ECE297 Milestone 4
Traveling Courier

"A person who never made a mistake never tried anything new."

–Albert Einstein

Assigned on Friday, April 1 — Autotester = 12/12

Due on April 14 or 16 (Poll) — **Total marks = 12/100**

1 Objectives

In this milestone you will find a path through your map that has multiple stops, so that you can find a good route for a courier company driver who has multiple deliveries to make. This is a variation on a classic optimization problem called the traveling salesman problem.

After completing this milestone you should be able to:

1	Find a valid path through a graph that passes through a set of vertices.

2	Develop heuristics to solve a computationally hard problem.

2 Overview and Background

In this milestone you will extend your project with `m4.h` and one or more .cpp files and use these files to implement a variation of the traveling salesman algorithm.

The traveling salesman problem is *computationally hard* which means that there is no known algorithm that (1) gives a guaranteed optimal (lowest travel time) result and (2) has a computational complexity that is a polynomial of N, where N is the number of vertices the salesman must visit. In practice this means that guaranteed optimal algorithms have computational complexity at least exponential, $O(2^N)$, in the problem size and become impractically slow for large enough N. Therefore we must resort to heuristic (i.e. "seems like a good idea") methods that do not guarantee a perfect answer, but which can run much faster. Most optimization problems that are actively researched are computationally hard; coming up with better heuristics for such problems is important in many fields including the design of computer chips, transportation systems, and new pharamaceutical drugs.

3 Detailed Specification

Figure 1 shows the input you are given.

Figure 1: Input to your courier delivery algorithm. In this example you are given 3 depot intersections, and 4 deliveries to make. Each delivery has a pick up intersection (shown here as a letter followed by an up arrow) and a corresponding drop off location (shown here as the same letter with a down arrow).

The courier company has a set of M depots for their delivery trucks, and you can start your day's deliveries from any one of those M depots. At the end of the day you need to return your truck to one of the depots, but it doesn't have to be the same one at which you started.

After picking up your delivery truck, you need to make N deliveries, where each delivery has an intersection at which you pick up the package, and another intersection where you drop off the corresponding package. You can visit these intersections in any order you like, but you must visit the intersections such that all the deliveries are made. A delivery is made when you visit the pick up intersection to pick the package, and then some time later visit the corresponding drop off intersection. It is possible that a single intersection may appear more than once as a pick up location; in that case you pick up all the packages automatically when you visit that intersection. Similarly, an intersection could appear more than once as a drop off location; in that case, when you visit the intersection you will drop off all the packages you have already picked up that need to go to that intersection.

Figure 2 shows the output your algorithm must generate – a delivery route that that begins at a depot, follows connected street_segments to reach intersections such that all N deliveries are made, and ends at a (possibly the same or possibly a different) depot.

We will unit test your algorithm by calling your `travelingCourier` function which must have the function prototype shown in Listing 1.

