
Lecture 21

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1 Softening Policy Elimination

In this lecture we are finally ready to discuss the state-of-the-art algorithm for the i.i.d. contextual bandit problem, which is both optimal and oracle-efficient [Agarwal et al., 2014]. Recall that the idea of Policy Elimination is to find $P_t \in \Delta(\Pi_t)$ such that $V(P_t, \pi) \leq 2K$ for all $\pi \in \Pi_t$ where

$$V(P, \pi) = \mathbb{E}_x \left[\frac{1}{P^\mu(\pi(x)|x)} \right]$$

is essentially the variance of the loss estimators that we want to control. To obtain an efficient algorithm, we need to forget about the idea of removing policies from Π . So is it possible to ensure $V(P_t, \pi) \leq 2K$ for all $\pi \in \Pi$ while at the same time P_t puts most of the weights on good policies?

Unfortunately this is too strong of a requirement. For example, if there is a bad policy π_{bad} which always picks a bad action a_{bad} with loss 1 and no other policy ever picks a_{bad} , then

$$2K \geq V(P_t, \pi_{\text{bad}}) = \mathbb{E}_x \left[\frac{1}{P_t^\mu(a_{\text{bad}}|x)} \right] = \frac{1}{(1 - K\mu)P_t(\pi_{\text{bad}}) + \mu}$$

which implies that $P_t(\pi_{\text{bad}})$ will be pretty large assuming μ is small. This is clearly not a good algorithm.

From this example, however, we can see that the condition $V(P_t, \pi) \leq 2K$ should be somehow relaxed for bad policies. Just as in Policy Elimination, whether a policy is good or bad can be roughly determined by its empirical performance compared to the empirically best policy. Specifically, recall the notation $\bar{\ell}_t(\pi) = \frac{1}{t} \sum_{\tau=1}^t \hat{\ell}_\tau(\pi(x_\tau))$ for the empirical average loss and $\pi_t^* = \operatorname{argmin}_{\pi \in \Pi} \bar{\ell}_t(\pi)$ for the empirically best policy up to time t . Define empirical average regret for a policy π to be

$$\text{Reg}_t(\pi) = \bar{\ell}_t(\pi) - \bar{\ell}_t(\pi_t^*).$$

We now relax the low-variance condition as: find P_t such that

$$V(P_t, \pi) \leq 2K + \beta \text{Reg}_{t-1}(\pi) \quad \forall \pi \in \Pi$$

for some parameter $\beta > 0$ to be specified later. Now there is hope to impose exploitation simultaneously. Specifically, we want $\sum_{\pi \in \Pi} P_t(\pi) \text{Reg}_{t-1}(\pi)$ to be as small as possible. How small can it be? The following lemma answers this question.

Lemma 1. *For any $\beta > 0$, there always exists a distribution $P \in \Delta(\Pi)$ such that*

$$\begin{aligned} \sum_{\pi \in \Pi} P(\pi) \text{Reg}_{t-1}(\pi) &\leq \frac{2K}{\beta} \\ V(P, \pi) &\leq 2K + \beta \text{Reg}_{t-1}(\pi) \quad \forall \pi \in \Pi. \end{aligned}$$

Proof. Define function $F_t : \Delta(\Pi) \rightarrow \mathbb{R}_+$ as

$$F_t(P) = \sum_{\pi \in \Pi} P(\pi) \text{Reg}_{t-1}(\pi) + \frac{2}{\beta} \mathbb{E}_x \left[\sum_{a=1}^K \ln \frac{1}{P^\mu(a|x)} \right].$$

The claim is that the minimizer of $F_t(P)$, which always exists due to compactness of $\Delta(\Pi)$ and continuousness of F_t , satisfies both conditions. To see this, first notice that we can extend the function to a set of “sub-distributions” $\Delta(\Pi)' = \{P \in \mathbb{R}_+^N : \sum_{\pi \in \Pi} P(\pi) \leq 1\}$ and still have

$$\min_{P \in \Delta(\Pi)} F_t(P) = \min_{P \in \Delta(\Pi)'} F_t(P).$$

This is because for any sub-distribution $P \in \Delta(\Pi)'$, one can make it a distribution by increasing the weight for policy π_{t-1}^* until the weights sum up to 1. This will only decrease the function value since $\text{Reg}_{t-1}(\pi_{t-1}^*) = 0$ and the second term of F_t is decreasing in any coordinate of P .

