# Data Structures and Algorithms for Competitive Programming

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# 1 Graphs

```
Search
   \mathbf{BFS}
   int dist[MXN];
   vector < int > adjList [MX.N];
3
   int main(){
        for (int i = 0; i < MX.N; i++)
            dist[i]=INF;
        queue<int> q;
        q.push(0);
        dist[0] = 0;
10
        while (!q.empty()) {
            int u = q.front(); q.pop();
11
12
            int d = dist[u];
            for(int i : adjList[u]) {
13
                 if (dist [i]==INF) {
14
                     dist[i]=d+1;
15
16
                     q.push(i);
17
18
19
20
        return 0;
   Dijkstras
   struct path {
2
        int u,d;
3
        bool operator < (const path& p) const {
            return d > p.d;
 4
5
6
    };
    for (int i =0; i < N; ++i)
        dist[i] = INF;
    dist[S] = 0;
    priority_queue < path > q;
   q.push(path(S,0));
12
   while (!q.empty()) {
13
        path p = q.top(); q.pop();
14
        u = p.u, d = p.d;
15
        if(dist[u] < d)
16
            continue;
17
        for (auto v : adjList[u]) {
            nd = d + v.second;
18
19
            if(nd < dist[v.first]){</pre>
20
                 dist[v.first] = nd;
21
                q.push(path{v.first,nd});
22
23
24
```

```
\mathbf{A}^*
    struct path {
        int u, wght;
        bool operator < (const path& rhs) const {
4
            return wght > rhs.wght;
5
6
    };
   int h(int x) { // heuristic function
        // Don't overestimate, should be 0 at sink
9
        return 0; // 0 for djikstras
10
11
    for (int i = 0; i < N; ++i)
12
        dist[i] = INF;
    dis[S] = 0;
13
    priority_queue < path > pq;
14
   pq.push({s, h(s)});
16
    while (!pq.empty()) {
17
        path p = pq.top(); pq.pop();
18
19
        if(p.u = t)
20
            return p.wght;
21
22
        int d = p.wght - h(p.u);
23
24
        if(dis[p.u] < d)
25
            continue:
26
27
        for(auto v : adjList[p.u]) {
28
            int nd = d + v.second;
29
30
            if(dis[v.first] \le nd)
31
                 continue;
32
33
            dis[v.first] = nd;
34
35
            pq.push(\{v.first, nd + h(v.first)\});
36
37
    Bidirectional BFS
1
   int bidirbfs(){
2
        queue < int > fq, rq;
3
        fq.push(S); dis[S] = 0;
4
        rq.push(T); disr[T] = 0;
        vector<int> proc;
7
        for (;;) {
8
            if (fq.empty() && rq.empty())
9
                 return -1; // disconnected
10
            if (! fq . empty()) {
11
                 int u = fq.front(); fq.pop();
12
                 for(int v : adjList[u]){
13
                     if(dis[v] \le dis[u] + 1)
```

```
continue;
14
                     dis[v] = dis[u] + 1;
15
16
                     fq.push(v);
17
18
                 vis[u] = true;
19
                 proc.pb(u);
20
21
            if (!rq.empty()){
22
                 int u = rq.front(); rq.pop();
23
                 for (int v : RadjList [u]) { // REVERSE adj
24
                     if(disr[v] \le disr[u] + 1)
25
                         continue;
26
                     disr[v] = disr[u] + 1;
27
                     rg.push(v);
28
29
                 if (vis [u])
30
                     break;
31
                 proc.pb(u);
32
33
        }
34
35
        int mn = INF;
36
        for(int u : proc)
37
            if(dis[u] != INF && disr[u] != INF)
38
                mn = min(mn, dis[u] + disr[u]);
39
        return mn;
40
   Trees
   MST
   struct edge {
2
        int x,y,w;
        bool operator < (edge e) const {
 4
            return w < e.w;
 5
 6
   int main(){
9
        vector < edge > eList; //Input
10
        for (int i = 0; i < N; i++)// Set up UFDS
11
            p[i]=i;
12
        vector<ii> treeList;
13
        sort (eList.begin(), eList.end());
14
        int cost = 0;
        int sz=N;
15
16
        int u, v, w;
17
        for (const auto &i : eList) {
18
            v=i.x; u=i.y; w=i.w;
19
            if (!connected(u,v)){
20
                 join (u, v);
21
                 treeList.push_back(\{min(u,v),max(u,v)\});
22
23
                 cost+=w;
```

```
24
25
26
        if(sz!=1)
27
            puts("Impossible");
28
   LCA
1
    *H/u is first visit of u
    *E[x] is vertex at time x
    *L/x is depth at time x
    * MEMSET H to -1
    * SET L to size of list * 2
7
    */
8
   int vind;
   void vis(int u, int d){
10
       H[u] = vind;
11
       E[vind] = u;
12
       L[vind++] = d;
        for (auto i : adjList[u]) {
13
14
            if(H[i]!=-1)
15
                continue;
16
            vis(i,d+1);
17
            E[vind] = u;
            L[vind++] = d;
18
19
20
21
22
   int LCA(int u, int v){
23
        if(H[u] > H[v])
24
            int t = u;
25
            u = v;
26
            v = t;
27
        //run some range min query on L
28
29
        //between H[u] and H[v]
30
        int ind = rmq(H[u],H[v]);
31
        return E[ind];
32
   }
33
   int dist(int u, int v){
34
35
        int a = H[u];
36
        int b = H[v];
        int ind = LCA(u, v);
37
38
        return abs(L[H[ind]]-L[a])
39
            + abs(L[H[ind]]-L[b]);
40 }
    Centroid Decomposition
   void fill_sz(int u, int p){
2
        sz[u] = 1;
3
        for(int v : adjList[u]) {
4
            if (v==p || mkd[v])
```

```
continue:
6
            fill_sz(v,u);
7
            sz[u]+=sz[v];
8
9
10
11
   int get_centroid(int u, int n, int p){
12
        for(int v : adjList[u]){
13
            if (v==p || mkd[v])
14
                continue;
15
            if(sz[v] > n/2)
16
                return get_centroid(v, n, u);
17
18
        return u:
19
20
21
   int decomp(int u){
22
        fill_sz(u, -1);
23
        int cent = get_centroid(u, sz[u], -1);
24
        mkd[cent] = true;
25
        for(int v : adjList[cent]){
26
            if (mkd [v])
27
                continue;
28
            int r = decomp(v);
29
            centP[r] = cent;
30
31
        return cent;
32
   SCC Tarjans
   stack<int> scc;
   int dfsCounter=1;
   int sccIdx=1;
   map<int, int> sccMap;
   void tarjans(int u){
        scc.push(u);
        vis [u]=true;
8
9
10
        dfs_low[u] = dfs_num[u] = dfsCounter++;
11
12
        for (int i = 0; i < adjList[u].size(); i++){
13
            int v = adjList[u][i];
            if(dfs_num[v]==0)
14
15
                tarjans(v);
                dfs_low[u] = min(dfs_low[u], dfs_low[v]);
16
17
            } else if (vis[v]) {
                dfs_low[u]=min(dfs_low[u],dfs_num[v]);
18
19
20
21
        if(dfs_low[u]==dfs_num[u])
22
            while(1){
23
                int v = scc.top(); scc.pop();
```

