Assignment 1: Introduction to Systems Programming

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1 Mathematical Analysis

1.1 Algorithm 1

[H] Integer array A of size N Greatest Sum of Subarray $i \leftarrow 0$ N $j \leftarrow i$ N $s \leftarrow 0$ k $\leftarrow i$ j $s \leftarrow s + A[k]$ s ξ maxmax $\leftarrow s$ Pseudocode for Basic Enumeration

1.2 Algorithm 2

[H] Integer array A of size N Greatest Sum of Subarray $i \leftarrow 0$ N $s \leftarrow 0$ j $\leftarrow i$ N $s \leftarrow A[j]$ s \downarrow maxmax $\leftarrow s$ Pseudocode for Better Enumeration

2 Theoretical Correctness

3 Testing

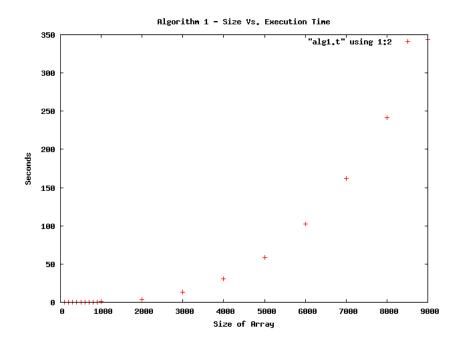
931678074 - 5703

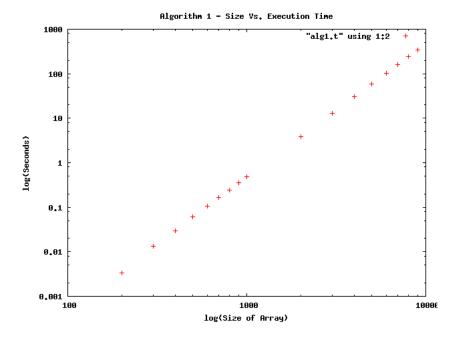
930569466 - 8184

932086449 - 4949

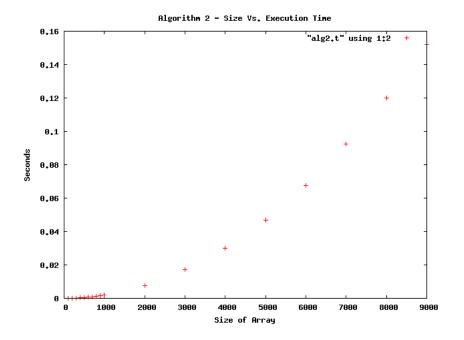
4 Experimental Analysis

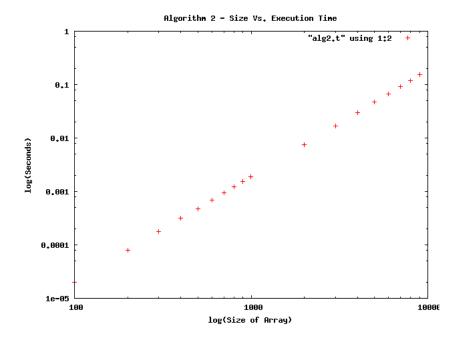
4.1 Algorithm 1



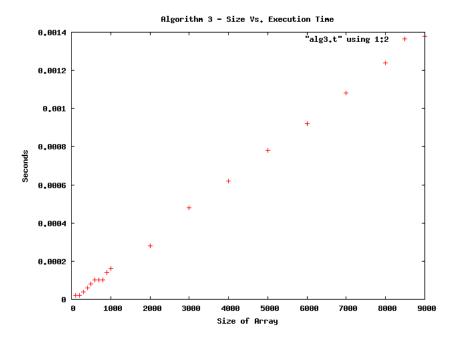


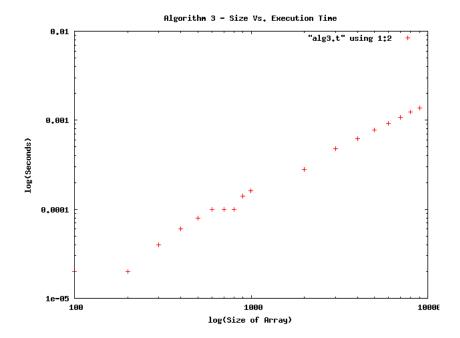
4.2 Algorithm 2





4.3 Algorithm 3





5 Extrapolation and Interpretation

6 Code

6.1 Algorithm 1

```
* Enumeration
    * Loop over each pair of indices i\,;\,j and compute the sum from k{=}i
           to j of a[k].
    * Keep the best sum you have found so far.
   using namespace std;
   int MaxSubarray(int a[], int n){
10
      \begin{array}{ll} \mbox{int} & i \;, j \;, k \,; \\ \mbox{int} & \max \; = \; a \; [ \; 0 \; ] \;; \end{array}
12
13
      int sum;
      for (i = 0; i < n; ++i)
15
        for (j = i; j < n; ++j){
sum = 0;
16
           for (k = i; k <=j; ++k) {
  sum += a[k];
18
19
           if(max < sum)
21
22
              \max = \sup;
23
        }
24
      }
25
26
      return max;
```

alg1.cpp

6.2 Algorithm 2

```
* Better Enumeration
   st Notice that in the previous algorithm, the same sum is computed
        many times.
   * In particular, notice that sum from k=i to j of a[k] can be
        computed from sum from k=i to j-1 of a[k] in O(1) time,
   rather than starting from scratch.

* Write a new version of the frst algorithm that takes advantage
