Part B-1

1 Excellence #@REL-EVA-RE@#

1.1 Quality and pertinence of the project's research and innovation objectives (and the extent to which they are ambitious, and go beyond the state of the art)

Proof assistants — also called *interactive theorem provers* (ITPs) — are software tools used to rigorously verify formal modelling and reasoning. Contemporary systems such as *Rocq*, *Lean*, and *Isabelle*[1] offer powerful frameworks for constructing formal specifications and proofs: they have been used successfully in various applications, ranging from the verification of advanced theorems in mathematics to the certification of complex software artifacts such as programming language compilers, operating system kernels and cryptographic protocols[2]. Yet they remain notoriously difficult to learn and use, limiting broader adoption in various settings where they could bring transformative societal impact, such as:

- **mathematics education**, where they could act as a unifying medium for interactive exploration and understanding of mathematical concepts, fostering closer collaboration amongst students and offloading some teaching burden (e.g. grading) through automation[3];

- **formal verification**, where they could bring higher standards of quality assurance (QA) to businesses that rely on complex hardware and software systems, especially in safety-critical industries such as healthcare, transportation and energy;

- **artificial intelligence** (AI), where the uncertainty inherent to current technologies based on probabilistic techniques such as large language models (LLMs) could be mitigated by the exact logical reasoning capabilities of ITPs, an approach sometimes termed *neurosymbolic AI*.

In view of this large potential for applications, it is natural to ask what exactly limits adoption of the current generation of ITPs. The recent surge of interest arising both in academia and industry — in great part due to the popularity of the Lean language and its Mathlib library — suggests that social factors such as public communication, community building, and vast amounts of expository content and learning resources all play an important role in the widespread appropriation of this complex technology. Even more recently, the promise of a new kind of generative AI free from so-called "hallucinations" that could aid in accelerating scientific discoveries has been a powerful narrative attracting much attention and funding[4].

However, many researchers in the field agree that current ITPs suffer from more *foundational* issues that affect directly their accessibility and ease of use, as well as their ability to scale to larger developments. While these issues are quite diverse in nature, a recurring theme since the birth of the technology in the 60s is the overwhelming **bureaucracy** involved in formalization efforts: formal proofs require a level of care and detail that is far superior and much more time consuming than what is expected from informal paper

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[1] Team, The Rocq Development. The Rocq Prover. Apr. 2025. https://doi.org/10.5281/zenodo.15149629; Moura, Leonardo de, and Sebastian Ullrich. *Automated Deduction – CADE 28*. Edited by André Platzer and Geoff Sutcliffe, 2021, pp. 625–35; Tobias Nipkow, et al., *Isabelle/HOL: a proof assistant for higher-order logic*, vol. 2283 (Springer Science & Business Media, 2002).

[2] Georges Gonthier, "Formal Proof—The Four- Color Theorem" vol. 55, no. 11, 2008, Xavier Leroy, "Formal Verification of a Realistic Compiler", *Commun. ACM* vol. 52, no. 7, July 2009, https://doi.org/10.1145/1538788.1538814, pp. 107–15; Klein, Gerwin, et al. "seL4: Formal Verification of an OS Kernel". *Proceedings of the ACM SIGOPS 22nd Symposium on Operating Systems Principles*, SOSP '09, Oct. 2009, pp. 207–20. https://doi.org/10.1145/1629575.1629596; Gilles Barthe, et al., "EasyCrypt: A Tutorial", *Foundations of Security Analysis and Design VII: FOSAD 2012/2013 Tutorial Lectures*, edited by Alessandro Aldini et al. (Cham: Springer International Publishing, 2014) pp. 146–66.

[3] Minh, Frédéric Tran, et al. *Research Report - Proof Assistants for Teaching: A Survey*. Apr. 2024. Accessed 4 Sept. 2025.

