Lab 2

Lexing and Parsing with ANTLR4

Objective

- Discover typing in python with Pyright.
- Understand the software architecture of ANTLR4.
- Be able to write simple grammars and correct grammar issues in ANTLR4.
- Write a first evaluator under the form of semantic actions in ANTLR4.
- Optionally, deposit your work on https://etudes.ens-lyon.fr and get feedback (see instructions at the end).

EXERCISE #1 **► Lab preparation**

In the cap-lab23 directory¹:

```
git commit -a -m "my changes to LAB1" #push is not allowed git pull
```

will provide you all the necessary files for this lab in TP02. You also have to install ANTLR4. For tests, we will use pytest, and for printing trees we will use graphviz, you may have to install them:

```
sudo apt install graphviz
python3 -m pip install pytest pytest-cov pytest-xdist graphviz
```

If the pip command fails, you need to install pip on your machine and re-launch the command, on Ubuntu this is done with sudo apt install python3-pip, see https://pip.pypa.io/en/stable/installation/otherwise.

2.1 Typing in python with Pyright

Python is originally a dynamically typed language, i.e. types are associated with values at runtime, but no check is done statically. Recent versions of Python, however, allow adding static typing annotations in the code. You don't *need* to add these annotations, but having them allows static typecheckers like Mypy or Pyright to find errors in the code before running it. In this course, we will use Pyright. You don't need to understand all the details of typing annotations, but you need to get familiar with them to understand the code provided to you.

If needed, install pyright. It is provided in the archive and the Docker image, so you probably have nothing to do.

EXERCISE #2 ► Basic type checking

Consider the following code (available as TP02/python/typecheck.py in your repository):

```
# Typing annotations for variables:
# name: type
int_variable: int
float_variable: float
int_variable = 4.2 # Static typing error, but no runtime error
float_variable = 42.0 # OK
float_variable = int_variable # OK
# Typing annotations for functions (-> means "returns")
```

if you don't have it already, get it from https://github.com/Drup/cap-lab23.git

```
def int_to_string(i: int) -> str:
   return str(i)
print(int_to_string('Hello')) # Static typing error, but no runtime error
print(int_to_string(42) / 5) # Both static and runtime error
Run the code in the Python interpreter:
python3 typecheck.py
Check for static typing errors using pyright:
pyright typecheck.py
Note how static typechecking finds more error in your code, and possibly saves you a lot of debbugging time.
Try modifying the code to get rid of typing errors.
EXERCISE #3 ► Type unions
Play with (execute, typecheck, read the comments) the following code (available as TP02/python/type_unions.py
in your repository):
from typing import List
# int | float means ``either an int or a float''.
NUMBER = int | float # or Union[int, float] with Union imported from typing
def add_numbers(a: NUMBER, b: NUMBER) -> NUMBER:
   return a + b
# Both int and floats can be passed to the function
print(add_numbers(1, 4.3))
def divide_numbers(a: NUMBER, b: NUMBER) -> float:
   return a / b
print(divide_numbers(1, 2))
# Declare the type of a list whose elements are numbers.
LIST_OF_NUMBERS = List[NUMBER]
def increment(a: LIST_OF_NUMBERS) -> LIST_OF_NUMBERS:
   return [x + 1 \text{ for } x \text{ in a}]
print(increment([1, 2, 3]))
# Skip the end if you are late.
# The type DEEP_LIST_OF_NUMBERS is a special case since it references itself.
# The identifier DEEP_LIST_OF_NUMBERS cannot be used before the end of its
# initialization, but the circular dependency can be broken using the string
# 'DEEP_LIST_OF_NUMBERS' instead.
DEEP_LIST_OF_NUMBERS = NUMBER | List['DEEP_LIST_OF_NUMBERS']
```

```
def deep_increment(d: DEEP_LIST_OF_NUMBERS) -> DEEP_LIST_OF_NUMBERS:
    if isinstance(d, list):
        # Note the unusual typing rule applied by Pyright here: because we are
        # in the 'isinstance(d, list)' branch, it knows that d is a list,
        # and accepts to iterate over it.
        return [deep_increment(e) for e in d]
    else:
        # ... and here, in the 'else' branch Pyright knows that d is not a list,
        # and can deduce that it is a NUMBER.
        return d + 1
```

The syntax int | float (or equivalently Union[int, float]) means "either an int or a float"; List[NUMBER] means "a list whose elements are numbers".

