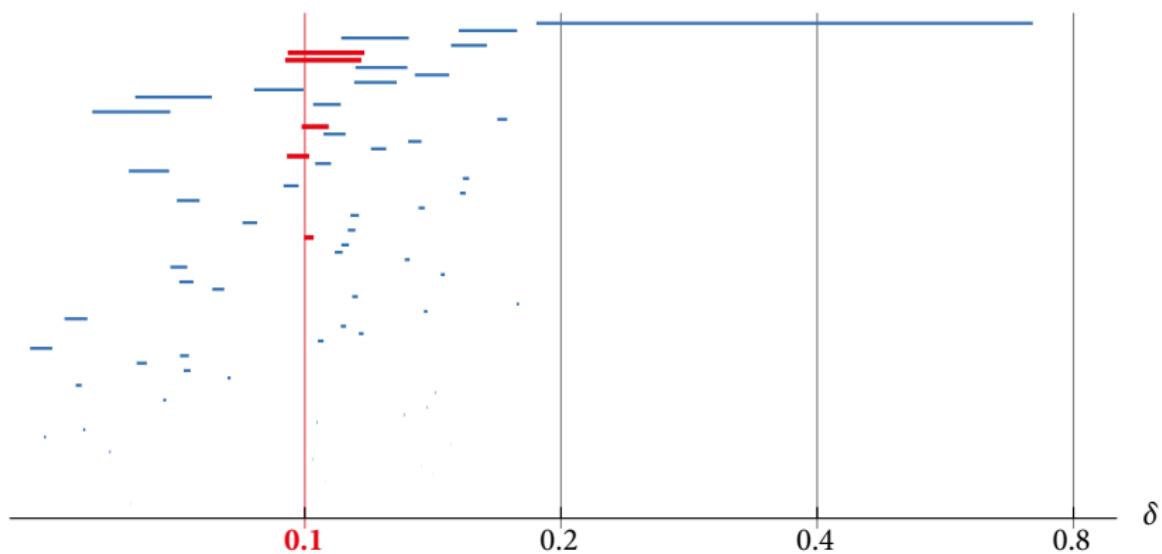
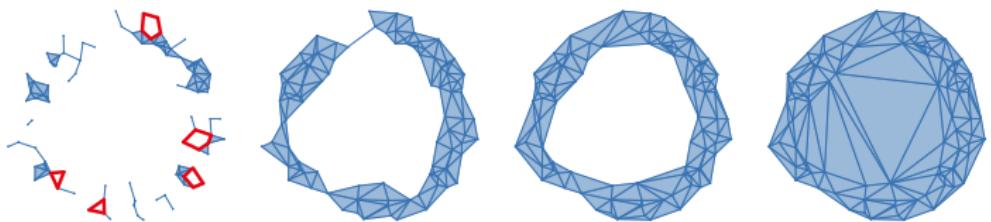
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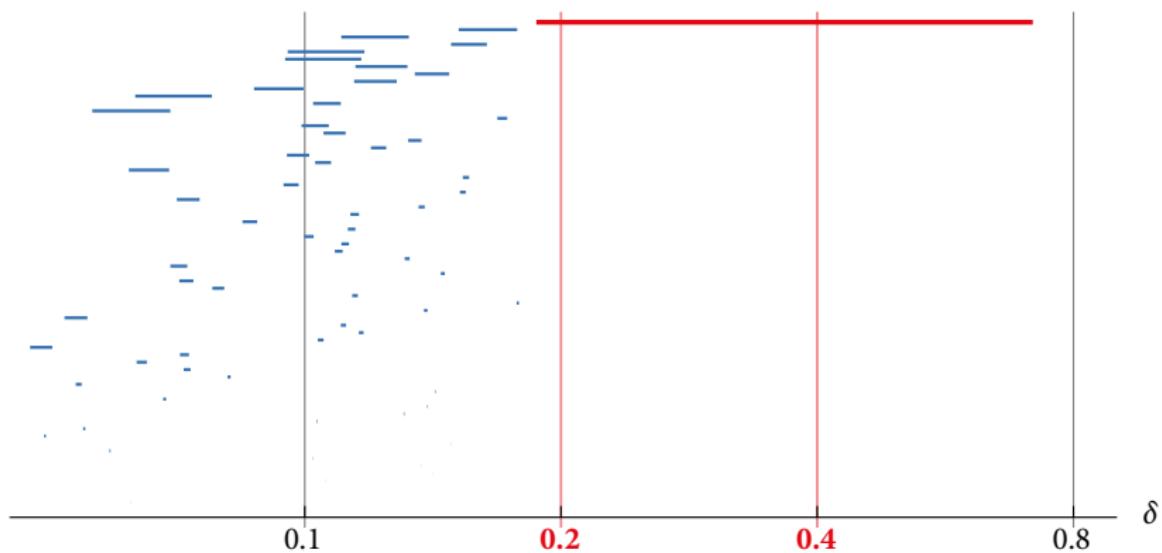
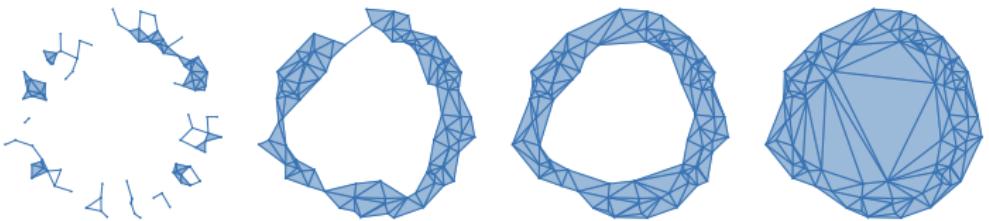
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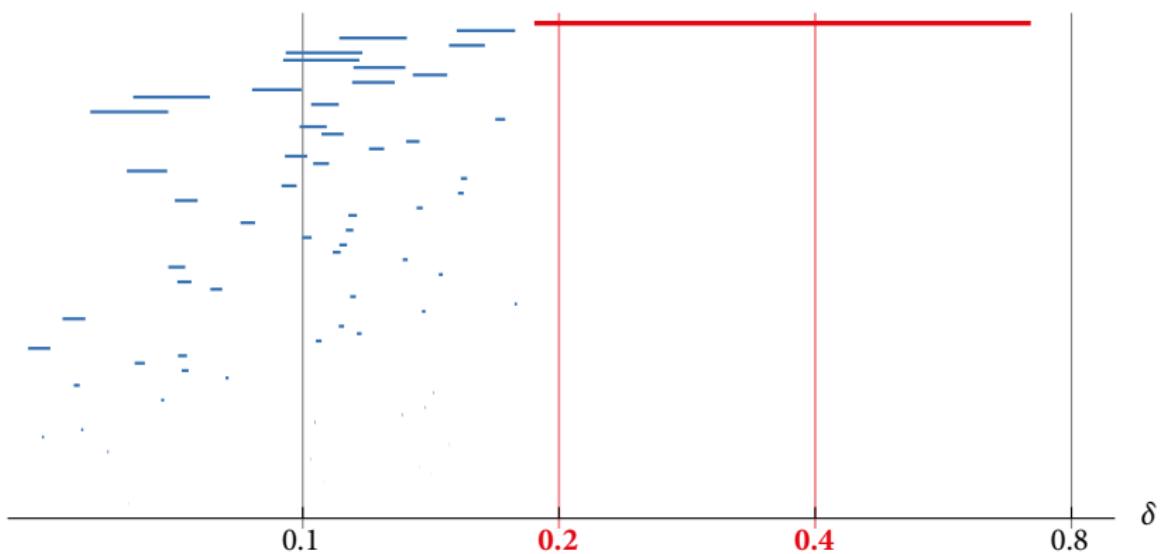
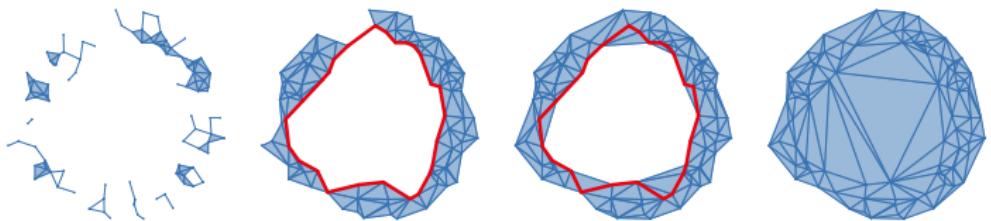
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 - A topological space K_t for each $t \in \mathbb{R}$

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- A filtration is a certain diagram $K : \mathbf{R} \rightarrow \mathbf{Top}$.
 - A topological space K_t for each $t \in \mathbb{R}$
 - An inclusion map $K_s \hookrightarrow K_t$ for each $s \leq t \in \mathbb{R}$

Inference

Homology inference

Problem (Homology inference)

Determine the homology $H_(\Omega)$ of a shape $\Omega \subset \mathbb{R}^d$ from a finite approximation P close to Ω .*

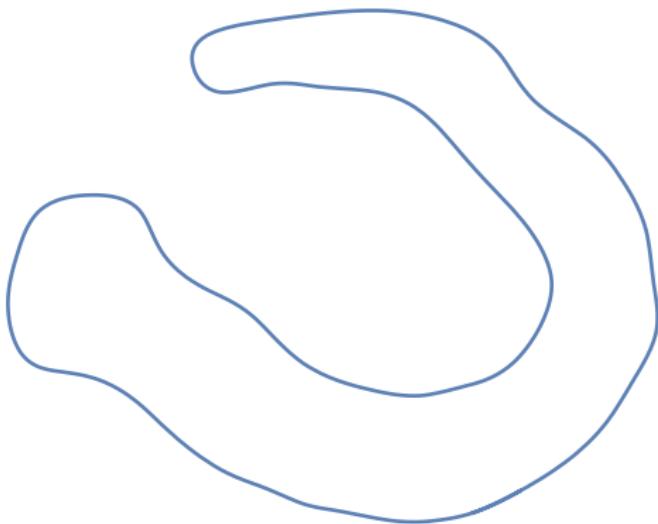
Homology inference

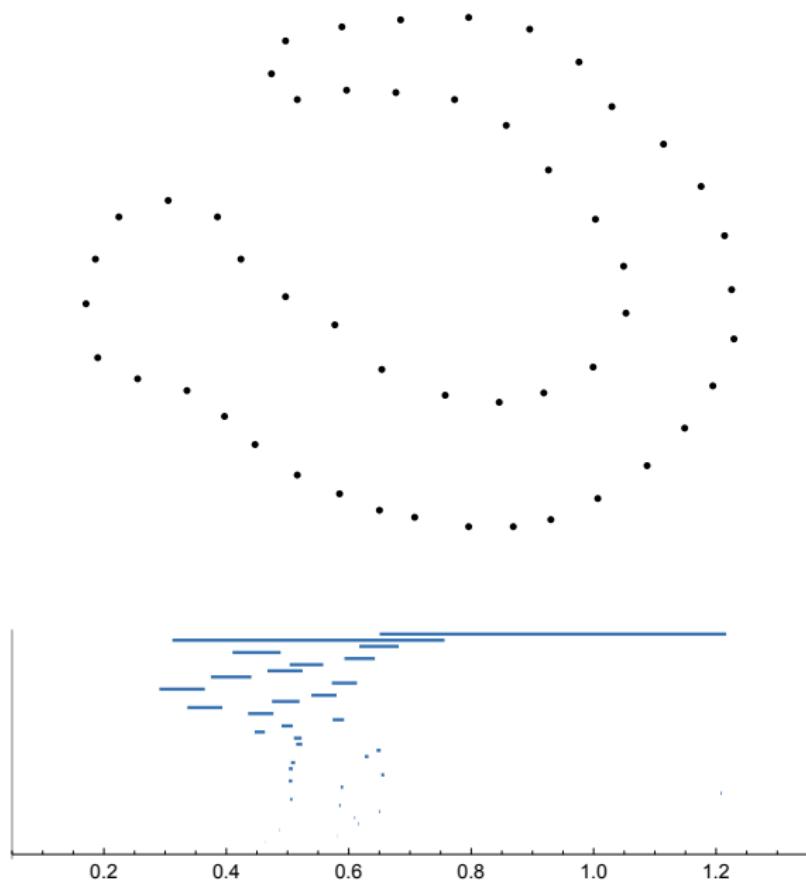
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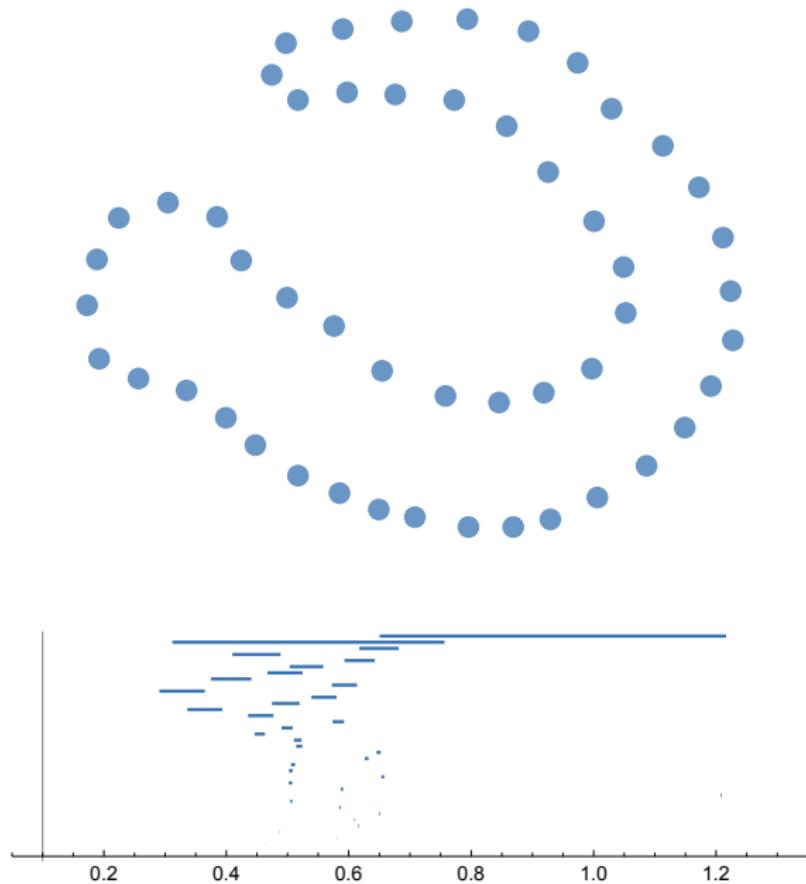
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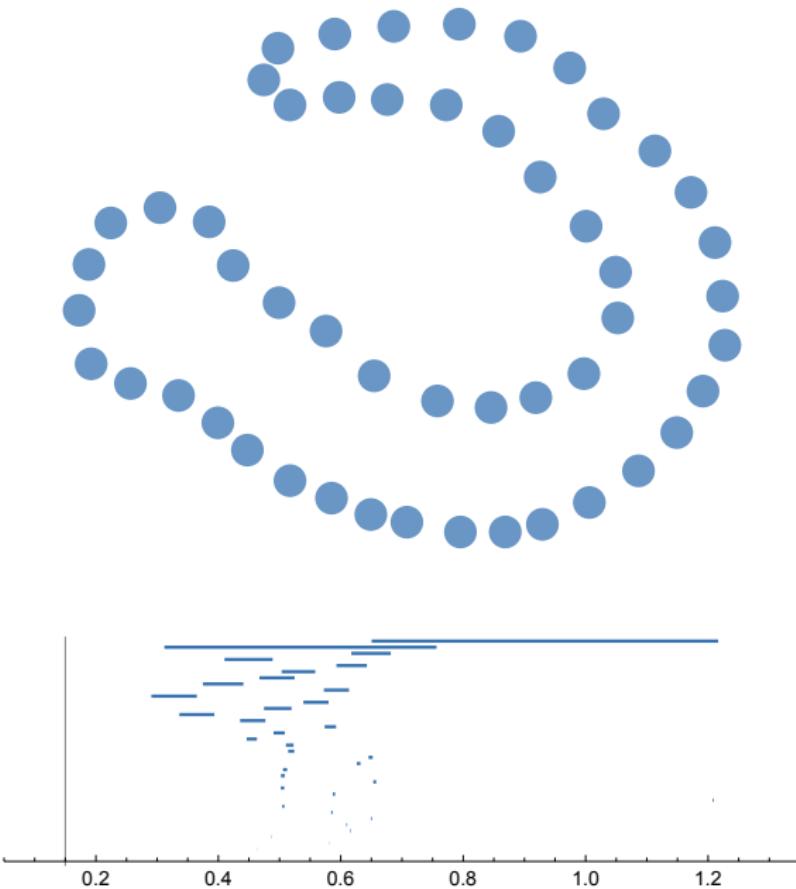
Idea:

