

# Basic Principles in Networking

## Digital signatures and End Point Authentication

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Version 1.0



Aalto University  
School of Electrical  
Engineering

# Video: INTERVIEW – Elliptic curve cryptography (10 min)

How did NSA read our emails?

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# Part I (15 min)

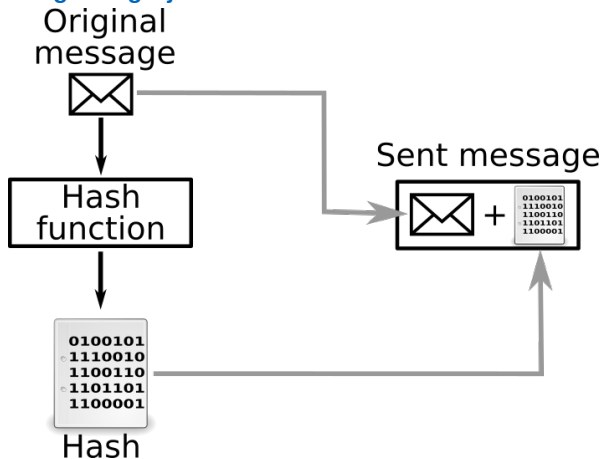
## Digital Signatures

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# Message Authentication codes

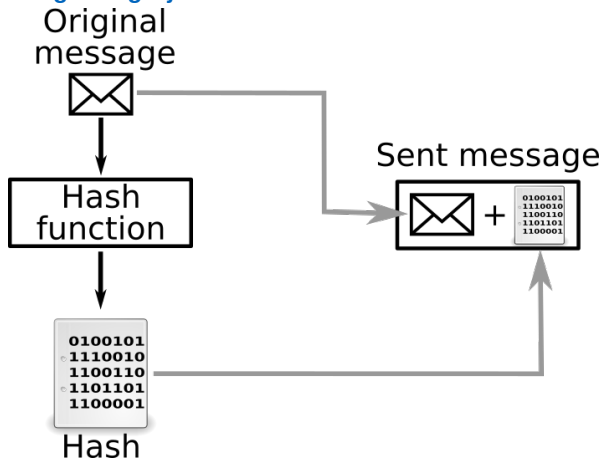
Use of hash functions to establish message integrity



# Message Authentication codes

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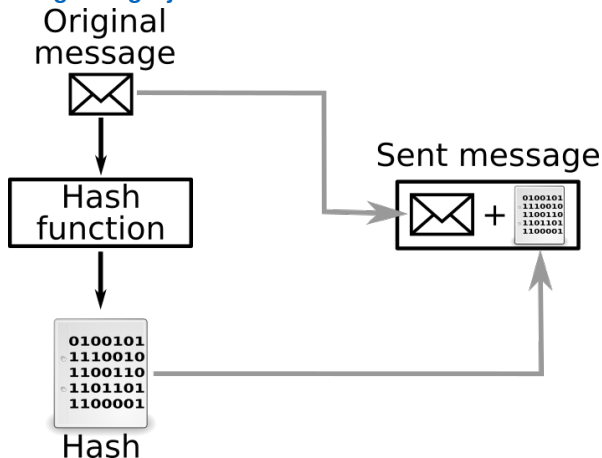
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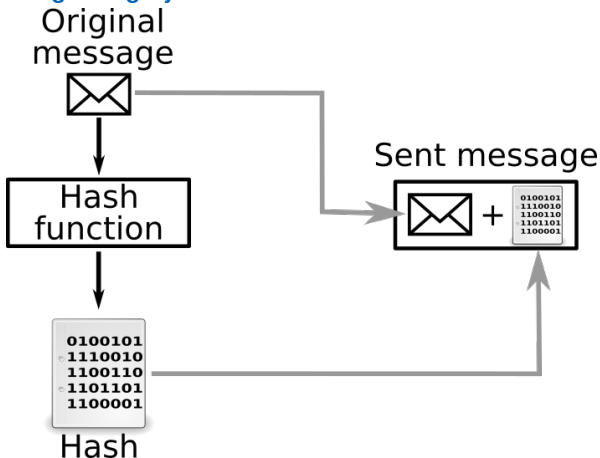
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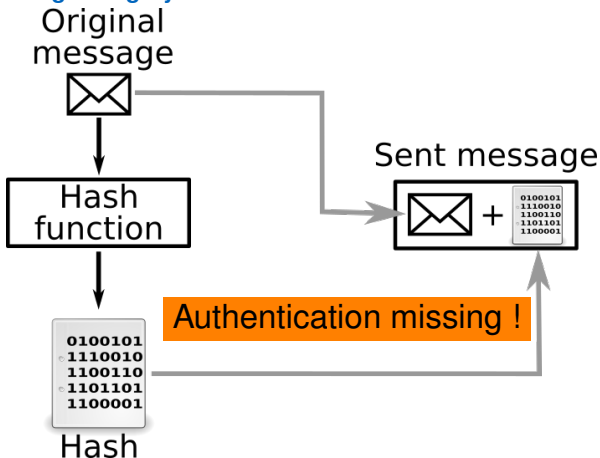
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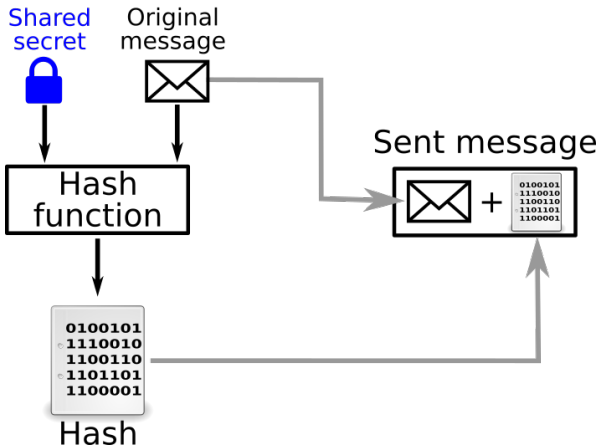