[4] Cade Metz, "Is Math the Path to Chatbots That Don't Make Stuff Up?", *The New York Times*, Sept. 2024,

33 proofs. This has been measured by the so-called *DeBruijn factor* comparing the size of formal and informal
34 proofs, often reaching a value of 4 on average[5].

The main approach to tame this complexity has been to *automate* the various processes involved in formaliza-
tion, which fall roughly within two categories: *elaboration*, concerned with the translation of requirements
expressed in natural language or mathematical notations into precise logical specifications; and *synthesis*,
where the system attempts to automatically generate (parts of) proofs and programs meeting these specifi-
cations. Progress on both fronts is currently being made with the help of state-of-the-art machine learning
techniques, including LLMs[6]. More theoretical research has also been pursued in **type theory**, the field
studying the logical formalisms at the foundation of all modern ITPs. They are the ultimate backbone on
which relies our *trust* in the output of these systems, and thus a key differentiator with respect to purely
probabilistic approaches to (generative) AI.

However, little attention has been devoted by the ITP community to another discipline closely related to
type theory: **proof theory**. In particular, *structural* proof theory is its branch concerned with the study of
combinatorial structures for representing and manipulating proofs. One can identify mainly three motivations
for this study:

- **Challenge 1** (C1) is the fundamental problem of **proof identity**, also known as Hilbert's 24th problem[7].
  It aims to answer the philosophical question "what is a proof?", and the mathematical question "when
  are two proofs equal?". It is thus intimately related to *homotopy type theory* which also investigates
  the structure of *equality*, a known weakness of many type theories.

- **Challenge 2** (C2) is to find proof systems for **non-standard logics** — such as *modal*, *intermediate*,
  substructural and *fixpoint* logics — satisfying good enough properties as to render an algorithmic
  treatment of these logics tractable. The most important property in this respect is that of **cut elim-
  ination**, which is essential both to reduce the complexity of *proof search* (proof *synthesis* in ITP
  terminology), and to ensure productivity of *program execution* through the **Curry-Howard corre-
  spondence** (CHC) between proofs and programs. The CHC is also at the heart of the *calculus of
  inductive constructions* (CoIC), which is the type theory used by the two leading proof assistants Rocq
  and Lean. Researchers are increasingly interested in type-theoretic formulations of these logics as
  they provide expressive languages for specifying behaviors of programs that go beyond pure functions,
  including effects (modal logics), resource-sensitivity (modal and substructural logics) and recursion
  (modal and fixpoint logics)[8].

- **Challenge 3** (C3) is to improve the **efficiency** of computational procedures on proofs, but also on pro-
  grams through the CHC. A well-established principle in computer science and software engineering
  is that choosing appropriate data structures for a problem can lead to orders-of-magnitude improve-
  ments in efficiency. Finding the right data structures for such general classes of objects as proofs and
  programs is thus a very enticing goal with wide implications, including faster automation in ITPs.

In the past decades, two families of proof formalisms have emerged to tackle these challenges:

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[5]Wiedijk, Freek. The De Bruijn Factor. 2000. www.cs.ru.nl/F.Wiedijk/factor/factor.pdf.

[6]Lasse Blaauwbroek, et al., "Learning Guided Automated Reasoning: A Brief Survey", vol. 14560 (2024) pp. 54–83.

[7]Lutz Straßburger, "The problem of proof identity, and why computer scientists should care about Hilbert's 24th problem",
Philosophical Transactions of the Royal Society A: Mathematical, Physical and Engineering Sciences vol. 377, no. 2140,
11 Mar. 2019, https://doi.org/10.1098/rsta.2018.0038, p. 20180038.

