Assignment 1

- 1. (5) Consider the matrix $H(v) = \hat{1} 2vv^T$, where v is a unit column vector. What is the rank of the matrix H(v)? Prove that it is orthogonal.
- 2. (10) Prove the following inequalities and provide examples of x and A when they turn into equalities:
 - $||x||_2 \leq \sqrt{m} ||x||_{\infty}$
 - $||A||_{\infty} \leq \sqrt{n} \, ||A||_{2}$

where x is a vector of m components and A is $m \times n$ matrix.

- 3. (5) Assuming u and v are m-vectors, consider the matrix $A=1+uv^T$ which is a rank-one perturbation of identity. Can it be singular? Assuming it is not, compute its inverse. You may look for it in a form of $A^{-1}=1+\alpha uv^T$ for some scalar α and evaluate α .
- 4. (5) Prove that for any unitary matrix U one has $||UA||_F = ||AU||_F = ||A||_F$.
- 5. (10) In this exercise your goal will be to study and speed up an implementation of K-means algorithm. In the notebook kmeans.ipynb, you can find a naive implementation. Explore the code, make sure you understand it. You will find there two functions dist_i and dist_ij which are (on purpose) implemented in a rather inefficient way. Improve them by getting rid of the loops in the favor of a proper numpy vectorized implementation and measure the speed-up of the full algorithm for N = 10000.
- 6. (10) Some things just can not be vectorized but still can be sped up compared to naive implementation. For example, consider computation of the Hofstadter-Conway sequence a(n) such that a(1) = 1, a(2) = 1 and

$$a(n) = a(a(n-1)) + a(n-a(n-1)), \quad n > 2$$
(1)

Write three functions, computing the sequence up to n-th element in three ways: i) pre-allocating numpy array and filling it using for loop, ii) cumulatively appending python list and converting it to numpy array, iii) same as i) but compiled (jit) version. Time the resulting implementations and conclude which is preferable. With the optimal one, compute $a(10^8)$.

7. (15*) Consider a function mapping six tensors to one tensor: $Z(\lambda^{(1)}, \lambda^{(2)}, \lambda^{(3)}, \Gamma^{(1)}, \Gamma^{(2)}, U)$, with

$$Z_{ahij} = \sum_{bcdefg} \lambda^{(1)}{}_{ab} \Gamma^{(1)}{}_{cbd} \lambda^{(2)}{}_{de} \Gamma^{(2)}{}_{feg} \lambda^{(3)}{}_{gh} U_{ijcf}. \tag{2}$$

Assume that all indices of the tensors appearing above take values from 1 to χ . Running the numerical experiments, explore the values of χ in the range 3–50 (from slowest to fastest implementation).

- In the notebook convolution.ipynb you may find implemented a *stupid* way to compute this convolution, which takes $\chi^4 \times \chi^6 = \chi^{10}$ flops. In fact, this can be computed much faster!
- Using the function numpy.einsum (its crucial to use the optimize argument), you can actually achieve a much faster implementation. In order to understand what it is doing under the hood, explore the function numpy.einsum_path. What is the minimal number of flops required for computation of \mathbb{Z} ?
- Using the understanding of the output of $numpy.einsum_path$, implement an algorithm to compute Z, which is as effective as numpy.einsum, but relying only on more elementary numpy.dot and $numpy.tensor_dot$.