

A Comparison of Parallel Workstation Sonar Beamforming Implementations

Gregory E. Allen

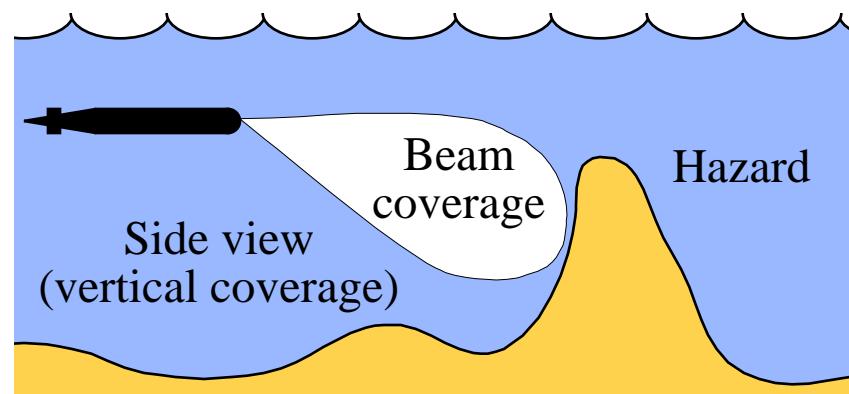
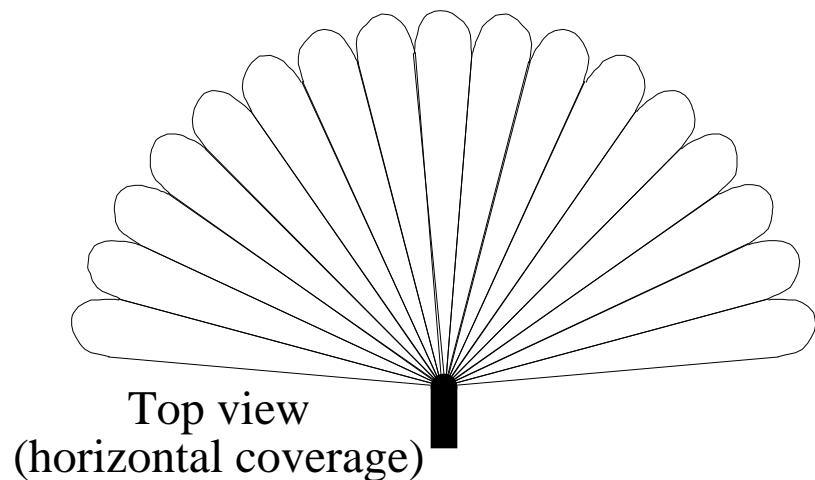
**Applied Research Laboratories
The University of Texas at Austin**



<http://www.ece.utexas.edu/~allen/>

What is Beamforming?

- A **beamformer** is a spatial filter that operates on the output of an array of sensors
 - Enhances sound from a desired direction, rejects others
 - Filter design chooses which direction a beam points
 - We can determine from which direction a sound is coming
- Many beams formed, each in a different direction



Time-Domain Beamforming

- Delay-and-sum weighted sensor outputs
- Geometrically project the sensor elements onto a line to compute the time delays

