


Algorithmics	Student information	Date	Number of session
	UO: 301949	27/02/2025	3
	Surname: Sánchez Menéndez	 Escuela de Ingeniería Informática Universidad de Oviedo	
	Name: Adrián		



Activity 1. Divide and Conquer by subtraction

Subtraction1: $O(n)$

For this implementation, the size of the problem is not big enough to get a significant result of the increase in the time it takes the algorithm to run. After using repetitions, we can calculate the time for each size and see that it doubles.

Subtraction2: $O(n^2)$

For this implementation, the size of the problem is increased by 2, and as the theoretical complexity expects, the times are increased by $2^2=4$

Subtraction3: $O(2^n)$

For this implementation, the size of the problem is increased by 1 linearly, and so to compute the time the next iteration will take, we use the expression:

$$t_i = \frac{n_i}{n_j} t_j = \frac{2^i}{2^j} t_j = 2^{i-j} \times t_j$$

which solving for $i = j+1$, will indicate that the time doubles when the size is increased by 1 and that matches the times obtained by measuring the algorithm.

- Above $n=8192$, Substraction1 and Substraction2 generate StackOverflowErrors, because the stack gets overloaded with the recursive calls from the method.
- To compute the time for Substraction3 with $n = 80$, the formula shown before is needed:

$$\begin{aligned}
t_{80} &= 2^{80-25} \times t_{25} = 2^{55} \times 1371 \text{ ms} = 4.94 \times 10^{19} \text{ ms} \\
4.94 \times 10^{19} \text{ ms} &\times \frac{1 \text{ s}}{1000 \text{ ms}} \times \frac{1 \text{ min}}{60 \text{ s}} \times \frac{1 \text{ h}}{60 \text{ min}} \times \frac{1 \text{ d}}{24 \text{ h}} \times \frac{1 \text{ y}}{365.25 \text{ d}} \\
&= 1,565,248,330 \text{ years}
\end{aligned}$$

- Subtraction4: $O(n^3)$

n	Substraction4
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100	LoR
200	LoR
400	119
800	913
1600	7106
3200	56909
6400	OoT

- Subtraction5: $O(3^{n/2})$

n	Subtraction5
30	457
32	1335
34	4029
36	11662
38	35151
40	OoT

- To compute the time for Subtraction5 with $n = 80$, the formula shown above is needed:

$$\begin{aligned}
 t_{80} &= 3^{\frac{80-34}{2}} \times t_{34} = 3^{23} \times 4029 \text{ ms} = 3.57 \times 10^{19} \text{ ms} \\
 3.57 \times 10^{19} \text{ ms} &\times \frac{1 \text{ s}}{1000 \text{ ms}} \times \frac{1 \text{ min}}{60 \text{ s}} \times \frac{1 \text{ h}}{60 \text{ min}} \times \frac{1 \text{ d}}{24 \text{ h}} \times \frac{1 \text{ y}}{365.25 \text{ d}} \\
 &= 1.13 \times 10^{15} \text{ years}
 \end{aligned}$$

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Activity 2. Divide and Conquer by division

Division1: $O(n)$

For this implementation, the size of the problem is increased by 2, and as the theoretical complexity expects, the times are double.

Division2: $O(n \log(n))$

For this implementation, the size of the problem is increased by 2, and as the theoretical complexity expects, the times are increased by $2^1 \log 2 = 0.603$. In reality, the times are much closer to a complexity of $O(n)$ because n is a very small number and the $\log(2)$ is not significant.

Division3: $O(n)$

For this implementation, the size of the problem is increased by 2, and as the theoretical complexity expects, the times are double.

$$n^{\log_b a} = n^{\log_2 2} = n$$

- Division4: $O(n^2)$ $a < b^k$

n	Division4
1000	LoR
2000	LoR
4000	121
8000	483
16000	1834
32000	7301
64000	28912
128000	OoT

- Division5: $O(n^2)$ $a > b^k$

n	Division5
1000	LoR

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2000	113
4000	446
8000	1749
16000	6812
32000	26680
64000	OoT

Activity 3. [TITLE OF THE ACTIVITY]

n	sum1 $O(n)$	sum2 $O(n^{k+1})=O(n)$	sum3 $O(n)$
3	LoR	LoR	LoR
6	LoR	LoR	LoR
12	LoR	LoR	LoR
24	LoR	LoR	91×10^{-5}
48	LoR	87×10^{-5}	181×10^{-5}
96	60×10^{-5}	172×10^{-5}	351×10^{-5}
192	116×10^{-5}	334×10^{-5}	698×10^{-5}
384	226×10^{-5}	667×10^{-5}	1425×10^{-5}
768	439×10^{-5}	1325×10^{-5}	2690×10^{-5}
1536	896×10^{-5}	2495×10^{-5}	5229×10^{-5}
3072	1805×10^{-5}	4904×10^{-5}	10764×10^{-5}
6144	3510×10^{-5}	9907×10^{-5}	21564×10^{-5}
12288	7113×10^{-5}	StackOverflow	

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n	fib1 $O(n)$	fib2 $O(n)$	fib3 $O(n)$	fib4 $O(1.6^n)$
10	101×10^{-6}	134×10^{-6}	202×10^{-6}	263×10^{-5}
20	186×10^{-6}	237×10^{-6}	360×10^{-6}	3148×10^{-4}
30	249×10^{-6}	343×10^{-6}	510×10^{-6}	37.96
40	344×10^{-6}	487×10^{-6}	661×10^{-6}	4655
50	418×10^{-6}	566×10^{-6}	831×10^{-6}	OoT
60	460×10^{-6}	696×10^{-6}	972×10^{-6}	OoT

Activity 4. Petanque championship organization

The algorithm used to solve the complexity of the problem has a complexity of $O(n^2)$, as it is implemented using divide and conquer and the number of recursive calls is 2, the size of the problem is reduced by half in each iteration and the complexity of the code without the recursive calls is $O(n^2)$; so $a=2$, $b=2$ and $k=2$ and as it is a division divide and conquer and $a < b^k$, then the complexity is $O(n^k)$, or $O(n^2)$.

n	t Calendar
2	LoR
4	LoR
8	LoR

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16	1.15×10^{-2}
32	4×10^{-2}
64	0.1344
128	0.5049
256	1.941
512	7.646
1024	31.2
2048	125.3
4096	528.8
8192	2201.9
16384	10533
32768	OoM

OoM = Out of Memory

The size of the problem is increasing by 2 in each iteration, and as the expected complexity of the algorithm predicts, the times increase by $2^2 = 4$.