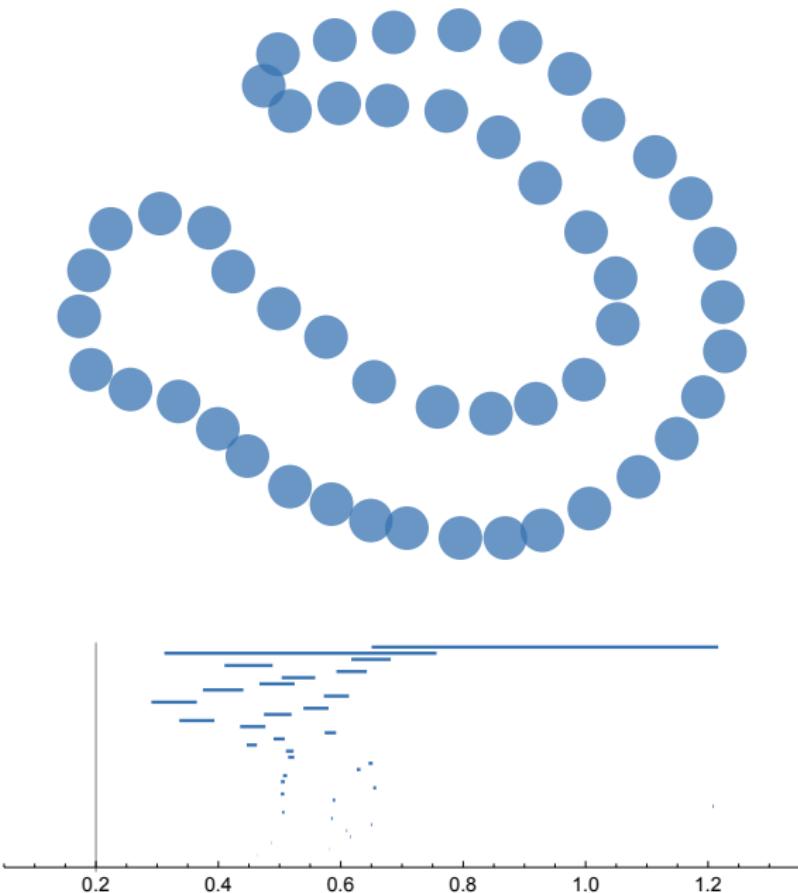
- approximate the shape by a thickening $B_\delta(P)$ covering Ω

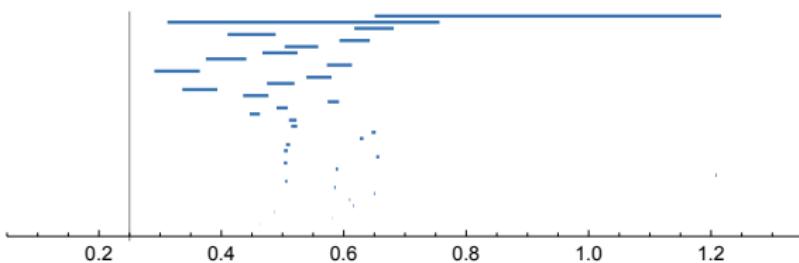
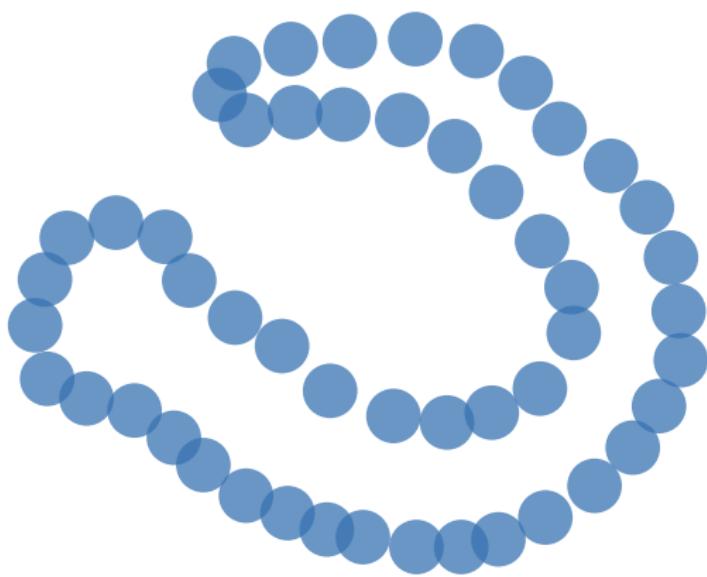


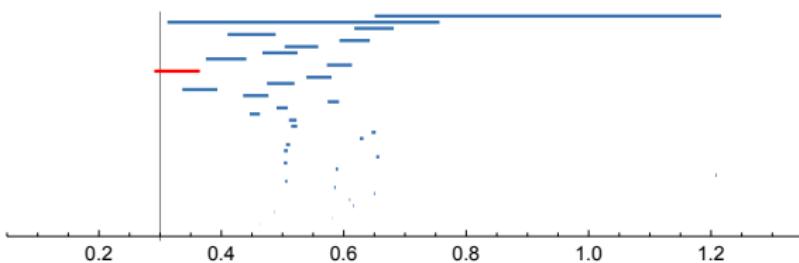
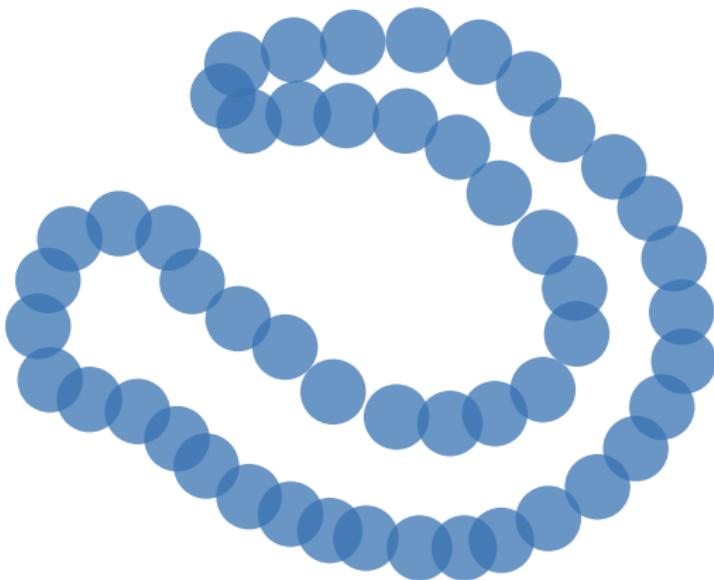


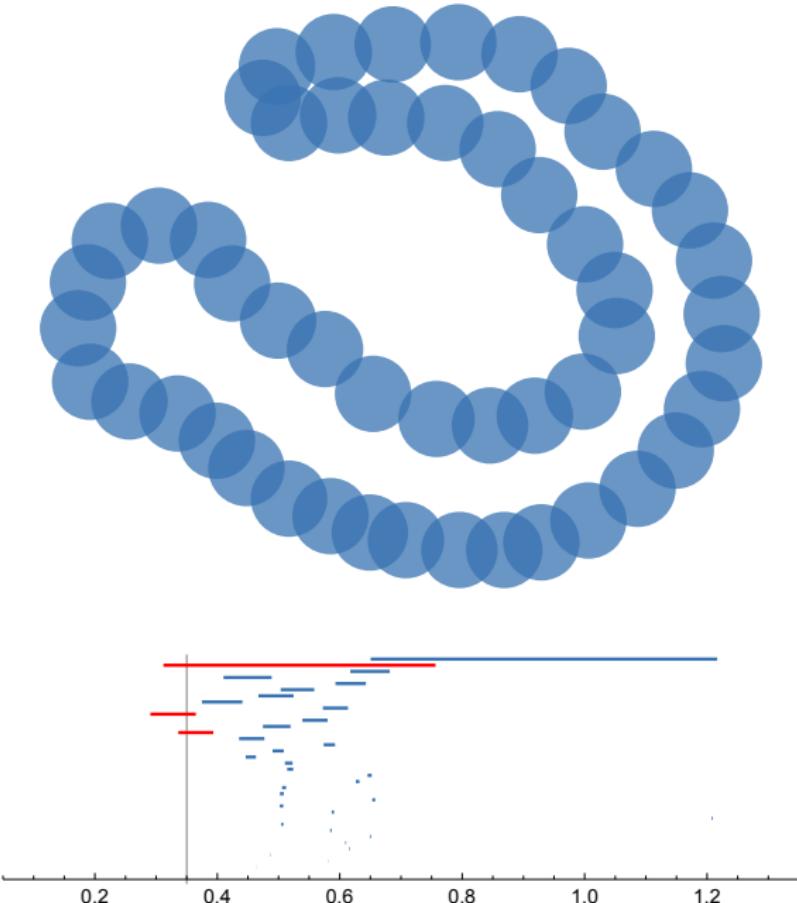


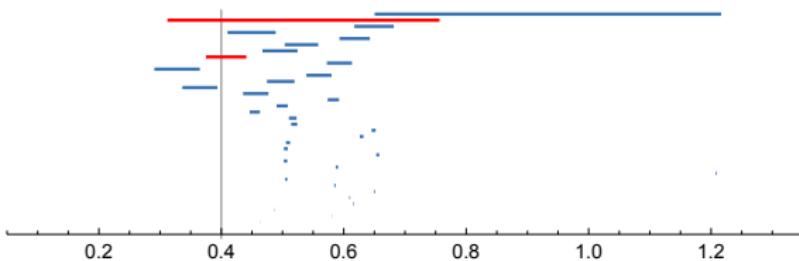
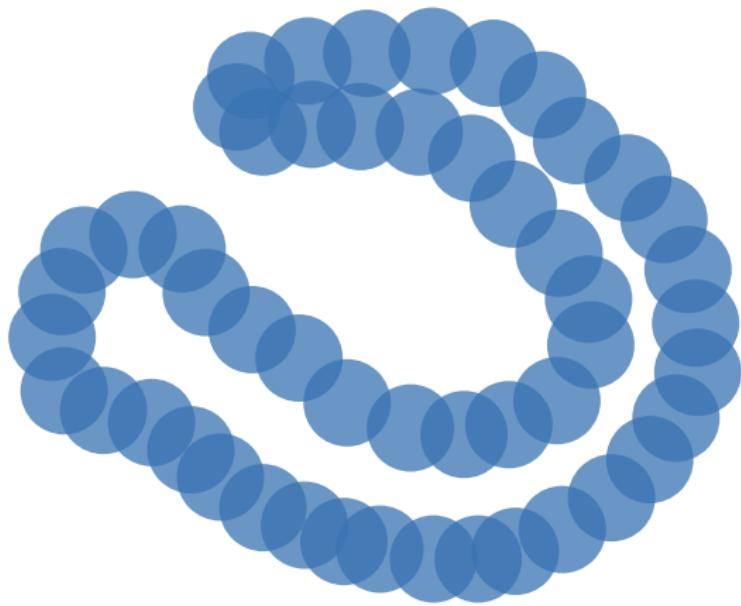


