

# Lecture 7 – First-Class Functions

## COSE212: Programming Languages

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2023 Fall

- **F1VAE – VAE with first-order functions**
  - Concrete and Abstract Syntax
  - Evaluation with Function Environments
  - Interpreters and Natural Semantics
  - Static Scoping vs Dynamic Scoping

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- **FVAE – VAE with first-class functions**
  - Concrete and Abstract Syntax
  - Interpreter and Natural Semantics

## 1. First-Class Functions

## 2. FVAE – VAE with First-Class Functions

- Concrete Syntax

- Abstract Syntax

## 3. Interpreter and Natural Semantics for F1VAE

- Closures – Functions as Values

- Addition and Multiplication

- Anonymous Functions

- Function Application

- Function Application (Dynamic Scoping)

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- ① **assigned** to a **variable**,
- ② **passed** as an **argument** to a function, and
- ③ **returned** from a function.



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- ① **assigned** to a **variable**,
- ② **passed** as an **argument** to a function, and
- ③ **returned** from a function.

For example, an integer is obviously a first-class citizen in Scala:

```
// 1. We can assign an integer to a variable.
val n: Int = 3

// 2. We can pass an integer as an argument to a function.
def square(n: Int): Int = n * n
square(3)           // 3 * 3 = 9

// 3. We can return an integer from a function.
def square(n: Int): Int = n * n
square(3)           // 3 * 3 = 9
```

In Scala, **functions** are also **first-class citizens**, and we call them **first-class functions**.

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```
def inc(n: Int): Int = n + 1

// 1. We can assign a function to a variable.
val f: Int => Int = inc
val g: Int => Int = x => x + 1           // anonymous (lambda) function

// 2. We can pass a function as an argument to a function.
def twice(f: Int => Int, n: Int): Int = f(f(n))
twice(inc, 3)                         // inc(inc(3)) = 3 + 1 + 1 = 5
List(1, 2, 3).map(inc)                 // List(2, 3, 4)

// 3. We can return a function from a function.
def addN(n: Int): Int => Int = m => n + m
val add3: Int => Int = addN(3)
add3(5)                               // 3 + 5 = 8
addN(3)(5)                            // 3 + 5 = 8 -- currying
```

In **functional programming**, functions are **first-class functions**.

- Scala

```
List(1, 2, 3).map(x => x * 2)           // List(2, 4, 6)
```

- Python

```
list(map(lambda x: x * 2, [1, 2, 3]))  # [2, 4, 6]
```

- Rust

```
[1, 2, 3].iter().map(|x| x * 2).collect() // [2, 4, 6]
```

- Haskell

```
map (\x -> x * 2) [1, 2, 3]           -- [2, 4, 6]
```

- ...

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Function Application (Dynamic Scoping)

Now, we want to extend VAE into FVAE with **first-class functions** rather than **first-order functions** in F1VAE.

```
/* FVAE */  
val addN = n => m => n + m;  
val add3 = addN(3);  
add3(5) // 3 + 5 = 8
```

```
/* FVAE */  
val inc = x => x + 1;  
val twice = f => n => f(f(n));  
twice(inc)(5) // 5 + 1 + 1 = 7
```

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val inc = x => x + 1;  
val twice = f => n => f(f(n));  
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```

- For FVAE, we need to extend **expressions** with
  - 1 anonymous (lambda) functions
  - 2 function applications

Let's define the **concrete syntax** of FVAE in BNF:

- For FVAE, we need to extend **expressions** with
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- For FVAE, we need to extend **expressions** with
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```
// expressions
<expr> ::= ...
        | "{" <expr> "}"
        | "val" <id> "=" <expr> ";" <expr>
        | <id>
        | <id> "=>" <expr>
        | <expr> "(" <expr> ")"
```

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- For FVAE, we need to extend **expressions** with
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Why not the following function application syntax?

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| <id> "(" <expr> ")"
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Let's define the **concrete syntax** of FVAE in BNF:

- For FVAE, we need to extend **expressions** with
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// expressions
<expr> ::= ...
        | "{" <expr> "}"
        | "val" <id> "=" <expr> ";" <expr>
        | <id>
        | <id> "=>" <expr>
        | <expr> "(" <expr> ")"
```

Why not the following function application syntax?

