

WARREN D. MACEVOY

# LUCK

PERSONAL NOTES

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# Contents

<i>Introduction</i>	9
<i>Normal Distribution</i>	15
<i>Computation</i>	21
<i>Proofs</i>	27



# List of Figures

- 1 Luck to get  $x$  heads on 8 fair coins. After the probabilities are arranged in decreasing order on the unit interval,  $L(x)$  is the center point of equally probable outcomes. 11
- 2 Exact (blue, equation 22) vs approximate (red, equation 23) luck for 1-d normal distribution. 13
- 3 Exact (blue) vs approximate (red) luck for normal distribution for  $n = 1, 10$ , and 100. 17
- 4 Histograms of 10,000 luck values in the same distributions as table 2. Upshot:  $y$  values in the  $x$  distribution are viewed as extremely lucky and vice-versa, while they are have uniform luck in their own respective distributions. 18
- 5 Scilab listing to compute the natural log of the probability of multivariate normal outcomes.  $x$  is a  $n_{\text{dim}} \times n_{\text{samps}}$  of outcomes,  $\mu$  is a  $n_{\text{dim}} \times 1$  column vector of the means, and  $\Sigma$  is a  $n_{\text{dim}} \times n_{\text{dim}}$  covariance matrix. The result  $\ln p$  is a  $1 \times n_{\text{samps}}$  row vector. 18
- 6 Scilab listing to efficiently compute the luck of multivariate normal outcomes.  $x$  is a  $n_{\text{dim}} \times n_{\text{samps}}$  of outcomes,  $\mu$  is a  $n_{\text{dim}} \times 1$  column vector of the means, and  $\Sigma$  is a  $n_{\text{dim}} \times n_{\text{dim}}$  covariance matrix. The result  $L$  is a  $1 \times n_{\text{samps}}$  row vector of the luck associated with each outcome. 19
- 7 Like `mnluck` in figure 6, but approximates luck via (39). 19
- 8 Scilab function to compute the natural log of the probability of multinomial outcomes.  $p$  is a  $n_{\text{probs}} \times 1$  column vector of category probabilities,  $x$  is a  $n_{\text{probs}} \times n_{\text{samps}}$  matrix of outcomes, and the result  $\ln p$  is a  $1 \times n_{\text{samps}}$  row vector. 22
- 9 Scilab listing to get a sample of outcomes from a multinomial distribution.  $n_{\text{samps}}$  is the number of desired samples,  $n_{\text{trials}}$  is the number of trials in each sample, and  $p$  is a  $n_{\text{probs}} \times 1$  column vector of category probabilities. The result  $x$  is a  $n_{\text{probs}} \times n_{\text{samps}}$  matrix of sample outcomes, where  $\sum_{j=1}^{n_{\text{probs}}} x(j, i) = n_{\text{trials}}$ . 22

- 10 Recursively compute luck of multinomial exactly using exhaustive  
sum.  $x$  is a  $n_{\text{probs}} \times n_{\text{samps}}$  matrix of outcomes, and  $p$  is a  $n_{\text{probs}} \times$   
1 column vector of category probabilities. The result is a  $1 \times n_{\text{samps}}$   
of luck values. 23
- 11 Given the log of the probabilities of a set of sample data, return a ta-  
ble used for quickly estimating luck. `problns` is a  $1 \times n_{\text{nsamps}}$  row  
vector of logs of probabilities, and `eps` is an optional parameter giv-  
ing the absolute error (in log space) for considering two probabili-  
ties to be equal. Returns `setup`, a  $2 \times N$  matrix giving  $-\log p(x)$  and  
estimates for  $L(x)$  for each unique probability in the sample. This 24  
function is useful generally (not just for multinomial distributions).
- 12 Estimate luck given log of probabilities and `setup`. `problns` is a row  
vector of log-probabilities, `setup` is the setup from a (possibly differ-  
ent) sample. 24
- 13 Use the above functions for estimating luck numerically (`nluck`) and  
estimating the error in luck (numerical standard deviation). The last  
lines optionally compute the exact values of luck (`luck`) and standard  
deviations to compare with. One run produced `nluck=0.63025` and  
`luck=0.62875` with error bound estimates `nsd=0.004827`, `sd=0.0048314`  
and `z=0.3105`. 25
- 14 Histogram of `z` error values for 10,000 numerical approximations of  
luck. The maximum value of `sd` was `max(sd)=0.005`. 25

## *List of Tables*

- 1 This is arranged in increasing luck (which is decreasing probability). Getting exactly  $x = 4$  heads is unlucky, requiring only  $L = 14\%$  luck, while getting  $x = 0$  or  $x = 8$  heads is almost 100% luck. 11
- 2 Luck from two randomly generated distributions  $\mu^{(x)}$  and  $\mu^{(y)}$  uniformly chosen in  $[0, 1]^{100}$ , and  $\Sigma^{(x)}, \Sigma^{(y)}$  are transposed squares of random  $100 \times 100$  matrices. In each row,  $x$  is a sample from the  $\mu^{(x)}, \Sigma^{(x)}$ , normal distribution, and  $y$  is from the  $\mu^{(y)}, \Sigma^{(y)}$  distribution. The actual values of  $x$  and  $y$  are not given, since they are very large (100 numbers each) and uninteresting. 17





# Introduction

The point of these notes is to introduce an idea of “luck” that connects mathematical probability with the everyday notion.

As a motivating problem, imagine walking along the beach and asking a random person to toss a tennis ball so that it lands in the sand. The probability that it lands at some point would depend on the habits of the thrower and the details of the beach, but we can summarize this as some probability distribution,  $p(x) = \rho(x)dx$ , where  $x \in \mathbb{R}^2$  is a suitable coordinate system for the beach in question. It would almost certainly not be a uniform distribution, and it would almost certainly not be particularly concentrated.

Traditional probability feels uncomfortable here. The chances of the ball landing at a given point is zero, and so miraculous. Yet anyone watching this process would only occasionally be surprised by the outcome.

As common (and mundane, not miraculous) such situations are, the language of statistics seems to have difficulty with the notion. Nor is it limited to continuous cases, just when there are a lot of possible outcomes. Such examples lead to non-zero but very small probabilities.

To distinguish from the more general notion of luck, note that that there is no extrinsic value on an outcome. To say something is “lucky” often means there is some value (different from the probability) associated with outcomes. However, outcomes that are the most valuable are often the least probable, and outcomes of equal probability ought to be equally lucky. In the most extreme case of all equally probable outcomes (uniform probability), every outcome should have a luck of  $\frac{1}{2}$ .

