COMP3011: Design and Analysis of Algorithms

Assignment 2

Due: 31 Oct 2024

The Hong Kong Polytechnic University

Problem 1 (25 points)

Our TA suggests the following variant on mergesort: instead of splitting the list into two halves, we split it into three thirds. Then we recursively sort each third and merge them.

Mergesort3($A[0, \cdots, n-1]$):

- 1. If $n \leq 1$, then return $A[0, \dots, n-1]$.
- 2. Let $k = \lceil n/3 \rceil$ and $m = \lfloor 2n/3 \rfloor$.
- 3. Return Merge3(Mergesort3($A[0, \dots, k-1]$), Mergesort3($A[k, \dots, m-1]$), Mergesort3($A[m, \dots, n-1]$)).

Merge3(L_0, L_1, L_2):

1. Return $Merge(L_0, Merge(L_1, L_2))$.

Assume that you have a subroutine Merge that merges two sorted lists of lengths l, l' in time O(l + l'). You may assume that n is a power of three, if you wish.

- (a) What is the asymptotic running time for executing Merge3(L_0, L_1, L_2), if L_0, L_1, L_2 are three sorted lists each of length n/3? Express your answer using O() notation.
- (b) Let T(n) denote the running time of Mergesort3 on an array of size n. Write a recurrence relation for T(n).
- (c) Solve the recurrence relation in part (b). Express your answer using O() notation.

Problem 2 (25 points)

Recall the Longest Common Subsequence problem, where we are given two strings, S_1 and S_2 , the task is to find the length of the Longest Common Subsequence, i.e. longest subsequence present in both of the strings. A longest common subsequence (LCS) is defined as the longest subsequence which is common in all given input sequences. Use the dynamic programming algorithm to compute the LCS between S_1 = "quantity" and S_2 = "inequality". You need to present the intermediate results.

Problem 3 (25 points)

There are n points on a road you are driving your taxi on. The n points on the road are labeled from 1 to n in the direction you are going, and you want to drive from point 1 to the direction of point n to make money by picking up passengers. You cannot change the direction of the taxi. The passengers are represented by an array rides, where $rides[i] = [start_i, end_i, tip_i]$ denotes the ith passenger requesting a ride from point $start_i$ to point end_i who is willing to give a tip_i dollar tip. We have $end_i \geq start_i$ for all i. For each passenger i you pick up, you earn $end_i - start_i + tip_i$ dollars. You may only drive at most one passenger at a time. All the numbers are positive integers.

Given n and the array rides, design an $O(n \log n)$ -time algorithm to return the maximum number of dollars you can earn by picking up the passengers optimally.

Note: You may drop off a passenger and pick up a different passenger at the same point.

Problem 4 (25 points)

We are given n rectangles, where each rectangle can be described by two integers a > 0 and b > 0 (i.e., the length and width respectively). Rectangle X(a, b) can be put inside of rectangle Y(c, d) if and only if (1) a < c and b < d or (2) b < c and a < d (i.e., the rectangle can be rotated 90°). For example, (1,5) can be put inside of (6,2) but not inside of (3,4). The task is to find the maximum number of rectangles that can be arranged in a line so that every rectangle (except the last one) can be put inside of the next one.

Design a polynomial-time algorithm for the above problem, and justify the correctness and the running time of the algorithm.