        of this observation.
  using namespace std;
11
  int MaxSubarray(int a[], int n){
12
13
     int i, j, k;
14
     int max = a[0];
15
16
     int sum;
     for (i = 0; i < n; ++i)
17
      sum = 0;
       for (j = i; j < n; ++j){
    sum += a[j];
19
20
21
         if(max < sum){
           \max = \sup;
22
23
       }
     }
25
26
     {\tt return}\ \max;
27
```

 ${\rm alg2.cpp}$

6.3 Algorithm 3

```
* Divide and Conquer
    st If we split the array into two halves, we know that the maximum
         subarray will either be
             * contained entirely in the frst half,
             \ast contained entirely in the second half, or
6
             * made of a suffix of the frst half of maximum sum and a
          prefix of the second half of maximum sum
    * The frst two cases can be found recursively. The last case can
         be found in linear time.
   using namespace std;
11
   void MaxSubarray_h(int a[], int n, int b[]) {
12
     if(n <= 1){
13
        b[0] = a[0];
14
        b[1] = a[0];
        b[2] = a[0];

b[3] = a[0];
16
17
        return;
18
19
      int i = n/2;
20
     \begin{array}{lll} \operatorname{int} & *\operatorname{left} &= & \operatorname{new} & \operatorname{int} \left[ \, 4 \, \right]; \end{array}
21
      MaxSubarray_h(a,i,left);
22
      int *right = new int[4];
23
     MaxSubarray_h(a+i,n-i,right);
24
     int sum = left[0] + right[0];
26
     int l = left[0] + right[1];
27
      int r = left[2] + right[0];
28
     int mid = left[2] + right[1];
29
30
     l = l > left[1] ? l : left[1];
31
     r = r > right[2] ? r : right[2];
32
     int m = left[3] > right[3]? left[3] : right[3];
33
     m = m > mid ? m : mid;
34
35
36
     b[0] = sum;
     b[1] = 1;
37
     b[2] = r;
38
     b[3] = m;
39
40
42
   int MaxSubarray(int a[], int n){
     int *p = new int [4];
     \begin{aligned} &\text{MaxSubarray.h(a,n,p)};\\ &\text{int } s1 = p[0] > p[1] ? p[0] : p[1];\\ &\text{int } s2 = p[2] > p[3] ? p[2] : p[3];\\ &s1 = s1 > s2 ? s1 : s2; \end{aligned}
45
46
47
48
49
      return s1;
   }
50
```

alg3.cpp