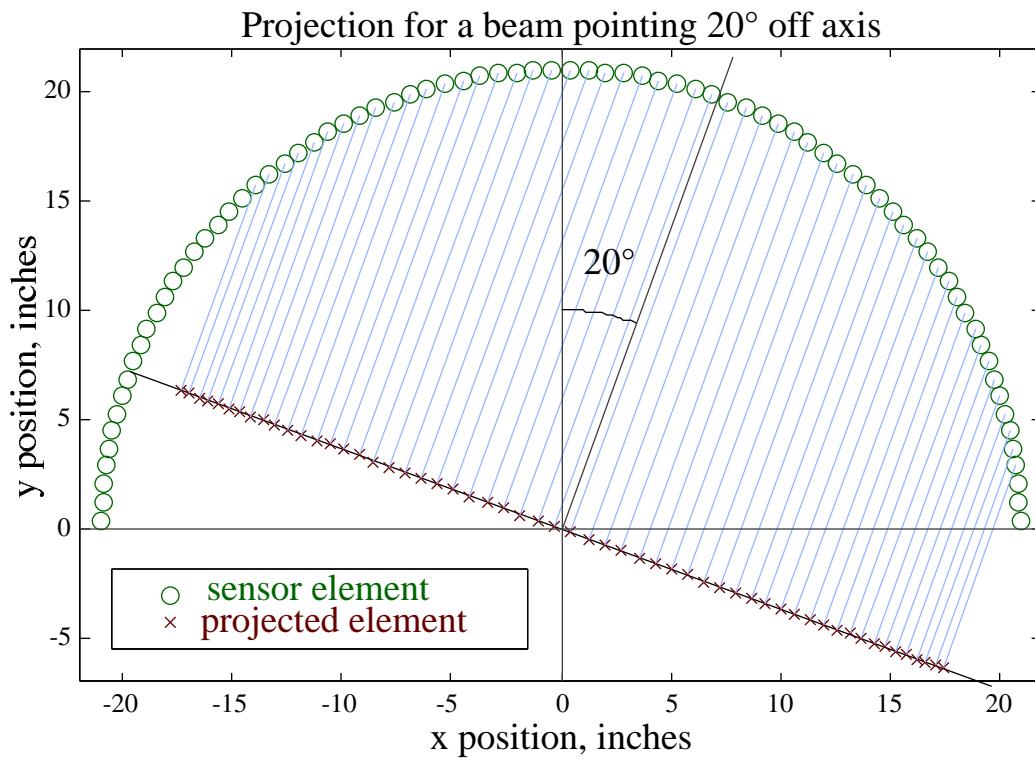
$$b(t) = \sum_{i=1}^M w_i x_i(t - \tau_i)$$

$b(t)$ beam output

$x_i(t)$ i^{th} sensor output

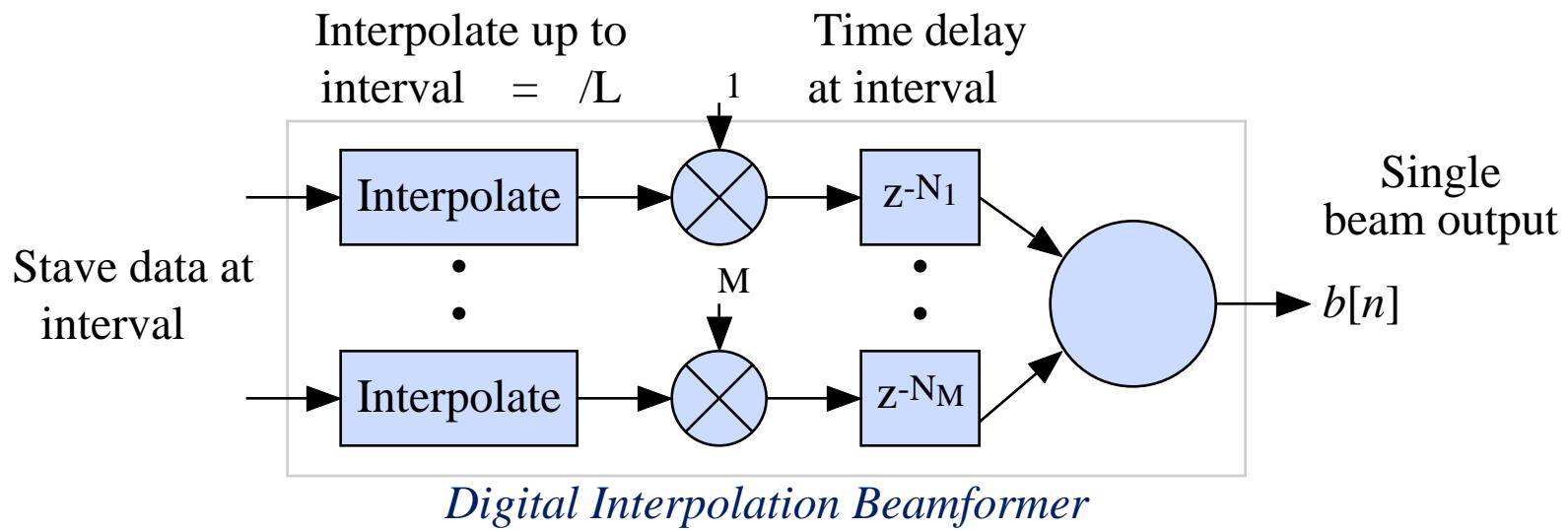
τ_i i^{th} sensor delay

w_i i^{th} sensor weight



Interpolation Beamforming

- Quantized time delays distort the beam pattern
- Sample at just above the Nyquist rate, interpolate to obtain desired time delay resolution