```
| <id> "(" <expr> ")"
```

We cannot support curried function application (e.g., `twice(inc)(5)`).

Let's define the **abstract syntax** of FVAE in BNF:

Expressions  $\mathbb{E} \ni e ::= \dots$

	val $x=e$ ; $e$	(Val)
	$x$	(Id)
	$\lambda x.e$	(Fun)
	$e(e)$	(App)

Let's define the **abstract syntax** of FVAE in BNF:

Expressions $\mathbb{E} \ni e$	$::=$	...
		val x=e; e (Val)
		x (Id)
		$\lambda x.e$ (Fun)
		$e(e)$ (App)

```
enum Expr:
  ...
  case Val(name: String, init: Expr, body: Expr)
  case Id(name: String)
  // anonymous (lambda) functions
  case Fun(param: String, body: Expr)
  // function applications
  case App(fun: Expr, arg: Expr)
```

For example, let's **parse** the following FVAE program:

```
/* FVAE */  
val addN = n => m => n + m;  
val add3 = addN(3);  
add3(5)
```

Then, the following **abstract syntax tree (AST)** is produced:

```
Val("addN",  
  Fun("n",  
    Fun("m",  
      Add(Id("n"), Id("m"))  
    )  
  ),  
  Val("add3",  
    App(Id("addN"), Const(3)),  
    App(Id("add3"), Const(5))  
  )  
)
```

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Function Application (Dynamic Scoping)

Let's evaluate the following FVAE program:

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How to evaluate the expression `addN(3)` and `add3(5)`?

$[\text{addN} \mapsto ???] \vdash \text{addN}(3) \Rightarrow ???$

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Environments map variables to values. What's values of `addN` and `add3`?

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**Functions!**

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**Functions!** Let's define **values** as either **numbers** or **functions**:

Values  $\forall \ni v ::= n \mid \lambda x.e$

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**Functions!** Let's define **values** as either **numbers** or **functions**:

Values  $\forall \ni v ::= n \mid \lambda x.e$

However, it is **NOT** what exactly we want to do. Why?

```
/* FVAE */  
val addN = n => m => n + m;  
val add3 = addN(3);  
add3(5) // 3 + 5 = 8
```

$[\text{addN} \mapsto \lambda n. \lambda m. (n + m)] \vdash \text{addN}(3) \Rightarrow \lambda m. (n + m)$

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We know that  $m$  represents 5, but what about  $n$ ?

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For **static scoping**, we need to remember the environment when the function is defined as values.



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For **static scoping**, we need to remember the environment when the function is defined as values.

Let's define function values as **closures**, a pair of 1) a **function** and 2) its **environment**:

Values  $\mathbb{V} \ni v ::= n \mid \langle \lambda x. e, \sigma \rangle$

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Let's define function values as **closures**, a pair of 1) a **function** and 2) its **environment**:

Values  $\mathbb{V} \ni v ::= n \mid \langle \lambda x. e, \sigma \rangle$

Then, it works!

$$[\text{addN} \mapsto \langle \lambda n. \lambda m. (n + m), \emptyset \rangle] \vdash \text{addN}(3) \Rightarrow \langle \lambda m. (n + m), [n \mapsto 3] \rangle$$

$$[\text{add3} \mapsto \langle \lambda m. (n + m), [n \mapsto 3] \rangle] \vdash \text{add3}(5) \Rightarrow 8$$

```
// values
enum Value:
  case NumV(n: BigInt)
  case CloV(param: String, body: Expr, env: Env)
```

$$\begin{array}{ll} \text{Values } \mathbb{V} \ni v ::= n & (\text{NumV}) \\ \quad \mid \langle \lambda x.e, \sigma \rangle & (\text{CloV}) \end{array}$$

For FVAE, a **value** is either 1) a **number**  $n$  or a **closure**  $\langle \lambda x.e, \sigma \rangle$ ,

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For FVAE, a **value** is either 1) a **number**  $n$  or a **closure**  $\langle \lambda x. e, \sigma \rangle$ ,

```
// interpreter
def interp(expr: Expr, env: Env): Value = ???
```

$$\sigma \vdash e \Rightarrow v$$

and the interpreter takes an **expression**  $e$  with an **environment**  $\sigma$  and returns a **value**  $v$  (either a number or a closure):

For FVAE, we need to 1) implement the **interpreter**:

```
def interp(expr: Expr, env: Env): Value = ???
```

For FVAE, we need to 1) implement the **interpreter**:

```
def interp(expr: Expr, env: Env): Value = ???
```

and 2) define the **natural semantics** with environments:

$$\sigma \vdash e \Rightarrow v$$

Expressions	$\mathbb{E} \ni e$	$::=$	...
			$\lambda x.e$ (Fun) $e(e)$ (App)
Values	$\mathbb{V} \ni v$	$::=$	$n$ (NumV) $\langle \lambda x.e, \sigma \rangle$ (CloV)

where

Environments	$\sigma \in \mathbb{X} \xrightarrow{\text{fin}} \mathbb{V}$	(Env)
Integers	$n \in \mathbb{Z}$	(BigInt)
Identifiers	$x \in \mathbb{X}$	(String)

