

Bridging Morse theory and persistent homology of geometric complexes

Ulrich Bauer

TUM

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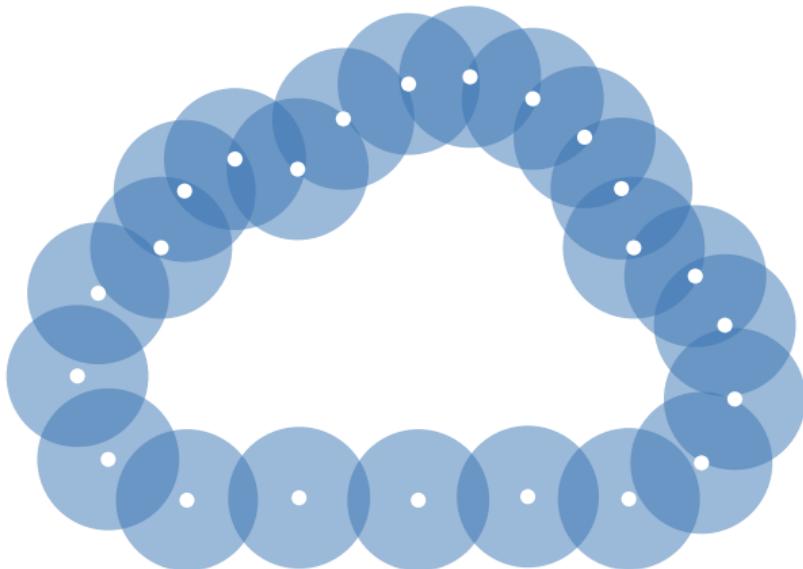
SFB
TRR
109
Discretization
in Geometry
and Dynamics

Technische
Universität
München

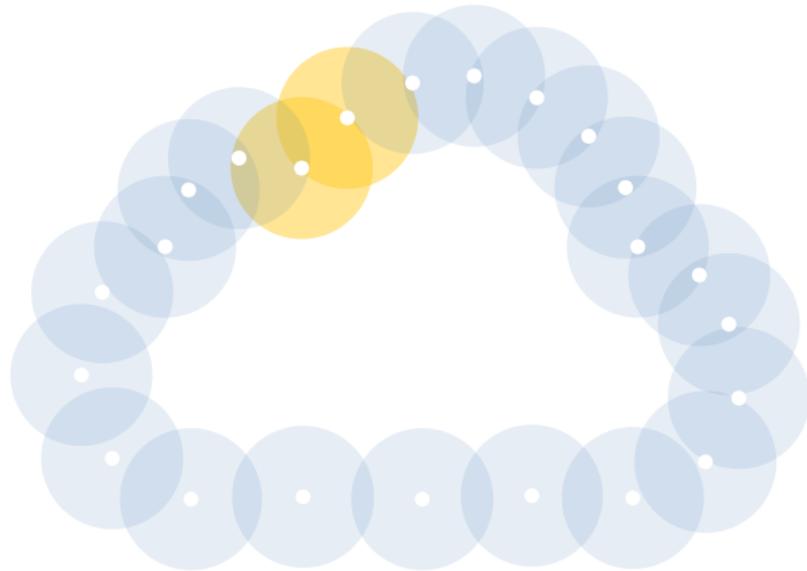


MCMUL
Munich Center for Machine Learning

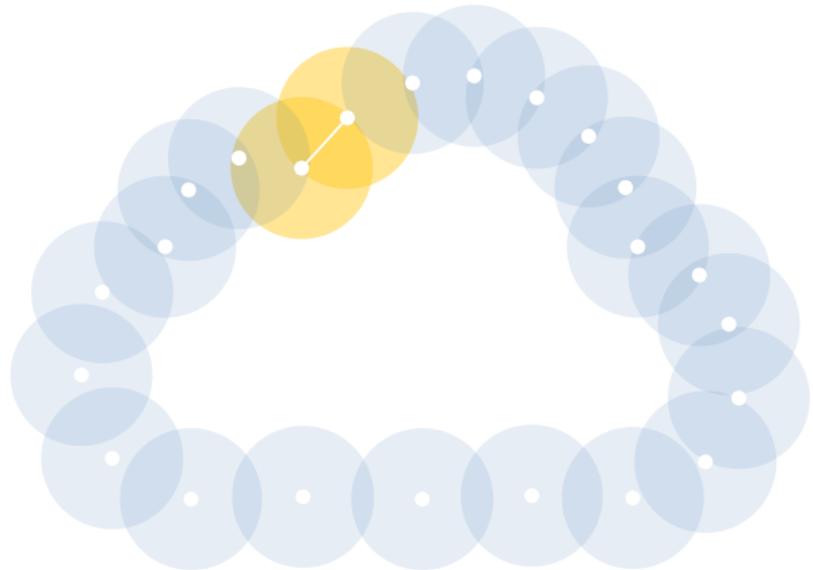
Čech complexes



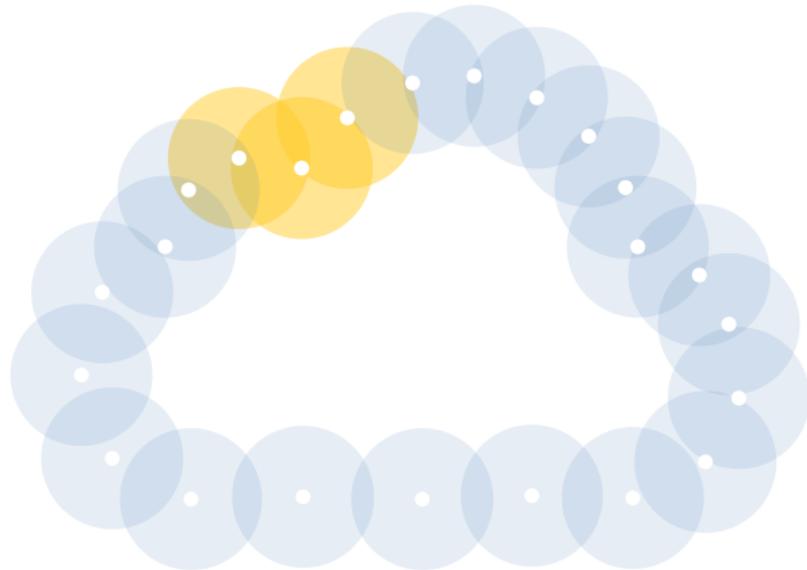
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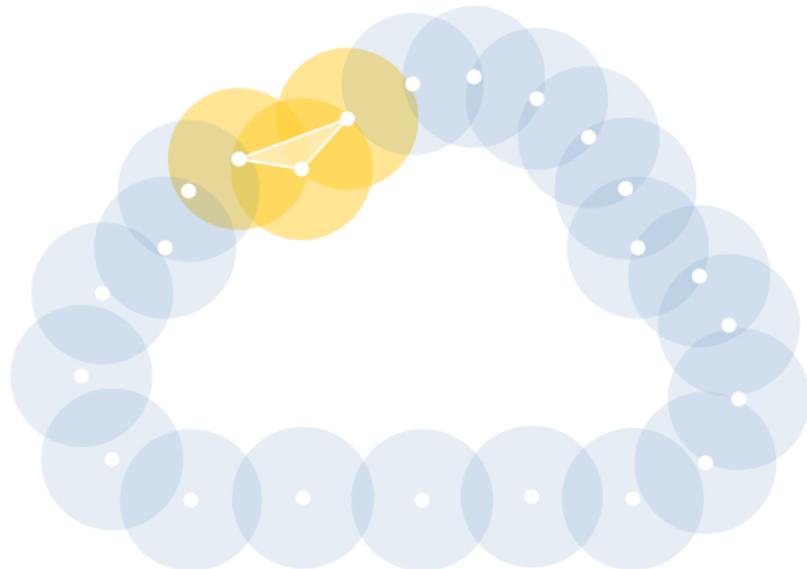
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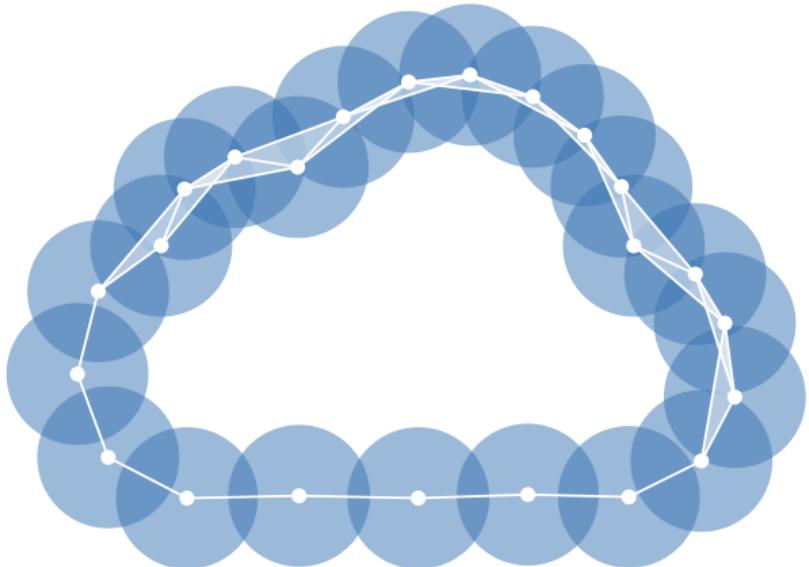
Čech complexes



Čech complexes



Čech complexes



Nerves

Definition (Alexandrov 1928)

Let $\mathcal{U} = (U_i)_{i \in I}$ be a cover of a space X . The *nerve* of \mathcal{U} is the simplicial complex

$$\text{Nrv}(\mathcal{U}) = \{J \subseteq I \mid 0 < |J| < \infty \text{ and } \bigcap_{i \in J} U_i \neq \emptyset\}$$

recording the nonempty intersections of cover elements.