2.2 User install for ANTLR4 and ANTLR4 Python runtime

2.2.1 User installation

EXERCISE #4 ► Install

To be able to use ANTLR4 for the next labs, download it and install the python runtime (not needed if you use the Docker image, and ANTLR is already in /opt/antlr-4.11.1-complete.jar if you use the precompiled RiscV toolchain):

```
mkdir ~/lib
cd ~/lib
wget https://www.antlr.org/download/antlr-4.11.1-complete.jar
python3 -m pip install antlr4-python3-runtime==4.11.1
    Then add to your ~/.bashrc:
export ANTLR4="java -jar $HOME/lib/antlr-4.11.1-complete.jar"
alias antlr4="java -jar $HOME/lib/antlr-4.11.1-complete.jar"
    Then source your .bashrc:
source ~/.bashrc
    You may need to install java on your machine:
sudo apt install default-jre
Tests will be done in Section 2.4.
```

2.3 Simple examples with ANTLR4

2.3.1 Structure of a .g4 file and compilation

Links to a bit of ANTLR4 syntax:

- Lexical rules (extended regular expressions): https://github.com/antlr/antlr4/blob/master/doc/lexer-rules.md
- Parser rules (grammars) https://github.com/antlr/antlr4/blob/master/doc/parser-rules.md

The compilation of a given .g4 (for the PYTHON back-end) is done by the following command line if you modified your .bashrc properly (note: antlr4, not antlr which may also exist but is not the one we want):

```
antlr4 -Dlanguage=Python3 filename.g4
```

If you did not define the alias or if you installed the . jar file to another location, you may also use:

```
java -jar /path/to/antlr-4.11.1-complete.jar -Dlanguage=Python3 filename.g4
```

2.3.2 **Up to you!**

EXERCISE #5 **▶ Demo files**

Work your way through the three examples (open them in your favorite editor!) in the directory demo_files:

ex1: lexer grammar and PYTHON **driver** A very simple lexical analysis² for simple arithmetic expressions of the form x+3. To compile, run:

```
antlr4 -Dlanguage=Python3 Example1.g4
```

(or type make).

This generates a lexer in Example1.py (you may look at its content, and be happy you didn't have to write it yourself) plus some auxiliary files. We provide you a simple main.py file that calls this lexer (this one is hand-written and readable):

```
python3 main.py
```

(or type make run, which re-generates the lexer as needed and runs main.py).

To signal the program you have finished entering the input, use **Control-D** (you may need to press it twice). Examples of runs: [^D means that I pressed Control-D]. What I typed is in boldface.

```
1+1
^D^D

[@0,0:0='1',<2>,1:0]
[@1,1:1='+',<1>,1:1]
[@2,2:2='1',<2>,1:2]
[@3,4:3='<EOF>',<EOF>,2:0]
)+
^D^D

line 1:0 token recognition error at: ')'
[@0,1:1='+',<1>,1:1]
[@1,3:2='<EOF>',<-1>,2:0]
%
```

Questions:

- Reproduce the above behavior.
- Read and understand the code in Example1.g4. Observe the outputs above.
- Modify the source file to allow for parentheses to appear in the input.
- What is an example of a recognized expression that looks odd (i.e. that is not a real arithmetic expression)? To fix this problem we need a syntactic analyzer (see later).
- Observe the PYTHON main.py file.

From now on you can alternatively use the commands make and make run instead of calling antlr4 and python3.

ex2: full grammar (**lexer + parser**) **and** PYTHON **driver** Now we write a grammar for valid expressions. Observe how we recover information from the lexing phase in the ID rule (for ID, the associated text is \$ID.text). More generally, we can associate to each production a piece of Python code that will be executed each time the production is reduced. This piece of code is called a *semantic action* and computes attributes of non-terminals.