 $\operatorname{sccMap}[v] = \operatorname{sccId}x;$ 

24

```
25
                vis[v] = false;
26
                i f ( v==u )
27
                    break;
28
29
            sccIdx++;
30
31
   AP & Bridges
   int dfs(int u,int p){
2
        dfs_num[u] = dfs_low[u] = ++dfs_counter;
3
        for(auto v : adjList[u]){
4
            if(dfs_num[v]==0)
                dfs(v,u);
5
                if(dfs_low[v] >= dfs_num[u]) {
6
                    articulation [u]=true;
8
9
                if(dfs_low[v] > dfs_num[u])
10
                    bridge = true;
                dfs_low[u] = min(dfs_low[u], dfs_low[v]);
11
12
            else if(v!=p)
13
                dfs_low[u] = min(dfs_low[u], dfs_num[v]);
14
15
```

#### **Network Flow**

#### Edmond Karp Max Flow

```
void aug(int u, int minE){
2
        if(u=S){ f=minE; return; }
3
        if(p[u]!=u){
            aug(p[u], min(minE, res[p[u]][u]));
            res[p[u]][u]-=f;
            res[u][p[u]] += f;
7
8
9
10
   int main(){
11
        int mf=0:
12
        for (;;) {
13
             f=0;//Global
14
            for (int i = 0; i < N; i++)
15
                 dist[i]=INF, p[i]==i;
            dist[S]=0;
16
17
            queue < int > q; q.push(S);
18
            while (!q.empty()) {
19
                 int u = q. front(); q.pop();
                 if(u==T) break;
20
21
                 for (int i = 0; i < N; i++)
22
                     if(res[u][i] > 0 \&\& dist[i] == INF)
23
                          dist[i]=dist[u]+1, p[i]=u, q.push(i);
24
```

```
25
            aug(T, INF);
26
            if(f==0) break;
27
            mf+=f;
28
29
        vector<ii> used;
30
        for (int i = 0; i < N; i++)
31
            for (int j = 0; j < N; j++)
32
                if(graph[i][j] > 0 && res[i][j] < graph[i][j])
33
                     used.push_back(make_pair(i,j));
34
   Ford Fulkerson Max Flow
   int ff(int u, int minE){
2
        i f ( u==T)
 3
            return minE;
 4
        vis[u]=true;
        for(auto i : adjList[u]){
 5
 6
            if (! vis [i] && res [u][i] > 0) {
                if(int f = ff(i, min(minE, res[u][i])))
                     res[u][i] -= f;
                     res[i][u] += f;
10
                     return f;
11
12
13
14
        return 0;
15
16
17
   int main(){
        int mf = 0:
18
19
        while(1){
20
            memset(vis,0,sizeof(vis));
21
            int f = ff(S, INF);
22
            if(f==0)
23
                break;
24
            mf+=f;
25
26
        printf("%d\n", mf);
27 }
```