These observations lead to the following definition of luck:

The luck  $L(x)$  of an outcome  $x$  is the probability of getting any outcome  $y$  that is more probable than  $x$ , plus one-half the probability of getting any outcome  $y$  that is equally probable to  $x$ .

From the perspective of discussing luck, it is convenient to have few sets:  $\Omega(x)$ , the outcomes more likely than  $x$ , and  $\omega(x)$ , the outcomes equally likely to  $x$ .

These might easily exist, perhaps as another name, in the literature; I just don't know what it is called.

For a continuous probability distribution such as this, the chance of the ball landing in some small area  $dx$  near  $x$  is  $p(x) = \rho(x)dx$ . But the ball lands at a point, so  $dx$  is zero, so the probability  $p(x) = \rho(x)dx$  is zero.

The real motivation of this came from the space of passwords a person might choose from, which is an effectively infinite discrete space.

**Definition 1.** Omega.  $\Omega(x)$  is set of outcomes more likely than  $x$ :

$$\Omega(x) = \{y \mid p(y) > p(x)\} . \quad (1)$$

We define  $|\Omega(x)|$  as the probability an outcome is in  $\Omega(x)$ ,  $|\Omega(x)| = P(y \in \Omega(x))$ . In the discrete case, this is

$$|\Omega(x)| = \sum_{y \in \Omega(x)} p(y), \quad (2)$$

and, in the continuous case,

$$|\Omega(x)| = \int_{\Omega(x)} \rho(y) dy . \quad (3)$$

**Definition 2.** omega.  $\omega(x)$  is the set of outcomes equally likely to  $x$ :

$$\omega(x) = \{y \mid p(y) = p(x)\} . \quad (4)$$

Similar to  $\Omega(x)$ , we define  $|\omega(x)|$  as the probability an outcome is in  $\omega(x)$ ,  $|\omega(x)| = P(y \in \omega(x))$ .

In the discrete case, this is

$$|\omega(x)| = \sum_{y \in \omega(x)} p(y), \quad (5)$$

and, in the continuous case,

$$|\omega(x)| = \int_{\omega(x)} \rho(y) dy . \quad (6)$$

With these definitions in place, we define luck mathematically as follows:

**Definition 3.** Luck. The luck of an outcome is the probability getting any more likely outcome, plus half the probability of getting any equally likely outcome:

$$L(x) = |\Omega(x)| + \frac{1}{2}|\omega(x)| . \quad (7)$$

*Properties of Luck.*

- Range of luck.  $0 \leq L(x) \leq 1$ . This ranges from no luck to perfect luck.
- Lucky outcomes. If  $L(x)$  is close to 1, then  $p(x)$  is comparatively small, and most outcomes would have a higher probability (you are lucky).
- Unlucky outcomes. If  $L(x)$  is close to 0, then  $p(x)$  is comparatively large, and most outcomes would have a lower probability (you are unlucky).

For the typical case of many outcomes with different probabilities,  $|\omega(x)|$  is small. For example,  $|\omega(x)| = 0$  for any multivariate normal distribution.

- Luck on average.  $E(L) = \frac{1}{2}$ . On average, luck is always 50:50.

We are interested in cases which have many possible outcomes with low but somewhat different probabilities (like the tennis ball on the beach). If the space is well divided (so  $\max |\omega| = \max_x |\omega(x)|$  is small), then there are other interesting properties of luck:

- $E(f(L)) = \int_0^1 f(L) dL - \varepsilon$ , where  $|\varepsilon| \leq \max |f''| \cdot \max |\omega|^2 / 24$ .
- For  $p \geq 1$ ,  $E(L^p) = 1/(p+1) - \varepsilon$ ,  $0 \leq \varepsilon \leq p \cdot (p-1) \max |\omega|^2 / 24$ .
- For  $0 \leq a \leq b \leq 1$ ,  $E(L \in [a, b]) = b - a - \varepsilon$ ,  $|\varepsilon| \leq \max |\omega|$ .

The proofs of these come from the midpoint integration rule and thinking about the general case for figure 1 below.

**Example 1. Coins.** Suppose we toss 8 fair coins. The probability of getting exactly  $x$  heads out of 8 tosses is given by the binomial distribution

$$p(x) = \frac{8!}{x!(8-x)!} \left(\frac{1}{2}\right)^8. \quad (8)$$

What is the luck associated with this distribution?

$x$	$p(x)$	$\Omega(x)$	$ \Omega(x) $	$\omega(x)$	$ \omega(x) $	$L(x)$
4	0.2734	{}	0.0000	{4}	0.2734	0.1367
3 or 5	0.2188	{4}	0.2734	{3,5}	0.4375	0.4922
2 or 6	0.1094	{3,4,5}	0.7109	{2,6}	0.2188	0.8203
1 or 7	0.0313	{2,3,4,5,6}	0.9297	{1,7}	0.0625	0.9609
0 or 8	0.0039	{1,2,3,4,5,6,7}	0.9922	{0,8}	0.0078	0.9961

In particular  $|\omega(x)| = 0$  for the normal, exponential, beta, and gamma distributions. As a worst-case counterexample, the flattest distribution is the uniform distribution, for which  $|\omega(x)| = 1$ .

Table 1: This is arranged in increasing luck (which is decreasing probability). Getting exactly  $x = 4$  heads is unlucky, requiring only  $L = 14\%$  luck, while getting  $x = 0$  or  $x = 8$  heads is almost 100% luck.

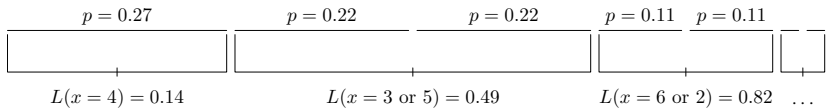


Figure 1: Luck to get  $x$  heads on 8 fair coins. After the probabilities are arranged in decreasing order on the unit interval,  $L(x)$  is the center point of equally probable outcomes.

Luck on average is  $\frac{1}{2}$ :

$$E(L) = \sum_{x=0}^8 p(x) \cdot L(x) = \frac{1}{2}. \quad (9)$$

The second moment should be close to  $\frac{1}{3}$ :

$$E(L^2) = \sum_{x=0}^8 p(x) \cdot L(x)^2 = \frac{1}{3} - 0.0096 \quad (10)$$

The probability luck is in the middle half is about  $\frac{1}{2}$ :

For any distribution,  $E(L) = \frac{1}{2}$ .

For any distribution,  $E(L^2) = \frac{1}{3} - \varepsilon$ , with  $0 \leq \varepsilon \leq \max |\omega|^2 / 12$ .