- Multiple different beams formed from same data

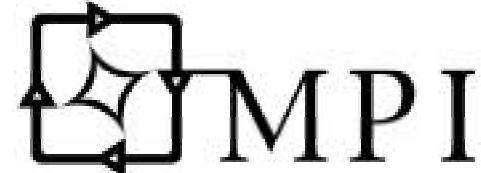
Motivation

- **High-performance, low production volume (~100 MB/s I/O; 1-20 GFLOPS; under 50 units)**
- **Current real-time implementation technologies**
 - **Custom hardware**
 - **Custom integration using commercial-off-the-shelf (COTS) processors (e.g. 100 digital signal processors in a VME chassis)**
- **Wish to target commodity workstations**
 - **Symmetric multiprocessing (SMP) operating systems**
 - **Leverage native signal processing (NSP) kernels**
 - **Development environment and target architecture are same**
 - **Concurrent development on less powerful workstations**
 - **Reduce development time and cost**

Objective

- Given an UltraSPARC-II beamforming kernel...
 - Highly optimized C++ (loop unrolling and SPARCompiler5.0)
 - Operates at 440 MFLOPS at 336 MHz (60% of peak)
- Compare beamforming performance using different frameworks for parallelism
 - Message Passing Interface (MPI)
 - Computational Process Networks (CPN)
 - Extend CPN implementation with MPI: a hybrid
- Measure performance on Sun Ultra Enterprise 4000
 - Eight 336 MHz UltraSPARC-II processors
 - 2 GB RAM, Solaris 2.6

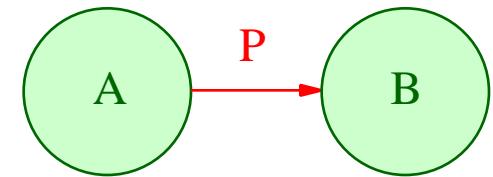
Message Passing Interface



- A standard interface for
 - Explicit message passing in application programs
 - MIMD distributed memory concurrent computers
- A library for C or Fortran, developed by about 80 people from 40 organizations (edu, gov, com)
- Intended to be portable and easy to use
- Many implementations exist, free and commercial
- Using MPI from Sun's HPC 2.0 Package
 - Based on MPICH (public from ANL)
 - Claims to be thread safe

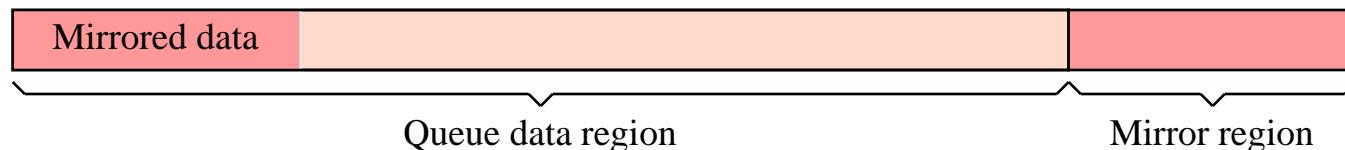
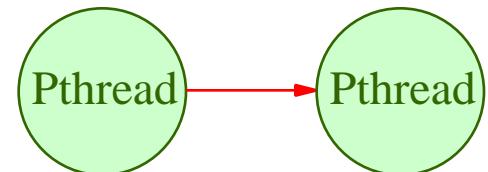
Computational Process Networks

- Based on formal **Process Network** model [Kahn, 1974]
 - Program is represented as a directed graph
 - Captures concurrency and parallelism
 - Provable model provides correctness and determinate execution
- Leverage **bounded scheduling** [Parks, 1995]
 - Permits realization in finite memory, regardless of scheduler
- Extend this model with **firing thresholds**, similar to **Computation Graphs** [Karp & Miller, 1966]
 - Models algorithms on overlapping continuous streams of data, e.g. digital filters and fast Fourier transforms (FFTs)
 - Decouples computation (node) from communication (queue)
 - Allows compositional parallel programming



High-Performance Implementation

- Designed for **real-time high-throughput** systems
 - Uses **POSIX lightweight Pthreads**
 - Each node corresponds to a thread
 - Portable to many different operating systems
 - Optional fixed-priority real-time scheduling
 - Operate **directly on queue memory** to avoid copying
 - Queues use mirroring to keep data contiguous



