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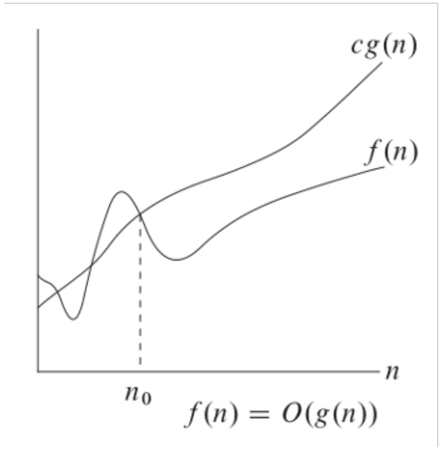

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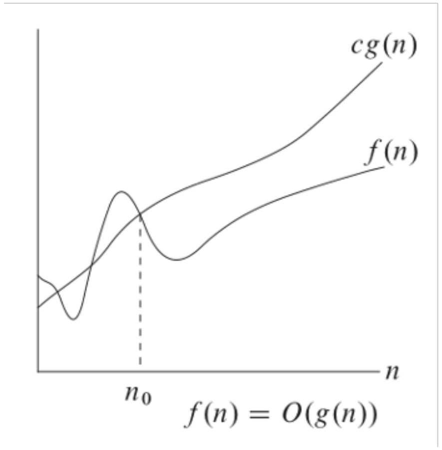

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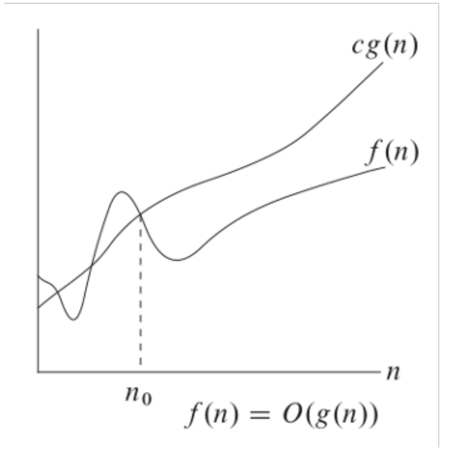

Insertion Sort as an $\boxed{\text{example}}$:

- Best Case: $T(n) \approx cn + c'n - c''$
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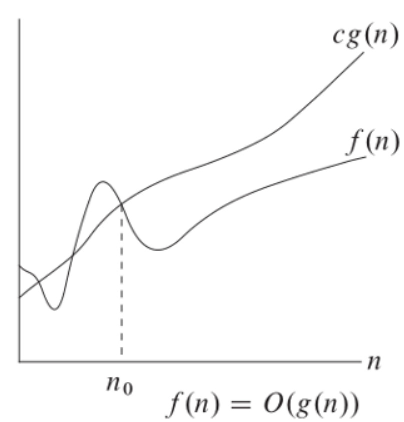

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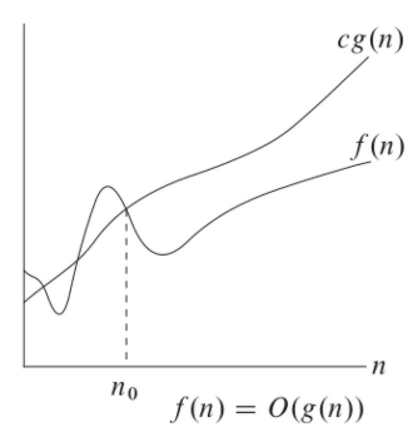

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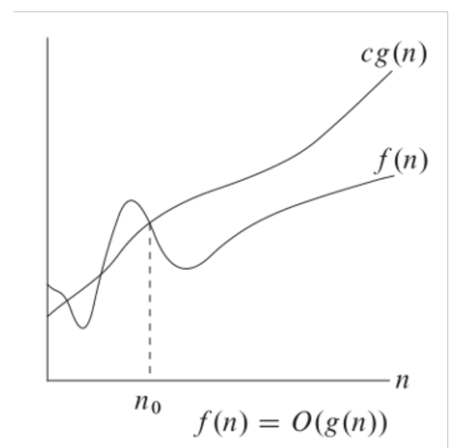

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No.

Show that the equation $\left(\frac{3}{2}\right)^n \geq c$
has infinitely many solutions for n.