

Starred 106k

Algorithm Techniques

Graph Theory - Basic Knowledge

Xian Su

Watch 569 ▼

Y Fork 13.3k

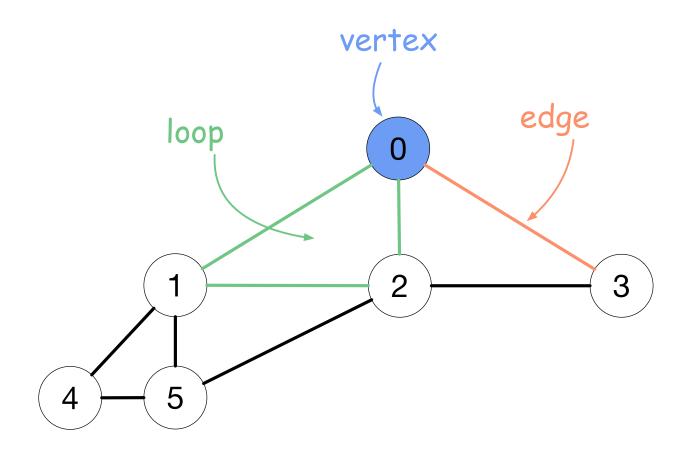
Recommended Reference Material:

- Github Open Source Project: Hello Algo
- Github Repo: <u>krahets/hello-algo</u>
- Animated Illustrations
- Support More than 10 Programming Languages
- Programming, Data Structure, Algorithm, and Algorithm Techniques

Graphs and Trees



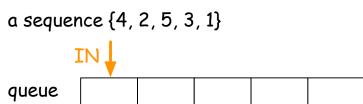
- A graph is a finite set of vertices that are connected by edges.
- A loop may exist in a graph where an edge leads back to the original vertex.



Necessary Recap: Queue & Stack

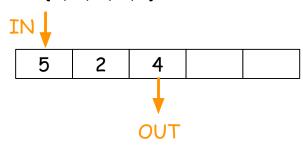


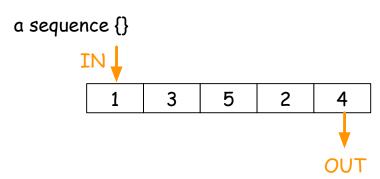
Queue



First IN First OUT

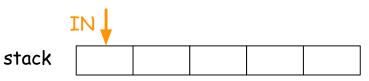
a sequence {4, 2, 5, 3, 1}





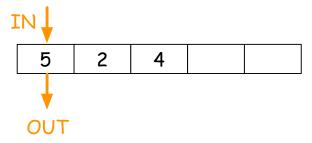
Stack

a sequence {4, 2, 5, 3, 1}

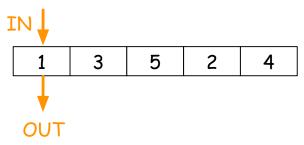


First IN Last OUT

a sequence {4, 2, 5, 3, 1}



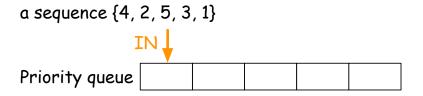
a sequence {}



Necessary Recap: Priority Queue



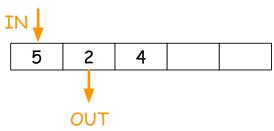
• The priority queue (PQ) is an abstract concept (a queue where elements are dequeued based on priority).



First IN PRIORITY OUT

e.g., priority = minimum value

a sequence {4, 2, 5, 3, 1}



a sequence {}

IN

1 3 5 2 4

OUT

Example 1: Vertex Traversal in an Undirected Graph

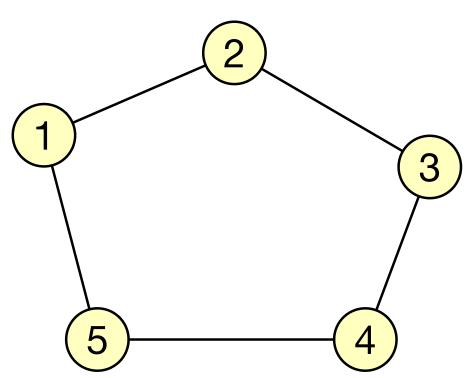


- $V = \{1,2,3,4,5\}$
- $E = \{(1,2), (1,5), (2,3), (3,4), (4,5)\}$
- 1. Use **Depth-First Search (DFS)** starting from vertex 1 and output the order of visited vertices.
- 2. Use Breadth-First Search (BFS) starting from vertex 1 and output the order of visited vertices.

Example 1: Vertex Traversal in an Undirected Graph



- $V = \{1,2,3,4,5\}$
- $E = \{(1,2), (1,5), (2,3), (3,4), (4,5)\}$
- 1. Use $Depth-First\ Search\ (DFS)$ starting from vertex 1 and output the order of visited vertices.
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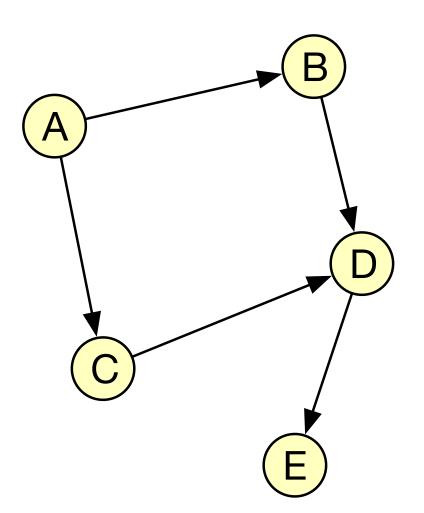




- $V = \{A, B, C, D, E\}$
- $E = \{(A, B), (A, C), (B, D), (C, D), (D, E)\}$
- 1. Use **BFS** starting from vertex A to find all paths to vertex E.
- 2. Use **DFS** starting from vertex A to list the sequence of all visited vertices during the traversal.

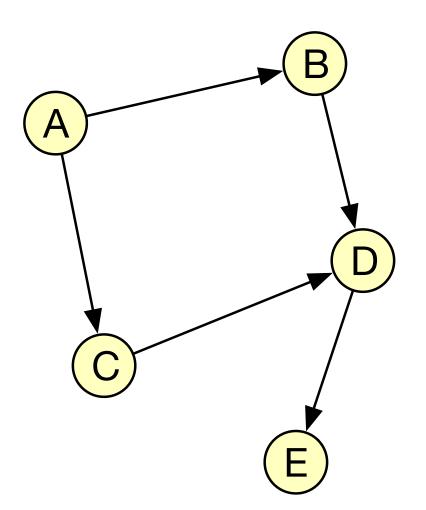


- $V = \{A, B, C, D, E\}$
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- $E = \{(A, B), (A, C), (B, D), (C, D), (D, E)\}$
- 1. Use **BFS** starting from vertex A to find all paths to vertex E.

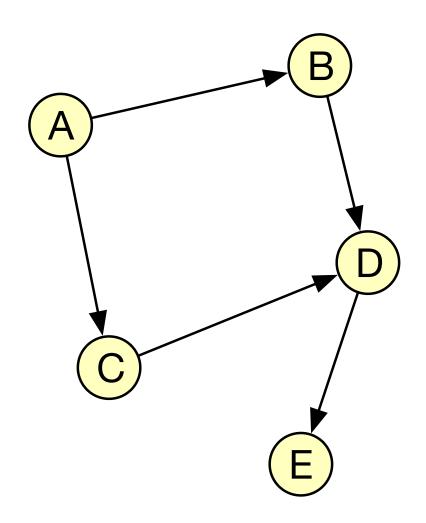




- $V = \{A, B, C, D, E\}$
- $E = \{(A, B), (A, C), (B, D), (C, D), (D, E)\}$
- 2. Use **DFS** starting from vertex A to list the sequence of all visited vertices during the traversal.

```
def dfs(graph, start, visited=None):
    if visited is None:
        visited = []
    visited.append(start)
    for neighbor in graph[start]:
        if neighbor not in visited:
            dfs(graph, neighbor, visited)
    return visited

graph = {1: [2, 5], 2: [1, 3], 3: [2, 4], 4: [3, 5], 5: [1, 4]}
print(dfs(graph, 1))
```

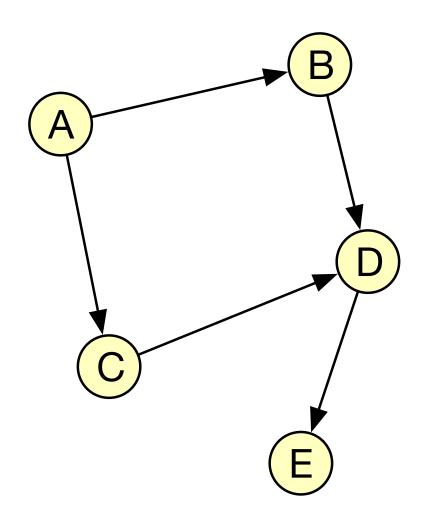




- $V = \{A, B, C, D, E\}$
- $E = \{(A, B), (A, C), (B, D), (C, D), (D, E)\}$
- 1. Use **BFS** starting from vertex A to find all paths to vertex E.

```
def bfs_all_paths(graph, start, target):
    queue = [[start]]
    paths = []
    while queue:
        path = queue.pop(0)
        node = path[-1]
        if node == target:
            paths.append(path)
        for neighbor in graph[node]:
            if neighbor not in path:
                queue.append(path + [neighbor])
    return paths

graph = {'A': ['B', 'C'], 'B': ['D'], 'C': ['D'], 'D': ['E'], 'E': []}
print(bfs_all_paths(graph, 'A', 'E'))
```

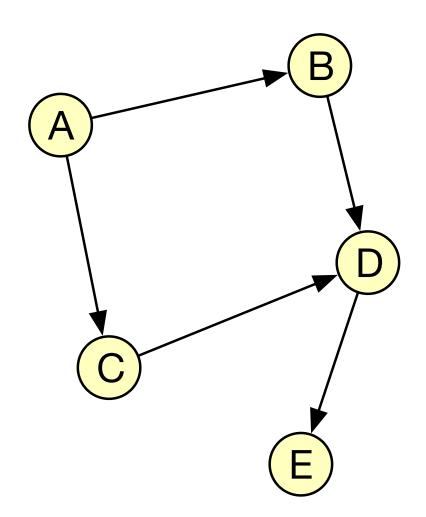




- $V = \{A, B, C, D, E\}$
- $E = \{(A, B), (A, C), (B, D), (C, D), (D, E)\}$
- 2. Use **DFS** starting from vertex A to list the sequence of all visited vertices during the traversal.

```
def dfs(graph, start, visited=None):
    if visited is None:
        visited = []
    visited.append(start)
    for neighbor in graph[start]:
        if neighbor not in visited:
            dfs(graph, neighbor, visited)
    return visited

graph = {'A': ['B', 'C'], 'B': ['D'], 'C': ['D'], 'D': ['E'], 'E': []}
print(dfs(graph, 'A'))
```





You are given a 2D grid of size 5×5 . Each cell in the grid represents a location, and you can either move up, down, left, or right between adjacent cells, provided that the cell is passable (0 and 1 represent road and obstacle, respectively).

