Mechanising Recursion Schemes with Magic-Free Coq Extraction

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Background

Hylomorphisms

Fold over Lists

One way to guarantee **recursive functions** are **well-defined** is via **Recursion Schemes**.

```haskell
foldr :: (a -> b -> b) -> b -> [a] -> b
foldr g b [] = b
foldr g b (x : xs) = g x (foldr g b xs)
```

There are many different kinds of Recursion Schemes (e.g. Folds, Paramorphisms, Unfolds, Apomorphisms, . . .)

Folds as Initial Algebras

```haskell
data Fix f = In { inOp :: f (Fix f) }

fold :: Functor f =>
          (f x -> x) ->
          Fix f ->
          x
fold a = f
    where f (In x) = (a . fmap f) x
```

$$\begin{tikzcd}
\texttt{f (Fix f)} \arrow[r, dotted] \arrow[d, "\texttt{In}"'] & \texttt{f x} \arrow[d, "\texttt{a}"] \\
\texttt{Fix f} \arrow[r, dotted] & \texttt{x}
\end{tikzcd}$$

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Least Fixed-Point

$\texttt{Fix f} \cong \texttt{f (Fix f)}$

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f-algebra

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\end{tikzcd}
$$

initial f-algebra

Hylomorphisms: Divide-and-conquer Computations

```haskell
hylo :: Functor f =>
          (f b -> b) ->
          (a -> f a) ->
          a -> b
hylo a c = a  . fmap (hylo a c) . c
```

$$
\begin{tikzcd}
\texttt{f a} \arrow[r, dotted] & \texttt{f b} \arrow[d, "\texttt{a}"] \\
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f-coalgebra
“divide”

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f-algebra
"conquer"

Folds as Hylomorphisms

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data Fix f = In { inOp :: f (Fix f) }

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          (f x -> x) ->
          Fix f ->
          x
fold a = a . fmap (fold a) . inOp
```

f-coalgebra

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f-algebra

Example: Nonstructural Recursion

```haskell
data TreeF a b = Leaf |  Node b a b

split [] = Leaf
split (h : t) = Node l h r
  where
    (l, r) = partition (\x -> x < h) t

