

2021 FALL OS Project 3 Help Document



Distributed Computing Systems Laboratory
Department of Computer Science and Engineering
Seoul National University, Korea
2021-10-14

NOTE

•과제 #0 Set up development environment

기한: 09/18 midnight

방법: etl 제출

•과제 #1

기한: 10/08 midnight

방법: 팀별 repository에 proj1 branch 생성 / 발표자료는 대표 한 명이 etl 제출

•과제 #2

기한: 11/06 midnight

방법: 팀별 repository에 proj2 branch 생성 / 발표자료는 대표 한 명이 etl 제출

•과제 #3

기한: 11/27 midnight

방법: 팀별 repository에 proj3 branch 생성 / 발표자료는 대표 한 명이 etl 제출

•과제 #4

기한: 12/15 midnight

방법: 팀별 repository에 proj4 branch 생성 / 발표자료는 대표 한 명이 etl 제출



Project 3 Overview

- Implement rotation-based read/write lock
- Each lock has a "rotation range"
 - o Every lock can be acquired when the current rotation is in the rotation range
 - o Or, it is blocked until the current rotation is located in the rotation range
- Read lock could be acquired when no acquired write lock range is overlapping with its range
- Write lock could be acquired when no acquired read/write lock range is overlapping with its range
 - Exclusive access



Rotation Range

- 1 axis: rotation(use daemon)
 - Actually, Tizen orientation has three axes! (Azimuth, Pitch, Roll)
- (degree range) <= range <= (degree + range)
- Rotation ranges are inclusive
 - o Ex) [30, 60], [60, 90] are overlapped
- Rotation ranges are circular
 - o Ex) [330, 30] and [30, 330]
 - o [330 ... 0 ... 30]
 - o [30 ... 180 ... 330]



Range Example (1)

- Rotation 1
 - \circ degree = 30
 - ∘ range = 30

Are they overlapped?

- Rotation 2
 - \circ degree = 45
 - ∘ range = 30



Range Example (1)

- Rotation 1
 - \circ degree = 30
 - ∘ range = 30

- Rotation 1
 - 0, 60]

Rotation 2

 \circ degree = 45 ∘ range = 30

- Rotation 2
 - o [15, 75]

Are they overlapped? YES!

Range Example (2)

- Rotation 1
 - \circ degree = 30
 - ∘ range = 60

Are they overlapped?

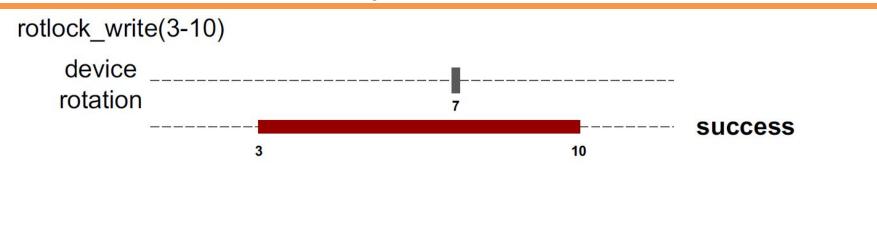
- Rotation 2
 - \circ degree = 315
 - ∘ range = 30

Range Example (2)

- Rotation 1
 - \circ degree = 30
 - ∘ range = 60
- Rotation 1
 - o [0, 90] + [330, 360)

- Rotation 2
 - degree = 315
 - ∘ range = 30
- Rotation 2
 - o [285, 345]

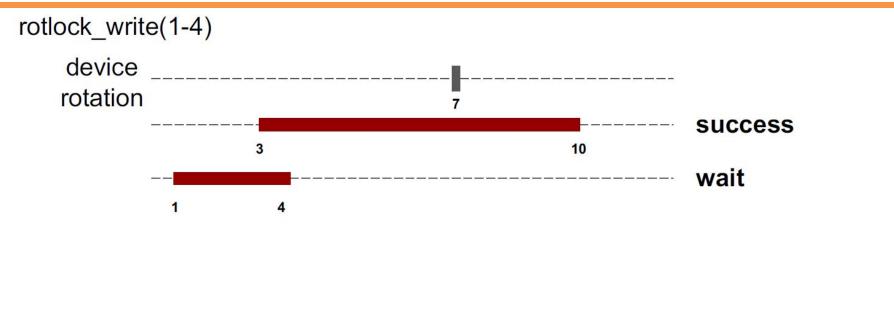
Are they overlapped? YES!