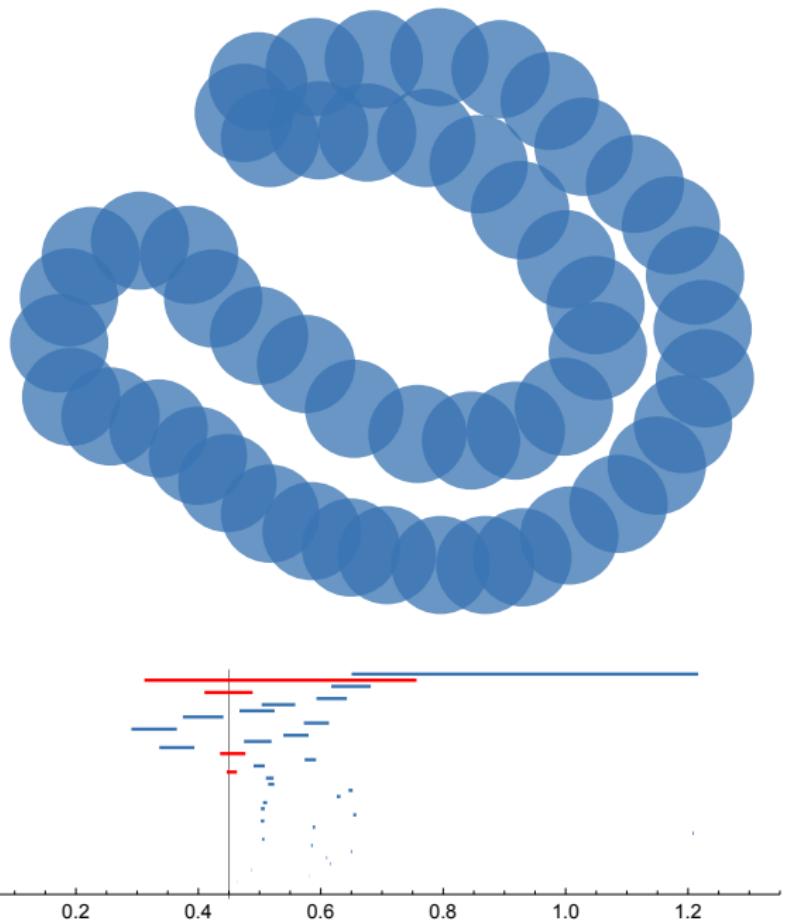


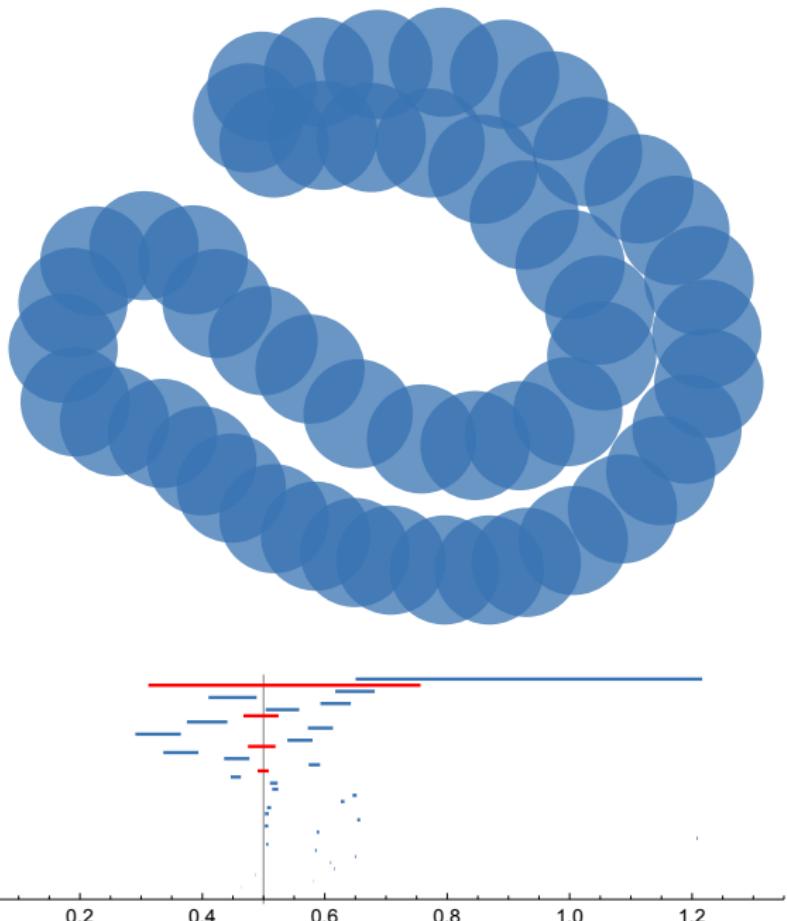


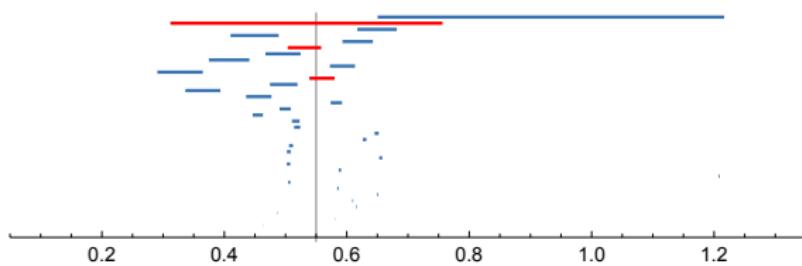
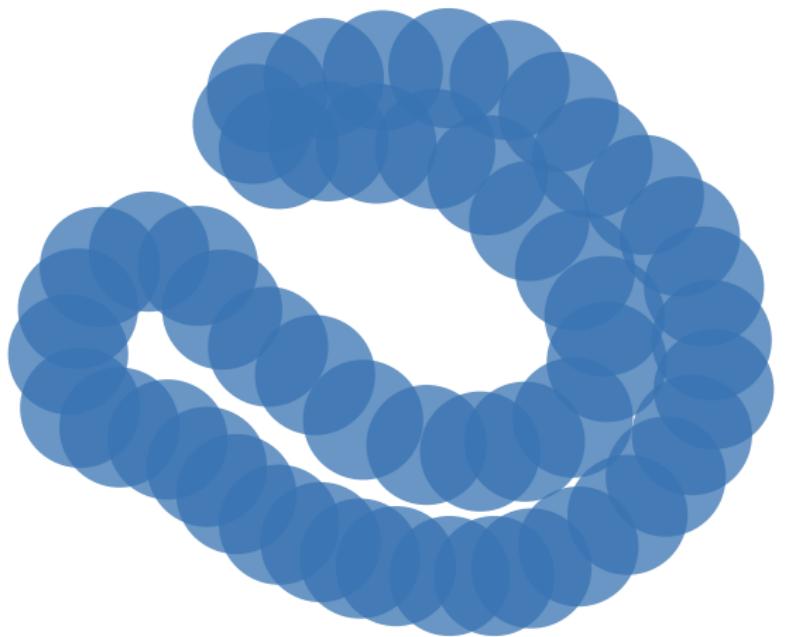


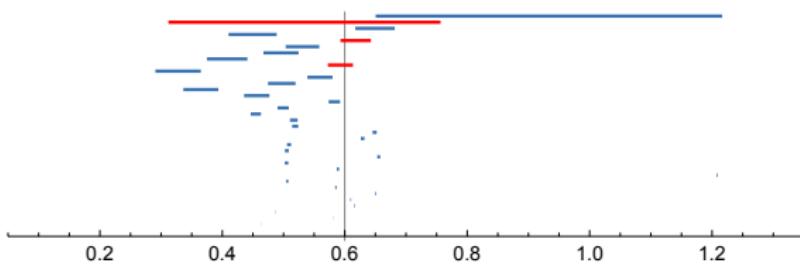
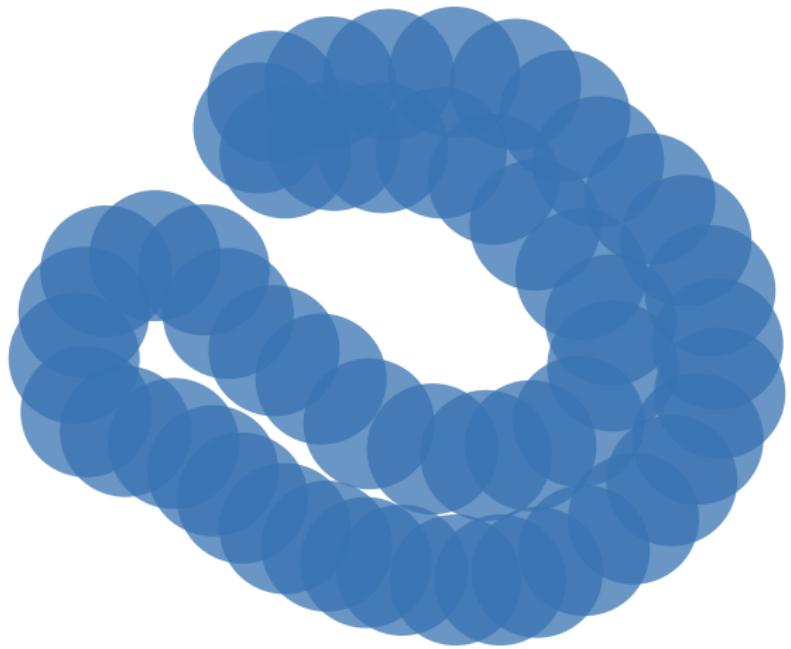


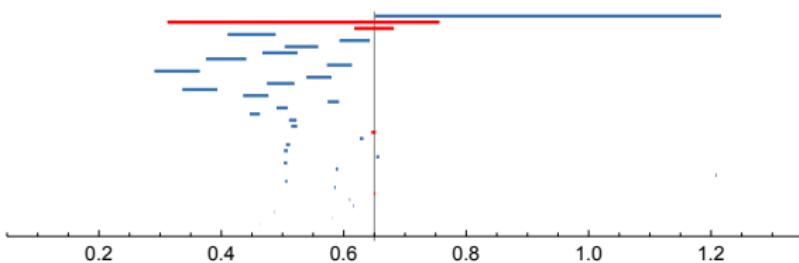
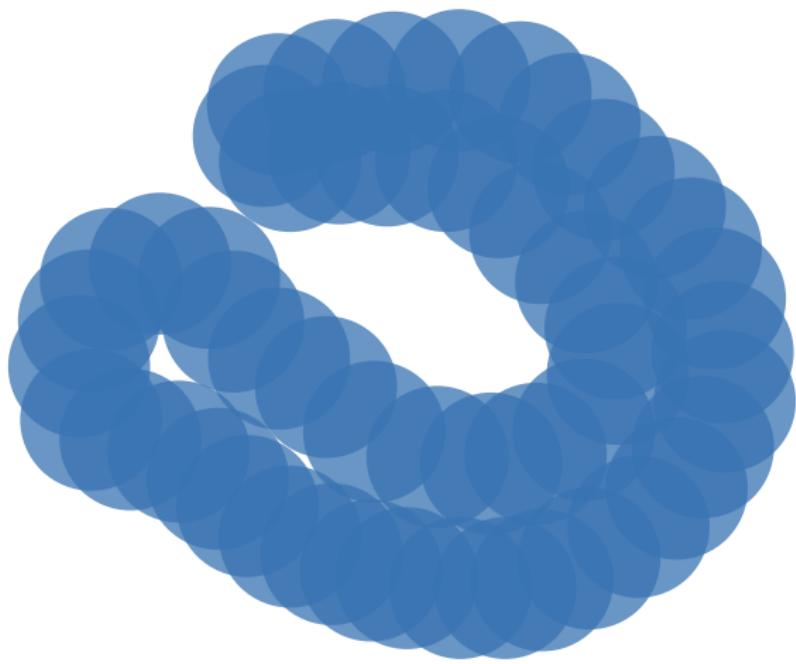