Next note that the derivate of F_t with respect to a policy π is

$$\nabla F_t(P)(\pi) = \text{Reg}_{t-1}(\pi) - \frac{2(1-K\mu)}{\beta} V(P, \pi).$$

Let P^* be a minimizer of F_t over $\Delta(\Pi)'$. By KKT conditions, we have

$$\text{Reg}_{t-1}(\pi) - \frac{2(1-K\mu)}{\beta} V(P^*, \pi) - \lambda_\pi + \lambda = 0 \quad (1)$$

for some Lagrangian multipliers $\lambda_\pi \geq 0$ and $\lambda \geq 0$. Multiply both sides by $P^*(\pi)$ and sum over $\pi \in \Pi$ gives

$$\begin{aligned} \sum_{\pi \in \Pi} P^*(\pi) \text{Reg}_{t-1}(\pi) &= \frac{2(1-K\mu)}{\beta} \sum_{\pi \in \Pi} P^*(\pi) V(P^*, \pi) + \sum_{\pi \in \Pi} P^*(\pi) \lambda_\pi - \lambda \quad (P^* \in \Delta(\Pi)) \\ &= \frac{2(1-K\mu)}{\beta} \sum_{\pi \in \Pi} P^*(\pi) V(P^*, \pi) - \lambda \quad (\text{complementary slackness}) \\ &\leq \frac{2}{\beta} \mathbb{E}_x \left[\sum_{\pi \in \Pi} \frac{P^*(\pi)}{P^*(\pi(x)|x)} \right] - \lambda = \frac{2K}{\beta} - \lambda \leq \frac{2K}{\beta}, \end{aligned}$$

showing that P^* satisfies the first condition. Moreover, the last equality above also implies $\lambda \leq \frac{2K}{\beta}$ since $\text{Reg}_{t-1}(\pi) \geq 0$. Rearranging Eq. (1) thus gives

$$V(P^*, \pi) \leq \frac{\beta}{2(1-K\mu)} (\text{Reg}_{t-1}(\pi) + \lambda) \leq 2K + \beta \text{Reg}_{t-1}(\pi),$$

where we assume $\mu \leq \frac{1}{2K}$ so that $2(1-K\mu) \geq 1$ (since otherwise we trivially have $V(P^*, \pi) \leq 1/\mu \leq 2K$). This shows that P^* satisfies the second condition too. \square

The question is now what β we should use. Assuming $\text{Reg}_{t-1}(\pi)$ concentrates well around the actual expected regret of π compared to π^* ,

$$\text{Reg}(\pi) \stackrel{\text{def}}{=} \bar{\ell}(\pi) - \bar{\ell}(\pi^*),$$

which is exactly what we hope for, $\text{Reg}_{t-1}(\pi)$ should be at most a constant. A reasonable choice of β would then be of order $1/\mu$, since $V(P, \pi)$ is trivially bounded by $1/\mu$. In other words, when a policy π is good, which means $\text{Reg}_{t-1}(\pi)$ is close to zero, we still require $V(P, \pi)$ to be close to $2K$, while when the policy is bad, which means $\text{Reg}_{t-1}(\pi)$ is a large constant, then there is almost no requirement on $V(P, \pi)$ with this choice of β .

On the other hand, this means that the exploitation constraint is $\sum_{\pi \in \Pi} P(\pi) \text{Reg}_{t-1}(\pi) = \mathcal{O}(K\mu)$, which also makes sense because μ should be of order $1/\sqrt{T}$, and if the per round regret is of order $1/\sqrt{T}$, then the overall regret over T rounds is of order \sqrt{T} . With some specific constant (chosen based on the analysis), this leads to the final algorithm called ILOVETOCONBANDITS (see Algorithm 1).