For the coin distribution, the bound is  $0 \leq \varepsilon \leq 0.016$ , and the actual error  $\varepsilon = 0.0096$ .

For any distribution,  $E(L \in [a, b]) = b - a + \varepsilon$ , with  $|\varepsilon| \leq \max |\omega|$ .

For the coin distribution and  $[a, b] = [\frac{1}{4}, \frac{3}{4}]$ ,  $E(L \in [a, b]) = \frac{1}{2} + \varepsilon$  with  $|\varepsilon| \leq 0.4375$  as the (poor) error bound, and the actual error of  $\varepsilon = -0.0625$

$$\begin{aligned}
E(L \in [\frac{1}{4}, \frac{3}{4}]) &= \sum_{x=0}^8 p(x) \cdot \left\{ \begin{array}{ll} 1 & \text{if } L(x) \in [\frac{1}{4}, \frac{3}{4}] \\ 0 & \text{otherwise} \end{array} \right\} \\
&= \frac{1}{2} - 0.0625.
\end{aligned} \tag{11}$$

**Example 2.** *Normal.* This is a special case of the multivariate normal we cover in the next section, but working out the details for the one-dimensional case can be illuminating. We define the one-dimensional normal distribution with mean  $\mu$  and variance  $\Sigma$  as

$$P_{\text{normal}}(x; \mu, \Sigma) = \frac{e^{-\frac{(x-\mu)^2}{2\Sigma}}}{\sqrt{2\pi\Sigma}}, \tag{12}$$

where

$$\mu = E(x), \tag{13}$$

and

$$\Sigma = E((x - \mu)^2). \tag{14}$$

First note that  $\Omega(x)$  is the open the interval between  $x$  and  $x$  reflected around  $\mu$ :

$$\Omega(x) = (\min(x, 2\mu - x), \max(x, 2\mu - x)), \tag{15}$$

and  $\omega(x)$  is the endpoints of that interval:

$$\omega(x) = \{x, 2\mu - x\}. \tag{16}$$

Since the 1-dimensional normal distribution is a continuous distribution and  $\omega(x)$  is a finite set,  $|\omega(x)| = 0$ , i.e.,

$$|\omega(x)| = \int_{\omega(x)} P_{\text{normal}}(y; \mu, \Sigma) dy = 0. \tag{17}$$

This means all the luck properties are exact (the error terms are zero), and the  $\frac{1}{2}|\omega(x)|$  contributes nothing to the luck of an outcome.

What remains is to calculate luck,

$$L(x) = \int_{\Omega(x)} P_{\text{normal}}(y; \mu, \Sigma) dy. \tag{18}$$

Changing variables to the normalized z-score:  $z = \sqrt{\Sigma^{-1}}(x - \mu)$ , this can be rewritten as

$$L(x) = \int_{-R}^R P_{\text{normal}}(y, 0, 1) dy, \tag{19}$$

where

$$R = |\sqrt{\Sigma^{-1}}(x - \mu)|. \tag{20}$$

Using  $\text{erf}(x)$ , defined as

$\text{erf}(x)$  is a normalized integral of the  $P_{\text{normal}}(x; \mu = 0, \Sigma = 1/2)$  so that  $\text{erf}(0) = 0$  and  $\text{erf}(\pm\infty) = \pm 1$ .

$$\operatorname{erf}(x) = \frac{2}{\sqrt{\pi}} \int_0^x e^{-y^2} dy, \quad (21)$$

the luck of an outcome can be written as

$$L(x) = \operatorname{erf} \left| \frac{x - \mu}{\sqrt{2\Sigma}} \right|. \quad (22)$$

In the next chapter, where we address the more general multivariate normal case, we obtain the approximation,

$$L(x) \approx \frac{1}{2} \left[ 1 + \operatorname{erf} \left( \left| \frac{x - \mu}{\sqrt{\Sigma}} \right| - \sqrt{\frac{1}{2}} \right) \right]. \quad (23)$$

Figure 2 compares the exact and approximate result in the 1-d case.

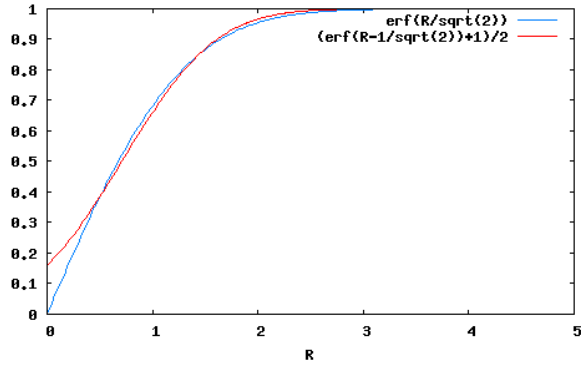


Figure 2: Exact (blue, equation 22) vs approximate (red, equation 23) luck for 1-d normal distribution.



# Normal Distribution

Suppose we are in a probability space well approximated by the multivariate normal (Gaussian) distribution of a random variable  $x \in \mathbb{R}^n$  with mean  $\mu$  and non-singular covariance  $\Sigma$ :

The  $n = 1$  example in the introduction is just where where  $\mu$  and  $\Sigma$  are scalars.

$$P_{\text{normal}}(x; \mu, \Sigma) = \frac{e^{-\frac{1}{2}(x-\mu)^T \Sigma^{-1}(x-\mu)}}{\sqrt{(2\pi)^n \det \Sigma}}, \quad (24)$$

where

$$\mu_i = E(x_i), \quad (25)$$

and

$$\Sigma_{ij} = E((x_i - \mu_i)(x_j - \mu_j)). \quad (26)$$

How lucky is some outcome  $x$ ? From the definition:

$$L(x) = |\Omega(x)| + \frac{1}{2}|\omega(x)|, \quad (27)$$

where

$$\Omega(x) = \{y \in \mathbb{R}^n | P_{\text{normal}}(y) > P_{\text{normal}}(x)\} \quad (28)$$

$$= \left\{y \in \mathbb{R}^n \left| |\sqrt{\Sigma^{-1}}(y - \mu)| < |\sqrt{\Sigma^{-1}}(x - \mu)| \right. \right\} \quad (29)$$

and

$$\omega(x) = \left\{y \in \mathbb{R}^n \left| |\sqrt{\Sigma^{-1}}(y - \mu)| = |\sqrt{\Sigma^{-1}}(x - \mu)| \right. \right\}. \quad (30)$$

Because  $\omega(x)$  has no volume in  $\mathbb{R}^n$ ,

$$|\omega(x)| = \int_{\omega(x)} P_{\text{normal}}(y; \mu, \Sigma) dy = 0. \quad (31)$$

So

$$L(x) = |\Omega(x)| \quad (32)$$

$$= \int_{\Omega(x)} P_{\text{normal}}(y; \mu, \Sigma) dy. \quad (33)$$

By changing variables to  $z = \sqrt{\Sigma^{-1}}(x - \mu)$ ,

$$L(x) = \int_{|z| < R} P_{\text{normal}}(z; 0, I) dz, \quad (34)$$

where  $R = |\sqrt{\Sigma^{-1}}(x - \mu)|$ .

