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module CPU (clock);
parameter LW = 6'b100011, SW = 6'b101011, BEQ = 6'b000100, no-op = 32'b00000000_00000000_00000000_00000000, ALUop = 6'b0;
input clock;
    reg[31:0] PC, Regs[0:31], IMemory[0:1023], DMemory[0:1023], // separate memories
        IFIDIR, IDEXA, IDEXB, IDEXIR, EXMEMIR, EXMEMB, // pipeline registers
        EXMEMALUOut, MEMWBValue, MEMWBIR; // pipeline registers
wire [4:0] IDEXrs, IDEXrt, EXMEMrd, MEMWBrd; //hold register fields
wire [5:0] EXMEMop, MEMWBop, IDEXop; Hold opcodes
wire [31:0] Ain, Bin;
// declare the bypass signals
wire takebranch, stall, bypassAfromMEM, bypassAfromALUinWB, bypassBfromMEM, bypassBfromALUinWB,
    bypassAfromLWinWB, bypassBfromLWinWB;
assign IDEXrs = IDEXIR[25:21]; assign IDEXrt = IDEXIR[15:11]; assign EXMEMrd = EXMEMIR[15:11];
assign MEMWBrd = MEMWBIR[20:16]; assign EXMEMop = EXMEMIR[31:26];
assign MEMWBop = MEMWBIR[31:26]; assign IDEXop = IDEXIR[31:26];
// The bypass to input A from the MEM stage for an ALU operation
assign bypassAfromMEM = (IDEXrs == EXMEMrd) & (IDEXrs!=0) & (EXMEMop==ALUop); // yes, bypass
// The bypass to input B from the MEM stage for an ALU operation
assign bypassBfromMEM = (IDEXrt == EXMEMrd)&(IDEXrt!=0) & (EXMEMop==ALUop); // yes, bypass
// The bypass to input A from the WB stage for an ALU operation
assign bypassAfromALUinWB = (IDEXrs == MEMWBrd) & (IDEXrs!=0) & (MEMWBop==ALUop);
// The bypass to input B from the WB stage for an ALU operation
assign bypassBfromALUinWB = (IDEXrt == MEMWBrd) & (IDEXrt!=0) & (MEMWBop==ALUop); /
// The bypass to input A from the WB stage for an LW operation
assign bypassAfromLWinWB = (IDEXrs == MEMWBIR[20:16]) & (IDEXrs!=0) & (MEMWBop==LW);
// The bypass to input B from the WB stage for an LW operation
assign bypassBfromLWinWB = (IDEXrt == MEMWBIR[20:16]) & (IDEXrt!=0) & (MEMWBop==LW);
// The A input to the ALU is bypassed from MEM if there is a bypass there,
// Otherwise from WB if there is a bypass there, and otherwise comes from the IDEX register
assign Ain = bypassAfromMEM? EXMEMALUOut :
    (bypassAfromALUinWB | bypassAfromLWinWB)? MEMWBValue : IDEXA;
// The B input to the ALU is bypassed from MEM if there is a bypass there,
// Otherwise from WB if there is a bypass there, and otherwise comes from the IDEX register
assign Bin = bypassBfromMEM? EXMEMALUOut :
    (bypassBfromALUinWB | bypassBfromLWinWB)? MEMWBValue: IDEXB;
// The signal for detecting a stall based on the use of a result from LW
assign stall = (MEMWBIR[31:26]==LW) && // source instruction is a load
    (((IDEXop==LW)|(IDEXop==SW)) && (IDEXrs==MEMWBrd)) | // stall for address calc
    ((IDEXop==ALUop) && ((IDEXrs==MEMWBrd)|(IDEXrt==MEMWBrd))); // ALU use

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**FIGURE e4.14.4** A behavioral definition of the five-stage MIPS pipeline with stalls for loads when the destination is an ALU instruction or effective address calculation. The changes from Figure e4.14.3 are highlighted.