

Algebraic Structures (08)

- 7.1 Algebraic structures with one binary operation: semigroup, monoids and groups**
- 7.2 Cyclic groups, Normal subgroups**
- 7.3 Hamming Code ,Minimum Distance**
- 7.4 Group codes ,encoding-decoding techniques**
- 7.5 Parity check Matrix ,Maximum Likelihood**

Algebraic systems

- $N = \{1, 2, 3, 4, \dots, \infty\}$ = Set of all natural numbers.
- $Z = \{0, \pm 1, \pm 2, \pm 3, \pm 4, \dots, \infty\}$ = Set of all integers.
- Q = Set of all rational numbers, R = Set of all real numbers.
- **Binary Operation:** The binary operator * is said to be a binary operation (closed operation) on a non empty set A, if
 $a * b \in A$ for all $a, b \in A$ (Closure property).

Ex: The set N is closed with respect to addition and multiplication

but not w.r.t subtraction and division.

- **Algebraic System:** A set 'A' with one or more binary(closed) operations defined on it is called an algebraic system.

Ex: $(N, +)$, $(Z, +, -)$, $(R, +, ., -)$ are algebraic systems.

Properties

- **Commutative:** Let $*$ be a binary operation on a set A.
The operation $*$ is said to be commutative in A if
 $a * b = b * a$ for all a, b in A
- **Associativity:** Let $*$ be a binary operation on a set A.
The operation $*$ is said to be associative in A if
 $(a * b) * c = a * (b * c)$ for all a, b, c in A

(Addition , Subtraction)

- **Idempotent :** Let $*$ be a binary operation on a set A.
The operation $*$ is said to be idempotent in A if
 $a * a = a$
- **Identity:** For an algebraic system $(A, *)$, an element 'e' in A is said to be an identity element of A if
 $a * e = e * a = a$ for all $a \in A$.
- **Inverse:** Let $(A, *)$ be an algebraic system with identity 'e'. Let a be an element in A. An element b is said to be inverse of A if
 $a * b = b * a = e$

Semi group

- **Semi Group:** An algebraic system $(A, *)$ is said to be a semi group if
 1. *** is closed operation on A.**
 2. *** is an associative operation, for all a, b, c in A.**
- Ex. $(N, +)$ is a semi group.
- Ex. (N, \cdot) is a semi group.
- Ex. $(N, -)$ is not a semi group.
- **Monoid:** An algebraic system $(A, *)$ is said to be a **monoid** if the following conditions are satisfied.
 - 1) *** is a closed operation in A.**
 - 2) *** is an associative operation in A.**
 - 3) **There is an identity in A.**

Monoid

- Ex. Show that the set ‘N’ is a monoid with respect to multiplication.
- Solution: Here, $N = \{1, 2, 3, 4, \dots\}$
 1. Closure property : We know that product of two natural numbers is again a natural number.
i.e., $a.b = b.a$ for all $a, b \in N$
 \therefore Multiplication is a closed operation.
 2. Associativity : Multiplication of natural numbers is associative.
i.e., $(a.b).c = a.(b.c)$ for all $a, b, c \in N$
 3. Identity : We have, $1 \in N$ such that
 $a.1 = 1.a = a$ for all $a \in N$.
 \therefore Identity element exists, and 1 is the identity element.

Hence, N is a monoid with respect to multiplication.

Subsemigroup & submonoid- “**Self read**”

Subsemigroup : Let $(S, *)$ be a semigroup and let **T be a subset of S**.

If T is closed under operation * , then $(T, *)$ is called a subsemigroup of $(S, *)$.

