

# Real-time Process Network Sonar Beamformer



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## Introduction

- **Problem:** Beamforming is computationally intensive
- **Example:** High-resolution sonar beamformers (GFLOPS)
  - Generation 1: expensive custom digital hardware (large state machines)
  - Generation 2: custom integration of programmable digital signal processors on commercial-off-the-shelf hardware (e.g. 120 DSPs in a VME rack)
- **Objective:** Unix workstation beamformers
  - Analysis: evaluate performance of beamforming kernels and systems
  - Modeling: capture parallelism, guarantee determinate bounded execution
  - Implementation: use portable, scalable software to achieve real-time performance on commodity hardware and lower development costs
- **Solution:** Real-time beamforming on workstations
  - Analysis: optimize kernels and profile beamformers to measure scalability
  - Modeling: Process Networks
  - Implementation: Real-time POSIX threads using C++ on symmetric multiprocessor UltraSPARC-II workstation with native signal processing

## Time-Domain Beamforming

- Delay and sum weighted sensor outputs
- Geometrically project the sensor elements onto a line to compute the time delays

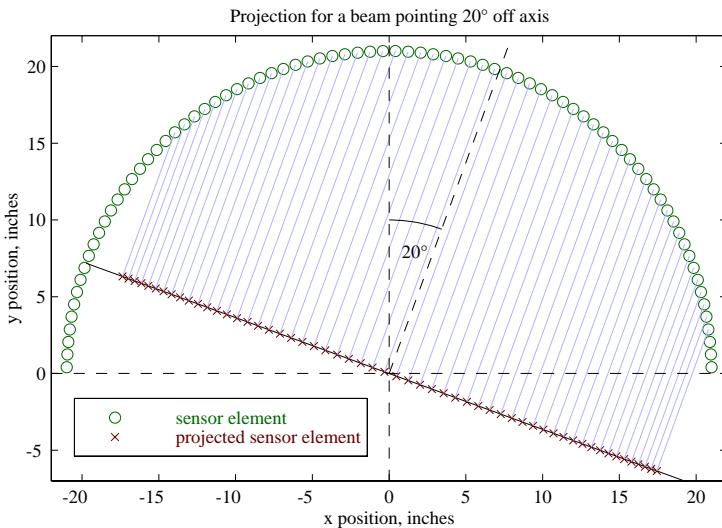
$$b(t) = \sum_{i=1}^M \alpha_i x_i(t - \tau_i)$$

$b(t)$  beam output

$x_i(t)$   $i$ th sensor output

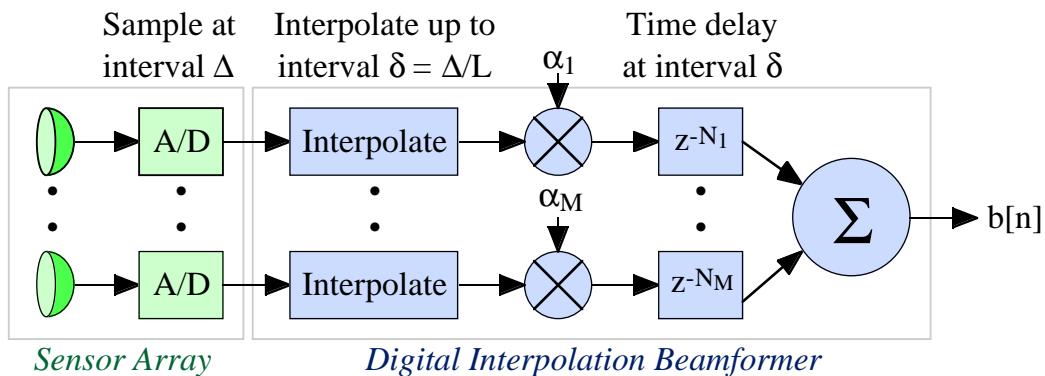
$\tau_i$   $i$ th sensor delay

$\alpha_i$   $i$ th sensor weight



## Interpolation Beamforming

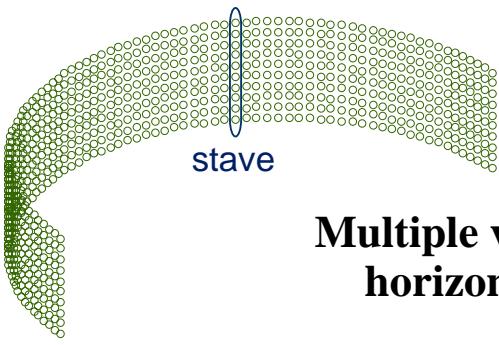
- Quantized time delays perturb beam pattern
- Sample at just above the Nyquist rate
- *Interpolate* to obtain desired time-delay resolution



- Kernel implementation on UltraSPARC-II

- Highly optimized C++ (loop unrolling and SPARCompiler5.0EA)
- Currently operating at 60% of peak, which is 2 FLOPs per cycle

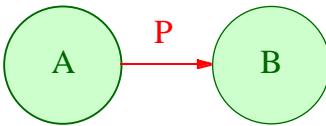
## Vertical Beamforming



Multiple vertical transducers for every horizontal position

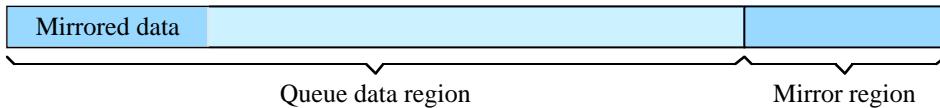
- Each vertical sensor column is combined into a *stave*
  - No time delay or interpolation is required
  - Staves are calculated by a simple *integer* dot product
  - Integer-to-float conversion must be performed
  - Output data must be interleaved
- Kernel implementation on UltraSPARC-II
  - Native signal processing with Visual Instruction Set (VIS)
  - Software data prefetching to hide memory latency and keep the pipeline full

