

Lecture 19 – Typed Languages

COSE212: Programming Languages

Jihyeok Park



2023 Fall

- Safe Language Systems
 - Dynamic vs Static Analysis for Detecting Run-Time Errors
 - Soundness vs Completeness of Analysis
- Type Systems
 - Types
 - Type Errors
 - Type Checking
 - Type Soundness

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- In this lecture, we will define our first **typed language**.
- **TFAE** – FAE with **type system**.
 - **Type Checker** and **Typing Rules**
 - Interpreter and Natural Semantics

1. TFAE – FAE with Type System

- Concrete Syntax

- Abstract Syntax

2. Type Checker and Typing Rules

- Recall: Type Checking

- Type Environment

- Numbers

- Addition and Multiplication

- Immutable Variable Definition and Identifier Lookup

- Function Definition and Application

- Examples

3. Interpreter with Type Checker

- Interpreter and Natural Semantics

- Dynamic Semantics vs Static Semantics

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Dynamic Semantics vs Static Semantics

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It should take a **number** type argument and return a **number**.

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Let's say its type is **Number => Number** called **arrow type**.

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```
/* FAE */ x => x
```

How about this? There is no information on the parameter x .

One simple solution is to explicitly add **type annotations**!

Let's extend FAE into TFAE with **type annotations** to specify the types of function parameters:

```
/* TFAE */  
(x: Number) => x           // x is `Number` type  
(f: Number => Number) => 42 // f is `Number => Number` type
```

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```

If we define immutable variable definitions as **syntactic sugar**, it requires the type annotations: $\mathcal{D}[\text{val } x:\tau = e; e'] = (\lambda x:\tau. e')(\mathcal{D}[e])$

```
/* TFAE */  
val x: Number = 42; x + 1 // == `((x: Number) => x + 1)(42)`
```


Let's extend FAE into TFAE with **type annotations** to specify the types of function parameters:

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/* TFAE */  
(x: Number) => x           // x is `Number` type  
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```
/* TFAE */  
val x: Number = 42; x + 1 // == `((x: Number) => x + 1)(42)`
```

However, if we **explicitly define** them rather than syntactic sugar, we can guess variable types from their initial values:

```
/* TFAE */  
val x = 42; x + 1 // x is `Number` type because of `42`
```

For TFAE, we need to extend **expressions** of FAE with

- ① **function definitions** with **type annotations**
- ② **immutable variable definitions** without **type annotations**
- ③ **types**

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We can extend the **concrete syntax** of FAE as follows:

```
// expressions
<expr> ::= ...
        | "(" <id> ":" <type> ")" "=>" <expr>
        | "val" <id> "=" <expr> ";" <expr>

// types
<type> ::= "Number"           // number type
        | <type> "=>" <type> // arrow type
```

For TFAE, we need to extend **expressions** of FAE with

- ① **function definitions** with **type annotations**
- ② **immutable variable definitions** without **type annotations**
- ③ **types**

We can extend the **concrete syntax** of FAE as follows:

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// expressions
<expr> ::= ...
    | "(" <id> ":" <type> ")" "=" <expr>
    | "val" <id> "=" <expr> ";" <expr>

// types
<type> ::= "Number"           // number type
        | <type> "=" <type>   // arrow type
```

Since functions are first-class values, the parameter and return types could be recursively arrow types.

We can extend the **abstract syntax** of FAE for TFAE as follows:

Expressions $\mathbb{E} \ni e ::= \dots$

$\lambda x:\tau. e$	(Fun)
$\text{val } x=e; e$	(Val)

Types $\mathbb{T} \ni \tau ::= \text{num}$ (NumT)

$\tau \rightarrow \tau$	(ArrowT)
-------------------------	----------

We can define the abstract syntax of TFAE in Scala as follows:

```
enum Expr:
  ...
  case Fun(param: String, ty: Type, body: Expr)
  case Val(name: String, init: Expr, body: Expr)

enum Type:
  case NumT
  case ArrowT(paramTy: Type, bodyTy: Type)
```

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Dynamic Semantics vs Static Semantics

If the following conditions hold, we say “**the expression e has type τ** ”:

- e does not cause any type error, and
- e evaluates to a value of type τ or does not terminate.

If so, we use the following notation and say that e is **well-typed**:

$$\boxed{\vdash e : \tau}$$

Definition (Type Checking)

Type checking is a kind of static analysis checking whether a given expression e is **well-typed**. A **type checker** returns the **type** of e if it is well-typed, or rejects it and reports the detected **type error** otherwise.

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We need to

- ① design **typing rules** to define when an expression is well-typed
- ② implement a **type checker** in Scala according to typing rules

Let's ① design **typing rules** of TFAE to define when an expression is well-typed in the form of:

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and ❷ implement a **type checker** in Scala according to typing rules:

```
def typeCheck(expr: Expr): Type = ???
```

The type checker returns the **type** of e if it is well-typed, or rejects it and throws a **type error** otherwise.

Let's ❶ design **typing rules** of TFAE to define when an expression is well-typed in the form of:

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In addition, we need to keep track of the **variable types**.

Let's ① design **typing rules** of TFAE to define when an expression is well-typed in the form of:

$$\Gamma \vdash e : \tau$$

and ② implement a **type checker** in Scala according to typing rules:

```
def typeCheck(expr: Expr, tenv: TypeEnv): Type = ???
```

The type checker returns the **type** of e if it is well-typed, or rejects it and throws a **type error** otherwise.

In addition, we need to keep track of the **variable types**.

Let's define a **type environment** Γ as a mapping from variable names to their types and pass it to the type checker.

$$\text{Type Environments} \quad \Gamma \in \mathbb{X} \xrightarrow{\text{fin}} \mathbb{T} \quad (\text{TypeEnv})$$

```
type TypeEnv = Map[String, Type]
```

```
def typeCheck(expr: Expr, tenv: TypeEnv): Type = expr match  
  case Num(_) => ???  
  ...
```

$$\boxed{\Gamma \vdash e : \tau}$$
$$\tau\text{-Num} \frac{}{\Gamma \vdash n : ???}$$

```
def typeCheck(expr: Expr, tenv: TypeEnv): Type = expr match
  case Num(_) => NumT
  ...
```

$$\boxed{\Gamma \vdash e : \tau}$$

$$\tau\text{-Num} \frac{}{\Gamma \vdash n : \text{num}}$$

The number literal n has `num` type in any type environment Γ .

```
def typeCheck(expr: Expr, tenv: TypeEnv): Type = expr match
  ...
  case Add(left, right) =>
    ???
```

$$\boxed{\Gamma \vdash e : \tau}$$

$$\tau\text{-Add} \frac{\text{???}}{\Gamma \vdash e_1 + e_2 : \text{???}}$$

```
def typeCheck(expr: Expr, tenv: TypeEnv): Type = expr match
  ...
  case Add(left, right) =>
    typeCheck(left, tenv)
    ???
```

$$\boxed{\Gamma \vdash e : \tau}$$

$$\tau\text{-Add} \frac{\Gamma \vdash e_1 : \tau \quad ???}{\Gamma \vdash e_1 + e_2 : ???}$$

Type checker should do

- 1 get the type of e_1 in Γ


```
def typeCheck(expr: Expr, tenv: TypeEnv): Type = expr match
  ...
  case Add(left, right) =>
    mustSame(typeCheck(left, tenv), NumT)
    ???

def mustSame(lty: Type, rty: Type): Unit =
  if (lty != rty) error(s"type error: ${lty.str} != ${rty.str}")
```

$$\boxed{\Gamma \vdash e : \tau}$$

$$\tau\text{-Add} \frac{\Gamma \vdash e_1 : \text{num} \quad ???}{\Gamma \vdash e_1 + e_2 : ???}$$

Type checker should do

- 1 check the type of e_1 is `num` in Γ