If these files read like a novel, go on with the other exercises. Otherwise, make sure that you understand what is going on. You can ask the Teaching Assistant, or another student, for advice.

²Lexer Grammar in ANTLR4 jargon

ex3: full grammar with attributes and PYTHON **driver** We take the exact same grammar as for ex2, but modify its semantic actions so as to compute the number of operators in an expression. Observe how to define an attribute in the grammar with **expr returns** [type name], and how to propagate it through the syntax with actions.

As before, once you understand these files, go on with the other exercises.

From now you will write your own grammars. Be careful the ANTLR4 syntax use unusual conventions: "Parser rules start with a lowercase letter and lexer rules with an upper case." a

 $^a \verb|https://stackoverflow.com/questions/11118539/antlr-combination-of-tokens|$

EXERCISE #6 ► Well-founded parentheses

Write a grammar and files to make an analyser that:

- skips all characters but '(', ')', '[', ']' (use the parser rule CHARS: ~[()[\]] -> skip; for it)
- · accepts well-formed parentheses.

Thus your analyser will accept "(hop)" or "[()](tagada)" but reject "plop]" or "[)". Test it on well-chosen examples. Begin with a proper copy of ex2 (usemake clean), change the name of the files, name of the grammar, do not forget the main and the Makefile, and THEN, change the grammar to answer the exercise.

EXERCISE #7 ► Towards analysis: If then else ambiguity³ - Skip if you are late

We give you the following grammar for nested "ifs" (ITE/ directory):

```
grammar ITE;
prog: stmt EOF;
stmt : ifStmt | ID ;
ifStmt : 'if' ID 'then' thenstmt=stmt ('else' elsestmt=stmt)?;

ID : [a-zA-Z]+;
WS : [ \t\r\n]+ -> skip ; // skip spaces, tabs, newlines

Find a way (with the right semantic actions) to test if:
if x then if y then a else b
is parsed as:
if x then (if y then a else b)

or
if x then (if y then a) else b
```

Is it the same with this grammar as with real-life C-like languages?

Notice the thenstmt=stmt syntax to give unique names to the two stmts of the above ifStmt grammar rule.

There are more explicit ways of getting around the dangling else problem by adapting the grammar. For more details see https://en.wikipedia.org/wiki/Dangling_else

2.4 Grammar Attributes (actions), ariteval/directory

Until now, our analyzers were passive oracles, i.e. language recognizers. Moving towards a "real compiler", a next step is to execute code during the analysis, using this code to produce an intermediate representation of the recognized program, typically ASTs. This representation can then be used to generate code or perform program analysis (see next labs). This is what *attribute grammars* are made for.

³Also known as "dangling else"

Important remark As Python is sensitive to indentation, there might be some issues when writing semantic actions on several lines. You can often avoid the problem by defining a function in the Python header and then calling it in the right-hand side of the rules.

The goal of this section is to implement a simple arithmetic evaluator, starting from the basic infrastructure given in the ariteval/directory.

EXERCISE #8 ► Test the provided code (ariteval/directory)

To test your installs:

```
1. Type
```

```
make ; python3 arit.py tests/hello01.txt
This should print:
prog = Hello
on the standard output.
```

2. Type:

make test

This should print:

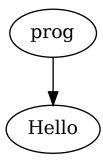
```
test_ariteval.py::TestEVAL::test_expect[./tests/hello01.txt] PASSED
```

To debug our grammar, we will use an interactive test (unlike make test which is fully automatic): display the parse tree in Lisp format, and check manually that it matches your expectation.

To help you, we provide a make print-tree target in the Makefile, which you can run:

```
make TESTFILE=tests/hello01.txt print-tree
```

It opens a window with a graphical representation of the parse tree:



The result is also stored in tree.dot.pdf.

