



BIT MANIPULATION 2

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Goal

- Cool tips for bit manipulation.
- Brute-forcing using bitmasks.
- Understanding bitsets.

Some Cool Tricks

- ✓ Swap 2 numbers without a temporary variable
- ✓ $(2^k - 1)$ has last k bits set
- ✓ __builtin_popcount and __builtin_popcountll → counting the no. of set bits.
- ✓ __builtin_clz / __builtin_clzll → count leading zeroes
- ✓ __builtin_ctz / __builtin_ctzll → counts # of trailing zeroes.

$$a = a^5$$

$$b = a^{-1}$$

$$\begin{array}{c} \frac{a}{\underline{b}} \\ \hline \boxed{\begin{array}{l} a = a \wedge b \\ b = (\cancel{a}) \wedge b \\ a = a \wedge b \end{array}} \end{array} \quad (a \wedge b) \wedge b = a \wedge (b \wedge b) = a \wedge 0 = \cancel{a}$$

$$\begin{aligned}
 a &= a \wedge b \\
 &= \overline{(a \wedge b)} \wedge (a) \\
 &= \overline{(a \wedge a)} \wedge \overbrace{b}^b = b
 \end{aligned}$$

$$\overline{L \wedge a = 0}$$

$$2^k - 1$$

$$2^k \rightarrow$$

1 0 0 0 0 0 0
 $\underbrace{\hspace{1cm}}_{k \text{ zeros}}$

0 | 0

$$\begin{array}{r} 1000000 \\ - 1 \\ \hline \boxed{9\ 9\ 9\ 9\ 9} \end{array}$$

$$1000000$$

~~0 1 1 1 1~~
 $\underbrace{\hspace{1cm}}_{k \text{ ones}}$

$2^k \rightarrow$

1 0 0 0 0 ... 0
k zeros

$$\begin{array}{r} 1 \overset{\sim}{\overbrace{0,000}} \\ - 1 \\ \hline 09999 \end{array}$$

1000000 - $\boxed{2^s}$

- 1



$\lceil \log n + \# \text{ of digits} \rceil$

$\lceil \log_2(10^5) \rceil$ bits

 $\overline{\text{00000}}$

$\text{clz}(16)$

0 0 0 0 0
0 0 0 0 0 0 0 1 0 0 0 0

Bitmasking

- ❖ A bitmask is nothing but a sequence of bits where every bit represents something.

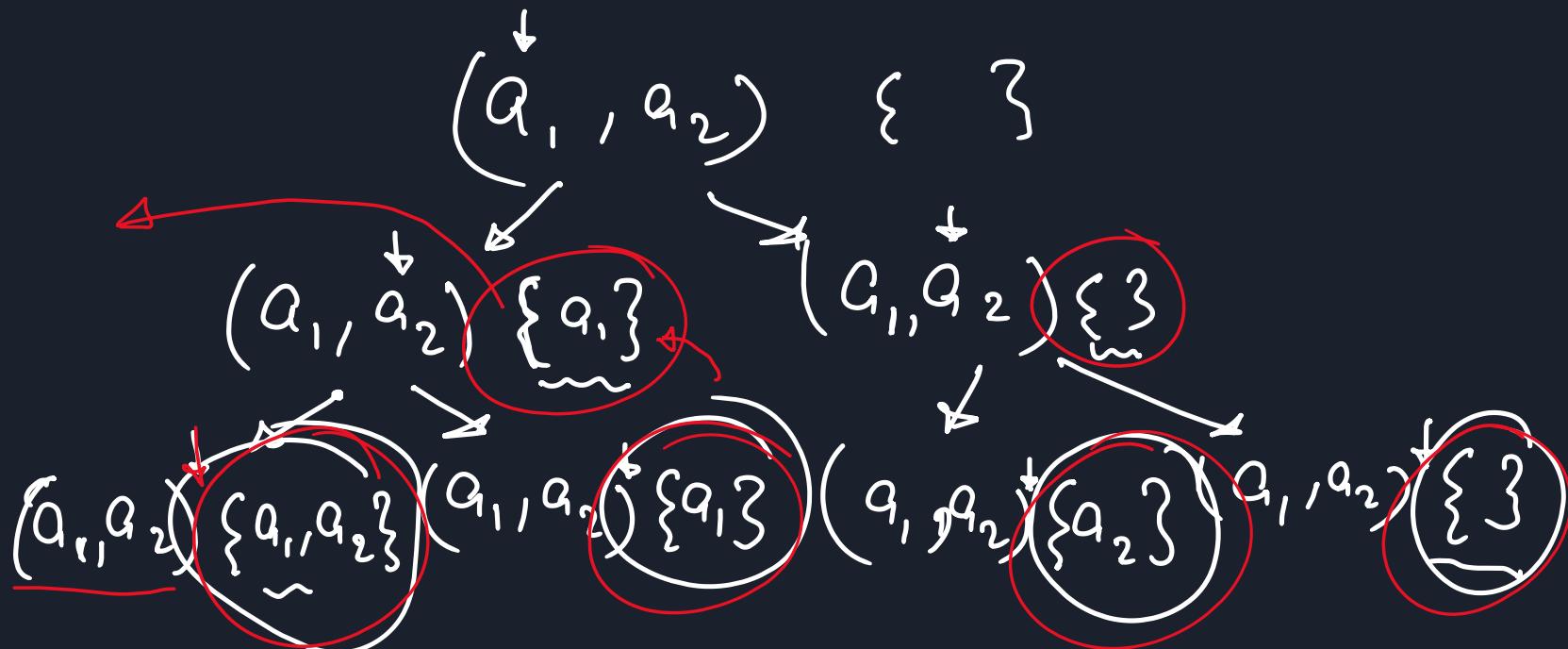
Think of this as a boolean array which takes less memory.