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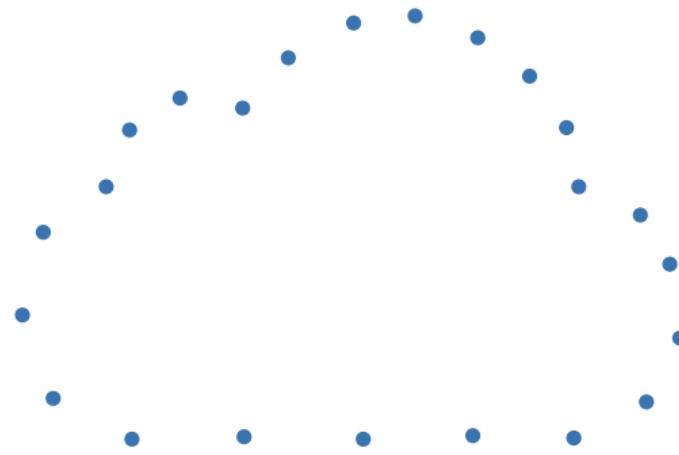


U. Bauer, M. Kerber, F. Roll, and A. Rolle

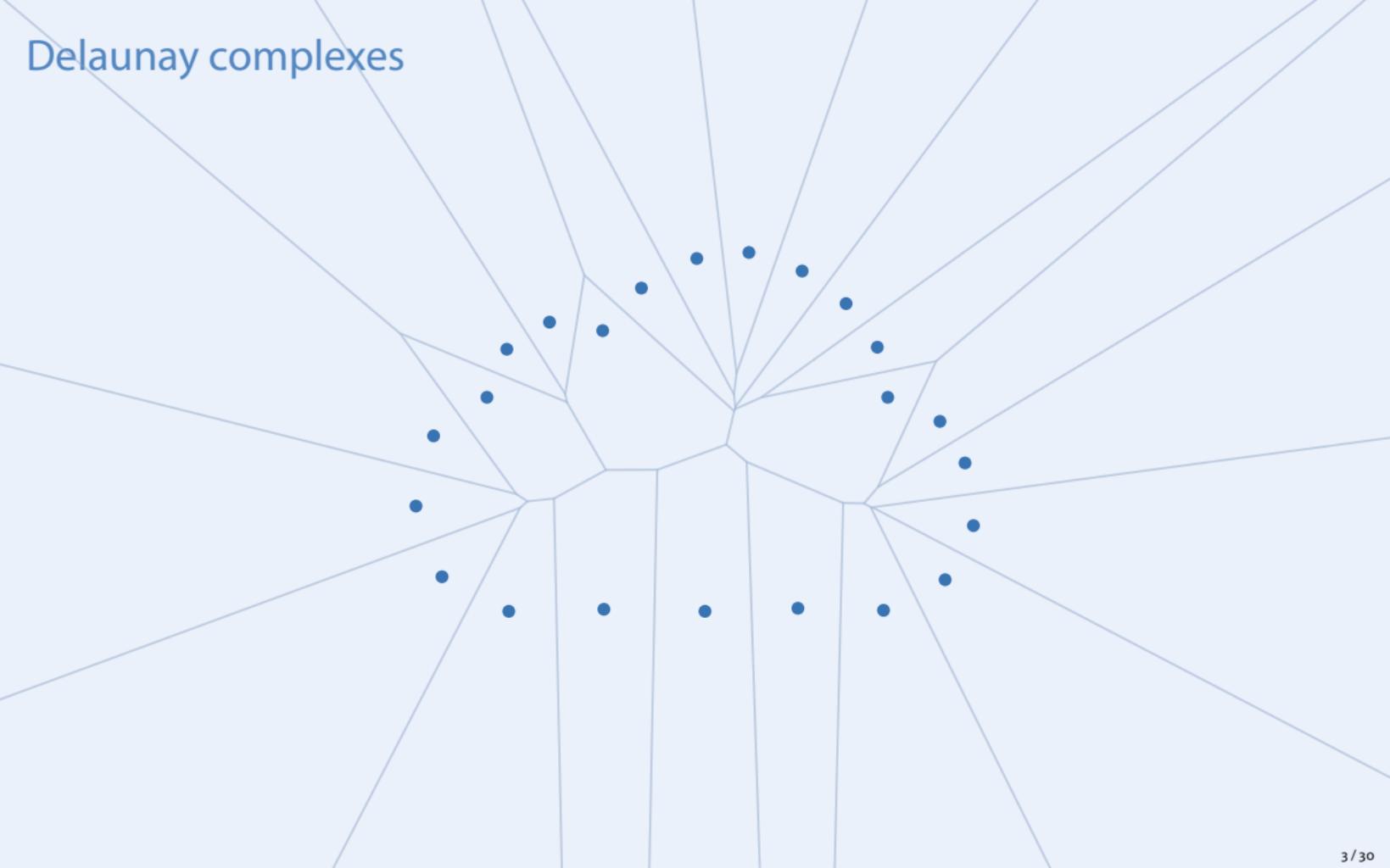
A Unified View on the Functorial Nerve Theorem and its Variations

