

Assignment 1: Introduction to Systems Programming

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1 Mathematical Analysis

1.1 Algorithm 1

Data: Integer array A of size N

Result: Greatest Sum of Subarray

```
1 for i ← 0 to N do
2   for j ← i to N do
3     s ← 0
4     for k ← i to j do
5       s ← s + A[k]
6     end
7     if s > max then
8       max ← s
9     end
10  end
11 end
```

Algorithm 1: Pseudocode for Basic Enumeration

1.2 Algorithm 2

Data: Integer array A of size N

Result: Greatest Sum of Subarray

```
1 for i ← 0 to N do
2   s ← 0
3   for j ← i to N do
4     s ← A[j]
5     if s > max then
6       max ← s
7     end
8   end
9 end
```

Algorithm 2: Pseudocode for Better Enumeration

1.3 Algorithm 3

Data: Integer array A of size N

Result: Greatest Sum of Subarray

Algorithm 3: Pseudocode for Divide and Conquer

2 Theoretical Correctness

3 Testing

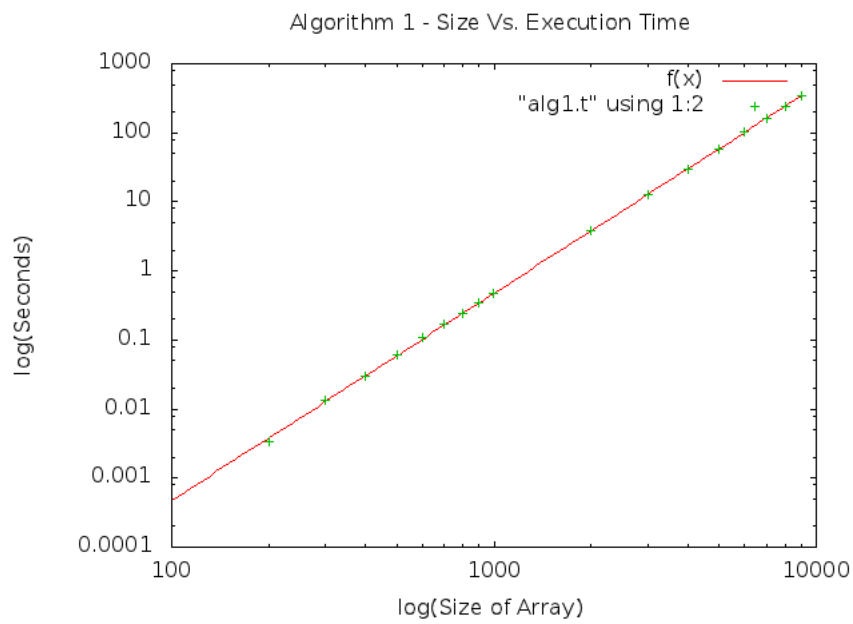
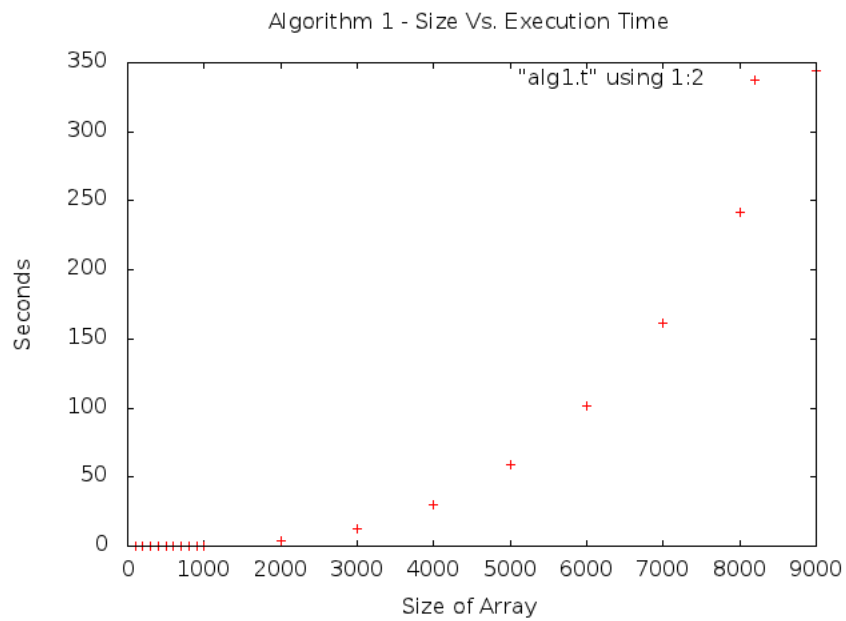
931678074 - 5703

