

Analyzing Bitcoin Blockchain Data DMQL

Project Milestone-2

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I. INTRODUCTION

The primary goal of this project is to analyze Bitcoin blockchain data to identify transaction patterns, understand the distribution of funds, and explore the flow of Bitcoin across the network. To facilitate this, a database system will be developed to efficiently store and query blockchain data, enabling faster insights and analyses.

II. PROBLEM STATEMENT

As the volume of Bitcoin transactions grows, efficiently storing, querying, and analyzing blockchain data has become increasingly challenging. The decentralized and cryptographically secure nature of blockchain adds complexity to data management and transaction analysis. Current tools often struggle with scalability and real-time processing, limiting the ability to extract meaningful insights from vast amounts of data.

The proposed database system will solve the problem of efficiently managing and analyzing large-scale Bitcoin blockchain data. Blockchain data is highly complex and rapidly growing, consisting of millions of transactions spread across thousands of blocks. Storing and processing such vast amounts of data in a structured way is essential for extracting insights, such as tracking transaction patterns, detecting anomalies, and analyzing the flow of Bitcoin.

A. Why do you need a database instead of an Excel file?

- Excel files are limited in their ability to handle large datasets. Bitcoin blockchain data involves gigabytes or terabytes of information, far exceeding Excel's capacity. A database system can store and manage vast amounts of data while ensuring efficient access and performance.
- Databases allow for complex querying using structured query languages to retrieve specific information quickly. Excel lacks such advanced query capabilities, making it cumbersome to search, filter, and analyze large datasets.
- As data size increases, Excel becomes slow and prone to crashes. A database, however, is optimized for performance and can handle large data volumes without compromising speed or reliability.
- Blockchain data is sensitive, and maintaining its integrity is crucial. Databases have features like constraints which are

crucial to maintain the integrity of complex blockchain data. Excel lacks such robust features.

- Databases can be automated and integrated with other systems, such as data analytics tools or machine learning models, allowing for seamless, real-time processing of blockchain data. Excel is not designed for these types of interactions.

III. TARGET USER

The target users of the proposed Bitcoin blockchain database system will be individuals and organizations involved in cryptocurrency analysis, financial services, and blockchain technology. These may include:

- Blockchain Analysts and Researchers
- Cryptocurrency Exchanges
- Regulatory and Compliance Teams
- Financial Institutions

A. Database Administrator

The database will likely be administered by Data Engineers or Database Administrators (DBAs) working within the IT departments of organizations like exchanges, financial firms, or blockchain research firms. They will ensure the database operates efficiently, and data is streamed appropriately from the blockchain. These administrators will also manage database backups, data updates, and user access controls.

B. Real-life Scenario

When there's a large cryptocurrency exchange, "CryptoTradeX," that handles millions of Bitcoin transactions daily, relies on fast data access to ensure real-time price calculations, monitor fraud patterns, and comply with financial regulations. Their blockchain analysts regularly analyze transaction trends, while compliance officers use the system to detect and investigate potentially fraudulent transactions. Data engineers at CryptoTradeX would administer the database, ensuring smooth operation and handling any technical issues that arise, keeping the data up-to-date for various departments across the organization.

In this scenario, the database serves multiple users—analysts, compliance officers, and financial executives—while being managed by an IT team skilled in handling large-scale blockchain data systems.

IV. DATABASE SCHEMA

1) The database contains the following four 5 tables. The table block reference is new addition created after milestone 1. The descriptions of each field are specified in the images.

1. Blocks Table

Field	Data Type	Description
hash	VARCHAR(64)	Unique identifier (hash) of the block.
size	BIGINT	Total size of the block in bytes.
stripped_size	BIGINT	Size of the block excluding witness data (used in SegWit transactions).
weight	BIGINT	Weighted size of the block, taking SegWit into account.
number	BIGINT	Block height or sequence number, starting from 0 for the genesis block.
version	BIGINT	Block version number, which indicates the rules for block validation.
merkle_root	VARCHAR(64)	Hash of the Merkle tree of transactions in the block, used to verify transactions.
timestamp	TIMESTAMP	Time when the block was mined (Unix time format).
nonce	VARCHAR(64)	Value that miners vary to create a valid block hash in proof-of-work.
bits	VARCHAR(64)	Target threshold for the block hash, used to adjust mining difficulty.
transaction_count	BIGINT	Number of transactions included in the block.

Fig. 1. Blocks Table Attributes

Field	Data Type	Description
block_hash	VARCHAR(64)	Unique identifier (hash) of the block.
block_timestamp	TIMESTAMP	The time when the block was mined (Unix time format).