Expositiones Mathematicae, 2023. doi:10.1016/j.exmath.2023.04.005

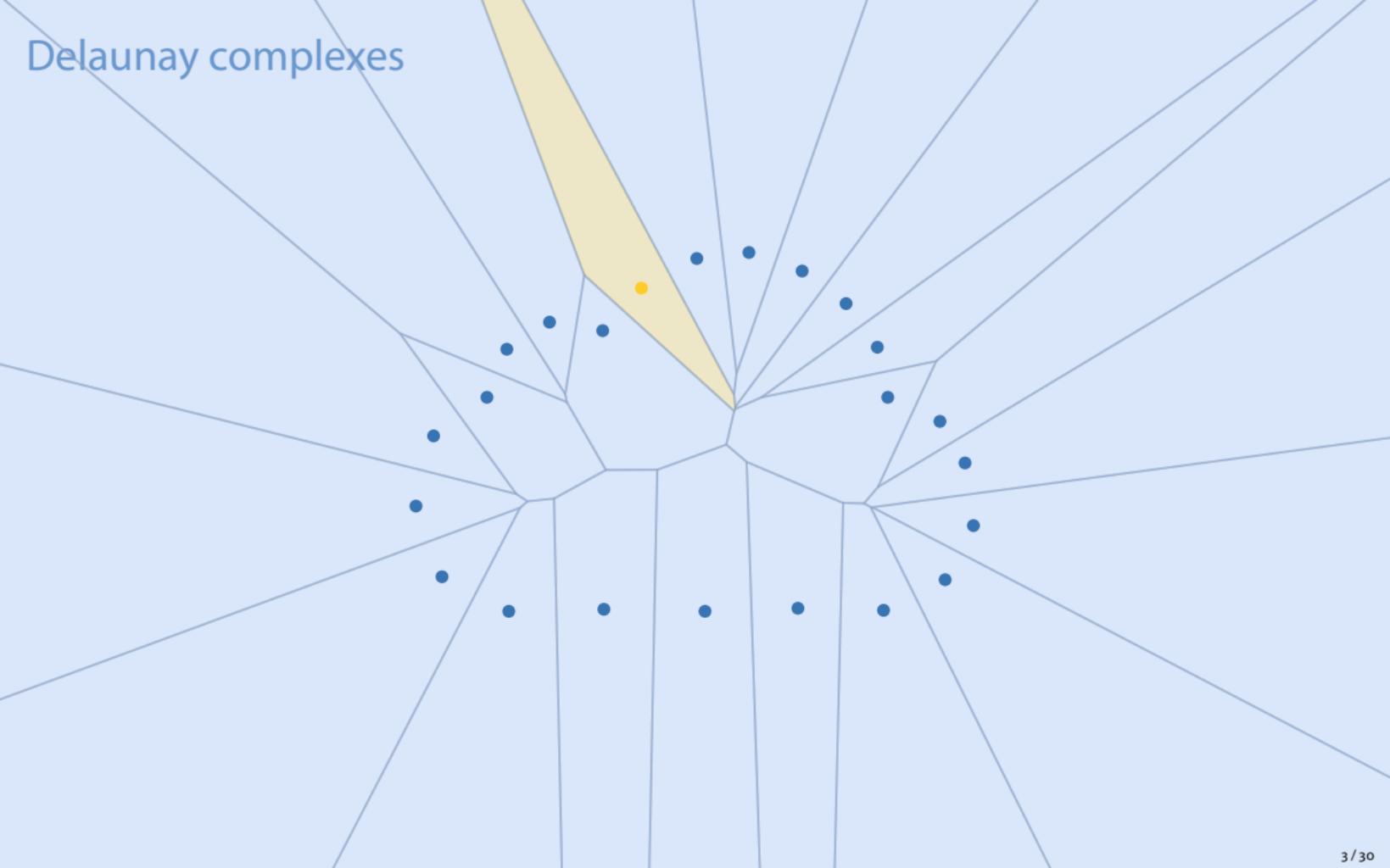
Delaunay complexes



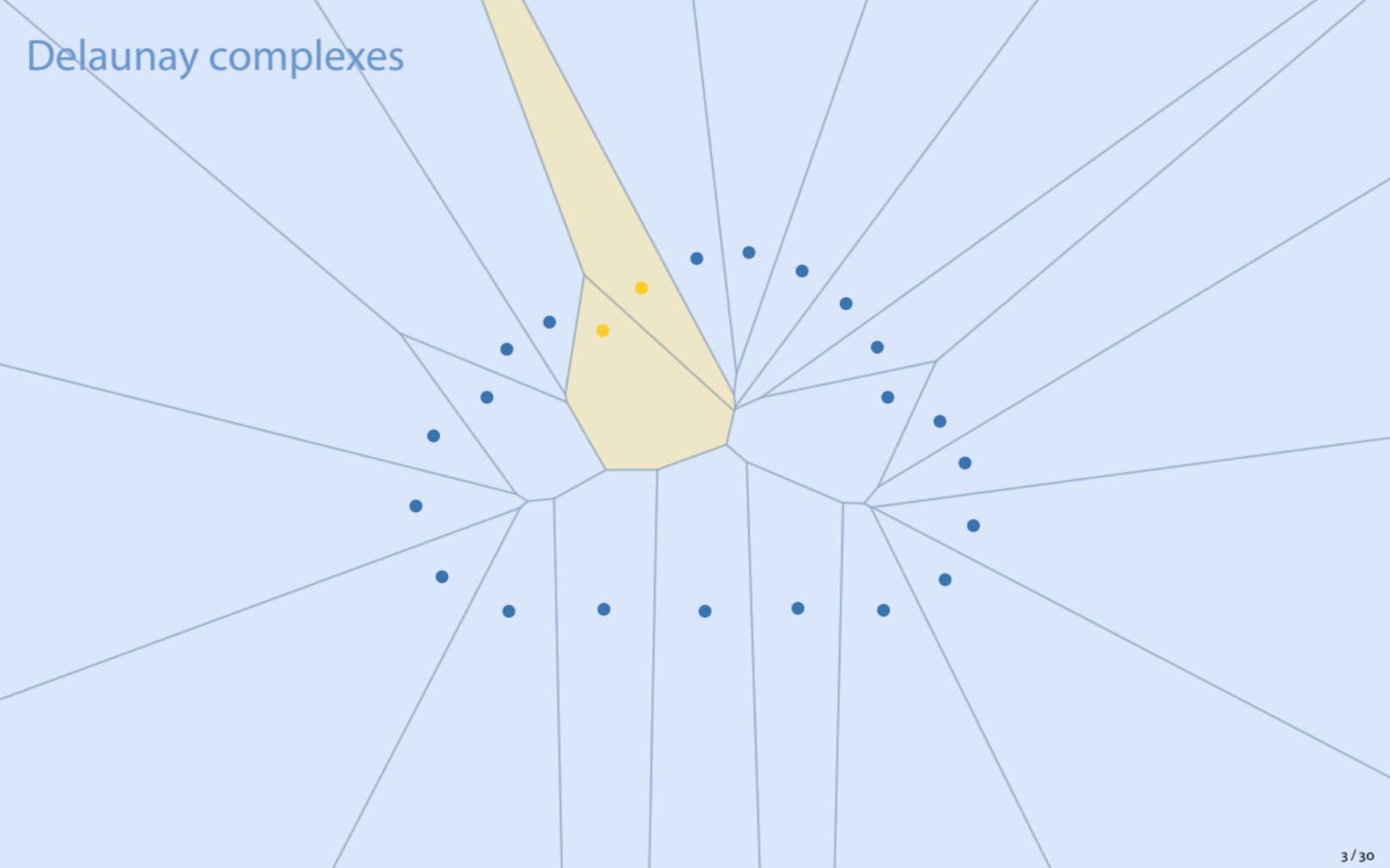
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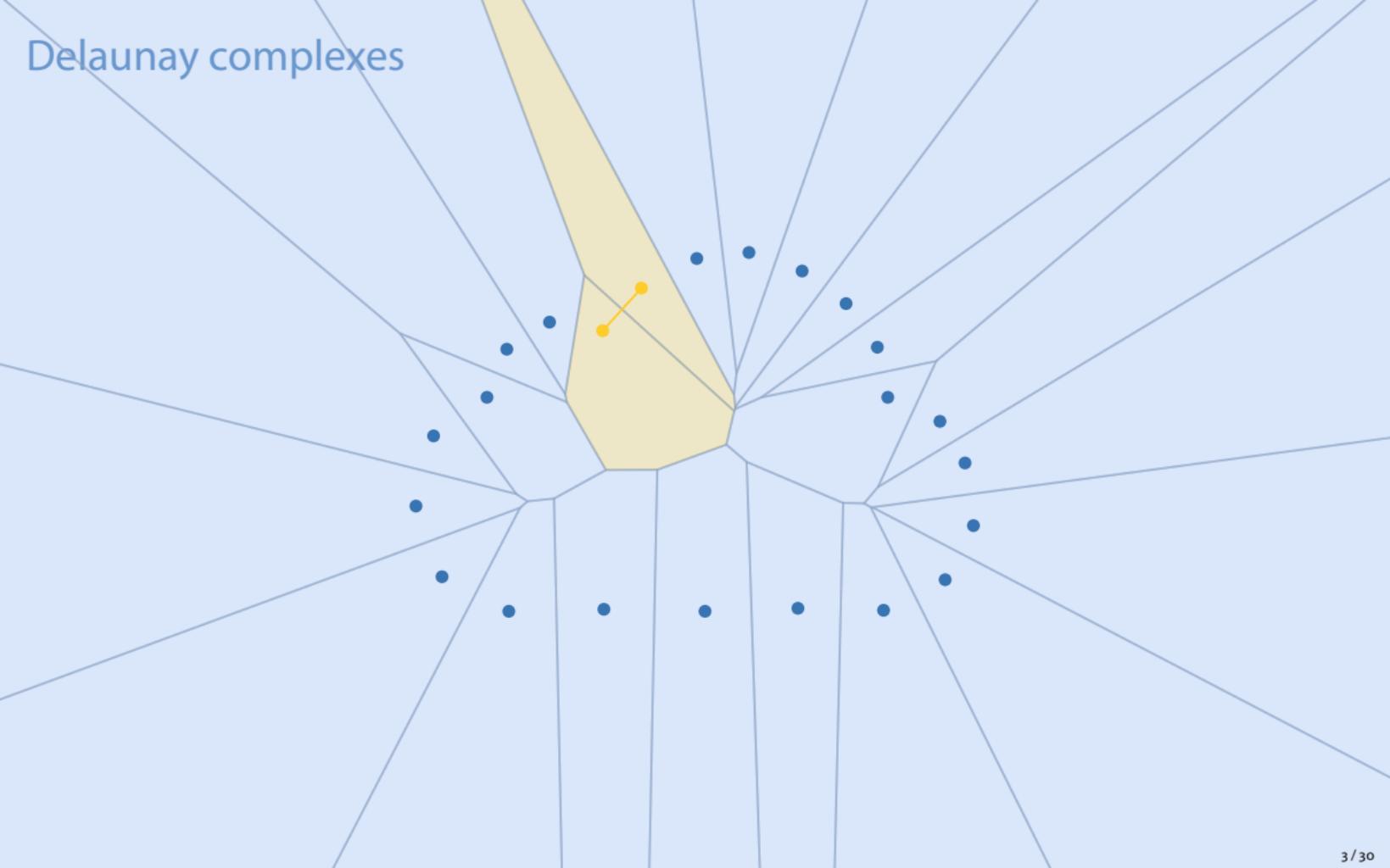
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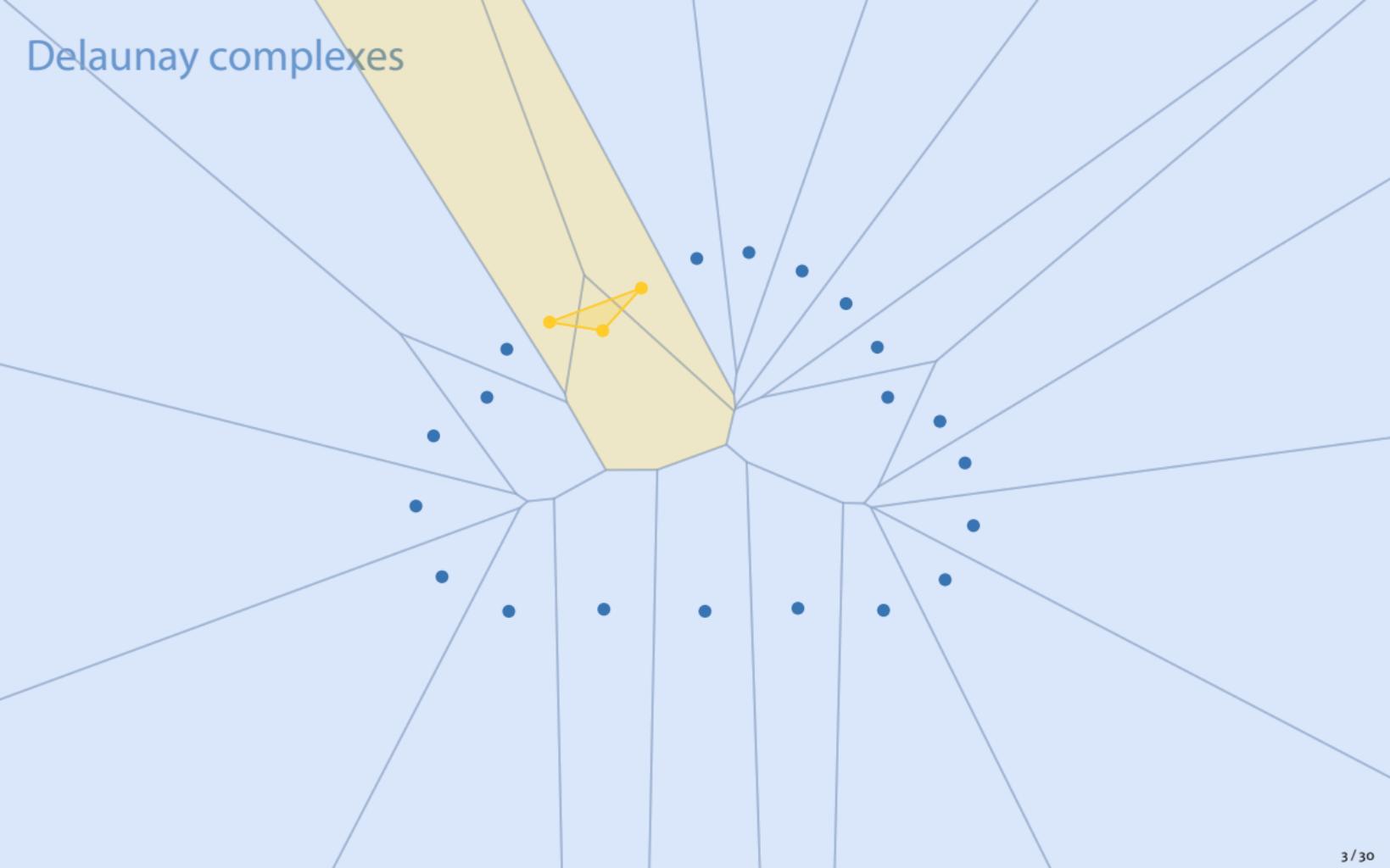
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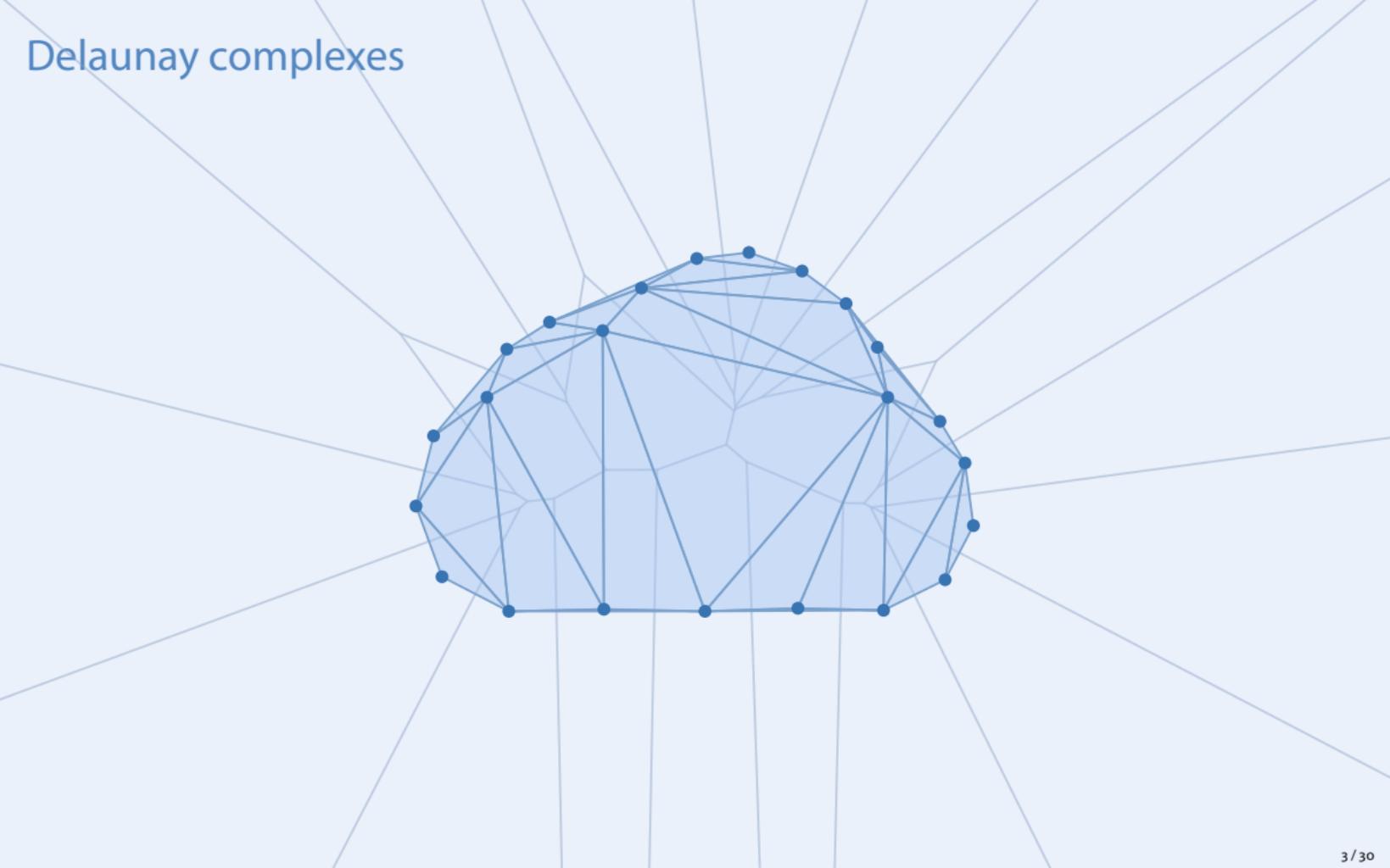
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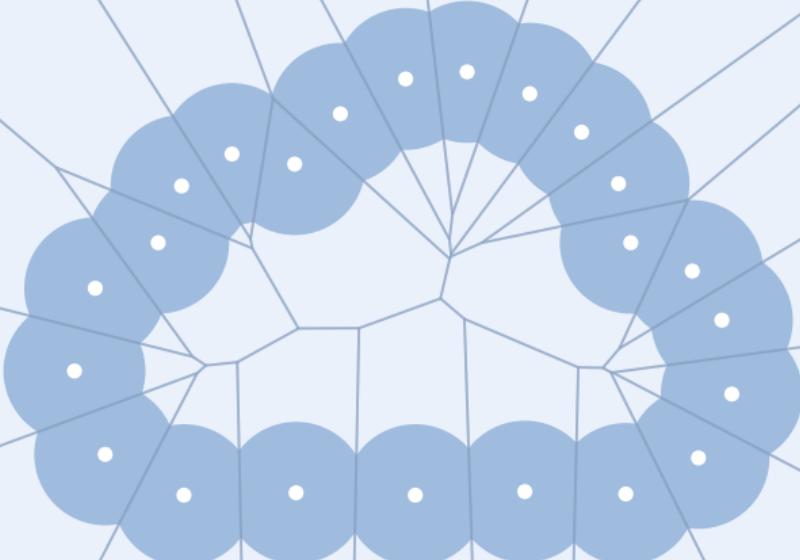
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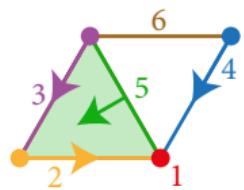
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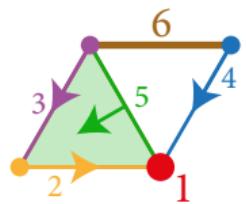
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Discrete Morse theory