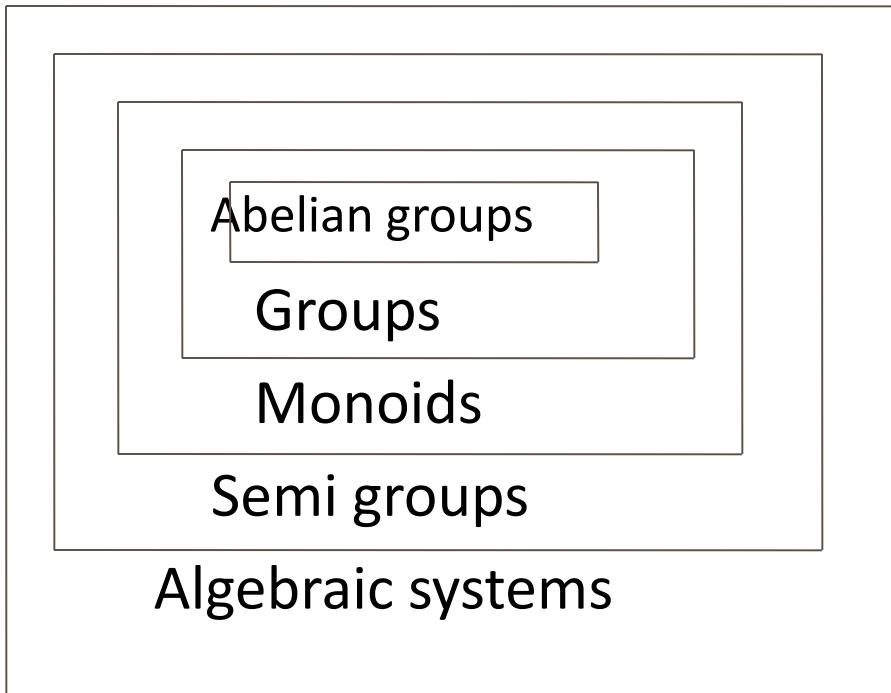
Ex: $(N, .)$ is semigroup and T is set of multiples of positive integer m then $(T,.)$ is a sub semigroup.

Submonoid : Let $(S, *)$ be a monoid with identity e, and let T be a non-empty subset of S. If T is closed under the operation * and $e \in T$, then $(T, *)$ is called a submonoid of $(S, *)$.

Group

- **Group:** An algebraic system $(G, *)$ is said to be a **group** if the following conditions are satisfied.
 - 1) * is a closed operation.
 - 2) * is an associative operation.
 - 3) There is an identity in G.
 - 4) Every element in G has inverse in G.
- **Abelian group (Commutative group):** A group $(G, *)$ is said to be **abelian** (or **commutative**) if
$$a * b = b * a \quad \text{for all } a, b \text{ belongs to } G.$$

Algebraic systems



Theorems –”**Self Study**”

- In a Group $(G, *)$ the following properties hold good

1. Identity element is unique.
2. Inverse of an element is unique.
3. Cancellation laws hold good

$$a * b = a * c \Rightarrow b = c \quad (\text{left cancellation law})$$

$$a * c = b * c \Rightarrow a = b \quad (\text{Right cancellation law})$$

$$4. \quad (a * b)^{-1} = b^{-1} * a^{-1}$$

- In a group, the identity element is its own inverse.

- **Order of a group** : The number of elements in a group is called order of the group.

- **Finite group**: If the order of a group G is finite, then G is called a finite group.

Ex. Show that, the set of all integers is a group with respect to **addition**.

- Solution: Let Z = set of all integers.

Let a, b, c are any three elements of Z .

1. **Closure property**: We know that, Sum of two integers is again an integer.

i.e., $a + b \in Z$ for all $a, b \in Z$

2. **Associativity**: We know that addition of integers is associative.

i.e., $(a+b)+c = a+(b+c)$ for all $a, b, c \in Z$.

3. **Identity**: We have $0 \in Z$ and $a + 0 = a$ for all $a \in Z$.

\therefore Identity element exists, and '0' is the identity element.

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Contd.,

4. **Inverse**: To each $a \in Z$, we have $-a \in Z$ such that

$$a + (-a) = 0$$

Each element in Z has an inverse

■ 5. **Commutativity**: We know that addition of integers is commutative.

i.e., $a + b = b + a$ for all $a, b \in Z$.