- Compensates for lack of hardware support for circular buffers (e.g. modulo addressing in DSPs)
 - Virtual memory manager keeps data circularity in hardware

Process Networks with MPI

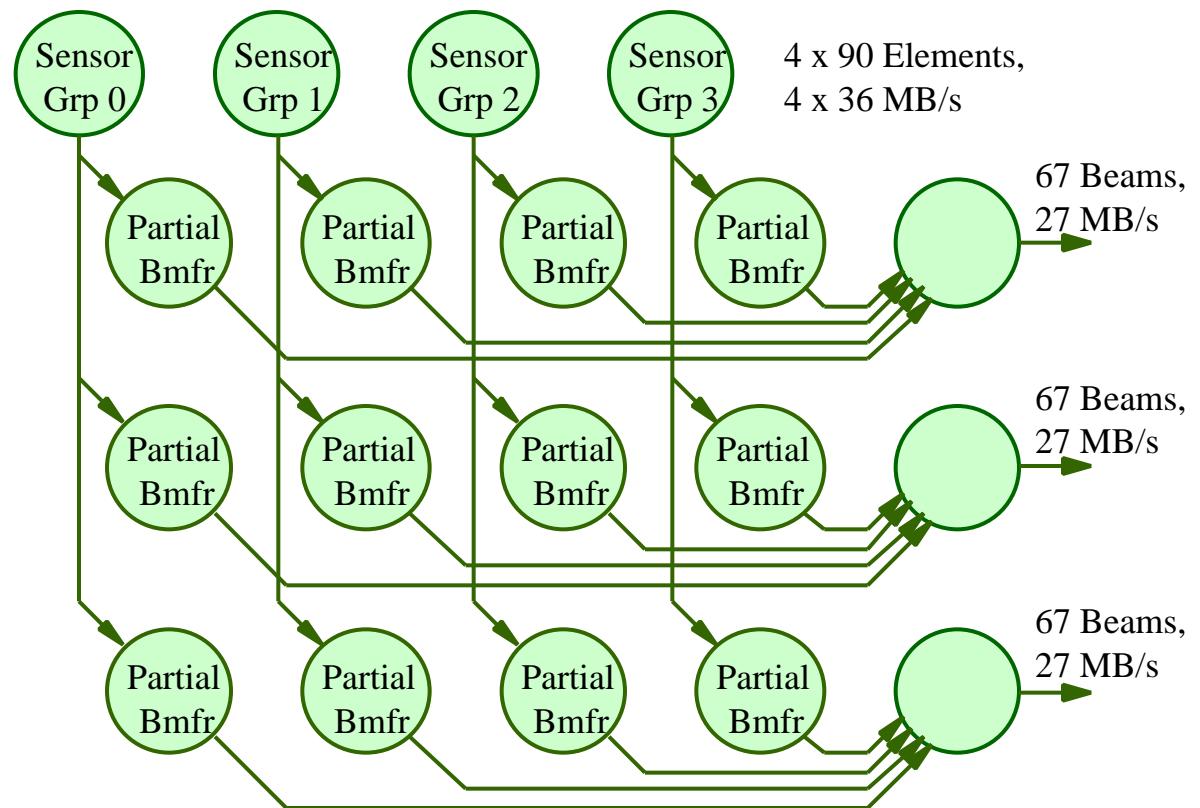
- Threads require shared memory, MPI does not
- Allow a queue to be implemented across MPI



- Easy to implement, but
- Programming style must change (threaded vs. SPMD)
- Mapping nodes across processes and matching up queues is difficult, and best left to automated tools
- Other point-to-point technologies may be a better fit for Process Networks (sockets, RACEway)