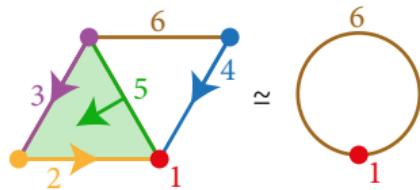
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Theorem (Forman 1998)

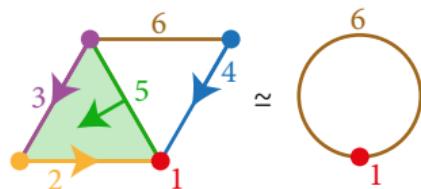
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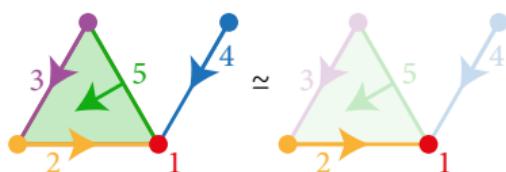
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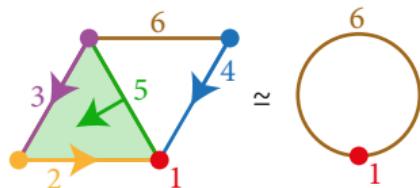
Discrete Morse functions – and their gradients – encode collapses of sublevel sets:



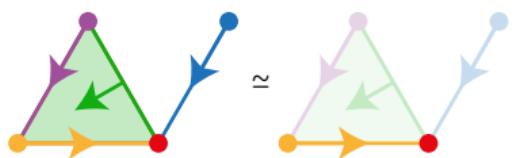
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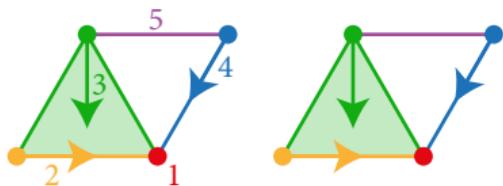


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Generalizing discrete Morse theory

Generalized gradients consist of intervals (in the face poset) instead of just facet pairs:



Morse theory for Čech and Delaunay complexes

Proposition (B, Edelsbrunner 2014)

The Čech complexes and the Delaunay complexes (alpha shapes) are sublevel sets of (generalized) discrete Morse functions.

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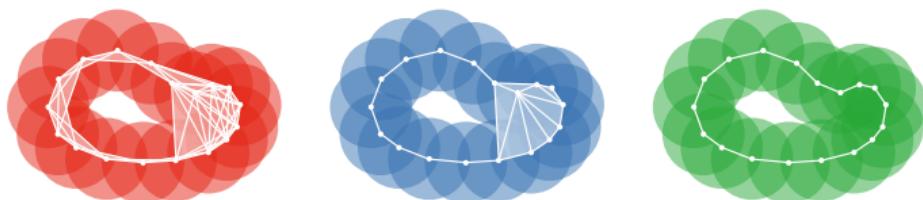
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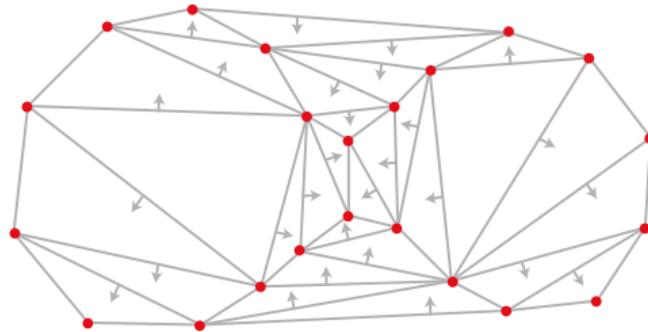
Čech, Delaunay, and Wrap complexes (at any scale r) are related by collapses encoded by a single discrete gradient field:

$$\text{Cech}_r X \searrow \text{Del}_r X \searrow \text{Wrap}_r X.$$



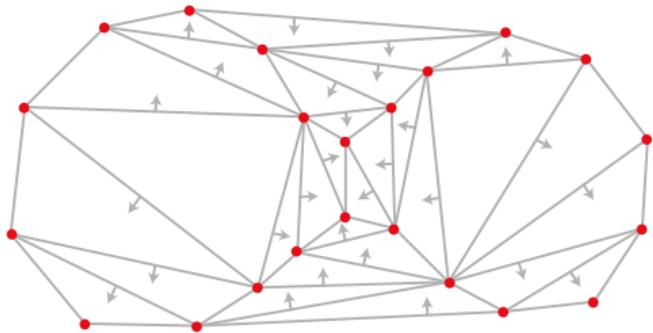
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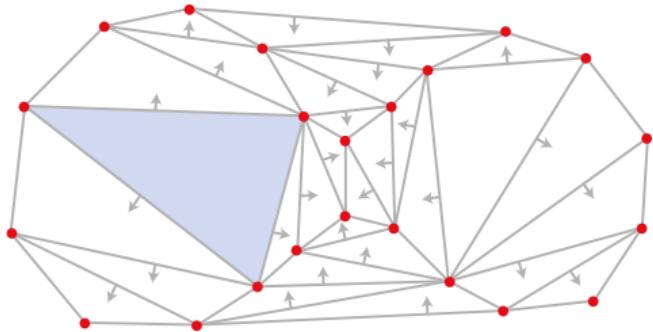
Definition (Edelsbrunner 1995; B, Edelsbrunner 2017)

$\text{Wrap}_r(X)$ is the *descending complex* of V on $\text{Del}_r X$: the smallest subcomplex of $\text{Del}_r X$ that