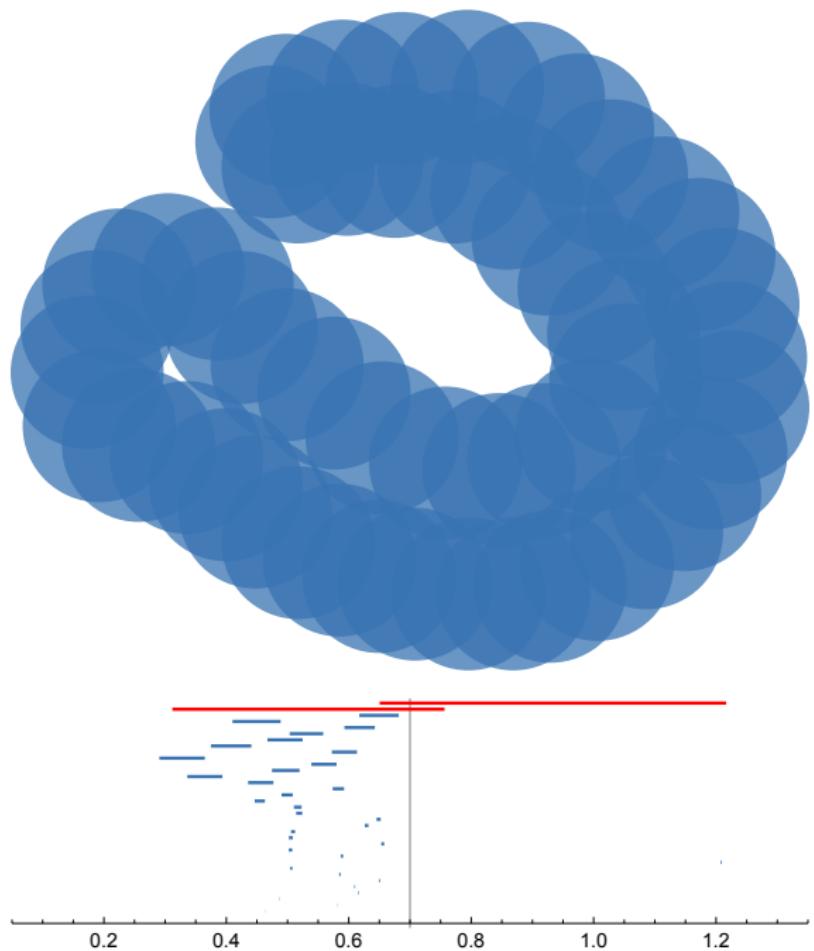


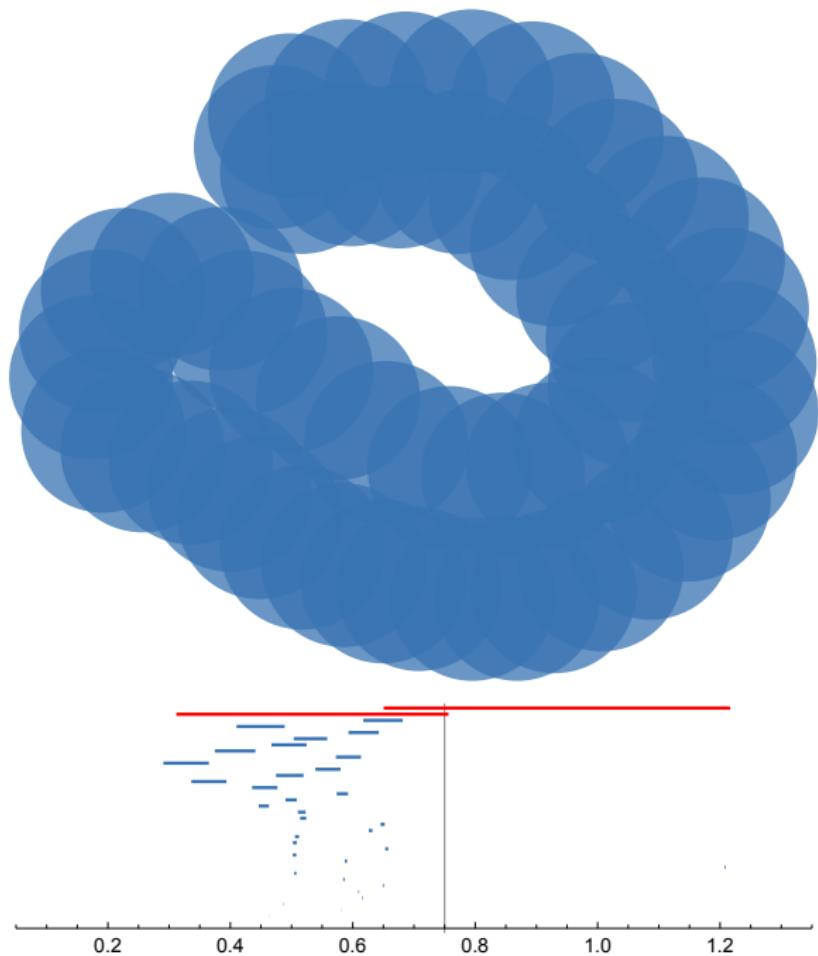


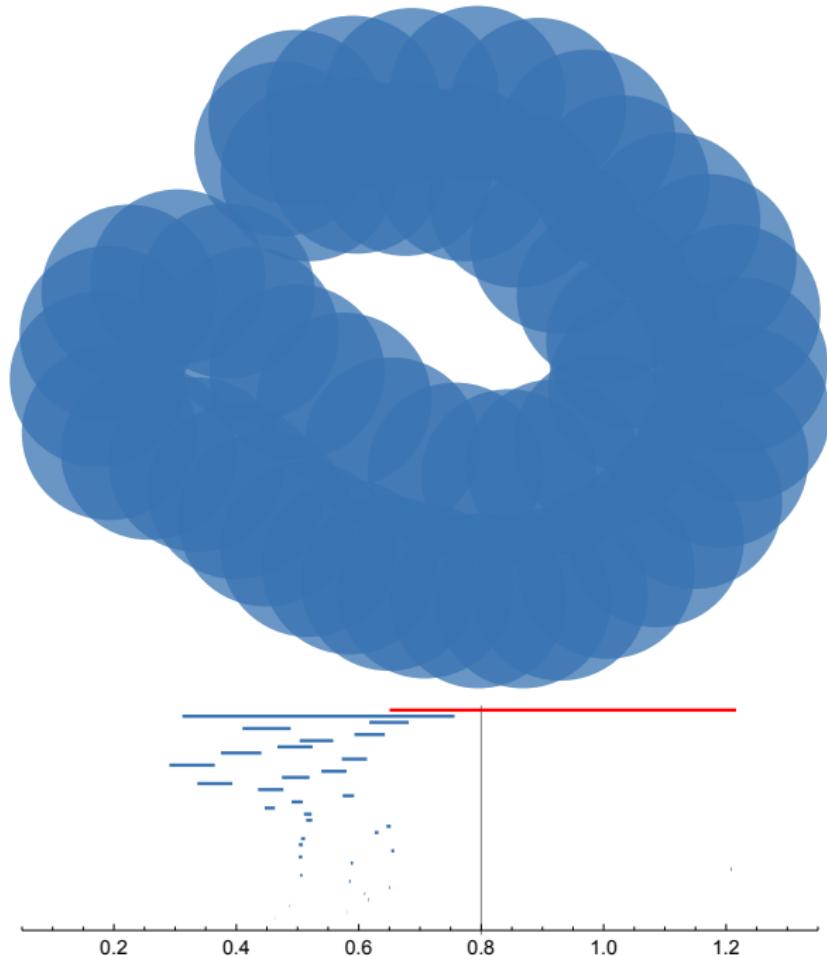


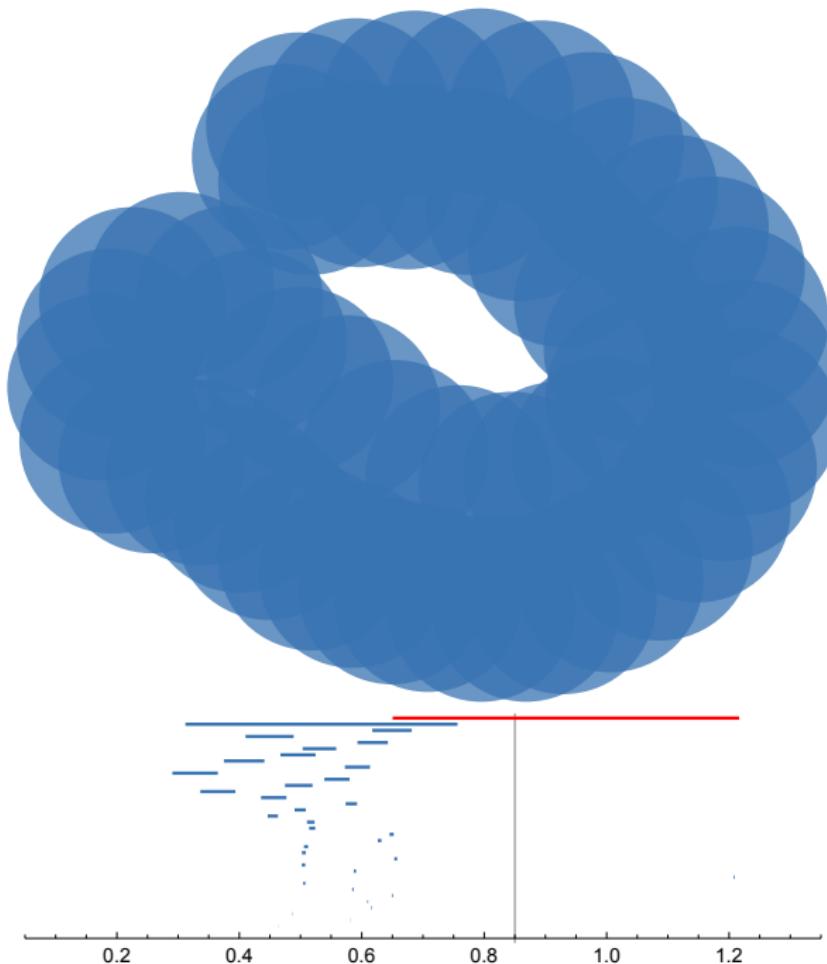


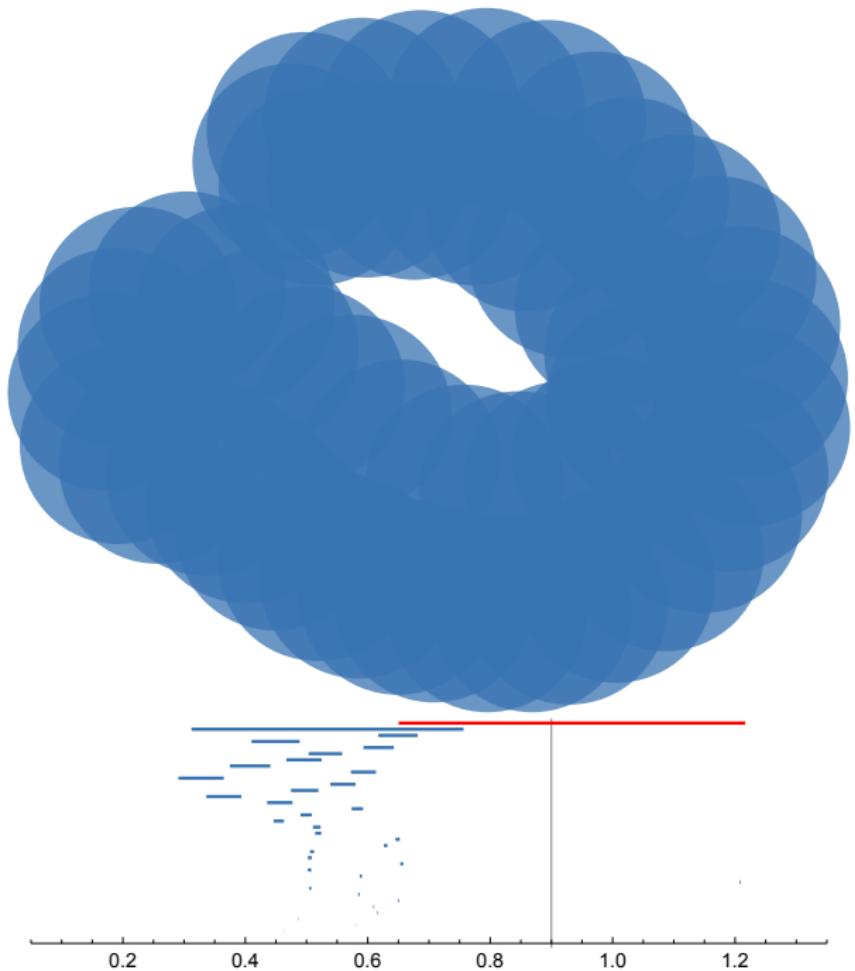












Homology inference using persistent homology

$P_\delta = B_\delta(P)$: δ -neighborhood (union of balls) around P

Theorem (Cohen-Steiner, Edelsbrunner, Harer 2005)

Let $\Omega \subset \mathbb{R}^d$. Let $P \subset \Omega$ be such that

- $\Omega \subseteq P_\delta$ for some $\delta > 0$,
- both $H_*(\Omega \hookrightarrow \Omega_\delta)$ and $H_*(\Omega_\delta \hookrightarrow \Omega_{2\delta})$ are isomorphisms.