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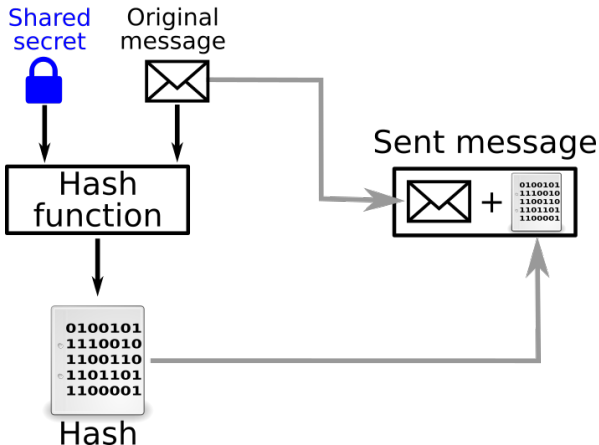
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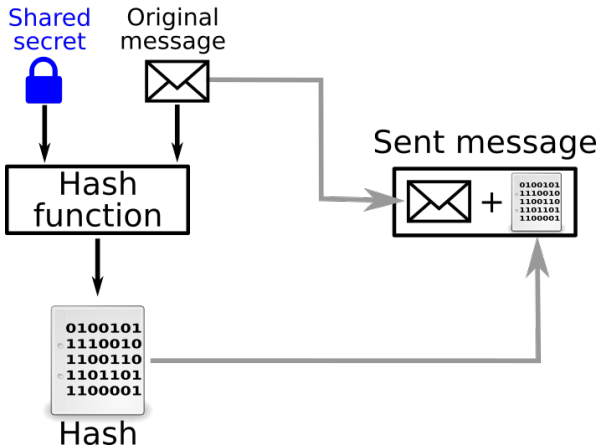
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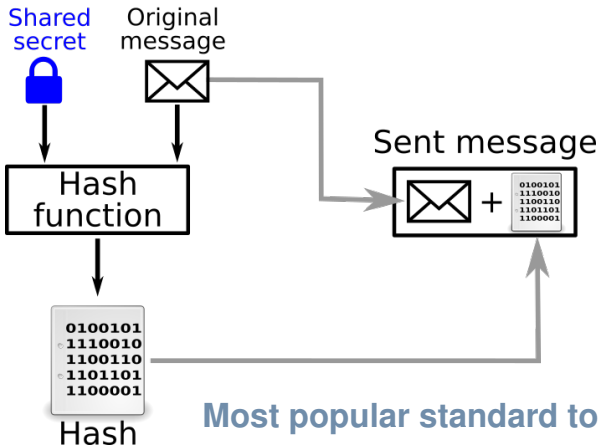
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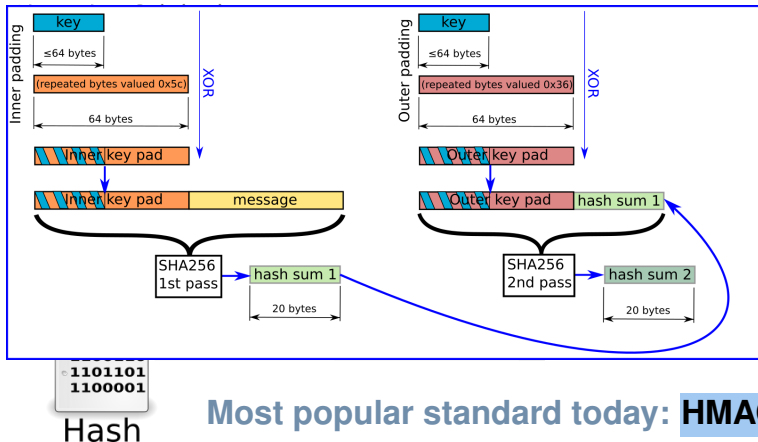
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Can be used with e.g. SHA-2/3

