Lecture 4 – Identifiers (1)

COSE212: Programming Languages

Jihyeok Park



2023 Fall





• ADT for Abstract Syntax of AE

```
enum Expr:
    case Num(number: BigInt)
    case Add(left: Expr, right: Expr)
    case Mul(left: Expr, right: Expr)
```

• Parser for Concrete Syntax of AE

```
lazy val expr: P[Expr] = ...
```

• Interpreter for Semantics of AE

```
def interp(expr: Expr): Value = ...
```





• ADT for Abstract Syntax of AE

```
enum Expr:
   case Num(number: BigInt)
   case Add(left: Expr, right: Expr)
   case Mul(left: Expr, right: Expr)
```

• Parser for Concrete Syntax of AE

```
lazy val expr: P[Expr] = ...
```

• Interpreter for Semantics of AE

```
def interp(expr: Expr): Value = ...
```

• In this lecture, we will learn identifiers.

Contents



1. Identifiers

Bound Identifiers Free Identifiers Shadowing

2. VAE - AE with Variables

Concrete Syntax Abstract Syntax Examples

Contents



1. Identifiers

Bound Identifiers Free Identifiers Shadowing

2. VAE – AE with Variables Concrete Syntax Abstract Syntax Examples

Identifiers



An **identifier** is a **name** for a certain element in a program.

In Scala, there are diverse kinds of identifiers:

- variable names
- function names
- parameter names
- class names
- field names
- etc.

```
val x = 42
def add(a: Int, b: Int): Int = a + b
case class Person(name: String, age: Int)
...
```

Bound Identifiers



```
val x = 3
def add(a: Int, b: Int): Int = {
  val x = a + b
  x + y
}
val y = x + 1
add(y, z)
```

A **bound identifier** is an identifier that is **defined** in a program.

- A binding occurrence of an identifier is the occurrence in its definition.
- A **scope** of an identifier is a code region where the identifier is usable.
- A bound occurrence of an identifier is an occurrence in its scope.

Bound Identifiers



```
val x = 3
def add(a: Int, b: Int): Int = {
  val x = a + b
  x + y
}
val y = x + 1
add(y, z)
```

A **bound identifier** is an identifier that is **defined** in a program.

- A binding occurrence of an identifier is the occurrence in its definition.
- A **scope** of an identifier is a code region where the identifier is usable.
- A **bound occurrence** of an identifier is an occurrence in its scope.

Bound Identifiers



```
val x = 3
def add(a: Int, b: Int): Int = {
  val x = a + b
  x + y
}
val y = x + 1
add(y, z)
```

A **bound identifier** is an identifier that is **defined** in a program.

- A binding occurrence of an identifier is the occurrence in its definition.
- A **scope** of an identifier is a code region where the identifier is usable.
- A bound occurrence of an identifier is an occurrence in its scope.

Free Identifiers



```
val x = 3
def add(a: Int, b: Int): Int = {
  val x = a + b
  x + y
}
val y = x + 1
add(y, z)
```

A free identifier is an identifier that is **not defined** in a program.

Shadowing



```
val x = 3
def add(a: Int, b: Int): Int = {
  val x = a + b
  x + y
}
val y = x + 1
add(y, z)
```

Shadowing means that the innermost binding occurrence shadows the outer binding occurrences of the same name.

- A shadowing identifier is an identifier that shadows another identifier.
- A shadowed identifier is an identifier that is shadowed by another identifier.

Contents



1. Identifiers

Bound Identifiers Free Identifiers Shadowing

2. VAE - AE with Variables

Concrete Syntax Abstract Syntax Examples

VAE – AE with Variables



Now, we want to extend AE into VAE with variables:

```
/* VAE */
val x = 1 + 2; // x = 1 + 2 = 3
val y = x + 3; // y = x + 3 = 3 + 3 = 6
y + 4 // 6 + 4 = 10
```

First, we need to define the **syntax** of VAE.

Concrete Syntax



Let's define the **concrete syntax** of VAE in BNF:





Let's define the **concrete syntax** of VAE in BNF:

For example, the following string is a valid VAE program:

```
/* VAE */
val x = 1 + 2;
val y = x + 3;
y + 4
```





Let's define the **concrete syntax** of VAE in BNF:

For example, the following string is a valid VAE program:

```
/* VAE */
val x = 1 + 2;
val y = x + 3;
y + 4
```

means

```
/* VAE */
val x = 1 + 2; {
 val y = x + 3; {
 y + 4
 }
}
```

Abstract Syntax



Let's define the abstract syntax of VAE in BNF:

Abstract Syntax



Let's define the **abstract syntax** of VAE in BNF:

Let's define Scala ADTs for VAE.

```
enum Expr:
    case Num(number: BigInt)
    case Add(left: Expr, right: Expr)
    case Mul(left: Expr, right: Expr)
    case Val(name: String, init: Expr, body: Expr)
    case Id(name: String)
```

Abstract Syntax



Let's define the abstract syntax of VAE in BNF:

Let's define Scala **ADTs** for VAE.

```
enum Expr:
   case Num(number: BigInt)
   case Add(left: Expr, right: Expr)
   case Mul(left: Expr, right: Expr)
   case Val(name: String, init: Expr, body: Expr)
   case Id(name: String)
```

For example, a VAE expression val x = 1; x + 2 is represented as:

```
Val("x", Num(1), Add(Id("x"), Num(2)))
```

Examples



For each VAE program, please draw:

- an arrow from each bound occurrence to its binding occurrence.
- a dotted arrow from each shadowing variable to its shadowed one.
- an X mark on each free identifier.

```
/* VAE */
val x = 1; x
```

```
/* VAE */
val x = x + 1;
val y = x * 2;
val x = y + 3;
x + y * z
```

```
/* VAE */
val x = 1;
val y = {
  val x = 2 * x;
  { val y = x; y } + { val y = 3; y }
};
x + y
```

Summary



1. Identifiers

Bound Identifiers Free Identifiers Shadowing

2. VAE - AE with Variables

Concrete Syntax Abstract Syntax Examples

Next Lecture



• Identifiers (2)

Jihyeok Park
jihyeok_park@korea.ac.kr
https://plrg.korea.ac.kr