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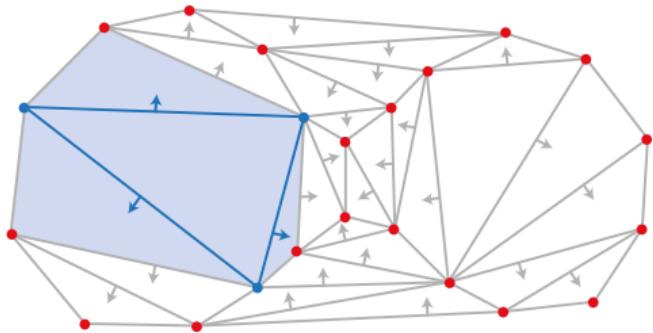
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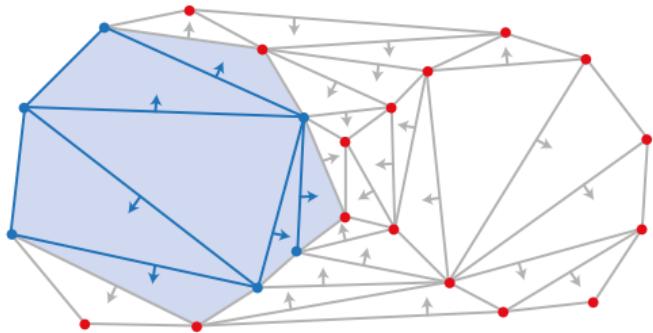
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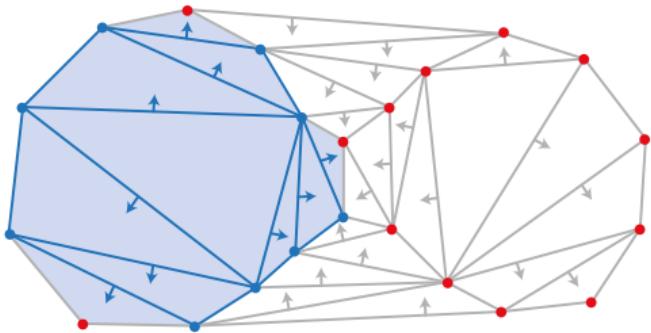
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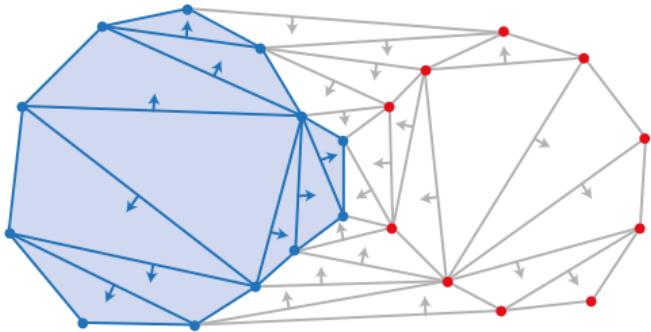
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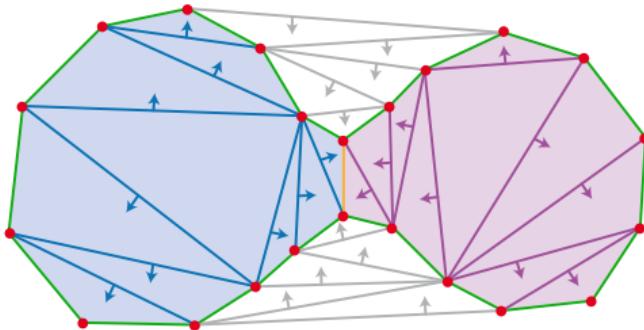
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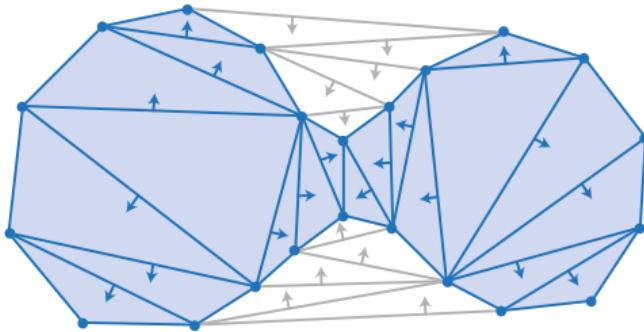
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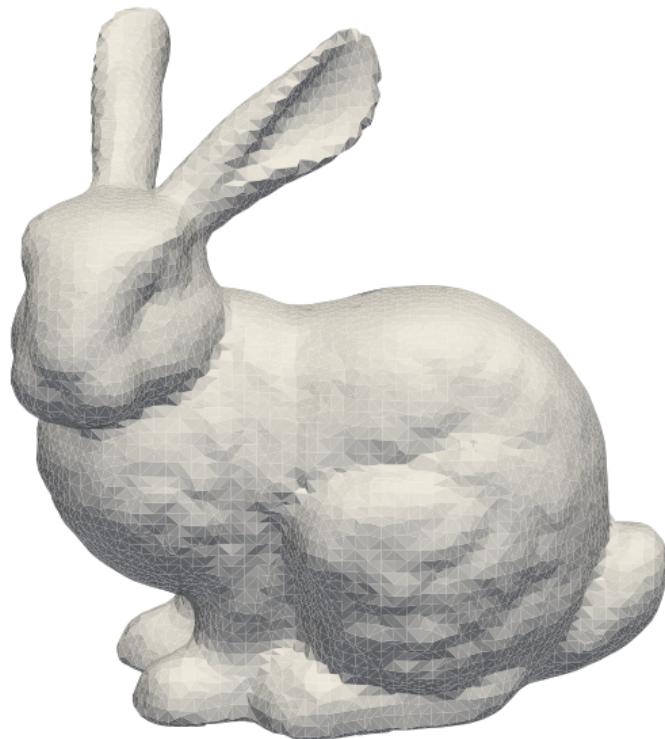
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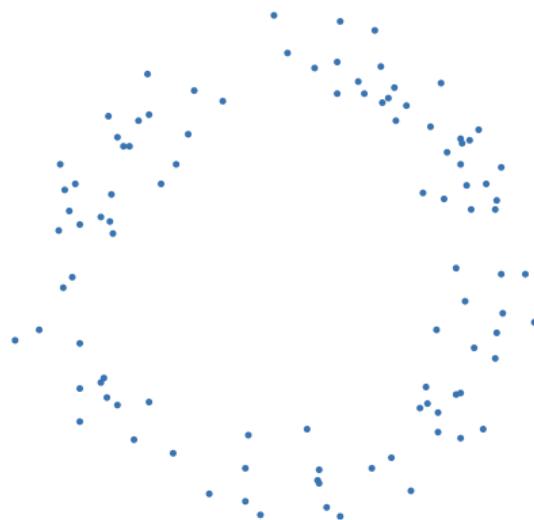
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Vietoris–Rips complexes

For a metric space X , the *Vietoris–Rips complex* at $t > 0$ is the simplicial complex

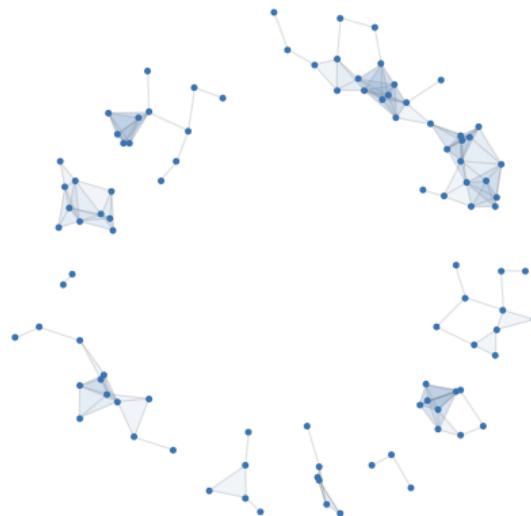
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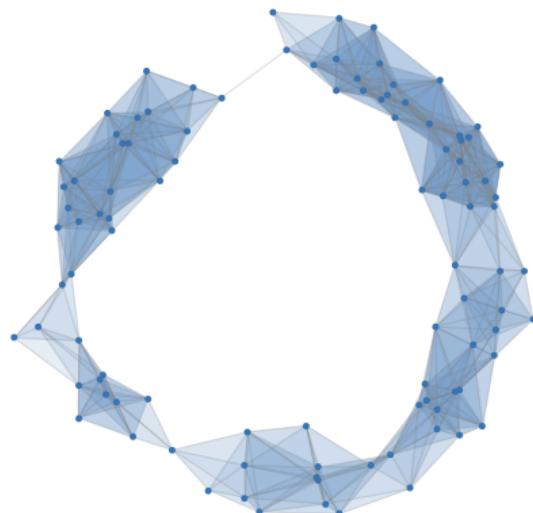
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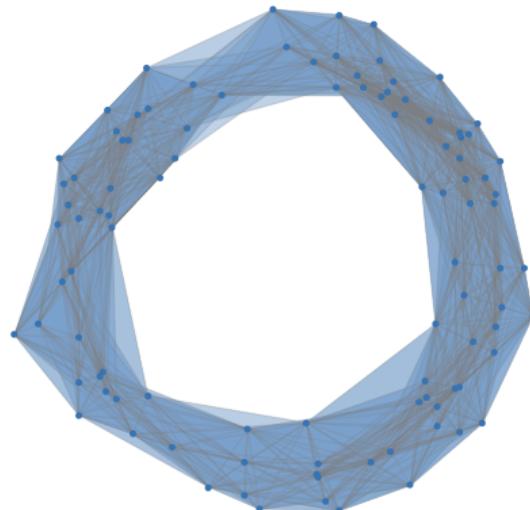
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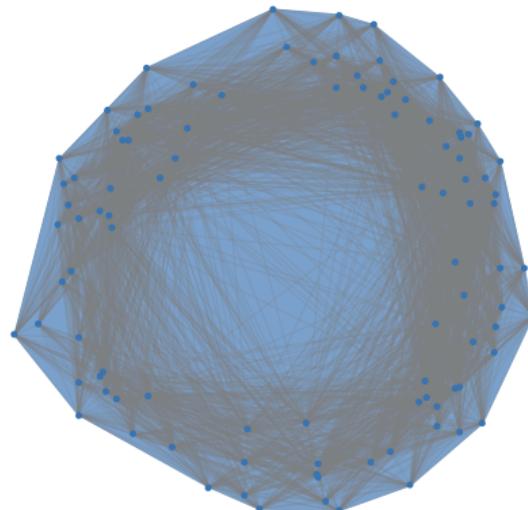
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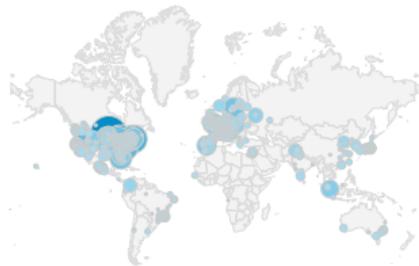
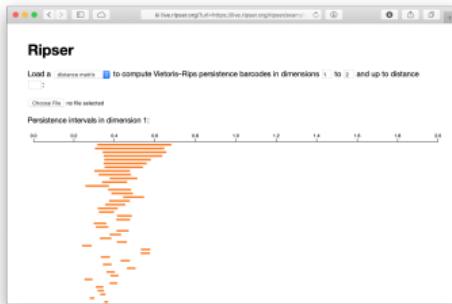
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Ripser: software for computing Vietoris–Rips persistence barcodes

Open source software (ripser.org)

- significantly faster / more memory efficient than previous codes
- de facto standard for topological data analysis applications
- most popular persistent homology project on GitHub



Ripser users from 935 different locations

Computational improvements based on

- *implicit matrix representations*
- *apparent pairs*, connecting persistence to discrete Morse theory

Apparent pairs

Ripser uses the following pairing of simplices (for the filtration, we break ties lexicographically):

Definition (B 2016, 2021)

In a simplexwise filtration $(K_i = \{\sigma_1, \dots, \sigma_i\})_i$, two simplices (σ_i, σ_j) form an *apparent pair* if

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The apparent pairs are both gradient pairs and persistence pairs.