Fig. 2. Block Reference Table Attributes

2. Transactions Table

Field	Data Type	Description
hash	VARCHAR(64)	Unique identifier (hash) of the transaction.
size	BIGINT	Total size of the transaction in bytes.
virtual_size	BIGINT	Size of the transaction when weighted by SegWit data.
version	BIGINT	Transaction version number, specifying transaction structure.
lock_time	BIGINT	Time or block number after which the transaction is valid.
block_hash	VARCHAR(64)	Hash of the block containing this transaction.
block_number	BIGINT	Block height or sequence number of the block containing this transaction.
block_timestamp	TIMESTAMP	Time when the block containing the transaction was mined.
input_count	BIGINT	Number of inputs in the transaction.
output_count	BIGINT	Number of outputs in the transaction.
input_value	BIGINT	Total value of inputs in the transaction (in satoshis).
output_value	BIGINT	Total value of outputs in the transaction (in satoshis).
fee	BIGINT	Transaction fee, calculated as input value minus output value (in satoshis).

Fig. 3. Transaction Table Attributes

3. Transaction Inputs Table

Field	Data Type	Description
transaction_hash	VARCHAR(64)	Hash of the transaction this input belongs to.
block_hash	VARCHAR(64)	Hash of the block containing this transaction.
block_number	BIGINT	Block height of the block containing this transaction.
block_timestamp	TIMESTAMP	Timestamp of the block containing this transaction.
index	BIGINT	Index of the input in the transaction.
spent_transaction_hash	VARCHAR(64)	Hash of the previous transaction where this input's value originated.
spent_output_index	BIGINT	Index of the output in the previous transaction that this input is spending.
script_asm	TEXT	Script in assembly language that unlocks the input for spending.
script_hex	TEXT	Script in hexadecimal format for unlocking the input.
sequence	BIGINT	Sequence number used for transaction replacement before lock time is reached.
required_signatures	BIGINT	Number of signatures required to authorize spending the input.
type	VARCHAR(255)	Type of input script (e.g., P2PKH, P2SH).
addresses	TEXT[]	List of addresses associated with the input.
value	BIGINT	Value of the input in satoshis.

Fig. 4. Transaction Inputs Table Attributes

4. Transaction Outputs Table

Field	Data Type	Description
transaction_hash	VARCHAR(64)	Hash of the transaction this output belongs to.
block_hash	VARCHAR(64)	Hash of the block containing this transaction.
block_number	BIGINT	Block height of the block containing this transaction.
block_timestamp	TIMESTAMP	Timestamp of the block containing this transaction.
index	BIGINT	Index of the output in the transaction.
script_asm	TEXT	Script in assembly language that locks the output for spending.
script_hex	TEXT	Script in hexadecimal format that locks the output.
required_signatures	BIGINT	Number of signatures required to unlock the output.
type	VARCHAR(255)	Type of output script (e.g., P2PKH, P2SH).
addresses	TEXT[]	List of addresses associated with the output.
value	BIGINT	Value of the output in satoshis.

Fig. 5. Transaction Outputs Table Attributes

V. DATABASE DESIGN

The database design for this project focuses on capturing the essential components of the Bitcoin blockchain, including blocks, transactions, and their associated inputs and outputs. Five tables were created: blocks, transactions, transaction inputs, and transaction outputs. Each table reflects a key aspect of the blockchain, with foreign key relationships linking transactions to blocks, and inputs/outputs to transactions. This structure ensures efficient data organization and allows for detailed queries between blocks, transactions, and value movements. The schema was implemented using PostgreSQL, with careful attention to indexing and relational integrity for performance and scalability. In addition to the database design, we have also provided an ER diagram to visually represent the relationships between the key entities (tables). This diagram highlights the one-to-many relationships between blocks and transactions, as well as between transactions and their respective inputs and outputs, providing a clear overview of the data structure and dependencies within the system. The script to create all the tables is available in the TableCreation.sql file and should be run when creating the instance of the db.

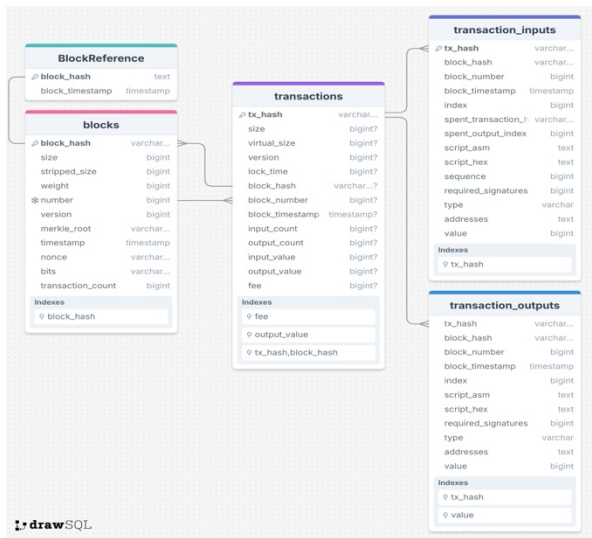


Fig. 6. Entity Relationship Diagram

IV. DATA ACQUISITION PROCESS

The blockchain data for this project was acquired using a free RPC endpoint provided by [Getblocks.io](https://getblocks.io) and the opensource library [bitcoin-etl](https://github.com/bitcoin-etl/bitcoin-etl). The following steps describe the data acquisition process:

