

COSC 2436: Final Exam Review

Hashing #1

**Insert the following values into a hash table using linear probing.
Assume the hash table is of size 10.**

{54, 75, 24, 45, 18, 10}

Hashing #1

Insert the following values into a hash table using linear probing.
Assume the hash table is of size 10.

{54, 75, 24, 45, 18, 10}

0	1	2	3	4	5	6	7	8	9
10				54	75	24	45	18	

Hashing #2

The code below shows double hashing. What is wrong with the code?


```
73 ▼ void doubleHashing(int table[], int x, int tableSize){  
74 ▼     for(int i = 0; i < tableSize; i++){  
75         int index = (hash1(x, tableSize) + (i * hash2(x, 7))) % tableSize;  
76 ▼         if(table[index] == -1){  
77             table[index] = x;  
78         }  
79     }  
80 }
```

Hashing #2

The code below shows double hashing. What is wrong with the code?

There should be a *break* statement after line 77. If there is no break statement, *x* will keep getting added to the table.

```
73 ▼ void doubleHashing(int table[], int x, int tableSize){  
74 ▼   for(int i = 0; i < tableSize; i++){  
75     int index = (hash1(x, tableSize) + (i * hash2(x, 7))) % tableSize;  
76 ▼     if(table[index] == -1){  
77       table[index] = x;  
78     }  
79   }  
80 }
```

 **add *break*;**

Hashing #3

Match the following hash function with its correct description.

- Direct Hashing _____
 - Linear Probing _____
 - Double Hashing _____
- a) data is overwritten during collision
 - b) clustering can occur
 - c) uses two hash functions

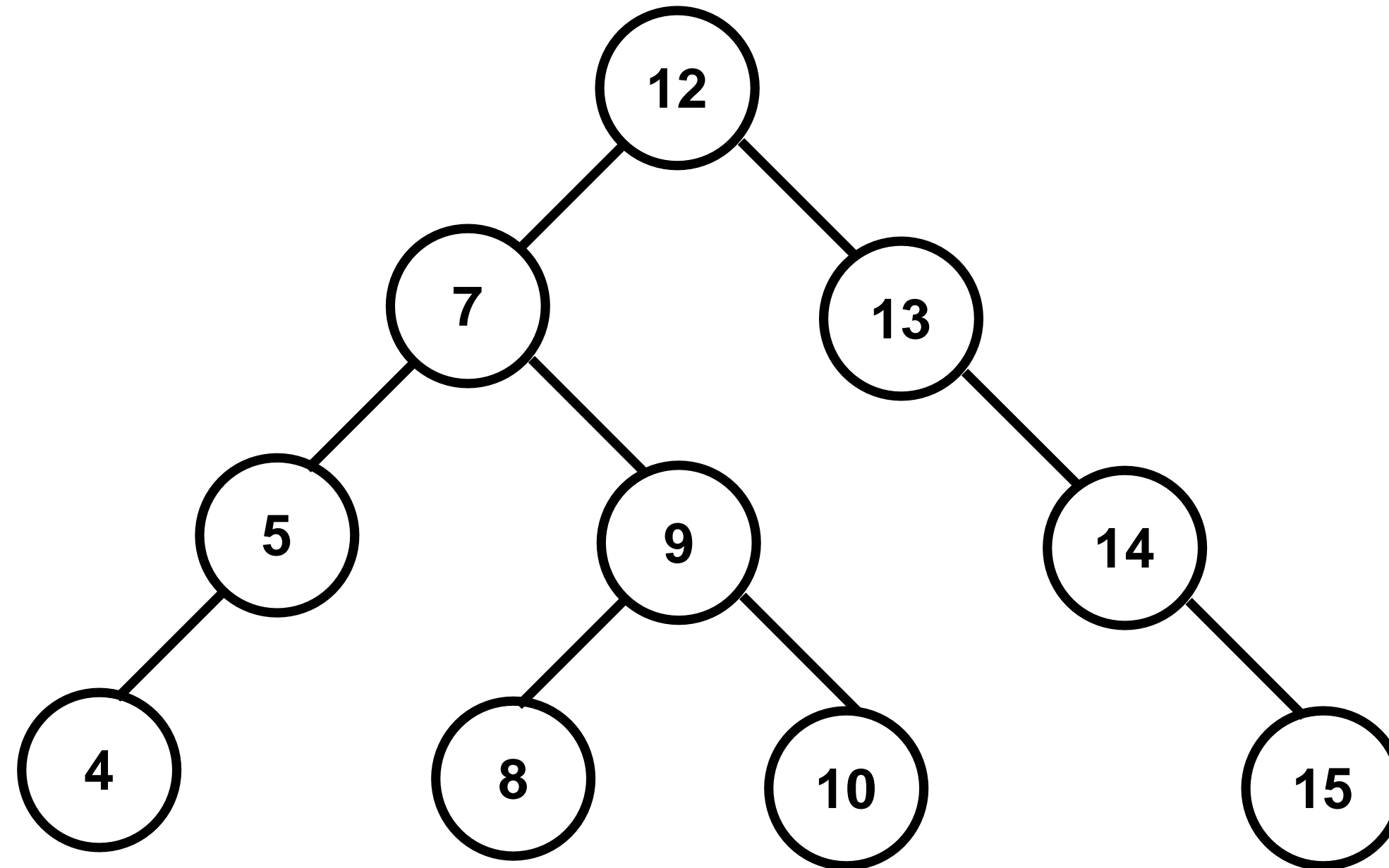
Hashing #3

Match the following hash function with its correct description.