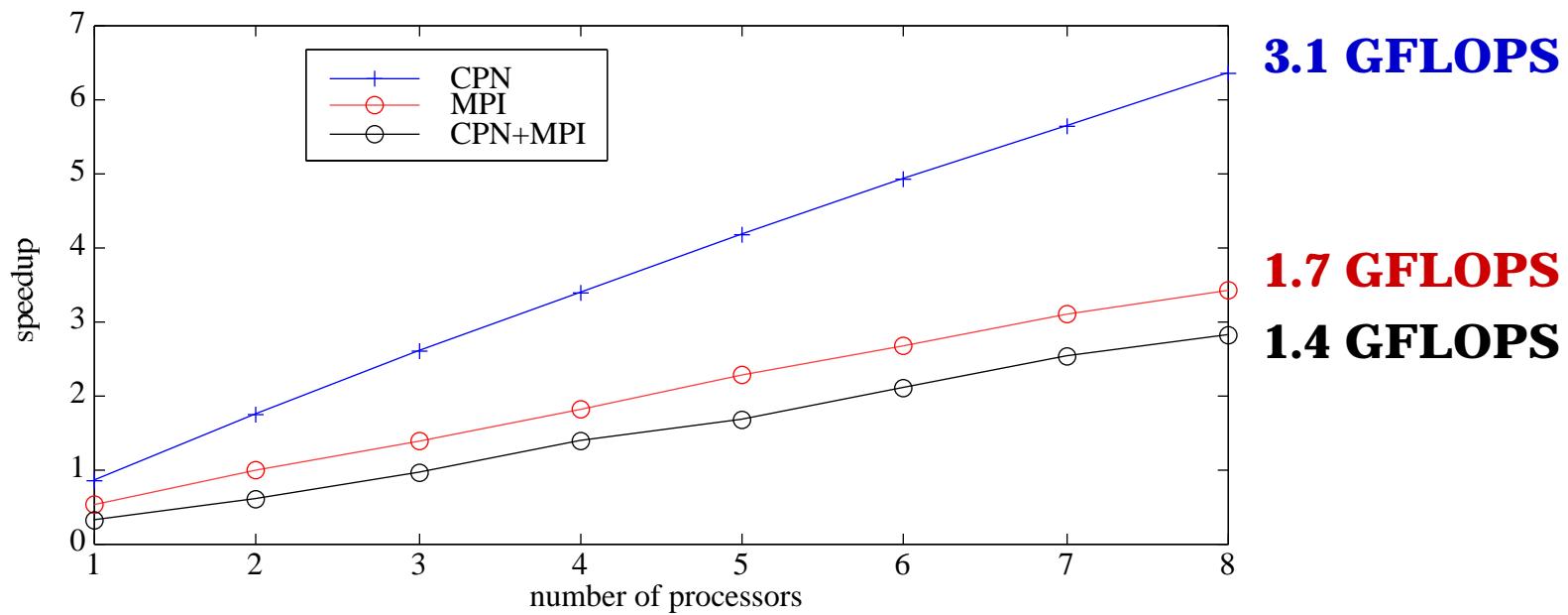
Beamformer Block Diagram

- Calculate 201 beams from 360 sensor elements
 - 21 GFLOPS total, requires about fifty 336 MHz UltraSPARC-IIIs
 - 144 MB/s sensor data in, 81 MB/s beam data out
- **Partial beamforming** divides the problem functionally
- Each node calculates different part, results are summed
- Average node is 1.75 GFLOPS (2.4 max)
- Each part needs to operate at real-time
- Workstation cluster



Performance Results

- Benchmarked on a single SMP machine
- Results as compared to sequential case (480 MFLOPS)
- Slowdown on one processor: **CPN 16%, MPI 84%**
- Speedup on eight processors: **CPN 6.5, MPI 3.5**



My Comments on MPI

- MPI is straightforward to use, and a step in the right direction, but
 - It lacks any formal methodology
 - It needs a C++ class interface (added in MPI-2)
 - Type handling is very messy
 - Does not leverage lightweight threads (smallest unit of computation is the process)
 - Strangely absent from the commercial embedded real-time community
- Sun's Implementation of MPI
 - \$375 run-time license per CPU (with educational discount)
 - Breaks CPN hardware data circularity (~10% penalty)

Conclusion

- Benchmark on eight processor 336 MHz SMP Sun
 - Use highly optimized beamforming kernel (single processor)
 - Compare performance within several parallel frameworks
- Process Network outperforms MPI significantly
 - Slowdown on one processor: CPN 16%, MPI 84%
 - Speedup on eight processors: CPN 6.5, MPI 3.5
 - CPN is 86% faster than MPI on eight processors
 - The hybrid CPN / MPI is 18% slower than MPI alone
 - CPN is 58% of peak performance (2 FLOPS x 336 MHz x 8 CPUs)
- If I needed to develop a real-time implementation today, I probably would not use MPI