# 2 Data Structures

### Fenwick Tree

```
int tree [MX_N];
   int N;
   int lsOne(int i) \{ return i \& (-i); \}
    void update(int k,int v){
        for (; k<MX.N; k+=lsOne(k))
6
            tree[k]+=v;
7
8
   int query(int k){
9
        int cnt = 0;
        for (; k; k-=lsOne(k)) {
10
11
            cnt+=tree [k];
12
13
        return cnt;
14 }
    UFDS
   int find(int u) { return p[u] = (p[u] = u ? u : find(p[u])); }
   inline void join (int a, int b) {
        pa = find(a);
        pb = find(b);
        if (pa!=pb) {
            if(rank[pa] < rank[pb])
                ni = pb;
                pb = pa;
10
                pa = ni;
11
12
            p[pb] = pa;
13
            if (rank [pa]==rank [pb])
14
                rank [pa]++;
15
16
   Sparse Table
```

```
inline int rmq(int u, int v){
2
        if(u > v)
3
            return -20000000000;
        int k=(int) floor(log2((double)(v-u+1)));
        if(r[mtable[u][k]]>
5
6
                r [ mtable [ v-(1 << k) + 1 ] [k] ] )
            return mtable [u][k];
        return mtable [v-(1 << k) + 1][k];
9
10
11
   for(int i = 0; i < N; i++)
        mtable[i][0] = i;
  for (int j = 1; (1 << j) <= N; j++)
```

```
Lazy Segment Tree
14
        for (int i = 0; i + (1 << j) - 1 < N; ++i)
15
            if (r [mtable [i] [j−1]]
16
                     >r [ mtable [ i+(1<<(j-1)) ] [ j-1] ])
                                                                                      11 tree [4 * MX N + 2];
17
                 mtable[i][j] = mtable[i][j-1];
                                                                                      11 d[4 * MX_N + 2];
18
            else
19
                 \text{mtable}[i][j] = \text{mtable}[i+(1 << (j-1))][j-1];
                                                                                      void prop(int p, int l, int r){
                                                                                           tree [p] += d[p] *(r-l+1);
                                                                                   6
                                                                                           if (r>l) {
   Segment Tree
                                                                                               d[2*p]+=d[p];
                                                                                               d[2*p + 1] + = d[p];
                                                                                   9
 1 int tree [MX.N*4 + 2];
                                                                                   10
                                                                                           d[p] = 0;
   int a [MX_N];
   int N;
                                                                                   11
                                                                                   12
                                                                                      ll rsq(int p, int ql, int qr, int l, int r) {
   void construct(int p, int L, int R){
                                                                                   13
                                                                                   14
                                                                                           prop(p, l, r);
 6
        if (L=R) {
                                                                                  15
                                                                                           if(r < ql \mid | l > qr)
            tree[p] = L;
                                                                                  16
                                                                                               return 0;
 8
            return:
                                                                                  17
                                                                                           if(l >= ql \&\& r <= qr)
9
        if (R<L)
                                                                                   18
                                                                                               return tree[p];
10
                                                                                   19
                                                                                           int md = (l+r)/2;
11
            return;
                                                                                   20
                                                                                           return rsq(2*p, ql, qr, l, md)
12
        int md = (L+R)/2;
                                                                                   21
13
        construct(2*p,L,md);
                                                                                               + rsq(2*p+1, ql, qr, md+1, r);
                                                                                   22
14
        construct (2*p+1,md+1,R);
                                                                                   23
15
        tree[p] = a[tree[2*p]] < a[tree[2*p+1]]? tree[2*p]: tree[2*p+1];
16
                                                                                   24
                                                                                       void rng_up(int p, int ql, int qr, ll v, int l, int r){
17
                                                                                   25
                                                                                           prop(p,l,r);
                                                                                   26
                                                                                           if(r < ql \mid | l > qr)
18
   void update(int p, int L, int R, int ind, int v){
19
                                                                                   27
                                                                                               return;
        if (L=R) {
                                                                                   28
20
            a[ind] = v;
                                                                                           if(l)=ql \&\& r <= qr)
21
            tree[p] = ind;
                                                                                   29
                                                                                               d[p]+=v;
                                                                                   30
22
                                                                                               prop(p, l, r);
            return;
23
                                                                                  31
                                                                                               return:
                                                                                   32
24
        int md = (L+R)/2;
                                                                                   33
                                                                                           int md=(l+r)/2;
25
        if (ind \le md)
                                                                                   34