- Direct Hashing **a**
 - Linear Probing **b**
 - Double Hashing **c**
- a) data is overwritten during collision
 - b) clustering can occur
 - c) uses two hash functions

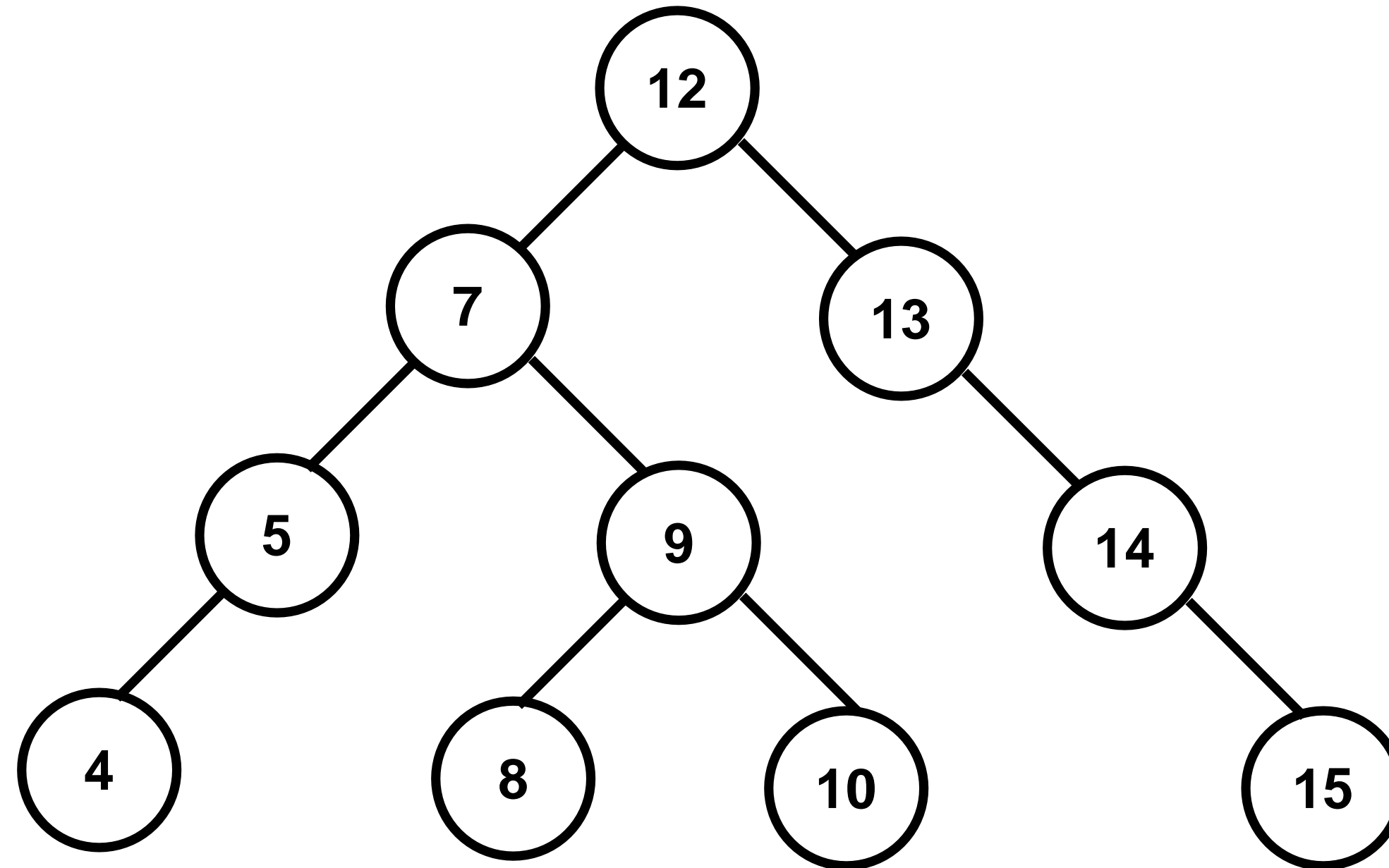
BST #1

Perform preorder traversal on the following BST.



BST #1

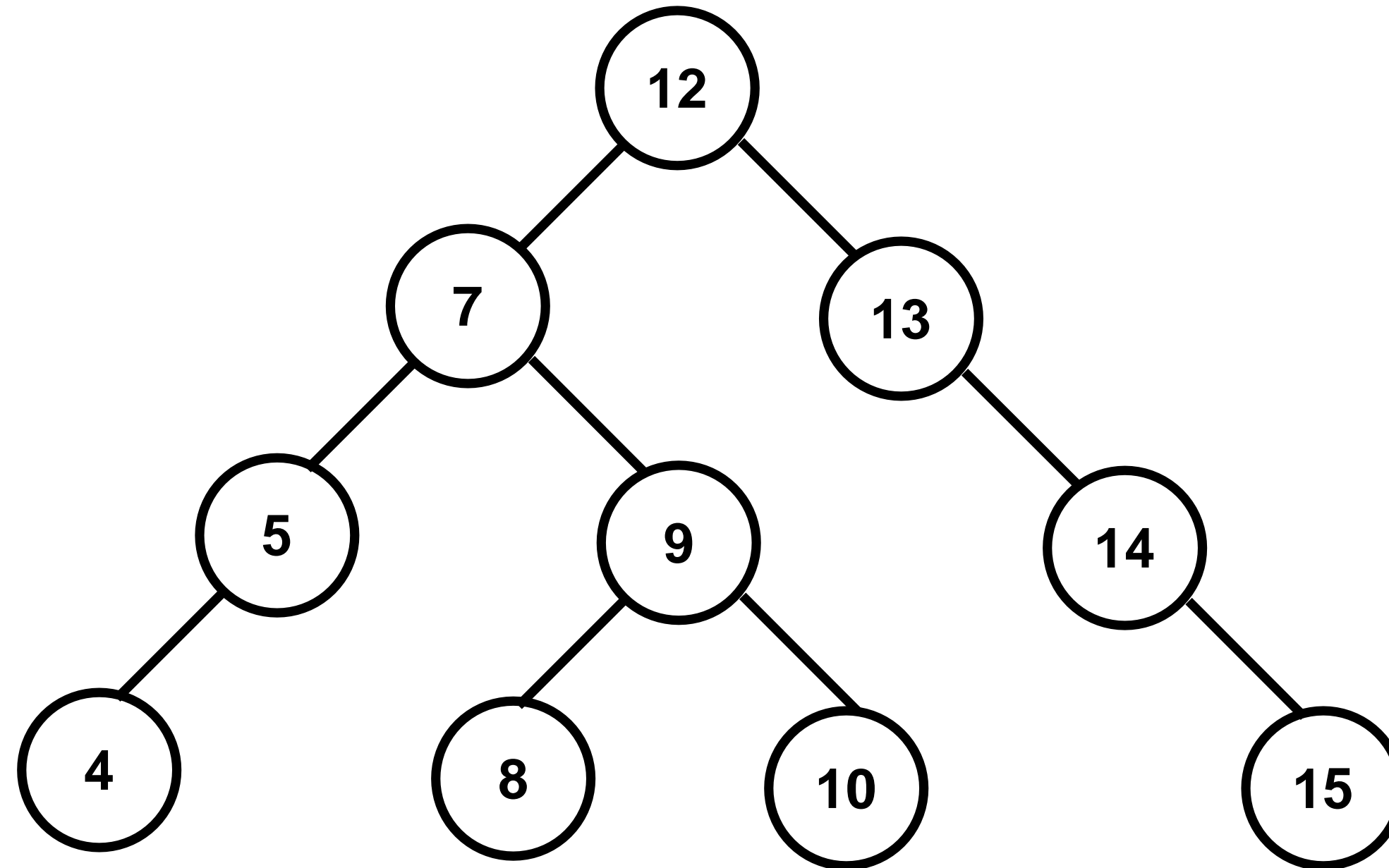
Perform preorder traversal on the following BST.



