

# DIJKSTRA'S ALGORITHM

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# Greedy algorithms

**Greedy algorithms:** constructs a solution through a sequence of **steps**, each **expanding** a partially constructed solution

On each step, the choice made **must be**:

- **Feasible**: satisfies the problem's constraints
- **Locally optimal**: best feasible local choice
- **Irrevocable**: cannot be changed later



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# Dijkstra's algorithm

Let  $G$  be a weighted graph and  $v \in V$  (*source*), finds the shortest path from  $v$  to all other nodes in  $V$

- Single-source shortest paths

## Applications

- Transport planning
- Communication networks
- Social networks
- Robotics
- Pathfinding
- Puzzles
- etc.

Dijkstra's algorithm: **cannot** be used on weighted graphs with **negative weights**



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# Dijkstra's algorithm

First, find the **closest node** to  $v$  (itself)

On the  $i$ -th step:

- Knows the  $(i - 1)$ -th closest nodes to  $v$  (they form a tree)
- Since there are no negative weights, the next closest one is adjacent to one of the  $i - 1$  closest nodes to  $v$
- After choosing the  $i - th$  closest node ( $w$ ), updates the possible shortest paths to yet unchosen nodes ( $u$ ) if  $d_w + \text{weight}(w, u) < d_u$



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# Dijkstra's algorithm

**Algorithm:** void

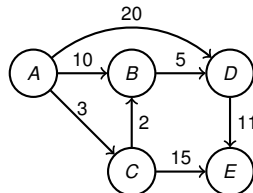
Dijkstra(Graph G, int s, int[] D)

```

1  for  $i \leftarrow 0$  to  $n(G) - 1$  do
2  |    $D[i] \leftarrow \infty$  ;  $P[i] \leftarrow -$ ;
3  |    $\text{setMark}(G, i, \text{UNVISITED})$ ;
4   $H[1] \leftarrow (s, s, 0)$  ;  $D[s] \leftarrow 0$ ;
5  for  $i \leftarrow 0$  to  $n(G) - 1$  do
6  |   repeat
7  |   |    $(p, v) \leftarrow \text{removemin}(H)$ ;
8  |   |   if  $v = \text{NULL}$  then return;
9  |   until  $\text{getMark}(G, v) = \text{UNVISITED}$ ;
10 |    $\text{setMark}(G, v, \text{VISITED})$  ;  $P[v] \leftarrow p$ ;
11 |    $w \leftarrow \text{first}(G, v)$ ;
12 |   while  $w < n(G)$  do
13 |   |   if  $\text{getMark}(G, w) \neq \text{VISITED} \wedge$ 
14 |   |   |    $D[w] > D[v] + \text{weight}(G, v, w)$  then
15 |   |   |   |    $D[w] \leftarrow D[v] + \text{weight}(G, v, w)$ ;
16 |   |   |   |    $\text{insert}(H, (v, w, D[w]))$ ;
17 |   |    $w \leftarrow \text{next}(G, v, w)$ ;

```

Let  $s = A$



	A	B	C	D	E
Mark	×	×	×	×	×
Distance	0	$\infty$	$\infty$	$\infty$	$\infty$
Parent	—	—	—	—	—

(A,A,0)



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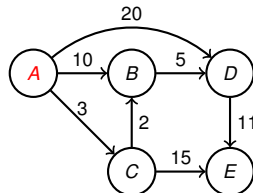
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Let  $s = A$



	A	B	C	D	E
Mark	✓	×	×	×	×
Distance	0	$\infty$	$\infty$	$\infty$	$\infty$
Parent	A	—	—	—	—

—



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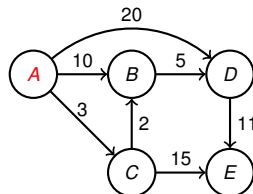
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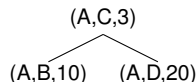
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```