Then

$$H_*(\Omega) \cong \text{im } H_*(P_\delta \hookrightarrow P_{2\delta}).$$

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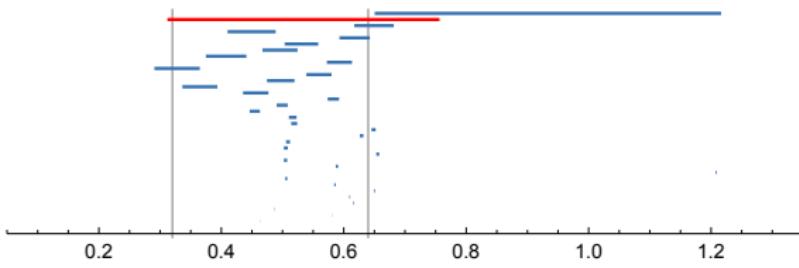
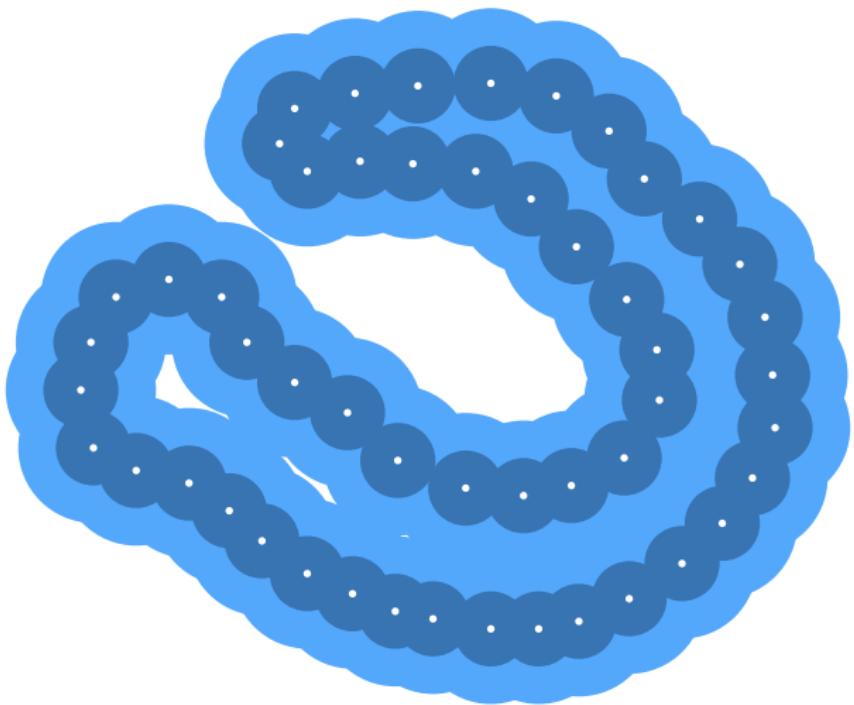
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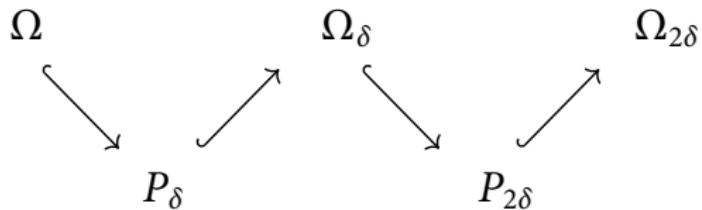
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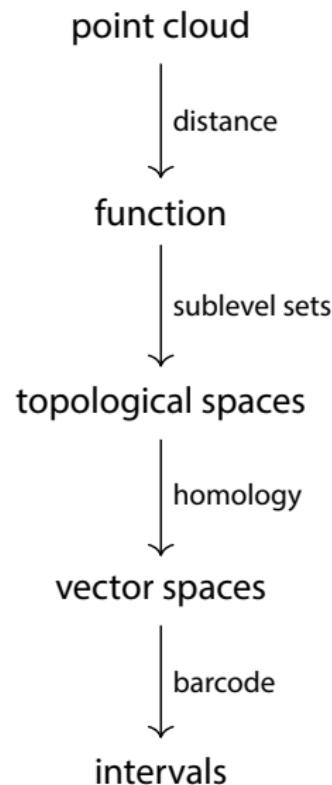
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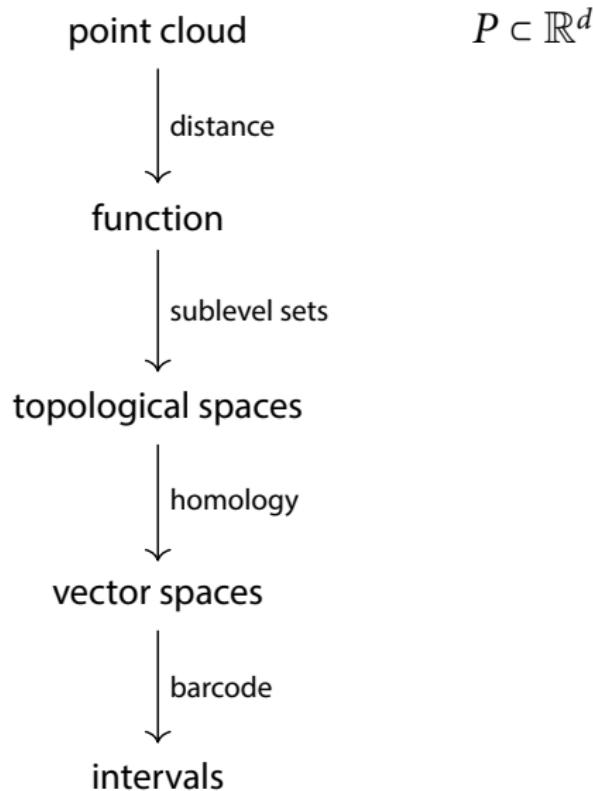


Stability

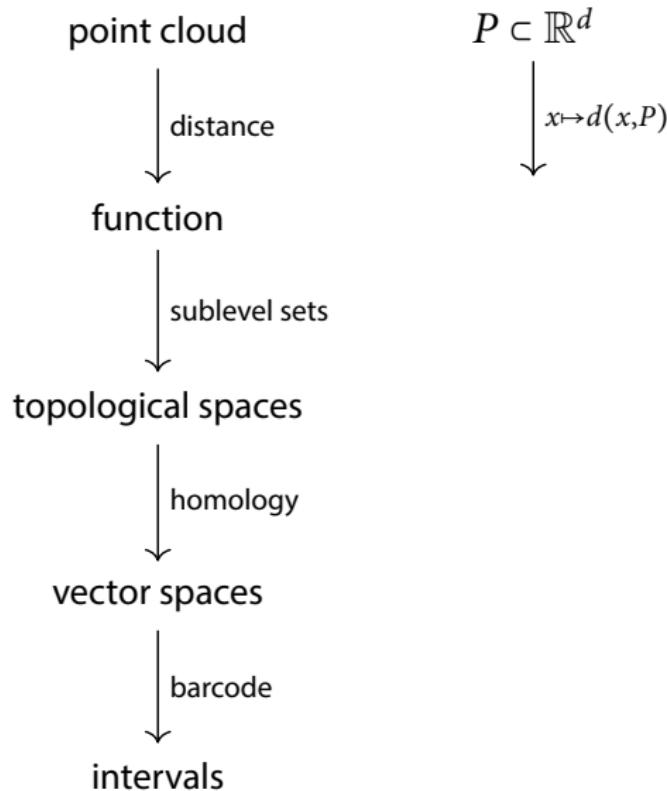
The pipeline of topological data analysis