Preorder: 12 7 5 4 9 8 10 13 14 15

BST #2

Write the function which produced the output below.



Output: 4 5 8 10 9 7 15 14 13 12

BST #2

Write the function which produced the output below.

```
void postorder(node *n){  
    if(n == nullptr)  
        return;  
    postorder(n->left);  
    postorder(n->right);  
    cout << n->value << " ";  
}
```

BST #3

Write the function `getSum()` which returns the sum of all the values in a BST.

```
struct node{  
    int value;  
    node *right;  
    node *left;  
};
```

```
int getSum(node *n) {
```

```
}
```

BST #3

Write the function `getSum()` which returns the sum of all the values in a BST.

```
int getSum(node *n){  
    if(n == nullptr)  
        return 0;  
    return (n->value + getSum(n->right) + getSum(n->left));  
}
```

BST #4

Write the function leafCount() which returns the number of leafs in a BST.

```
struct node{
    int value;
    node *left;
    node *right;
};
```

```
int leafCount(node *root) {
```

BST #4

Write the function `leafCount()` which returns the number of leafs in a BST.

```
int leafCount(node *n){  
    if(n == nullptr)  
        return 0;  
    else if(n->left == nullptr && n->right == nullptr)  
        return 1;  
    else  
        return leafCount(n->left) + leafCount(n->right);  
}
```

BST #5

Which of the following is NOT a property of a BST:

- A) Best case time complexity is $O(\log(n))$
- B) Left-Child $<$ Root $<$ Right-Child
- C) Only contains unique values
- D) Can have more than 2 children
- E) None of the above

BST #5

Which of the following is NOT a property of a BST:

- A) Best case time complexity is $O(\log(n))$
- B) Left-Child $<$ Root $<$ Right-Child
- C) Only contains unique values
- D) Can have more than 2 children
- E) None of the above

D) Can have more than 2 children

AVL Tree #1

Perform the following AVL Tree commands.

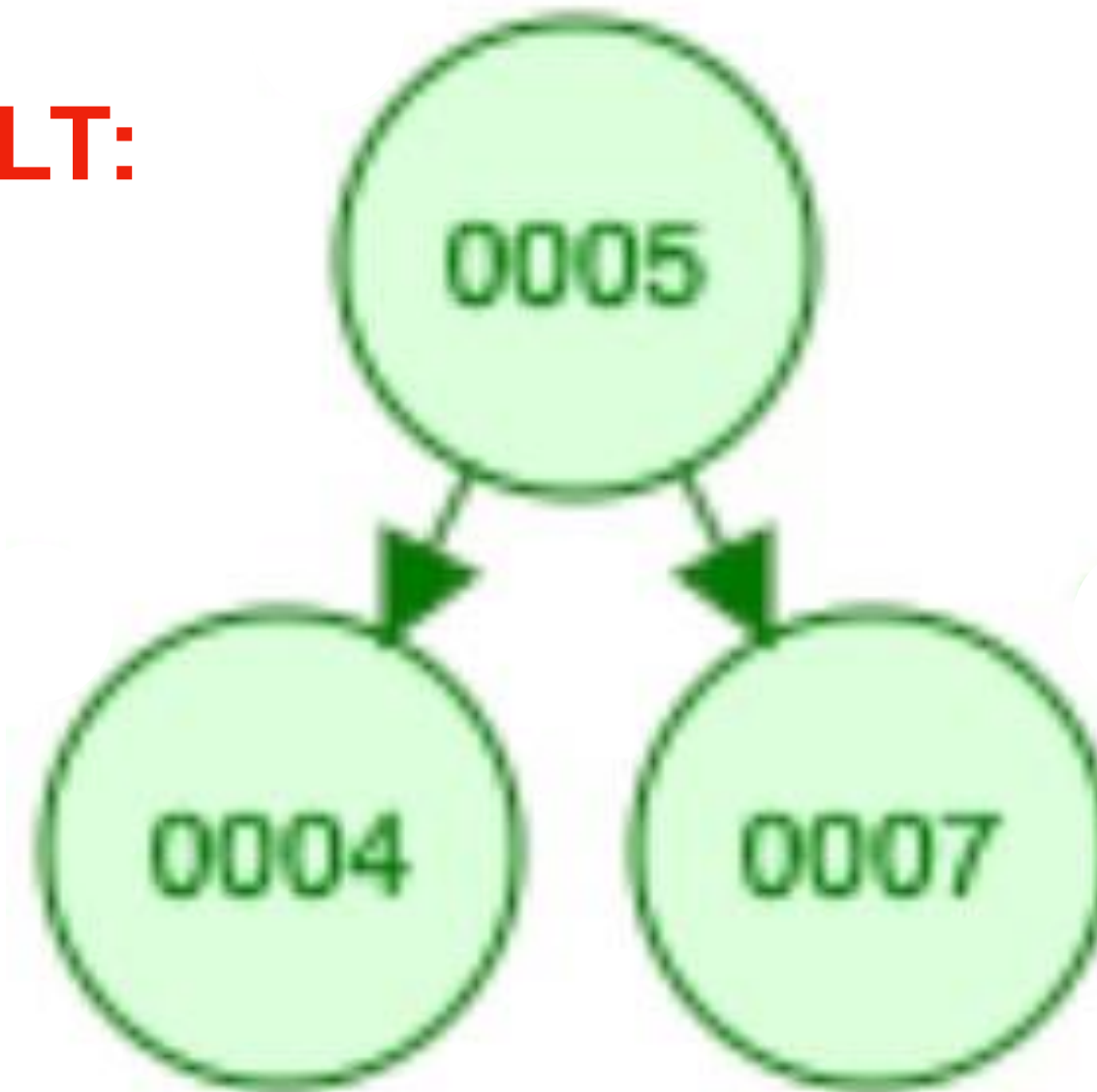
- **Insert(5)**
- **Insert(6)**
- **Insert(7)**
- **Insert(2)**
- **Insert(3)**
- **Insert(4)**
- **Delete(6)**
- **Delete(3)**
- **Delete(2)**