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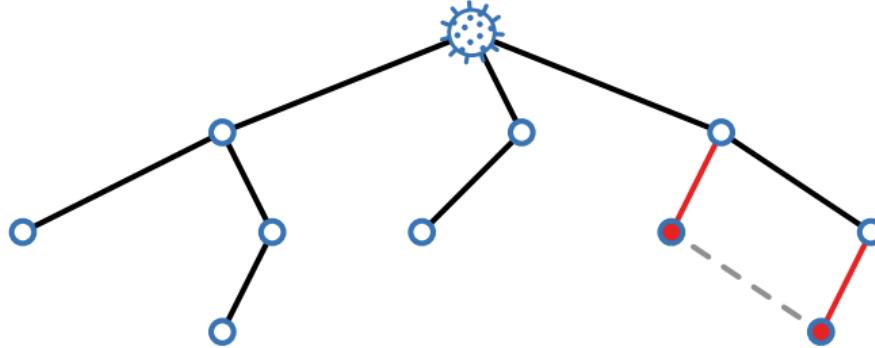
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Proposition (B, Roll 2022)

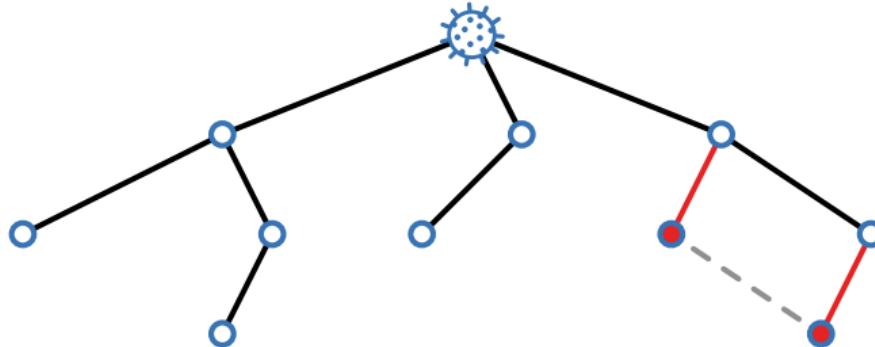
Consider a generalized discrete Morse function f , and the simplexwise filtration from lexicographic refinement. Then the zero-persistence apparent pairs form a gradient that refines the gradient of f .

Topology of viral evolution



Joint work with: A. Ott, M. Bleher, L. Hahn (Heidelberg), R. Rabadañ, J. Patiño-Galindo (Columbia), M. Carrière (INRIA)

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Observation: Ripser runs unusually fast on genetic distance data

- SARS-CoV2 RNA sequences (spike protein)
- 25556 data points (2.8×10^{12} simplices in 2-skeleton)
- 120 s computation time (with data points ordered appropriately)

The Rips Contractibility Lemma

Theorem (Rips; Gromov 1988)

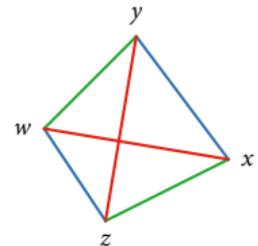
Let X be a δ -hyperbolic geodesic metric space. Then $\text{Rips}_t(X)$ is contractible for all $t \geq 4\delta$.

Gromov-hyperbolicity

Definition (Gromov 1988)

A metric space X is δ -hyperbolic (for $\delta \geq 0$) if for all $w, x, y, z \in X$ we have

$$d(w, x) + d(y, z) \leq \max\{d(w, y) + d(x, z), d(w, z) + d(x, y)\} + 2\delta.$$



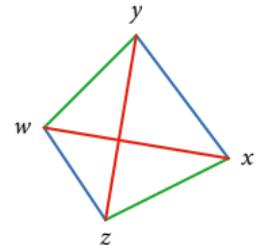
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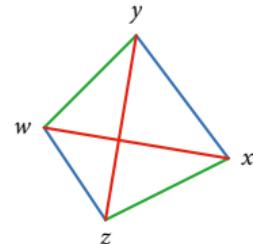
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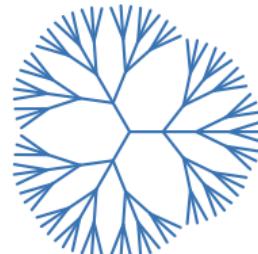
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- The 0-hyperbolic spaces are precisely the metric trees and their subspaces.



Rips contractibility for non-geodesic spaces

- What about non-geodesic spaces? In particular, finite metric spaces?
- What about collapsiblility?
- What about the filtration?
- Connection to apparent pairs/Ripser?

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- What about the filtration?
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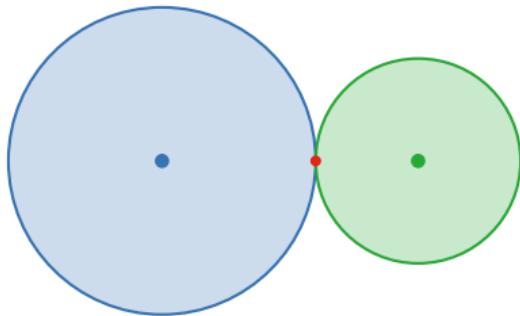
Theorem (B, Roll 2022)

Let X be a finite δ -hyperbolic space. Then there is a single discrete gradient encoding the collapses

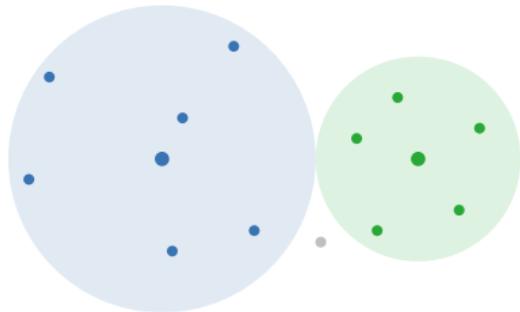
$$\text{Rips}_u(X) \searrow \text{Rips}_t(X) \searrow \{\ast\}$$

for all $u > t \geq 4\delta + 2\nu$, where ν is the geodesic defect of X .

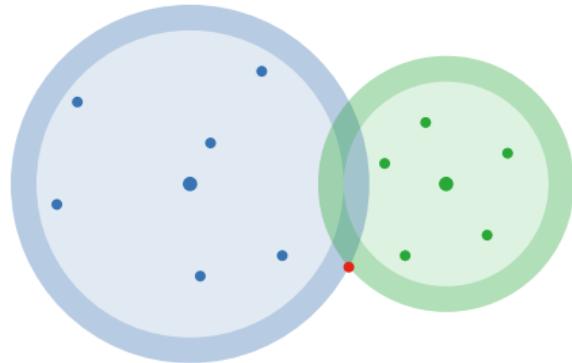
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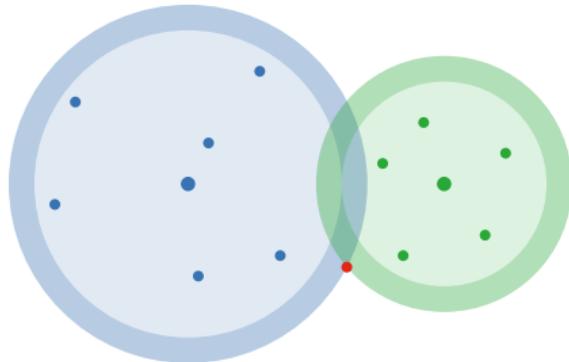
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Geodesic defect



Definition (Bonk, Schramm 2000)

A metric space X is ν -geodesic if for all points $x, y \in X$ and all $r, s \geq 0$ with $r + s = d(x, y)$ we have

$$B_{r+\nu}(x) \cap B_{s+\nu}(y) \neq \emptyset.$$

The infimum of all such ν is the *geodesic defect* of X .

The diameter function of generic trees

Proposition (B, Roll 2022)

Consider a finite weighted tree (V, E) with a generic path length metric (distinct pairwise distances). Then the diameter function $\text{diam}: \Delta(V) \rightarrow \mathbb{R}$ is a generalized discrete Morse function.

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- The apparent pairs refine this gradient.

Collapsing Vietoris–Rips complexes of generic trees

Theorem (B, Roll 2022)

The discrete gradient of the diameter function for a generic tree metric space X induces the collapses

$$\text{Rips}_t(X) \searrow T_t$$

for all $t \in \mathbb{R}$, with

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for $t \geq \max l(E)$, and

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whenever $(t, u] \cap l(E) = \emptyset$.

In particular, the persistent homology is trivial in degrees > 0 .

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Example: phylogenetic trees

- non-generic tree metric
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- The apparent gradient in turn refines the perturbed gradient (for a suitable vertex ordering).

Collapsing Rips complexes of trees with apparent pairs

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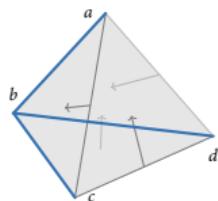
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The apparent pairs gradient for this total order induces the same collapses as in the generic case:

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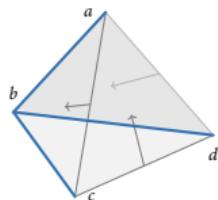
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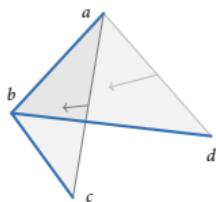
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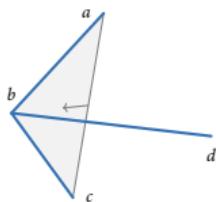
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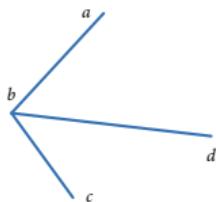
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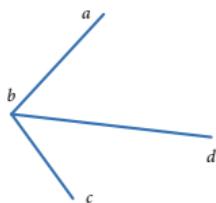
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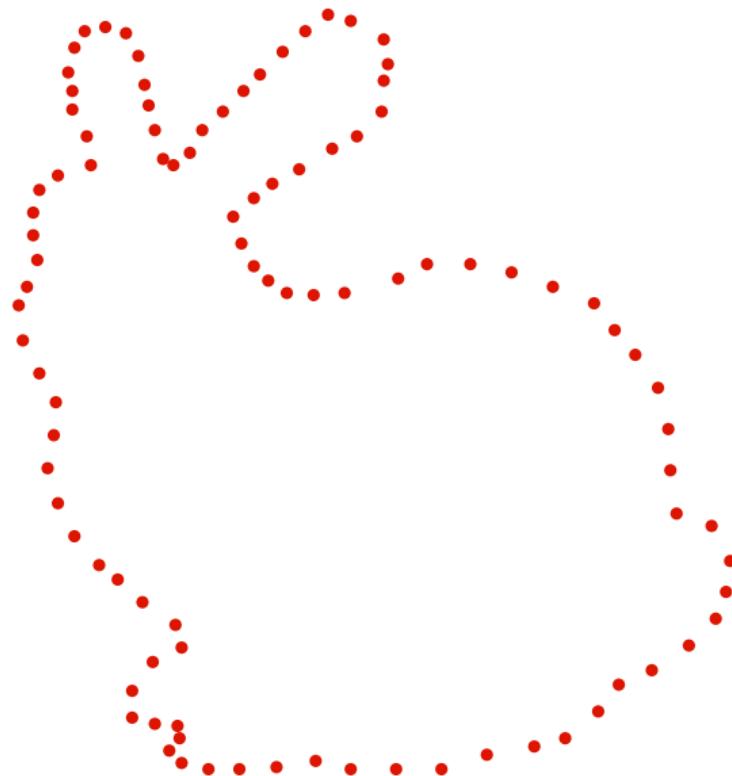
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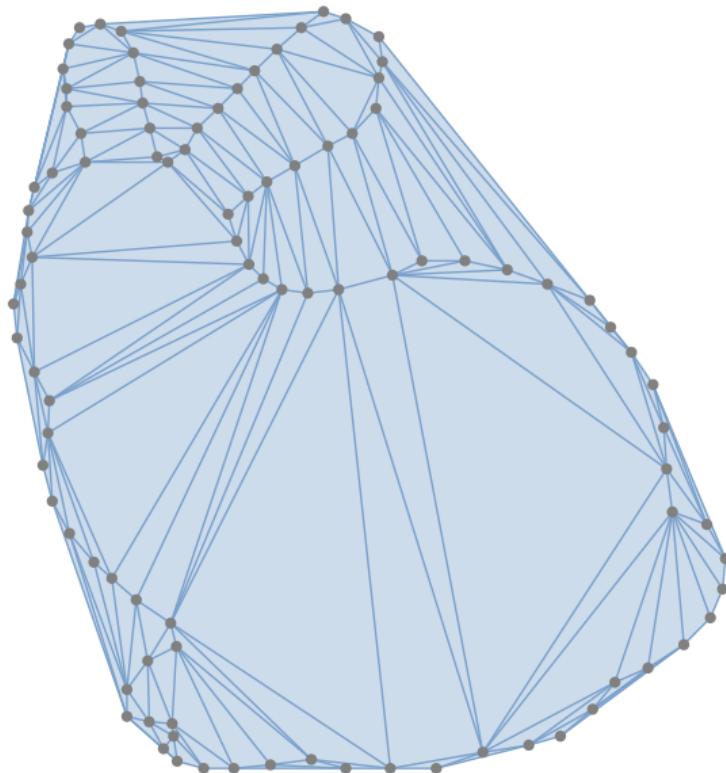