A. Setup

A free RPC endpoint from [GetBlocks.io](https://getblocks.io) was used to connect to the Bitcoin network. The *bitcoin-etl* library was used to facilitate the extraction of blockchain data in JSON format.

B. Data Collection Command

The following command was used to export blocks and transactions for block numbers ranging from 861700 to 871668:

```

!bitcoinetl export_blocks_and_transactions \
--start-block 861700 --end-block 871668 \
--provider-uri <RPC-uri> \
--blocks-output blocks.json \
--transactions-output transactions.json
  
```

Parameters:

- `--start-block`: Specifies the starting block number for data collection.
- `--end-block`: Specifies the ending block number for data collection.
- `--provider-uri`: Provides the RPC endpoint URI for the Bitcoin network.
- `--blocks-output`: The output file for block data (blocks.json).
- `--transactions-output`: The output file for transaction data (transactions.json).

Output:

- `blocks.json`: A JSON file containing details about each block, such as block hash, size, timestamp, and transaction count.

C. Data Validation

After exporting the data, it was reviewed for integrity to ensure:

- All blocks within the specified range were exported.
- Transaction data matched the block data in terms of transaction counts.

D. Data Insertion into the Database

The exported JSON files (blocks.json and transactions.json) were parsed and loaded into the database in batches. This is because the transactions.json is a total of 68 GB and batching the data line by line significantly improved performance. The data insertion scripts are available in the `collection_loading.ipynb` file.

Batch Processing: Transactions were inserted into the database in batches of 100,000 lines for efficiency using Python scripts.

```

def load_transactions(batch_size=100000):
    transactions_batch = []
    inputs_batch = []
    outputs_batch = []

    with open(transactions_file, 'r') as f:
        with engine.connect() as conn:
            for line_num, line in enumerate(f, \
                start=1):
                transaction = json.loads(line) #
                # Insert batches into the database if line_num %
                # batch_size == 0:
                pd.DataFrame(transactions_batch)
                .to_sql('transactions', conn, if_exists='append',
                    index=False)
                pd.DataFrame(inputs_batch)
                .to_sql('transactioninputs', conn, if_exists='append',
                    index=False)
                pd.DataFrame(outputs_batch)
                .to_sql('transactionoutputs', conn, if_exists='append',
                    index=False)

                # Clear batches
                transactions_batch.clear()
                inputs_batch.clear()
                outputs_batch.clear()
                print(f"Processed {line_num} transactions..")
  
```

VII. NORMALIZATION TO BOYCE-CODD NORMAL FORM (BCNF)

- {block timestamp, block nonce, block bits, block transaction count}

- BlockReference: {block hash, block timestamp}
- Transactions: {tx hash, tx size, tx virtual size, tx version, tx lock time, block hash, tx_is coinbase, tx index, tx input count, tx output count, txinput value, tx-output value, tx fee}
- TransactionInput: {tx hash, input index, prev tx hash, prev output index, script_asm, script_hex, sequence}
- TransactionOutput: {tx_hash, output index, value, script asm, script hex}

- prev_output_index
- script_asm
- script_hex
- sequence

B. Functional Dependencies

The functional dependencies for each relation are listed below:

C. Functional Dependencies

The functional dependencies for each relation are as follows:

Blocks:

- block_hash determines:
 - block_size
 - block_stripped_size
 - block_weight
 - block_number
 - block_version
 - block_merkle_root
 - block_timestamp
 - block_nonce
 - block_bits
 - block_transaction_count

BlockReference:

- block_hash determines:
 - block_timestamp

Transactions:

- tx_hash determines:
 - tx_size
 - tx_virtual_size
 - tx_version
 - tx_lock_time
 - block_hash
 - tx_is_coinbase – tx_index
 - tx_input_count
 - tx_output_count
 - tx_input_value
 - tx_output_value
 - tx_fee

TransactionInputs:

- {tx_hash, input_index} determines:
 - prev_tx_hash

TransactionOutputs:

- {tx_hash, output_index} determines:
 - value
 - script_asm
 - script_hex

D. Boyce-Codd Normal Form (BCNF) Verification

The schema was analyzed to ensure all relations are in BCNF. A relation is in BCNF if, for every non-trivial functional dependency $X \rightarrow Y$, the determinant X is a superkey.