                                                                                           rng_up(2*p, ql, qr, v, l, md);
26
            update (2*p, L, md, ind, v);
                                                                                   35
27
                                                                                           rng_up(2*p+1, ql, qr, v, md+1,r);
        else
28
                                                                                   36
                                                                                           tree[p] = tree[2*p] + tree[2*p+1];
            update(2*p+1,md+1,R,ind,v);
29
        tree[p] = a[tree[2*p]] < a[tree[2*p+1]]? tree[2*p]: tree[2*p+1];
                                                                                  37
30
                                                                                       Convex Hull Trick
31
   int rmg(int p, int L, int R, int l, int r){
33
        if(r < L \mid \mid l > R)
                                                                                      // Convex hull trick, simpler if can sort insertions and queries
34
            return INF;
                                                                                      struct CvxHullTrickSimple {
35
                                                                                   3
        if(l>=L && r<=R)
                                                                                           vector < ll > A;
36
            return tree[p];
                                                                                           vector<ll> B;
                                                                                   4
37
        int md = (1+r)/2;
38
        int lf = rmq(2*p, L, R, l, md);
                                                                                           int ptr;
39
        int rf = rmq(2*p+1,L,R,md+1,r);
40
        if(lf >= INF)
                                                                                           cvxH():ptr(0)\{\}
41
            return rf;
                                                                                           // insert a descending
42
        if(rf >= INF)
                                                                                   10
43
                                                                                   11
                                                                                           void addLine(ll\ a, ll\ b) { // intersection\ of\ (A[len-2],B[len-2])\ with
            return lf;
44
        return a[lf] < a[rf]? lf : rf;
                                                                                               (A/len-1), B/len-1) must lie to the left of intersection of (A/len-1)
                                                                                               -1/B[len-1] with (a,b)
45 }
```

```
12
             while (sz(A) \ge 2 \&\& (B[sz(B) - 2] - B[sz(B) - 1]) * (a - A[sz(A)65])
                 [-1] >= [B[sz(B) - 1] - b] * (A[sz(A) - 1] - A[sz(A) - 2])) 66
13
                 A. pop_back();
                                                                                     67
14
                 B. pop_back();
                                                                                     68
15
                                                                                     69
            A. pb(a);
                                                                                     70
16
17
            B.pb(b);
                                                                                     71
18
                                                                                     72
19
                                                                                     73
20
        // query x ascending
                                                                                     74
21
        ll minValue(ll x) {
                                                                                     75
22
          ptr = min(ptr, sz(A) - 1);
                                                                                     76
23
          while (ptr + 1 < sz(A) \&\& A[ptr + 1] * x + B[ptr + 1] <= A[ptr] * x77
                                                                                     79
24
            ++ptr;
25
                                                                                     80
26
          return A[ptr] * x + B[ptr];
                                                                                     81
27
                                                                                     82
28
    };
                                                                                     83
29
                                                                                     84
   struct CvxHullOpt {
                                                                                     85
31
        // Fully dynamic variant for use if can't quarantee
                                                                                     86
32
        // insertion or query order
                                                                                     87
33
        // stores lower envelope, negate lines and function to get upper
                                                                                     88
                                                                                     89
34
        static const 11 qV = -(1LL<<50); //hacky special value
35
        struct line {
36
             ll a,b;
37
             mutable double xLeft;
38
             bool operator < (const line & l) const {
39
                 if(l.a != qV)
                     return a < l.a;
41
                 return xLeft > l.xLeft;
42
        };
43
44
        multiset < line > hull;
45
        ll xcomp(const line& a, const line& b, const line& c){
46
             // < \theta \Rightarrow AxB < AxC
47
             // > 0 \Rightarrow AxB > AxC
48
49
             return (a.a-c.a)*(b.b-a.b) - (c.b-a.b)*(a.a-b.a);
50
51
52
        double xin(const line& a, const line& b){
53
             return (b.b - a.b)/(1.0*(a.a-b.a));
54
        }
55
56
        bool bad(set<line>::iterator y){
57
             auto z = next(y);
58
             if (v=hull.begin()) {
59
                 if (z=hull.end())
60
                     return false;
61
                 return y\rightarrow a == z\rightarrow a \&\& y\rightarrow b >= z\rightarrow b;
62
63
             auto x = prev(y);
64
             if (z=hull.end())
```

```
return x->a == v->a && x->b <= v->b:
    return xcomp(*x, *y, *z) < 0;
void addLine(ll a, ll b) { // add line a*x + b
    CvxHullOpt::line l{a,b};
    auto y = hull.insert(1);
    if(bad(y)) { hull.erase(y); return;}
    while (next (y) != hull.end() && bad(next (y)))
        hull.erase(next(y));
    while (y!=hull.begin() && bad(prev(y)))
         hull.erase(prev(y));
    if(next(y)=hull.end())
        v \rightarrow x Left = -INF;
    else
        y \rightarrow x Left = xin(*y, *next(y));
    if(y != hull.begin())
         prev(y) \rightarrow xLeft = xin(*prev(y), *y);
11 \text{ eval}(11 \text{ x})
    auto l = hull.lower_bound(\{qV, 0, x\});
    return l\rightarrow a*x + l\rightarrow b;
```