This can be evaluated in spherical coordinates:

$$L(x) = \frac{1}{\sqrt{(2\pi)^n}} \int_0^R \frac{n\pi^{n/2}}{\Gamma(\frac{n}{2} + 1)} r^{n-1} e^{-\frac{1}{2}r^2} dr, \quad (35)$$

$$= \frac{\gamma(n/2, R^2/2)}{\Gamma(n/2)}. \quad (36)$$

The last form uses the lower incomplete gamma function and gamma function, defined to be

$$\gamma(s, x) = \int_0^x t^{s-1} e^{-t} dt \quad (37)$$

$$\Gamma(s) = \gamma(s, \infty). \quad (38)$$

For any value of  $n$ , but particularly for large values, we find the following approximation to be very good:

$$L(x) \approx \frac{1}{2} \left[ 1 + \text{erf}(|\sqrt{\Sigma^{-1}}(x - \mu)| - \sqrt{n - 1/2}) \right]. \quad (39)$$

Not only is this result pretty, it is very useful. Suppose we have distribution parameters  $\mu$  and  $\Sigma$ , and would like to know if they fit actual observations. A traditional approach requires a large sample to estimate  $\mu$  and  $\Sigma$ , but we don't need this, nor do we need to assume that the distribution is normal. We just need to ask if the observations are surprising (lucky or unlucky). In large dimensions, numerical experiments suggest one sample is in most cases sufficient to establish practical certainty (probability of error less than  $10^{-15}$ ).

### Combining Normal Luck

The approximate result (39) leads to a rule for combining (normal) luck: Suppose there are two independent normal distributions parameterized by  $\mu^{(x)}, \Sigma^{(x)}$  of dimension  $n_x$ , and  $\mu^{(y)}, \Sigma^{(y)}$  of dimension  $n_y$ . What is the luck of a single combined observation  $(x, y)$ ?

$$L(x, y) \approx \frac{1}{2} \left[ 1 + \text{erf}(\sqrt{R_x(x)^2 + R_y(y)^2} - \sqrt{n_x + n_y - \frac{1}{2}}) \right], \quad (40)$$

$\Sigma$  is symmetric and positive definite, and so is its inverse. The square-root can be computed as a Cholesky decomposition. In Scilab,  $z = \text{chol}(\Sigma)^{\backslash}(x - \mu)$

$R = \text{norm}(z)$

$L = \text{cdfgam}("PQ", R^2/2, n/2, 1)$

This comes from a Taylor expansion of the log of the integrand in (35), which is specifically invalid for  $n = 1$ , hence the difference between the general case and the  $n = 1$  case (22).

$L_{\text{approx}} = 0.5 * (1 + \text{erf}(R - \sqrt{n - 1/2}))$





Figure 3: Exact (blue) vs approximate (red) luck for normal distribution for  $n = 1, 10$ , and  $100$ .

$L^{(x)}(x)$	$L^{(y)}(x)$	$L^{(x)}(y)$	$L^{(y)}(y)$
0.501 417 202 0	1.000 000 000 0	1.000 000 000 0	0.838 180 264 1
0.731 421 266 5	1.000 000 000 0	1.000 000 000 0	0.239 258 143 2
0.982 563 033 9	1.000 000 000 0	1.000 000 000 0	0.271 695 512 7
0.033 455 080 7	1.000 000 000 0	1.000 000 000 0	0.421 320 625 9
0.689 429 934 0	1.000 000 000 0	1.000 000 000 0	0.074 461 655 7
0.736 397 593 7	1.000 000 000 0	1.000 000 000 0	0.294 050 728 4
0.304 521 296 7	1.000 000 000 0	1.000 000 000 0	0.707 849 014 7
0.231 111 574 4	1.000 000 000 0	1.000 000 000 0	0.290 313 093 2
0.585 247 719 9	1.000 000 000 0	1.000 000 000 0	0.636 902 202 8
0.214 552 926 1	1.000 000 000 0	1.000 000 000 0	0.268 989 787 4

Table 2: Luck from two randomly generated distributions  $\mu^{(x)}$  and  $\mu^{(y)}$  uniformly chosen in  $[0, 1]^{100}$ , and  $\Sigma^{(x)}, \Sigma^{(y)}$  are transposed squares of random  $100 \times 100$  matrices. In each row,  $x$  is a sample from the  $\mu^{(x)}, \Sigma^{(x)}$ , normal distribution, and  $y$  is from the  $\mu^{(y)}, \Sigma^{(y)}$  distribution. The actual values of  $x$  and  $y$  are not given, since they are very large (100 numbers each) and uninteresting.

where

$$R_x(x) = \text{erf}^{-1}(2L_x - 1) - \sqrt{n_x - \frac{1}{2}} = \sqrt{(\Sigma^{(x)})^{-1}} \left( x - \mu^{(x)} \right), \quad (41)$$

$$R_y(y) = \text{erf}^{-1}(2L_y - 1) - \sqrt{n_y - \frac{1}{2}} = \sqrt{(\Sigma^{(y)})^{-1}} \left( y - \mu^{(y)} \right). \quad (42)$$

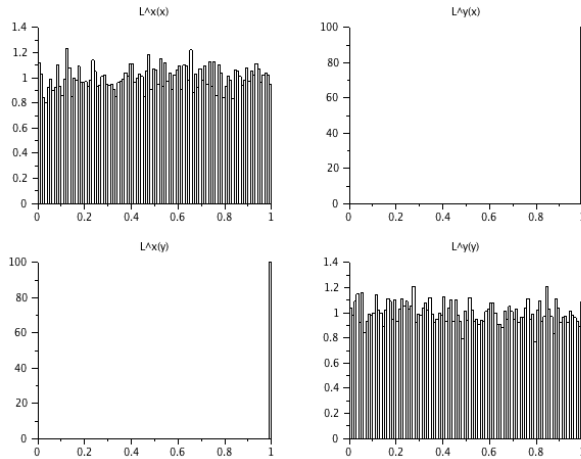


Figure 4: Histograms of 10,000 luck values in the same distributions as table 2. Upshot:  $y$  values in the  $x$  distribution are viewed as extremely lucky and vice-versa, while they are have uniform luck in their own respective distributions.