930569466 - 8184

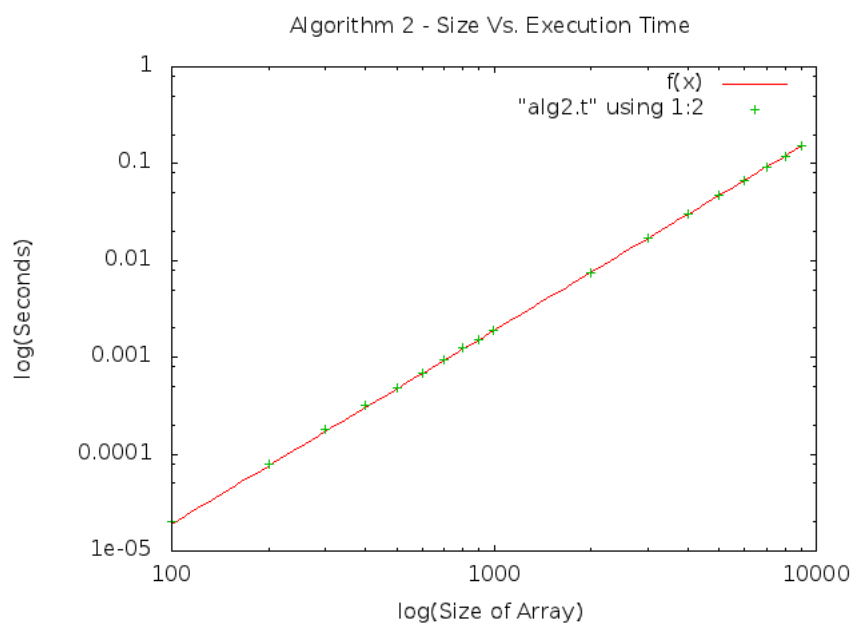
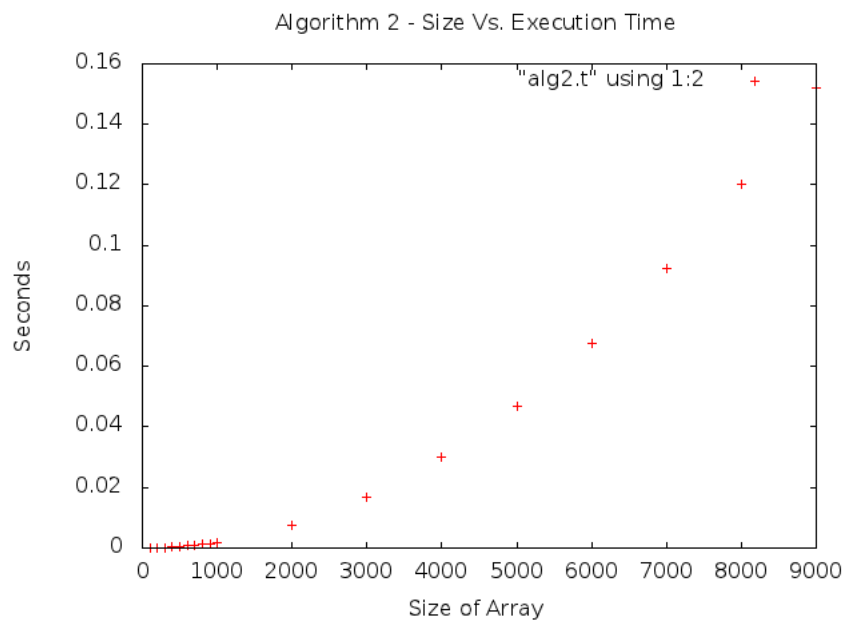
932086449 - 4949