};

# 3 Geometry

#### Convex Hull

```
int main() {
2
        for (int i = 0; i < N; i++){
3
            perm[i]=i;
4
 5
        sort (perm, perm+N,
6
                 [](int a, int b){
                     const point &pa = V[a];
                     const point &pb = V[b];
9
                     if (real (pa)!=real (pb))
10
                          return real(pa) < real(pb);
11
                     return imag(pa) < imag(pb);
12
13
        vector<int> L; vector<int> U;
14
        for (int i = 0; i < N;) {
15
            int t = L. size();
16
             if(t \ge 2 \&\& ! ccw(V[L[t-2]],V[L[t-1]],V[perm[i]]))
17
                 L. pop_back();
18
             else
19
                 L. push_back (perm [i++]);
20
21
        for (int i = N-1; i >=0;) {
22
            int t = U. size();
23
             if(t \ge 2 \&\& ! ccw(V[U[t-2]],V[U[t-1]],V[perm[i]]))
24
                 U.pop_back();
25
             else
26
                 U. push_back (perm [i --]);
27
28
        vector<int> hull;
29
        for (int i = 0; i < L. size() -1; ++i)
             hull.push_back(L[i]);
30
31
        for (int i = 0; i < U. size() -1; ++i)
32
             hull.push_back(U[i]);
33
        return 0;
34
```