2 Oracle-Efficiency

To discuss oracle-efficiency, keep in mind that as in Policy Elimination, the true context distribution in the definition of V can be replaced by the empirical distribution of observed contexts, that is, a

Algorithm 1: ILOVETOCONBANDITS (colloquially referred as Mini-monster)

Input: failure probability $\delta \in (0, 1)$

Initialization: let $\mu = \min \left\{ \frac{1}{K}, \sqrt{\frac{\ln(TN/\delta)}{TK}} \ln T \right\}$

for $t = 1, \dots, T$ **do**

find P_t such that

$$\sum_{\pi \in \Pi} P_t(\pi) \text{Reg}_{t-1}(\pi) \leq 20K\mu$$

$$V(P_t, \pi) \leq 2K + \frac{\text{Reg}_{t-1}(\pi)}{10\mu} \quad \forall \pi \in \Pi.$$

play $a_t \sim P_t^\mu(\cdot | x_t)$

uniform distribution over x_1, \dots, x_{t-1} at time t . (For simplicity, the analysis of next section will assume that the true context distribution is known instead.)

According to the proof of Lemma 1, to find distribution P_t it suffices to solve the optimization problem $\operatorname{argmin}_{P \in \Delta(\Pi)} F_t(P)$ (in fact an approximate solution is enough). This is in fact very similar to FTRL with a special regularizer. To see how to solve it efficiently with the oracle, notice that the derivative of $F_t(P)$ with respect to a policy π can be written as (with $\beta = 1/(10\mu)$)

$$\nabla F_t(P)(\pi) = \frac{1}{t-1} \sum_{\tau=1}^{t-1} \left(\hat{\ell}_\tau(\pi(x_\tau)) - \hat{\ell}_\tau(\pi_{t-1}^*(x_\tau)) \right) - \frac{20\mu(1-K\mu)}{(t-1)} \sum_{\tau=1}^{t-1} \frac{1}{P^\mu(\pi(x_\tau) | x_\tau)}.$$

Since the part involving π_{t-1}^* is independent of π , if we feed the oracle with a training set

$$\mathcal{S} = \left\{ \left(x_1, \hat{\ell}_1 - \frac{20\mu(1-K\mu)}{P^\mu(\cdot | x_1)} \right), \dots, \left(x_{t-1}, \hat{\ell}_{t-1} - \frac{20\mu(1-K\mu)}{P^\mu(\cdot | x_{t-1})} \right) \right\},$$

we have $\text{ERM}(\mathcal{S}) = \operatorname{argmin}_{\pi \in \Pi} \nabla F_t(P)(\pi)$. In other words, the oracle can tell us the minimum coordinate of the gradient of $F_t(P)$ for any P , which opens up many possibilities to utilize the theory of optimization to find P_t . For example, one can directly apply the Frank-Wolfe algorithm (also known as conditional gradient method). Specifically, for a constraint convex optimization problem $\min_{w \in \Omega} f(w)$, the Frank-Wolfe algorithm performs the following iterative updates (staring with an arbitrary $w_1 \in \Omega$):

$$v_k = \operatorname{argmin}_{v \in \Omega} \langle v, \nabla f(w_k) \rangle$$

$$w_{k+1} = (1 - \gamma_k)w_k + \gamma_k v_k$$

for some step-size γ_k (default choice is $2/(k+1)$). When Ω is the simplex, the first step is exactly to find the minimum coordinate of the gradient. Therefore, with the oracle we can implement the Frank-Wolfe algorithm to solve $\operatorname{argmin}_{P \in \Delta(\Pi)} F_t(P)$ efficiently. We omit the details on how many iterations are needed but it will be polynomial in T, K , and $\ln N$.

Importantly, notice that unlike gradient descent, Frank-Wolfe leads to a sparse solution: when Ω is the simplex, after k rounds w_k has only k non-zero coordinates (assuming w_1 concentrates on one element to start with). This means that P_t 's are all sparse distributions and operations involving P_t , such as constructing the training set \mathcal{S} and sampling a_t , are all efficient.