1) *Block Table:* - Attributes: {block hash, block_size, block_stripped_size, block_weight, block_number, block_version, block_merkle_root, block_timestamp, block_nonce, block_bits, block_transaction_count}. Analysis: The primary key block_hash is a superkey, and all non-trivial dependencies satisfy the BCNF condition. - Conclusion: The Block table is in BCNF.

2) *BlockReference Table:* - Attributes: {block hash, block timestamp}. - Analysis: The primary key block hash is a superkey, ensuring BCNF compliance. Conclusion: The BlockReference table is in BCNF. 3) *Transactions Table:* - Attributes: {tx_hash, tx_size, tx_virtual_size, tx_version, tx_lock_time, block_hash, tx_is coinbase, tx_index, tx input count, tx output count, tx input value, tx output value, tx fee}. Analysis: The primary key tx hash is a superkey, and no further decomposition is required. - Conclusion: The Transactions table is in BCNF.

4) *TransactionInput Table:* - Attributes: {tx hash, input index, prev tx hash, prev output index, script asm, script hex, sequence}. - Analysis: The composite primary key {tx hash, input index} is a superkey, there are only two attributes in the table, hence the table is automatically in BCNF form. - Conclusion: The TransactionInput table is in BCNF.

5) *TransactionOutput Table:* - Attributes: {tx hash, output index, value, script asm, script hex}. - Analysis: The composite primary key {tx

hash, output index} is a superkey, ensuring BCNF compliance. - Conclusion: The TransactionOutput table is in BCNF.

E. Conclusion

The final schema remains the same as the initial schema. All relations are in BCNF, ensuring that there is no redundancy or update anomalies.

VIII. CHALLENGES AND SOLUTIONS FOR HANDLING LARGE DATASET (Indexing)

While handling the 70 GB blockchain dataset, a performance issue was identified during query execution:

A. Challenge: Query Performance

A query filtering transactions by block_hash took approximately 40 seconds to execute. The reasons for the poor performance were:

The Transactions table contained millions of rows.

No index was present on the block_hash column. • A full table scan was required to retrieve matching rows.

B. Solution: Indexing with Clustered Index

To address the performance issue, a clustered index was implemented on the block_hash column. A clustered index physically organizes the table rows to store data with the same key value close together. This significantly improved query performance.

The steps involved in implementing the clustered index included:

- Creating an index on the block_hash column.
- Clustering the table using the created index.

C. Results

After implementing the clustered index:

- Query execution time reduced from 40 seconds to less than 1 second.
- Queries filtering by block_hash became significantly faster.

D. Conclusion

Various other indexes were also implemented but none provided large significant performance improvements. All index creation code is available in the indexes.sql script.

The clustered index optimized the physical data storage in the Transactions table, enabling efficient query execution for large datasets. This solution demonstrates the importance of indexing and storage optimization, particularly for:

- Filtering queries, such as those using block_hash. • Join operations, where block_hash is used as a foreign key.

IX. DATABASE CONSTRAINTS

A. Database Constraints

To ensure data integrity and consistency, the following constraints were implemented in the database:

- Foreign Key Constraints with Cascading Deletes:

- In the Transactions table, the block_hash column references Block(block_hash). The ON DELETE CASCADE clause ensures that when a block is deleted, all associated transactions are automatically removed. This maintains referential integrity by preventing orphaned transactions.

- In the TransactionOutput table, the tx_hash column references Transactions(tx_hash). The ON DELETE CASCADE clause ensures that when a transaction is deleted, all its outputs are also removed, avoiding inconsistencies in the output data.

- Check Constraints:

- In the Transactions table, the tx_fee must be greater than or equal to 0. This constraint ensures that transaction fees are always non-negative, which aligns with blockchain principles where fees cannot be negative.

- In the Block table, the block_weight must be greater than 0. This constraint ensures that blocks have a valid positive weight, reflecting their size and significance in the blockchain network. Blocks with zero or negative weight would not be valid in real-world blockchain systems.

B. Validation with Real-World Data

The schema and constraints were validated using realworld blockchain data samples. The validation process ensured that:

- The foreign key constraints maintained referential integrity between related tables, ensuring no orphaned records in the database.
- The check constraints enforced logical consistency, preventing invalid entries such as negative fees or invalid block weights.
- The schema adhered to database normalization principles, minimizing redundancy and ensuring scalability for large datasets.

C. Conclusion

The finalized schema, supported by the described constraints, ensures that the database is robust, consistent, and capable of handling large-scale blockchain data efficiently. The constraints enforce data integrity by:

- Automatically maintaining relationships between blocks, transactions, and outputs through cascading deletes.
- Ensuring logical validity by preventing negative fees and invalid block weights.

