Lecture 4 – Identifiers (1)

COSE212: Programming Languages

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• ADT for Abstract Syntax of AE

```
enum Expr:
    case Num(number: BigInt)
    case Add(left: Expr, right: Expr)
    case Mul(left: Expr, right: Expr)
```

• Parser for Concrete Syntax of AE

```
lazy val expr: P[Expr] = ...
```

• Interpreter for Semantics of AE

```
def interp(expr: Expr): Value = ...
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• ADT for Abstract Syntax of AE

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• Parser for Concrete Syntax of AE

```
lazy val expr: P[Expr] = ...
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• Interpreter for Semantics of AE

```
def interp(expr: Expr): Value = ...
```

• In this lecture, we will learn identifiers.

Contents



1. Identifiers

Bound Identifiers Free Identifiers Shadowing

2. VAE - AE with Variables

Concrete Syntax Abstract Syntax Exercise

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1. Identifiers

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An identifier is a name for a certain element in a program.

In Scala, there are diverse kinds of identifiers:

```
// variable names
val x: Int = 42

// function and parameter names
def f(a: Int, b: Int): Int = a + b

// class and field names
case class Person(name: String, age: Int)
...
```



```
val x: Int = 3
val y: Int = x + 1
def f(a: Int, b: Int): Int = {
  val x: Int = a + b
  x + y + z
}
f(x, b)
```

- A binding occurrence of an identifier is the occurrence in its definition position.
- A scope of an identifier is a code region where the identifier is usable.
- A bound occurrence of an identifier is an occurrence in a lookup position in its scope.



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Free Identifiers



```
val x: Int = 3
val y: Int = x + 1
def f(a: Int, b: Int): Int = {
  val x: Int = a + b
  x + y + z
}
f(x, b)
```

A **free identifier** is an identifier that is **not defined** in the current scope of the program.

Shadowing



```
val x: Int = 3
val y: Int = x + 1
def f(a: Int, b: Int): Int = {
  val x: Int = a + b
  x + y + z
}
f(x, b)
```

Shadowing means that the innermost binding occurrence shadows the outer binding occurrences of the same name.

- A shadowing identifier is an identifier that shadows another identifier.
- A shadowed identifier is an identifier that is shadowed by another identifier.

Shadowing



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val x: Int = 3
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def f(a: Int, b: Int): Int = {
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Shadowing means that the innermost binding occurrence shadows the outer binding occurrences of the same name.

- A shadowing identifier is an identifier that shadows another identifier.
- A shadowed identifier is an identifier that is shadowed by another identifier.

Note that this is **NOT** a mutation because the value stored in the shadowed identifier is unchanged.

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2. VAE – AE with Variables

Concrete Syntax Abstract Syntax Exercise

VAE – AE with Variables



Now, we want to extend AE into VAE with variables:

```
/* VAE */
val x = 1 + 2; // x = 1 + 2 = 3
val y = x + 3; // y = x + 3 = 3 + 3 = 6
y + 4 // 6 + 4 = 10
```





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/* VAE */
val x = 1 + 2; // x = 1 + 2 = 3
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y + 4 // 6 + 4 = 10
```

First, we define the **concrete syntax** of **identifiers** used in VAE:

```
<alphabet> ::= "A" | "B" | "C" | ... | "Z" | "a" | "b" | "c" | ... | "z"
<idstart> ::= <alphabet> | "_"
<idcont> ::= <alphabet> | "_" | <digit>
<id> ::= <idstart> <idcont>*
```



Now, we want to extend AE into VAE with variables:

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/* VAE */
val x = 1 + 2; // x = 1 + 2 = 3
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<idstart> ::= <alphabet> | "_"
<idcont> ::= <alphabet> | "_" | <digit>
<id> ::= <idstart> <idcont>*
```

For example, the following are valid identifiers:

x y get_name getName add42

Concrete Syntax



Now, let's define the **concrete syntax** of VAE in BNF:

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Note that each variable definition creates a **new scope**.





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Note that each variable definition creates a **new scope**. For example:

```
/* VAE */
val x = 1 + 2;
val y = x + 3;
y + 4
```

means

```
/* VAE */
val x = 1 + 2;
{
  val y = x + 3;
  {
    y + 4
  }
}
```

Abstract Syntax



Let's define the abstract syntax of VAE in BNF:

$$e ::= n$$
 (Num)
 $| e + e$ (Add)
 $| e \times e$ (Mul)
 $| val x = e; e$ (Val)
 $| x$ (Id)

Abstract Syntax



Let's define the **abstract syntax** of VAE in BNF:

```
e ::= n (Num)

| e + e (Add)

| e \times e (Mul)

| val x = e; e (Val)

| x (Id)
```

```
enum Expr:
    case Num(number: BigInt)
    case Add(left: Expr, right: Expr)
    case Mul(left: Expr, right: Expr)
    // variable definition
    case Val(name: String, init: Expr, body: Expr)
    // variable lookup
    case Id(name: String)
```

Abstract Syntax



Let's define the **abstract syntax** of VAE in BNF:

```
enum Expr:
    case Num(number: BigInt)
    case Add(left: Expr, right: Expr)
    case Mul(left: Expr, right: Expr)
    // variable definition
    case Val(name: String, init: Expr, body: Expr)
    // variable lookup
    case Id(name: String)
```

```
Expr("val x = 1; x + 2")  // Val("x", Num(1), Add(Id("x"), Num(2)))
```





For each VAE program, please draw:

- an arrow from each bound occurrence to its binding occurrence.
- a dotted arrow from each shadowing variable to its shadowed one.
- an X mark on each free variable.

```
/* VAE */
val x = 1; x
```

```
/* VAE */
val x = x + 1;
val y = x * 2;
val x = y + x;
x * z
```

```
/* VAE */
val x = 1;
val y = {
  val x = 2 * x;
  { val y = x; y } + { val y = 3; y }
};
x + y
```

Summary



1. Identifiers

Bound Identifiers Free Identifiers Shadowing

2. VAE – AE with Variables

Concrete Syntax Abstract Syntax Exercise

Next Lecture



• Identifiers (2)

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