Instead of using Frank-Wolfe, another possibility is to do some kind of coordinate descent: iteratively use the oracle to fine the coordinate with minimum derivative and adjust the weight for this coordinate appropriately. This is exactly the method taken in [Agarwal et al., 2014]. In fact, with additional tricks that are specialized for this task, it was shown that over T rounds only $\mathcal{O}(\sqrt{T})$ oracle calls are needed, which also implies that all P_t 's are $\mathcal{O}(\sqrt{T})$ -sparse.

3 Regret Analysis

Finally in this section we prove that ILOVETOCONBANDITS enjoys optimal regret. The key is to show the following concentration results on regret.

Lemma 2. *With probability $1 - \delta/2$, Algorithm 1 ensures that for all $t \in [T]$ and all $\pi \in \Pi$,*

$$\text{Reg}(\pi) \leq 2\text{Reg}_t(\pi) + \epsilon_t \quad \text{and} \quad \text{Reg}_t(\pi) \leq 2\text{Reg}(\pi) + \epsilon_t$$

where $\epsilon_t = \frac{20C}{\mu t} + 15K\mu$ and $C = \ln\left(\frac{4NT}{\delta}\right)\ln T$.

Proof. By Freedman's inequality and a union bound, we have with probability $1 - \delta/2$, for all $t \in [T]$, all $\pi \in \Pi$, and any $\lambda \in [0, \mu]$,

$$|\bar{\ell}_t(\pi) - \bar{\ell}(\pi)| \leq \frac{\lambda}{t} \sum_{\tau=1}^t V(P_\tau, \pi) + \frac{\ln\left(\frac{4NT}{\delta}\right)}{\lambda t}.$$

Specifically picking $\lambda = \frac{\mu}{\ln T}$ gives

$$|\bar{\ell}_t(\pi) - \bar{\ell}(\pi)| \leq \frac{\mu}{t \ln T} \sum_{\tau=1}^t V(P_\tau, \pi) + \frac{C}{\mu t}. \quad (2)$$

Now we use induction to prove the lemma. The base case $t = 0$ is trivial. Assuming the statement holds for all rounds before time t , we have by the algorithm

$$V(P_\tau, \pi) \leq 2K + \frac{\text{Reg}_{\tau-1}(\pi)}{10\mu} \leq 2K + \frac{\text{Reg}(\pi)}{5\mu} + \frac{\epsilon_{\tau-1}}{10\mu} \quad (3)$$

for all $\tau = 2, \dots, t$ and $V(P_1, \pi) \leq 2K$. Therefore, we have

$$\begin{aligned} \text{Reg}(\pi) - \text{Reg}_t(\pi) &= \bar{\ell}(\pi) - \bar{\ell}(\pi^*) - \bar{\ell}_t(\pi) + \bar{\ell}_t(\pi_t^*) \\ &\leq \bar{\ell}(\pi) - \bar{\ell}(\pi^*) - \bar{\ell}_t(\pi) + \bar{\ell}_t(\pi^*) && (\text{by optimality of } \pi_t^*) \\ &\leq \frac{2C}{\mu t} + \frac{\mu}{t \ln T} \sum_{\tau=1}^t (V(P_\tau, \pi) + V(P_\tau, \pi^*)) && (\text{by Eq. (2)}) \\ &\leq \frac{2C}{\mu t} + \frac{4K\mu}{\ln T} + \frac{\text{Reg}(\pi)}{5 \ln T} + \frac{1}{5t \ln T} \sum_{\tau=2}^t \epsilon_{\tau-1} && (\text{by Eq. (3) and } \text{Reg}(\pi^*) = 0) \\ &\leq \frac{2C}{\mu t} + \frac{4K\mu}{\ln T} + \frac{\text{Reg}(\pi)}{5 \ln T} + \frac{8C}{\mu t} + \frac{3K\mu}{\ln T} && (\text{by plugging in } \epsilon_\tau) \\ &\leq \frac{10C}{\mu t} + 7K\mu + \frac{\text{Reg}(\pi)}{2} \leq \frac{\epsilon_t}{2} + \frac{\text{Reg}(\pi)}{2}. \end{aligned}$$