[8]Wenhao Tang, et al., "Modal Effect Types", *Proc. ACM Program. Lang.* vol. 9, OOPSLA1, Apr. 2025, https://doi.org/10.1145/
3720476, Marshall, Danielle, et al. "Linearity and Uniqueness: An Entente Cordiale". *Programming Languages and Systems*,
edited by Ilya Sergey, 2022, pp. 346–75. https://doi.org/10.1007/978-3-030-99336-8_13; Ranald Clouston, et al., "The
Guarded Lambda-Calculus: Programming and Reasoning with Guarded Recursion for Coinductive Types", *Logical Methods
in Computer Science* vol. Volume 12, Issue 3. *arXiv*, arxiv.org/abs/1606.09455, Apr. 2017, https://doi.org/10.2168/LMCS-
12(3:7)2016, p. 2019.

- **Graphical proof systems** represent proof objects as *graphs* instead of *trees* of inference rules. *Proof nets*[9] are one of the first graphical proof formalisms, meant to reduce the bureauceacy of proofs in linear logic (C2) in order to identify their essence (C1). Further developments stemming from proof nets like the *geometry of interaction* and *interaction nets* have found applications in program optimization for hardware synthesis and parallelized execution[10] (C3). The *combinatorial proofs* of Hughes are direct decendants of proof nets mainly concerned with C1[11], and *string diagrams* from category theory have been applied to logic and programming languages, with well-known connections to proof nets[12].

- **Deep inference** generalizes Gentzen-style proof systems by allowing inference rules to be applied at any depth inside of a formula, rather than being restricted to its top-level logical connective. The terminology of "deep inference" was proposed by Alessio Guglielmi, who invented the *calculus of structures* to overcome the inability of sequent calculus to capture the substructural logic MV[13]. Since then, calculi of structures and so-called *nested sequent calculi* have been introduced to give proof systems enjoying cut-elimination to most substructural, modal and intermediate logics[14] (C2). Deep inference has also been used in the study of proof complexity, providing in some cases exponential speedup over sequent calculus with respect to proof size, as well as quasipolynomial-time cut-elimination[15] (C3). Lastly, many deep inference formalisms enjoy CHC-style interpretations with variants of λ-calculus, which also improve space efficiency[16].

The experienced researcher (ER) has accumulated significant knowledge of graphical and deep inference proof systems, as well as their applications to ITPs. This expertise was developed during his PhD thesis[17] titled "Deep Inference for Graphical Theorem Proving", where he designed various proof formalisms that enable a novel approach to interactive theorem proving based on **direct manipulation** of logical statements in a graphical user interface (GUI). This extends earlier works on *Proof-by-Pointing*[18] and *Proof-by-Linking*[19]

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[9] Jean-Yves Girard, "Linear logic", *Theoretical Computer Science* vol. 50, no. 1, 1987, https://doi.org/https://doi.org/10.1016/0304-3975(87)90045-4, pp. 1–101.

[10] Ghica, Dan R. "Geometry of Synthesis: A Structured Approach to VLSI Design". *Proceedings of the 34th Annual ACM SIGPLAN-SIGACT Symposium on Principles of Programming Languages*, POPL '07, Jan. 2007, pp. 363–75. https://doi.org/10.1145/1190216.1190269; Mackie, Ian. "An Interaction Net Implementation of Closed Reduction". *Implementation and Application of Functional Languages*, edited by Sven-Bodo Scholz and Olaf Chitil, 2011, pp. 43–59. https://doi.org/10.1007/978-3-642-24452-0_3.

[11] Dominic Hughes, "Proofs without syntax", *Annals of Mathematics* vol. 164, no. 3, Nov. 2006, https://doi.org/10.4007/annals.2006.164.1065, pp. 1065–76.

[12] Robin Piedeleu and Fabio Zanasi, "An Introduction to String Diagrams for Computer Scientists", *Elements in Applied Category Theory*, May 2025, https://doi.org/10.1017/9781009625715,

[13] Guglielmi, Alessio. *A Calculus of Order and Interaction*. 1999. www.researchgate.net/publication/2807151_A_Calculus_of_Order_and_Interaction.