The listings in figures 5-7 give Scilab functions for basic luck calculations for multivariate normal distributions.

```
function lnp=mnprobln(x,mu,Sigma)
    [dim,nsamps]=size(x);
    sigma=chol(Sigma)';
    z=sigma \ (x-mu*ones(1,nsamps));
    R2=sum(z.^2,'r');
    lnp=-R2/2-(dim/2*log(2*%pi)+...
        sum(log(diag(sigma))));
endfunction
```

Figure 5: Scilab listing to compute the natural log of the probability of multivariate normal outcomes.  $x$  is a  $n_{\text{dim}} \times n_{\text{samps}}$  of outcomes,  $\mu$  is a  $n_{\text{dim}} \times 1$  column vector of the means, and  $\Sigma$  is a  $n_{\text{dim}} \times n_{\text{dim}}$  covariance matrix. The result `lnp` is a  $1 \times n_{\text{samps}}$  row vector.

```

function L=mnluck(x,mu,Sigma)
    [dim,nsamps]=size(x);
    one=ones(1,nsamps);
    sigma=chol(Sigma)';
    z=sigma\x-mu*one;
    R2=sum(z.^2,'r');
    L=cdfgam("PQ",R2/2,dim/2*one,one);
endfunction

```

Figure 6: Scilab listing to efficiently compute the luck of multivariate normal outcomes.  $x$  is a  $n_{\text{dim}} \times n_{\text{samps}}$  of outcomes,  $\mu$  is a  $n_{\text{dim}} \times 1$  column vector of the means, and  $\Sigma$  is a  $n_{\text{dim}} \times n_{\text{dim}}$  covariance matrix. The result  $L$  is a  $1 \times n_{\text{samps}}$  row vector of the luck associated with each outcome.

```

function L=mnapproxluck(x,mu,Sigma)
    [dim,nsamps]=size(x);
    sigma=chol(Sigma)';
    z=sigma\x-mu*ones(1,nsamps);
    R=sqrt(sum(z.^2,'r'));
    L=0.5*(1+erf(R-sqrt(dim-0.5)));
endfunction

```

Figure 7: Like mnluck in figure 6, but approximates luck via (39).



## Computation

Unfortunately, the explicit computations for luck may be impractical. However, if  $n$  independent samples are taken  $S = (x_1, \dots, x_n)$ , then it can be estimated as:

$$\ell(x) = \frac{1}{n} \{\# \text{ of outcomes in } S \text{ more probable than } x\} \\ + \frac{1}{2n} \{\# \text{ of outcomes in } S \text{ equally probable to } x\}$$

It is a reasonably straightforward calculation to show that

$$E(\ell(x)) = L(x), \quad (43)$$

and

$$E((\ell(x) - L(x))^2) \leq \frac{1}{n} L(x) \cdot (1 - L(x)) \leq \frac{1}{4n}. \quad (44)$$

**Example 3. Multinomial.** We are aware of no computationally efficient way to exactly compute the luck for a multinomial distribution other than the explicit sum. Suppose there is a questionnaire with 4 possible answers with category probabilities 0.1, 0.2, 0.3, and 0.4. How lucky would it be to get 13, 15, 27 and 45 responses of each respective answer in 100 samples?

The log-probability in this case is the multinomial probability (see figure 8):

$$y = \text{mulprobln}(x, p) = \log \left[ \left( \sum_{i=1}^{n_p} x_i \right)! \prod_{i=1}^{n_p} \frac{p_i^{x_i}}{x_i!} \right]. \quad (45)$$

Numerical estimates of luck requires samples from the given probability distribution. Figure 9 uses a built-in function in scilab to generate such a sample set.

For smaller spaces, luck for a multinomial distribution can be computed explicitly. Figure 10 computes this as an inefficient testing reference.

```

function lnp=mulprobln(x,p)
    [nprobs,nsamps]=size(x);
    pln=log(p);
    y=zeros(1,nsamps);
    for i=1:nsamps
        xi=x(:,i);
        ni=sum(xi);
        lnp(i)=gammaln(ni+1)+...
            sum(xi.*pln-gammaln(xi+1));
    end
endfunction

```

Figure 8: Scilab function to compute the natural log of the probability of multinomial outcomes.  $p$  is a  $n_{\text{probs}} \times 1$  column vector of category probabilities,  $x$  is a  $n_{\text{probs}} \times n_{\text{samps}}$  matrix of outcomes, and the result  $\text{lnp}$  is a  $1 \times n_{\text{samps}}$  row vector.

```

function x=mulsamp(nsamps,ntrials,p)
    x=grand(nsamps,"mul",ntrials,p(1:(length(p)-1)));
endfunction

```

Figure 9: Scilab listing to get a sample of outcomes from a multinomial distribution.  $n_{\text{samps}}$  is the number of desired samples,  $n_{\text{trials}}$  is the number of trials in each sample, and  $p$  is a  $n_{\text{probs}} \times 1$  column vector of category probabilities. The result  $x$  is a  $n_{\text{probs}} \times n_{\text{samps}}$  matrix of sample outcomes, where  $\sum_{j=1}^{n_{\text{probs}}} x(j,i) = n_{\text{trials}}$ .

```

// Used by mulluck below to
// recursively compute luck
function el=mulluckrec(nb,mb,...
    lpx,k,p,y,eps)
    [n,m]=size(lpx);
    el=zeros(n,m);
    for yk=0:nb-mb
        y(k)=yk;
        if k < length(p)-1 then
            el=el+mulluckrec(nb,mb+yk,...
                lpx,k+1,p,y,eps);
        else
            y(length(p))=nb-mb-yk;
            lpy=mulprobln(y,p);
            c=0.5*bool2s(lpy > lpx-eps)+...
                0.5*bool2s(lpy > lpx+eps);
            el = el + c .* exp(lpy);
        end
    end
endfunction

function L=mulluck(x,p,eps)
    if ~exists("eps","local") then
        eps=sqrt(%eps);
    end
    [nprobs,nsamps]=size(x);
    ntrials=sum(x,'r');
    no=min(ntrials);
    n1=max(ntrials);
    assert_checkequal(no,n1);
    L=mulluckrec(no,0,mulprobln(x,p),...
        1,p,zeros(nprobs,1),eps);
endfunction