```
def typeCheck(expr: Expr, tenv: TypeEnv): Type = expr match
  ...
  case Add(left, right) =>
    mustSame(typeCheck(left, tenv), NumT)
    mustSame(typeCheck(right, tenv), NumT)
    ???

def mustSame(lty: Type, rty: Type): Unit =
  if (lty != rty) error(s"type error: ${lty.str} != ${rty.str}")
```

$$\boxed{\Gamma \vdash e : \tau}$$

$$\tau\text{-Add} \frac{\Gamma \vdash e_1 : \text{num} \quad \Gamma \vdash e_2 : \text{num}}{\Gamma \vdash e_1 + e_2 : ???}$$

Type checker should do

- 1 check the types of e_1 and e_2 are num in Γ

```
def typeCheck(expr: Expr, tenv: TypeEnv): Type = expr match
  ...
  case Add(left, right) =>
    mustSame(typeCheck(left, tenv), NumT)
    mustSame(typeCheck(right, tenv), NumT)
    NumT

def mustSame(lty: Type, rty: Type): Unit =
  if (lty != rty) error(s"type error: ${lty.str} != ${rty.str}")
```

$$\boxed{\Gamma \vdash e : \tau}$$

$$\tau\text{-Add} \frac{\Gamma \vdash e_1 : \text{num} \quad \Gamma \vdash e_2 : \text{num}}{\Gamma \vdash e_1 + e_2 : \text{num}}$$

Type checker should do

- ① check the types of e_1 and e_2 are num in Γ
- ② return num as the type of $e_1 + e_2$

```
def typeCheck(expr: Expr, tenv: TypeEnv): Type = expr match
  ...
  case Mul(left, right) =>
    mustSame(typeCheck(left, tenv), NumT)
    mustSame(typeCheck(right, tenv), NumT)
    NumT

def mustSame(lty: Type, rty: Type): Unit =
  if (lty != rty) error(s"type error: ${lty.str} != ${rty.str}")
```

$$\boxed{\Gamma \vdash e : \tau}$$

$$\tau\text{-Mul} \quad \frac{\Gamma \vdash e_1 : \text{num} \quad \Gamma \vdash e_2 : \text{num}}{\Gamma \vdash e_1 \times e_2 : \text{num}}$$

Type checker should do

- ① check the types of e_1 and e_2 are num in Γ
- ② return num as the type of $e_1 \times e_2$

```
def typeCheck(expr: Expr, tenv: TypeEnv): Type = expr match
...
case Val(x, init, body) =>
  val initTy = typeCheck(init, tenv)
  typeCheck(body, tenv + (x -> initTy))
```

$$\boxed{\Gamma \vdash e : \tau}$$

$$\tau\text{-Val} \frac{\Gamma \vdash e_1 : \tau_1 \quad \Gamma[x \mapsto \tau_1] \vdash e_2 : \tau_2}{\Gamma \vdash \text{val } x = e_1; e_2 : \tau_2}$$

This rule stores the type of x in Γ inferred from the initial value.

```
/* TFAE */ val x = 1; x + 2      // `x: Number` in `tenv` <- `1: Number`
```

```
def typeCheck(expr: Expr, tenv: TypeEnv): Type = expr match
...
case Id(x) =>
  tenv.getOrElse(x, error(s"free identifier: $x"))
```

$$\boxed{\Gamma \vdash e : \tau}$$

$$\tau\text{-Id} \frac{x \in \text{Domain}(\Gamma)}{\Gamma \vdash x : \Gamma(x)}$$

This rule looks up the type of x in Γ .