[14] Roman Kuznets and Lutz Straßburger, "Maehara-style modal nested calculi", *Archive for Mathematical Logic* vol. 58, nos. 3–4, May 2019, https://doi.org/10.1007/s00153-018-0636-1, pp. 359–85; Postniece, Linda. "Proof theory and proof search of bi-intuitionistic and tense logic". 2010, Artwork Size: vii, 228 leaves., vii, 228 leaves. https://doi.org/10.25911/5D5FCC3A4DB33.
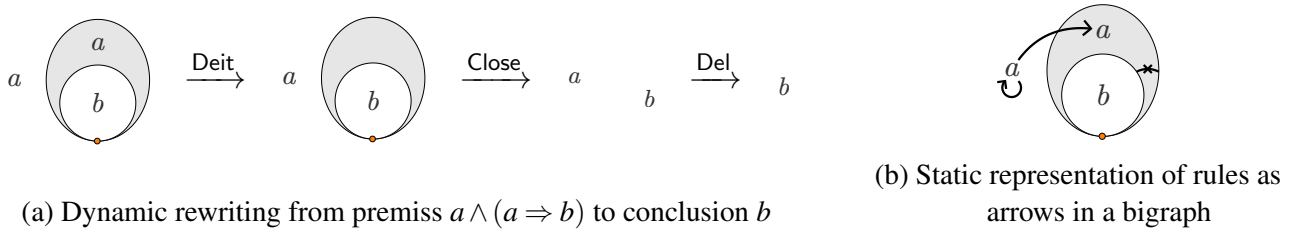
[15] Anupam Das, "On the Relative Proof Complexity of Deep Inference via Atomic Flows", *Logical Methods in Computer Science* vol. Volume 11, Issue 1, Mar. 2015, https://doi.org/10.2168/LMCS-11(1:4)2015, p. 735; Paola Bruscoli, et al., "Quasipolynomial Normalisation in Deep Inference via Atomic Flows and Threshold Formulae", *Logical Methods in Computer Science* vol. Volume 12, Issue 2, May 2016, https://doi.org/10.2168/LMCS-12(2:5)2016, p. 1637.

[16] Nicolas Guenot, "Nested Deduction in Logical Foundations for Computation", PhD dissertation, Ecole Polytechnique X, 2013; Gundersen, Tom, et al. "Atomic Lambda Calculus: A Typed Lambda-Calculus with Explicit Sharing". *28th Annual ACM/IEEE Symposium on Logic in Computer Science*, June 2013, pp. 311–20. https://doi.org/10.1109/LICS.2013.37.

[17] Pablo Donato, "Deep Inference for Graphical Theorem Proving", PhD dissertation, Institut Polytechnique de Paris, 2024.

[18] Yves Bertot, et al., "Proof by pointing", *Theoretical Aspects of Computer Software*, edited by Masami Hagiya and John C. Mitchell, redacted by Gerhard Goos and Juris Hartmanis, vol. 789 (Berlin, Heidelberg: Springer Berlin Heidelberg, 1994) pp. 141–60.

[19] Kaustuv Chaudhuri, "Subformula Linking as an Interaction Method", *Interactive Theorem Proving*, edited by Sandrine Blazy et al., redacted by David Hutchison et al., vol. 7998 (Berlin, Heidelberg: Springer Berlin Heidelberg, 2013) pp. 386–401.

(a) Dynamic rewriting from premiss $a \wedge (a \Rightarrow b)$ to conclusion b

(b) Static representation of rules as arrows in a bigraph

Figure 1: Proof of *modus ponens* $a \wedge (a \Rightarrow b) \vdash b$ in scroll nets

where proofs are constructed through *click* and *drag-and-drop* gestures on formulas, to a more encompassing paradigm termed *Proof-by-Action* (PbA). The goal is to improve accessibility and usability of ITPs by focusing on better principles for *human interaction*, complementing more mainstream research around machine automation. Applying a mixture of graphical and deep inference proof theory to that effect is a highly original endeavor, with no similar efforts in the contemporary research landscape.