0	0	0	1	0
1	1	0	1	0
0	0	0	0	0
0	1	1	1	0
0	0	0	1	0

Your task is to find the shortest path from a given start cell (green) to a goal cell (red) while avoiding obstacles. Specifically, we use the following grid to represent the grid and (x, y) to represent the location.



You are given a 2D grid of size 5×5 . Each cell in the grid represents a location, and you can either move up, down, left, or right between adjacent cells, provided that the cell is passable (0 and 1 represent road and obstacle, respectively).

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0	0	0	1	0
1	1	0	1	0
0	0	0	0	0
0	1	1	1	0
0	0	0	1	0

- Vertex:
 - state/position
- Edge:
 - move between positions



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0	0	0	1	0
1	1	0	1	0
0	0	0	0	0
0	1	1	1	0
0	0	0	1	0

- Vertex:
 - state/position
- Edge:
 - move between positions x

					y _
	0	1	2	3	4
0	(0,0)	(0,1)	(0,2)	(0,3)	(0,4)
1	(1,0)	(1,1)	(1,2)	(1,3)	(1,4)
2	(2,0)	(2,1)	(2,2)	(2,3)	(2,4)
3	(3,0)	(3,1)	(3,2)	(3,3)	(3,4)
√ 4	(4,0)	(4,1)	(4,2)	(4,3)	(4,4)
∇					

$$G = \begin{bmatrix} 0 & 0 & 0 & 1 & 0 \\ 1 & 1 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 0 \\ 0 & 1 & 1 & 1 & 0 \\ 0 & 0 & 0 & 1 & 0 \end{bmatrix}_{5 \times 5}$$

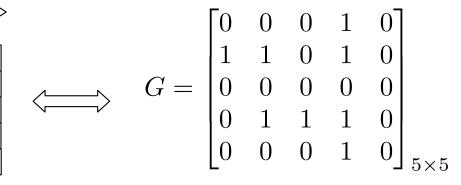


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0	0	0	1	0
1	1	0	1	0
0	0	0	0	0
0	1	1	1	0
0	0	0	1	0

- Vertex:
 - States/positions: (x, y)
 - Root node ???
 - Goal state ???
- Edge:
 - move between positions: ???

					У
	0	1	2	3	4
0	(0,0)	(0,1)	(0,2)	(0,3)	(0,4)
1	(1,0)	(1,1)	(1,2)	(1,3)	(1,4)
2	(2,0)	(2,1)	(2,2)	(2,3)	(2,4)
3	(3,0)	(3,1)	(3,2)	(3,3)	(3,4)
4	(4,0)	(4,1)	(4,2)	(4,3)	(4,4)



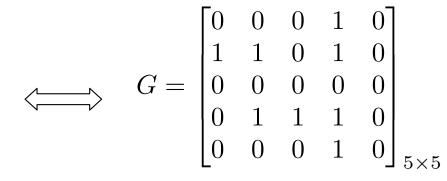


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0	0	0	1	0
1	1	0	1	0
0	0	0	0	0
0	1	1	1	0
0	0	0	1	0

- Vertex:
 - States/positions: (x, y)
 - Root node --> initial state --> (0, 0)
 - Goal state --> (4, 4)
- Edge:
 - move between positions: ???

					У
	0	1	2	3	4
0	(0,0)	(0,1)	(0,2)	(0,3)	(0,4)
1	(1,0)	(1,1)	(1,2)	(1,3)	(1,4)
2	(2,0)	(2,1)	(2,2)	(2,3)	(2,4)
3	(3,0)	(3,1)	(3,2)	(3,3)	(3,4)
4	(4,0)	(4,1)	(4,2)	(4,3)	(4,4)



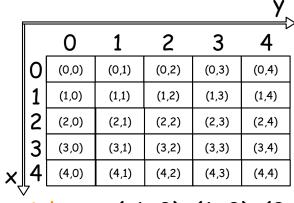


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0	0	0	1	0
1	1	0	1	0
0	0	0	0	0
0	1	1	1	0
0	0	0	1	0

Required to use graph algorithm, but where is the graph?

- Vertex:
 - States/positions: (x, y)
 - Root node --> initial state --> (0, 0)
 - Goal state --> (4, 4)
- Edge:



$$G = \begin{bmatrix} 0 & 0 & 0 & 1 & 0 \\ 1 & 1 & 0 & 1 & 0 \\ 0 & 0 & 0 & 0 & 0 \\ 0 & 1 & 1 & 1 & 0 \\ 0 & 0 & 0 & 1 & 0 \end{bmatrix}$$

move between positions: Up, down, left, right --> (-1, 0), (1, 0), (0, -1), (0, 1)



You are given a 2D grid of size 5×5 . Each cell in the grid represents a location, and you can either move up, down, left, or right between adjacent cells, provided that the cell is passable (0 and 1 represent road and obstacle, respectively).

0	0	0	1	0
1	1	0	1	0
0	0	0	0	0
0	1	1	1	0
0	0	0	1	0

Can we move arbitrarily?

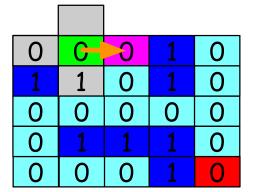
• Edge: move between positions: Up, down, left, right --> (-1, 0), (1, 0), (0, -1), (0, 1)

0	0	0	1	0
1	1	0	1	0
0	0	0	0	0
0	1	1	1	0
0	0	0	1	0



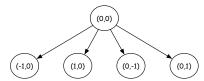
Only one option:

$$(0,1) \rightarrow (0,2)$$





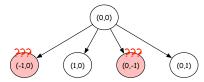
0	0	0	1	0
1	1	0	1	0
0	0	0	0	0
0	1	1	1	0
0	0	0	1	0



- Vertex:
 - States/positions: (x, y)
 - Root node --> initial state --> (0, 0)
 - Goal state --> (4, 4)
- Edge:
 - move between positions:
 - Up, down, left, right
 - (-1, 0), (1, 0), (0, -1), (0, 1)



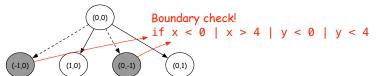
0	0	0	1	0
1	1	0	1	0
0	0	0	0	0
0	1	1	1	0
0	0	0	1	0



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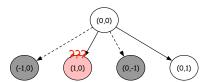
0	0	0	1	0
1	1	0	1	0
0	0	0	0	0
0	1	1	1	0
0	0	0	1	0



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 - Goal state --> (4, 4)
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 - Up, down, left, right
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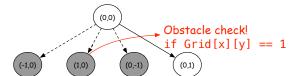
0	0	0	1	0
1	1	0	1	0
0	0	0	0	0
0	1	1	1	0
0	0	0	1	0



- Vertex:
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 - (-1, 0), (1, 0), (0, -1), (0, 1)



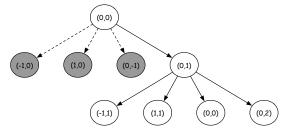
0	0	0	1	0
1	1	0	1	0
0	0	0	0	0
0	1	1	1	0
0	0	0	1	0



- Vertex:
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 - Up, down, left, right
 - (-1, 0), (1, 0), (0, -1), (0, 1)



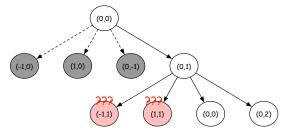
0	0	0	1	0
1	1	0	1	0
0	0	0	0	0
0	1	1	1	0
0	0	0	1	0



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 - Up, down, left, right
 - (-1, 0), (1, 0), (0, -1), (0, 1)



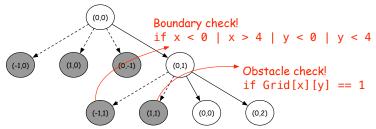
0	0	0	1	0
1	1	0	1	0
0	0	0	0	0
0	1	1	1	0
0	0	0	1	0



- Vertex:
 - States/positions: (x, y)
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- Edge:
 - move between positions:
 - Up, down, left, right
 - (-1, 0), (1, 0), (0, -1), (0, 1)



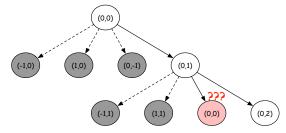
0	0	0	1	0
1	1	0	1	0
0	0	0	0	0
0	1	1	1	0
0	0	0	1	0



- Vertex:
 - States/positions: (x, y)
 - Root node --> initial state --> (0, 0)
 - Goal state --> (4, 4)
- Edge:
 - move between positions:
 - Up, down, left, right
 - (-1, 0), (1, 0), (0, -1), (0, 1)



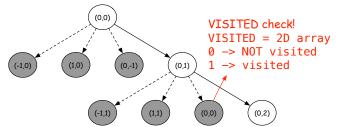
0	0	0	1	0
1	1	0	1	0
0	0	0	0	0
0	1	1	1	0
0	0	0	1	0



- Vertex:
 - States/positions: (x, y)
 - Root node --> initial state --> (0, 0)
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 - Up, down, left, right
 - (-1, 0), (1, 0), (0, -1), (0, 1)



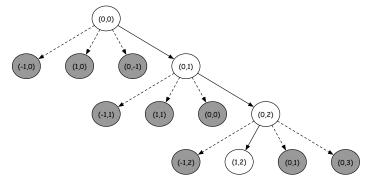
0	0	0	1	0
1	1	0	1	0
0	0	0	0	0
0	1	1	1	0
0	0	0	1	0



- Vertex:
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 - Up, down, left, right
 - (-1, 0), (1, 0), (0, -1), (0, 1)



0	0	0	1	0
1	1	0	1	0
0	0	0	0	0
0	1	1	1	0
0	0	0	1	0

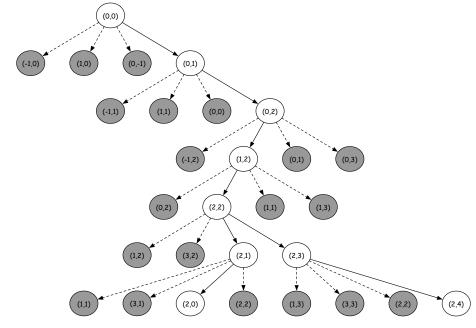


- Vertex:
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 - (-1, 0), (1, 0), (0, -1), (0, 1)



0	0	0	1	0
1	1	0	1	0
0	0	0	0	0
0	1	1	1	0
0	0	0	1	0

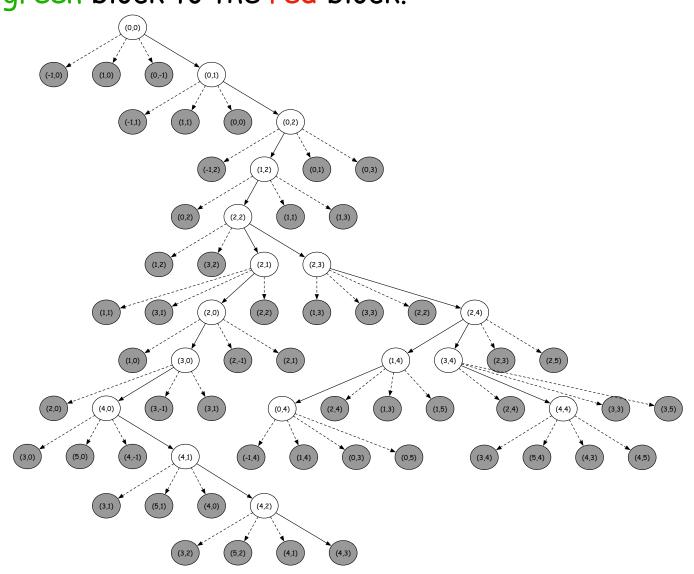
- Vertex:
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 - Root node --> initial state --> (0, 0)
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 - Up, down, left, right
 - (-1, 0), (1, 0), (0, -1), (0, 1)