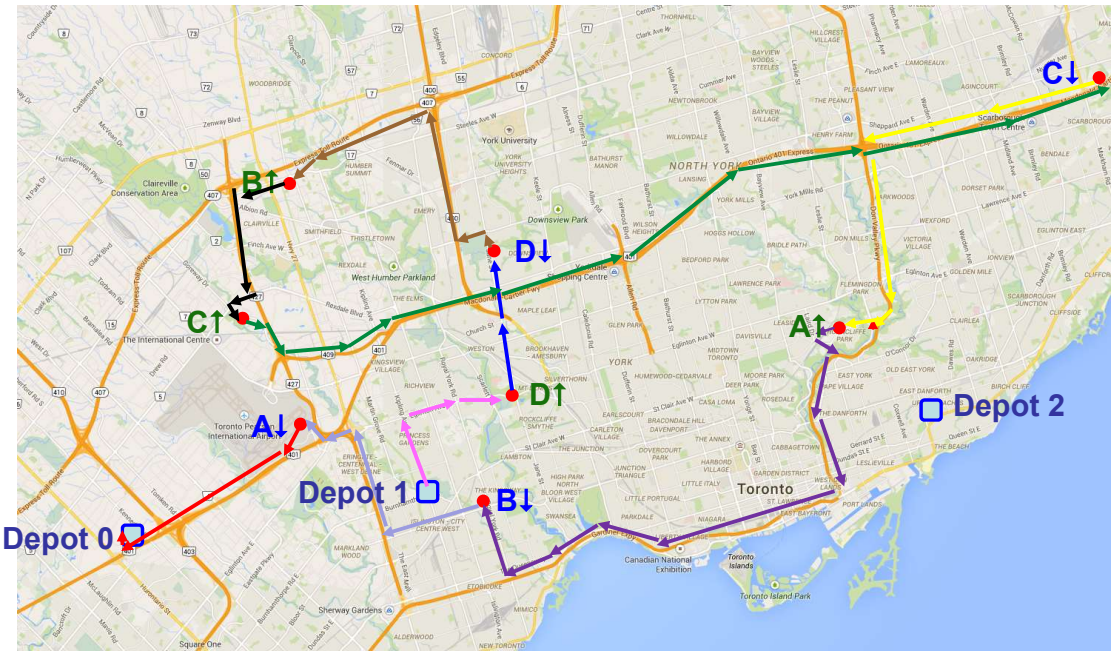

Figure 2: A possible solution generated by your courier delivery algorithm. This example solution contains 9 CourierSubpaths. It starts at depot 1, picks up D, drops off D, picks up B, picks up C, drops off C, picks up A, drops off B, drops off A, and finally ends at depot 0.

```
/*
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 *
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 * SOFTWARE.
 */
```

```
#pragma once
#include <vector>
#include "StreetsDatabaseAPI.h"


//Specifies a delivery order (input to your algorithm).
//To satisfy the order the item-to-be-delivered must have been picked-up
//from the pickUp intersection before visiting the dropOff intersection.
struct DeliveryInf {

    // Constructor
    DeliveryInf(IntersectionIdx pick_up, IntersectionIdx drop_off)
        : pickUp(pick_up), dropOff(drop_off) {}

    //The intersection id where the item-to-be-delivered is picked-up.
    IntersectionIdx pickUp;

    //The intersection id where the item-to-be-delivered is dropped-off.
    IntersectionIdx dropOff;
};


// Specifies one subpath of the courier truck route
struct CourierSubPath {

    // The intersection id where a start depot, pick-up intersection or
    // drop-off intersection is located
    IntersectionIdx start_intersection;

    // The intersection id where this subpath ends. This must be the
    // start_intersection of the next subpath or the intersection of an end depot
    IntersectionIdx end_intersection;

    // Street segment ids of the path between start_intersection and end_intersection
    // They form a connected path (see m3.h)
    std::vector<StreetSegmentIdx> subpath;
};


// This routine takes in a vector of D deliveries (pickUp, dropOff
// intersection pairs), another vector of N intersections that
// are legal start and end points for the path (depots), and a turn
// penalty in seconds (see m3.h for details on turn penalties).
//
// The first vector 'deliveries' gives the delivery information.  Each delivery
// in this vector has pickUp and dropOff intersection ids.
// A delivery can only be dropped-off after the associated item has been picked-up.
//
// The second vector 'depots' gives the intersection ids of courier company
// depots containing trucks; you start at any one of these depots and end at
// any one of the depots.
//
// This routine returns a vector of CourierSubPath objects that form a delivery route.
// The CourierSubPath is as defined above. The first street segment id in the
```

```
// first subpath is connected to a depot intersection, and the last street
// segment id of the last subpath also connects to a depot intersection.
// A package will not be dropped off if you haven't picked it up yet.
//
// The start_intersection of each subpath in the returned vector should be
// at least one of the following (a pick-up and/or drop-off can only happen at
// the start_intersection of a CourierSubPath object):
//      1- A start depot.
//      2- A pick-up location
//      3- A drop-off location.
//
// You can assume that D is always at least one and N is always at least one
// (i.e. both input vectors are non-empty).
//
// It is legal for the same intersection to appear multiple times in the pickUp
// or dropOff list (e.g. you might have two deliveries with a pickUp
// intersection id of #50). The same intersection can also appear as both a
// pickUp location and a dropOff location.
//
// If you have two pickUps to make at an intersection, traversing the
// intersection once is sufficient to pick up both packages. Additionally,
// one traversal of an intersection is sufficient to drop off all the
// (already picked up) packages that need to be dropped off at that intersection.
//
// Depots will never appear as pickUp or dropOff locations for deliveries.
//
// If no valid route to make *all* the deliveries exists, this routine must
// return an empty (size == 0) vector.
std::vector<CourierSubPath> travelingCourier(
                            const float turn_penalty,
                            const std::vector<DeliveryInf>& deliveries,
                            const std::vector<IntersectionIdx>& depots);

/**
 * @brief A simple pathfinding function with no turn penalty
 *
 * If you do not have a working findPathBetweenIntersections from
 * milestone 3, this file includes a Djikstra algorithm with no turn
 * penalties that you can use for milestone 4.
 *
 * The function prototype is:
 * std::vector<StreetSegmentIdx> findSimplePath(
 *           const std::pair<IntersectionIdx, IntersectionIdx> intersect_ids);
 *
 * The function is calling 2 of your own milestone 1 functions:
 *  - findStreetSegmentsOfIntersection
 *  - findStreetSegmentTravelTime
 * so those functions must be implemented and correct.
 *
 * @note If you have a working solution from milestone 3, it is better to use
 *       it instead of findSimplePath, as it will give smaller travel times
 *       due to handling turn penalties, and it may also have lower cpu time.
 */
#include "FindSimplePath.h"
```