In continuation of this programme, the ER has introduced in his last preprint[20] a new graphical framework called **scroll nets**. It is based on a long forgotten diagrammatic proof system called *existential graphs* (EGs), invented by the famous philosopher and logician Charles Sanders Peirce at the dusk of the 19th century — thus predating the very existence of proof theory and computer science. Proofs in EGs are defined by a small set of inference rules that *dynamically* rewrite diagrams in contexts of arbitrary depth (Figure 1a), thus combining features of both deep inference and string diagrams[21]. Scroll nets provide a *static* way to represent proofs in EGs by recording explicitly applications of inference rules in a *directed graph* (Figure 1b), similarly to proof nets and combinatorial proofs. Crucially, this graph shares the same nodes as the (tree-shaped) statements involved in the proof, making scroll nets a more compact representation than other graphical proof structures, but also surprisingly a variant of the notion of *bigraph*. Bigraphs were introduced by Milner as a foundational combinatorial structure encompassing most models of concurrent/parallel computation, including Petri nets and his own CCS and π-calculus[22]. In his preprint, the ER shows how scroll nets naturally subsume the simply-typed λ-calculus — the common kernel of all type theories used in ITPs — by identifying a generalization of the notion of *detour* from natural deduction that abstracts away the shape of formulas.

All these discoveries hint toward the potential of scroll nets as a very expressive framework for both proof theory and type theory, unifying features found in most of the formalisms that have emerged from the two disciplines in a natural and efficient way. Moreover, the ER's vision is to exploit the diagrammatic nature of scroll nets to redesign not only the *backend* but also the *frontend* of ITPs, making them the *interaction substrate*[23] of a new kind of GUI in the PbA paradigm.

Given the early stage of the theory of scroll nets and the expertise of the Theory of Computation group at Birmingham hosting the ER, the SCROLLNETS project will progress towards this vision by focusing on the following key objectives:

[20]Pablo Donato, "Scroll Nets", July 2025, https://doi.org/10.48550/arXiv.2507.19689,

[21]Bonchi, Filippo, et al. "Diagrammatic Algebra of First Order Logic". *Proceedings of the 39th Annual ACM/IEEE Symposium on Logic in Computer Science*, LICS '24, 2024. https://doi.org/10.1145/3661814.3662078.

[22]Milner, Robin. "Bigraphical Reactive Systems". *Proceedings of the 12th International Conference on Concurrency Theory*, CONCUR '01, Aug. 2001, pp. 16–35. Accessed 1 July 2025.

[23]Mackay, Wendy E., and Michel Beaudouin-Lafon. "Interaction Substrates: Combining Power and Simplicity in Interactive Systems". *Proceedings of the 2025 CHI Conference on Human Factors in Computing Systems*, Apr. 2025, pp. 1–16. https://doi.org/10.1145/3706598.3714006.

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> (K1) extend the theory of scroll nets to account for **richer logics**, beyond minimal implicative logic;
>
> (K2) find natural and efficient **translations** of state-of-the-art proof systems and typed programming models into scroll nets.

121   This should support the following very ambitious long-term goal of the project:

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> *Establish **scroll nets** as the foundation for a new generation of ITPs that support **interactive refinement** of formal specifications, proofs and programs through **diagrammatic manipulations**.*

1.2   *Soundness of the proposed methodology (including interdisciplinary approaches, consideration of the gender dimension and other diversity aspects if relevant for the research project, and the quality of open science practices)*

- WP1 Study of sequentialization and normalization in implicative scroll nets

- WP2 Horizontal and vertical generalization to account for disjunction/sum types, bi-intuitionistic and classical logic

- The Functional Machine Calculus (FMC) and Relational Machine Calculus (RMC) are two foundational models of computation introduced very recently by W. Heijltjes, also co-inventor of intuitionistic combinatorial proofs[24]. They achieve respectively a unification of the two main paradigms of declarative programming, *functional* and *logic* programming, with the more mainstream paradigm of *imperative* programming. An open problem is to find a suitable combination of the FMC and RMC that could subsume all three paradigms. WP3 will explore the potential of scroll nets as a solution in the well-typed fragment, by devising translations of the FMC and RMC that preserve both its denotational and operational semantics. This will build on WP1 for the operational semantics given by detour elimination, and WP2 to account for non-determinism with $(n, 1)$-ary scrolls for sum and empty types.