AVL Tree #1

Perform the following AVL Tree commands.

- Insert(5)
- Insert(6)
- Insert(7)
- Insert(2)
- Insert(3)
- Insert(4)
- Delete(6)
- Delete(3)
- Delete(2)

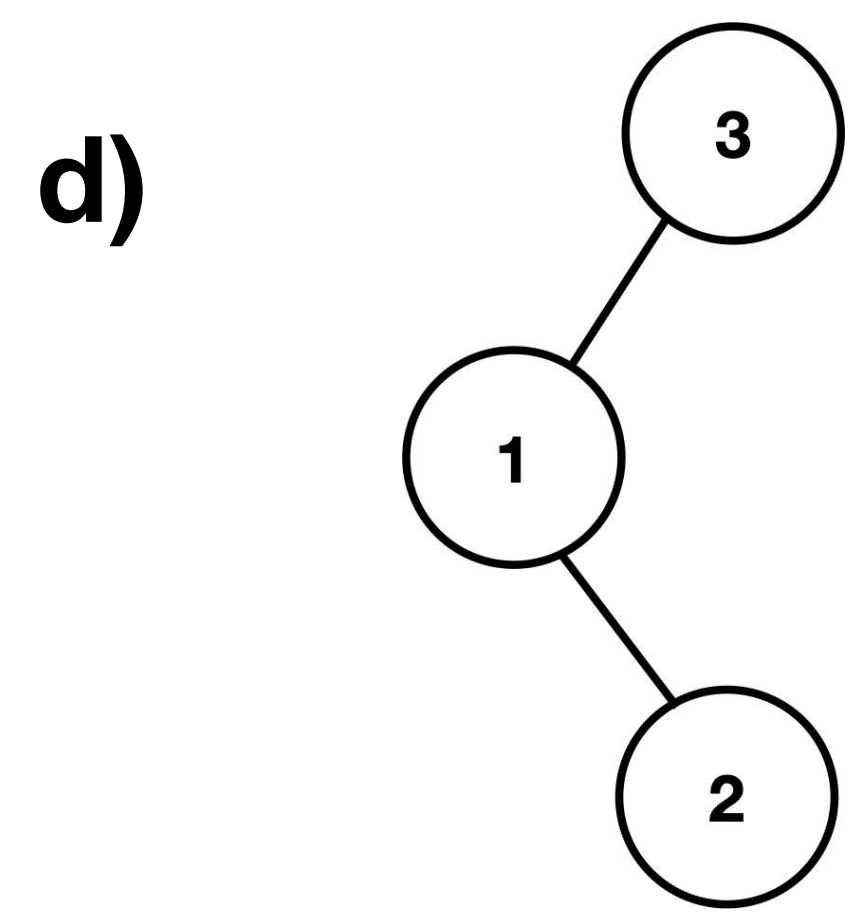
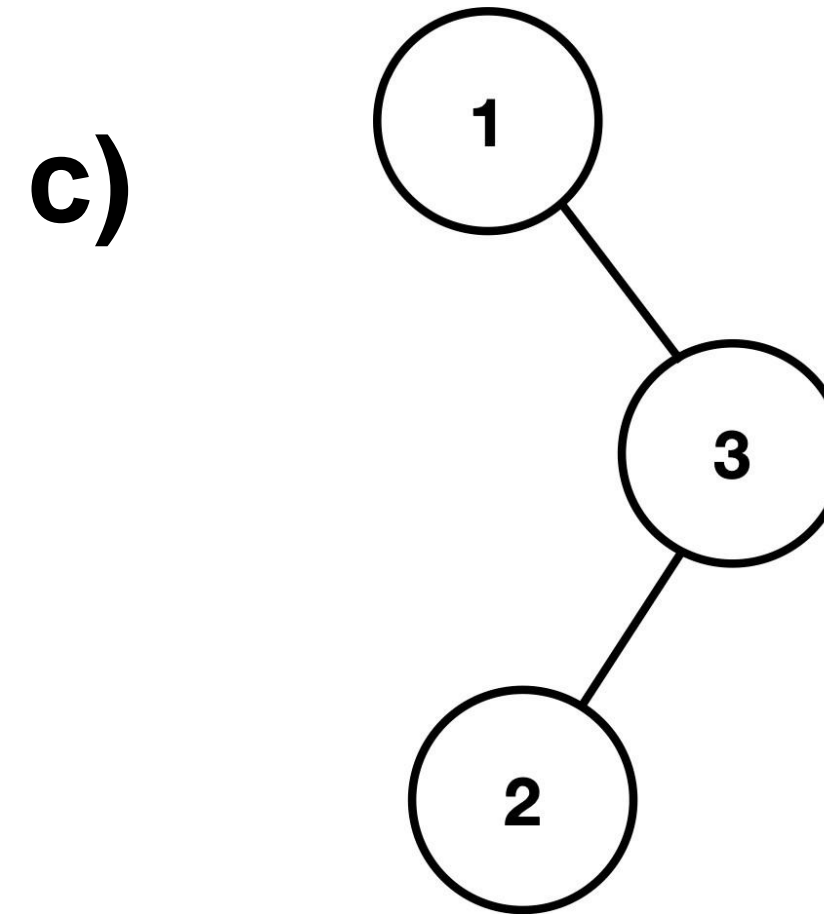
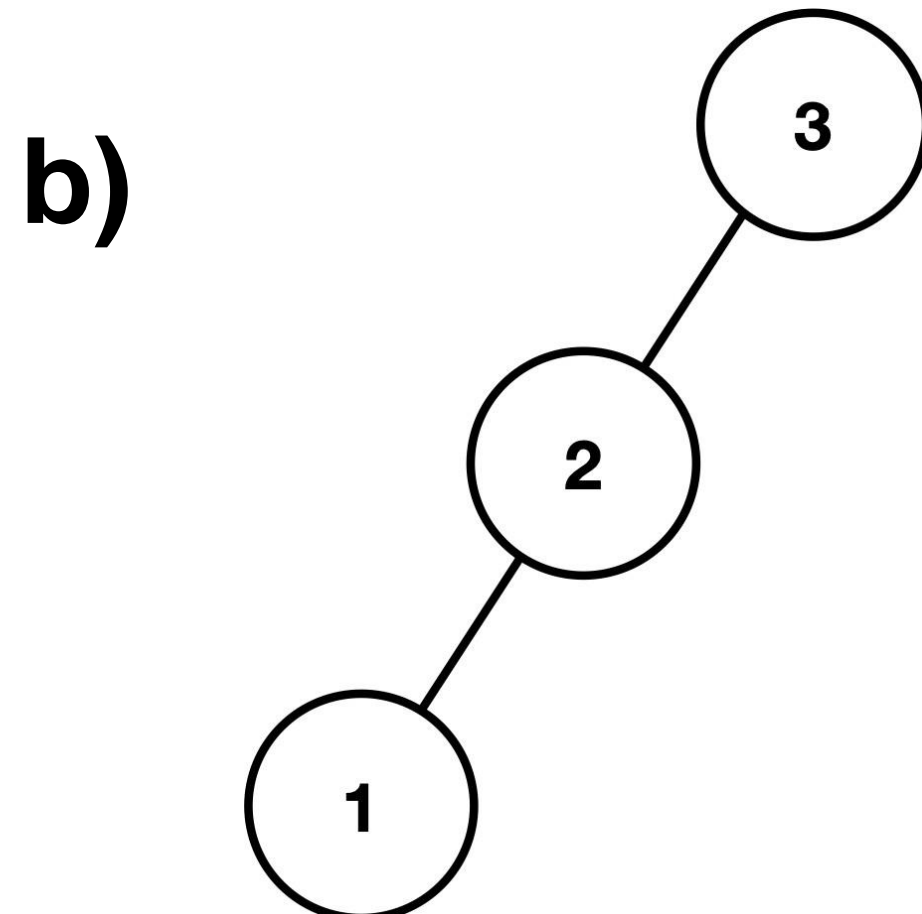
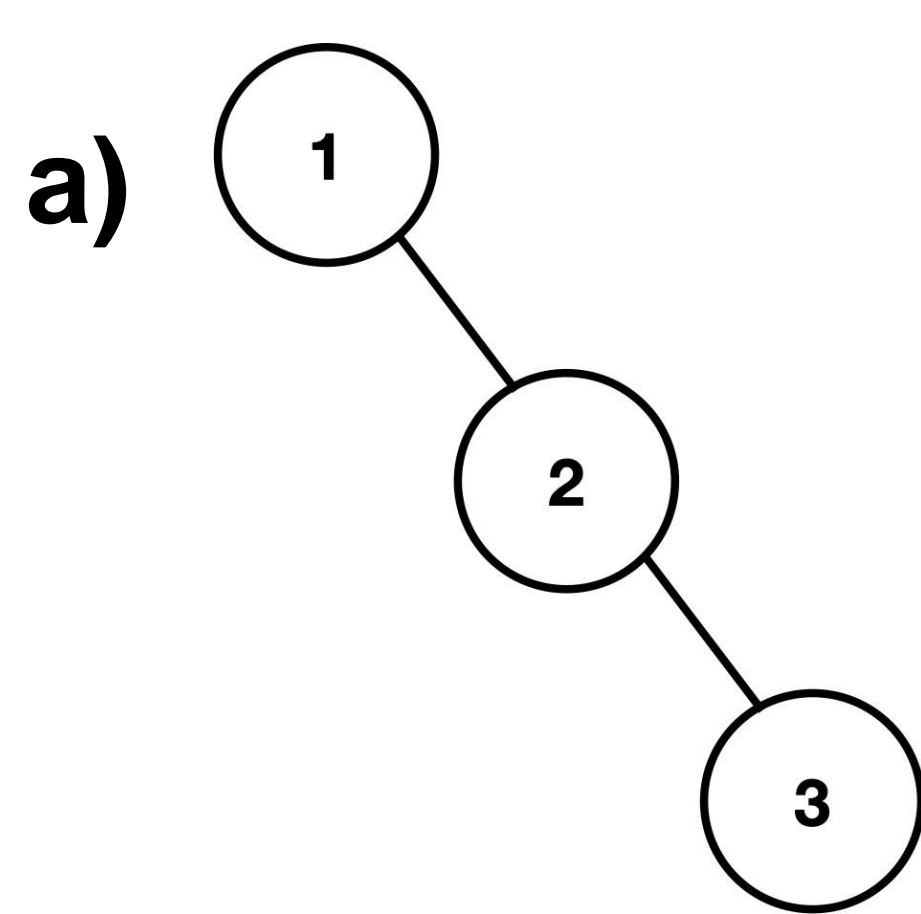
RESULT:



AVL Tree #2

Match the following picture with the rotation that should be performed:

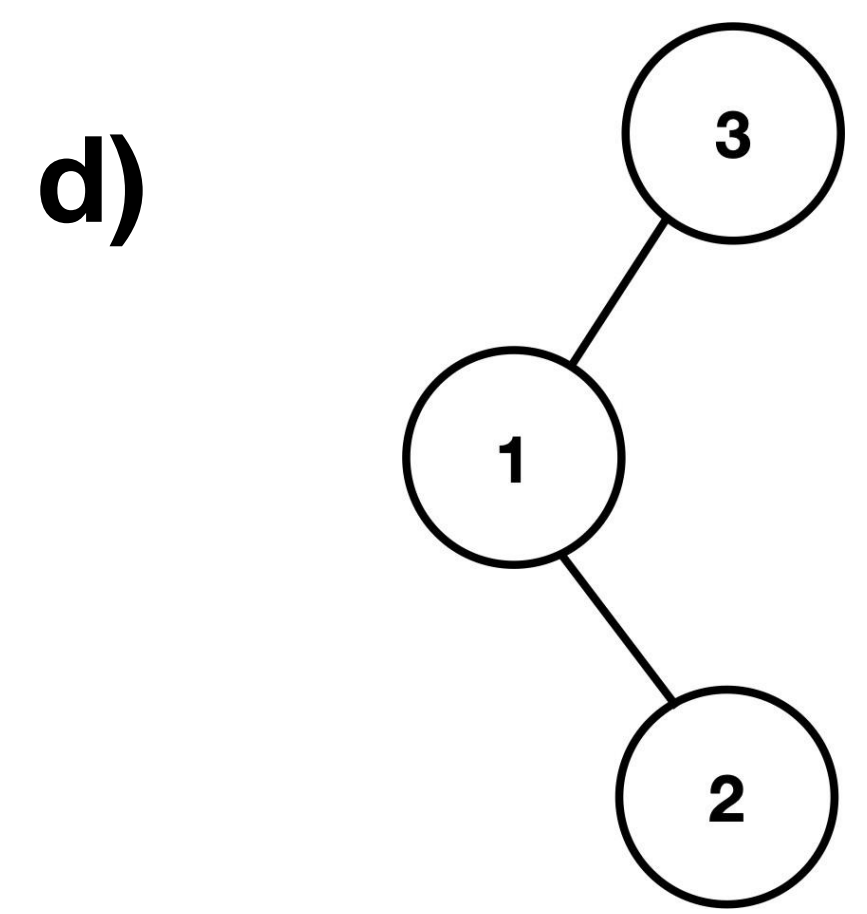
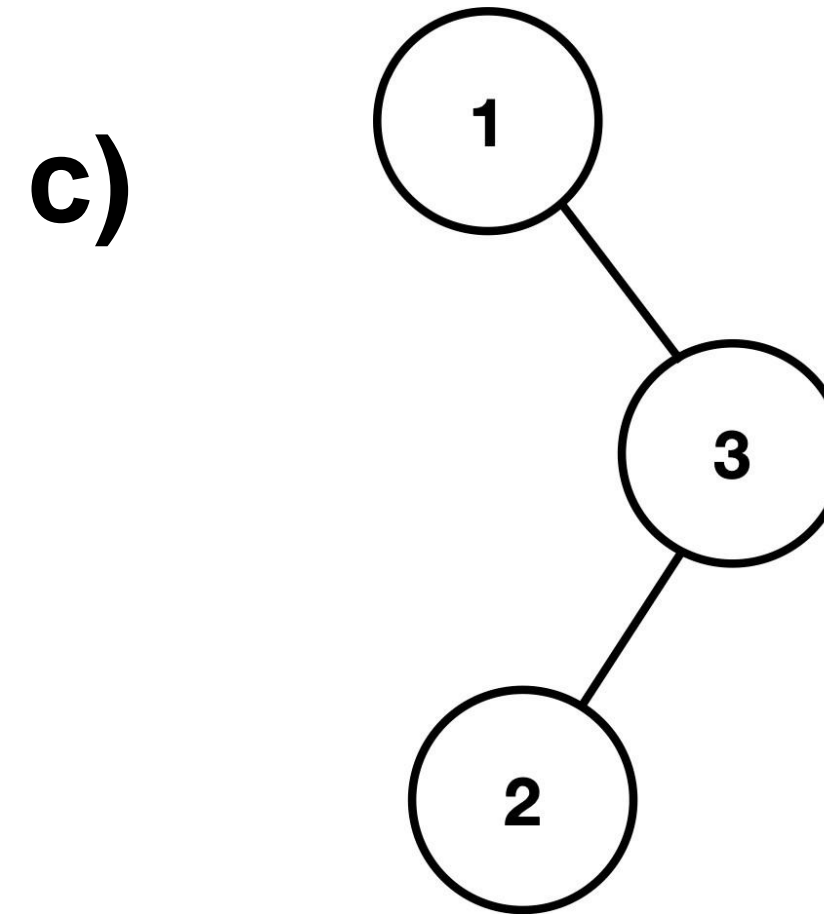