```
/* TFAE */ val x = 1; x + 2      // `x: Number` in `tenv` <- `1: Number`
```

```
def typeCheck(expr: Expr, tenv: TypeEnv): Type = expr match
...
case Fun(param, paramTy, body) =>
  val bodyTy = typeCheck(body, tenv + (param -> paramTy))
  ArrowT(paramTy, bodyTy)
```

$$\boxed{\Gamma \vdash e : \tau}$$

$$\tau\text{-Fun} \frac{\Gamma[x \mapsto \tau] \vdash e : \tau'}{\Gamma \vdash \lambda x:\tau. e : \tau \rightarrow \tau'}$$

We can check the body of a function with the its parameter type.

```
/* TFAE */ (x: Number) => x      // Number => Number
```

```
def typeCheck(expr: Expr, tenv: TypeEnv): Type = expr match
...
case App(fun, arg) => typeCheck(fun, tenv) match
  case ArrowT(paramTy, bodyTy) =>
    mustSame(typeCheck(arg, tenv), paramTy)
    bodyTy
  case ty => error(s"not a function type: ${ty.str}")
```

$$\boxed{\Gamma \vdash e : \tau}$$

$$\tau\text{-App} \frac{\Gamma \vdash e_0 : \tau_1 \rightarrow \tau_2 \quad \Gamma \vdash e_1 : \tau_1}{\Gamma \vdash e_0(e_1) : \tau_2}$$

We don't have to check the type of the function body because it is already checked when the function is defined.

```
/* TFAE */ ((x: Number) => x)(1) // Number
```



```
/* TFAE */ val x = 1; x + 2           // 3: Number
```

```
/* TFAE */ val x = 1; x + 2           // 3: Number
```

$$\frac{\frac{\frac{}{\emptyset \vdash 1 : \text{num}}}{x \in \text{Domain}([x \mapsto \text{num}])} \quad \frac{[x \mapsto \text{num}] \vdash x : \text{num} \quad [x \mapsto \text{num}] \vdash 2 : \text{num}}{[x \mapsto \text{num}] \vdash x + 2 : \text{num}}}{\emptyset \vdash \text{val } x=1; x + 2 : n}$$

```
/* TFAE */ val x = 1; x + 2           // 3: Number
```

$$\frac{\frac{\frac{}{\emptyset \vdash 1 : \text{num}}}{x \in \text{Domain}([x \mapsto \text{num}])} \quad \frac{\frac{[x \mapsto \text{num}] \vdash x : \text{num} \quad [x \mapsto \text{num}] \vdash 2 : \text{num}}{[x \mapsto \text{num}] \vdash x + 2 : \text{num}}}{\emptyset \vdash \text{val } x=1; x + 2 : n}$$

```
/* TFAE */ ((x: Number) => x)(2) * 3   // 3: Number
```

```
/* TFAE */ val x = 1; x + 2           // 3: Number
```

$$\frac{\frac{}{\emptyset \vdash 1 : \text{num}} \quad \frac{\frac{x \in \text{Domain}([x \mapsto \text{num}])}{[x \mapsto \text{num}] \vdash x : \text{num}} \quad \frac{}{[x \mapsto \text{num}] \vdash 2 : \text{num}}}{[x \mapsto \text{num}] \vdash x + 2 : \text{num}}}{\emptyset \vdash \text{val } x=1; x + 2 : n}$$