Listing 1: m4.h

We will always call your `load_map` function before calling `travelingCourier`. Your algorithm must return a valid path within **50 seconds** of *wall clock* (real) time; if your code takes more than **50 seconds** we will consider it a failure for that test case. The UG machines on which we measure your program have 4 cores (CPUs) so you can use multi-threading to reduce your wall clock time if you wish.

You will always be given at least one delivery location and at least one depot, and you can assume that intersections in the depot vector will not appear as pick up or drop off locations in the deliveries vector. Note that it is also possible that no path to make all the deliveries exists; in this case you should return an empty (size == 0) vector.

If you didn't get your `findPathBetweenIntersections` algorithm fully working in milestone 3, you can use the findSimplePath function mentioned in Listing 1 as part of your algorithm. The object code (but not source code) for this function is in the libstreetmap library which is already in your mapper project.

4 Grading

Your entire mark in this milestone will depend on the correctness and performance of your code. The unit tests for this milestone are arranged in increasing difficulty (problem size), with the smallest problems being 5 or less delivery locations, and the largest having approximately 200 deliveries. For each size of problem, we will:

- Test that you can find and return a legal solution within the 50 second wall clock time limit.
- If you find a legal solution within the time limit, we will also evaluate its quality (travel time) and compare it to the quality of our reference solutions. You can see how your quality compares to our reference solutions and to the solutions of your classmates using the leaderboard web page described below.

Your grade will be determined from both of these components – some marks will be given for finding legal solutions, and most marks will depend on your solution quality. Some unit tests will be public, and others that are similar but use different intersections and/or different city maps will be run during final grading.

Note also that we have made some changes to the usual traveling salesman formulation, so while the ideas you find in the traveling salesman literature will be helpful, the exact code and algorithms required will not be the same.

5 Contest

In addition to the grades above, bonus marks are available for this lab as shown below:

- 1st place: 4% bonus on total course mark
- 2nd place: 3% bonus on total course mark
- 3rd place: 2% bonus on total course mark
- 4th place: 1% bonus on total course mark

The quality of each team's solution will be assessed by computing the geometric average of the travel time for a set of test inputs, all of which will be fairly large (on the order of 60 or more deliveries). To qualify, a team must compute a legal solution for each test within the 50 second wall clock time limit (i.e. any failed test means the team cannot compete). If two teams tie on the average travel time metric, the tie will be broken by lowest geometric average wall clock time required.

An anonymized team id will be posted on your team wiki page. Each time you submit milestone 4 with `ece297submit 4` your average travel time score and anonymized id will be posted to the leaderboard web site at http://ug251.eecg.utoronto.ca/ece297s/contest_2022/ . This leaderboard will use data from the public test cases run by exercise and submit. However, the final contest results will also include private test cases that are similar to the public ones but use different intersections and/or cities to prevent any hard-coding or overtuning.