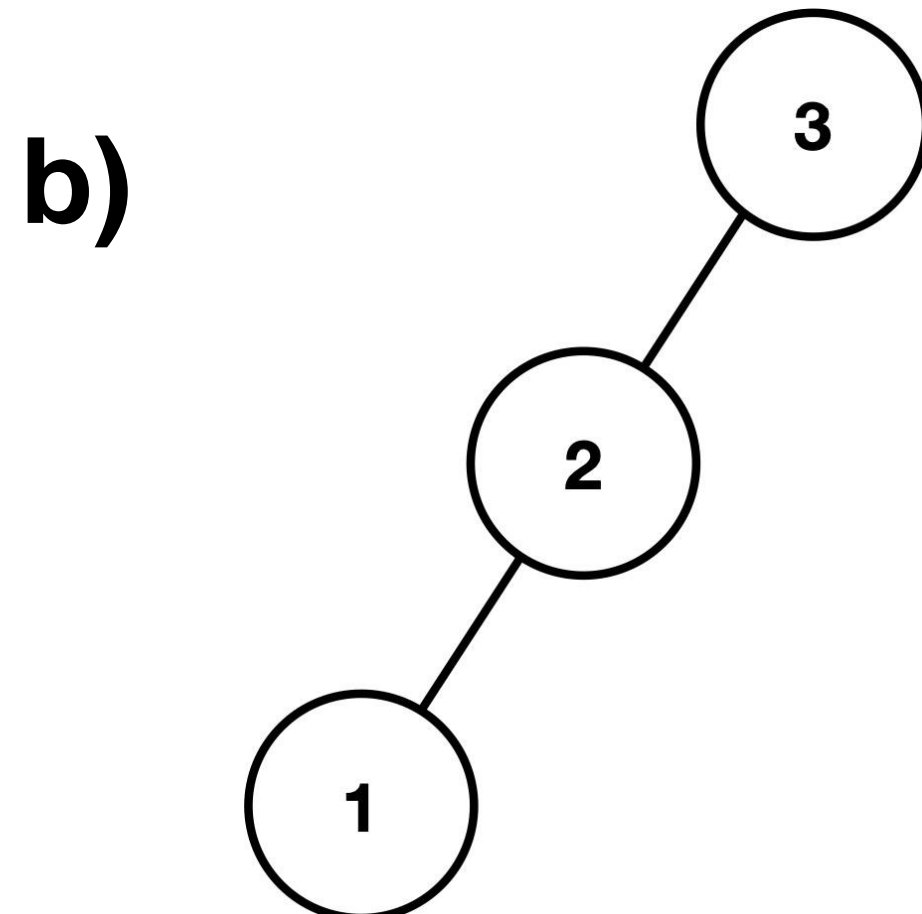
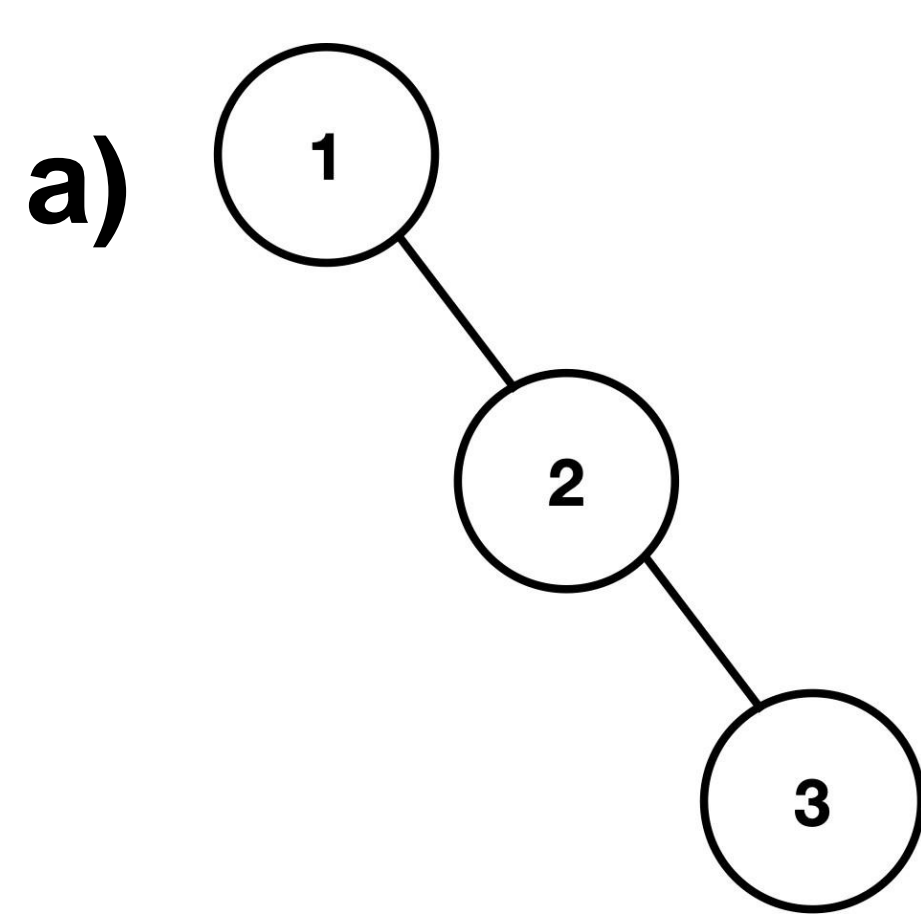
- Single Right Rotation _____
- Single Left Rotation _____
- Right Left Rotation _____
- Left Right Rotation _____



