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# Karnaugh Maps (K-Maps)

# Gray Code

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<i>Decimal number</i>	<i>Gray</i>				<i>Binary</i>			
	$g_3$	$g_2$	$g_1$	$g_0$	$b_3$	$b_2$	$b_1$	$b_0$
0	0	0	0	0	0	0	0	0
1	0	0	0	1	0	0	0	1
2	0	0	1	1	0	0	1	0
3	0	0	1	0	0	0	1	1
4	0	1	1	0	0	1	0	0
5	0	1	1	1	0	1	0	1
6	0	1	0	1	0	1	1	0
7	0	1	0	0	0	1	1	1
8	1	1	0	0	1	0	0	0
9	1	1	0	1	1	0	0	1
10	1	1	1	1	1	0	1	0
11	1	1	1	0	1	0	1	1
12	1	0	1	0	1	1	0	0
13	1	0	1	1	1	1	0	1
14	1	0	0	1	1	1	1	0
15	1	0	0	0	1	1	1	1

# K-Maps

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- A visual way to simplify logic expressions
- It gives the most simplified form of the expression
- A Karnaugh map is a graphical method used to obtain the most simplified form of an expression in a standard form (Sum-of-Products or Product-of-Sums).
- The simplest form of an expression is the one that has the minimum number of terms with the least number of literals (variables) in each term, thus the function will be implemented with the minimum number of gates with minimum number of input.