- Explains why Ripser is extraordinarily fast on genetic distances (tree-like)

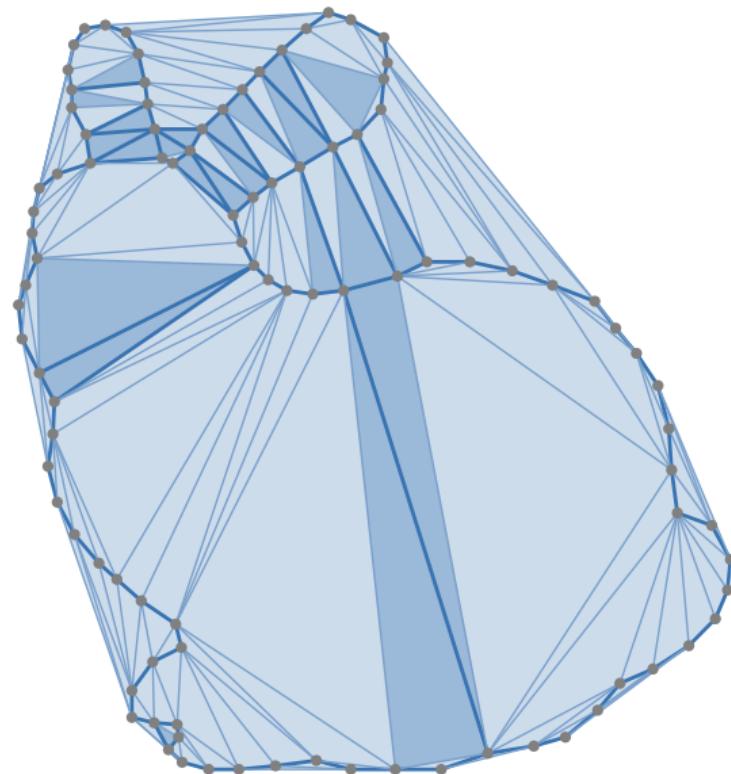
Wrap complexes and lexicographically minimal cycles



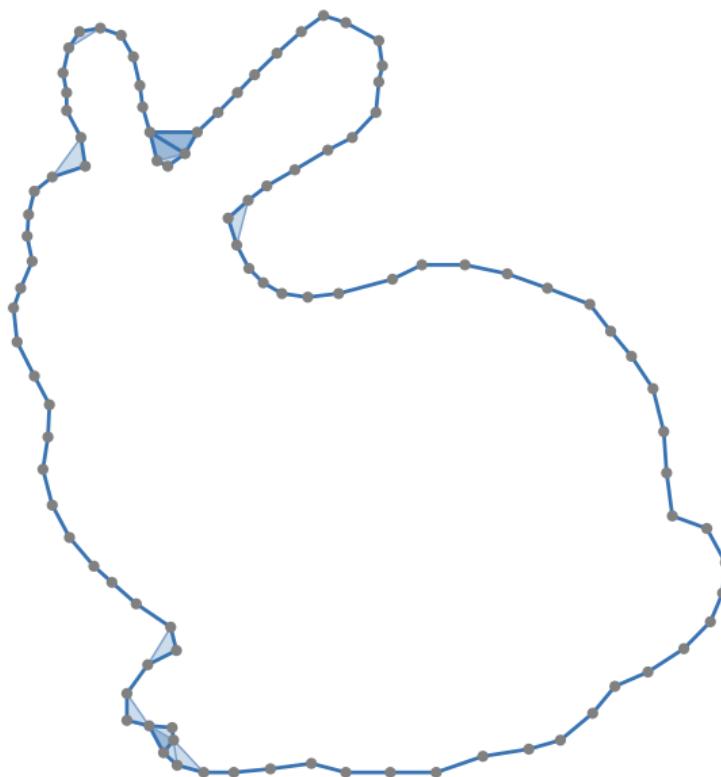
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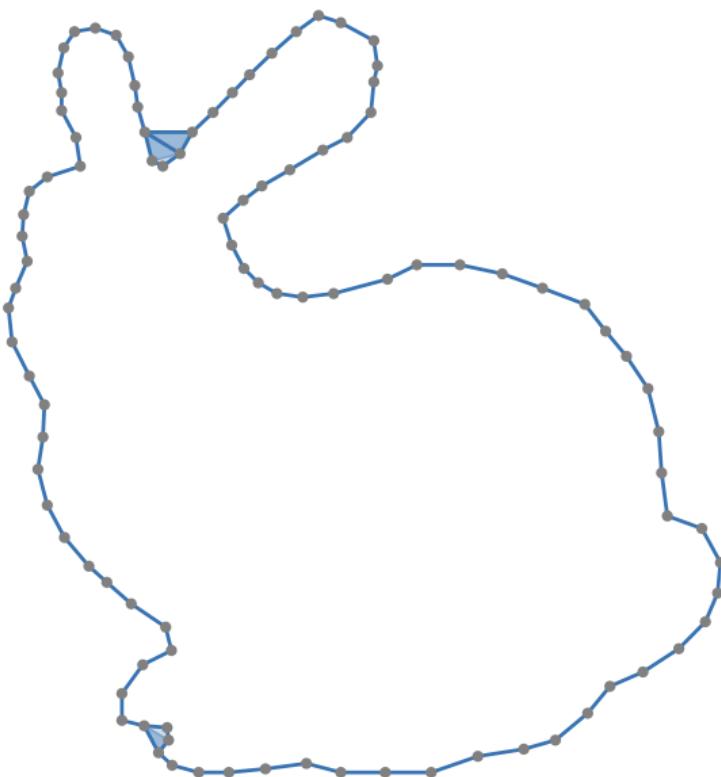
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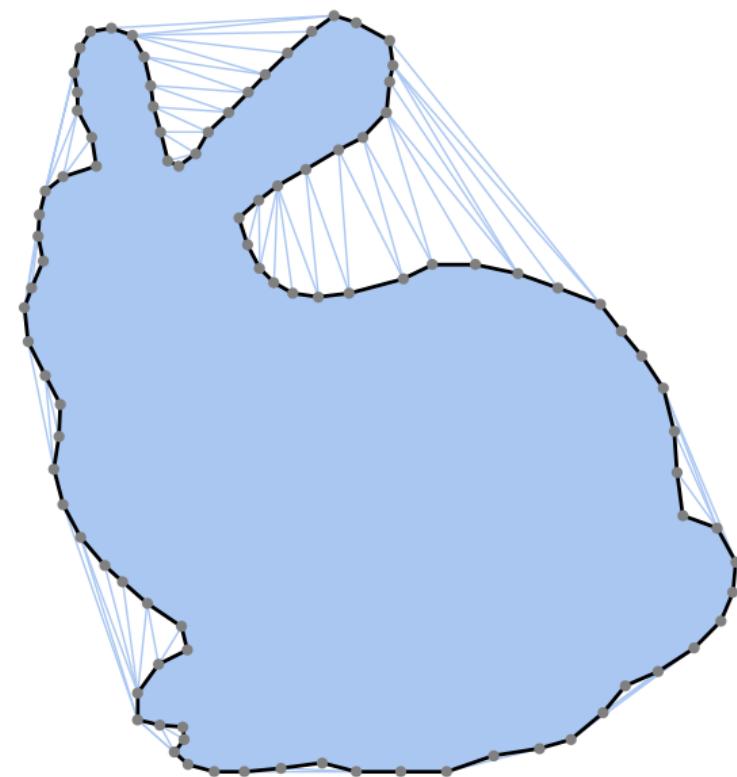
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Algebraic gradient flows and persistent homology

Loose ends in the literature:

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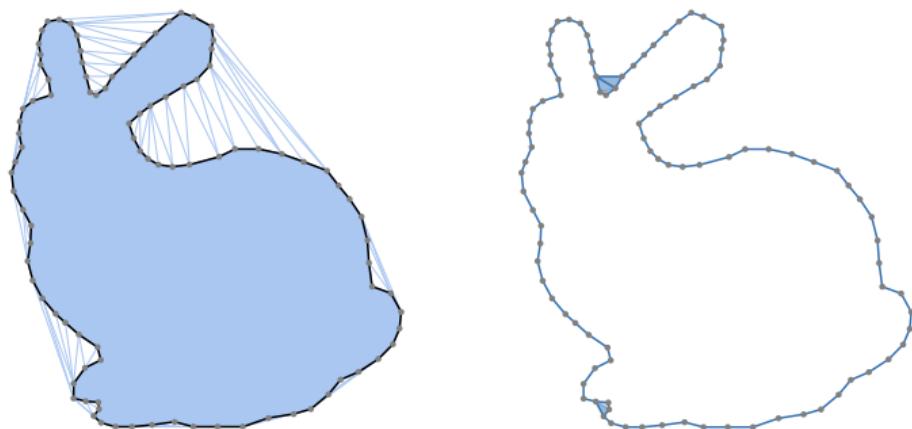
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- The resulting representative cycles are lexicographically minimal

Minimal cycles and Wrap complexes

Theorem (B, Roll 2022)

Let $X \subset \mathbb{R}$ be a finite subset in general position and let $r \in \mathbb{R}$. Then the lexicographically minimal cycles of $\text{Del}_r(X)$ are supported on $\text{Wrap}_r(X)$.