```
/* TFAE */ ((x: Number) => x)(2) * 3   // 3: Number
```

$$\frac{\frac{\frac{x \in \text{Domain}([x \mapsto \text{num}])}{[x \mapsto \text{num}] \vdash x : \text{num}}}{\emptyset \vdash \lambda x : \text{num}. x : \text{num} \rightarrow \text{num}} \quad \frac{}{\emptyset \vdash 2 : \text{num}}}{\emptyset \vdash (\lambda x : \text{num}. x)(2) : \text{num}} \quad \frac{}{\emptyset \vdash 3 : \text{num}}}{\emptyset \vdash (\lambda x : \text{num}. x)(2) \times 3 : \text{num}}$$

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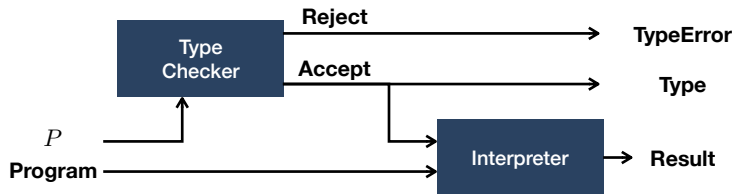
Examples

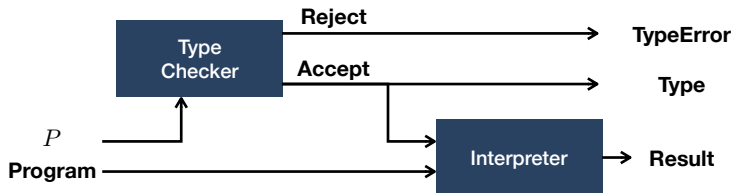
3. Interpreter with Type Checker

Interpreter and Natural Semantics

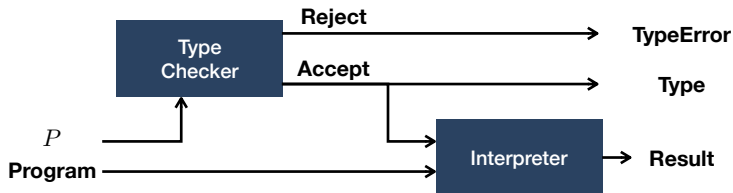
Dynamic Semantics vs Static Semantics

Interpreter with Type Checker





```
def eval(str: String): String =  
  val expr = Expr(str)  
  val ty = typeCheck(expr, Map.empty)  
  val v = interp(expr, Map.empty)  
  s"${v.str}: ${ty.str}"
```



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  val expr = Expr(str)  
  val ty = typeCheck(expr, Map.empty)  
  val v = interp(expr, Map.empty)  
  s"${v.str}: ${ty.str}"
```

```
eval("val x = 1; x + 2")           // 3: Number  
eval("((x: Number) => x)(2) * 3") // 6: Number  
eval("1 + (x: Number) => x")      // a type error thrown by `typeCheck`
```


For interpreter and natural semantics for TFAE, it is just enough to extend the those for function definitions in FAE.

```
def interp(expr: Expr, env: Env): Value = expr match
  ...
  case Fun(p, t, b) => CloV(p, b, env)
```

$$\sigma \vdash e \Rightarrow v$$

$$\text{Fun} \frac{}{\sigma \vdash \lambda x:\tau. e \Rightarrow \langle \lambda x. e, \sigma \rangle}$$

The type annotation is ignored in the interpreter and natural semantics.

What is the difference between **operational semantics** and **typing rules**?

$$\boxed{\sigma \vdash e \Rightarrow v} \quad \text{vs} \quad \boxed{\Gamma \vdash e : \tau}$$

What is the difference between **operational semantics** and **typing rules**?

$$\boxed{\sigma \vdash e \Rightarrow v} \quad \text{vs} \quad \boxed{\Gamma \vdash e : \tau}$$

We call the former **dynamic semantics** and the latter **static semantics**!

See the table below for the comparison.

	Dynamic Semantics	Static Semantics
Rule	$\sigma \vdash e \Rightarrow v$	$\Gamma \vdash e : \tau$
Purpose	Evaluation	Type Checking
Implementation	Interpreter	Type Checker
Result	Value	Type

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- Dynamic Semantics vs Static Semantics

- Please see this document¹ on GitHub.
 - Implement `typeCheck` function.
 - Implement `interp` function.
- It is just an exercise, and you **don't need to submit** anything.
- However, some exam questions might be related to this exercise.

¹<https://github.com/ku-plrg-classroom/docs/tree/main/cose212/tfae>.

- Typing Recursive Functions

Jihyeok Park

jihyeok_park@korea.ac.kr

<https://plrg.korea.ac.kr>