# Geometry Axioms

```
1 typedef complex < double > pt;
2 typedef complex < double > vec;
3 typedef vector < pt > pgon;
4 typedef struct { pt p,q; } lseg;
5 struct circ{ pt c; double r; };
6 struct rect{ pt p,q;}; // X(p) <= X(q) and Y(p) <= Y(q)
7
8 double cross(const vec& a, const vec &b){
9 return x(a)*y(b)-y(a)*x(b);
10 }
11 // cross product of (b-a) and (c-b), 0 is collinear
12 int orientation(const pt& a,
13 const pt& b, const pt& c){</pre>
```

```
14
        double v = cross(b-a, c-b);
15
        if (abs (v-0.0) < EPS)
16
             return 0;
17
        return v > 0 ? 1 : 2;
18
19
    //Line segment intersection
    bool intersects (const lseg& a, const lseg& b) {
21
        if(a.q = b.p \mid\mid b.q = a.p)
22
             return false;
23
        if (orientation (a.p,a.q,b.p)
24
                 != orientation (a.p, a.q, b.q)
25
                 && orientation (b.p,b.q,a.p)
26
                 != orientation(b.p,b.q,a.q))
27
             return true;
28
        return false;
29
30
   //Area of polygon
31
    double area (const pgon& p) {
32
        double area = 0.0;
33
        for(int i = 1; i < p.size(); ++i)
34
             area += cross(p[i-1],p[i]);
35
        return abs(area)/2.0;
36
37
    //If a \rightarrow b \rightarrow c is a counterclockwise turn
38
    double ccw(const point& a, const point& b,
39
             const point& c){
40
        if (a==b | b==c | a==c)
41
             return false:
42
        point relA = b-a;
43
        point relC = b-c;
44
        return cross (relA, relC) >= 0.0;
45
    //Returns if point p is in the polygon poly
    bool in Poly (const pgon& poly, const pt& p) {
47
        for (int i = 0; i < poly.size()-1; i++){
48
49
             if (! ccw ( poly [ i ] , p , poly [ i +1] ) )
                 return false:
50
51
52
        return true:
53
    //Distance from p to line (a,b)
    double distToLine(const pt& p, const pt& a,
56
             const pt &b){
        vec ap = p-a;
57
        vec ab = b-a:
58
59
        double u = dot(ap, ab)/dot(ab, ab);
60
        //Ignore for non-line segment
        if(u < 0.0) //Closer to a
61
62
             return abs(a-p);
63
        if(u > 1.0) //Closer to b
64
             return abs(b-p);
65
        pt c = a+ab*u; // This is the point
66
        return abs(c-p);
67
    //intersection pts of two circles
```

```
vector<pt> insct(const circ& a, const circ& b){
 70
         vector<pt> o;
 71
         double dist = abs(a.c - b.c);
 72
         if(dist > a.r + b.r)
 73
             return o; //none, don't touch
         if(abs(a.r-b.r) > dist)
 74
 75
             return o; //none, inside
 76
         if(abs(dist - (a.r + b.r)) < EPS) \{ // one intersect \}
 77
             pt p = a.c + (a.r/dist)*(b.c-a.c);
 78
             o.pb(p);
 79
             return o;
 80
         double delta = (sq(dist) + (a.r-b.r)*(a.r+b.r))/(2.0*dist);
 81
 82
         pt cent = a.c + (delta/dist)*(b.c-a.c);
 83
 84
         double h = sqrt(sq(a.r) - sq(delta));
 85
 86
         pt dVec = (b.c - a.c)/dist;
 87
         o.pb(cent + h*pt(0,1)*dVec);
 88
         o.pb(cent + h*pt(0,-1)*dVec);
 89
         return o:
 90
 91
    // intersection of two rectangles, sets none to true if no overlap
    rect overlap (const rect& a, const rect& b, bool& none) {
 94
         rect r:
 95
         if(X(a.p) > X(b.q) | | Y(a.p) > Y(b.q)
 96
             | | X(b.p) > X(a.q) | | Y(b.p) > Y(a.q) ) |
 97
             none=true;
 98
             return r;
99
100
         r.p = \{ \max(X(a.p), X(b.p)), \max(Y(a.p), Y(b.p)) \};
         r.q = \{ min(X(a.q), X(b.q)), min(Y(a.q), Y(b.q)) \};
101
102
         return r;
103
```

# 4 Strings

### **Suffix Array**

```
void countingSort(int k){
        int i, sum, maxi=\max(300,N);
3
        memset(c, 0, sizeof(c));
4
        for (i = 0; i < N; i++)
            c[i+k < N ? RA[i+k] : 0]++;
        for (i=sum=0; i < maxi; i++)
            int t = c[i];
8
            c[i]=sum;
9
            sum+=t;
10
11
        for (i = 0; i < N; i++)
12
            tempSA [c [SA [i]+k < N
                 ? RA[SA[i]+k]: 0]++] = SA[i];
13
14
        for (i=0; i < N; i++)
15
            SA[i] = tempSA[i];
16
17
18
    int main(){
19
        for (int i = 0; i < N; i++)
20
            SA[i]=i, RA[i]=input[i];
21
22
        for (int k = 1; k < N; k <<= 1) {
23
            countingSort(k);
24
            countingSort(0);
25
            tempRA[SA[0]] = r = 0;
            for (int i = 1; i < N; i++){
26
27
                 tempRA [SA [ i ] ]
                     =(RA[SA[i]])==RA[SA[i-1]]
28
29
                     && RA[SA[i]+k] = RA[SA[i-1]+k]
30
                     ? r:++r);
31
32
            for (int i = 0; i < N; i++)
33
                 RA[i] = tempRA[i];
34
35
        return 0;
36
    Trie
    struct node {
2
        node * children [26];
3
        int count;
4
        node(){
            memset(children, 0, sizeof(children));
5
            count = 0;
8
    };
10
   void insert(node* nd, char *s){
        if(*s){
11
```

```
if (!nd->children[*s-'a'])
12
                 nd \rightarrow children [*s-'a'] = new node();
13
14
             insert(nd \rightarrow children[*s-'a'], s+1);
15
16
        nd \rightarrow count ++;
17
18
19
   int count(node* nd, char *s){
20
        if(*s){
21
             if (!nd->children[*s-'a'])
22
                 return 0;
23
             return count (nd->children [*s-'a'], s+1);
24
        } else {
25
             return nd->count;
26
27 }
   KMP
1 vector < int > buildFailure(string s){
2
        vector < int > T(n+1,0);
3
        T[0] = -1;
        int j = 0;
4
 5
        for (int i = 1; i < s.size();++i){
6
             if (s[i]==s[j]) {
                 T[i]=T[j];
 8
                 j++;
9
             } else{
10
                 T[i] = j;
11
                 j = T[j];
12
                 while (j >= 0 \&\& s[i]! = s[j])
13
                     j = T[j];
14
                 j++;
15
16
        T[s.size()] = j;
17
18
        return T;
19
20
   vector < int > search (string W, string S) {
21
        auto T=buildFailure(W);
22
        vector <int> p;
23
        int k = 0;
24
        int j = 0;
25
        \mathbf{while}(j < S.size())
            if(W[k]==S[j])
26
27
                 k++; j++;
28
                 if (k=₩. size ()) {
29
                     p.push_back(j-k);
30
                     k = T[k];
31
32
             }else{
33
                 k = T[k];
34
                 if(k < 0)
35
                     j+=1, k+=1;
36
```

```
37
38
        return p;
39 }
```