0	0	0	1	0
1	1	0	1	0
0	0	0	0	0
0	1	1	1	0
0	0	0	1	0

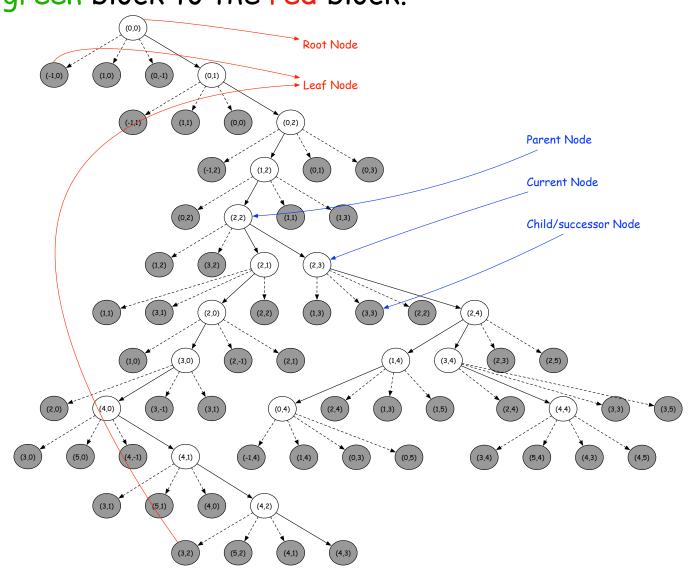
- Vertex:
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0	0	0	1	0
1	1	0	1	0
0	0	0	0	0
0	1	1	1	0
0	0	0	1	0

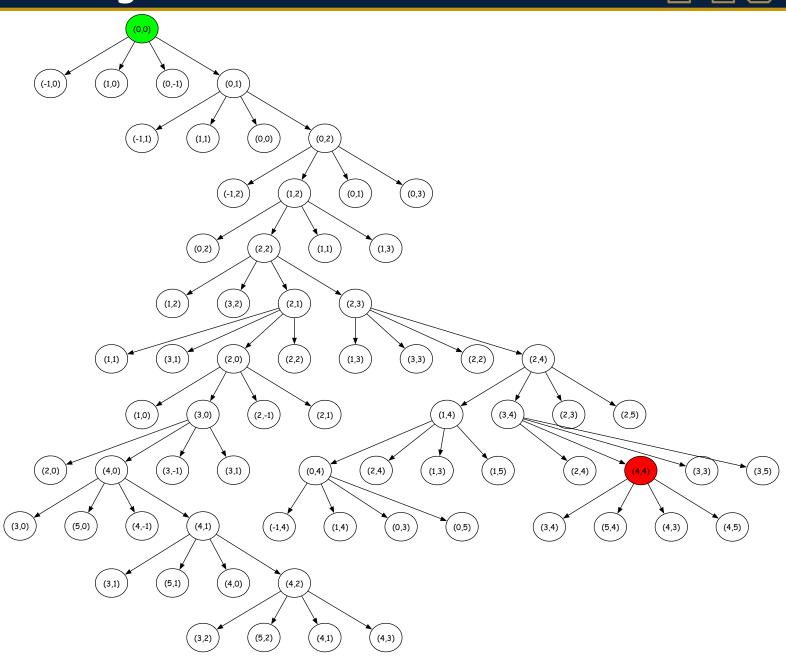
- Representing the state/search space is the first step toward solving the problem.
- The state space comprises all possible states and can be represented as a graph.





Given an directed graph G = (V, E):

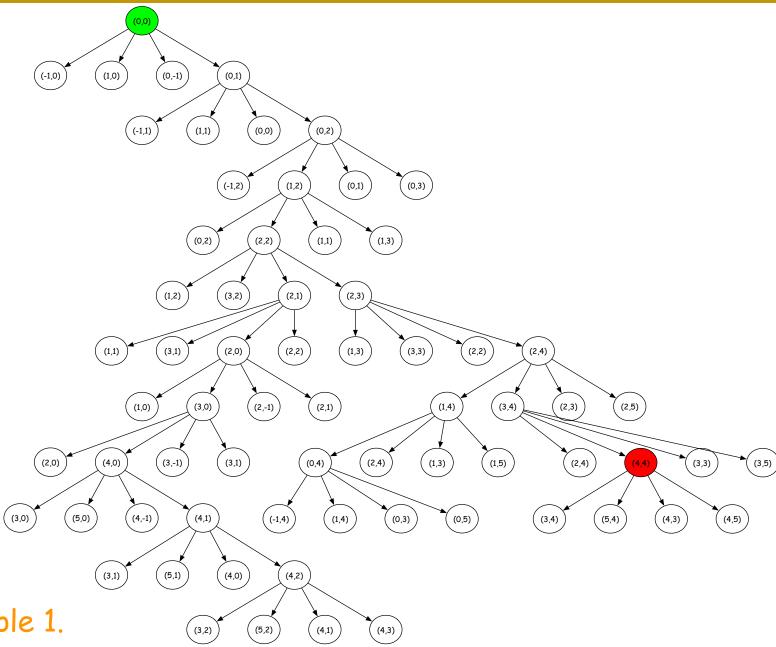
- 1. Use Depth-First Search (DFS) starting from vertex (0,0) and output the order of visited vertices.
- Use Breadth-First Search
 (BFS) starting from vertex
 (4,4) and output the order of visited vertices.





Given an directed graph G = (V, E):

- 1. Use **Depth-First Search (DFS)**starting from vertex (0,0) and output the order of visited vertices.
- 2. Use Breadth-First Search
 (BFS) starting from vertex
 (4,4) and output the order of visited vertices.



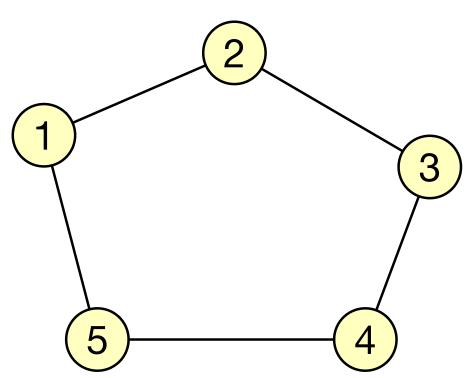
Basically, it's the same as Example 1.

Example 1: Vertex Traversal in an Undirected Graph



Given an undirected graph G = (V, E), where:

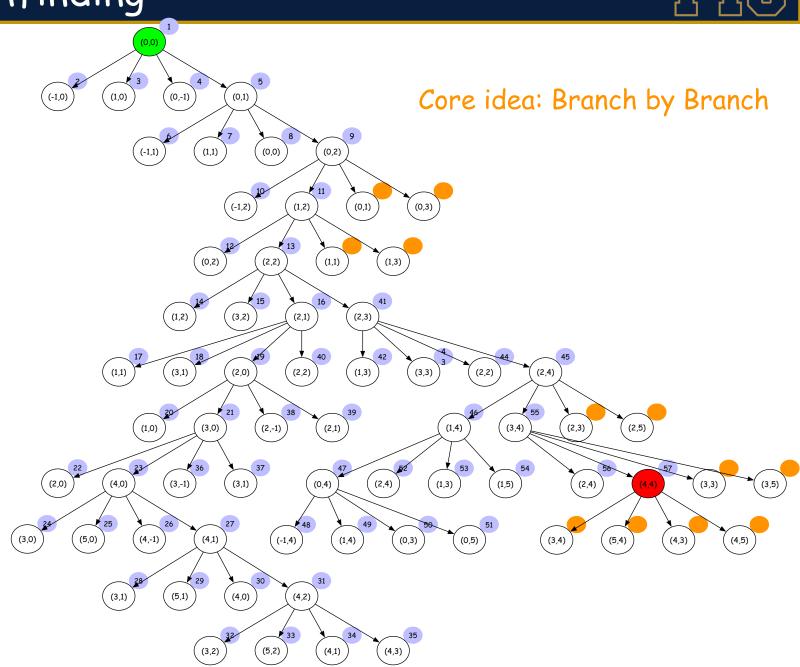
- $V = \{1,2,3,4,5\}$
- $E = \{(1,2), (1,5), (2,3), (3,4), (4,5)\}$
- 1. Use $Depth-First\ Search\ (DFS)$ starting from vertex 1 and output the order of visited vertices.
- 2. Use Breadth-First Search (BFS) starting from vertex 1 and output the order of visited vertices.





Given an directed graph G = (V, E):

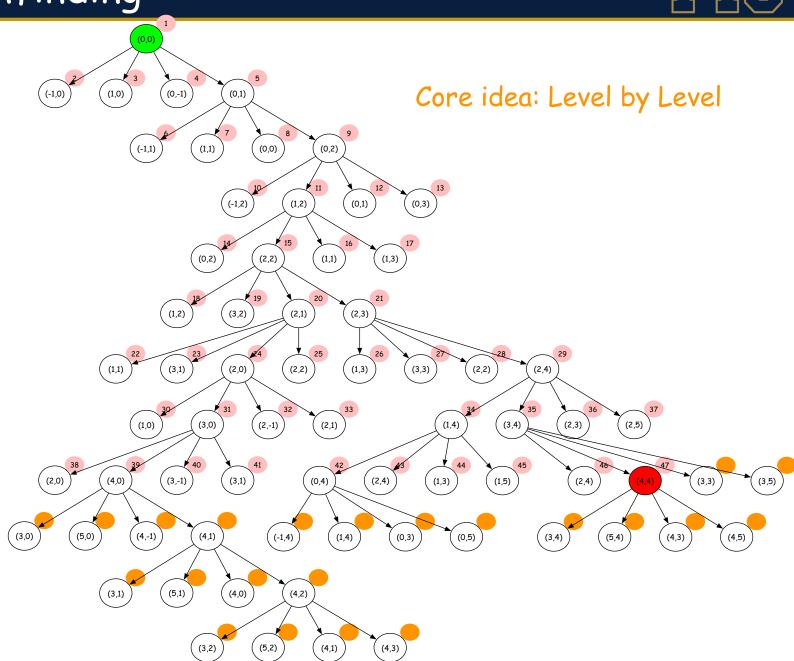
- 1. Use Depth-First Search (DFS) starting from vertex (0,0) and output the order of visited vertices.
- 2. Use Breadth-First Search
 (BFS) starting from vertex
 (4,4) and output the order of visited vertices.





Given an directed graph G = (V, E):

- 1. Use **Depth-First Search (DFS)**starting from vertex (0,0) and output the order of visited vertices.
- Use Breadth-First Search
 (BFS) starting from vertex
 (4,4) and output the order of visited vertices.

