

Sum of  $n$ -integers

$A_{s,b}$   $s$ : jumps  $b$ : # items per block  $B$ : page

$$\Rightarrow A_s = A[s \cdot b + 1, (s+1) \cdot b]$$

- Scan all items in each block before moving to the next one

$s \cdot b = 1 \Rightarrow$  normal sequential access

$b$  large  $\Rightarrow$  block wide access

$s$  large  $\Rightarrow$  random-like access

- Not all values of  $s$  takes account all blocks

- If  $s$  is coprime with  $\frac{M}{b} \Rightarrow$  then  $\left\{ s \cdot i \bmod \left( \frac{M}{b} \right) \right\}$  generates a permutation  
 $\Rightarrow$  thus  $A_{s,b}$  touches all block  
 $\Rightarrow$  sum of  $n$  integers

Cost

- cpu time =  $n$

- space occupancy =  $\frac{M}{B}$  pages

1) Case  $(s=1) \Rightarrow A_{1,b} \Rightarrow \forall b \neq 1$  it scans all block with word  
 $\Rightarrow$  takes  $O\left(\frac{M}{b}\right)$  I/Os

2) Case  $(s=2, b < B)$

$\Rightarrow$  touches  $\frac{1}{2}$  the half of the  $\frac{M}{b}$  logically divided small blocks  $b$   
 $\Rightarrow$  jumping by 2

$$\Rightarrow \left( \frac{2M}{b} \right) \text{ I/Os}$$

$\Rightarrow$  total cost:

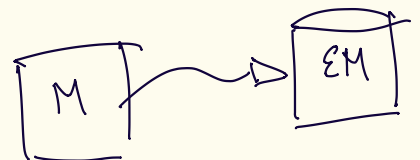
$$\min \left\{ s, \frac{B}{b} \right\} \cdot \frac{M}{b} \text{ I/Os}$$

Memory: assume input is size

$$m = (1 + \epsilon) M$$

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$\Rightarrow$



want evaluate  $p(\epsilon)$  = "probability I'm going to access the disk"

$$\frac{\epsilon M}{m} = \frac{\epsilon M}{(1 + \epsilon) M} = \frac{\epsilon}{1 + \epsilon}$$

if  $p(\epsilon) = 1 \Rightarrow$  share  $p$  to disk  
 $p(\epsilon) = 0 \Leftrightarrow$  stay in memory  
 $p(\epsilon) = \frac{\epsilon}{1 + \epsilon}$  random behaviour

EVALUATE AVG TIME FOR A STEP

Let be:  $\alpha$ : frequent of steps (load / write)  $\rightarrow$  I/O  
 $(1-\alpha)$ : " " " (computation)  $\rightarrow$  CPU

$$(1-\alpha) \cdot \overset{\text{computation}}{\alpha} + \alpha \cdot \overset{\text{I/O cost}}{\beta} [ \beta (1-p(\epsilon)) + c \cdot (p(\epsilon)) ] =$$

$$\alpha = \beta = 1$$

$$= \underbrace{(1-\alpha) \cdot 1}_{\leq 1} + \alpha \underbrace{[ 1 \cdot (1-p(\epsilon)) + c \cdot (p(\epsilon)) ]}_{\leq 1} =$$

$$= \alpha \cdot c \cdot p(\epsilon)$$

$$\alpha \approx 30\%$$

$$10^5 \leq \epsilon \leq 10^6$$

$$\hookrightarrow 300000 \cdot p(\epsilon)$$