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Definition

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- The basis elements Σ_* take the role of the simplices in discrete Morse theory
- All further definitions of discrete Morse theory apply verbatim

Algebraic gradients from persistent homology

Let $R = D \cdot V$ be the result of exhaustive reduction of the boundary matrix D (for the Delaunay filtration, breaking ties lexicographically).

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- The apparent pairs of persistence zero refine the gradient of the Delaunay radius function.

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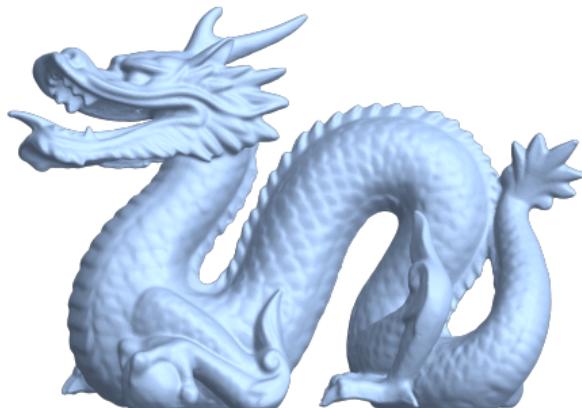
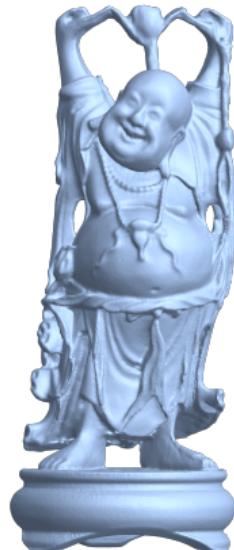
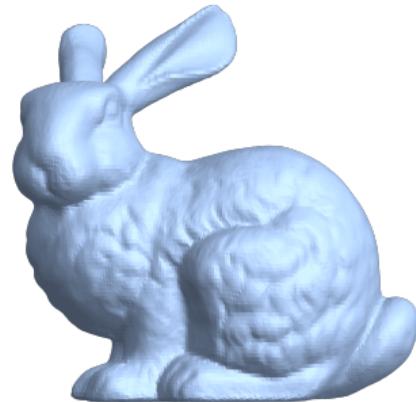
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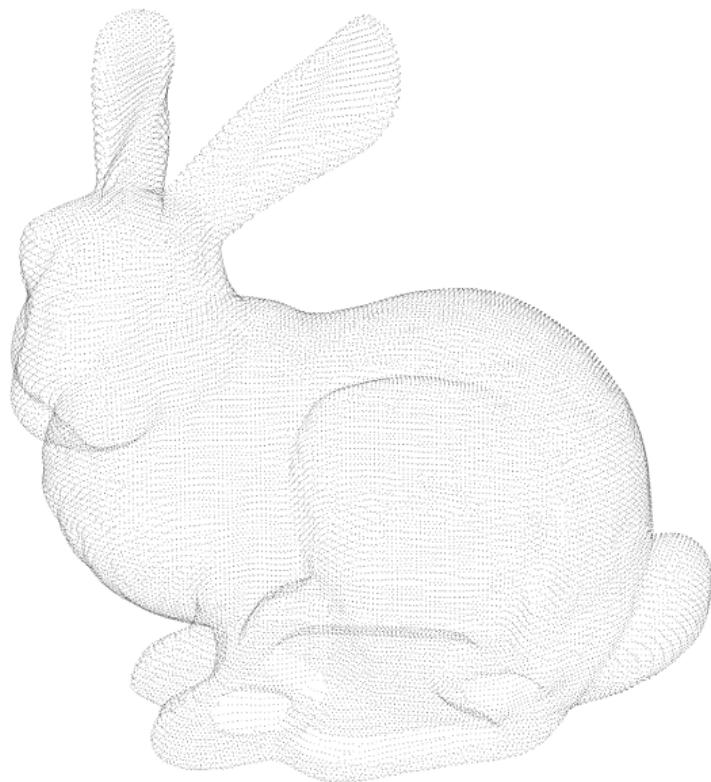
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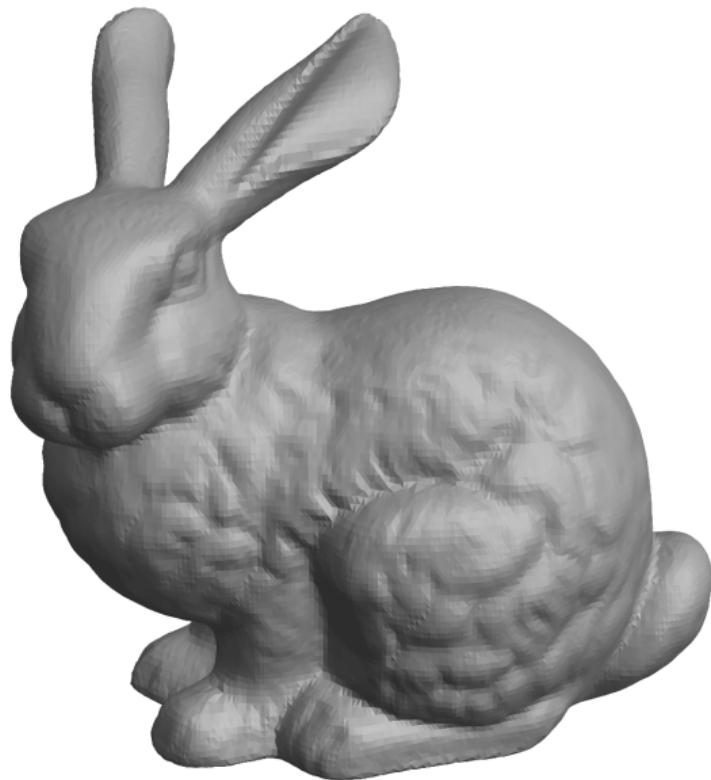
Point cloud reconstruction with lexicographically minimal cycles



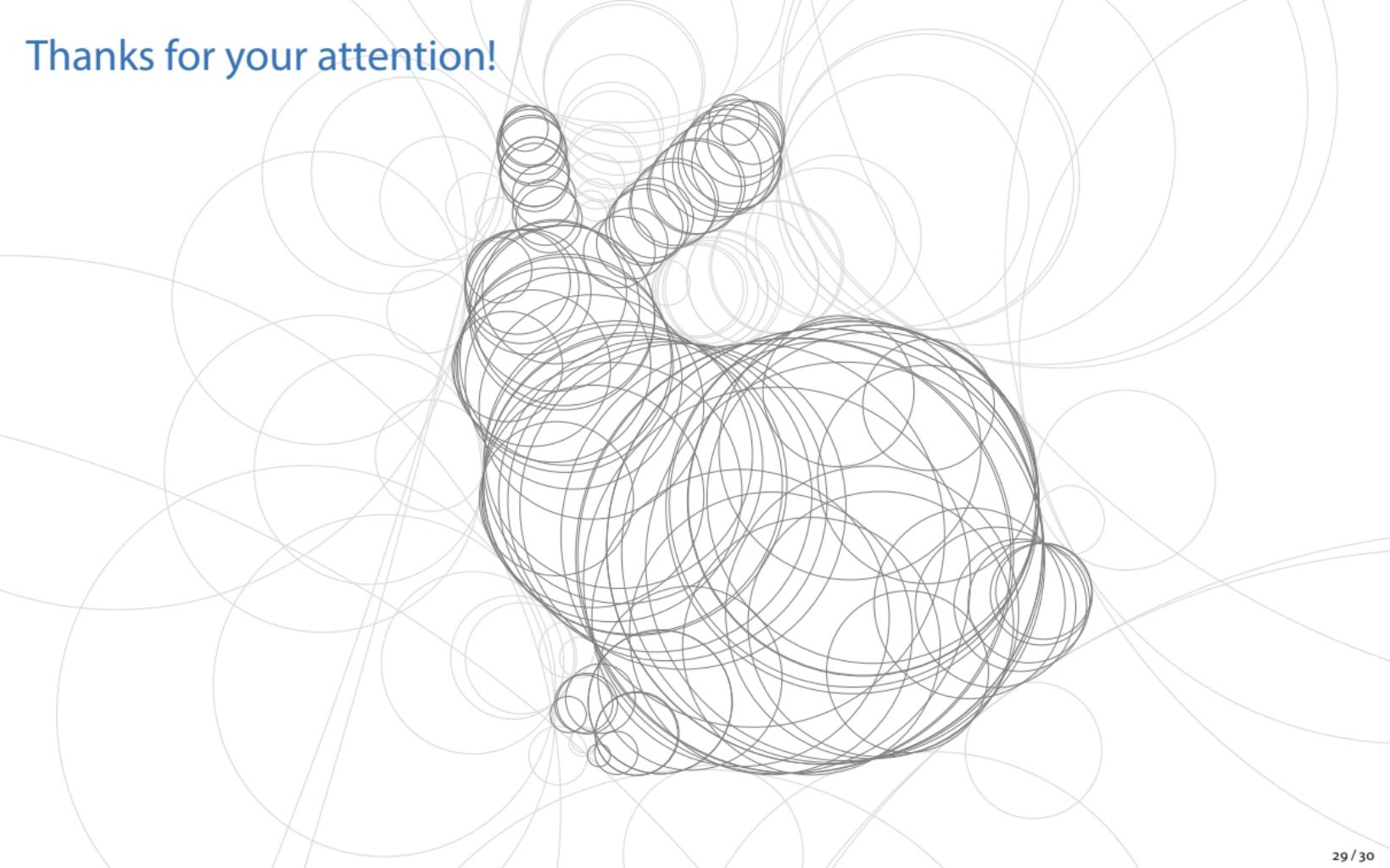
Computing minimal cycles with exhaustive reduction



Computing minimal cycles with exhaustive reduction



Thanks for your attention!



Further reading



[U. Bauer, F. Roll](#)

Gromov hyperbolicity, geodesic defect, and apparent pairs in Vietoris-Rips filtrations
Symposium on Computational Geometry, 2022. doi:10.4230/LIPIcs.SoCG.2022.15



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Connecting Discrete Morse Theory and Persistence: Wrap Complexes and Lexicographic Optimal Cycles
Preprint, 2022. arXiv:2212.02345



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