# Three-Variable K-Maps

$$f = \sum(0,4) = \overline{B} \overline{C}$$

A \ BC	00	01	11	10
	0	1	1	0
0	1	0	0	0
1	1	0	0	0

$$f = \sum(4,5) = A \overline{B}$$

A \ BC	00	01	11	10
	0	1	1	0
0	0	0	0	0
1	1	1	0	0

$$f = \sum(0,1,4,5) = \overline{B}$$

A \ BC	00	01	11	10
	0	1	1	0
0	1	1	0	0
1	1	1	0	0

$$f = \sum(0,1,2,3) = \overline{A}$$

A \ BC	00	01	11	10
	0	1	1	0
0	1	1	1	1
1	0	0	0	0

$$f = \sum(0,4) = \overline{A} C$$

A \ BC	00	01	11	10
	0	1	1	0
0	0	1	1	0
1	0	0	0	0

$$f = \sum(4,6) = A \overline{C}$$

A \ BC	00	01	11	10
	0	1	1	0
0	0	0	0	0
1	1	0	0	1

$$f = \sum(0,2) = \overline{A} \overline{C}$$

A \ BC	00	01	11	10
	0	1	1	0
0	1	0	0	1
1	0	0	0	0

$$f = \sum(0,2,4,6) = \overline{C}$$

A \ BC	00	01	11	10
	0	1	1	0
0	1	0	0	1
1	1	0	0	1

# Three-Variable K-Maps

A \ BC	BC			
	00	01	11	10
0		1		
1	1		1	1

A \ BC	BC			
	00	01	11	10
0	1		1	1
1	1			1

A \ BC	BC			
	00	01	11	10
0			1	1
1	1	1		

A \ BC	BC			
	00	01	11	10
0			1	
1	1		1	1

A \ BC	BC			
	00	01	11	10
0		1	1	1
1		1	1	

A \ BC	BC			
	00	01	11	10
0				
1				

# Four-Variable K-Maps

AB \ CD				
	00	01	11	10
00	1	0	0	0
01	0	0	0	0
11	0	0	0	0
10	1	0	0	0

$$f = \sum(0,8) = \bar{B} \cdot \bar{C} \cdot \bar{D}$$

AB \ CD				
	00	01	11	10
00	0	0	0	0
01	0	1	0	0
11	0	1	0	0
10	0	0	0	0

$$f = \sum(5,13) = B \cdot \bar{C} \cdot D$$

AB \ CD				
	00	01	11	10
00	0	0	0	0
01	0	0	0	0
11	0	1	1	0
10	0	0	0	0

$$f = \sum(13,15) = A \cdot B \cdot D$$

AB \ CD				
	00	01	11	10
00	0	0	0	0
01	1	0	0	1
11	0	0	0	0
10	0	0	0	0

$$f = \sum(4,6) = \bar{A} \cdot B \cdot \bar{D}$$

AB \ CD				
	00	01	11	10
00	0	0	1	1
01	0	0	1	1
11	0	0	0	0
10	0	0	0	0

$$f = \sum(2,3,6,7) = \bar{A} \cdot C$$

AB \ CD				
	00	01	11	10
00	0	0	0	0
01	1	0	0	1
11	1	0	0	1
10	0	0	0	0

$$f = \sum(4,6,12,14) = B \cdot \bar{D}$$

AB \ CD				
	00	01	11	10
00	0	0	1	1
01	0	0	0	0
11	0	0	0	0
10	0	0	1	1

$$f = \sum(2,3,10,11) = \bar{B} \cdot C$$

AB \ CD				
	00	01	11	10
00	1	0	0	1
01	0	0	0	0
11	0	0	0	0
10	1	0	0	1

$$f = \sum(0,2,8,10) = \bar{B} \cdot \bar{D}$$

# Four-Variable K-Maps

AB \ CD	CD			
	00	01	11	10
00	0	0	0	0
01	1	1	1	1
11	0	0	0	0
10	0	0	0	0

$$f = \sum(4,5,6,7) = \bar{A} \bullet B$$

AB \ CD	CD			
	00	01	11	10
00	0	0	1	0
01	0	0	1	0
11	0	0	1	0
10	0	0	1	0

$$f = \sum(3,7,11,15) = C \bullet D$$

AB \ CD	CD			
	00	01	11	10
00	1	0	1	0
01	0	1	0	1
11	1	0	1	0
10	0	1	0	1

$$f = \sum(0,3,5,6,9,10,12,15)$$

$$f = A \otimes B \otimes C \otimes D$$

AB \ CD	CD			
	00	01	11	10
00	0	1	0	1
01	1	0	1	0
11	0	1	0	1
10	1	0	1	0

$$f = \sum(1,2,4,7,8,11,13,14)$$

$$f = A \oplus B \oplus C \oplus D$$

AB \ CD	CD			
	00	01	11	10
00	0	1	1	0
01	0	1	1	0
11	0	1	1	0
10	0	1	1	0

$$f = \sum(1,3,5,7,9,11,13,15)$$

$$f = D$$

AB \ CD	CD			
	00	01	11	10
00	1	0	0	1
01	1	0	0	1
11	1	0	0	1
10	1	0	0	1

$$f = \sum(0,2,4,6,8,10,12,14)$$

$$f = \bar{D}$$

AB \ CD	CD			
	00	01	11	10
00	0	0	0	0
01	1	1	1	1
11	1	1	1	1
10	0	0	0	0

$$f = \sum(4,5,6,7,12,13,14,15)$$

$$f = B$$

AB \ CD	CD			
	00	01	11	10
00	1	1	1	1
01	0	0	0	0
11	0	0	0	0
10	1	1	1	1

$$f = \sum(0,1,2,3,8,9,10,11)$$

$$f = \bar{B}$$

# Design of combinational digital circuits

- Steps to design a combinational digital circuit:
  - From the problem statement derive the truth table
  - From the truth table derive the unsimplified logic expression
  - Simplify the logic expression
  - From the simplified expression draw the logic circuit
- Example: Design a 3-input (A,B,C) digital circuit that will give at its output (X) a logic 1 only if the binary number formed at the input has more ones than zeros.

	Inputs			Output
	A	B	C	X
0	0	0	0	0
1	0	0	1	0
2	0	1	0	0
3	0	1	1	1
4	1	0	0	0
5	1	0	1	1
6	1	1	0	1
7	1	1	1	1

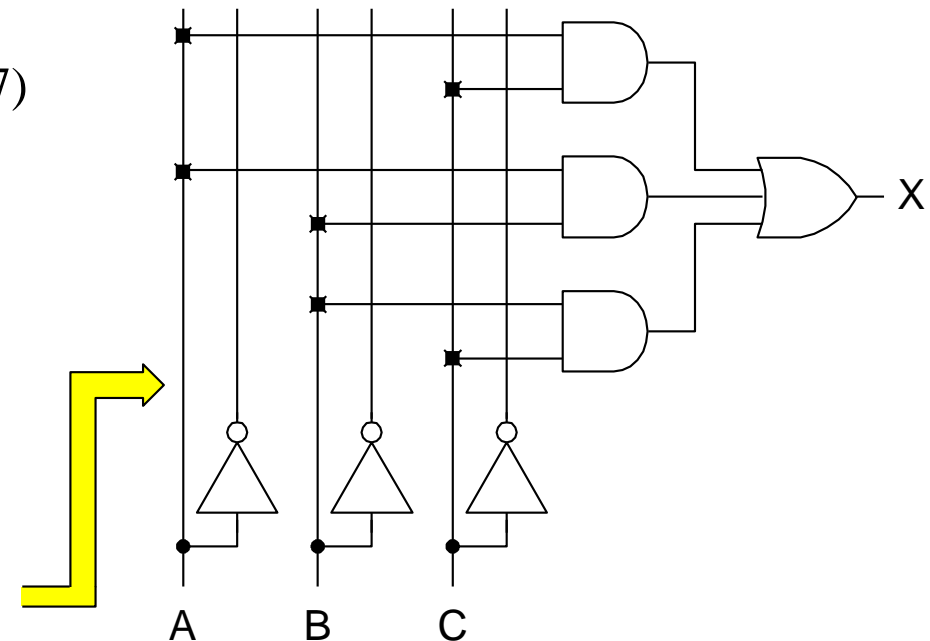
→  $X = \sum (3, 5, 6, 7)$

↓

A \ BC	00	01	11	10
0	0	0	1	0
1	0	1	1	1

↓

$X = AC + AB + BC$





# Design of combinational digital circuits

- Example: Design a 4-input (A,B,C,D) digital circuit that will give at its output (X) a logic 1 only if the binary number formed at the input is between 2 and 9 (including).

	Inputs				Output
	A	B	C	D	X
0	0	0	0	0	0
1	0	0	0	1	0
2	0	0	1	0	1
3	0	0	1	1	1
4	0	1	0	0	1
5	0	1	0	1	1
6	0	1	1	0	1
7	0	1	1	1	1
8	1	0	0	0	1
9	1	0	0	1	1
10	1	0	1	0	0
11	1	0	1	1	0
12	1	1	0	0	0
13	1	1	0	1	0
14	1	1	1	0	0
15	1	1	1	1	0

→  $X = \sum (2,3,4,5,6,7,8,9)$

↓

CD \ AB				
	00	01	11	10
00	0	0	1	1
01	1	1	1	1
11	0	0	0	0
10	1	1	0	0

↓

$$X = \bar{A}C + \bar{A}B + A\bar{B}\bar{C}$$