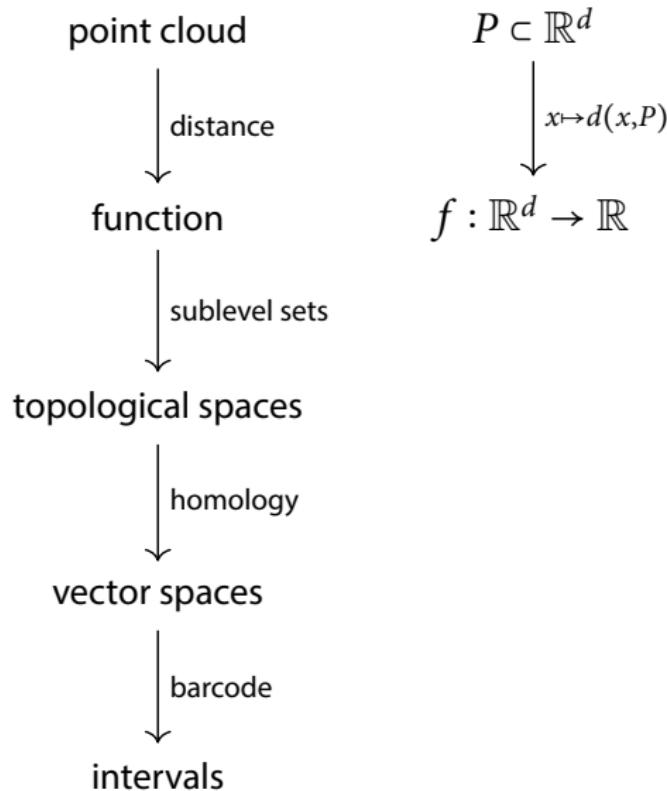
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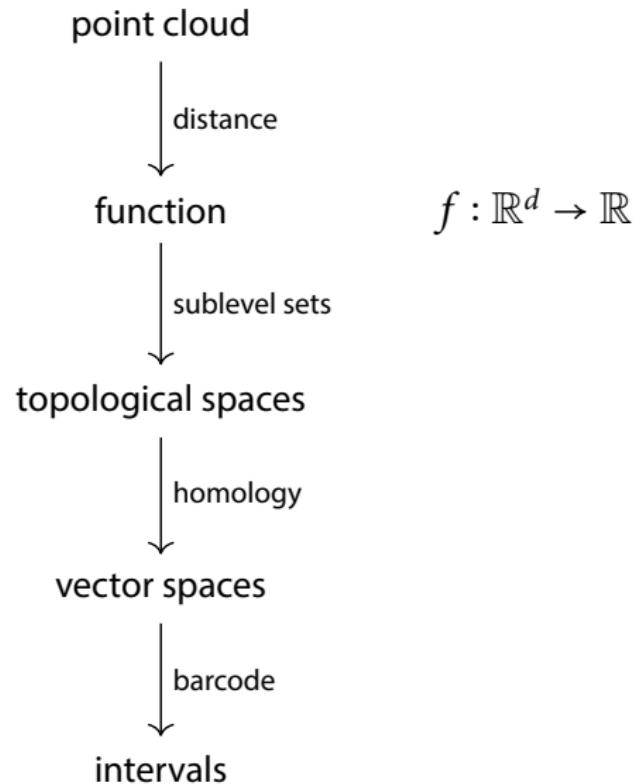
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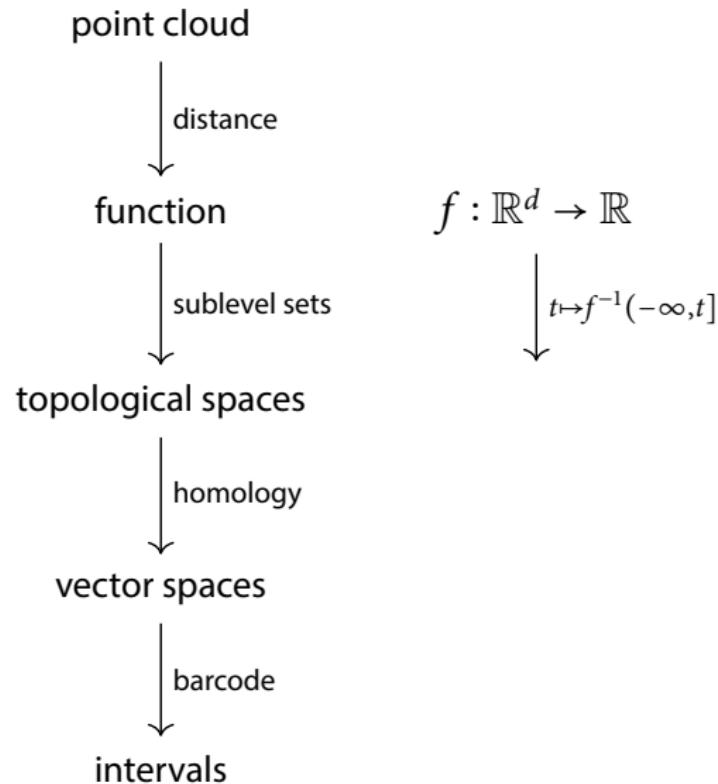
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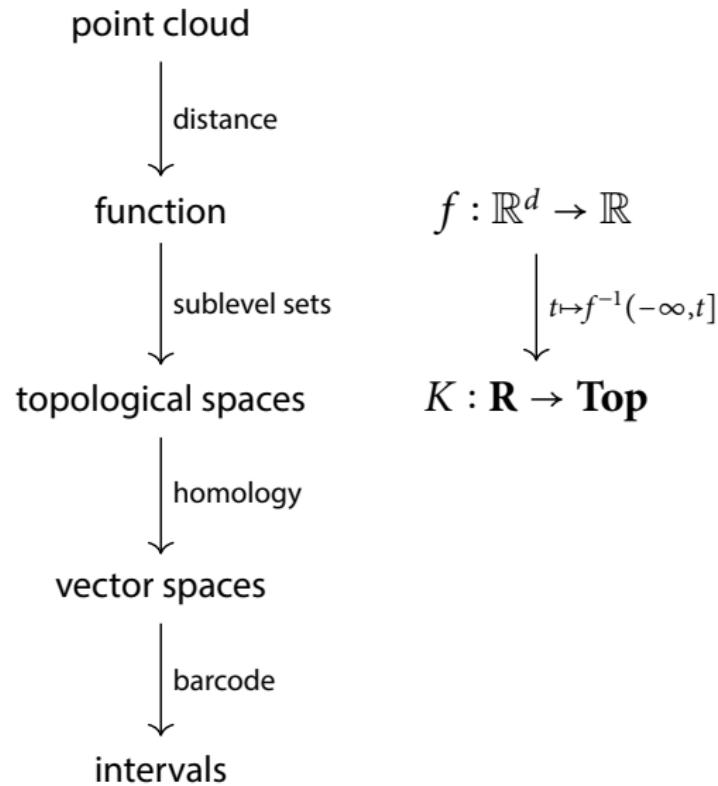
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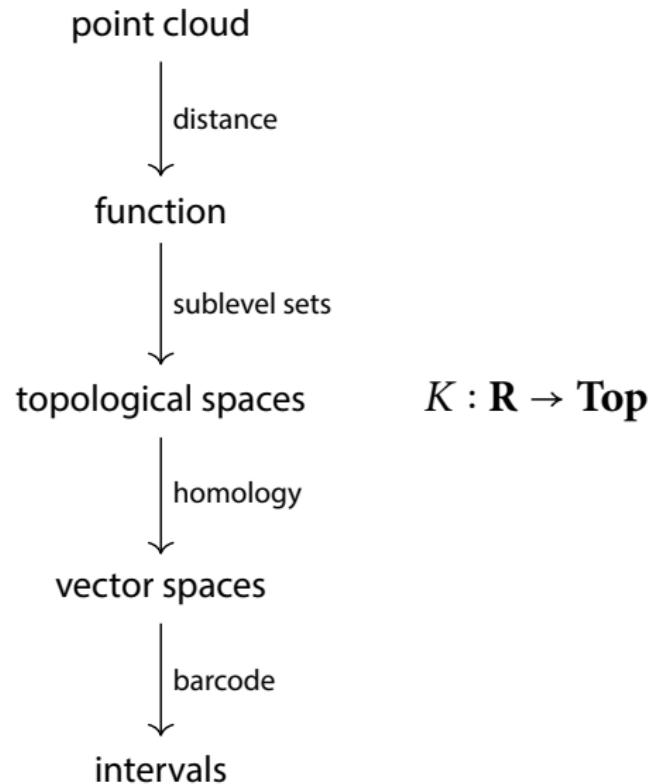
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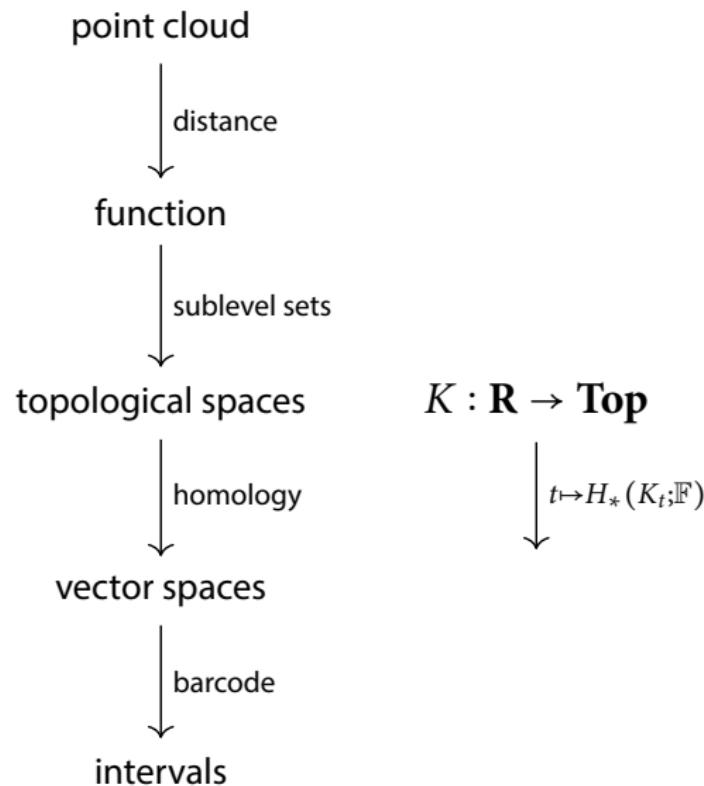
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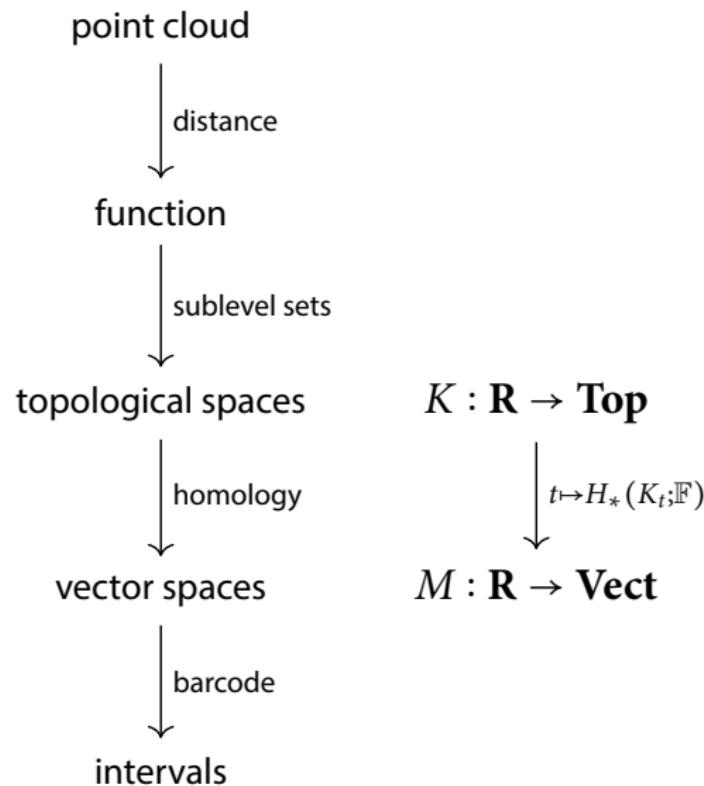
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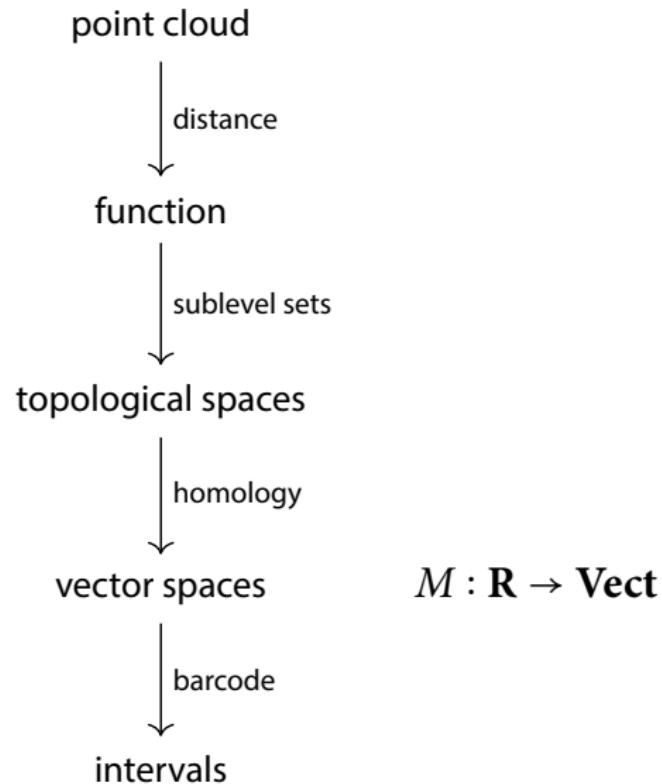
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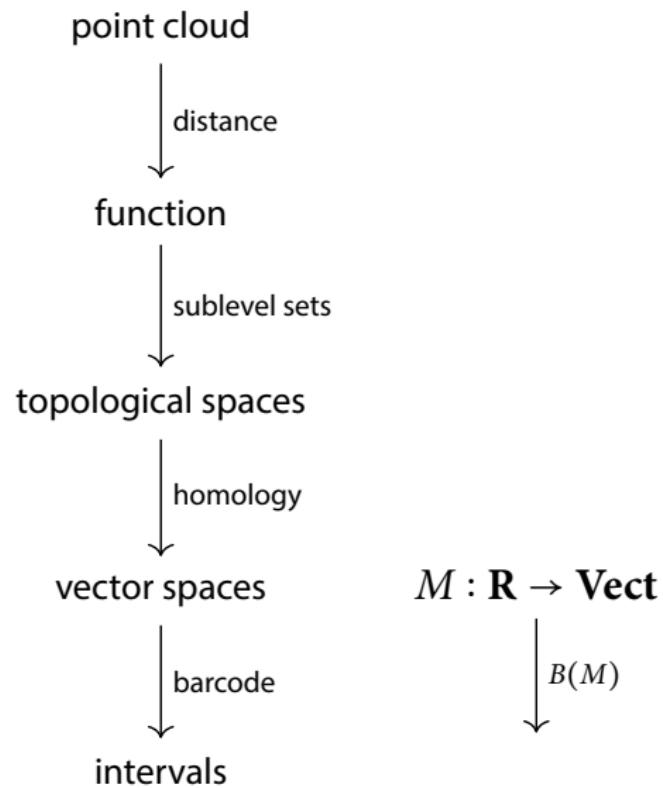
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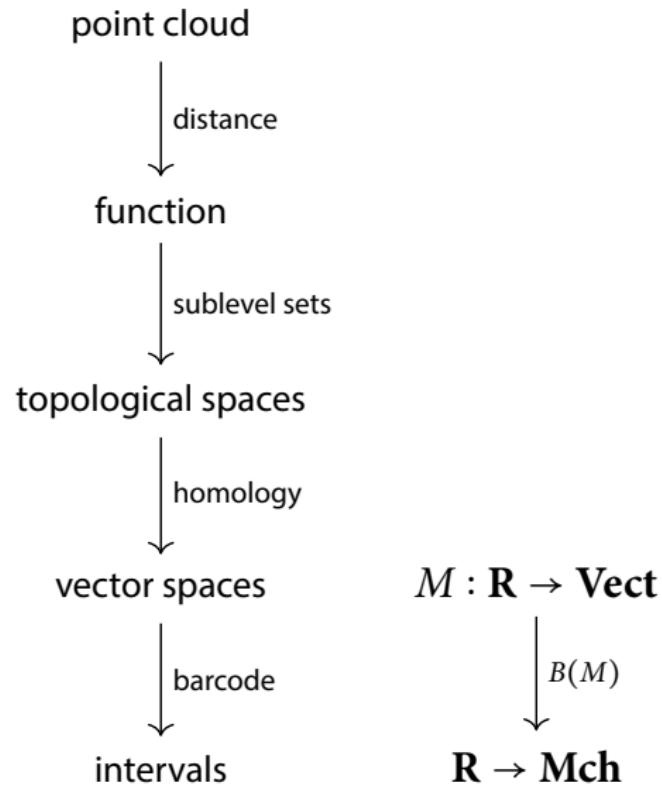
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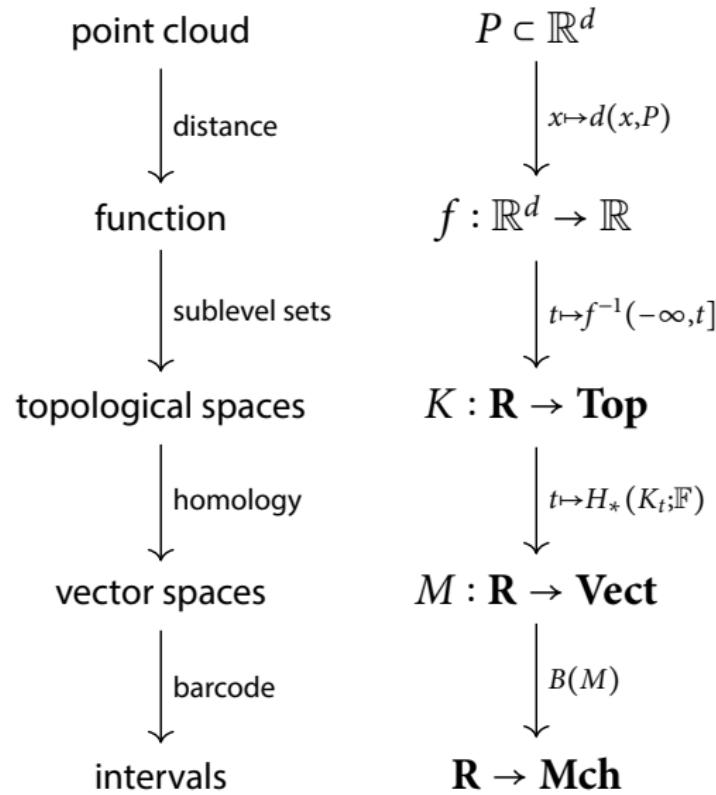
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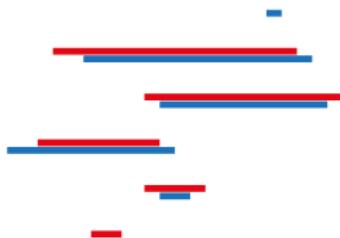
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Stability of persistence barcodes for functions

Theorem (Cohen-Steiner, Edelsbrunner, Harer 2005)

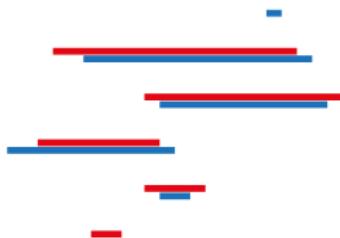
If two functions $f, g : K \rightarrow \mathbb{R}$ have distance $\|f - g\|_\infty \leq \delta$,
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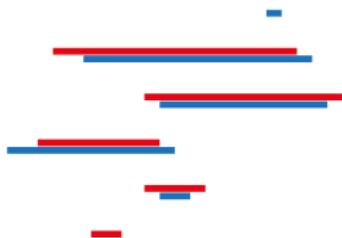