```
def interp(expr: Expr, env: Env): Value = expr match
  ...

  case Add(l, r) => interp(l, env) + interp(r, env)

  case Mul(l, r) => interp(l, env) * interp(r, env)
```

$$\sigma \vdash e \Rightarrow v$$

$$\text{Add} \frac{\sigma \vdash e_1 \Rightarrow v_1 \quad \sigma \vdash e_2 \Rightarrow v_2}{\sigma \vdash e_1 + e_2 \Rightarrow v_1 + v_2}$$

$$\text{Mul} \frac{\sigma \vdash e_1 \Rightarrow v_1 \quad \sigma \vdash e_2 \Rightarrow v_2}{\sigma \vdash e_1 \times e_2 \Rightarrow v_1 \times v_2}$$

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Is it correct?



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def interp(expr: Expr, env: Env): Value = expr match
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Is it correct? **No!**

We can only add or multiply **numbers** rather than arbitrary **values**.

```
def interp(expr: Expr, env: Env): Value = expr match
  ...

  case Add(l, r) => (interp(l, env), interp(r, env)) match
    case (NumV(l), NumV(r)) => NumV(l + r)
    case (l, r) => error(s"invalid operation: ${l.str} + ${r.str}")

  case Mul(l, r) => (interp(l, env), interp(r, env)) match
    case (NumV(l), NumV(r)) => NumV(l * r)
    case (l, r) => error(s"invalid operation: ${l.str} * ${r.str}")
```

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def interp(expr: Expr, env: Env): Value = expr match
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    case (l, r) => error(s"invalid operation: ${l.str} + ${r.str}")

  case Mul(l, r) => (interp(l, env), interp(r, env)) match
    case (NumV(l), NumV(r)) => NumV(l * r)
    case (l, r) => error(s"invalid operation: ${l.str} * ${r.str}")
```

$$\sigma \vdash e \Rightarrow v$$

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$$\text{Mul} \frac{\sigma \vdash e_1 \Rightarrow n_1 \quad \sigma \vdash e_2 \Rightarrow n_2}{\sigma \vdash e_1 \times e_2 \Rightarrow n_1 \times n_2}$$

Let's refactor the code to avoid duplication using a helper function.

```
def numBOp(name: String)(
  op: (BigInt, BigInt) => BigInt
)(l: Value, r: Value): Value = (l, r) match
  case (NumV(l), NumV(r)) => NumV(op(l, r))
  case (l, r) => error(s"invalid operation: ${l.str} $name ${r.str}")

val numAdd: (Value, Value) => Value = numBOp("+")(_ + _)
val numMul: (Value, Value) => Value = numBOp("*")(_ * _)

def interp(expr: Expr, env: Env): Value = expr match
  ...
  case Add(l, r) => numAdd(interp(l, env), interp(r, env))
  case Mul(l, r) => numMul(interp(l, env), interp(r, env))
```

$$\sigma \vdash e \Rightarrow v$$

$$\text{Add} \frac{\sigma \vdash e_1 \Rightarrow n_1 \quad \sigma \vdash e_2 \Rightarrow n_2}{\sigma \vdash e_1 + e_2 \Rightarrow n_1 + n_2}$$

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```
def interp(expr: Expr, env: Env): Value = expr match
  ...
  case Fun(p, b) => ???
```

$$\sigma \vdash e \Rightarrow v$$

$$\text{Fun} \frac{\text{???}}{\sigma \vdash \lambda x. e \Rightarrow \text{???}}$$

```
def interp(expr: Expr, env: Env): Value = expr match
  ...
  case Fun(p, b) => CloV(p, b, env)
```

$$\sigma \vdash e \Rightarrow v$$

$$\text{Fun} \frac{}{\sigma \vdash \lambda x.e \Rightarrow \langle \lambda x.e, \sigma \rangle}$$

Construct a closure  $\langle \lambda x.e, \sigma \rangle$  from the function  $\lambda x.e$  with the current environment  $\sigma$  for **static scoping**.

```
def interp(expr: Expr, env: Env): Value = expr match
  ...
  case App(f, e) => ???
```

$$\sigma \vdash e \Rightarrow v$$

$$\text{App} \frac{\text{???}}{\sigma \vdash e_0(e_1) \Rightarrow \text{???}}$$

```
def interp(expr: Expr, env: Env): Value = expr match
...
case App(f, e) => interp(f, env) match
  case CloV(p, b, fenv) => ...
  case v                => error(s"not a function: ${v.str}")
```

$$\sigma \vdash e \Rightarrow v$$

$$\text{App} \frac{\sigma \vdash e_0 \Rightarrow \langle \lambda x. e_2, \sigma' \rangle \quad \dots}{\sigma \vdash e_0(e_1) \Rightarrow ???}$$

First, evaluate the function expression  $e_0$ , check that it is a closure, and let  $\langle \lambda x. e_2, \sigma' \rangle$  be the closure.