- WP4 Recursive generalization to account for (co)inductive types (cyclic proofs)

- WP5 Study of Beta and Gamma to account for modal logic, higher-order logic and dependent types

1.3   *Quality of the supervision, training and of the two-way transfer of knowledge between the researcher and the host*

143   At a minimum, address the following aspects:

- Describe the qualifications and experience of the supervisor(s). Provide information regarding the supervisors' level of experience on the research topic proposed and their track record of work, including main international collaborations, as well as the level of experience in supervising/training, especially at advanced level (i.e. PhD and postdoctoral researchers).

- Planned training activities for the researcher (scientific aspects, management/organisation, horizontal and key transferrable skills...).

- For *European Fellowships*: two-way transfer of knowledge between the researcher and host organisation.

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[24]Willem Heijltjes, "The Functional Machine Calculus", *Electronic Notes in Theoretical Informatics and Computer Science* vol. Volume 1 - Proceedings of... *arXiv*, arxiv.org/abs/2212.08177, 22 Feb. 2023, https://doi.org/10.46298/entics.10513, p. 10513; Barrett, Chris, et al. "The Relational Machine Calculus". *Proceedings of the 39th Annual ACM/IEEE Symposium on Logic in Computer Science*, 8 July 2024, pp. 1–15. https://doi.org/10.1145/3661814.3662091; Heijltjes, Willem B., et al. "Intuitionistic proofs without syntax". *2019 34th Annual ACM/IEEE Symposium on Logic in Computer Science (LICS)*, June 2019, pp. 1–13. https://doi.org/10.1109/LICS.2019.8785827.

- For *Global Fellowships*: three-way transfer of knowledge between the researcher, host organisation, and associated partner for outgoing phase.

- Rationale and added-value of the non-academic placement (if applicable) and secondment (if applicable).

Employers and/or funders should ensure that a person is clearly identified to whom researchers can refer for the performance of their professional duties, and should inform the researchers accordingly.

Such arrangements should clearly define that the proposed supervisors are sufficiently expert in supervising research, have the time, knowledge, experience, expertise and commitment to be able to offer the postdoctoral researcher appropriate support and provide for the necessary progress and review procedures, as well as the necessary feedback mechanisms.

Supervision is one of the crucial elements of successful research. Guiding, supporting, directing, advising and mentoring are key factors for a researcher to pursue his/her career path. In this context, all MSCA-funded projects are encouraged to follow the recommendations outlined in the MSCA Guidelines on Supervision [25].

1.4 Quality and appropriateness of the researcher's professional experience, competences and skills

Discuss the quality and appropriateness of the researcher's existing professional experience in relation to the proposed research project.

2 Impact #@IMP-ACT-IA@#

2.1 Credibility of the measures to enhance the career perspectives and employability of the researcher and contribution to his/her skills development

At a minimum, address the following aspects:

- Specific measures to enhance career perspectives and employability of the researcher inside and/or outside academia.

- *Expected* contribution of proposed skills development to the future career of the researcher.

2.2 Suitability and quality of the measures to maximise expected outcomes and impacts, as set out in the dissemination and exploitation plan, including communication activities #@COM-DIS-VIS-CDV@#

At a minimum, address the following aspects:

- *Plan for the dissemination and exploitation activities, including communication activities*[26]: Describe the planned measures to maximize the impact of your project by providing a first version of your 'plan for the dissemination and exploitation including communication activities'. Describe the dissemination, exploitation measures that are planned, and the target group(s) addressed (e.g. scientific community, end users, financial actors, public at large). Regarding communication measures and public engagement strategy, the aim is to inform and reach out to society and show the activities performed, and the use and the benefits the project will have for citizens. Activities must be strategically planned, with clear objectives, start at the outset and continue through the lifetime of the project. The

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[25]While the MSCA Guidelines on Supervision are non-binding, funded-projects are strongly encouraged to take them into account.

description of the communication activities needs to state the main messages as well as the tools and channels that will be used to reach out to each of the chosen target groups.