```

Figure 10: Recursively compute luck of multinomial exactly using exhaustive sum.  $x$  is a  $n_{\text{probs}} \times n_{\text{samps}}$  matrix of outcomes, and  $p$  is a  $n_{\text{probs}} \times 1$  column vector of category probabilities. The result is a  $1 \times n_{\text{samps}}$  of luck values.

```

function setup=numlucksetup(problns,eps)
    if ~exists("eps","local") then
        eps=sqrt(%eps);
    end

    n=length(problns);
    problns=gsort(problns);

    for pass=1:2
        if pass == 2 then
            setup=zeros(2,count);
        end
        i=1;
        count=0;
        while i<=n
            j=i;
            while j <= n-1 & ...
                problns(i)-problns(j+1) < eps
                j=j+1;
            end
            count=count+1;
            if pass == 2 then
                setup(1,count)= ...
                    -0.5*(problns(i)+problns(j));
                setup(2,count)= ...
                    (i-1)/n+0.5*(j-i+1)/n;
            end
            i=j+1;
        end
    end
endfunction

```

Figure 11: Given the log of the probabilities of a set of sample data, return a table used for quickly estimating luck. `problns` is a  $1 \times n_{\text{samps}}$  row vector of logs of probabilities, and `eps` is an optional parameter giving the absolute error (in log space) for considering two probabilities to be equal. Returns `setup`, a  $2 \times N$  matrix giving  $-\log p(x)$  and estimates for  $L(x)$  for each unique probability in the sample. This function is useful generally (not just for multinomial distributions).

```

function L=numluck(problns,setup)
    L=max(0,min(1,interpnl(setup,-problns)));
endfunction

```

Figure 12: Estimate luck given log of probabilities and setup. `problns` is a row vector of log-probabilities, `setup` is the setup from a (possibly different) sample.



```

// get sample set
nsamps=10000;
ntrials=100;
p=[0.1;0.2;0.3;0.4];
x=mulsamp(nsamps,ntrials,p);
problns=mulprobln(x,p);

// setup for luck estimates
setup=numlucksetup(problns);

// estimate luck numerically
xo=[13; 15; 27; 45];
problnso=mulprobln(xo,p);
nluck=numluck(problnso,setup);
nsd=sqrt(nluck .* (1-nluck) ./ nsamps);

// optional exact luck
luck=mulluck(xo,p);
sd=sqrt(luck .* (1-luck) ./ nsamps);

// z is approximately normally distributed
z=(nluck-luck) ./ sd;

```

Figure 13: Use the above functions for estimating luck numerically (nluck) and estimating the error in luck (numerical standard deviation). The last lines optionally compute the exact values of luck (luck) and standard deviations to compare with. One run produced  $nluck=0.63025$  and  $luck=0.62875$  with error bound estimates  $nsd=0.004827$ ,  $sd=0.0048314$  and  $z=0.3105$ .

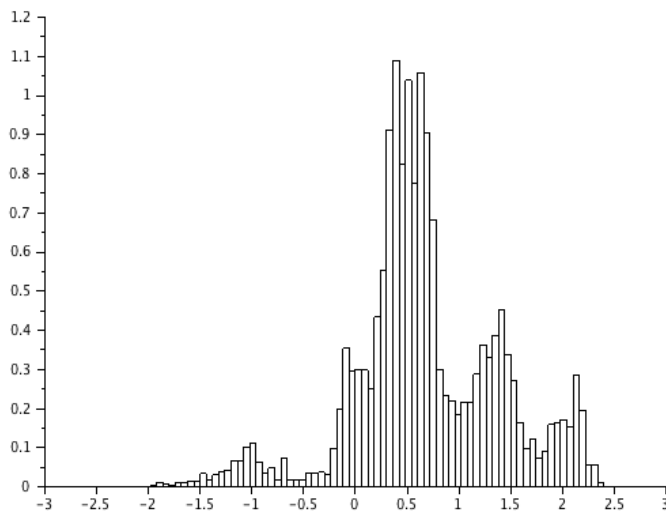


Figure 14: Histogram of z error values for 10,000 numerical approximations of luck. The maximum value of sd was  $\max(sd)=0.005$ .



# Proofs

*Luck in general.*

Fundamentally, defining luck depends on a partial ordering of the elements of the probability space in question. In a finite or countable space, this can be done just from the probability. In larger spaces more topological structure is required:

$$y < x \text{ iff } \lim_{r \downarrow 0} \frac{p(B_r(x))}{p(B_r(y))} < 1. \quad (46)$$

Here,  $p(S)$  denotes the probability of an outcome in  $S \subseteq X$ , and  $B_r(x) \subseteq X$  is the open ball of radius  $r$  in the topology of  $X$ .

Using this ordering, luck can be generalized as follows:

$$\Omega(x) = \{y \mid y < x\}, \quad (47)$$

$$|\Omega(x)| = p(\Omega(x)), \quad (48)$$

$$\omega(x) = \{y \mid y \not< x \text{ and } x \not< y\}, \quad (49)$$

$$|\omega(x)| = p(\omega(x)), \quad (50)$$

and

$$L(x) = |\Omega(x)| + \frac{1}{2}|\omega(x)|. \quad (51)$$

*From the introduction*

**Theorem 1.** *Range of Luck. For any probability space,*

$$0 \leq L(x) \leq 1. \quad (52)$$

*Proof.* From the definition,  $L(x) = |\Omega(x)| + \frac{1}{2}|\omega(x)|$ , which is clearly non-negative, and  $\Omega(x)$  and  $\omega(x)$  are disjoint subsets of the probability space. Here the  $|\cdot|$  notation is the measure of these sets in the probability space, and their union is at most the whole space, so  $L(x) \leq |\Omega(x) \cup \omega(x)| \leq 1$ .

**Theorem 2.** *Lucky values. If  $L(x)$  is close to 1, then  $p(x)$  is relatively small and most outcomes have a higher probability (you are lucky).*

*Proof.* Since  $L(x)$  is close to 1, write  $L(x) = 1 - \varepsilon(x)$ , where  $\varepsilon(x)$  is a small non-negative number.

First, since  $\Omega(x)$  and  $\omega(x)$  are disjoint subsets of the space of outcomes  $X$ ,  $|\Omega(x)| + |\omega(x)| \leq 1$ , or

$$|\Omega(x)| \leq 1 - |\omega(x)|.$$

Second,  $L(x) = |\Omega(x)| + \frac{1}{2}|\omega(x)|$ , which by the first inequality can be bounded as  $1 - \varepsilon(x) \leq 1 - |\omega(x)| + \frac{1}{2}|\omega(x)|$ , which can be rearranged as

$$|\omega(x)| \leq 2\varepsilon(x).$$

Third,  $|\Omega(x)| + \frac{1}{2}|\omega(x)| = 1 - \varepsilon(x)$ , or  $|\Omega(x)| = 1 - \varepsilon(x) - \frac{1}{2}|\omega(x)|$ , which by the second inequality,

$$|\Omega(x)| \geq 1 - 2\varepsilon(x). \quad (53)$$

Thus at least a  $1 - 2\varepsilon(x)$  fraction of the probability space have a higher probability of occurring.

