Lecture 16 – First-Class Continuations
COSE212: Programming Languages

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⚫PLRG

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- We will learn about **continuations** with the following topics:
  - **Continuations** (Lecture 14 & 15)
  - **First-Class Continuations** (Lecture 16)
  - **Compiling with continuations** (Lecture 17)

- A **continuation** represents the **rest of the computation**.
  - Continuation Passing Style (CPS)
  - Interpreter of FAE in CPS
  - Small-step operational (reduction) semantics of FAE

- In this lecture, let's learn **first-class continuations**.

- **KFAE** – FAE with **first-class continuations**
  - Interpreter and Reduction semantics

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    Interpreter and Reduction Semantics for KFAE
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4. Control Statements

In a programming language, an entity is said to be **first-class citizen** if it is treated as a **value**. In other words, it can be

1 **assigned** to a **variable**,

2 **passed** as an **argument** to a function, and

3 **returned** from a function.

For example, Scala supports **first-class functions**.

```scala
def inc(n: Int): Int = n + 1
// 1. We can assign a function to a variable.
val f: Int => Int =
// 2. We can pass a function as an argument to a function.
List(1, 2, 3).map(inc)                    // List(2, 3, 4)
// 3. We can return a function from a function.
def addN(n: Int): Int => Int = m => n + m
val add3: Int => Int = addN(3)
add3(5)                                   // 3 + 5 = 8
```

First-class Continuations

Similarly, a **first-class continuation** means that a continuation is treated as a **value**.

For example, **Racket** supports `let`/`cc` to create a first-class continuation.

```
; prefix notation              ; infix notation
(* 2 (+ 3 5))                  ; 2 * (3 + 5) = 16
```

❶ Evaluate 2.   (Result: 2)

❷ Evaluate 3.   (Result: 3)

❸ Evaluate 5.   (Result: 5)

❹ Add the results of step ❷ and ❸.   (Result: 3 + 5 = 8)

❺ Multiply the results of step ❶ and ❷ – ❹.   (Result: 2 * 8 = 16)

```
; first-class continuation with `let/cc`
(* 2 (let/cc k (+ 3 (k 5))))            ; 2 * 5 = 10
```

For example, Racket supports let/cc to create a first-class continuation.

```
; first-class continuation with `let/cc`
(* 2 (let/cc k (+ 3 (k 5))))          ; 2 * 5 = 10
```

① Evaluate 2.   (Result: 2)

② Let k be the continuation of ② – ⑥.   (k is x => 2 * x)

③ Evaluate 3.   (Result: 3)

④ Evaluate 5.   (Result: 5)

⑤ Apply the result of step ④ to k.   (Result of ② – ⑥: 5)

⑥ Add the results of step ③ and ④ – ⑤.   (Unreachable)

⑦ Multiply the results of step ① and ② – ⑥.   (Result: 2 * 5 = 10)

- The step ② defines the continuation of ② – ⑥ as a value in k.
- The step ⑤ changes the control flow using k with the result of ④.

First-class Continuations

Some functional languages support **first-class continuations**.

- Rakcet

  ```
  (* 2 (let/cc k (+ 3 (k 5))))          ; 2 * 5 = 10
  ```

- Ruby

  ```
  2 * (callcc { |k| 3 + k.call(5)})     # 2 * 5 = 10
  ```

- Haskell

  ```
  do
    x <- callCC $ \k -> do
      y <- k 5
      return $ 3 + y
    return $ 2 * x                      -- 2 * 5 = 10
  ```

- ...

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Now, let's extend FAE into KFAE to support **first-class continuations**.
(Assume that val is supported in FAE as syntactic sugar.)

```
/* KFAE */
2 * { vcc k; 3 + k(5) } // 2 * 5 = 10
```

```
/* KFAE */
{
  vcc done;        // done = x => x
  val f = {
    vcc exit;      // exit = x => val f = x; f(3) * 5
    2 * done(1 + {
      vcc k;       // k   = x => val f = { 2 * done(1 + x) }; f(3) * 5
      exit(k)
    })
  };
  f(3) * 5
}
// 1 + 3
```

For KFAE, we need to extend **expressions** of FAE with

1 first-class continuations (vcc)