- *Strategy for the management of intellectual property, foreseen protection measures*: if relevant, discuss the strategy for the management of intellectual property, foreseen protection measures, such as patents, design rights, copyright, trade secrets, etc., and how these would be used to support exploitation.

- All measures should be proportionate to the scale of the project, and should contain concrete actions to be implemented both during and after the end of the project.

2.3 The magnitude and importance of the project's contribution to the expected scientific, societal and economic impacts

- Provide a narrative explaining how the project's results are expected to make a difference in terms of impact, beyond the immediate scope and duration of the project. The narrative should include the components below, tailored to your project.

- Be specific, referring to the effects of your project, and not R&I in general in this field. State the target groups that would benefit.

- The impacts of your project may be:

  - *Scientific*: e.g. contributing to specific scientific advances, across and within disciplines, creating new knowledge, reinforcing scientific equipment and instruments, computing systems (i.e. research infrastructures);

  - *Economic/technological*: e.g. bringing new products, services, business processes to the market, increasing efficiency, decreasing costs, increasing profits, contributing to standards' setting, etc.

  - *Societal*: e.g. decreasing CO_2 emissions, decreasing avoidable mortality, improving policies and decision-making, raising consumer awareness.

- Only include such outcomes and impacts where your project would make a significant and direct contribution. Avoid describing very tenuous links to wider impacts.

- Give an indication of the magnitude and importance of the project's contribution to the expected outcomes and impacts, should the project be successful. Provide credible quantified estimates where possible and meaningful.

"Magnitude" refers to how widespread the outcomes and impacts are likely to be. For example, in terms of the size of the target group, or the proportion of that group, that should benefit over time.

"Importance" refers to the value of those benefits. For example, number of additional healthy life years; efficiency savings in energy supply.

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[26]In case your proposal is selected for funding, a more detailed Dissemination and Exploitation plan will need to be provided as a mandatory project deliverable during project implementation.

3 Quality and Efficiency of the Implementation #@QUA-LIT-QL@# #@WRK-PLA-WP@# #@CON-SOR-CS@#
#@PRJ-MGT-PM@#

3.1 Quality and effectiveness of the work plan, assessment of risks and appropriateness of the effort assigned to work packages

At a minimum, address the following aspects:

- Brief presentation of the overall structure of the work plan, including deliverables and milestones.

- Timing of the different work packages and their components;

- Mechanisms in place to assess and mitigate risks (of research and/or administrative nature).

A Gantt chart must be included and should indicate the proposed Work Packages (WP), major deliverables, milestones, secondments, placements, if applicable. This Gantt chart counts towards the 10-page limit.

The schedule in the Gantt chart should indicate the number of months elapsed from the start of the action (Month 1, Month 2, etc.), but no actual dates.

3.2 Quality and capacity of the host institutions and participating organisations, including hosting arrangements

At a minimum, address the following aspects:

- Hosting arrangements, including integration in the team/institution(s) and support services available to the researcher.

- Quality and capacity of the participating organisations, including infrastructure, logistics and facilities. Additional information should be outlined in Part B-2 Section 5 ("Capacity of the Participating Organisations").

Note that for GF, both the quality and capacity of the outgoing Third Country host and the return host should be outlined.

Associated partners linked to a beneficiary [27]

If applicable, outline here the involvement of any 'associated partners linked to a beneficiary' (in particular, the name of the entity, the type of link with the beneficiary and the tasks to be carried out).

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[27] See the definitions section of the MSCA Work Programme for further information.