```
def interp(expr: Expr, env: Env): Value = expr match
...
case App(f, e) => interp(f, env) match
  case CloV(p, b, fenv) => ... interp(e, env) ...
  case v                => error(s"not a function: ${v.str}")
```

$$\sigma \vdash e \Rightarrow v$$

$$\text{App} \frac{\sigma \vdash e_0 \Rightarrow \langle \lambda x. e_2, \sigma' \rangle \quad \sigma \vdash e_1 \Rightarrow v_1 \quad \dots}{\sigma \vdash e_0(e_1) \Rightarrow ???}$$

Then, evaluate the argument expression  $e_1$  and let  $v_1$  be the resulting value.

```
def interp(expr: Expr, env: Env): Value = expr match
  ...
  case App(f, e) => interp(f, env) match
    case CloV(p, b, fenv) => interp(b, fenv + (p -> interp(e, env)))
    case v                  => error(s"not a function: ${v.str}")
```

$$\sigma \vdash e \Rightarrow v$$

$$\text{App} \frac{\sigma \vdash e_0 \Rightarrow \langle \lambda x. e_2, \sigma' \rangle \quad \sigma \vdash e_1 \Rightarrow v_1 \quad \sigma'[x \mapsto v_1] \vdash e_2 \Rightarrow v_2}{\sigma \vdash e_0(e_1) \Rightarrow v_2}$$

Finally, evaluate the body expression  $e_2$  in the environment  $\sigma'[x \mapsto v_1]$ , where  $\sigma'$  denotes the environment captured at the **definition** site of the function for **static scoping**.

```
def interp(expr: Expr, env: Env): Value = expr match
...
case App(f, e) => interp(f, env) match
  case CloV(p, b, fenv) => interp(b, env + (p -> interp(e, env)))
  case v                => error(s"not a function: ${v.str}")
```

$$\sigma \vdash e \Rightarrow v$$

$$\text{App} \frac{\sigma \vdash e_0 \Rightarrow \langle \lambda x. e_2, \sigma' \rangle \quad \sigma \vdash e_1 \Rightarrow v_1 \quad \sigma[x \mapsto v_1] \vdash e_2 \Rightarrow v_2}{\sigma \vdash e_0(e_1) \Rightarrow v_2}$$

We can define **dynamic scoping** by using the current environment  $\sigma$  instead of the environment  $\sigma'$  captured at the definition site of the function.

```
/* FVAE (static scoping) */
val x = 3;
val f = y => x * y;
val x = 4;
f(5) // 3 * 5 = 15
```

```
/* FVAE (dynamic scoping) */
val x = 3;
val f = y => x * y;
val x = 4;
f(5) // 4 * 5 = 20
```

$$\text{APP} \frac{\sigma_1 \vdash f \Rightarrow \langle \lambda y. (x \times y), \sigma_0 \rangle \quad \sigma_1 \vdash 5 \Rightarrow 5 \quad \sigma_0[y \mapsto 5] \vdash x \times y \Rightarrow 15}{\sigma_1 \vdash f(5) \Rightarrow 15}$$

$$\text{APP} \frac{\sigma_1 \vdash f \Rightarrow \langle \lambda y. (x \times y), \sigma_0 \rangle \quad \sigma_1 \vdash 5 \Rightarrow 5 \quad \sigma_1[y \mapsto 5] \vdash x \times y \Rightarrow 20}{\sigma_1 \vdash f(5) \Rightarrow 20}$$

where

$$\begin{aligned} \sigma_0 &= [x \mapsto 3] \\ \sigma_1 &= [x \mapsto 4, f \mapsto \langle \lambda y. (x \times y), \sigma_0 \rangle] \end{aligned}$$

## 1. First-Class Functions

## 2. FVAE – VAE with First-Class Functions

- Concrete Syntax

- Abstract Syntax

## 3. Interpreter and Natural Semantics for F1VAE

- Closures – Functions as Values

- Addition and Multiplication

- Anonymous Functions

- Function Application

- Function Application (Dynamic Scoping)

- Please see this document<sup>1</sup> on GitHub.
  - Implement `interp` function.
  - Implement `interpDS` function.
- It is just an exercise, and you **don't need to submit** anything.
- However, some exam questions might be related to this exercise.

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<sup>1</sup><https://github.com/ku-plrg-classroom/docs/tree/main/cose212/fvae>.

- Lambda Calculus

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