Hence, $(Z, +)$ is an abelian group.

Ex. Show that set of all non zero real numbers is a group with respect to multiplication .("Self Study")

- Solution: Let R^* = set of all non zero real numbers.

Let a, b, c are any three elements of R^* .

1. Closure property : We know that, product of two nonzero real numbers is again a nonzero real number .

i.e., $a \cdot b \in R^*$ for all $a,b \in R^*$.

2. Associativity: We know that multiplication of real numbers is associative.

i.e., $(a.b).c = a.(b.c)$ for all $a,b,c \in R^*$.

3. Identity: We have $1 \in R^*$ and $a \cdot 1 = a$ for all $a \in R^*$.

\therefore Identity element exists, and '1' is the identity element.

4. Inverse: To each $a \in R^*$, we have $1/a \in R^*$ such that
 $a \cdot (1/a) = 1$ i.e., Each element in R^* has an inverse.

Contd.,

- **5. Commutativity:** We know that multiplication of real numbers is commutative.
i.e., $a \cdot b = b \cdot a$ for all $a, b \in R^*$.
Hence, (R^*, \cdot) is an abelian group.

- **Ex: Show that set of all real numbers 'R' is not a group with respect to multiplication.**
■ Solution: We have $0 \in R$.
The multiplicative inverse of 0 does not exist.
Hence, R is not a group.

MODULO SYSTEMS

Addition modulo m (+_m)

let m be a positive integer. For any two positive integers a and b

$$a +_m b = a + b \quad \text{if } a + b < m$$

$a +_m b = r \quad \text{if } a + b \geq m$ where r is the remainder obtained
by dividing $(a+b)$ with m .

Ex $14 +_6 8 = 22 \% 6 = 4$; **Ex** $9 +_{12} 3 = 12 \% 12 = 0$

Multiplication modulo p (×_p)

let p be a positive integer. For any two positive integers a and b

$$a \times_p b = ab \quad \text{if } ab < p$$

$a \times_p b = r \quad \text{if } ab \geq p$ where r is the remainder obtained
by dividing (ab) with p .

Ex. $3 \times_5 4 = 2$, $5 \times_5 4 = 0$, $2 \times_5 2 = 4$

Ex.The set $G = \{0,1,2,3,4,5\}$ is a group with respect to addition modulo 6.

Solution: The composition table of G is

$+_6$	0	1	2	3	4	5
0	0	1	2	3	4	5
1	1	2	3	4	5	0
2	2	3	4	5	0	1
3	3	4	5	0	1	2
4	4	5	0	1	2	3
5	5	0	1	2	3	4

1. Closure property: Since all the entries of the composition table are the elements of the given set, the set G is closed under $+_6$.

Contd.,

2. Associativity: The binary operation $+_6$ is associative in G.

for ex. $(2 +_6 3) +_6 4 = 5 +_6 4 = 3$ and

$$2 +_6 (3 +_6 4) = 2 +_6 1 = 3$$

3. Identity : Here, The first row of the table coincides with the top row.

The element heading that row , i.e., 0 is the identity element.

4. . Inverse: From the composition table, we see that the inverse

elements of 0, 1, 2, 3, 4, 5 are 0, 5, 4, 3, 2, 1 respectively.

5. Commutativity: The corresponding rows and columns of the table are

identical. Therefore the binary operation $+_6$ is commutative.

Hence, $(G, +_6)$ is an abelian group.

Ex.The set $G = \{1,2,3,4,5,6\}$ is a group with respect to multiplication modulo 7.

Solution: The composition table of G is

\times_7	1	2	3	4	5	6
1	1	2	3	4	5	6
2	2	4	6	1	3	5
3	3	6	2	5	1	4
4	4	1	5	2	6	3
5	5	3	1	6	4	2
6	6	5	4	3	2	1

1. Closure property: Since all the entries of the composition table are the elements of the given set, the set G is closed under \times_7 .