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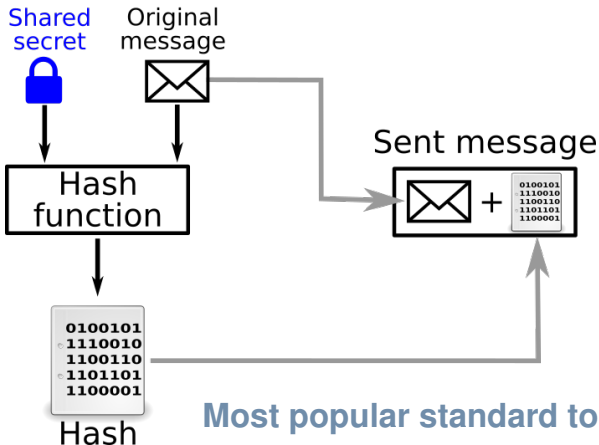


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# Digital signatures

A digital signature is a **cryptographic technique** to prove the identity of the owner/creator of a document or to signify one's agreement with a document's content

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# Digital signatures

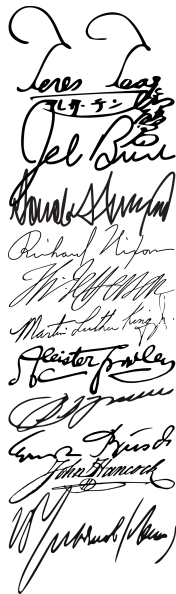
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## Requirements for a digital signature

The signature must be unique to an entity in the network, verifiable and non-forgable

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Handwritten signatures (from top to bottom):  
Loren Loag  
Jel Brue  
Richard Wipon  
Martin Luther King  
Reister Fowler  
John Hancock  
Stephan Sigg



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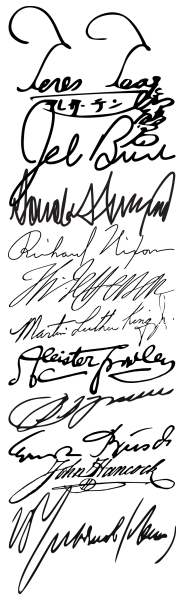
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Can we use the message authentication code for the signature?



Loren Loag  
Jel Brue  
Richard Nipon  
W. J. H. H. H.  
Martin Luther King  
Reister F. J. J.  
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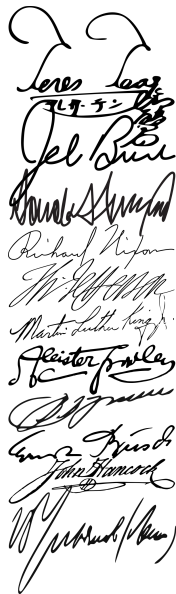
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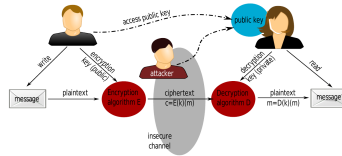
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**No:** To verify the signature, the receiver needs a copy → not unique



# Digital signatures

## Asymmetric encryption for Digital Signatures

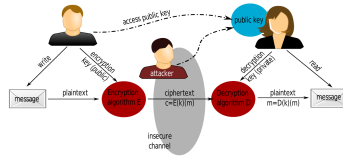


*Loren Loag*  
*Jel Brui*  
*Richard Nipon*  
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# Digital signatures

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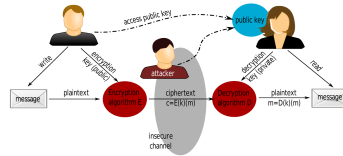
Signature is created by encrypting the message  $m$  with the private key:  $K_B^-(m)$



Handwritten signatures of various individuals, including "Loren Loag", "Jel Brui", "Richard Nipon", "Martin Luther King", "Steinbock", "John Hancock", and "John Hancock".