## Formal Design Methodology

- The *Process Network* model [Kahn, 1974]
  - Superset of dataflow models of computation
  - Captures concurrency and parallelism
  - Provides correctness
  - Guarantees determinate execution of the program
- A program is represented as a directed graph
  - Each **node** is an independent process
  - Each **edge** is a one-way queue of dataA diagram showing two green circular nodes labeled 'A' and 'B'. A red arrow labeled 'P' points from node A to node B, representing a directed edge in a process network graph.
- Blocking reads, non-blocking writes, infinite queues
- Scheduling for bounded queues is possible [Parks, 1995]
  - Blocking reads and writes
  - Dynamically increase queue capacities to prevent *artificial deadlock*
- Fits the thread model of concurrent programming

## Process Network Implementation

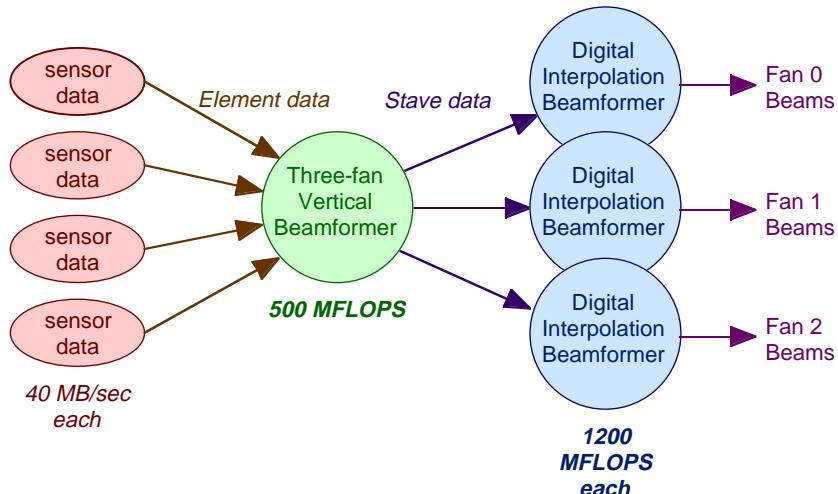
- Each node corresponds to a thread
  - Implemented in C++ using POSIX Pthreads
  - Low-overhead, high-performance, scalable
  - Granularity larger than a thread context switch (~10 us)
  - Symmetric multiprocessing operating system dynamically schedules threads
  - Efficient utilization of multiple processors
- Optimize queues for high-throughput signal processing
  - Nodes operate directly on queue memory, avoiding copying
  - Queues use mirroring to keep data contiguous



- Compensates for the lack of circular address buffers
- Queues trade-off memory usage for overhead
- Virtual memory manager maintains data circularity

## System Implementation

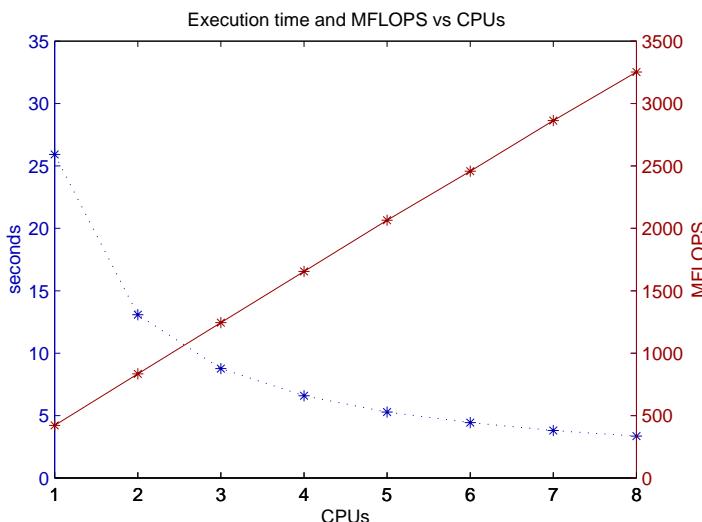
- Vertical beamformer forms 3 sets of 80 staves from 10 vertical elements per stave
- Each horizontal beamformer forms 61 beams from the 80 staves, using a two-point interpolation filter
- 4 GFLOPS total computation



## Performance Results

- 100 trial mean execution time for 2.6 seconds of data
- Sun Ultra Enterprise 4000 with eight 336-MHz UltraSPARC-IIIs, 2 Gb RAM, running Solaris 2.6

Implementation	Time (s)	MFLOPS	Mbytes
thread pool	3.607	3024.8	832
process network	3.354	3252.8	654



- Process network is 7% faster than thread pool, overhead is small
- Process network uses 20% less memory with lower latency
- Process network scalability is nearly linear
- Will continue to scale with additional CPUs
- Real-time performance achievable with 12 CPUs

## Conclusion

- Third generation beamformers: *Workstation* hardware
  - Commodity hardware saves development/manufacturing costs
  - Multiprocessor servers, native signal processing
  - Upgradable hardware, Moore's Law
- Software model: *Process Networks*
  - Captures parallelism, guarantees determinate bounded execution
  - Portable, reusable, scalable C++ code
  - High-performance, low overhead POSIX threads
  - Symmetric multiprocessing operating system
- The example 4-GFLOPS 1-Gb 3-D sonar beamforming system does execute in real time using a Sun Ultra Enterprise 4000 server with twelve 336 MHz UltraSPARC-II CPUs with 14.5% to spare