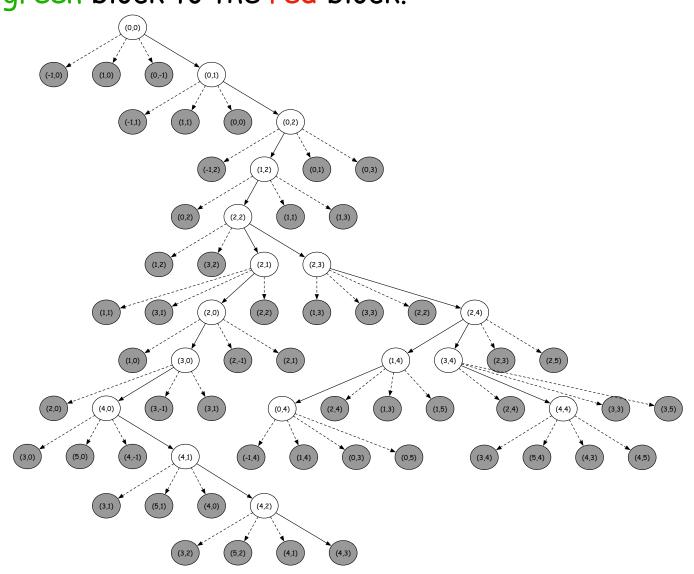




• Given the following Grid, where 0 and 1 represent road and obstacle, respectively, try to find the shortest path from the green block to the red block.

0	0	0	1	0
1	1	0	1	0
0	0	0	0	0
0	1	1	1	0
0	0	0	1	0

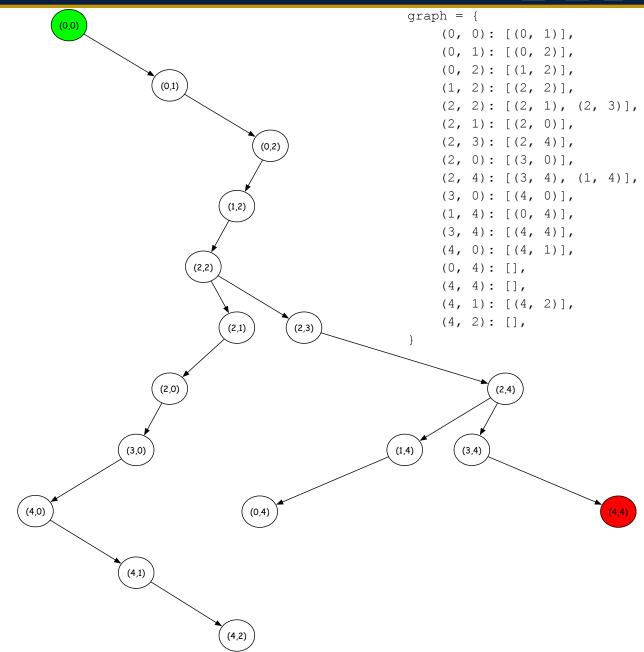
- Vertex:
 - States/positions: (x, y)
 - Root node --> initial state --> (0, 0)
 - Goal state --> (4, 4)
- Edge:
 - move between positions:
 - Up, down, left, right
 - (-1, 0), (1, 0), (0, -1), (0, 1)





$$G = (V, E)$$
:

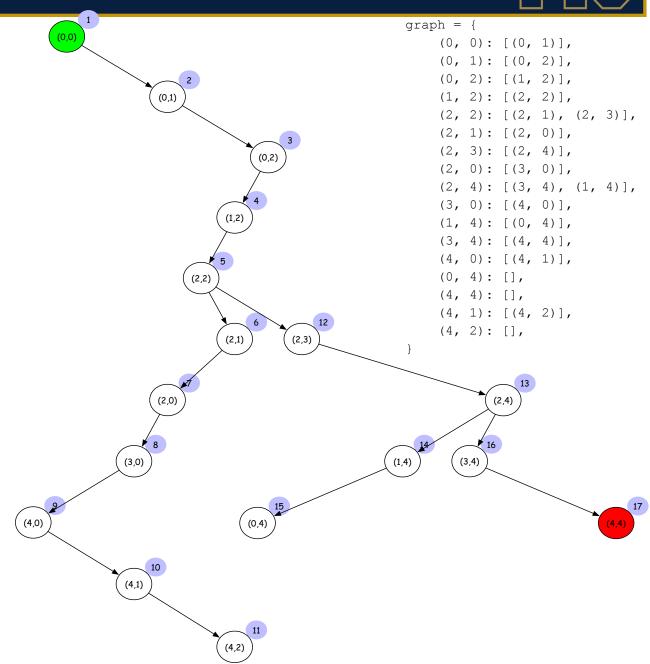
- 1. Use **Depth-First Search (DFS)**starting from vertex (0,0) and output the order of visited vertices.
- 2. Use Breadth-First Search
 (BFS) starting from vertex
 (4,4) and output the order of visited vertices.





$$G = (V, E)$$
:

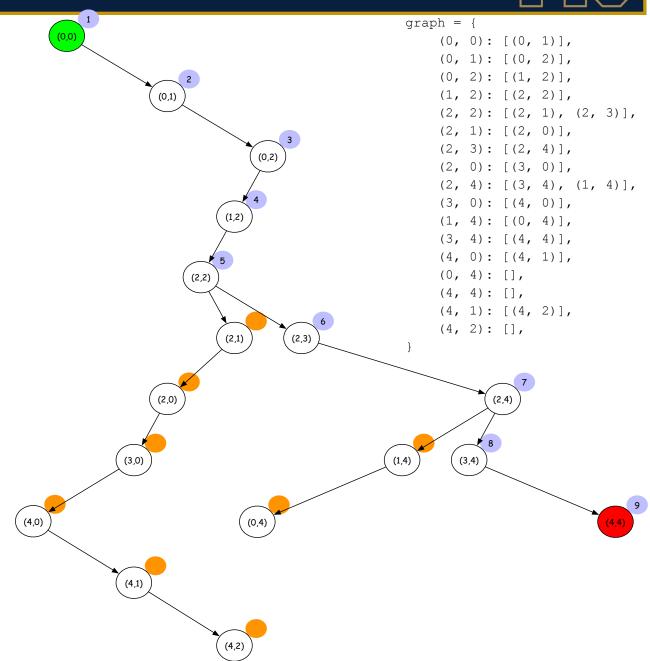
- 1. Use Depth-First Search (DFS)
 - starting from vertex (0,0) and output the order of visited vertices.
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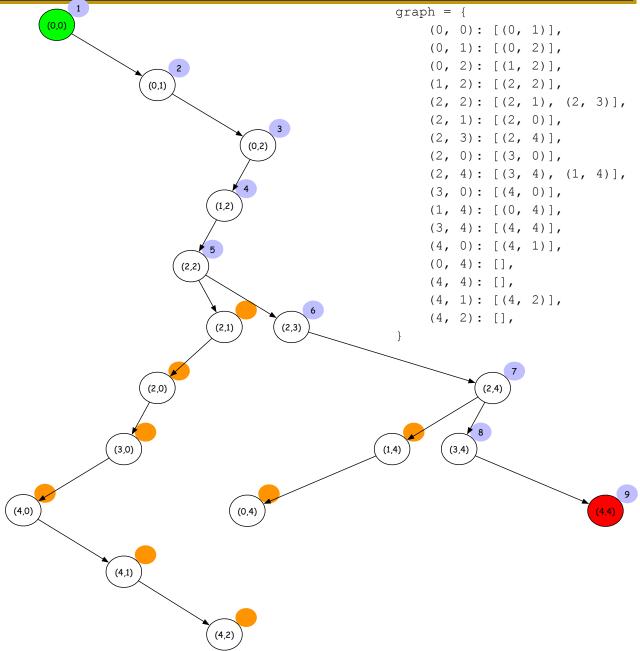


$$G = (V, E)$$
:

- 1. Use Depth-First Search (DFS)
 - starting from vertex (0,0) and output the order of visited vertices.
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 (BFS) starting from vertex
 (4,4) and output the order of visited vertices.



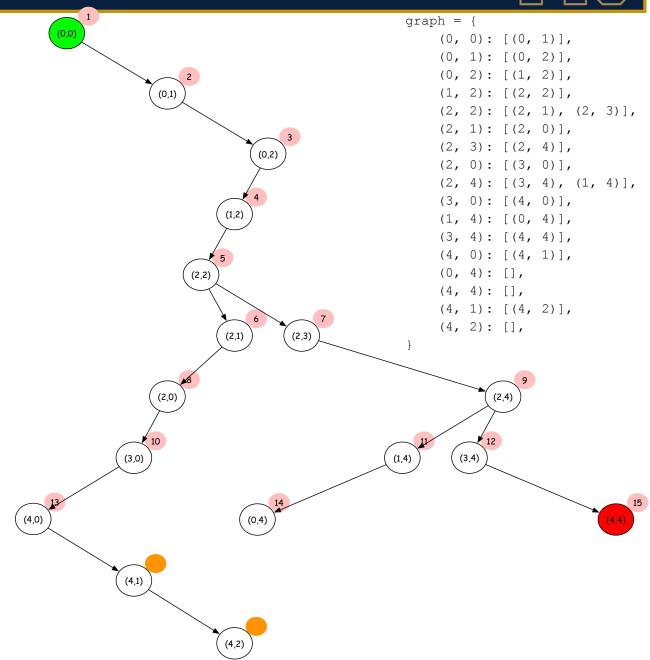
```
def dfs (graph, start, goal, path=None,
   visited=None):
    if path is None:
        path = [start]
    if visited is None:
        visited = set()
    if start == goal:
        return path
    visited.add(start)
    for neighbor in graph.get(start, []):
        if neighbor not in visited:
            result = dfs(graph, neighbor,
      goal, path + [neighbor], visited)
            if result:
                return result
    return None
```





$$G = (V, E)$$
:

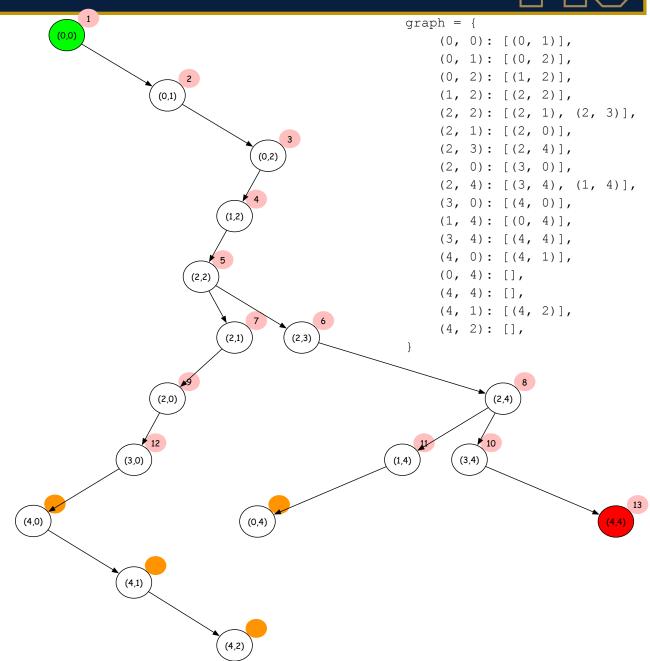
- 1. Use Depth-First Search (DFS) starting from vertex (0,0) and
 - output the order of visited vertices.
- 2. Use Breadth-First Search
 (BFS) starting from vertex
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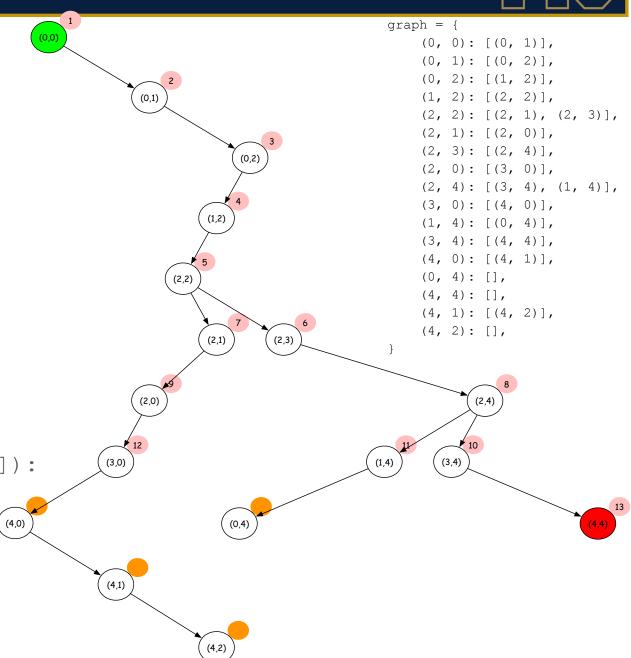


$$G = (V, E)$$
:

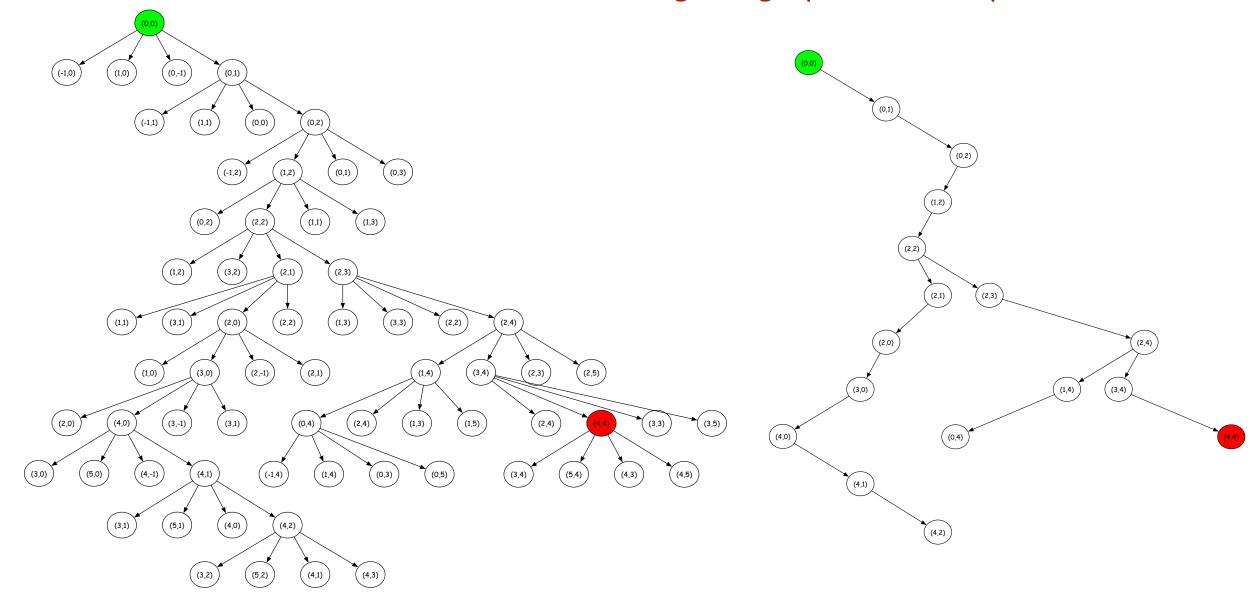
- 1. Use Depth-First Search (DFS)
 - starting from vertex (0,0) and output the order of visited vertices.
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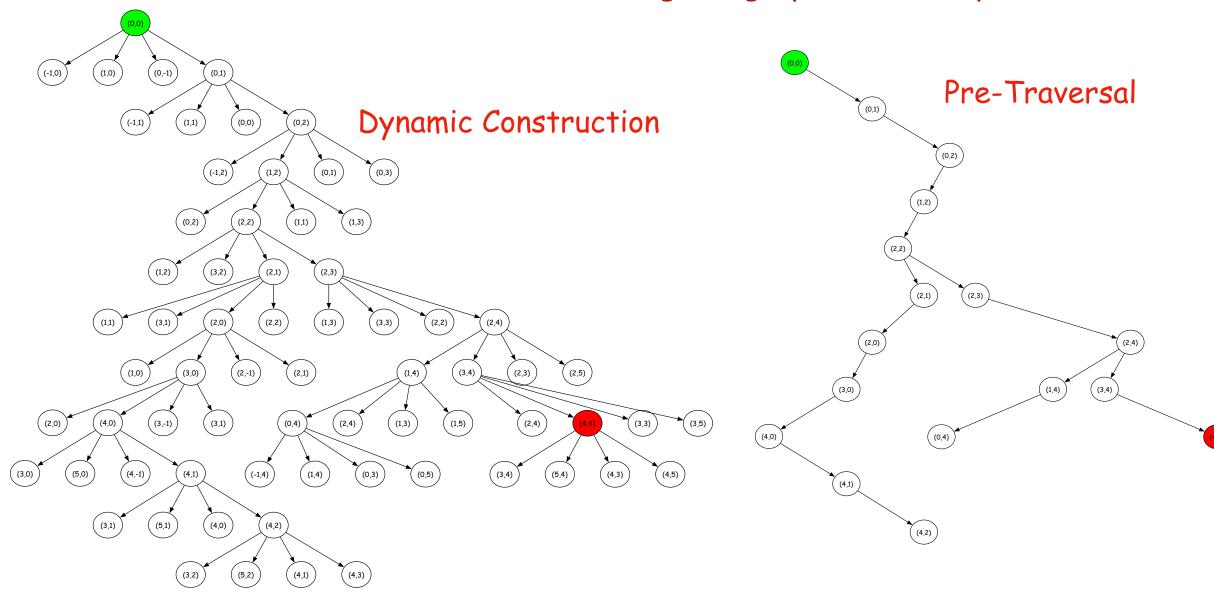
```
from collections import deque
def bfs(graph, start, goal):
    queue = deque([[start]])
    visited = set()
    while queue:
        path = queue.popleft()
        node = path[-1]
        if node == goal:
            return path
        if node not in visited:
            visited.add(node)
            for neighbor in graph.get(node, []):
                new path = path + [neighbor]
                queue.append(new path)
    return None
```



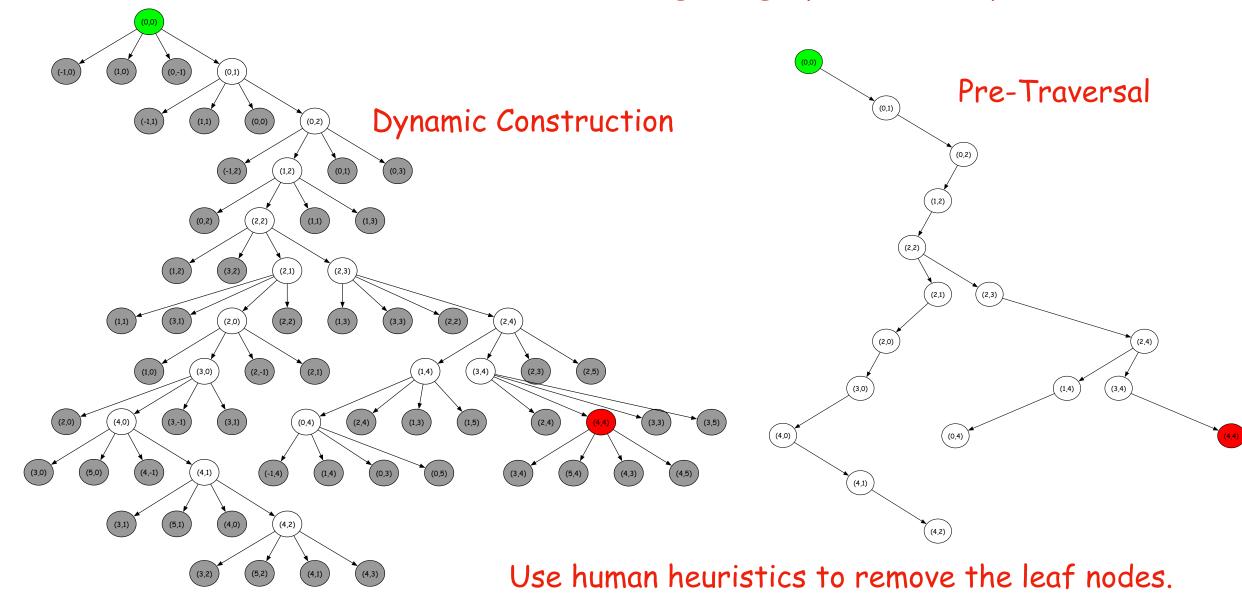








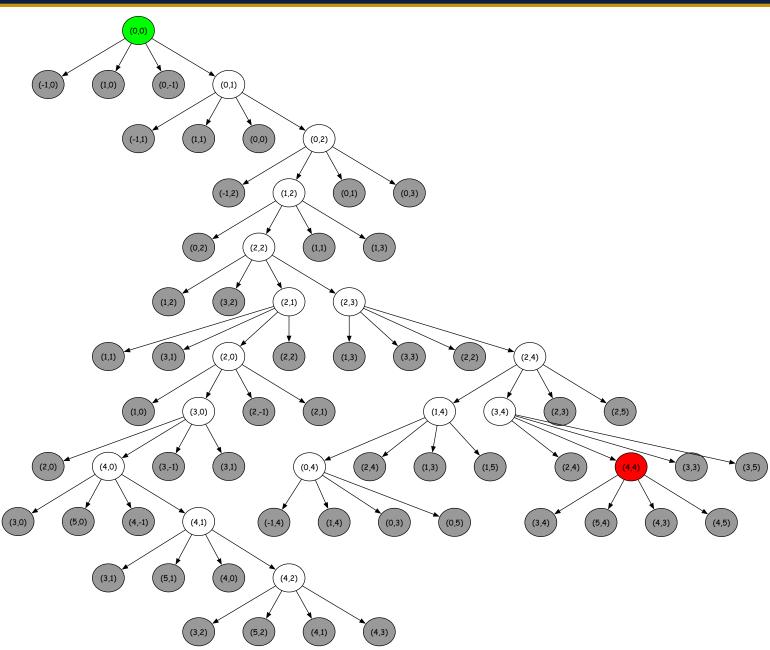






In practice, we could only follow the logic of the graph/tree structure instead of constructing the graph/tree explicitly.