- matching $A \rightarrow B$: bijection of subsets $A' \subseteq A, B' \subseteq B$

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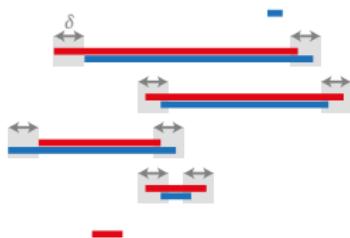


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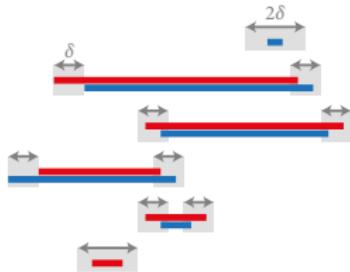


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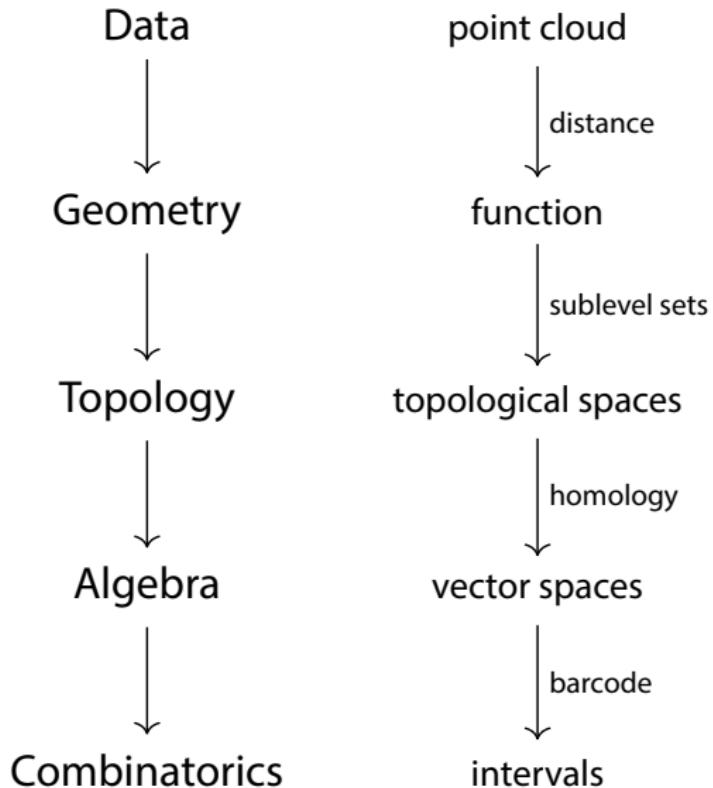
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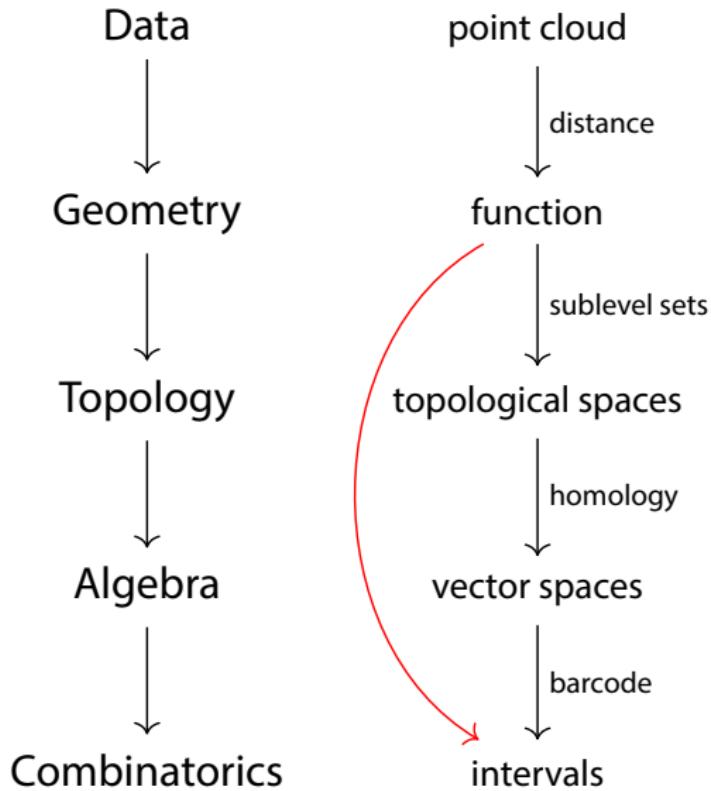


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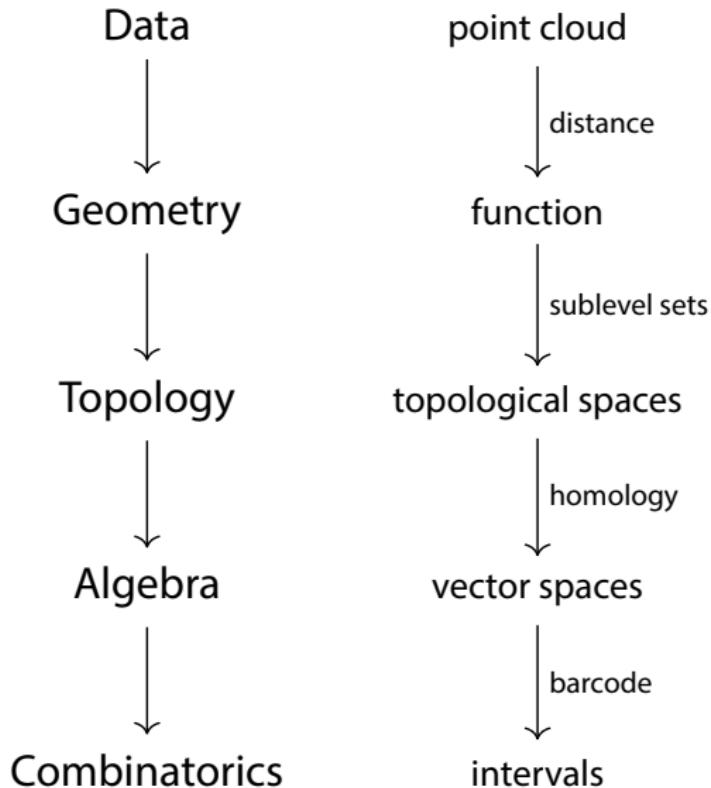
Stability for functions in the big picture



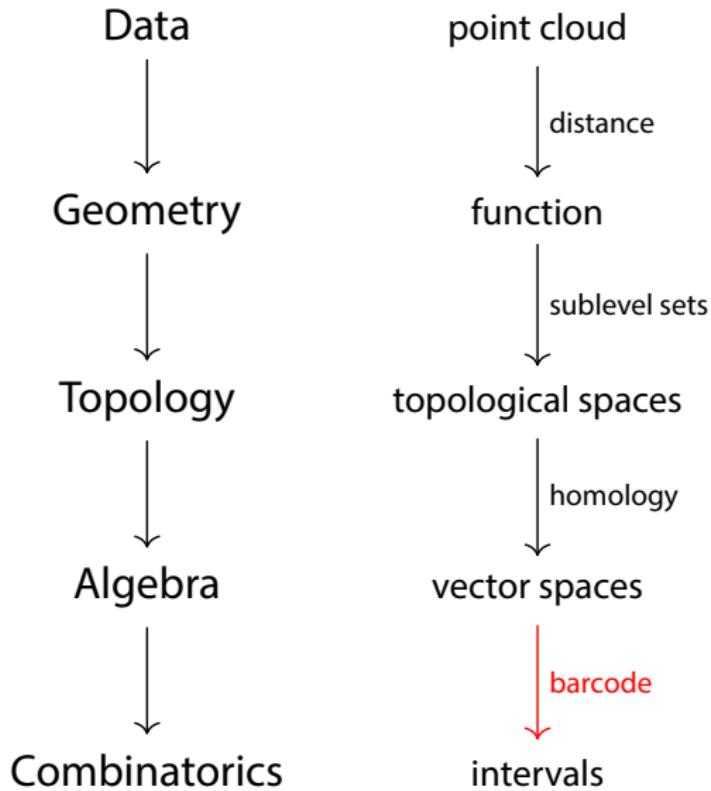
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Interleavings of sublevel sets

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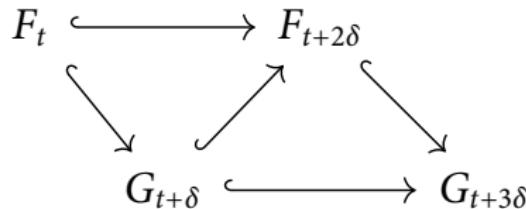
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$$\begin{array}{ccccc} H_*(F_t) & \longrightarrow & H_*(F_{t+2\delta}) & & \\ \searrow & & \nearrow & & \searrow \\ & & H_*(G_{t+\delta}) & \longrightarrow & H_*(G_{t+3\delta}) \end{array}$$

Homology is a *functor*: homology groups are interleaved too.

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Interval Persistence Modules

Let \mathbb{K} be a field. For an arbitrary interval $I \subseteq \mathbb{R}$, define the *interval persistence module* $C(I)$ by

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(shift barcode to the left by δ)



Algebraic stability of persistence barcodes

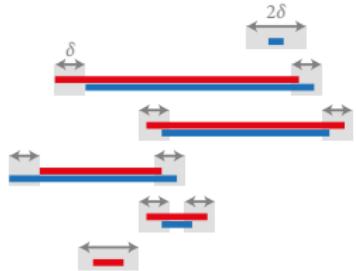
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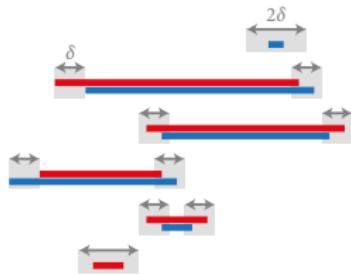
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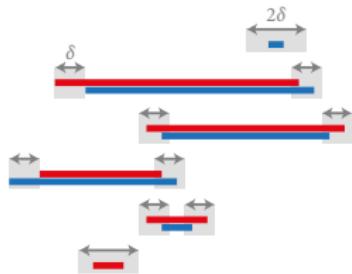


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- converse statement also holds (isometry theorem)
- indirect proof, following interpolation argument of original stability theorem

Induced matchings

Structure of submodules and quotient modules

Proposition (B, Lesnick 2014)

For a persistence submodule $K \subseteq M$:

- $B(K)$ is obtained from $B(M)$ by moving left endpoints to the right,

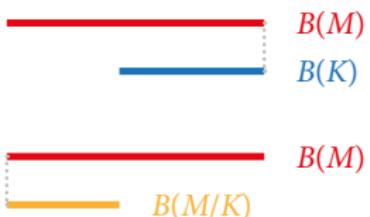


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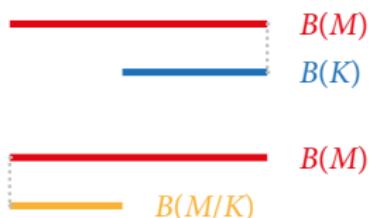


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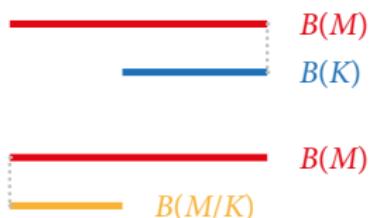
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- If multiple bars have same endpoint:
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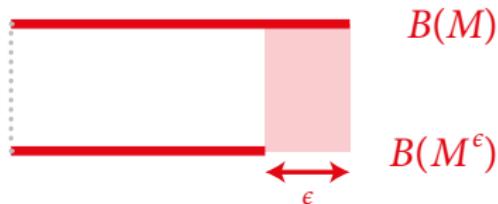
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Define M^ϵ by shrinking bars of $B(M)$ from the right by ϵ .



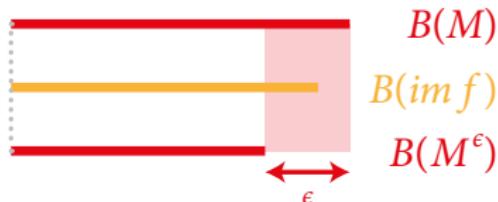
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Let $f : M \rightarrow N$ be such that $\ker f$ is ϵ -trivial: $(\ker f)^\epsilon = 0$.

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Define ${}^\epsilon N$ by shrinking bars of $B(N)$ from the left by ϵ .