# 5 Algorithms

NlogN LIS

#### int ls [MX\_N]; int L[MX\_N]; int I[MX.N]; 5 void nlogn(){ for (int i = 1; i < N+1; ++i) I[i]=INF;8 I[0] = -INF;9 int mx = 1;for (int i = 0; i < N; ++i){ 10 11 int ind = lower\_bound(I, I+N+1, ls[i]) - I; 12 I[ind] = ls[i];13 L[i] = ind;14 mx = max(mx, ind);15 16 int prv = INF;17 vector < int > out; 18 for (int i = N-1; i >= 0; —i){ 19 if(ls[i] < prv && L[i]==mx){ 20 out.push\_back(ls[i]); 21 prv = ls[i];22 mx--;23 24

# RectInHist

25

```
1 int R,C;
   char board [MX_RC] [MX_RC];
   int h[MX_RC][MX_RC];
5
   int perim(int 1, int w){
        if ( l==0 | | w==0)
6
            return 0:
8
        return 2*1 + 2*w;
9
10
11 int main(){
        for (int i = 0; i < R; i++){
12
13
            int run=0;
14
            for (int j = 0; j < C; j++){
15
                run = (board[i][j] == '. '?run + 1:0);
16
                h[i][j] = run;
17
18
        int mx = 0;
19
20
        for (int j = 0; j < C; j++){
21
            stack<int> s;
22
            for (int i = 0; i < R; i++){
```

```
23
                  if (s.empty()
24
                           | | h [ i ] [ j ] > h [ s.top() ] [ j ])
25
                      s.push(i);
26
                 else if(h[i][j]<h[s.top()][j]){
27
                      while (!s.empty()
                      &&h[i][j]<h[s.top()][j]){
28
                          int l = h[s.top()][j];
29
30
                           s.pop();
31
                          int pm = perim(1,
32
                               (s.empty()?
33
                                i : i - s \cdot top() - 1);
34
                          mx = max(mx,pm);
35
36
                      s.push(i);
37
                 \} else if (h[i][j]==h[s.top()][j])
38
                      s.pop();
39
                      s.push(i);
40
41
42
             while (! s . empty()) {
43
                 int l = h[s.top()][j]; s.pop();
44
                 int pm = perim(1, s.empty() ? R : R - s.top()-1);
45
                 mx = max(mx, pm);
46
47
48
         printf("%d\n",mx-1);
49
```

## 6 Maths

#### Miller Rabin

```
void factor(ll x, ll& e, ll& k){
2
         while (x%2LL==0LL) {
3
             x/=2LL;
 4
             ++e;
 6
        k = x:
    //increase x for higher certainty, 5 works well
    bool is_prime(ll n, int x){
10
11
         if (n&2LL==0 | | n==1LL)
12
             return false;
13
         if (n==2 || n==3 || n==5 || n==7)
14
             return true;
15
         ll e, k;
16
         factor (n-1,e,k);
17
         while (x-->0) {
18
             11 \ a = (rand())\%(n-5LL) + 2LL;
19
             ll p = mod_exp(a,k,n);
20
             if (p==1LL | p==n-1LL)
21
                  continue;
22
             bool all_fail = true;
23
             for (int i = 0; i < e-1; ++i){
24
                  p = mod_exp(p, 2, n);
25
                  \mathbf{i} \mathbf{f} (p=n-1LL) 
26
                       all_fail = false;
27
                       break:
28
29
30
             if (all_fail)
31
                  return false;
32
33
         return true;
34
   Binomial Coefficients
   ll ncrmem [MX_N] [MX_N];
2
   ll ncr(int n, int r){
             if(n==0)
 5
                       return r==0;
 6
             if(r==0)
                       return 1;
             \mathbf{if} (\operatorname{ncrmem} [n] [r] != -1)
9
                       return ncrmem[n][r];
10
             return \operatorname{ncrmem}[n][r] = \operatorname{ncr}(n-1, r-1) + \operatorname{ncr}(n-1, r);
11
```