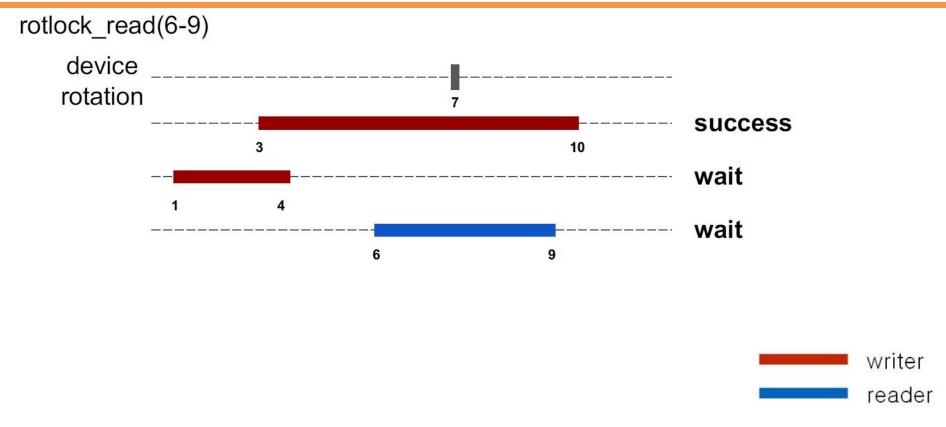


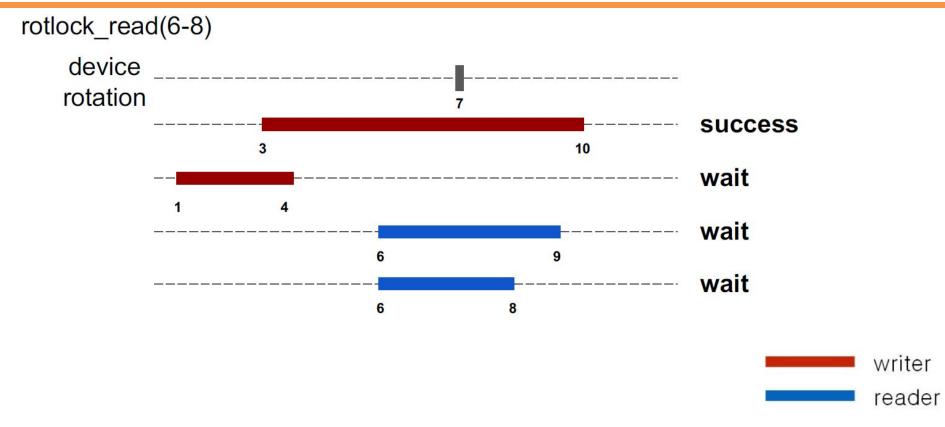




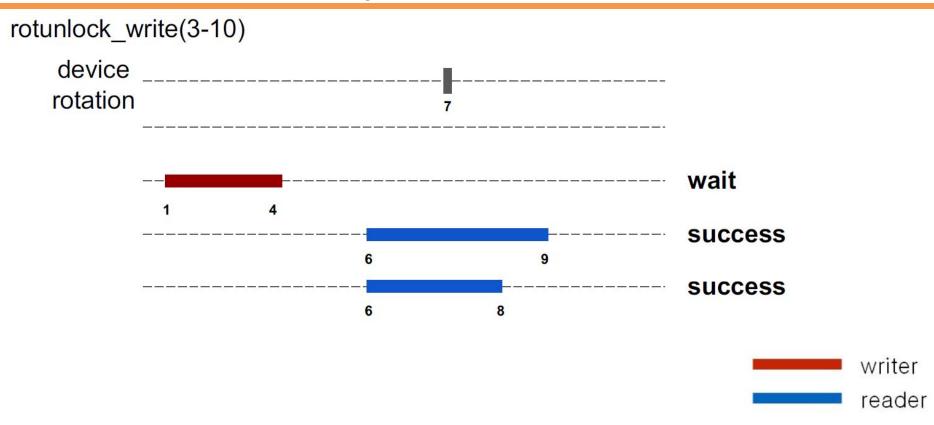




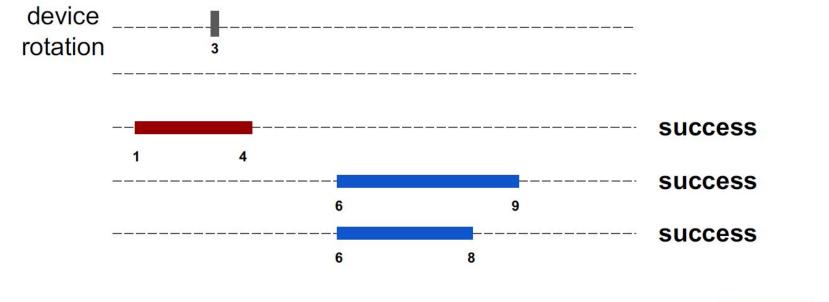








Device rotation changed to 3

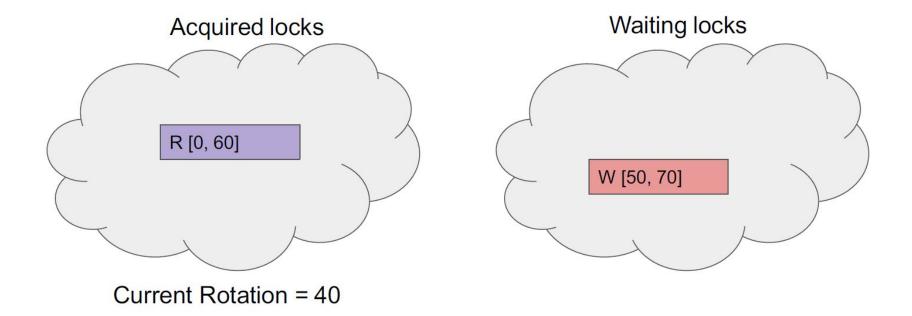




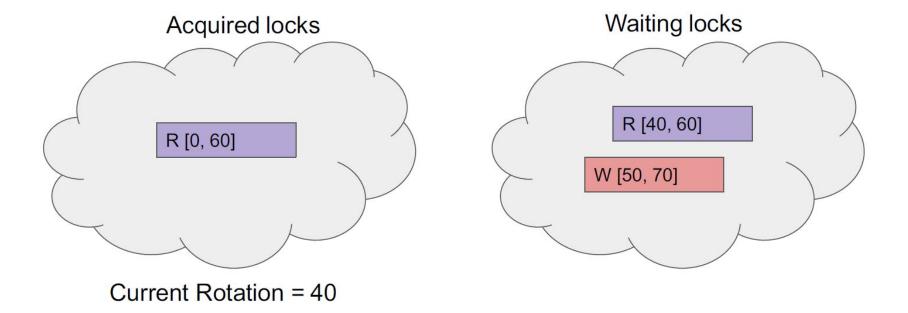
Preventing writer starvation

- You should implement a policy for preventing starvation of writers
 - Why? If reader comes in infinitely, a writer could wait forever!
- If a reader holds a lock and a writer wants to take the lock, no more readers can take the lock
- If you design your own additional policy, explain that in your README.md file, report, and slides!

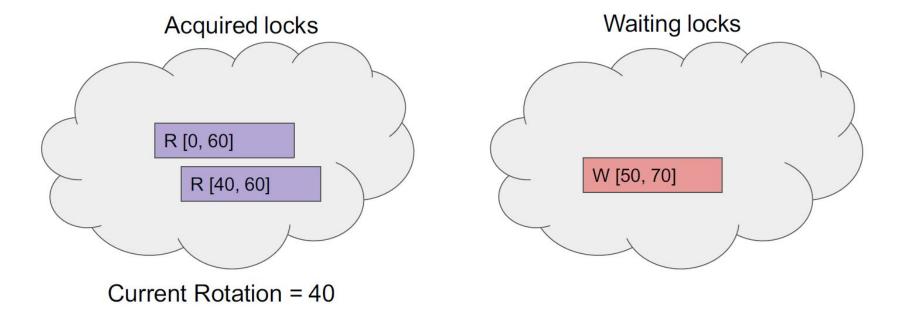
• A write lock is waiting for the rotation changes & the reader to release its lock