This approach not only enforces data integrity but also provides a flexible framework for future extensions and scalability.

D. Conclusion

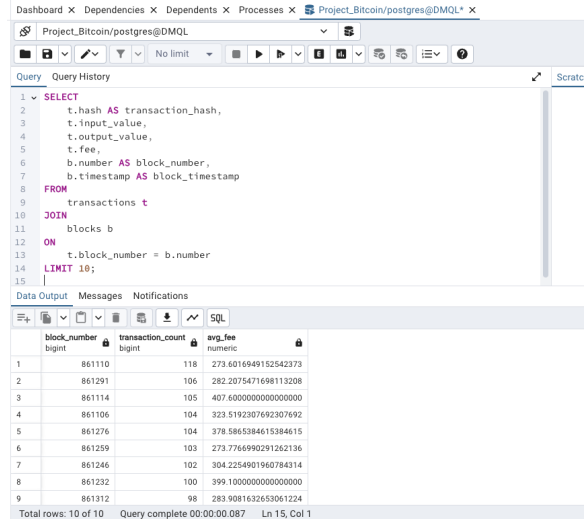
The finalized schema, supported by the described constraints, ensures that the database is robust, consistent, and capable of handling large-scale blockchain data efficiently. The constraints enforce data integrity while maintaining flexibility for future extensions.

X. QUERY EXECUTION

We can use various queries in the Bitcoin Database to uncover insights into transactions occurring on the blockchain. This Bitcoin blockchain database allows for a deep analysis of transactional and block-level insights. By querying the data, we can gather information on transaction patterns, such as the average transaction fees and input/output values. We can also track the frequency and size of blocks over time, understand trends in mining difficulty (using block size and weight), and analyze wallet activity by exploring transaction inputs and outputs. Furthermore, insights into the distribution of BTC across addresses, transaction volumes, and correlations between block times and transaction types can provide valuable knowledge about the network's performance and behavior. Here we show four examples of queries on the subset data by using various SQL queries that demonstrate the use of Group BY, subqueries, Join, and aggregate functions.

Note: Insertion queries are performed during data insertion. Since fundamentally blockchains are immutable we cannot perform any update or delete operations. Updates and deletes are prohibited to prevent tampering or fraud, ensuring the integrity of historical data. As a result, any changes must be handled by appending new data, rather than modifying existing records.

A. Using Joins Retrieve details of transactions along with their corresponding block information.

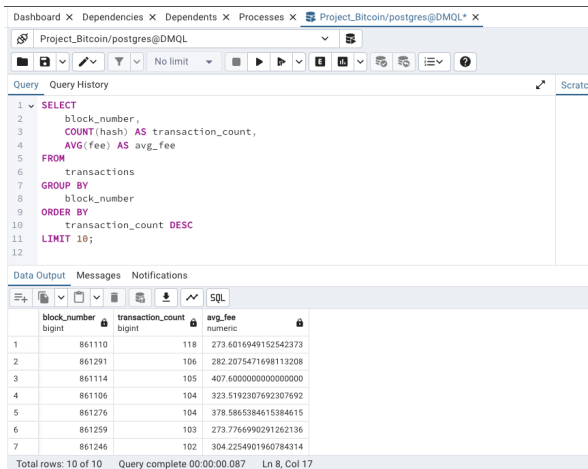


The screenshot shows a PostgreSQL query editor with a query that joins the 'transactions' table (t) and the 'blocks' table (b) on the condition 't.block_number = b.number'. The query is limited to 10 rows. The results table has three columns: 'block_number' (bigint), 'transaction_count' (bigint), and 'avg_fee' (numeric). The data is as follows:

block_number	transaction_count	avg_fee
861110	118	273.6016949152542373
861291	106	282.2075471698113208
861114	105	407.600000000000000000
861106	104	323.5192307692307692
861276	104	378.5865384615384615
861259	103	273.7766990291262136
861246	102	304.2254901960784314
861232	100	399.1000000000000000
861312	98	283.9081632653061224

Fig. 7. Using simple SELECT JOIN

B. Using Group By finding the total number of transactions and the average fee for each block.



The screenshot shows a PostgreSQL query editor with a query that groups transactions by 'block_number' and orders them by 'transaction_count' in descending order. The query is limited to 10 rows. The results table has three columns: 'block_number' (bigint), 'transaction_count' (bigint), and 'avg_fee' (numeric). The data is as follows:

block_number	transaction_count	avg_fee
861110	118	273.6016949152542373
861291	106	282.2075471698113208
861114	105	407.6000000000000000
861106	104	323.5192307692307692
861276	104	378.5865384615384615
861259	103	273.7766990291262136
861246	102	304.2254901960784314