4 Experimental Analysis

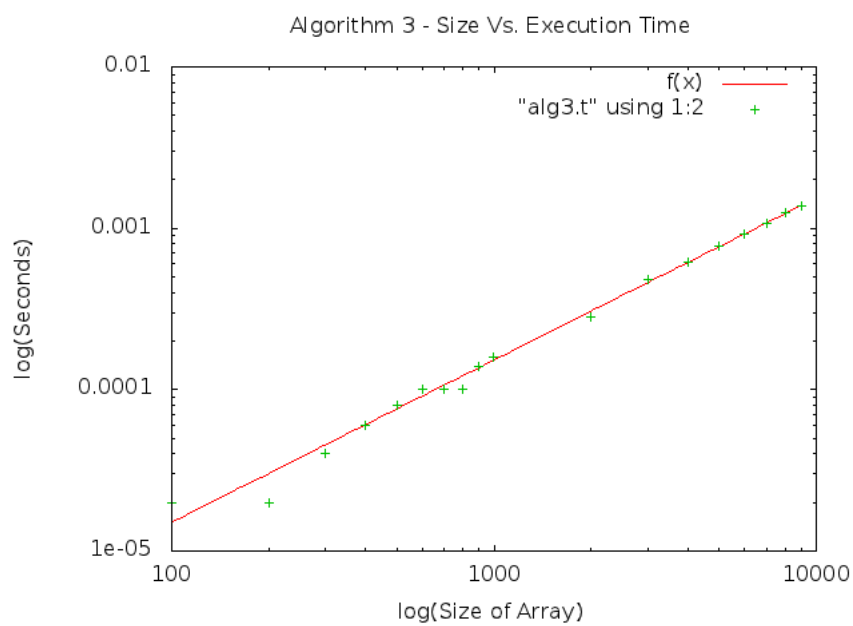
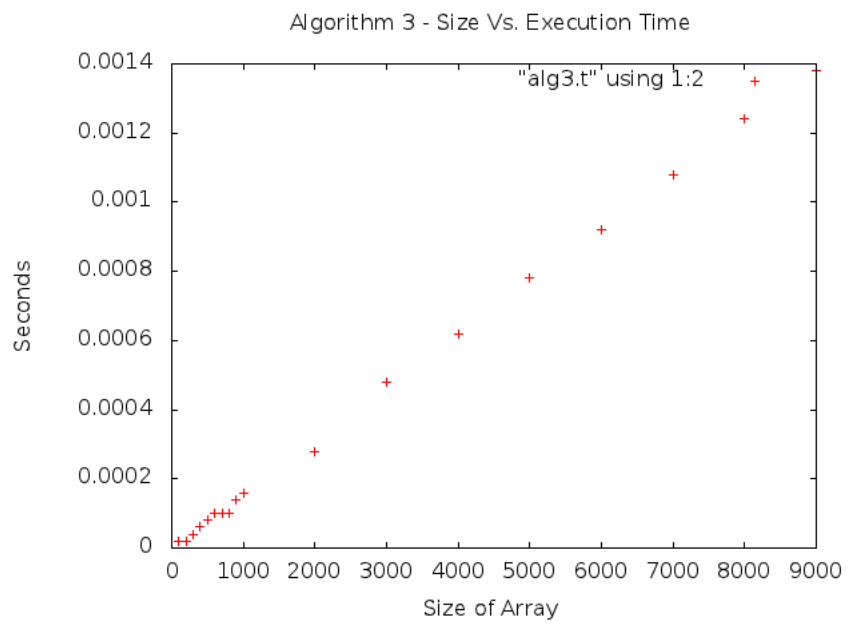
4.1 Algorithm 1