We can extend the **concrete syntax** of FAE as follows:

```
// expressions
<expr> ::= ... | "vcc" <id> ";" <expr>
```

and the **abstract syntax** of FAE as follows:

$$\text{Expressions} \quad \mathbb{E} \ni e ::= \ldots \mid \text{vcc } x;\ e \quad (\text{Vcc})$$

```
enum Expr:
  ...
  // first-class continuations
  case Vcc(name: String, body: Expr)
```

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In the previous lecture, we have defined the **first-order representation** of
continuations with **value stack**:

```
enum Cont:
  case EmptyK
  case EvalK(env: Env, expr: Expr, k: Cont)
  case AddK(k: Cont)
  case MulK(k: Cont)
  case AppK(k: Cont)

type Stack = List[Value]
```

$$
\begin{array}{llll}
\text{Continuations} & \mathbb{K} \ni \kappa ::= & \square & (\texttt{EmptyK}) \\
& & | \ (\sigma \vdash e) :: \kappa & (\texttt{EvalK}) \\
& & | \ (+) :: \kappa & (\texttt{AddK}) \\
& & | \ (\times) :: \kappa & (\texttt{MulK}) \\
& & | \ (@) :: \kappa & (\texttt{AppK}) \\
\text{Value Stacks} & \mathbb{S} \ni s ::= & \blacksquare \ | \ v :: s & (\texttt{List[Value]})
\end{array}
$$

Then, we have defined the **reduction relation** $\rightarrow \in (\mathbb{K} \times \mathbb{S}) \times (\mathbb{K} \times \mathbb{S})$ between **states** consisting of pairs of **continuations** and **value stacks**:

```
def reduce(k: Cont, s: Stack): (Cont, Stack) = ???
```

$$\langle \kappa \mid\mid s \rangle \rightarrow \langle \kappa' \mid\mid s' \rangle$$

And the eval function **iteratively reduces** the state until it reaches the empty continuation $\square$ and returns the single value in the value stack:

```
def eval(str: String): String =
  import Cont.*
  def aux(k: Cont, s: Stack): Value = reduce(k, s) match
    case (EmptyK, List(v)) => v
    case (k, s) => aux(k, s)
  aux(EvalK(Map.empty, Expr(str), EmptyK), List.empty).str
```

$$\langle (\varnothing \vdash e) :: \square \mid\mid \blacksquare \rangle \rightarrow^* \langle \square \mid\mid v :: \blacksquare \rangle$$

Now, let's extend the interpreter and reduction semantics for FAE to KFAE by adding the **first-class continuations**.

First, we need to extend the values of FAE with **continuation values** consisting of pairs of continuations and value stacks:

```
// values
enum Value:
  case NumV(number: BigInt)
  case CloV(param: String, body: Expr, env: Env)
  case ContV(cont: Cont, stack: Stack)
```

$$\text{Values} \quad \mathbb{V} \ni v ::= n \qquad\qquad (\texttt{NumV})$$
$$| \; \langle \lambda x.e, \sigma \rangle \quad (\texttt{CloV})$$
$$| \; \langle \kappa \parallel s \rangle \qquad (\texttt{ContV})$$

Then, let's fill out the missing cases in the `reduce` function and reduction rules for $\rightarrow$ in the reduction semantics of KFAE.

```
def reduce(k: Cont, s: Stack): (Cont, Stack) = (k, s) match
  case (EvalK(env, expr, k), s) => expr match
    ...
    case Vcc(x, b) => (EvalK(env + (x -> ContV(k, s)), b, k), s)
```

$$\langle \kappa \mid\mid s \rangle \to \langle \kappa \mid\mid s \rangle$$

$$\text{Vcc} \quad \langle (\sigma \vdash \texttt{vcc } x; \; e) :: \kappa \mid\mid s \rangle \;\; \to \;\; \langle (\sigma[x \mapsto \langle \kappa \mid\mid s \rangle] \vdash e) :: \kappa \mid\mid s \rangle$$

