

MINIMAL ORDERED PERFECT HASHING (graph)

- given a set of strings: $\{aa, ba, bb, bc, ca\}$
 $\begin{matrix} & & & & & \\ & 0 & 1 & 2 & 3 & 4 \end{matrix}$

=> assign codes that are rankings of the strings starting from 0

- given $h(x)_s$: \nearrow 2 chars of the strings

$$h_1(c'c'') = 2 \cdot \text{rank}(c') \cdot \text{rank}(c'') \bmod 11$$

$$h_2(c'c'') = 5 \cdot \text{rank}(c') + \text{rank}(c'') \bmod 11$$


	h_1	h_2	h_3
aa	0	8	2
ba	1	1	7
bb	2	7	8
bc	3	2	9
ca	4	5	1

rank	a	b	c
	2	3	4

$$h_1(cc) = 2 \cdot 4 \cdot 4 = 32 \\ \bmod 11 = 10$$

$$h_2(cc) = 5 \cdot 4 + 5 = 25 \\ \bmod 11 = 3$$

compute function g (represented by an array)

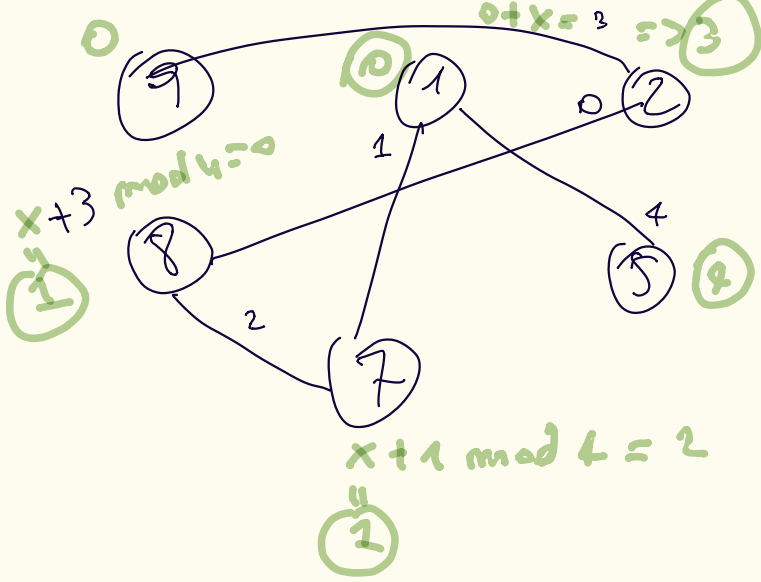
g 

=> build the graph with 11 nodes $[0, 10]$

\hookrightarrow domain from
max values

=> $[0, 10]$

cut
 $\bmod 11$



\Rightarrow ACYCLIC \Rightarrow solvable

(if cyclic \Rightarrow is solvable)

\Rightarrow START from the last node & propagate (5) =

final result

$$h(t) = g(h_1(t)) + g(h_2(t)) \pmod{5}$$

of strings

	0	1	2	3	4	5	6	7	8	9	10
g	0	0	3	0	0	4	0	2	1	0	0
	↓		↓	↓	↓					↓	

any values. as 0

o) Show

$$h(BA) = g(\underbrace{h_1(BA)}_1) + g(\underbrace{h_2(BA)}_7) \pmod{5} = 1$$

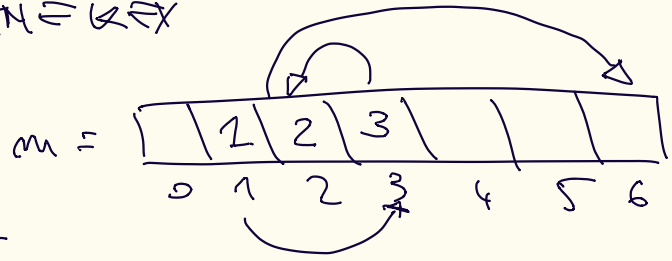
$$h(CC) = g(\underbrace{h_1(CC)}_{10}) + g(\underbrace{h_2(CC)}_3) \pmod{4} = 0$$

Cuckoo Hashing

$$h_1(x) = x \bmod 7$$

$$h_2(x) = 3x \bmod 7$$

: table of $m=7 \Rightarrow 7$ positions
 \Rightarrow SHOW CONTENT OF THE TABLE INSERTING THE KEY



keys to insert:

{3, 2, 1, 0, 6, 15}

$m=7$

• When both cells after h_1 & h_2 are occupied:

\Rightarrow PREFER ALWAYS h_1 TO INSERT

$$h_1(3) = 3 \bmod 7 = 3$$

$$h_2(3) = 3 \cdot 3 \bmod 7 = 2$$

$$h_1(2) = 2 \bmod 7 = 2$$

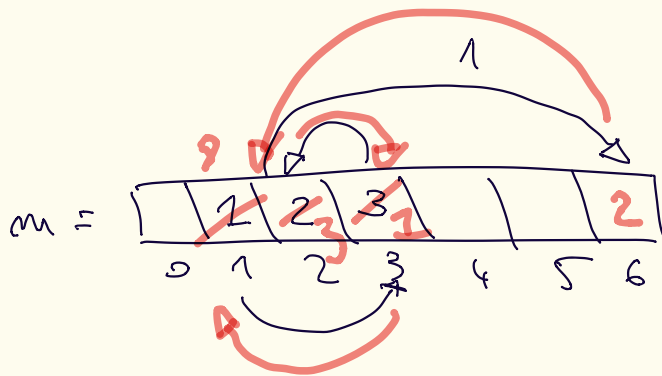
$$h_1(2) = 6 \bmod 7 = 6$$

$$h_1(1) = 1 \bmod 7 = 1$$

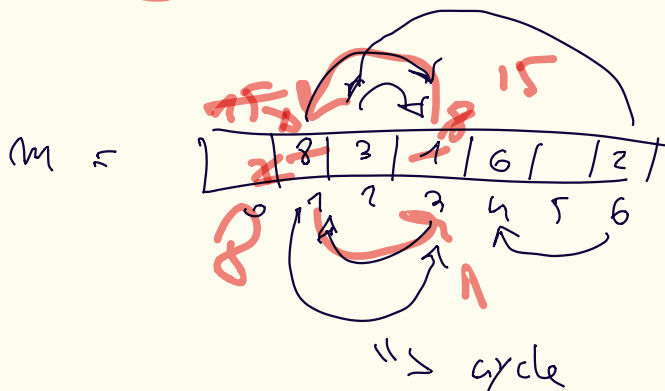
$$h_2(1) = 3 \cdot 1 \bmod 7 = 3$$

$$h_1(0) = 0 \text{ NOT KICK OUT}$$

$$h_2(0) = 0 \bmod 7 = 0$$



- kick out
 - insert the new



"> cycle

\Rightarrow loop with 15

involving 1, 2, 15 TRY h_2

\Rightarrow 15 CANNOT BE INSERTED

$$h_1(6) = 6 \text{ X}$$

$$h_2(6) = 18 \bmod 7 = 4$$

$$h_1(15) = 1 \text{ NOT}$$

$$h_2(15) = 15 \cdot 3 \bmod 7 = 45 \bmod 7 = 3$$

=> Should take other
2 $h(x)$

STRINGS: $S = \{AA, AC, BB, CC\}$ compute MOPE

given $h_1(xy) = x + y \mod 7$

$h_2(xy) = x + 2 \cdot y \mod 7$

CODES are

$A = 1$

$B = 2$

$C = 3$

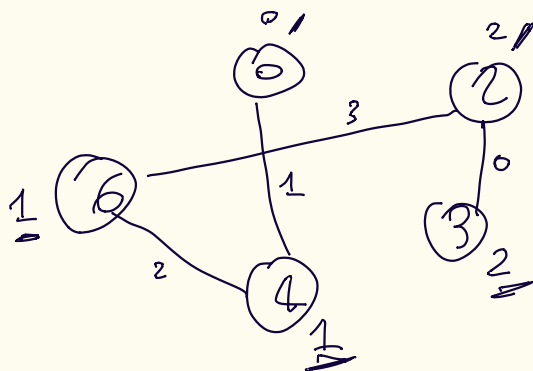
$AA = 1 + 1 = 2$

$AC = 1 + 3 = 4$

$BB = 2 + 2 = 4$

$CC = 3 + 3 = 6$

key	Edge label		NODES LABEL	
	h	h_1	h_2	
AA	0	<u>2</u>	$1 + 2 \cdot 1 \text{ m } 7 =$	<u>3</u>
AC	1	<u>4</u>	$1 + 2 \cdot 3 \text{ m } 7 =$	<u>0</u>
BB	2	<u>4</u>	$2 + 2 \cdot 2 \text{ m } 7 =$	<u>6</u>
CC	3	<u>6</u>	$3 + 2 \cdot 3 \text{ m } 7 =$	<u>2</u>