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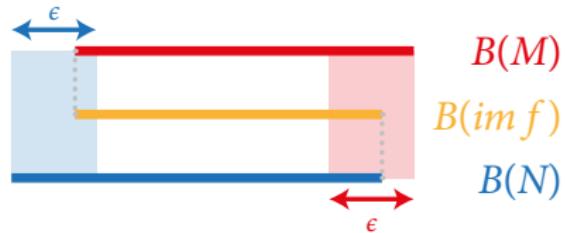
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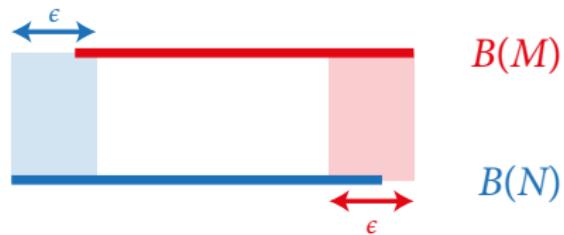
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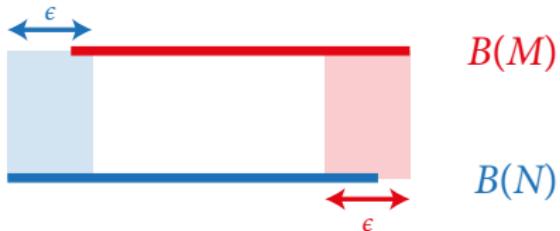
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Theorem (B, Lesnick 2014)

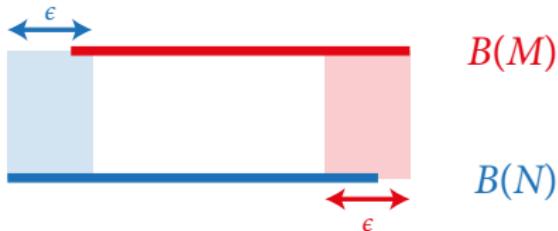
Let $f : M \rightarrow N$ be a morphism with $\ker f$ and $\text{coker } f$ ϵ -trivial.



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Theorem (B, Lesnick 2014)

Let $f : M \rightarrow N$ be a morphism with $\ker f$ and $\text{coker } f$ ϵ -trivial.
Then each interval of length $\geq \epsilon$ is matched by $B(f)$.



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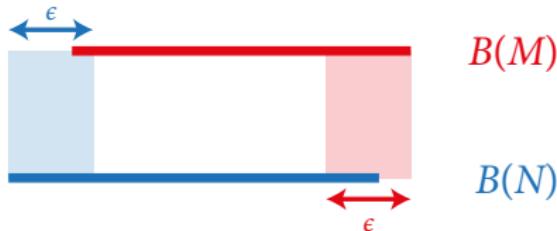
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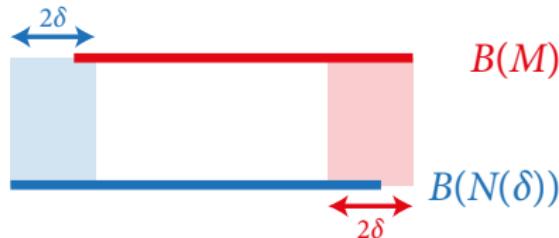
If $B(f)$ matches $[b, d) \in B(M)$ to $[b', d') \in B(N)$, then

$$b' \leq b \leq b' + \epsilon \text{ and } d - \epsilon \leq d' \leq d.$$



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Let $f : M \rightarrow N(\delta)$ be an interleaving morphism.
Then $\ker f$ and $\text{coker } f$ are 2δ -trivial.



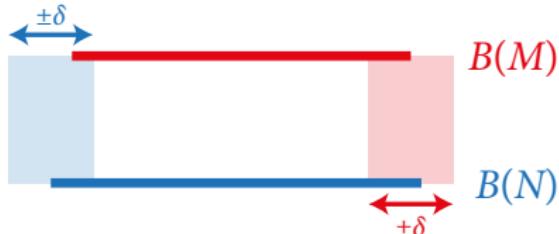
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Corollary (Algebraic stability via induced matchings)

A δ -interleaving between persistence modules induces a δ -matching of their persistence barcodes.



Stability via induced matchings



Stability via induced matchings



$B(M)$

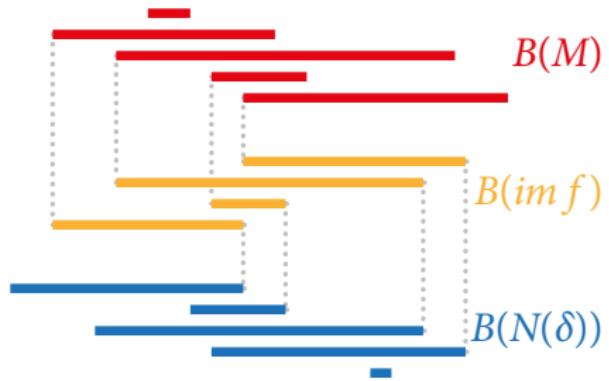


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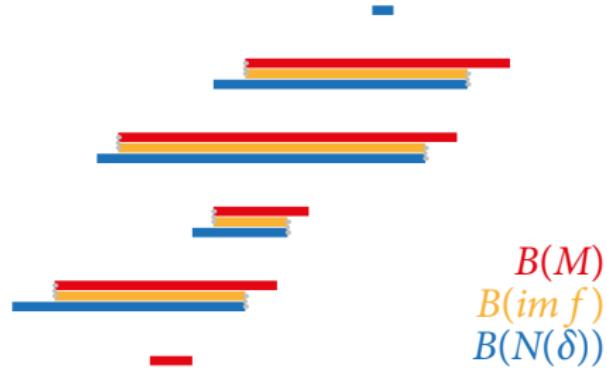
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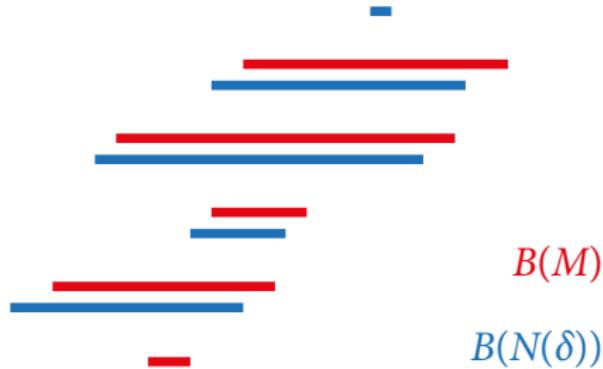
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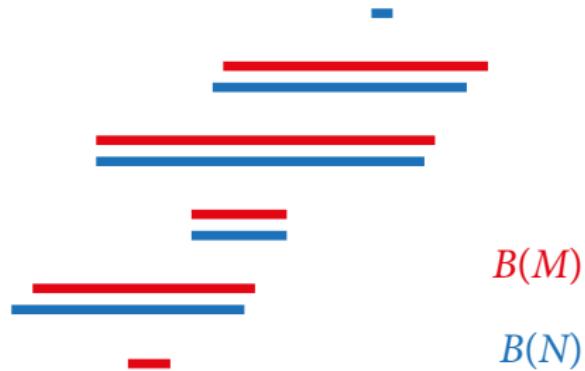
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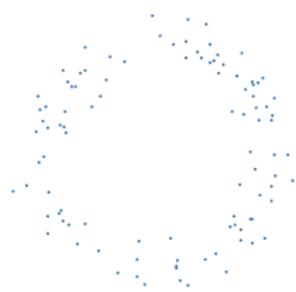


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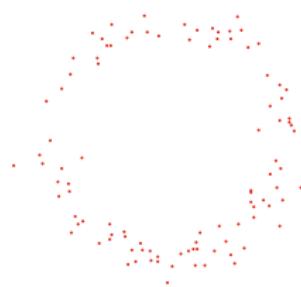


Application: connecting different point clouds

X

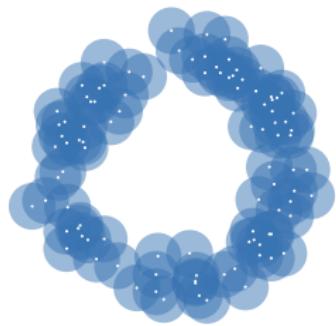


Y



Application: connecting different point clouds

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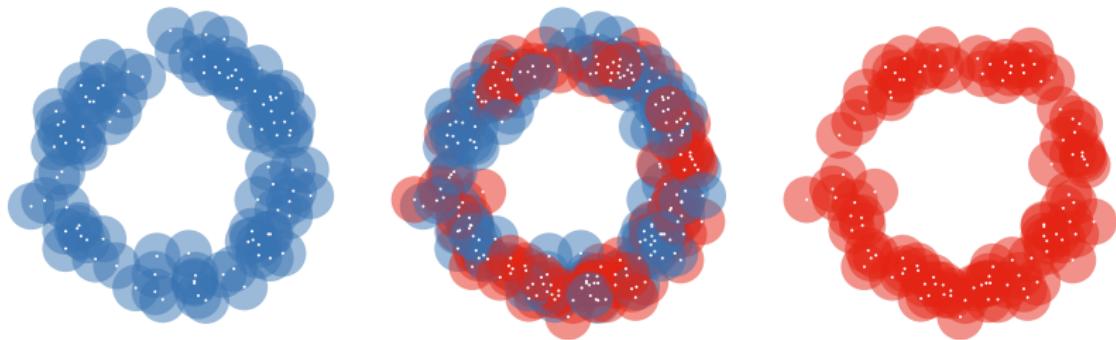


$B_r(Y)$



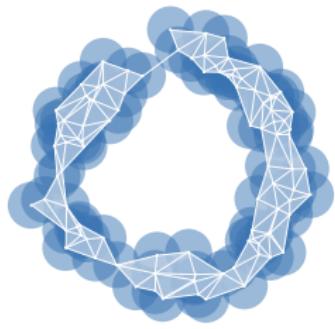
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$$B_r(X) \longleftrightarrow B_r(X \cup Y) \longleftrightarrow B_r(Y)$$

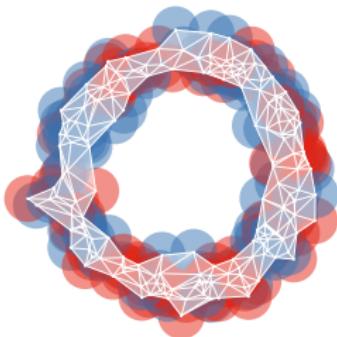


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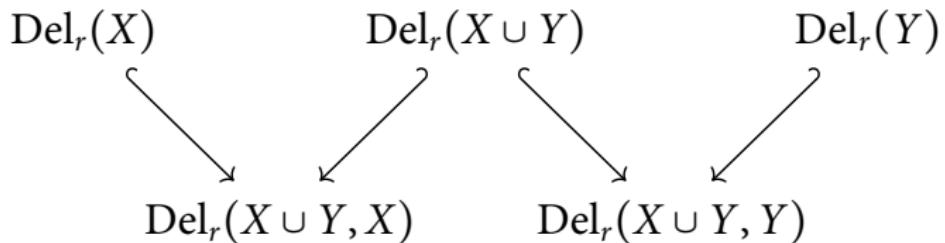
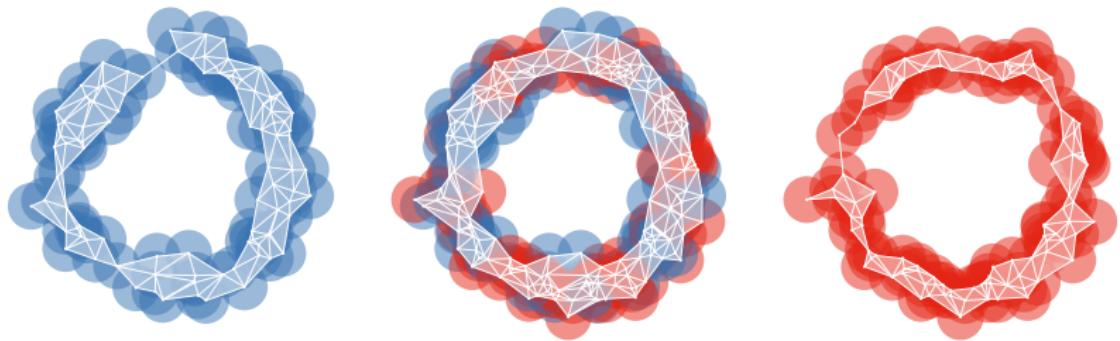
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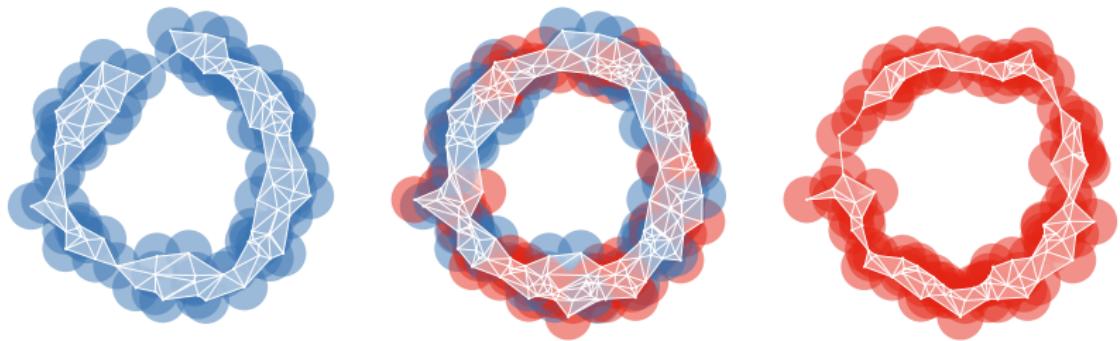
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$$B_r(X) \longleftrightarrow B_r(X \cup Y) \longleftrightarrow B_r(Y)$$



Application: connecting different point clouds

$$B_r(X) \longleftrightarrow B_r(X \cup Y) \longleftrightarrow B_r(Y)$$



$$\begin{array}{ccc} \text{Del}_r(X) & \text{Del}_r(X \cup Y) & \text{Del}_r(Y) \\ \searrow & \swarrow^{\simeq} & \searrow \\ \text{Del}_r(X \cup Y, X) & & \text{Del}_r(X \cup Y, Y) \end{array}$$