**Theorem 3.** *Unlucky values. If  $L(x)$  is close to 0, then  $p(x)$  is comparatively large, and most outcomes would have a lower probability (you are unlucky).*

*Proof.* Since  $L(x)$  is close to 0, write  $L(x) = \varepsilon(x)$ , where  $\varepsilon(x)$  is a small non-negative number.

First,  $L(x) = |\Omega(x)| + \frac{1}{2}|\omega(x)|$ , so

$$|\omega(x)| \leq 2\varepsilon(x).$$

Second,  $|\Omega(x)| + |\omega(x)| = |\Omega(x)| + \frac{1}{2}|\omega(x)| + \frac{1}{2}|\omega(x)| \leq 2\varepsilon(x)$ , so

$$|\Omega(x)| + |\omega(x)| \leq 2\varepsilon(x).$$

Third,  $|X - [\Omega(x) \cup \omega(x)]| \geq 1 - |\Omega(x) \cup \omega(x)| \geq 1 - 2\varepsilon(x)$ , or

$$|X - [\Omega(x) \cup \omega(x)]| \geq 1 - 2\varepsilon(x).$$

Thus at least a  $1 - 2\varepsilon(x)$  fraction of the probability space have a lower probability of occurring.

**Theorem 4.** *On average, luck is always 50:50.*

$$E(L) = \frac{1}{2}.$$

*Proof.* This is an application of the next theorem where  $f(L) = L$ .

**Theorem 5.** *Smooth uniformity - finite space X.*

$E(f(L)) = \int_0^1 f(L) dL - \varepsilon$ , where  $|\varepsilon| \leq \max |f''| \cdot \max |\omega|^2 / 24$ .

Since the definition of luck only depends on the probabilities of outcomes, it is natural to consider the set of equivalence classes  $[x]$  from  $[X]$  defined by equal probabilities:  $[x] = \{y \mid p(y) = p(x)\}$ .

We also use the midpoint integration estimate:

$$\int_a^b f(L) dL = (b-a)f\left(\frac{a+b}{2}\right) + \frac{(b-a)^3}{24} f''(\xi), \text{ where } a < \xi < b.$$

$$\int_0^1 f(L) dL = \sum_{[x]} \int_{L([x]) - \frac{1}{2}|\omega([x])|}^{L([x]) + \frac{1}{2}|\omega([x])|} f(L) dL, \quad (54)$$

$$= \sum_{[x]} \left\{ |\omega([x])| f(L([x])) + \frac{|\omega([x])|^3}{24} f''(\xi_{[x]}) \right\}, \quad (55)$$

$$= \sum_x p(x) f(L(x)) - \varepsilon, \quad (56)$$

$$= E(f(L)) - \varepsilon, \quad (57)$$

$$\varepsilon = \sum_{[x]} \frac{|\omega([x])|^3}{24} f''(\xi_{[x]}), \quad (58)$$

$$|\varepsilon| \leq \frac{1}{24} \max_x |\omega(x)|^2 \cdot \max_{0 \leq L \leq 1} |f''(L)|. \quad (59)$$

**Theorem 6.** *Moments - finite space X.*

For  $p \geq 1$ ,  $E(L^p) = 1/(p+1) - \varepsilon$ ,  $0 \leq \varepsilon \leq p \cdot (p-1) \max |\omega|^2 / 24$ .

*Proof.* This is an example of  $f(L) = L^p$  with  $p \geq 1$  in the theorem above, and noting that  $f''(L)$  is non-negative, so  $\varepsilon$  in (58) must be non-negative.

**Theorem 7.** *Interval uniformity - finite space X.*

For  $0 \leq a \leq b \leq 1$ ,  $E(L \in [a, b]) = b - a - \varepsilon$ ,  $|\varepsilon| \leq \max |\omega|$ .

*Proof.* Let  $f(L)$  be the characteristic function for the closed interval  $[a, b]$ .

$$b - a = \int_0^1 f(L) dL \quad (60)$$

$$= \sum_{[x]} \int_{L([x]) - \frac{1}{2}|\omega([x])|}^{L([x]) + \frac{1}{2}|\omega([x])|} f(L) dL, \quad (61)$$

$$= E(f(L)) + \varepsilon, \quad (62)$$

$$\varepsilon = \sum_{[x]} \int_{L([x]) - \frac{1}{2}|\omega([x])|}^{L([x]) + \frac{1}{2}|\omega([x])|} (f(L) - f(L([x]))) dL. \quad (63)$$

Now in the sum that defines  $\varepsilon$ , there are at most 2 discontinuities in an otherwise constant 0, 1 or -1 integrand. If none occur in an interval, that term is exactly zero. If one occurs, the integrand is zero for at least  $\frac{1}{2}$  of the interval, so that error is at most  $\frac{1}{2}|\omega([x])|$ , and can be only once more in one other interval (with the same kind of error). And if 2 occur, they occur

nowhere else and the error is bounded by  $|\omega([x])|$ . In each case we have our theorem.

### From Computation

In this section, we again consider only a finite probability space  $X$  with  $|X|$  elements  $x_k, k = 1, \dots, |X|$  and probabilities  $p_k = p(x_k)$ .

Let  $S = (s_1, \dots, s_n)$  be  $|S|$  independent samples taken from  $X$ . We define an estimator for  $L(x)$  as:

$$\begin{aligned} \ell(x) &= \frac{1}{|S|} \{\# \text{ of outcomes in } S \text{ more probable than } x\} \\ &\quad + \frac{1}{2|S|} \{\# \text{ of outcomes in } S \text{ equally probable to } x\} \end{aligned}$$

We asserted that

$$E(\ell(x)) = L(x). \quad (64)$$

$$E([\ell(x) - L(x)]^2) = \frac{1}{|S|} \left[ L(x) \cdot (1 - L(x)) - \frac{1}{4} |\omega(x)| \right]. \quad (65)$$

For the remainder of this section, we consider  $x \in X$  to be a fixed choice, and so drop the  $\cdot(x)$  functional notation, which would otherwise pepper every equation. This way (64-65) becomes

$$E(\ell) = L. \quad (66)$$

$$E([\ell - L]^2) = \frac{1}{|S|} \left[ L \cdot (1 - L) - \frac{1}{4} |\omega| \right] \quad (67)$$

The point of this section is to work out the details of these facts.