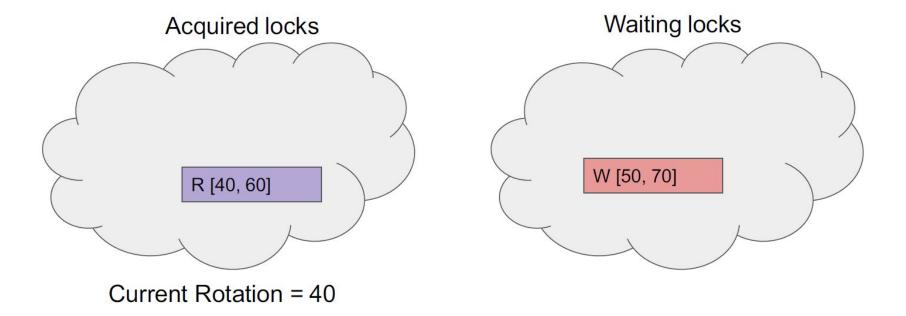
• A read lock [40, 60] came



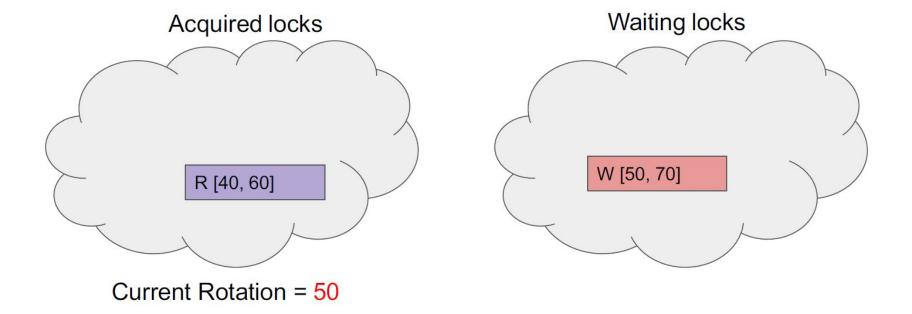
• $40 \in [40, 60] \rightarrow \text{Acquires its lock immediately}$



• R [0, 60] releases its lock

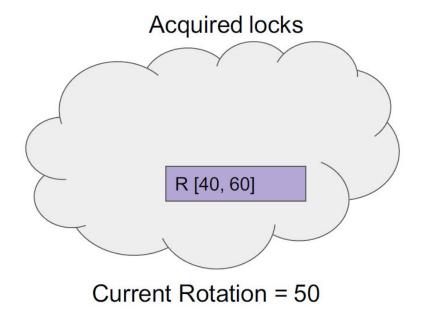


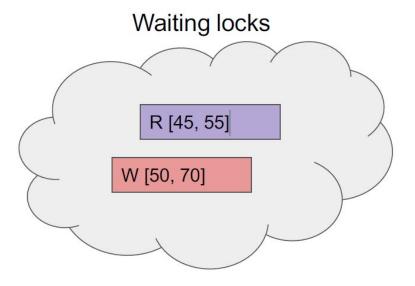
• Rotation changes 40 →50



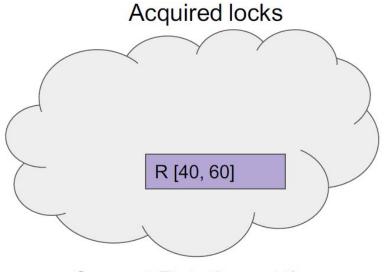
How to prevent writer starvation?

Acquired locks Waiting locks

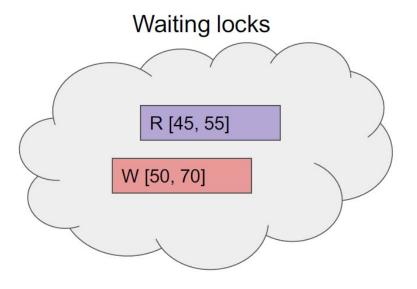




- W [50, 70] is waiting and 50 ∈ [50, 70]
- W [50, 70] is waiting for R [40, 60] to release its lock
- [45, 55] overlaps with [50, 70]