Let  $s = A$



	A	B	C	D	E
Mark	✓	×	×	×	×
Distance	0	10	3	20	$\infty$
Parent	A	—	—	—	—



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# Dijkstra's algorithm

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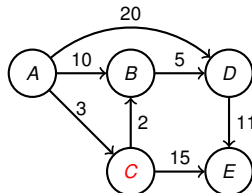
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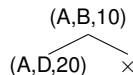
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```

Let  $s = A$



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Mark	✓	×	✓	×	×
Distance	0	10	3	20	$\infty$
Parent	A	—	A	—	—



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# Dijkstra's algorithm

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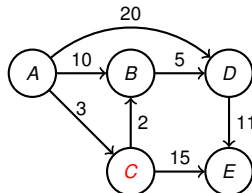
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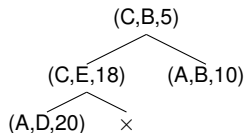
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	A	B	C	D	E
Mark	✓	×	✓	×	×
Distance	0	5	3	20	18
Parent	A	—	A	—	—



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# Dijkstra's algorithm

**Algorithm:** void

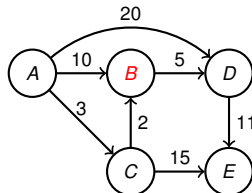
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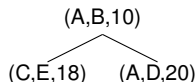
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Let  $s = A$



	A	B	C	D	E
Mark	✓	✓	✓	×	×
Distance	0	5	3	20	18
Parent	A	C	A	—	—



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# Dijkstra's algorithm

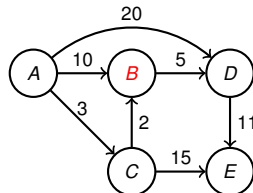
**Algorithm:** void

Dijkstra(Graph G, int s, int[] D)

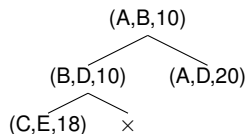
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	A	B	C	D	E
Mark	✓	✓	✓	×	×
Distance	0	5	3	10	18
Parent	A	C	A	—	—



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# Dijkstra's algorithm

**Algorithm:** void

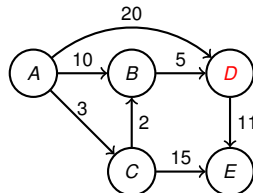
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```

Let  $s = A$



	A	B	C	D	E
Mark	✓	✓	✓	✓	×
Distance	0	5	3	10	18
Parent	A	C	A	B	—

(C,E,18)  
(A,D,20) ×



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# Dijkstra's algorithm

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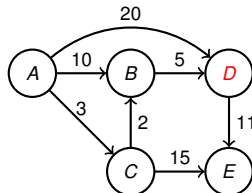
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Let  $s = A$



	A	B	C	D	E
Mark	✓	✓	✓	✓	×
Distance	0	5	3	10	18
Parent	A	C	A	B	—

(C,E,18)  
(A,D,20) ×



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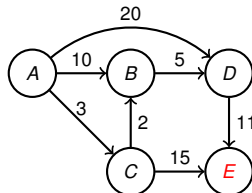
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Let  $s = A$



	A	B	C	D	E
Mark	✓	✓	✓	✓	✓
Distance	0	5	3	10	18
Parent	A	C	A	B	C

(A,D,20)



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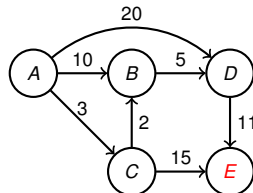
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Mark	✓	✓	✓	✓	✓
Distance	0	5	3	10	18
Parent	A	C	A	B	C

(A,D,20)