# Digital signatures

## Asymmetric encryption for Digital Signatures



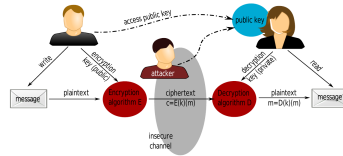
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Private key  $K_B^-$  unique to sender and can be verified by anyone with the public key:  $K_B^+(K_B^-(m)) = m$

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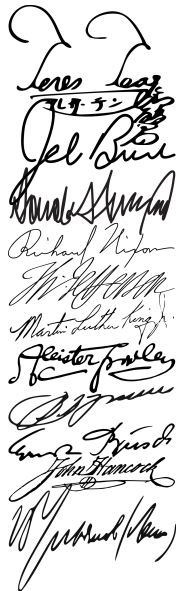
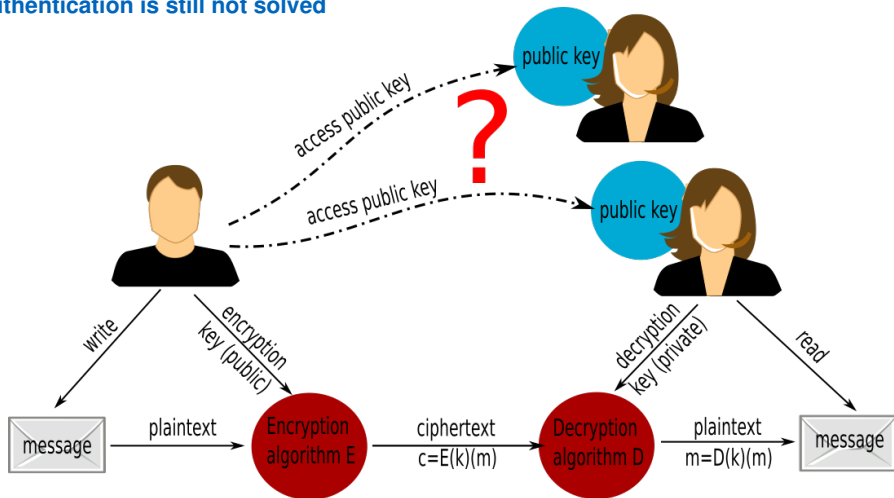
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Non-forgable since it is not computationally feasible to find  $m'$  with  $K_B^+(K_B^-(m)) = m'$