The definition of  $\ell$  does not depend on the ordering of the sample  $S$ , so the vector of counts  $(c_1, \dots, c_{|X|})$  is equivalent information about the sample for our purposes. We define  $c$  as the vector of counts of elements in  $S$ , so that  $x_k$  occurs  $c_k$  times in the sample  $S$ .

The probability of obtaining a sample counts  $c$  is given by the multinomial distribution

$$P(c) = \binom{n}{c} p^c = \frac{n!}{c_1! \cdots c_{|X|}!} p_1^{c_1} \cdots p_{|X|}^{c_{|X|}}, \text{ where } n = \sum c = |S|. \quad (68)$$

Here are some well-known moments of the multinomial distribution (the sum is over all samples with  $\sum c = n$ ):

$$\sum_c P(c) = 1, \quad (69)$$

$$\sum_c P(c) c = np, \quad (70)$$

$$\sum_c P(c) cc^T = n \text{diag}(p) + (n^2 - n)pp^T. \quad (71)$$

Here  $c$  and  $p$  are  $|X| \times 1$  column vectors, and  $\text{diag}(p)$  is the diagonal  $|X| \times |X|$  matrix with the probabilities  $p$  along the diagonal.

It is useful to introduce  $A$  and  $a$ , which are analogous to  $\Omega$  and  $\omega$  (recall that  $x$  is fixed in this discussion):

$$A = \{k \mid p(x_k) > p(x)\} , \quad (72)$$

$$= \{k \mid x_k \in \Omega\} , \quad (73)$$

$$|A| = \frac{1}{n} \sum_{k \in A} c_k . \quad (74)$$

Note that  $A$  is the index set of  $\Omega$ , and so independent of  $c$ , but  $|A|$  depends on the sample counts. By contrast, both  $\Omega$  and  $|\Omega|$  are constant (properties of the entire probability space) for a fixed choice of  $x$ .

Similarly,

$$a = \{k \mid p(x_k) = p(x)\} , \quad (75)$$

$$= \{k \mid x_k \in \omega\} , \quad (76)$$

$$|a| = \frac{1}{n} \sum_{k \in a} c_k . \quad (77)$$

Using the well-known multinomial moment results, we can compute some useful expected values of  $|A|$  and  $|a|$ :

First  $E(|A|) = |\Omega|$  :

Similarly  $E(|a|) = |\omega|$ .

$$E(|A|) = \sum_c P(c) \frac{1}{n} \sum_{k \in A} c_k \quad (78)$$

$$= \frac{1}{n} \sum_{k \in A} \sum_c P(c) c_k \quad (79)$$

$$= \frac{1}{n} \sum_{k \in A} n p_k \quad (80)$$

$$= |\Omega| . \quad (81)$$

Next  $E(|A|^2) = \frac{1}{n} |\Omega| \cdot (1 - |\Omega|) + |\Omega|^2$  :

Similarly,  $E(|a|^2) = \frac{1}{n} |\omega| \cdot (1 - |\omega|) + |\omega|^2$ .

$$E(|A|^2) = \frac{1}{n^2} \sum_{k,k' \in A} \sum_c P(c) c_k c_{k'} , \quad (82)$$

$$= \frac{1}{n^2} \sum_{k,k' \in A} \left[ n \text{diag}(p) + (n^2 + n) p p^T \right]_{k,k'} , \quad (83)$$

$$= \frac{1}{n} |\Omega| + \frac{n-1}{n} |\Omega|^2 , \quad (84)$$

$$= \frac{1}{n} |\Omega| \cdot (1 - |\Omega|) + |\Omega|^2 . \quad (85)$$

We also will need the cross term  $E(|A||a|)$ :

$$E(|A||a|) = \frac{1}{n^2} \sum_{k \in A, k' \in a} \sum_c P(c) c_k c_{k'}, \quad (86)$$

$$= \frac{1}{n^2} \sum_{k \in A, k' \in a} \left[ n \operatorname{diag}(p) + (n^2 + n) p p^T \right]_{k, k'}, \quad (87)$$

$$= \frac{n-1}{n} |\Omega| |\omega|. \quad (88)$$

Note that the diagonal term  $\operatorname{diag}(p)$  contributes nothing in this case because  $A$  and  $a$  are disjoint sets.

With these preliminaries,

$$\ell = |A| + \frac{1}{2} |a|, \quad (89)$$

And so

$$E(\ell) = |\Omega| + \frac{1}{2} |\omega| = L.$$

which is the first result.

For the variance,

$$E([\ell - L]^2) = E\left(\left(|A| - |\Omega| + \frac{1}{2}(|a| - |\omega|)\right)^2\right), \quad (90)$$

$$= E([|A| - |\Omega|]^2) + E([|A| - |\Omega|][|a| - |\omega|]) + \frac{1}{4} E([|a| - |\omega|]^2), \quad (91)$$

$$= E(|A|^2) - |\Omega|^2 + E(|A||a|) - |\Omega||\omega| + \frac{1}{4} \left\{ E(|a|^2) - |\omega|^2 \right\}, \quad (92)$$

$$= \frac{1}{n} \left\{ |\Omega|(1 - |\Omega|) - |\Omega||\omega| + \frac{1}{4} |\omega|(1 - |\omega|) \right\}. \quad (93)$$

A little algebra will show (93) is equal to

$$E([\ell - L]^2) = \frac{1}{n} \left\{ L \cdot (1 - L) - \frac{1}{4} |\omega| \right\} \leq \frac{L \cdot (1 - L)}{n}. \quad (94)$$

which is the other thing we wanted to prove.

### *Infinite and continuous spaces.*

Theorems 4-6 have been proven for any finite probability space  $X$ .

$$T_4(X) : E(L(x)) = \frac{1}{2}, \quad (95)$$

$$T_5(X) : \left| E(f(L)) - \int_0^1 f(L) dL \right| \leq \frac{\max_L |f''(L)| \cdot \sup_x |\omega(x)|^2}{24}, \quad (96)$$

$$T_6(X) : 0 \leq \frac{1}{p+1} - E(L^p) \leq \frac{p \cdot (p-1) \cdot \sup_x |\omega(x)|^2}{24}. \quad (97)$$



But expectation  $E(\cdot)$  and measure  $|\cdot|$  are continuous in the space of probability measures, and the inequalities are closed, so by a cauchy sequence of measures, these statements are true in general.

