


CSC3113

Theory Of Computation

Pre-Requisite

Basic Mathematical Concepts

-  Sets, Graphs, Relations, and Languages
-  Definitions, Theorems, and Proofs

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Assumed Background

- ✚ Sets / Sequences
- ✚ Functions / Relations
- ✚ Equivalence relations / Partitions
- ✚ Graphs
- ✚ Types of proof
 - ✚ Proof by construction
 - ✚ Proof by contradiction
 - ✚ Proof by induction
- ✚ Next we will go through the basic knowledge on the above topics.

Sets

- ✚ The symbols \in and \notin denote set membership and non membership, respectively.
example: $7 \in \{7, 21, 57\}$ and $8 \notin \{7, 21, 57\}$
- ✚ *Subset*: $A \subseteq B$, Every element of A is an element of B .
 - ✚ *Proper Subset*: If A is a subset of B and not equal to B .
- ✚ *Multiset*: $\{7\}$ and $\{7, 7\}$ are different as multisets but identical as sets.
- ✚ *Infinite set*: natural numbers $N = \{1, 2, 3, \dots\}$ and integers $Z = \{\dots, -2, -1, 0, 1, 2, \dots\}$, contains infinitely many elements.
- ✚ *Empty set*: Set with 0 members, written as \emptyset .
- ✚ $\{n \mid \text{rule about } n\}$: A set containing elements according to some rule.
 - ✚ $\{n \mid n = m^2 \text{ for some } m \in N\}$ means the set of perfect squares.
- ✚ *Cardinality* of a set: the number of elements in it.
- ✚ *Set Operations*:
 - ✚ *Compliment*: \bar{A} , is the set of all elements under consideration that are not in A .
 - ✚ *Union*: $A \cup B = \{x : x \in A \text{ or } x \in B\}$, the set we get by combining all the elements of in A and B . example: $\{7, 21\} \cup \{9, 5, 7\} = \{7, 21, 9, 5\}$.
 $\bigcup S = \{x : x \in P \text{ for some set } P \in S\}$ is the set whose elements are the elements of all the sets in S . example, $\bigcup S = \{a, b, c, d\}$ if $S = \{\{a, b\}, \{b, c\}, \{c, d\}\}$.

Sets

- **Intersection:** $A \cap B = \{x : x \in A \text{ and } x \in B\}$, the set of elements that are in both A and B . example: $\{7, 21\} \cap \{9, 5, 7\} = \{7\}$.
 $\bigcap S = \{x : x \in P \text{ for each set } P \in S\}$ is the set whose elements are the elements of all the sets in S . example, $\bigcap S = \{c, d\}$ if $S = \{\{a, c, d\}, \{c, d\}, \{b, c, d\}\}$.
- Two sets A and B are equal, written as $A = B$, if $A \subseteq B$ and $B \subseteq A$.
- **Difference** of two sets A and B , written $A - B$, is the set of all elements of A that are not elements of B . That is, $A - B = \{x : x \in A \text{ and } x \notin B\}$.
- Two sets are *disjoint* if they have no element in common. That is, $A \cap B = \phi$.
- **Power Set:** Power set of a set A is the set of all subsets of A .
if $A = \{0, 1\}$, then the power set of $A = \{\phi, \{0\}, \{1\}, \{0, 1\}\}$.
- A partition of a nonempty set A is a subset Π of 2^A such that,
 - Each element of Π is empty.
 - Distinct numbers of Π are disjoint.
 - $\bigcup \Pi = A$.
 - Example, $\{\{a, b\}, \{c\}, \{d\}\}$ is a partition of $\{a, b, c, d\}$.

Sequences

✚ *Sequence*: a list of object in some order.

▣ $(7, 21, 57)$ is a sequence of 7, 21, and 57.

✚ Order matters, so $(7, 21, 57)$ is not the same as $(21, 7, 57)$.

✚ Repetition is allowed, so $(7, 21, 57)$ is not the same as $(7, 21, 7, 57)$.

✚ *Tuple*: Finite sequence.

✚ *K-Tuple*: A sequence with k elements.

✚ *Pair*: A 2-tuple is called a *pair*.

✚ *Cartesian product/cross product* of A and B , written $A \times B$, is the set of all pairs wherein the first element is a member of A and the second element is a member of B .