It defines a new immutable binding x in the environment σ that maps to a **continuation value** $\langle \kappa \mid\mid s \rangle$, and then evaluates the body expression e in the extended environment $\sigma[x \mapsto \langle \kappa \mid\mid s \rangle]$.

```scala
def reduce(k: Cont, s: Stack): (Cont, Stack) = (k, s) match
  case (EvalK(env, expr, k), s) => expr match
    ...
    case App(f, e) => (EvalK(env, f, EvalK(env, e, AppK(k))), s)
  ...
  case (AppK(k), a :: f :: s) => f match
    case CloV(p, b, fenv) => (EvalK(fenv + (p -> a), b, k), s)
    case ContV(k1, s1)    => (k1, a :: s1)
    case v                => error(s"not a function: ${v.str}")
```

$$\boxed{\langle \kappa \ || \ s \rangle \to \langle \kappa \ || \ s \rangle}$$

$$
\begin{array}{lll}
\texttt{App}_1 & \langle (\sigma \vdash e_1(e_2)) :: \kappa \ || \ s \rangle & \to & \langle (\sigma \vdash e_1) :: (\sigma \vdash e_2) :: (@) :: \kappa \ || \ s \rangle \\
\texttt{App}_{2,\lambda} & \langle (@) :: \kappa \ || \ v_2 :: \langle \lambda x.e, \sigma \rangle :: s \rangle & \to & \langle (\sigma[x \mapsto v_2] \vdash e) :: \kappa \ || \ s \rangle \\
\texttt{App}_{2,\kappa} & \langle (@) :: \kappa \ || \ v_2 :: \langle \kappa' \ || \ s' \rangle :: s \rangle & \to & \langle \kappa' \ || \ v_2 :: s' \rangle \\
\end{array}
$$

The new $\texttt{App}_{2,\kappa}$ rule handles when the function expression evaluates to a continuation value $\langle \kappa' \ || \ s' \rangle$. It changes the control flow to the continuation κ' with the given argument value v_2 and the value stack s'.

Example 1

Let's interpret the expression $2 \times (\text{vcc } k; \ (3 + k(5)))$:

$$
\begin{array}{lll}
& \langle \ (\varnothing \vdash 2 \times (\text{vcc } k; \ (3 + k(5)))) :: \square & || \ \blacksquare & \rangle \\
(\text{Mul}_1) \\
\rightarrow & \langle \ (\varnothing \vdash 2) :: (\varnothing \vdash (\text{vcc } k; \ (3 + k(5)))) :: (\times) :: \square & || \ \blacksquare & \rangle \\
(\text{Num}) \\
\rightarrow & \langle \ (\varnothing \vdash (\text{vcc } k; \ (3 + k(5)))) :: (\times) :: \square & || \ 2 :: \blacksquare & \rangle \\
(\text{Vcc}) \\
\rightarrow & \langle \ (\sigma_0 \vdash (3 + k(5))) :: (\times) :: \square & || \ 2 :: \blacksquare & \rangle \\
(\text{Add}_1) \\
\rightarrow & \langle \ (\sigma_0 \vdash 3) :: (\sigma_0 \vdash k(5)) :: (+) :: (\times) :: \square & || \ 2 :: \blacksquare & \rangle \\
(\text{Num}) \\
\rightarrow & \langle \ (\sigma_0 \vdash k(5)) :: (+) :: (\times) :: \square & || \ 3 :: 2 :: \blacksquare & \rangle \\
(\text{App}_1) \\
\rightarrow & \langle \ (\sigma_0 \vdash k) :: (\sigma_0 \vdash 5) :: (@) :: (+) :: (\times) :: \square & || \ 3 :: 2 :: \blacksquare & \rangle \\
(\text{Id}) \\
\rightarrow & \langle \ (\sigma_0 \vdash 5) :: (@) :: (+) :: (\times) :: \square & || \ \langle \kappa_0 \ || \ s_1 \rangle :: 3 :: 2 :: \blacksquare & \rangle \\
(\text{Num}) \\
\rightarrow & \langle \ (@) :: (+) :: (\times) :: \square & || \ 5 :: \langle \kappa_0 \ || \ s_1 \rangle :: 3 :: 2 :: \blacksquare & \rangle \\
(\text{App}_{2,\kappa}) \\
\rightarrow & \langle \ (\times) :: \square & || \ 5 :: 2 :: \blacksquare & \rangle \\
(\text{Mul}_2) \\
\rightarrow & \langle \ \square & || \ 10 :: \blacksquare & \rangle
\end{array}
$$

$$
\text{where } \begin{cases} \sigma_0 = [k \mapsto \langle \kappa_0 \ || \ s_0 \rangle] \\ \kappa_0 = (\times) :: \square \\ s_0 = 2 :: \blacksquare \end{cases}
$$