Alternatively, replace print-tree with print-lisp in the above command to display a textual (LISP) representation of the parse tree:

(prog Hello)

We will now implement the following grammar (an expression followed by a semicolon):

```
Z \rightarrow E;

E \rightarrow E + T

E \rightarrow T

T \rightarrow T * F

T \rightarrow F

F \rightarrow id

F \rightarrow int

F \rightarrow (E)
```

Figure 2.1: A grammar for arithmetic expressions

EXERCISE #9 ► Implement!

In the ariteval/Arit.g4 file, implement the grammar from figure 2.1 in ANTLR4. Write test files and test them wrt. the grammar with make print-tree. In particular, verify the fact that '*' has higher priority on '+.' Is '+' left or right associative?

Important! Before you begin to implement, it is MANDATORY to read carefully until the end of the lab subject.

EXERCISE #10 ► Evaluating arithmetic expressions with ANTLR4 and PYTHON

Attribute the grammar you just wrote to evaluate arithmetic expressions with binary operators in $\{+,-,*,/\}$ with the usual priorities and print the result; / will be the integer division such that 4/3 = 1. For the moment, throw an error for all uses of variables:

```
if $ID.text not in idTab: # Always true, for now
  raise UnknownIdentifier($ID.text)
```

To test your grammar, run it against expressions such as 1+(2*3);. In this case the expected output is 1+(2*3) = 7.

In case of division by zero, do what you want but explain (see the next exercise for how to handle automated test for errors).

EXERCISE #11 ► Test infrastructure

We provide to you a test infrastructure. In the repository, we provide you a script that enables you to test your code. For the moment it only tests files of the form tests/hello*.txt.

To test your code on these files, just type:

make test

To test on more relevant files, you should open test_ariteval.py and change the value of ALL_FILES.

We will use the exact same script to test your code (but with our own test cases!).

A given test has the following behavior: if the pragma // EXPECTED is present in the file, it compares the actual output with the list of expected values (see tests/test01.txt for instance). There is also a special case for errors, with the pragma // EXITCODE n, that also checks the (non zero) return code n if there has been an error followed by an exit (see tests/bad01.txt).

EXERCISE #12 ► Adding features!

• Augment the grammar to treat lists of assignments (and expressions). You will use Python dictionaries to store values of ids when they are defined:

idTab[\$ID.text]=...

Line breaks should now be allowed between assignments (and expressions).

- Write tests for lists of assignements such as x=1;2+x;. When you read a variable that is not (yet) defined, you have to raise the UnknownIdentifier exception.
- Optionally, implement binary and unary minus. Be careful with operators precedence and associativity.
 Unary minus can apply to any expression (--1 is accepted, with the same meaning as in Python).
 Here are examples of expected outputs ⁴:

Input	Output (on stdout)
1;	1 = 1
-12;	-12 = -12
12;	12 = 12
1 + 2;	1+2 = 3
3 - 2;	3-2 = 1
1 + 2 * 3 + 4;	1+2*3+4 = 11
(1+2)*(3+4);	(1+2)*(3+4) = 21
a=1+4/1;	a now equals 5
b + 1;	b+1 = 43
a + 8;	a+8 = 13
-1 + x;	-1+x = 41
-(5+a);	-(5+a) = -10
3 - (-4);	3-(-4) = 7
34;	34 = 7
4/3;	4/3 = 1

The parsed expression can be printed from an expression, for instance with:

```
rule :
expr ... {print($expr.text)}
```

EXERCISE #13 ► Archive and test deposit

We do not ask you for a mandatory deposit. However we strongly encourage you to have an operational evaluator and above all, to understand the test infrastructure. This will be useful for the future labs!

If you want feedback, you can do a test deposit before 2023-10-04 23:59 (this will not be graded):

- Type make tar to obtain the archive to send. Your archive must also contain test cases in the tests/directory and a README.md with your name, the functionality of the code, how to use it, your design choices, and known bugs (probably less than 5 lines of text to add unless you did anything special).
- Go to https://etudes.ens-lyon.fr/course/view.php?id=5673 and deposit your archive.za

⁴The expected behavior of your evaluator may not be completely specified. If you make a design choice, explain it in the README.md file (division by 0, for instance).