Fig. 8. Using Group By

C. Using Sub-query to retrieve transactions where the fee is greater than the average fee across all transactions.

```

1 SELECT
2   hash,
3   input_value,
4   output_value,
5   fee
6 FROM
7   transactions
8 WHERE
9   fee > (SELECT AVG(fee) FROM transactions)
10 ORDER BY
11   fee DESC
12 LIMIT 10;
13

```

hash	input_value	output_value	fee
1	100713770	100000000	713770
2	1179914582	1179398458	516124
3	13012615	12606211	406404
4	275710906	275391142	319764
5	39236722	38948391	288331
6	133519511	133247788	271723
7	209947453	209723121	224332
8	749500	527500	222000

Fig. 9. Using Sub-queries

D. Using Joins with filtering to list all transactions with an output value greater than 1 BTC and show the associated block details.

```

1 SELECT
2   t.hash AS transaction_hash,
3   o.value AS output_value,
4   b.number AS block_number,
5   b.timestamp AS block_timestamp
6 FROM
7   transaction_outputs o
8 JOIN
9   transactions t
10 ON
11   o.transaction_hash = t.hash
12 JOIN
13   blocks b
14 ON
15   t.block_number = b.number
16 WHERE
17   o.value > 100000000 -- 1 BTC (1 BTC = 100,000,000 satoshis)
18 ORDER BY
19   o.value DESC
20 LIMIT 10;
21

```

transaction_hash	output_value	block_number	block_timestamp
1	2186503629	860597	2024-09-09 11:21:58
2	657374595	860938	2024-09-12 01:07:57
3	531759186	860586	2024-09-09 09:28:22
4	296144122	860806	2024-09-10 22:02:27
5	118899619	860088	2024-09-06 04:13:51

Fig. 10. Using Joins with Filtering

E. Daily Transaction Volume in Satoshis.

```

1 SELECT block_timestamp::date AS date, SUM(tx_output_value) AS total_volume
2 FROM Transactions
3 JOIN Blocks ON Transactions.block_hash = Blocks.block_hash
4 GROUP BY date
5 ORDER BY date;

```

date	total_volume
1	42929882823744
2	82406047729407
3	87861791817664
4	98001047985418
5	41589639732938
6	40955678358140

Fig. 11. Daily Transaction Volume

F. Largest Transaction by output value in Satoshis

```

1 SELECT tx_hash, MAX(value) AS largest_output
2 FROM transaction_outputs
3 GROUP BY tx_hash
4 ORDER BY largest_output DESC
5 LIMIT 10;

```

tx_hash	largest_output
1	16801007647900
2	15601007578200
3	7535408374933
4	4000000000000
5	3989999995460
6	3037132118904

Fig. 12. Largest Transaction by output value

G. Transactions with the highest fees

```

1 SELECT block_timestamp::date AS date, SUM(tx_output_value) AS total_volume
2 FROM Transactions
3 JOIN Blocks ON Transactions.block_hash = Blocks.block_hash
4 GROUP BY date
5 ORDER BY date;

```

date	total_volume
1	42929882823744
2	82406047729407
3	87861791817664
4	98001047985418
5	41589639732938
6	40955678358140

Fig. 13. Largest Transaction by output value

H. Top 10 Blocks by Total Transaction Value

Dashboard
X
Dependencies
X
Depends on
X
bitcoin_db/postgres@DMQL
X

bitcoin_db/postgres@DMQL

Query
Query History

```

1  SELECT  Blocks.block_hash, SUM(TransactionOutputs.value) AS total_value
2        FROM Transactions
3        JOIN Blocks ON Transactions.block_hash = Blocks.block_hash
4        JOIN TransactionOutputs ON Transactions.tx_hash = TransactionOutputs.tx_hash
5        GROUP BY Blocks.block_hash
6        ORDER BY total_value DESC
7        LIMIT 10;

```

Data Output
Messages
Notifications

	block_hash [PK] text	total_value numeric
1	00000000000000000000024352eb39119004ae033942c205b6781028b6d...	19234932674756
2	00000000000000000000020b4559264af7d4b5a103cf71bded81eefebf1a4c	18062493511070
3	000000000000000000000001dca15b0cf3785142392355a199f447e782ad...	19013955614682
4	00000000000000000000000263145254d0ee3ff4b671f62aa3280698c83225f7	10783146855070
5	00000000000000000000000000008042798121ba8eac7faced4036875d3142c8c85...	955085722242
6	000000000000000000000000792402d285b3c28698049f34748d159c099277...	934051248315
7	000000000000000000000000000004014702c722645109238a77f996ad6a0dd155f5...	870831409091
8	0000000000000000000000073c4d85645616082b70b2e9bf1c7cd4ade5081...	7234089761685
9	0000000000000000000000023b1122c0e53f3a214c38c949f7883611bc05...	7032041289015
10	00000000000000000000020f0a58e772642aeaa8f1d6636a2d0828f5629f...	6889332761964

Total rows: 10
Query complete 00.02.06.225

Fig. 14. Largest Transaction by output value

XI. QUERY EXECUTION ANALYSIS AND OPTIMIZATION

This section analyzes the execution plans for three queries to identify performance bottlenecks and proposes optimizations to improve query efficiency.