Contd.,

2. Associativity: The binary operation \times_7 is associative in G.

for ex. $(2 \times_7 3) \times_7 4 = 6 \times_7 4 = 3$ and

$$2 \times_7 (3 \times_7 4) = 2 \times_7 5 = 3$$

3. Identity: Here, The first row of the table coincides with the top row.

The element heading that row , i.e., 1 is the identity element.

4. . Inverse: From the composition table, we see that the inverse elements of 1, 2, 3, 4, 5, 6 are 1, 4, 5, 2, 3, 6 respectively.

5. Commutativity: The corresponding rows and columns of the table are identical. Therefore the binary operation \times_7 is commutative.

Hence, (G, \times_7) is an abelian group.

Normal Subgroup

A subgroup is called a **normal subgroup** if for any $a \in G$,
 $aH = Ha$.

Note 1:

$aH = Ha$ does not necessarily mean that $a * h = h * a$ for every $h \in H$.
It only means that $a * h_i = h_j * a$ for some $h_i, h_j \in H$.

Note2:

Every subgroup of an abelian group is normal.
 $Hg = gH$, for all $g \in G$, if and only if H is a normal subgroup of G .

Let $H = \{[0]_6, [3]_6\}$, Find left and right cosets in group Z_6
Is it a normal subgroup

- It is abelian group, $a +_6 b = b +_6 a$

+	0	1	2	3	4	5
0	0	1	2	3	4	5
1	1	2	3	4	5	0
2	2	3	4	5	0	1
3	3	4	5	0	1	2
4	4	5	0	1	2	3
5	5	0	1	2	3	4

Left coset(Right is fixed) of H , $aH = \{ a * h \mid h \in H \}$

$$0H = \{ 0 +_6 0, 0 +_6 3 \} = \{ 0, 3 \}$$

$$1H = \{ 1 +_6 0, 1 +_6 3 \} = \{ 1, 4 \}$$

$$2H = \{ 2 +_6 0, 2 +_6 3 \} = \{ 2, 5 \}$$

$$3H = \{ 3 +_6 0, 3 +_6 3 \} = \{ 3, 0 \}$$

$$4H = \{ 4 +_6 0, 4 +_6 3 \} = \{ 4, 1 \}$$

$$5H = \{ 5 +_6 0, 5 +_6 3 \} = \{ 5, 2 \}$$

Given

$$H = \{ [0]_6, [3]_6 \} \quad \text{w.k.t} \rightarrow a +_6 b = b +_6 a$$

Right coset(Left is fixed) of H , $H a = \{ h * a \mid h \in H \}$

$$H 0 = \{ 0 +_6 0, 3 +_6 0 \} = \{ 0, 3 \}$$

$$H 1 = \{ 0 +_6 1, 3 +_6 1 \} = \{ \quad \}$$

$$H 2 = \{ 0 +_6 2, 3 +_6 2 \} = \{ \quad \}$$

$$H 3 = \{ 0 +_6 3, 3 +_6 3 \} = \{ \quad \}$$

$$H 4 = \{ 0 +_6 4, 3 +_6 4 \} = \{ \quad \}$$

$$H 5 = \{ 0 +_6 5, 3 +_6 5 \} = \{ \quad \}$$

$H_0, H_1, H_2, H_3, H_4, H_5 = 0H, 1H, 2H, 3H, 4H, 5H$

Hamming distance

The Hamming distance $d(x, y)$ between two words x, y is the weight $|x \oplus y|$ of $x \oplus y$, (bits in which they differ)

Eg. $d(00111, 11001) = 4$

Find the distance between x and y

$x = 110110 ; y = 000101$

$x = 001100 ; y = 010110$

$x = 0100100 ; y = 0011010$

Theorems

- The minimum weight of all non zero words in a group code is equal to its minimum distance
- A code can **detect** all combinations of k or fewer iff the minimum distance between any two code words is at least $k + 1$
- A code can **correct** all combinations of k or fewer errors iff the minimum distance between any two code words is at least $2k + 1$

- Consider the (2,4) encoding function , how many errors will 'e' detect ($k+1$)?