In short, the "graph/tree" is a virtual logical structure only.





Pseudocode of BFS

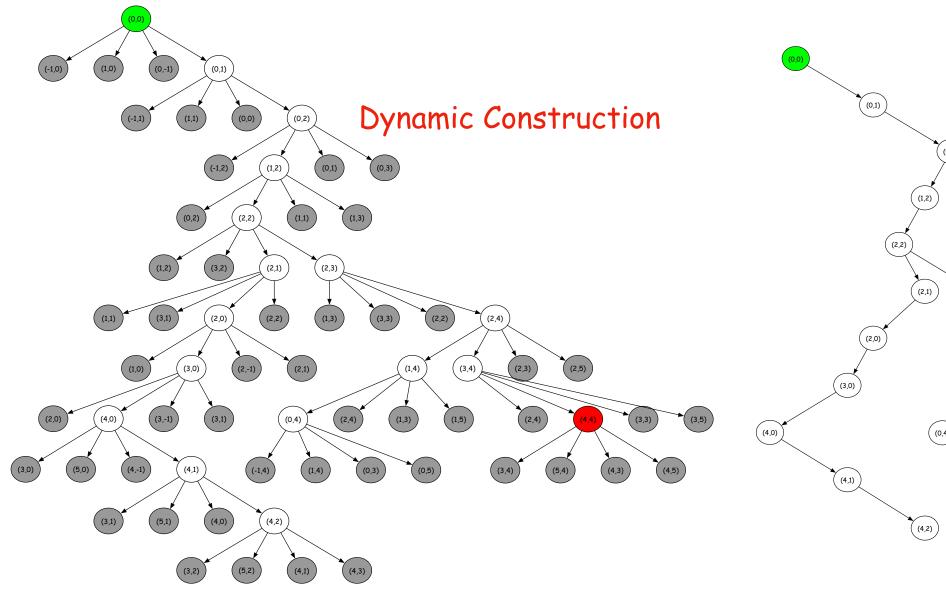
```
BFS Shortest Path (Grid, Start, Goal):
   Initialize a queue Q
    Initialize a 2D array 'visited' with False values
   Add (Start x, Start y, 0) to Q // (x, y, steps)
   Mark Start as visited
   while Q is not empty:
        (x, y, steps) = Q.dequeue()
        // If goal is reached, return steps
        if (x, y) == Goal:
           return steps
        // Explore neighbors in all directions (up, down, left, right)
        for each (dx, dy) in [(0, 1), (0, -1), (1, 0), (-1, 0)]:
            new x = x + dx
           new y = y + dy
            // Check if the move is valid (within grid bounds, not visited, not blocked)
            if is valid(new x, new y) and not visited[new x][new y]:
                Mark (new x, new y) as visited
                Add (new x, new y, steps + 1) to Q
   // If the goal cannot be reached, return -1
   return -1
```

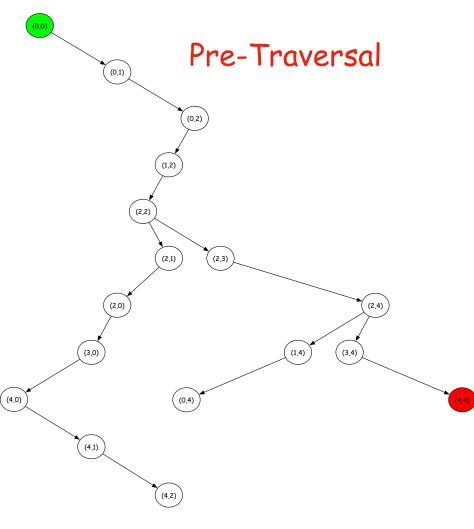


Pseudocode of DFS

```
DFS Shortest Path Stack(Grid, Start, Goal):
   Initialize a stack S
   Initialize a 2D array 'visited' with False values
   Add (Start x, Start y, 0) to S // (x, y, steps)
   Set min steps = infinity
   Set found = False
   while S is not empty:
        (x, y, steps) = S.pop()
       // If goal is reached, update min steps
       if (x, y) == Goal:
           min steps = min(min steps, steps)
           found = True
            continue
       // Mark the current cell as visited
       Mark (x, y) as visited
       // Explore neighbors in all directions (up, down, left, right)
       for each (dx, dy) in [(0, 1), (0, -1), (1, 0), (-1, 0)]:
           new x = x + dx
           new y = y + dy
           // Check if the move is valid (within grid bounds, not visited, not blocked)
           if is valid(new x, new y) and not visited[new x][new y]:
                Add (new x, new y, steps + 1) to S
   // If the goal was reached, return min steps
   if found:
       return min steps
   else:
         // use -1 to denote there is no solution
        return -1
```









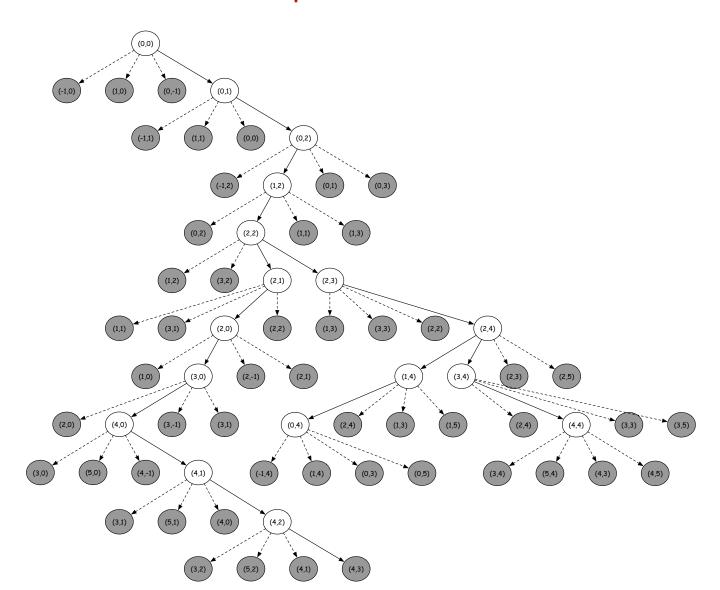
```
BFS Shortest Path (Grid, Start, Goal):
                                                                 from collections import deque
   Initialize a queue Q
   Initialize a 2D array 'visited' with False values
                                                                 def bfs(graph, start, goal):
   Add (Start x, Start y, 0) to Q // (x, y, steps)
                                                                     queue = deque([[start]])
   Mark Start as visited
                                                                     visited = set()
   while Q is not empty:
                                                                     while queue:
       (x, y, steps) = Q.dequeue()
                                                                          path = queue.popleft()
                                                                          node = path[-1]
      if (x, y) == Goal:
           return steps
                                                                          if node == goal:
      for each (dx, dy) in [(0, 1), (0, -1), (1, 0), (-1, 0)]:
                                                                              return path
           new x = x + dx
           new y = y + dy
                                                                          if node not in visited:
                                                                              visited.add(node)
          if is valid(new x, new y) and not visited[new x][new y]:
                                                                               for neighbor in graph.get(node, []):
               Mark (new x, new y) as visited
                                                                                   new path = path + [neighbor]
              Add (new x, new y, steps + 1) to Q
                                                                                   queue.append(new path)
                                                                      return None
   return -1
```



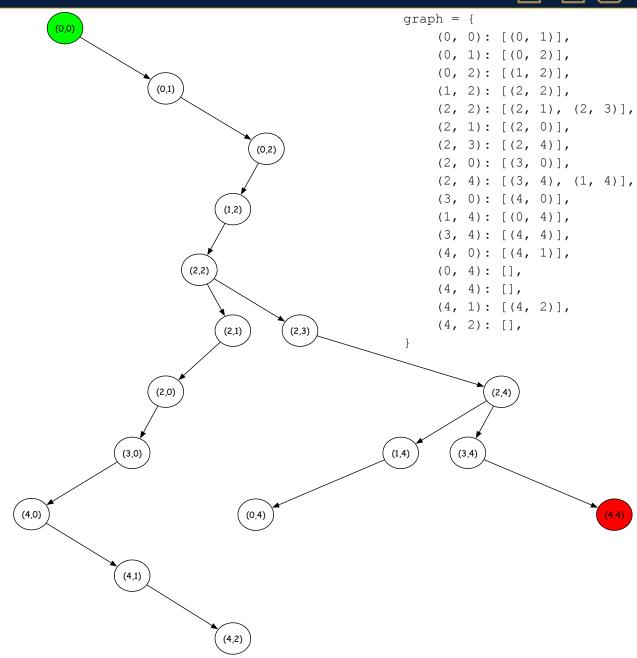
Do we have other possible graph/tree structure for this problem?

0	0	0	1	0
1	1	0	1	0
0	0	0	0	0
0	1	1	1	0
0	0	0	1	0

- Vertex:
 - States/positions: (x, y)
 - Root node --> initial state --> (0, 0)
 - Goal state --> (4, 4)
- Edge:
 - move between positions:
 - Up, down, left, right
 - (-1, 0), (1, 0), (0, -1), (0, 1)



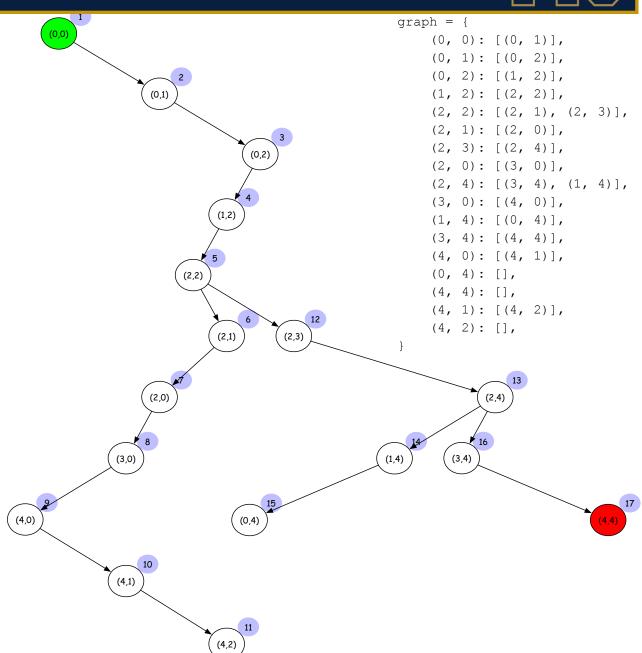
For easy to show the logic, let's start from the constructed simplified tree.





For easy to show the logic, let's start from the constructed simplified tree.