Memory consumed by a boolean array of size 30 = 30 bytes = 240 bits

Here we are using 1 byte for storing true/false which could have been represented by a single bit.

Can you think of some better way which consumes less memory?

$\{q_1, q_2, q_3, q_4, q_5, q_6, q_7\}$



0 1 2 3 4 5 6
a₁, a₂ a₃ a_n a₅ a₆ a₇

either including / excluding

↓
bitmasking

6 5 4 3 2 1 0

· - - - - - 0
6 5 4 3 2 1 0

 0 · - - - - -
 ↑ |

this particular index is
not taken

 - - - - - 1 -
 | ↑
this particular index
is taken

bool vector

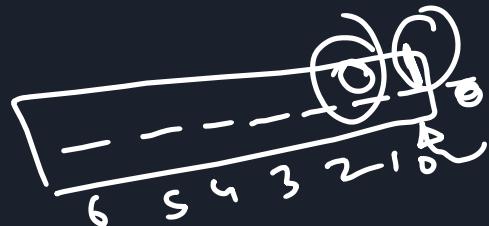
false

true



7

subset having element
with index 3rd and
6th



all the subsets

vector <bool>

T/F → 8 bits

$$30 \times 8 = 240 \text{ bits}$$



Instead of using an array of 30 boolean variables

We can use a sequence of 30 bits where

ith bit will be set if a[i] was true

else ith bit will be unset

We can use a 32-bit integer for allocating this memory so total memory used now is just 32 bits.

dp with bitmasking

32 bit numbers

n numbers in the array

$$2^N - 1$$

$\rightarrow 0$

0 0 0 0 0 0 0 0

10^2

100

mask

0 → empty set

(minimum) 99

0 (maximum) $2^N - 1$ → when we include all the
 $(\ll n) - 1$ last n bits
as 1

$2^N - 1$

$T \approx (20 - 25)$ constant $\frac{1}{N}$ bits T/F 1000
 n bits 16 bits

Bitmasking – Brute Force Code

$n \rightarrow$ size of array $2^N \times N$ T.C
 2^n for (int mask = 0; mask < ($1 \ll n$); mask++) {
 n for (int i = 0; i < n; i++)
 if ((mask & ($1 \ll i$))) set
 cout << a[i] << " ";
 } cout << endl;
 }

mask gives you one possible subset

worst case 1 byte



$n = 3$

for (mask \rightarrow 0 to 7)

→ 2nd bit

$\begin{array}{r} 10100 \\ 00100 \\ \hline 00100 \end{array}$

mask → 0

mask → 1

mask → 2

mask → 3

mask → 4

mask → 5

mask → 6

mask → 7

$\begin{array}{c} 000 \\ 001 \\ \hline 010 \end{array}$

$\{ 1, 2, 3 \}$

$\{ 3 \}$

$\{ 1, 3 \}$

$\{ 2 \}$

$\{ 1, 2 \}$

$\{ 3 \}$

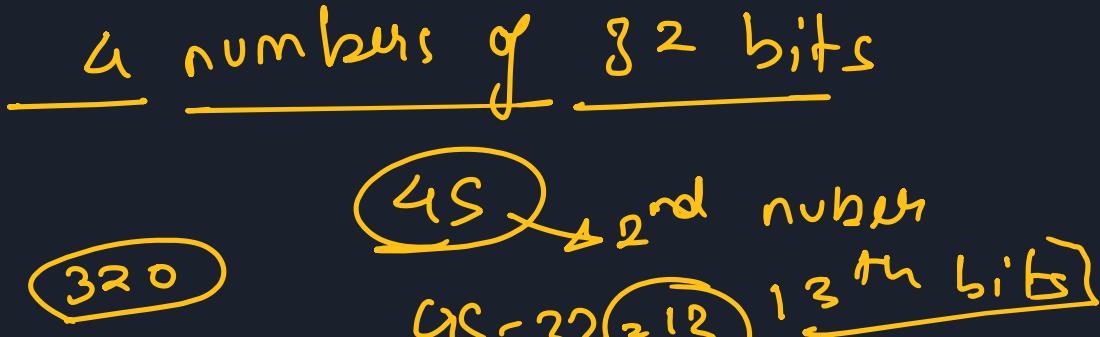
$\{ 1, 3 \}$

$\{ 2, 3 \}$

$\{ 1, 2, 3 \}$

2³

THINK



But what if the array size was 100?

We cannot create a single variable with that many bits.

But we can create, 4 32-bit integers or 2 64-bit integers.

Looks like it is becoming complicated.

If we required 320 bits then it might become complicated for us.



Bitsets

(interview
stuff)

In C++ STL, we have a data structure for this - **Bitsets**

Syntax:-

Defining bitset:

keyword `bitset<320> bit;` variable name

Accessing and setting ith bit:
bits you want to store

`bit[i]` can be used to access and set ith bit just like we use `arr[i]` to access or set the ith element of an array

Bitsets (CONTD...)

- Binary Operations can be performed on Bitsets
- We can initialize bitsets with some binary string or integer also.

Some Useful Functions:-

- bit.count() -> returns the number of set bits
- bit.flip() -> flips all the bits
- any(), none() and all()
- to_string() -> returns binary string







Bit Independent Problems

In these problems, whether the current bit is set or not doesn't affect the other bits, so we can just solve for every bit individually.

Problem:

- XORwice

Bit Dependent Problems



Here the bits are interdependent, whether we set a bit or unset it can affect other bits.

Example - Find the largest number less than 2^8 where at most 3 bits are set.

Here the whether we can set a bit or not depends upon the other bits as we cannot set more than 3 bits.

Problem:

- Gardener and the Array

largest $< 2^8$ — almost 3 bits set



bit dependent