### Gaussian Elimination

```
1
    * mat is augmented matrix
3
    * e.g 3x + 4y = 2 is [3,4,2]
4
    void gauss (double mat [MX_N] [MX_N+1], double ans [MX_N], int n) {
        int i, j, k, l; double t;
6
8
        for (j = 0; j < n-1; ++j)
9
            1 = j;
10
            for (i = j+1; i < n; ++i)
11
                 if (fabs (mat[i][j]) > fabs (mat[l][j]))
12
                     l=i:
13
            for(k = j; k \le n; ++k)
14
15
                t=mat[j][k], mat[j][k]=mat[l][k], mat[l][k]=t;
16
            for (i = j+1; i < n; ++i)
17
                for(k = n; k >= j; ---k)
18
                     mat[i][k] = mat[j][k] * mat[i][j] / mat[j][j];
19
20
21
        for (j = n-1; j >= 0; ---j)
22
            for (t=0.0, k=j+1; k< n; ++k)
23
                t += mat[j][k] * ans[k];
24
            ans[j] = (mat[j][n] - t) / mat[j][j];
25
26
    Ternary Search
1
    double ternary_search (double 1, double r) { //maximises
2
        while (r - l > EPS) {
3
            double m1 = 1 + (r - 1) / 3;
4
            double m2 = r - (r - 1) / 3;
            double f1 = f(m1);
            double f2 = f(m2);
            if (f1 < f2)
                l = m1;
            else
10
                r = m2:
11
12
        return f(l);
13
   Matrix Exponential
   /* c = a * b * /
   void mu(11 a[][NMAT], 11 b[][NMAT], 11 c[][NMAT], int _n) {
            for (int i = 0; i < n; i++)
                     for (int j=0; j<_n; j++) {
                             c[i][j]=0;
                             for (int h=0;h<_n;h++) {
                                      c[i][j]+=(a[i][h]*b[h][j])%mod;
8
                                      c [ i ] [ j]%=mod;
9
```

```
10
11
12
    /* returns ans=mat^b*/
   void power(ll ans[][NMAT], ll mat[][NMAT], ll b, int _n) {
             11 tmp[NMAT][NMAT];
15
16
             for (int i = 0; i < n; i++)
17
                      for (int j = 0; j < n; j + +)
18
                               ans [i][j]=i=j;
             while(b) {
19
20
                      if (b&1) {
21
                               mu(ans, mat, tmp, _n);
22
                               for (int i = 0; i < n; i++)
23
                                        for (int j=0; j<_n; j++)
24
                                                 ans[i][j]=tmp[i][j];
25
26
                      mu(mat, mat, tmp, _n);
27
                      for (int i=0; i<_n; i++)
28
                               for (int j=0; j<_n; j++)
29
                                        mat[i][j]=tmp[i][j];
30
                      b >> = 1;
31
32
    Exponents
 1 // a^b % md
    ll modExp(ll a, int b, ll mod){
        \mathbf{i} \mathbf{f} (b==0)
             return 1LL;
```

```
1  // a*b % md
2  ll modExp(ll a, int b, ll mod){
3     if (b==0)
4         return 1LL;
5     if (b%2==0){
6         ll y = modExp(a, b/2, mod);
7         return (y*y)%mod;
8     }
9     return (a*modExp(a,b-1,mod))%mod;
10  }
11
12  // a*b^-1 % mod
13  // mod should be prime
14  ll modDiv(ll a, ll b, ll mod) {
15     return (a * modExp(b, mod-2, mod))%mod;
16  }
```

### Theorems

#### Burnside's Lemma

Let G be a finite group that acts on a set X. For each g let  $X^g$  denote the set of elements in X by g.

$$X^g = \{ x \in X \mid g.x = x \}$$

Then the number of orbits is as follows

$$|X/G| = \frac{1}{|G|} \sum_{g \in G} |X^g|$$

#### Polya's enumeration theorem

Let X be a finite set and let G be a group of permutations of X. Let Y be a finite set of colors so that  $Y^X$  is the set of colored arrangements of X. Then the number of orbits of  $Y^X$  under G is

$$|Y^X/G| = \frac{1}{|G|} \sum_{g \in G} |Y|^{c(g)}$$

where c(q) is the number of cycles in permutation g.