merge Leaf = \acc -> acc
merge (Node l x r) = \acc -> l (x : r acc)
```

$$\begin{tikzcd}
\texttt{TreeF Int [Int]} \arrow[r, dotted, "\texttt{fmap qsort}"] & \texttt{TreeF Int ([Int] -> [Int])} \arrow[d, "\texttt{merge}"] \\
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Example: Nonstructural Recursion

TreeF Int-coalgebra

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TreeF Int-algebra

Conjugate Hylomorphisms

Every recursion scheme is a conjugate hylomorphism

recursion scheme	*adjunction*	*conjugates*	*para-hylo equation*	*algebra*
(hylo-shift law)	$\mathsf{Id} \dashv \mathsf{Id}$	$\alpha \dashv \alpha$	$x = a \cdot (id \vartriangle \mathsf{D}\, x \cdot \alpha\, C \cdot c) \;:\; A \leftarrow C$	$a : C \times \mathsf{D}\, A \to A$
mutual recursion	$\Delta \dashv (\times)$	ccf	$x_1 = a_1 \cdot (id \vartriangle \mathsf{D}\,(x_1 \vartriangle x_2) \cdot c) \;:\; A_1 \leftarrow C$ $x_2 = a_2 \cdot (id \vartriangle \mathsf{D}\,(x_1 \vartriangle x_2) \cdot c) \;:\; A_2 \leftarrow C$	$a_1 : C \times \mathsf{D}\,(A_1 \times A_2) \to A_1$ $a_2 : C \times \mathsf{D}\,(A_1 \times A_2) \to A_2$
accumulator	$- \times P \dashv (-)^P$	ccf	$x = a \cdot (outl \vartriangle ((\mathsf{D}\,(\Lambda\, x) \cdot c) \times P)) \;:\; A \leftarrow C \times P$	$a : C \times \mathsf{D}\,(A^P) \times P \to A$
course-of-values (§5.6)	$\mathsf{U}_{\mathsf{D}} \dashv \mathsf{Cofree}_{\mathsf{D}}$	ccf	$x = a \cdot (id \vartriangle \mathsf{D}\,(\mathsf{D}_\infty\, x \cdot [\![c]\!]) \cdot c) \;:\; A \leftarrow C$	$a : C \times \mathsf{D}\,(\mathsf{D}_\infty\, A) \to A$
finite memo-table (§5.6)	$\mathsf{U}_* \dashv \mathsf{Cofree}_*$	ccf	$x = a \cdot (id \vartriangle \mathsf{D}\,(\mathsf{D}_*\, x \cdot [\![c]\!]_*) \cdot c) \;:\; A \leftarrow C$	$a : C \times \mathsf{D}\,(\mathsf{D}_*\, A) \to A$

Table 1. Different types of para-hylos building on the canonical control functor (ccf); the coalgebra is $c : C \to \mathsf{D}\, C$ in each case.

R. Hinze, N. Wu, J. Gibbons: **Conjugate Hylomorphisms - Or: The Mother of All Structured Recursion Schemes**. POPL 2015.

Conjugate Hylomorphisms

- Every complex recursion scheme is an hylomorphism via its associated adjunction/conjugate pair
- (e.g) folds with parameters (accumulators) use the curry/uncurry adjunction
- A recursion scheme from comonads (RSFCs, Uustalu, Vene, Pardo, 2001) is an conjugate hylomorphism via the coEilemberg-Moore category for the cofree comonad

recursio				
(hylo-sh				
mutual r				
accumul				
course-of-values (§5.6)	$\mathsf{U}_{\mathsf{D}} \dashv \mathsf{Cofree}_{\mathsf{D}}$	ccf	$x = a \cdot (id \vartriangle \mathsf{D}\ (\mathsf{D}_\infty\ x \cdot [\![c]\!]) \cdot c)\ :\ A \leftarrow C$	$a : C \times \mathsf{D}\ (\mathsf{D}_\infty\ A) \rightarrow A$
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Why Mechanising Hylomorphisms in Coq?

- Structured Recursion Schemes have been used in Haskell to structure functional programs, but they do not ensure termination/productivity
- On the other hand, Coq does not capture all recursive definitions
- The benefits of formalising hylos in Coq is three fold:
  - Giving the Coq programmer a **library** where for most recursion schemes they do not have to prove termination properties
  - **Extracting code** into ML/Haskell to provide termination guarantees even in languages with non-termination
  - Using the laws of hylomorphisms as tactics for **program calculation** and **optimisation**

Challenges

1. Avoiding axioms: functional extensionality, heterogeneous equality, ...
2. Extracting "clean" code: close to what a programmer would have written directly in OCaml.
3. Fixed-points of functors, non-termination, etc.

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Our solutions (the remainder of this talk):

1. Machinery for building setoids, use of decidable predicates, . . .
2. Avoiding type families and indexed types.
3. **Containers** & **recursive coalgebras**

Roadmap

Part I: Extractable Containers in Coq
Part II: Recursive Coalgebras & Coq Hylomorphisms
Part III: Code Extraction & Examples

Part I

Extractable Containers in Coq

Setoids and Morphisms

To avoid the functional extensionality axiom, we use:

- **setoids**: types with an associated equivalence
- **proper morphisms** of the respectfulness relation: functions that map related inputs to related outputs

Setoids: Given `setoid A`, and `x y : A`, we write `x =e y : Prop`.

Morphisms: Given `setoid A` and `setoid B`, we write `f : A ~> B`.

Code Extraction for Setoids and Morphisms

We add wrappers on top of Coq's standard Setoids and Proper Morphisms.

Every type must have **exactly one** associated equivalence.

Morphisms are records with a function, and a proof that it respects the relations.

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- We provide automatic coercion to functions.
- Coq's extraction mechanism ignores the `Prop` field.
- We provide a (very basic!) mechanism to help building morphisms.
- We allow the use of Coq's **generalised rewriting** on any morphism or morphism input.

Containers

Containers are defined by a pair $S \lhd P$:

- a type of **shapes** $S : \mathrm{Type}$
- a **family** of positions, indexed by shape $P : S \to \mathrm{Type}$

Containers

Containers are defined by a pair $S \triangleleft P$:

- a type of **shapes** S : Type
- a **family** of positions, indexed by shape $P : S \to$ Type

A **container extension** is a functor defined as follows

$$
\begin{aligned}
[\![S \triangleleft P]\!]\ X &= \Sigma_{s:S} P\ s \to X \\
[\![S \triangleleft P]\!]\ f &= \lambda(s, p).\ (s, f \circ p)
\end{aligned}
$$

Example

Consider the functor $F\ X = 1 + X \times X$

S_F and P_F define a container that is isomorphic to F

$$S_F = 1 + 1 \qquad \begin{array}{l} P_F\ (\mathsf{inl}\ \bullet) = 0 \\ P_F\ (\mathsf{inl}\ \bullet) = 1 + 1 \end{array}$$

Examples of objects of types $F\ \mathbb{N}$ (left) and $[\![S_F \triangleleft P_F]\!]\ \mathbb{N}$ (right):

$$\begin{array}{rcl} \mathsf{inl}\ \bullet & \cong & (\mathsf{inl}\ \bullet, !_{\mathbb{N}}) \\ \mathsf{inr}\ (7, 9) & \cong & (\mathsf{inr}\ \bullet, \lambda x, \mathsf{case}\ x\ \{\ \mathsf{inl}\ \bullet \Rightarrow 7;\ \ \mathsf{inr}\ \bullet \Rightarrow 9\ \}) \end{array}$$

Example

Consider the functor $F\ X = \boxed{1} + \boxed{X \times X}$

Two cases ("shapes")

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No positions on the left shape

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Example

Two positions on the right shape

Consider the functor $F\ X = 1 + X \times X$

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Containers in Coq: A Bad Attempt

Assume a `Shape : Type` and `Pos : Shape -> Type`.

We can define a container extension in the straightforward way:

```
Record App (X : Type) :=
  MkCont { shape : Shape; contents : Pos shape -> X }.
```

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- The above definition forces us to use dependent equality and UIP/Axiom K/... E.g.: dealing with `eq_dep s1 p1 s2 p2` if `p1 : Pos s1` and `p2 : Pos s2`.
- Type families lead to OCaml code with `Obj.magic`.

Extractable Containers in Coq (I)

Solutions:

1. UIP is **not an axiom** in Coq for types with a **decidable equality**.
2. If a type family is defined as a **predicate subtype**, Coq can erase the predicate and extract code that is equivalent to the supertype. E.g. `{x | P x}` for some `P : X -> Prop`.

Extractable Containers in Coq (and II)

Our containers are defined by:

- `Sh : Type`: type of shapes
- `Po : Type`: type of **all** positions
- `valid : Sh * Po ~> bool`

  decidable predicate stating when a pair shape/position is valid

Container extensions that lead to "clean" code extraction:

```
Record App (X : Type)
:= MkCont { shape : Sh;
            contents : {p | valid (shape, p)} -> X
          }.
```

Extractable Containers in Coq (and II)

Our conta

- `Sh :`
- `Po :`
- `vali`

- All proofs of the form `V1 V2 : valid(s,p) = true` are provably equal in Coq to `eq_refl`.
- Given `p1 p2 : {p | valid(s, p)}`, `p1 = p2` iff `proj1_sig p1 = proj1_sig p2`.
- Extraction will treat the contents of container extensions equivalently to `contents : Po -> X`

  (no unsafe coercions).

Container

```
Record App (X : Type)
:= MkCont { shape : Sh;
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          }.
```

Example: $F\ X = 1 + X \times X$

Container definition:

```
Inductive ShapeF := Lbranch | Rbranch.
Inductive PosF := Lpos | Rpos.

Definition validF (x : ShapeF * PosF) : bool
  := match fst x with | Lbranch => false | Rbranch => true end.
```

Example: $F\ X = 1 + X \times X$

Example object equivalent to inr $(7, 8)$

```
Example e1 : App nat :=
  MkCont Rbranch (fun p => match elem p with
                            | Lpos => 7 | Rpos => 8
                            end).
```

The argument of container extensions occurs in strictly positive positions:

We can define least/greatest fixed points of container extensions.

We provide a library of polynomial functors as containers, as well as custom shapes (e.g. binary trees) that we use in our examples.

Not discussed:

- Container morphisms and natural transformations
- Container composition $S \lhd P = (S_1 \lhd P_1) \circ (S_2 \lhd P_2)$
- Container equality

Part II

Recursive Coalgebras & Coq Hylomorphisms

Container Initial Algebras

The least fixed-point of a container extension `App C` is:

```
Inductive LFix C := Lin { lin_op : App C (LFix C) }.
```

Algebras are of type `Alg C X = App C X ~> X`.

Cartamorphisms:

```
cata : Alg C X ~> LFix C ~> X

cata_univ : forall (a : Alg C X) (f : LFix C ~> X),
  f \o Lin =e a \o fmap f <-> f =e cata a
```

Container Terminal Coalgebras

The greatest fixed-point of a container extension `App C` is:

```
CoInductive GFix C := Gin { gin_op : App C (GFix C) }.
```

Coalgebras are of type `CoAlg C X = X ~> App C X`.

Anamorphisms:

```
ana : CoAlg C X ~> X ~> GFix C
```

```
ana_univ : forall (c : CoAlg C X) (f : X ~> GFix C),
  gin_op \o f =e fmap f \o c <-> f =e ana c
```

Recursive Coalgebras (I)

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Recursive coalgebras: coalgebras (`c : CoAlg C X`) that terminate in all inputs.

- i.e. their anamorphisms only produce finite trees.

Recursive Coalgebras (I)

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Recursive coalgebras: coalgebras (`c : CoAlg C X`) that terminate in all inputs.

- i.e. their anamorphisms only produce finite trees.
- i.e. they decompose inputs into "smaller" values of type `X`

Recursive Coalgebras (and II)

We define a predicate `RecF c x` that states that `c : CoAlg C X` terminates on `x : X`.

Using `RecF`, we define:

- Recursive coalgebras:
  ```
  RCoAlg C X = {c | forall x, RecF c x}
  ```
- Given a well-founded relation `R`, well-founded coalgebras
  ```
  WfCoalg C X = {c | forall x p, R x (contents (c x) p)}
  ```

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- Definitions (1) and (2) are equivalent

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- Our mechanisation represents (2) in terms of (1)

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  ```

- Termination proofs may be easier using (1) or (2), depending on the use case

Recursive Hylomorphisms

Recall: hylomorphisms are solutions to the equation $f = a \circ \text{fmap } f \circ c$.

But, due to termination, this solution may not exist, or may not be unique.

However, if c is recursive, then the solution **is unique, and guaranteed to exist**.

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However, if c is recursive, then the solution **is unique, and guaranteed to exist**.

```
Definition hylo_def (a : Alg F B) (c : Coalg F A)
  : forall (x : A), RecF c x -> B :=
  fix f x H :=
    match c x as cx
          return (forall e : Pos (shape cx), RecF c (cont cx e)) -> B
    with
    | MkCont sx cx => fun H => a (MkCont sx (fun e => f (cx e) (H e)))
    end (RecF_inv H).
```

Universal Property of Recursive Hylomorphisms

We define wrappers over `hylo_def`:

```
hylo : Alg C B ~> RCoAlg C A ~> A ~> B
```

From this definition, we can prove the universal property of hylomorphisms. Given `a : Alg C B` and `c : RCoAlg C A`:

```
hylo_univ : forall f : A ~> B,
  f =e a \o fmap f \o c <-> f = hylo a c
```

A Note on Recursive Anamorphisms

For simplicity, we define recursive anamorphisms as `rana c = hylo Lin c`.

- This way we avoid the need to convert `GFix` to `LFix`.
- We prove (straightforward) that `rana c` is equal to `ana c`, followed by converting the result to `LFix`.

Proving the Laws of Hylomorphisms

The following `hylo_fusion` laws are straightforward consequences of `hylo_univ`.

The proofs are exact copies of the pen-and-paper proofs.

```
Lemma hylo_fusion_l
    : f2 \o g1 =e g2 \o fmap f2 ->
          f2 \o hylo g1 h1 =e hylo g2 h1.

Lemma hylo_fusion_r
    : h1 \o f1 =e fmap f1 \o h2 ->
          hylo g1 h1 \o f1 =e hylo g1 h2.

Lemma deforest : cata a \o rana c =e hylo a c.
```

Part III

Code Extraction & Examples

A Tree Container for Divide &Conquer

Quicksort Definition

Quicksort Extraction

Using Hylo-fusion for Program Optimisation

Optimized Code Extraction

Dynamorphisms

Knapsack

Knapsack Extraction

Wrap-up