

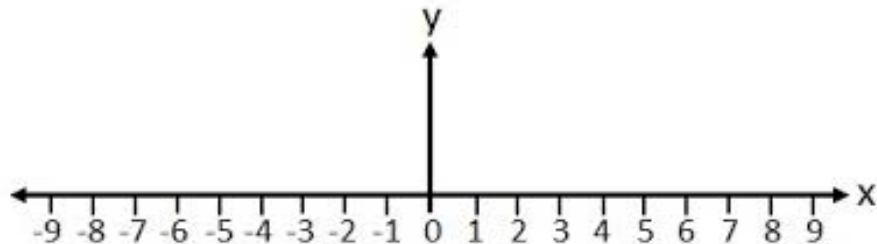
# Project: closest pair problem vs data stream

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# Simplify

If we can solve the closest pair problem, we can solve the problem “find minimum difference between any two elements in an unsorted array”.

Method: just let the elements in the array have y-axis value 0.  
Then the distance of the closest pair is the minimum difference.



# Hardness proof

Claim: A 1-pass deterministic algorithm to the **minimum difference problem** (and closest pair problem) can't have space complexity under  $\Omega(n)$ .

Proof: similar to the indexing problem.

Alice is the input, Bob is the program.

Alice send  $n$  distinct number + 1 repeat number to Bob.

Bob would answer the index(1~ $n$ ) of the repeated number.

Bob must use at least  $\Omega(n)$  space before he know the repeat number.

# approximate of minimum difference problem

Find  $r$  points:  $\frac{n}{r+1}$ -th,  $\frac{2n}{r+1}$ -th, ...,  $\frac{rn}{r+1}$ -th largest number by using find  $k$ -th number in streaming model.

Then at least a difference of two point

less than  $\frac{\max - \min}{r}$  can be found.

$O(r)$  pass,  $O(rn^{\frac{1}{r}} \log n)$  space.