AVL Tree #2

Match the following picture with the rotation that should be performed:

- Single Right Rotation **__b__**
- Single Left Rotation **__a__**
- Left Right Rotation **__c__**
- Right Left Rotation **__d__**



AVL Tree #3

Indicate whether the following statements are true or false

- The acceptable balance factor values for an AVL Tree are: -1, 0, and 1.**
- The worst case time complexity for an AVL Tree is $O(n)$.**
- The maximum height of an AVL Tree with 7 nodes is 2.**
- An AVL Tree is a BST with a self balancing property.**

AVL Tree #3

Indicate whether the following statements are true or false

- The acceptable balance factor values for an AVL Tree are: -1, 0, and 1.

true

- The worst case time complexity for an AVL Tree is $O(n)$.

false

- The maximum height of an AVL Tree with 7 nodes is 2.

false

- An AVL Tree is a BST with a self balancing property.

true

Stack & Queue #1

Implement an enqueue function using two stacks. No other data structure is allowed.

```
class Queue{
    private:
        stack<int> s1;
        stack<int> s2;
    public:
        void push(int x) ;
};

void Queue::enqueue(int x) {

}
```


Stack & Queue #1

```
void Queue::enqueue(int x) {  
    while(!s1.empty()) {  
        s2.push(s1.top());  
        s1.pop();  
    }  
    s1.push(x);  
    while(!s2.empty()) {  
        s1.push(s2.top());  
        s2.pop();  
    }  
}
```

Stack & Queue #2

Write a function to delete all of the occurrences of a certain value from a queue. You are only allowed one additional variable.

```
void deleteAll(queue<int> &q, int x) {
```

```
}
```

Stack & Queues #2

```
7 ▼ void deleteAll(queue<int> &q, int x){  
8     int size = q.size();  
9 ▼ while(size > 0){  
10         if(q.front() != x)  
11             q.push(q.front());  
12         q.pop();  
13         size--;  
14     }  
15 }
```

Stack & Queue #3

Write a function that returns the minimum value of a stack and has a time complexity of $O(1)$.

```
class Stack{
    private:
        stack<int> s;
    public:
        void push(int x) ;
        int getMin() ;
};

void Stack::push(int x) {

}

int Stack::getMin() {

}
```

Stack & Queue #3

```
int Stack::getMin(){  
    if(!s.empty())  
        return s.top();  
    else  
        return -1000;  
}
```

```
void Stack::push(int x){  
    if(s.empty())  
        s.push(x);  
    else{  
        if(x <= s.top())  
            s.push(x);  
        else{  
            stack<int> tempStack;  
            while(!s.empty() && x > s.top()){  
                tempStack.push(s.top());  
                s.pop();  
            }  
            s.push(x);  
            while(!tempStack.empty()){  
                s.push(tempStack.top());  
                tempStack.pop();  
            }  
        }  
    }  
}
```

Stack & Queue #4

Implement a function to find the n-th smallest element in an array using priority queue. You must implement both this function and enqueue function for priority queue.

```
void enqueue(int value){}
```

```
int nthSmallest(int arr[], int size, int n){}
```


Stack & Queue #4

```
int nthSmallest(int arr[], int size, int n) {
    priority_queue q;
    for (int i=0; i<size; i++) {
        q.enqueue(arr[i]);
    }
    for (int i=1; i<n; i++) {
        q.dequeue();
    }
    return q.front();
}
```

```
void enqueue(int value) {
    node* temp = new node();
    temp->data = value;
    temp->next = NULL;

    if (isEmpty()) {
        front=temp;
        rear=temp;
    }
    else {
        node* cu = front;
        node* prev = NULL;
        if (temp->data<cu->data) {
            temp->next = front;
            front = temp;
        }
        else {
            while (cu!=NULL && temp->data>=cu->data) {
                prev = cu;
                cu = cu->next;
            }
            prev->next = temp;
            temp->next = cu;
            if (temp->next==NULL)
                rear = temp;
        }
    }
}
```

Sorting #1

Perform heap sort on the array below. Sort the numbers in ascending order. Show all steps.

{4, 1, 2, 9, 5, 8, 3}

Sorting #1

Perform heap sort on the array below. Sort the numbers in ascending order.

{4, 1, 2, 9, 5, 8, 3} <- heapify

{9, 5, 8, 1, 4, 2, 3} <- swap

{3, 5, 8, 1, 4, 2, 9} <- heapify

{8, 5, 3, 1, 4, 2, 9} <- swap

{2, 5, 3, 1, 4, 8, 9} <- heapify

{5, 4, 3, 1, 2, 8, 9} <- swap

{2, 4, 3, 1, 5, 8, 9} <- heapify

{4, 2, 3, 1, 5, 8, 9} <- swap

{1, 2, 3, 4, 5, 8, 9} <- heapify

{3, 2, 1, 4, 5, 8, 9} <- swap

{1, 2, 3, 4, 5, 8, 9} <- heapify

{2, 1, 3, 4, 5, 8, 9} <- swap

{1, 2, 3, 4, 5, 8, 9}

Sorting #2

Implement a sorting function that have the time complexity as follow

Best: $O(n \log(n))$

Average: $O(n \log(n))$

Worse: $O(n^2)$

Sorting #2

Quick Sort:

Write the code for quick sort below as well, look at Exam 2 Basics-Code for the code.

Sorting #3

Describe the worse case scenario for quick sort.

Sorting #3

Worse case $O(n^2)$ When the pivot is at the beginning, or the end of the array and the array is already sorted

Sorting #4

**Implement a function that check if an array is a min heap.
Return true if the array represent a min heap, return false
otherwise.**

Sorting #4

```
bool isMinHeap(int arr[], int size) {  
    for (int i = 0; i < (size - 2) / 2; i++)  
        if (arr[i] > arr[2 * i + 1] || arr[i] > arr[2 * i +  
2])  
            return false;  
    return true;  
}
```