$e(00)=0000$

$e(10)=0111$

$e(01)=1011$

$e(11)=1100$

Soln: Find the Hamming distance between all pairs

Since $2 \geq k+1$

$k \leq 1$, Will detect 1 or fewer errors

- Consider the encoding function $B^2 \rightarrow B^6$ defined as follows

$e(00) = 0010000$

$e(10) = 100010$

$e(01) = 010100$

$e(11) = 110001$

How many errors can it correct and detect?

Error detection $3 >= k+1$; $k \leq 2$ or fewer errors

Error correction $3 >= 2k+1$; $k \leq 1$ or fewer errors

Group Codes- “Self Read”

An (m,n) encoding function $e:B^m \rightarrow B^n$ is called

a group code if $e(B^m) = \{e(b) | b \in B^m\} = \text{Ran}$

(e) is a subgroup of B^n

Subgroup if:

Identity element of B^n is in N

If x and y belong to N , then $x \oplus y \in N$

If x is in N , then its inverse in N

Consider the encoding function $B^2 \rightarrow B^5$ defined as follows

$$e(00) = 00000$$

$$e(10) = 10101$$

$$e(01) = 01110$$

$$e(11) = 11011$$

is a group code

Soln: Let $N = \{ 00000, 10101, 01110, 11011 \}$ be set of code words

\oplus	00000	01110	10101	11011
00000	00000	01110	10101	11011
01110	01110	00000	11011	10101
10101	10101	11011	00000	01110
11011	11011	10101	01110	00000

$a \oplus b \in N$ which is closed operation, associative, identity, inverse

1. Closed operation : For any $a, b \in N$, $a \oplus b \in N$, So N is closed under \oplus operation

2. Identity element of B^5 i.e 00000 also belongs to N

$$\text{00000} \oplus 00000 = 00000 \oplus 00000$$

$$\text{01110} \oplus 00000 = 00000 \oplus 01110$$

$$\text{10101} \oplus 00000 = 00000 \oplus 10101$$

$$\text{11011} \oplus 00000 = 00000 \oplus 11011$$

3. \oplus Associative operation

$$01110 \oplus (00000 \oplus 10101) = (01110 \oplus 00000) \oplus 10101$$

$$01110 \oplus 10101 = 01110 \oplus 10101$$

$$11011 = 11011$$

4 . Inverse $a * b = b * a = e$

$$\text{Ex: } 01110 \oplus 01110 = 01110 \oplus 01110 = 00000$$

Show that (3,5)encoding function $e: B^3 \rightarrow B^6$
defined as follows

$$e(000)=000000$$

$$e(001)=000110$$

$$e(100)=100101$$

$$e(110)=110111$$

$$e(010)=010010$$

$$e(011)=010100$$

$$e(101)=100011$$

$$e(111)=110001$$

PARITY CHECK MATRIX

Consider the parity check matrix given by H;

$$H = \begin{pmatrix} 1 & 1 & 0 \\ 0 & 1 & 1 \\ 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{pmatrix}$$

Determine the group code $e_H : B^2 \rightarrow B^5$

Soln: $B^2 = \{00, 01, 10, 11\}$

Then $e(00) = 00 x_1 x_2 x_3 = B^5$

$$x_1 = 0 \cdot 1 + 0 \cdot 0 = 0$$

$$x_2 = 0 \cdot 1 + 0 \cdot 1 = 0$$

$$x_3 = 0 \cdot 0 + 0 \cdot 1 = 0$$

$$e(00) = 00000$$

Next $e(01) = 01 x_1 x_2 x_3 = B^5$

$$x_1 = 0 \cdot 1 + 1 \cdot 0 = 0$$

$$x_2 = 0 \cdot 1 + 1 \cdot 1 = 1$$

$$x_3 = 0 \cdot 0 + 1 \cdot 1 = 1$$

$$e(01) = 01011$$

$$H = \begin{pmatrix} 1 & 1 & 0 \\ 0 & 1 & 1 \\ 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{pmatrix}$$