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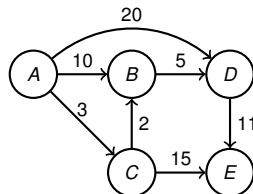
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10      $setMark(G, v, VISITED)$  ;  $P[v] \leftarrow p$ ;
11      $w \leftarrow first(G, v)$ ;
12     while  $w < n(G)$  do
13         if  $getMark(G, w) \neq VISITED \wedge$ 
14              $D[w] > D[v] + weight(G, v, w)$  then
15              $D[w] \leftarrow D[v] + weight(G, v, w)$ ;
16              $insert(H, (v, w, D[w]))$ ;
17          $w \leftarrow next(G, v, w)$ ;

```

Let  $s = A$



	A	B	C	D	E
Mark	✓	✓	✓	✓	✓
Distance	0	5	3	10	18
Parent	A	C	A	B	C

(A,D,20)



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# Dijkstra's algorithm

**Algorithm:** void

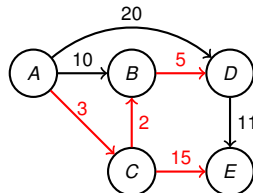
Dijkstra(Graph G, int s, int[] D)

```

1  for  $i \leftarrow 0$  to  $n(G) - 1$  do
2     $D[i] \leftarrow \infty$  ;  $P[i] \leftarrow -$ ;
3     $\text{setMark}(G, i, \text{UNVISITED})$ ;
4   $H[1] \leftarrow (s, s, 0)$  ;  $D[s] \leftarrow 0$ ;
5  for  $i \leftarrow 0$  to  $n(G) - 1$  do
6    repeat
7       $(p, v) \leftarrow \text{removemin}(H)$ ;
8      if  $v = \text{NULL}$  then return;
9    until  $\text{getMark}(G, v) = \text{UNVISITED}$ ;
10    $\text{setMark}(G, v, \text{VISITED})$  ;  $P[v] \leftarrow p$ ;
11    $w \leftarrow \text{first}(G, v)$ ;
12   while  $w < n(G)$  do
13     if  $\text{getMark}(G, w) \neq \text{VISITED} \wedge$ 
14        $D[w] > D[v] + \text{weight}(G, v, w)$  then
15        $D[w] \leftarrow D[v] + \text{weight}(G, v, w)$ ;
16        $\text{insert}(H, (v, w, D[w]))$ ;
17    $w \leftarrow \text{next}(G, v, w)$ ;

```

Let  $s = A$



	A	B	C	D	E
Mark	✓	✓	✓	✓	✓
Distance	0	5	3	10	18
Parent	A	C	A	B	C

(A,D,20)



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# Dijkstra's algorithm

**Algorithm:** void

Dijkstra(Graph G, int s, int[] D)

```

1  for i ← 0 to n(G) − 1 do
2      D[i] ← ∞ ; P[i] ← −;
3      setMark(G, i, UNVISITED);
4  H[1] ← (s, s, 0) ; D[s] ← 0;
5  for i ← 0 to n(G) − 1 do
6      repeat
7          (p, v) ← removemin(H);
8          if v = NULL then return;
9      until getMark(G, v) = UNVISITED;
10     setMark(G, v, VISITED) ; P[v] ← p;
11     w ← first(G, v);
12     while w < n(G) do
13         if getMark(G, w) ≠ VISITED ∧
14            D[w] > D[v] + weight(G, v, w) then
15             D[w] ← D[v] + weight(G, v, w);
16             insert(H, (v, w, D[w]));
17         w ← next(G, v, w);

```

Time efficiency

■ **Matrix and no heap**

$\Theta(|V|^2 + |E|) = \Theta(|V|^2)$ ,  
since  $|E| \in O(|V|^2)$

■ Better for dense graphs

■ **List and heap**

$\Theta((|V| + |E|) \log |V|)$

■ Better for sparse graphs



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# Agenda

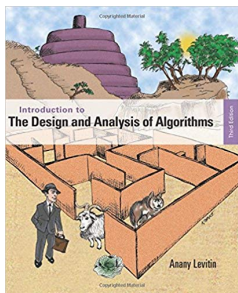


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# DIJKSTRA'S ALGORITHM

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