Rearranging proves $\text{Reg}(\pi) \leq 2\text{Reg}_t(\pi) + \epsilon_t$. Similarly, we also have

$$\begin{aligned} \text{Reg}_t(\pi) - \text{Reg}(\pi) &= \bar{\ell}_t(\pi) - \bar{\ell}_t(\pi_t^*) - \bar{\ell}(\pi) + \bar{\ell}(\pi^*) \\ &\leq \bar{\ell}_t(\pi) - \bar{\ell}_t(\pi_t^*) - \bar{\ell}(\pi) + \bar{\ell}(\pi_t^*) && (\text{by optimality of } \pi^*) \\ &\leq \frac{2C}{\mu t} + \frac{\mu}{t \ln T} \sum_{\tau=1}^t (V(P_\tau, \pi) + V(P_\tau, \pi_t^*)) && (\text{by Eq. (2)}) \\ &\leq \frac{2C}{\mu t} + \frac{4K\mu}{\ln T} + \frac{\text{Reg}(\pi)}{5 \ln T} + \frac{\text{Reg}(\pi_t^*)}{5 \ln T} + \frac{1}{5t \ln T} \sum_{\tau=2}^t \epsilon_{\tau-1} && (\text{by Eq. (3)}) \\ &\leq \frac{2C}{\mu t} + \frac{4K\mu}{\ln T} + \frac{\text{Reg}(\pi)}{5 \ln T} + \frac{\epsilon_t}{5 \ln T} + \frac{8C}{\mu t} + \frac{3K\mu}{\ln T} \\ &\leq \frac{14C}{\mu t} + 10K\mu + \text{Reg}(\pi) \leq \epsilon_t + \text{Reg}(\pi), \end{aligned} \quad (4)$$

where Step (4) uses the fact $\text{Reg}(\pi) \leq 2\text{Reg}_t(\pi) + \epsilon_t$ just proven above with π set to π_t^* , and also the fact $\text{Reg}_t(\pi_t^*) = 0$. Rearranging then proves $\text{Reg}_t(\pi) \leq 2\text{Reg}(\pi) + \epsilon_t$ as well. \square

The final regret bound is now a simple application of this lemma and the exploitation constraint of the algorithm.

Theorem 1. *Algorithm 1 ensures that with probability $1 - \delta$, we have $\mathcal{R}_T = \tilde{\mathcal{O}}\left(\sqrt{TK \ln(N/\delta)}\right)$.*

Proof. The first step is exactly the same as analyzing Policy Elimination: by Azuma's inequality we have with probability $1 - \delta/2$,

$$\sum_{t=1}^T \ell_t(a_t) \leq \sum_{t=1}^T \sum_{\pi \in \Pi} P_t(\pi) \bar{\ell}(\pi) + TK\mu + \mathcal{O}\left(\sqrt{T \ln(1/\delta)}\right).$$

Conditioning on this event and the event stated in Lemma 2, which happen simultaneously with probability $1 - \delta$, we have

$$\begin{aligned} \mathcal{R}_T &\leq \sum_{t=1}^T \sum_{\pi \in \Pi} P_t(\pi) \text{Reg}(\pi) + TK\mu + \mathcal{O}\left(\sqrt{T \ln(1/\delta)}\right) \\ &\leq 2 \sum_{t=1}^T \sum_{\pi \in \Pi} P_t(\pi) \text{Reg}_{t-1}(\pi) + \sum_{t=2}^T \epsilon_{t-1} + TK\mu + \mathcal{O}\left(\sqrt{T \ln(1/\delta)}\right) \\ &\leq 56TK\mu + \frac{40C \ln T}{\mu} + \mathcal{O}\left(\sqrt{T \ln(1/\delta)}\right), \quad (\text{by the exploitation constraint}) \end{aligned}$$

which is of order $\tilde{\mathcal{O}}\left(\sqrt{TK \ln(N/\delta)}\right)$ with the optimal tuning of μ . \square

References

Alekh Agarwal, Daniel Hsu, Satyen Kale, John Langford, Lihong Li, and Robert Schapire. Taming the monster: A fast and simple algorithm for contextual bandits. In *Proceedings of the 31st International Conference on Machine Learning*, 2014.