

op(31:26):								
28–26	0(000)	1(001)	2(010)	3(011)	4(100)	5(101)	6(110)	7(111)
31–29								
0(000)	<u>Rfmt</u>	<u>Bltz/gez</u>	j	jal	beq	bne	blez	bgtz
1(001)	addi	addiu	slti	sltiu	ANDi	ORi	xORi	lui
2(010)	<u>ILB</u>	<u>FlPt</u>						
3(011)								
4(100)	lb	lh	lwl	lw	lbu	lhu	lwr	
5(101)	sb	sh	swl	sw			swr	
6(110)	lwc0	lwc1						
7(111)	swc0	swc1						

op(31:26) = 010001 (FlPt), (rt(16:16) = 0 => c = f, rt(16:16) = 1 => c = t), rs(25:21):								
23–21	0(000)	1(001)	2(010)	3(011)	4(100)	5(101)	6(110)	7(111)
25–24								
0(00)	mfc1		cfcl		mtcl		ctcl	
1(01)	bc1. <i>c</i>							
2(10)	<i>f</i> = single	<i>f</i> = double						
3(11)								

op(31:26) = 010001 (FlPt), (f above: 10000 => f = s, 10001 => f = d), funct(5:0):								
2–0	0(000)	1(001)	2(010)	3(011)	4(100)	5(101)	6(110)	7(111)
5–3								
0(000)	add. <i>f</i>	sub. <i>f</i>	mul. <i>f</i>	div. <i>f</i>		abs. <i>f</i>	mov. <i>f</i>	neg. <i>f</i>
1(001)								
2(010)								
3(011)								
4(100)	cvt.s. <i>f</i>	cvt.d. <i>f</i>			cvt.w. <i>f</i>			
5(101)								
6(110)	c.f. <i>f</i>	c.un. <i>f</i>	c.eq. <i>f</i>	c.ueq. <i>f</i>	c.olt. <i>f</i>	c.ult. <i>f</i>	c.ole. <i>f</i>	c.ule. <i>f</i>
7(111)	c.sf. <i>f</i>	c.ngle. <i>f</i>	c.seq. <i>f</i>	c.ngl. <i>f</i>	c.lt. <i>f</i>	c.nge. <i>f</i>	c.le. <i>f</i>	c.ngt. <i>f</i>

FIGURE 3.18 MIPS floating-point instruction encoding. This notation gives the value of a field by row and by column. For example, in the top portion of the figure, lw is found in row number 4 (100_{two} for bits 31–29 of the instruction) and column number 3 (011_{two} for bits 28–26 of the instruction), so the corresponding value of the op field (bits 31–26) is 100011_{two}. Underlined text means the field is used elsewhere. For example, FlPt in row 2 and column 1 (op = 010001_{two}) is defined in the bottom part of the figure. Hence sub.f in row 0 and column 1 of the bottom section means that the funct field (bits 5–0) of the instruction is 000001_{two} and the op field (bits 31–26) is 010001_{two}. Note that the 5-bit rs field, specified in the middle portion of the figure, determines whether the operation is single precision (*f* = s, so rs = 10000) or double precision (*f* = d, so rs = 10001). Similarly, bit 16 of the instruction determines if the bc1.c instruction tests for true (bit 16 = 1 => bc1.t) or false (bit 16 = 0 => bc1.f). This information is also found in column 2 of the MIPS Reference Data Card at the front of this book.