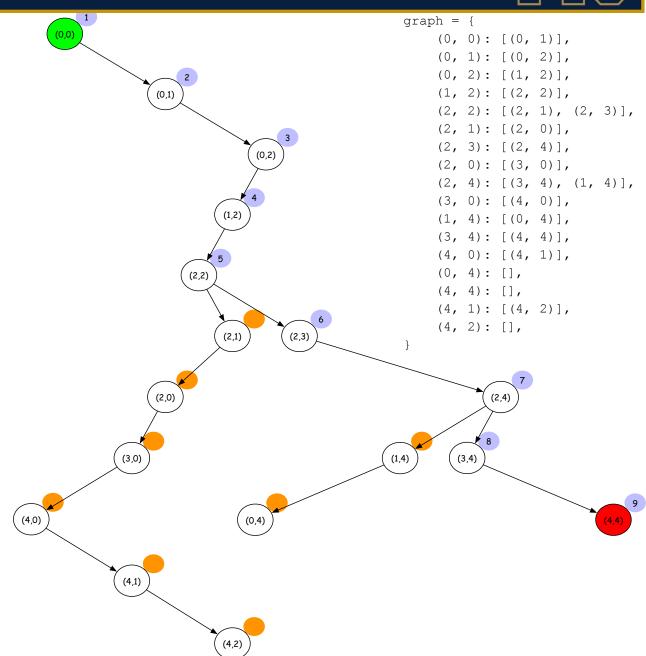
- 1. DFS (Push Right into the Stack first)
 - 17 steps





For easy to show the logic, let's start from the constructed simplified tree.

- 1. DFS (Push Right into the Stack first)
 - 17 steps
- 2. DFS (Push Left into the Stack first)
 - 9 steps





For easy to show the logic, let's start from the constructed simplified tree.

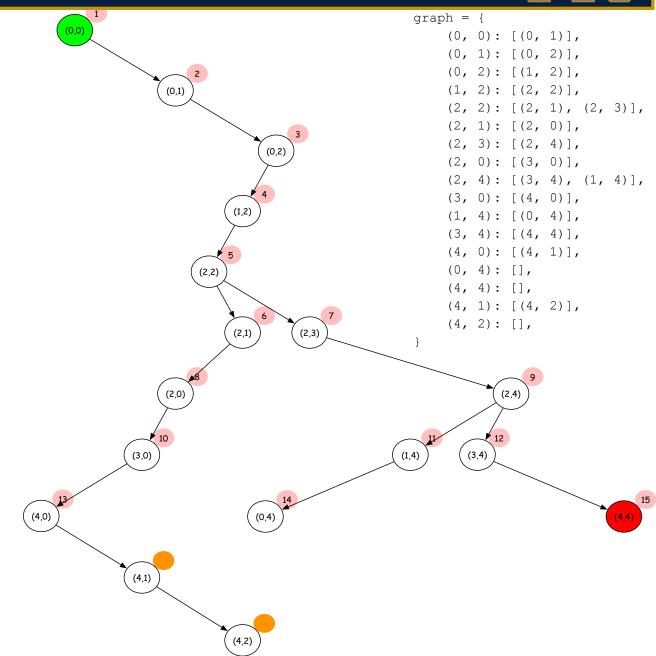
- 1. DFS (Push Right into the Stack first)
 - 17 steps
- 2. DFS (Push Left into the Stack first)
 - 9 steps
- 3. BFS (Push Right into the Queue first)
 - 13 steps

```
graph = {
                                      (0, 0): [(0, 1)],
                                      (0, 1): [(0, 2)],
                                      (0, 2): [(1, 2)],
                                      (1, 2): [(2, 2)],
(0,1)
                                      (2, 2): [(2, 1), (2, 3)],
                                      (2, 1): [(2, 0)],
                                      (2, 3): [(2, 4)],
             (0,2)
                                      (2, 0): [(3, 0)],
                                      (2, 4): [(3, 4), (1, 4)],
                                      (3, 0): [(4, 0)],
                                      (1, 4): [(0, 4)],
                                      (3, 4): [(4, 4)],
                                      (4, 0): [(4, 1)],
                                      (0, 4): [],
                                      (4, 4): [],
                                      (4, 1): [(4, 2)],
                                      (4, 2): [],
                  (2,3)
                                        (3,4)
```



For easy to show the logic, let's start from the constructed simplified tree.

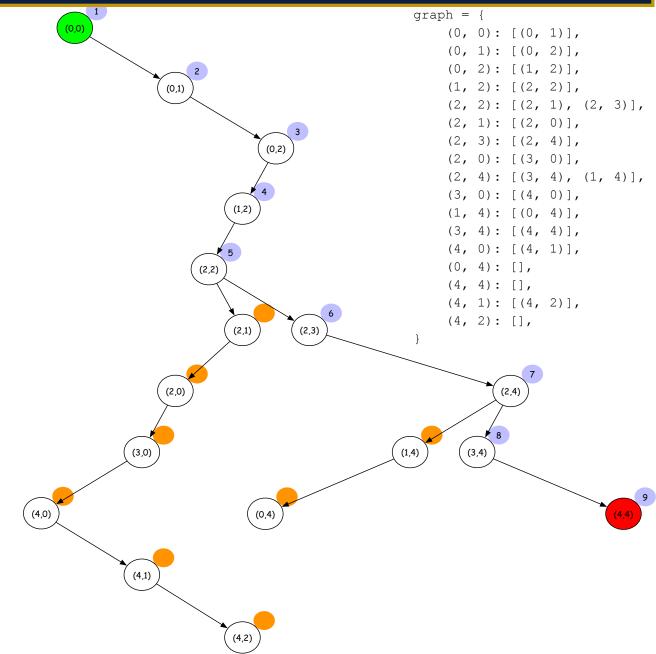
- 1. DFS (Push Right into the Stack first)
 - 17 steps
- 2. DFS (Push Left into the Stack first)
 - 9 steps
- 3. BFS (Push Right into the Queue first)
 - 13 steps
- 4. BFS (Push Left into the Queue first)
 - 15 steps





For easy to show the logic, let's start from the constructed simplified tree.

- 1. DFS (Push Right into the Stack first)
 - 17 steps
- 2. DFS (Push Left into the Stack first)
 - 9 steps
- 3. BFS (Push Right into the Queue first)
 - 13 steps
- 4. BFS (Push Left into the Queue first)
 - 15 steps



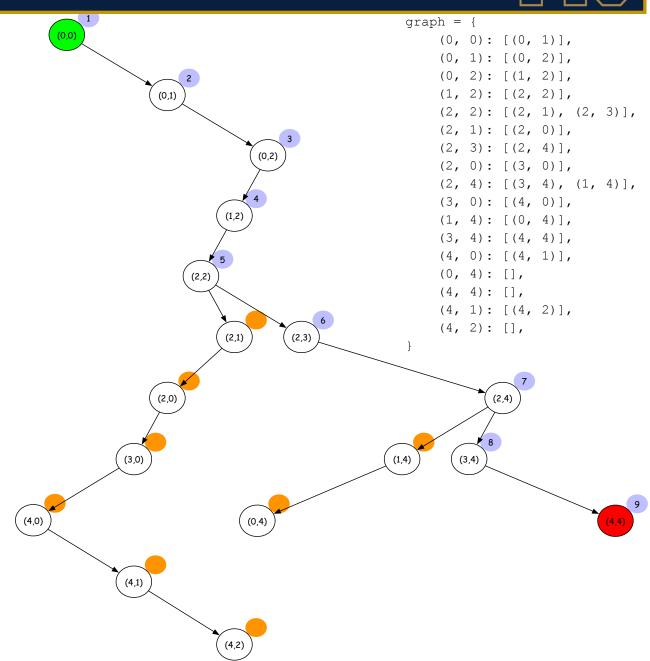


For easy to show the logic, let's start from the constructed simplified tree.

We have tried 4 traversal orders:

- 1. DFS (Push Right into the Stack first)
 - 17 steps
- 2. DFS (Push Left into the Stack first)
 - 9 steps
- 3. BFS (Push Right into the Queue first)
 - 13 steps
- 4. BFS (Push Left into the Queue first)
 - 15 steps

Why we save the running time (with the same time complexity)? Can we gunarantee to reproduce the best case?



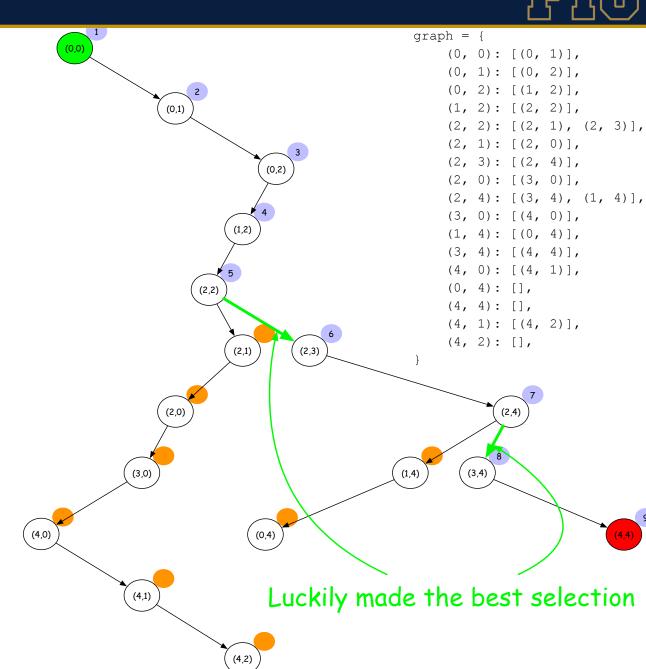


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- 1. DFS (Push Right into the Stack first)
 - 17 steps
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 - 15 steps

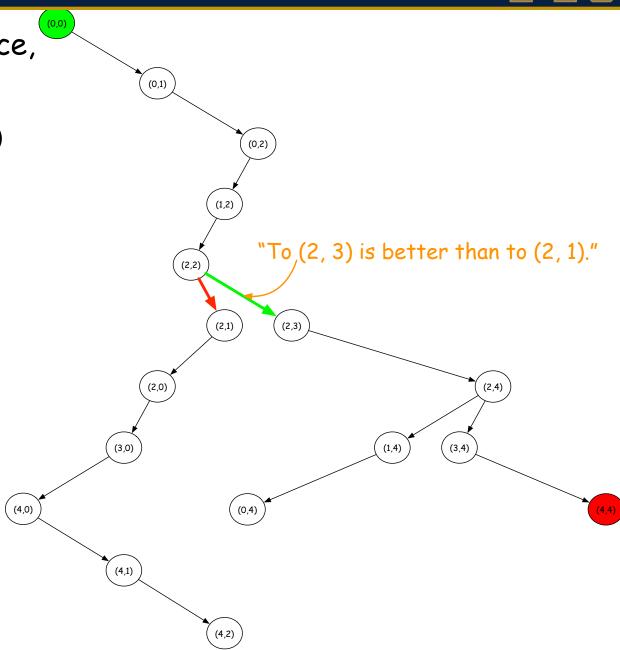
Why we save the running time (with the same time complexity)? Can we gunarantee to reproduce the best case?