Next $e(10) = 10 x_1 x_2 x_3 = B^5$

$$x_1 = 1 \cdot 1 + 0 \cdot 0 = 1$$

$$x_2 = 1 \cdot 1 + 0 \cdot 1 = 1$$

$$X_3 = 1 \cdot 0 + 0 \cdot 1 = 0$$

$e(10) = 10110$

Next $e(11) = 11 x_1 x_2 x_3 = B^5$

$$x_1 = 1 \cdot 1 + 1 \cdot 0 = 1$$

$$x_2 = 1 \cdot 1 + 1 \cdot 1 = 0$$

$$X_3 = 1 \cdot 0 + 1 \cdot 1 = 1$$

$e(11) = 11101$

$e_H : B^2 \rightarrow B^5$ is as above for $e(00)$, $e(01)$, $e(10)$, $e(11)$

$e(00) = 00000$, $e(01) = 01011$, $e(10) = 10110$, $e(11) = 11101$

Problem 1

Consider the parity check matrix given by H ;

$$H = \begin{pmatrix} 0 & 1 & 1 \\ 0 & 1 & 1 \\ 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{pmatrix}$$

Determine the group code $e_H : B^2 \rightarrow B^5$

$e(00) = 00000$

$e(01) = 01011$

$e(00) = 10011$

$e(00) = 11000$

Problem 2

Consider the parity check matrix given by H ;

$$H = \begin{pmatrix} 1 & 0 & 0 \\ 0 & 1 & 1 \\ 1 & 1 & 1 \\ 1 & 0 & 0 \\ 0 & 1 & 0 \\ 0 & 0 & 1 \end{pmatrix}$$

Determine the group code $e_H : B^3 \rightarrow B^6$

$e(000) = 000000$

$e(001) = 001111$

$e(010) = 010011$

$e(011) = 011100$

$e(100) = 100100$

$e(101) = 101011$

$e(110) = 110111$

$e(111) = 111000$

MAXIMUM LIKELIHOOD DECODING TECHNIQUE

Consider the encoding function $B^2 \rightarrow B^4$ defined as follows

$$e(00) = 0000$$

$$e(10) = 1011$$

$$e(01) = 0110$$

$$e(11) = 1101$$

Decode the foll words relative to MLD function,

- (i) 0101 e(01) (ii) 1010 e(10) (iii) 1101 e(11)

Step 1: Construct Decoding Table(Taking various combinations of 4 bit numbers such as 0000[1],then making MSB as 1 thus we get 0001[2] then shift($R \rightarrow L$) to next bit as 1 we get 0010[3] ,etc. BUT before every next combination we need to check if its in the table already like 0100 [4] . As 0100 is already considered we move to 1000[5]accordingly we proceed for unique values till we decode all words in the question.

	0000	0110	1011	1101
0000[1]	0000	0110	1011	1101
0001[2]	0001	0111	1010	1100
0010[3]	0010	0100 [4]	1001	1111
1000[5]	1000	1110	0011	0101

Consider the encoding function $B^2 \rightarrow B^5$
defined as follows

$$e(00) = \textcolor{teal}{00000}$$

$$e(10) = \textcolor{teal}{10101}$$

$$e(01) = \textcolor{teal}{01110}$$

$$e(11) = \textcolor{teal}{11011}$$

Decode the foll words relative to MLD function,

- (i) 11110 (ii) 10011 (iii) 10100

	e (00)	e (01)	e (10)	e (11)
	00000	01110	10101	11011
00000	00000	01110	10101	11011
0000<u>1</u>	00001	01111	<u>10100</u>	11010
000<u>1</u>0	00010	01100	10111	11001
00<u>1</u>00	00100	01010	10001	11111
<u>0</u>1000	01000	00110	11101	<u>10011</u>
<u>1</u>0000	10000	<u>11110</u>	00101	01011