# Digital signatures

Authentication is still not solved



## Part II (10 min)

End-point authentication

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# End-point authentication

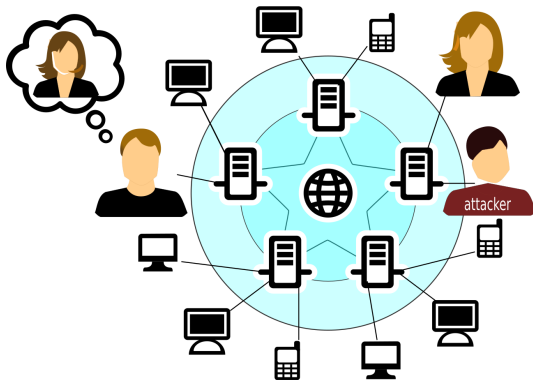
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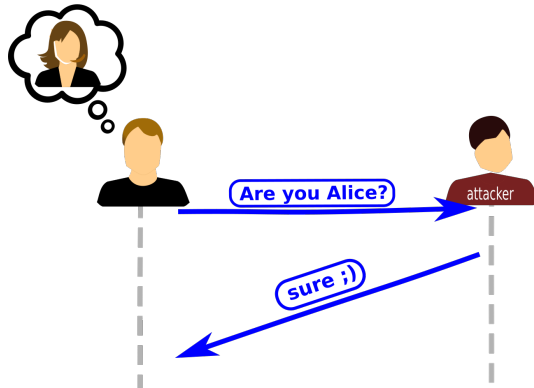
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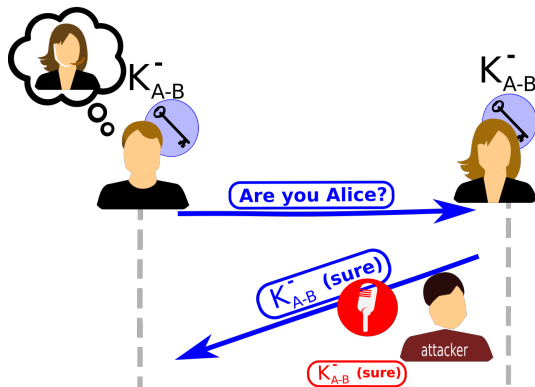
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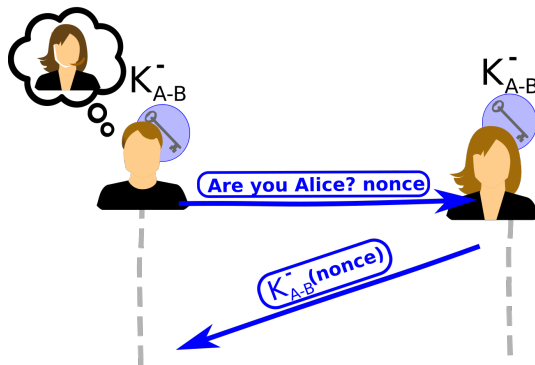
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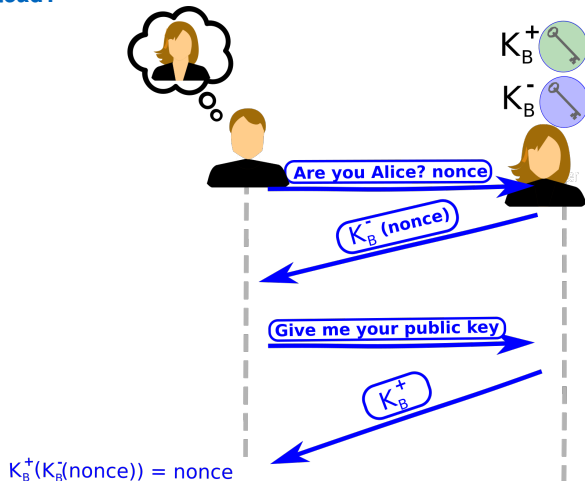
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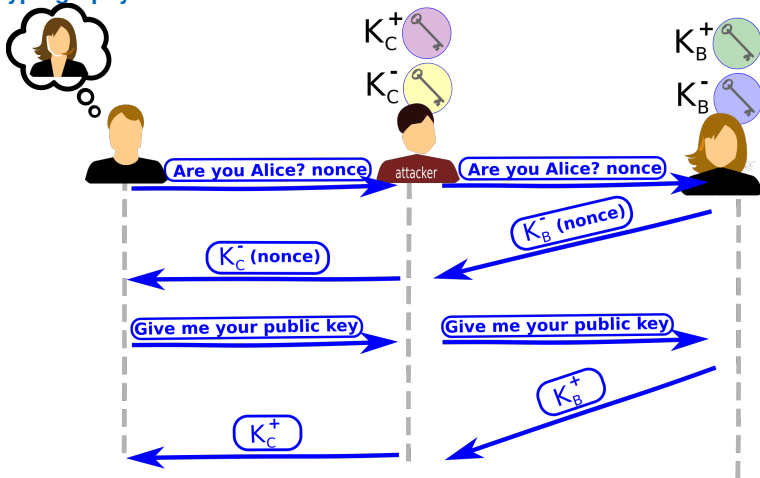
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Use asymmetric cryptography instead?



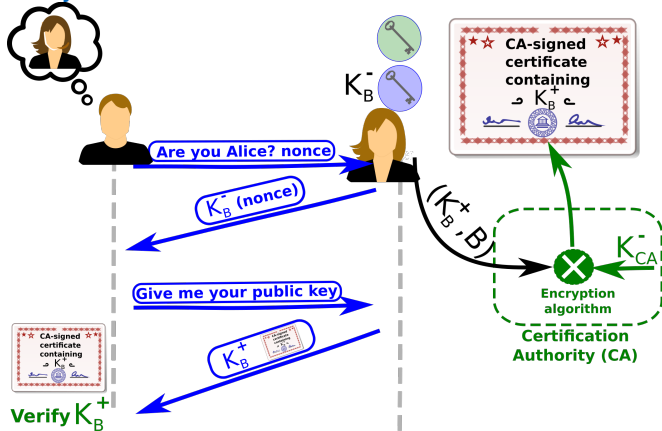
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# End-point authentication

## Solution: Certified Authority



⇒ CA issued certificate contains  $K_B^+$ ,  $B$ ; digitally signed by  $K_{CA}^-$



# Video: RSA 2016 – cryptographers panel (10 min)

Impact of the NSA-incident

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# Questions?

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Esa Vikberg

`esa.vikberg@aalto.fi`

Leo Lazar

`leo.lazar@aalto.fi`

# Literature

- J.F. Kurose, K.W. Ross: Computer Networking: A Top-Down approach (7th edition), Pearson, 2016.
- J.F. Kurose, K.W. Ross: Computer Networking: A Top-Down approach (6th edition), Addison-Wesley, 2012.