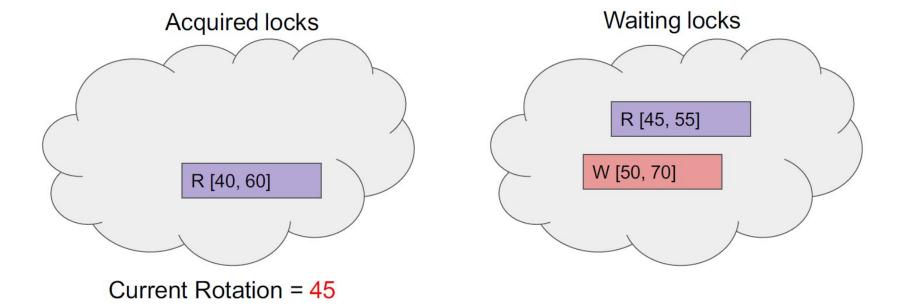
Current Rotation = 50



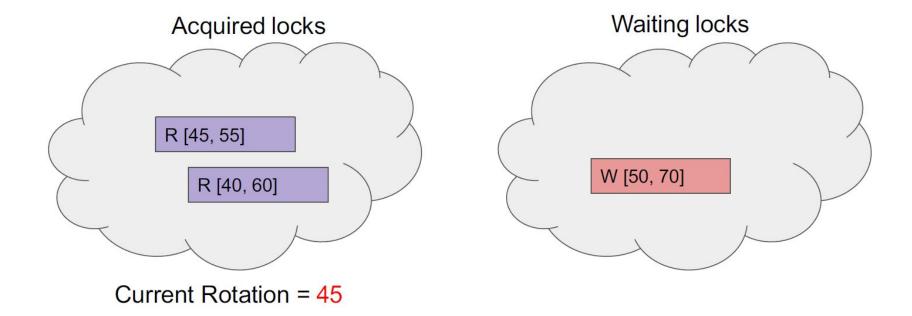
• W [50, 70] is waiting and 50 ∈ [50, 70] If a reader holds a lock • W [50, 70] is waitin and a writer wants to take the lock, • [45, 55] overlaps w ocks no more readers can take the lock R [45, 55] W [50, 70] R [40, 60] Current Rotation = 50

• W [50, 70] is waiting and $50 \in [50, 70]$ If a reader(R[40, 60]) holds a lock • W [50, 70] is waitin and a writer(W[50, 70]) wants to take($50 \in [50, 70]$) • [45, 55] overlaps w ocks A the lock([40, 60] overlaps with [50, 70], waiting), no more readers(R[45, 55]) can take the lock([45, 55] overlaps with [50, 70]) W [50, 70] R [40, 60] Current Rotation = 50

Rotation changes 50 →45



- W [50, 70] cannot grab its lock anymore (*.* 45 € [50, 70])
- Starvation prevention policy is no more applied → R [45, 55] acquires its lock

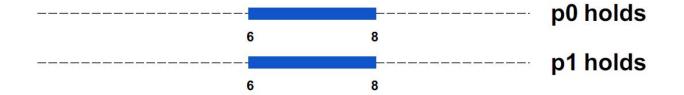


Terminating routine

- When a thread that has some holding or waiting locks is terminating, the remaining locks should be released (holding) or removed (waiting).
- Hints
 - o exit_rotlock() in kernel/rotation.c
 - Release holding locks
 - Remove waiting locks
 - Inject exit_rotlock() to do_exit() in kernel/exit.c

Isolation

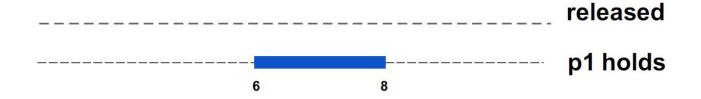
- Multiple processes shares same rotation lock system.
 - o You have to identify which process (thread) the lock belongs to.
- A process can't release locks that other processes hold.



Isolation

• A process can't release locks that other processes hold.

p0 calls unlock(6-8)



Isolation

• A process can't release locks that other processes hold.

p0 calls unlock(6-8) again



- You will have create things like acquired lock list and waiting lock list
 - Different approaches are also possible :)
 - Ex) Bitmap, linked list, ...
- Accessing those lists could result in race conditions
 - o if you are working on multicore machine
 - You should carefully design your code to prevent possible race conditions
 - Ex) One thread is removing a lock from waiting list, but another thread can access to the waiting list at the same time (inconsistency issues)

- Possible approaches
 - Global
 - Ex) Manage acquired lock list and waiting lock list using a same lock
 - Fine-grained
 - Ex) Manage acquired lock list and waiting lock list using separated locks
 - Better concurrency, more complicated :)
- Synchronization mechanism
 - Spin Lock
 - Eligible for short sleep (e.g. short list iteration)
 - RCU (Read-Copy-Update)
 - 0 ...

- How to block processes | How to wake up blocked processes
 - You may use not explained in this ppt :)
 - Wait Queue (Starts with DECLARE_WAIT_QUEUE_HEAD)
 - Condition Variable (Define your own CV)
 - Mutex
 - ...

- Lists could be changed during iteration
 - list_for_each_entry_safe could be useful to you
- Please remember that ...
 - The rotation range is circular!
 - You should implement a logic for determining two circular ranges are overlapping or not
 - Observe of deadlocks!

Selector and trial

- Selector & Trial require a same lock (0 <= range <= 180)
 - o If current rotation is 240, both selector & trial wait.
- When the device rotation is out of that range, both Selector & Trial stop working
- When the device rotation gets inside that range, Selector & Trial start to work

Selector	Trial
write_lock	
	read_lock & wait
10	
write_unlock	acquire_lock
	10 = 2 * 5
acquire lock	read_unlock

About submission

Same rules as previous projects

- Make sure your branch name: proj3
- Save your C program as: test/selector.c, test/trial.c
- Check for format : slides title / demo name / test file names / branch name and directory name
- Please aggregate your demo videos (=submit only one video!)
- Deadline
 - Due:11/27 midnight + 3 days of late submission allowed (10% will be deducted for every day)