A. Query 1: Daily Transaction Volume

Query History

```

1 - SELECT block_timestamp::date AS date, SUM(x_output_value) AS total_value
2       FROM Transactions
3       JOIN Blocks ON Transactions.block_hash = Blocks.block_hash
4       GROUP BY date
5       ORDER BY date]
```

Data Output Messages Explain X Notifications

Graphical

Analysis

Statistics

#	Node	Rows	Actual	Plan	Loops
1	→ Aggregate (cost=1596825.11..1599450.68 rows=9969 width=36) (no...	→ 1546.61	68		9969
2	→ Gather Merge (cost=1596825.11..159951.37 rows=19938 width=...	→ 97.74	204	19938	
3	→ Sort (cost=1595825.08..1595850.01 rows=9969 width=36) (no...	→ 1546.61	68	9969	
4	→ Aggregate (cost=1595013.44..1595162.98 rows=9969 width=36)...	→ 48.67	68	9969	
5	Buffers: Buffers: Memory Usage: 417 KB → Hash Joiner: Join (cost=1716.3..151212.34 rows=7... Hash Cond: (transactions.block_hash = blocks.block_ha...	→ 0.42	13407859	16759823	
6	→ Seq Scan on public.transactions as transacti...	→ 0.42	13407859	16759823	
7	→ Hash (cost=591.69..591.69 rows=9969 width=...	→ 0.34	9969	9969	
8	Buffers: 16384 batches / 1 Memory Usage: 1733 KB → Seq Scan on public.blocks as blocks (co...	→ 0.34	9969	9969	

Query Objective: Calculate the total transaction output value for each day. **Issues Identified:**

- **Sequential Scans:** Sequential scans on the Blocks and Transactions tables increase the cost significantly.
- **Hash Joins:** The hash join between Blocks and Transactions consumes memory and processing resources due to the large number of rows.
- **Sorting Overhead:** Sorting by date before grouping adds to the computational cost.

Proposed Optimizations:

- Create indexes on `block_hash` in the Transactions and Blocks tables to speed up joins.
- Use a materialized view to precompute and store daily transaction volumes.

- Apply filters earlier to reduce the dataset size before joining and grouping.

B. Query 2: Transactions with Largest Output

Query History

```

1 -> SELECT tx_hash, MAX(value) AS largest_output
2       FROM TransactionOutputs
3       GROUP BY tx_hash
4       ORDER BY largest_output DESC
5       LIMIT 10;

```

Data Output

Messages

Explain X

Notifications

Graphical

Analysis

Statistics

#	Note	Rows Rows X	Actual	Plan	Loops
1	→ Limit (cost=9524360.87, 9524360.87 rows=10 width=73) (rows=1)		7 1	10	10
2	→ Sort (cost=9524360.87, 9524917.34 rows=222830 width=73) (rows=222830)		7 222830		222830
3	→ Aggregate (cost=9463091.69, 9519545.59 rows=222830 width=73) (rows=222830)		6 18035	40223576	222830
4	→ Gather Merge (cost=9463091.69, 9519508.59 rows=4) (rows=4)		6 9031	40334889	445660
5	→ Sort (cost=9462091.67, 9462648.78 rows=222830 width=73) (rows=222830)		6 6034	13644963	222830
6	→ Aggregate (cost=8260404.15, 9422594.64 rows=222830, Buckets: Buckets, Memory Usage: 677318) (rows=222830)		6 2032	13644963	222830
7	→ Seq Scan on public.transactionoutputs a		6 0.42	3645082	4181335

Query Objective: Retrieve the top 10 transactions with the highest output value.

Issues Identified:

- **Sequential Scans:** The query performs a sequential scan on the TransactionOutputs table, which is inefficient given its large size.
- **Aggregation Overhead:** The MAX function applied on a large dataset increases the computation time.
- **Sorting Cost:** Sorting all transaction outputs by value before applying the LIMIT adds unnecessary overhead.

Proposed Optimizations:

- Create an index on the value column in the TransactionOutputs table to reduce sorting time.
- Partition the table by tx_hash to improve the efficiency of grouping and aggregation.