If $A = \{1, 2\}$ and $B = \{x, y, z\}$,

$$A \times B = \{(1, x), (1, y), (1, z), (2, x), (2, y), (2, z)\}$$

$$A \times A = A^2 = \{(1, 1), (1, 2), (2, 1), (2, 2)\}$$

Functions

- ✚ A *function* maps an input to an output.
- ✚ Also called *mapping*, written as $f(a) = b$, meaning, if f is a function whose output value is b when the input value is a .
- ✚ *Domain*: the set of possible inputs.
- ✚ *Range*: the set of outputs.
- ✚ The notation for saying that f is a function with domain D and range R is $f : D \rightarrow R$.
- ✚ *k-ary function*: a function with k arguments (*arity* of a function).
 - ▣ Input: (a_1, a_2, \dots, a_k) , a k -tuple (*argument*).
 - ▣ unary function if $k = 1$
 - ▣ binary function if $k = 2$

Relations

- ✚ *Predicate (property)* : a function whose range is $\{\text{TRUE}, \text{FALSE}\}$.
- ✚ *Relation*: a property whose domain is a set of k tuples, A^k for a set A .
- ✚ Relation, k -ary relation or k -ary relation on A is written as $R(a_1, a_2, \dots, a_k)$.
- ✚ *Binary relation*: 2-ary relation. Customary infix notation aRb , where R is the relation between the elements a and b .
- ✚ *Inverse* of a binary relation $R \subseteq A \times B$, denoted $R^{-1} \subseteq B \times A$ is simply the relation $\{(b, a) : (a, b) \in R\}$.
- ✚ *Equivalence relation*: two objects being equal
 - ✚ reflexive: $\forall x, xRx$.
 - ✚ symmetric: $\forall xy, xRy$ iff yRx
 - ✚ transitive: $\forall xyz, xRy$ and $yRz \Rightarrow xRz$

Graphs

- A graph consists of a finite set of vertices (nodes) with lines connecting some of them (edges). $G = (V, E)$.
- *Undirected graph*:
 - *degree* of a node: the number of edges at a particular node.
 - *path*: a sequence of nodes connected by edges.
 - ◆ *simple path*: a path that doesn't repeat any nodes.
 - *cycle*: a path starts and ends in the same node
 - *tree*: no cycle
 - ◆ *leaves*: nodes of degree 1 in a tree.
 - ◆ *root*: special designated node.
- *Directed graph*:
 - *in-degree* and *out-degree*
 - *directed path*
 - *directed acyclic graph* (DAG)
- *Sub Graph*: Graph G is a subgraph of graph H , if the nodes of G are a subset of the nodes of H (i.e. $G.V \subseteq H.V$).
- *connected*: every two nodes of a graph have a path between them.
 - *strongly connected*: every 2 nodes of a di-graph have a path between them.

Strings

- ✚ Strings of characters.
- ✚ *Alphabet*: any finite set, Σ and Γ designate alphabets and a typewriter font for symbols from an alphabet. Example: $\Sigma_1 = \{0,1\}$, $\Sigma_2 = \{a, x, y, z\}$, $\Gamma = \{0,1, x, z\}$.
- ✚ A *string over an alphabet*: a finite sequence of symbols from the alphabet. If $\Sigma_1 = \{0,1\}$, then 01001 is a string over Σ_1 .
- ✚ *Length* of a string w : $|w|$.
- ✚ *Empty* string: ε .
- ✚ String z is a *substring* of string w if z appears consecutively within w . Example: $z=\text{cad}$, $w=\text{abracadabra}$.
- ✚ If $w=xv$ for some x , then v is a *suffix* of w ; If $w=vy$ for some y , then v is a *prefix* of w .
- ✚ If w has length n , we can write $w = w_1w_2 \dots w_n$ where each $w_i \in \Sigma$. *Reverse* of w , written w^R , is the string obtained by writing w in the opposite order (i.e. $w_nw_{n-1} \dots w_1$).
- ✚ *Concatenation* of two strings x and y , written xy , is the string obtained by appending y to the end of x , as in $x_1 \dots x_n y_1 \dots y_n$. To concatenate a string with itself many times we use the superscript notation $\overbrace{xx \dots x}^k = x^k$.

Languages

- # A *language* is a set of strings.
- # The set of all strings of all lengths, including the empty string, over an alphabet Σ is denoted by Σ^* .
- # *Lexicographic ordering* of strings is the same as the familiar dictionary ordering, expect that shorter strings precede longer strings. Example: Lexicographic ordering of all strings over the alphabet $\Sigma = \{0,1\}$ is $(\varepsilon, 0, 1, 00, 01, 10, 11, 000, \dots)$.
- # A language L over the alphabet A is a subset of A^* . $L \subseteq A^*$.

Proofs

✚ Proof: a convincing logical argument that a statement is true.

✚ convincing in an absolute sense

✚ Methods of proof

✚ *The pigeonhole principle*: there are n pigeonholes, $n + 1$ or more pigeons, and every pigeon occupies a hole, then some hole must have at least two pigeons.

✚ *Proof by construction*: Prove a particular type of objects exists by constructing the object.

✚ *Proof by contradiction*: Assume a theorem is false and then show that this assumption leads to a false consequence.

✚ *Proof by induction*: A proof by induction has –

◆ A predicate: P ,

◆ A basis: $\exists k, P(k)$ is true,

◆ An induction hypothesis: for some $n \geq k, P(k), P(k + 1), \dots, P(n)$ are true.

◆ An inductive step: $P(n + 1)$ is true given the induction hypothesis.