Assignment 2: Dynamic Programming project

Francis Vo, Soo-Hyun Yoo

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1 Recursive function

```cpp
struct maxS {
  int current;    // Current sum
  int max;        // Overall max
}

struct maxS MaxSubarray(int array, int size) {
  struct maxS ms;

  if (size == 1) {    // Base case
    ms.current = 0;
    ms.max = array[0];
  } else {    // Recurse on array excluding last element.
    ms = MaxSubarray(array, size-1);
  }

  // Find maximum.
  ms.current += array[size-1];
  ms.current = (ms.current > 0) ? ms.current : 0;
  ms.max = (ms.current > ms.max) ? ms.current : 0;

  return ms;
}
```
rec.cpp

Where `maxS` is a struct holding the running sum and the overall maximum sum and `MaxSubarray` is the recursive function. For an array A of size n, we find the maximum subarray with `MaxSubarray(A, n)`.
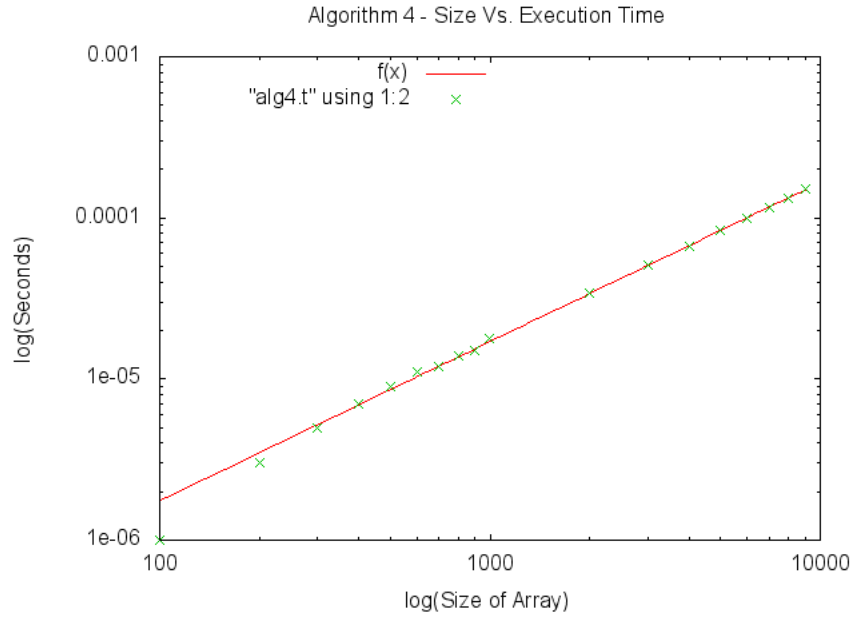
2 Pseudocode

3 Running time

Using the Log-Log plot will give us a good hint for the asymptotic run times. The slope for the graph is 0.987956, meaning that the asymptotic run times is around $\Omega(n)$. Then looking at the code, it only looks at each element once, making it $\Omega(n)$

4 Theoretical correctness

Induction Proof. $MS(k)$ will return the maximum subarray sum for the array $A[0:k]$
Base case: If $n = -1$ then $max = current = 0$
Inductive Step: $maxSubarray(n-1).current + A[n]$ or 0 is the current largest sum starting from the left
Proof:
Case if $A[n] > 0$ then $current = MS(n-1).current + A[n] > MS(n-1).current$.
        This number might also be the max value. So $max = Greater(max, current)$
Case if $A[n] > -maxSubarray(n-1)$ then $maxSubarray(n-1) + A[n] < 0$ making the Null set greater.
        $max = MS(n-1).max$ and $current = 0$
Case else making $A[n]$ negative but $maxSubarray(n-1) + A[n] >= 0$
        so it is still good to use for the next current: $current + A[n+1] > A[n+1]$
        $max = MS(n-1).max$ and $current = maxSubarray(n-1) + A[n]$
$MS(n).max = max$ and $MS(n).current = current$

$\square$

5 Implement

5.1 Algorithm 4

```cpp
/*
 * Enumeration
 * Loop over each pair of indices i; j and compute the sum from k=i to j of a[k].
 * Keep the best sum you have found so far.
 */

using namespace std;


int MaxSubarray(int a[], int n){
    int current = 0;
    int max = 0;
    int i;
    for(i = 0;i < n; i++){
        current += a[i];
        if(current <= 0){
            current = 0;
        }else if(current > max){
            max = current;
        }
    }
    return max;
}
```

alg4.cpp

6 Test

Test were run on the ms_test.txt file given last project and large arrays given by student ids.

7 Compare

Well, there is a huge difference as seen on the Compare Plot. Algorithm 4, Dynamic Programming, is great because doesn't use any recurvsive calls and doesn't need to hold much data. Algorithm 4 only needs to hold onto 2 integers (max and current) and the input integer array. Whereas algorithm 3, divide & conquer, needs to use memory on the stack for each recurvsive call and needs to pass 4 integers back to the parent function.
WHAT THE HELL IS GOOD ABOUT D&C???

Algorithm 3 and 4 Compare