C. Query 3: Blocks with the Most Transactions

Query

Query History

```

1 -> SELECT block_hash, COUNT(*) AS transaction_count
2       FROM Transactions
3       GROUP BY block_hash
4       ORDER BY transaction_count DESC
5       LIMIT 10;

```

Data Output

Messages

Explain X

Notifications

<

Graphical

Analysis

Statistics

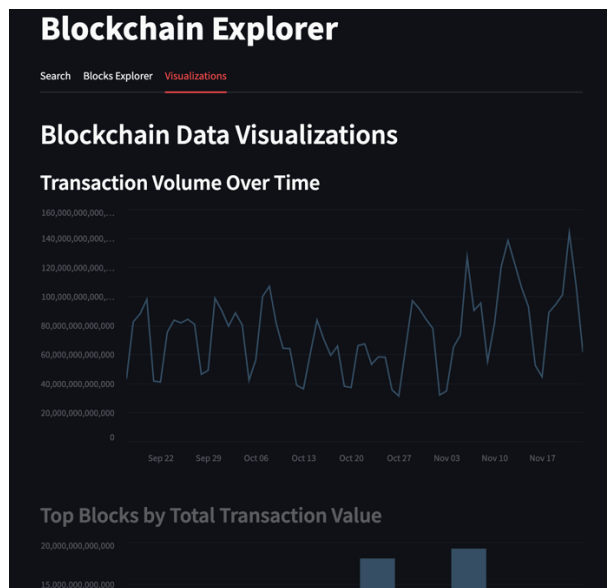
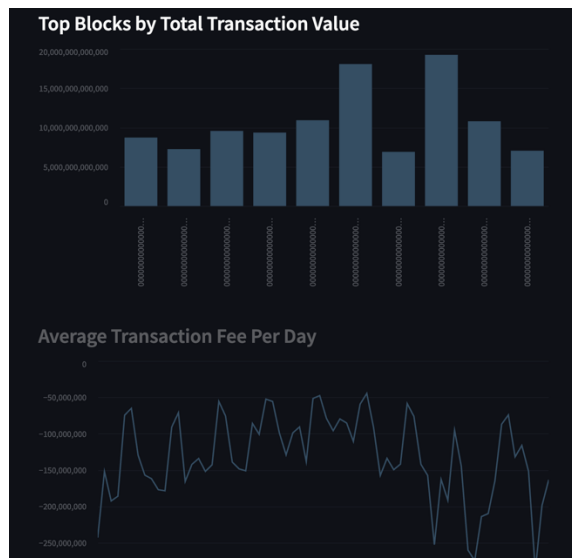
#	Node	Rows			Loops
		Rows X	Actual	Plan	
1.	→ Limit (cost=1512401.33..1512401.35 rows=10 width=73) (rows=10 ..		↑ 1	10	10
2.	→ Sort (cost=1512401.33..1512422.86 rows=8612 width=73) (row..		↑ 861.2	10	8612
3.	→ Aggregate (cost=1510033.38..1512215.23 rows=8612 widt..		4 1.16	9969	8612
4.	→ Gather Merge (cost=1510033.38..1512042.99 rows=1..		↑ 1.51	11413	17224
5.	→ Sort (cost=1509033.35..1509054.88 rows=8612 =		↑ 2.27	3804	8612
6.	→ Aggregate (cost=1508384.35..1508470.47 ro..		↑ 0.76	3804	8612
7.	→ Seq Scan on public.transactions as tran..		↑ 0.42	13407859	16759823

Query Objective: Identify the top 10 blocks with the highest transaction count.

Issues Identified:

- **Sequential Scans:** Sequential scans on the Transactions table increase cost.
- **Aggregation and Sorting Overhead:** Aggregating and sorting the transaction counts for all blocks add significant computation time.

This section includes interactive visualizations of blockchain data. For example, the "Transaction Volume Over Time" line chart illustrates daily transaction volumes, highlighting trends and spikes in activity over a given period. This bar chart ranks blocks by their total transaction output value, allowing users to identify the most valuable blocks in the blockchain. This line chart visualizes the average transaction fee per day, providing insights into the cost dynamics of blockchain transactions over time.



Implementation Details

- **Framework Used:** The website is built using using streamlit
- **Database Connection:** The database connection is created with sql alchemy module.
- **Visualization Tools:** Interactive charts are presented with the help of pandas and visualization functions from streamlit framework.

XII. PROJECT CONTRIBUTIONS

Work/Student	Rohan Shrikant Thorat	Gouri Ramdas Menon
Data Set Collection	Planned and Collected Data	
DB Design	Designed and Populated the Database	
Queries		Created queries for various SQL operations
Documentation and Project Planning		Documented and Planned project along with finding use cases

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- [5] Claude LLM for generating descriptions of attributes in Tables.