4.2 Algorithm 2



4.3 Algorithm 3



5 Extrapolation and Interpretation

5.1 Extrapolation

5.2 Interpretation

5.3 Algorithm 1

5.3.1 Extrapolation

$$f(n) = 4.71599 \times 10^{-10} \times n^3$$

$$f(n) = 3600 \rightarrow n = 19690$$

5.3.2 Interpretation

$$\text{Slope} = 2.99734$$

5.4 Algorithm 2

5.4.1 Extrapolation

$$f(n) = 1.87761 \times 10^{-9} \times n^2$$

$$f(n) = 3600 \rightarrow n = 1384678$$

5.4.2 Interpretation

$$\text{Slope} = 1.99602$$

5.5 Algorithm 3

5.5.1 Extrapolation

$$f(n) = 1.74832 \times 10^{-8} \times n \times \log(n)$$

$$f(n) = 3600 \rightarrow n = 8984428998 = 8.98 \times 10^9$$

5.5.2 Interpretation

$$\text{Slope} = 1.00506$$

6 Code

6.1 Files

alg1.cpp - function for algorithm 1
alg2.cpp - function for algorithm 1
alg3.cpp - function for algorithm 1
analysis.cpp - code to run algorithm and measure times for the number of array then outputs .t file
makefile - to compile files
maxSubarray.pdf -
maxSubarray.tex - to create pdf filename
test.cpp - allows input of file, and runs algorithm on input file
analysis/
test/
timingfiles/ - holds files for creating plots
timingfiles/*.t - files that holds run times for different array sizes
timingfiles/*.gp - code for gnuplot. 2 plots of each algorithm: 1 normal plot, and 1 log-log plot

6.2 Algorithm 1

```
1  /*
2  * Enumeration
3  * Loop over each pair of indices i; j and compute the sum from k=i to j of a[k].
4  * Keep the best sum you have found so far.
5  */
6
7  using namespace std;
8
9
10 int MaxSubarray(int a[], int n){
11
12     int i, j, k;
13     int max = a[0];
14     int sum;
15     for (i = 0; i < n; ++i) {
16         for (j = i; j < n; ++j) {
17             sum = 0;
18             for (k = i; k <= j; ++k) {
19                 sum += a[k];
20             }
21             if (max < sum) {
22                 max = sum;
23             }
24         }
25     }
26     return max;
27 }
```

alg1.cpp

6.3 Algorithm 2

```
1  /*
2  * Better Enumeration
3  * Notice that in the previous algorithm, the same sum is computed many times.
4  * In particular, notice that sum from k=i to j of a[k] can be computed from sum from k=
5  * i to j - 1 of a[k] in O(1) time, rather than starting from scratch.
6  * Write a new version of the first algorithm that takes advantage of this observation.
7  */
8
9  using namespace std;
10
11
12  int MaxSubarray(int a[], int n){
13
14      int i, j, k;
15      int max = a[0];
16      int sum;
17      for (i = 0; i < n; ++i) {
18          sum = 0;
19          for (j = i; j < n; ++j) {
20              sum += a[j];
21              if (max < sum) {
22                  max = sum;
23              }
24          }
25      }
26      return max;
27  }
```

alg2.cpp

6.4 Algorithm 3

```
1  /*
2  * Divide and Conquer
3  * If we split the array into two halves, we know that the maximum subarray will either
4  *   be
5  *   * contained entirely in the first half,
6  *   * contained entirely in the second half, or
7  *   * made of a suffix of the first half of maximum sum and a prefix of the second
8  *     half of maximum sum
9  * The first two cases can be found recursively. The last case can be found in linear
10 * time.
11 */
12 using namespace std;
13 void MaxSubarray_h(int a[], int n, int b[]) {
14     if(n <= 1){
15         b[0] = a[0];
16         b[1] = a[0];
17         b[2] = a[0];
18         b[3] = a[0];
19         return;
20     }
21     int i = n/2;
22     int *left = new int[4];
23     MaxSubarray_h(a, i, left);
24     int *right = new int[4];
25     MaxSubarray_h(a+i, n-i, right);
26
27     int sum = left[0] + right[0];
28     int l = left[0]+right[1];
29     int r = left[2]+right[0];
30     int mid = left[2]+right[1];
31
32     l = l > left[1] ? l : left[1];
33     r = r > right[2] ? r : right[2];
34     int m = left[3] > right[3] ? left[3] : right[3];
35     m = m > mid ? m : mid;
36
37     b[0] = sum;
38     b[1] = l;
39     b[2] = r;
40     b[3] = m;
41 }
42
43 int MaxSubarray(int a[], int n){
44     int *p = new int[4];
45     MaxSubarray_h(a, n, p);
46     int s1 = p[0] > p[1] ? p[0] : p[1];
47     int s2 = p[2] > p[3] ? p[2] : p[3];
48     s1 = s1 > s2 ? s1 : s2;
49     return s1;
50 }
```

alg3.cpp