Example 2

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Let's interpret the expression $(\lambda x.(\text{vcc } r;\ r(x+1) \times 2))(3)$:

$$
\begin{array}{ll}
(\text{App}_1) & \langle\ (\varnothing \vdash (\lambda x.(\text{vcc } r;\ r(x+1) \times 2))(3)) :: \square & || \ \blacksquare \ \rangle \\
\xrightarrow{} & \\
(\text{Fun}) & \langle\ (\varnothing \vdash (\lambda x.(\text{vcc } r;\ r(x+1) \times 2))) :: (\varnothing \vdash 3) :: (@) :: \square\ || \ \blacksquare \ \rangle \\
\xrightarrow{} & \\
(\text{Num}) & \langle\ (\varnothing \vdash 3) :: (@) :: \square & || \ \langle \lambda x.e_0, \varnothing \rangle :: \blacksquare \ \rangle \\
\xrightarrow{} & \\
(\text{App}_{2,\lambda}) & \langle\ (@) :: \square & || \ 3 :: \langle \lambda x.e_0, \varnothing \rangle :: \blacksquare \ \rangle \\
\xrightarrow{} & \\
(\text{Vcc}) & \langle\ (\sigma_0 \vdash \text{vcc } r;\ r(x+1) \times 2) :: \square & || \ \blacksquare \ \rangle \\
\xrightarrow{} & \\
(\text{Mul}_1) & \langle\ (\sigma_1 \vdash r(x+1) \times 2) :: \square & || \ \blacksquare \ \rangle \\
\xrightarrow{} & \\
(\text{App}_1) & \langle\ (\sigma_1 \vdash r(x+1)) :: (\sigma_1 \vdash 2) :: (@) :: (\times) :: \square & || \ \blacksquare \ \rangle \\
\xrightarrow{} & \\
 & \langle\ (\sigma_1 \vdash r) :: (\sigma_1 \vdash x+1) :: (@) :: (\sigma_1 \vdash 2) :: (\times) :: \square\ || \ \blacksquare \ \rangle \\
 & \cdots \\
\xrightarrow{*} & \langle\ (@) :: (\sigma_1 \vdash 2) :: (\times) :: \square & || \ 4 :: \langle \square\ ||\ \blacksquare \rangle :: \blacksquare \ \rangle \\
(\text{App}_{2,\kappa}) & \\
\xrightarrow{} & \langle\ \square & || \ 4 :: \blacksquare \ \rangle
\end{array}
$$

$$
\text{where} \begin{cases}
e_0 & = \ \text{vcc } r;\ r(x+1) \times 2 \\
\sigma_0 & = \ [x \mapsto 3] \\
\sigma_1 & = \ \sigma_0[r \mapsto \langle \square\ ||\ \blacksquare \rangle]
\end{cases}
$$

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Control Statements

Many real-world programming languages support **control statements** to change the **control-flow** of a program.

For example, C++ supports break, continue, and return statements:

```cpp
int sumEvenUntilZero(int xs[], int len) {
  if (len <= 0) return 0;          // directly return 0 if len <= 0
  int sum = 0;
  for (int i = 0; i < len; i++) {
    if (xs[i] == 0) break;         // stop the loop if xs[i] == 0
    if (xs[i] % 2 == 1) continue;  // skip the rest if xs[i] is odd
    sum += xs[i];
  }
  return sum;                      // finally return the sum
}
int xs[] = {4, 1, 3, 2, 0, 6, 5, 8};
sumEvenUntilZero(xs, 8);          // 4 + 2 = 6
```

Let's represent them using **first-class continuations**!

Control Statements

- `return` statement:

```
x => body
```

means

```
x => { vcc return;
  body // return(e) directly returns e to the caller
}
```

- `break` and `continue` statements:

```
while (cond) body
```

means

```
{ vcc break;
  while (cond) { vcc continue;
    body // continue(e)/break(e) jumps to the next/end of the loop
  }
}
```

Control Statements

We can represent other control statements similarly, but think for yourself!

- exception in Python

```python
try:
    x = y / z
except ZeroDivisionError:
    x = 0
```

- generator in JavaScript

```javascript
const foo = function* () { yield 'a'; yield 'b'; yield 'c'; };
let str = '';
for (const c of foo()) { str = str + c; }
str // 'abc'
```

- coroutines in Kotlin
- async/await in C#
- ...

- Please see this document[1] on GitHub.
  - Implement reduce function.
- It is just an exercise, and you **don't need to submit** anything.
- However, some exam questions might be related to this exercise.

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[1]https://github.com/ku-plrg-classroom/docs/tree/main/cose212/kfae.

- Compiling with Continuations

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