mod 4

$m = 7$

$[0, 6]$

$h(xy) = g(h_1(xy)) + g(h_2(xy)) \mod 4$

g

0	0	2	2	1	0	1
0	1	2	3	4	5	6

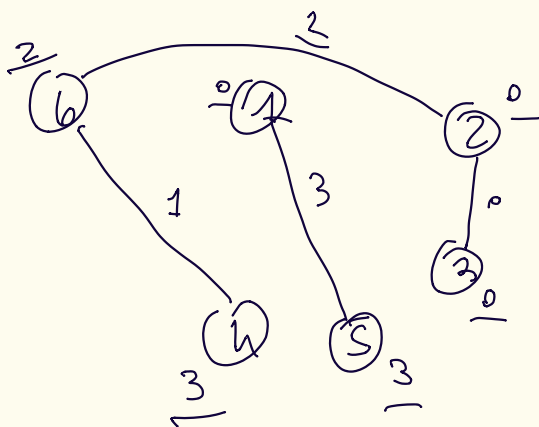
MORF

$$S = \{ \underset{0}{11}, \underset{1}{22}, \underset{2}{33}, \underset{3}{44} \}$$

$$h_1(xy) = x + y \mod 7$$

$$h_2(xy) = x + 2y \mod 7$$

keys	h	h_1	h_2
11	0	2	$1 + 2 \cdot 1 \mod 7 = 3$
22	1	4	$2 + 2 \cdot 2 \mod 7 = 6$
33	2	6	$3 + 2 \cdot 3 \mod 7 = 2$
44	3	1	$4 + 2 \cdot 4 \mod 7 = 5$



remember
mod 4

$$h(t) = g(t(h_1)) + g(t(h_2)) \mod x$$

0	0	0	0	3	3	2
0	1	2	3	4	5	6

MOPF

$$S = \{ \underset{0}{AA}, \underset{1}{AC}, \underset{2}{BC}, \underset{3}{CC} \}$$

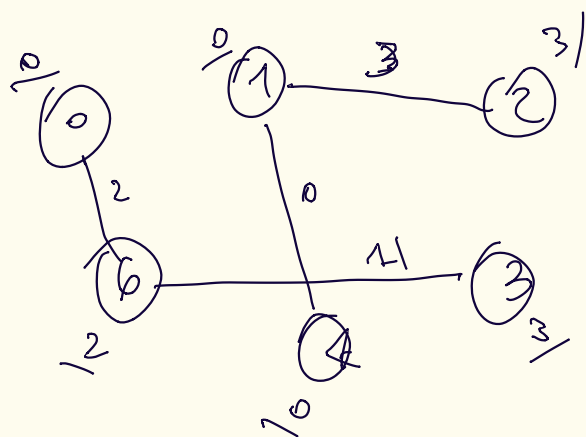
$$h_1(xy) = 3 \text{rank}(x) + \text{rank}(y) \bmod 7$$

$$\text{rank}(x) = 2, 3, 4$$

$$x = a, b, c$$

$$h_2(xy) = \text{rank}(x) + \text{rank}(y) \bmod 7$$

keys	h_1	h_1	h_2
AA	0	$3 \cdot 2 + 2 \bmod 7 = 1$	4
AC	1	$3 \cdot 2 + 4 \bmod 7 = 3$	$2 + 4 = 6$
BC	2	$3 \cdot 3 + 4 \bmod 7 = 6$	$3 + 4 = 7 \bmod 7 = 0$
CC	3	$3 \cdot 4 + 4 \bmod 7 = 2$	$4 + 4 \bmod 7 = 1$



$\bmod 4$

$$h(t) = g(h_1(t)) + g(h_2(t)) \bmod 4$$

$$x = 4$$

0	0	3	3	0	0	2
0	1	2	3	4	5	6