During the search, we may need some guidance, e.g.,

((2,2),(2,3)) is better than ((2,2),(2,1))

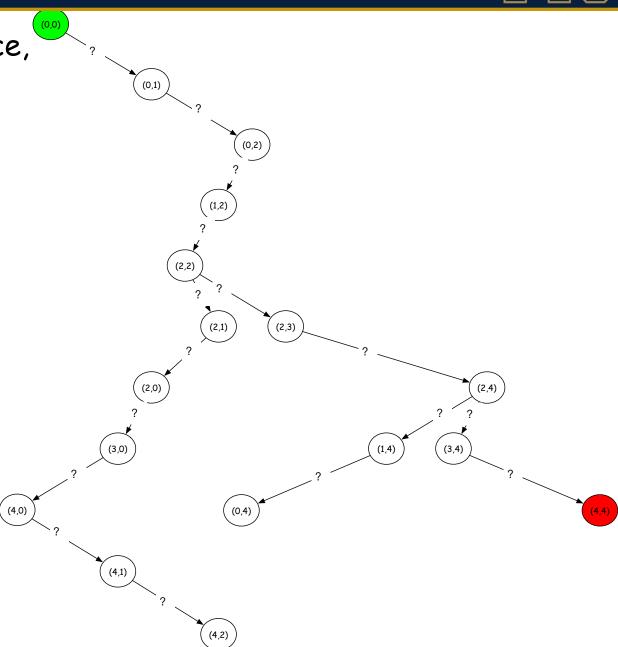




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We may need one WEIGHTED graph.





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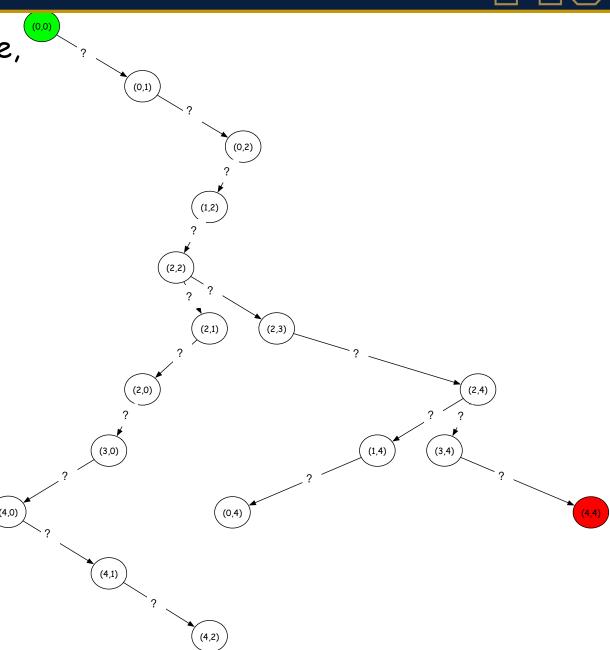
((2,2),(2,3)) is better than ((2,2),(2,1))

We may need one WEIGHTED graph.

To get the proper weight/guidance, we need to find some useful knowledge/heuristic from the specific problem.

0	0	0	1	0
1	1	0	1	0
0	0	0	0	0
0	1	1	1	0
0	0	0	1	0

Any intuition?





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((2,2),(2,3)) is better than ((2,2),(2,1))

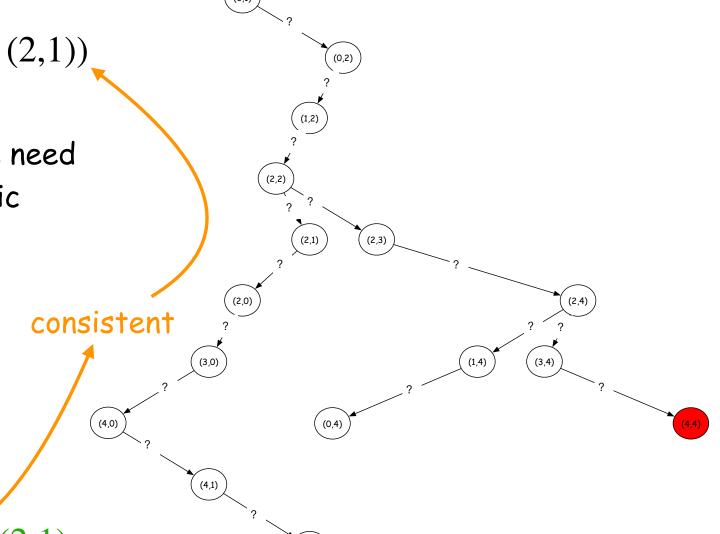
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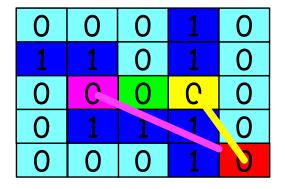
0	0	0	1	0
1	1	0	1	0
0	0	0	0	0
0	1	1	1	0
0	0	0	1	0

Any intuition?

For (2,2), (2,3) should be better than (2,1).





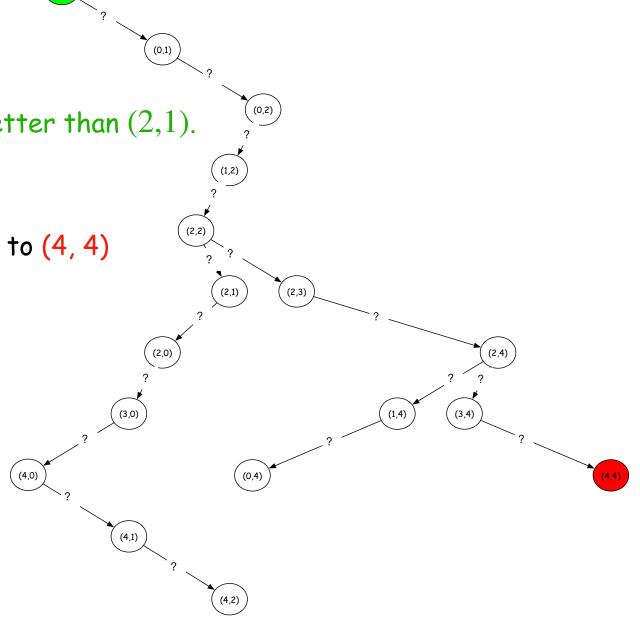


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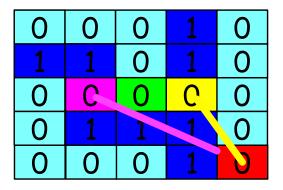
We has this intuition, as the distance from (2, 3) to (4, 4) is shorter than the distance from (2, 1) to (4, 4).

Manhattan distance

- A way to measure the distance between two points in a grid where you can only move along the grid lines, either horizontally or vertically.
- Formula:
 - If you have two points (x_1,y_1) and (x_2,y_2) on a grid, the Manhattan distance d is
 - $d = |x_2 x_1| + |y_2 y_1|$







For (2,2), (2,3) should be better than (2,1).

We has this intuition, as the distance from (2, 3) to (4, 4)is shorter than the distance from (2, 1) to (4, 4).

Manhattan distance

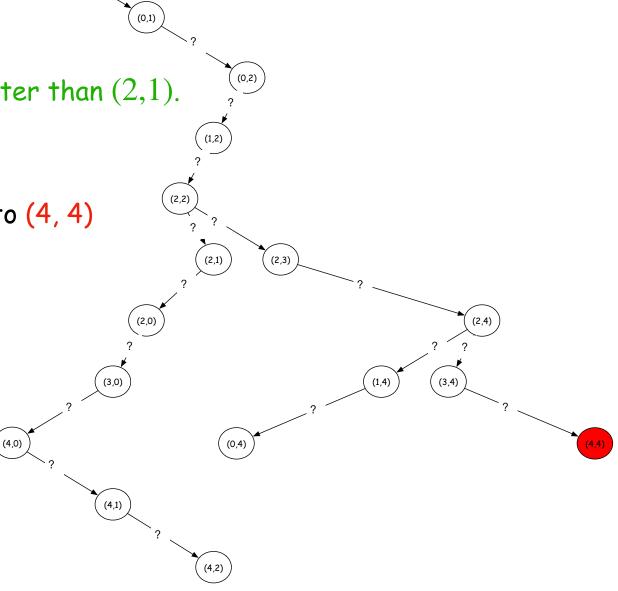
$$d = |x_2 - x_1| + |y_2 - y_1|$$

(0,0)	(0,1)	(0,2)	(0,3)	(0,4)
(1,0)	(1,1)	(1,2)	(1,3)	(1,4)
(2,0)	(2,1)	(2,2)	(2,3)	(2,4)
(3,0)	(3,1)	(3,2)	(3,3)	(3,4)
(4,0)	(4,1)	(4,2)	(4,3)	(4,4)

_					
	(0,0)	(0,1)	(0,2)	(0,3)	(0,4)
	(1,0)	(1,1)	(1,2)	(1,3)	(1,4)
	(2,0)	(2,1)	(2,2)	(2,3)	(2,4)
	(3,0)	(3,1)	(3,2)	(3,3)	(3,4)
	(4,0)	(4,1)	(4,2)	(4,3)	(4,4)

$$d_1 = |3 - 0| + |1 - 0| = 4$$
 $d_2 = |4 - 1| + |3 - 2| = 4$

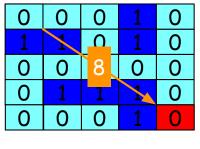
$$d_2 = |4 - 1| + |3 - 2| = 4$$

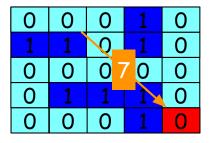


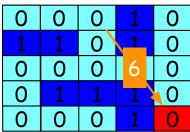


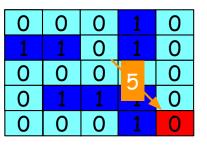
Calculate the distance for all positions as the

weight.









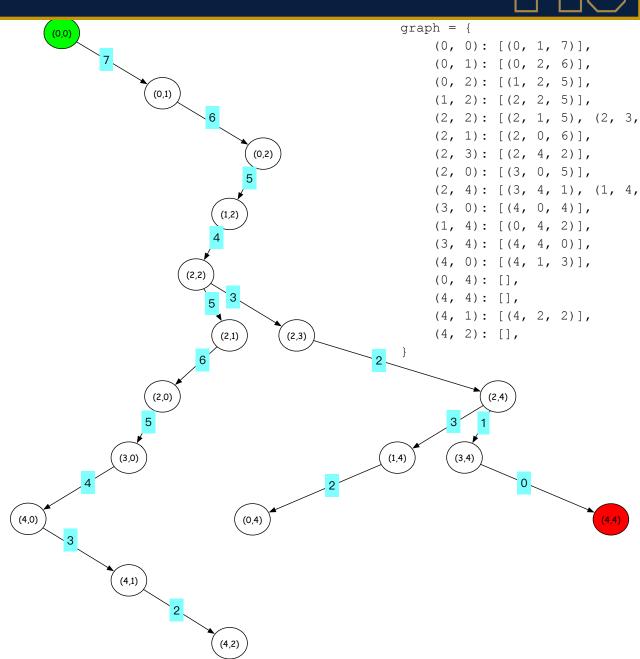
Heuristic function h(n) can return An rational estimate of the cost from node n to the goal. In this example, manhattan distance is h(n).

```
def manhattan_distance(curent_node, target):
    x1, y1 = curent_node
    x2, y2 = target

return abs(x1 - x2) + abs(y1 - y2)
```

Greedy Best-First Search:

• In given weighted graph, just pick the best local option every time.



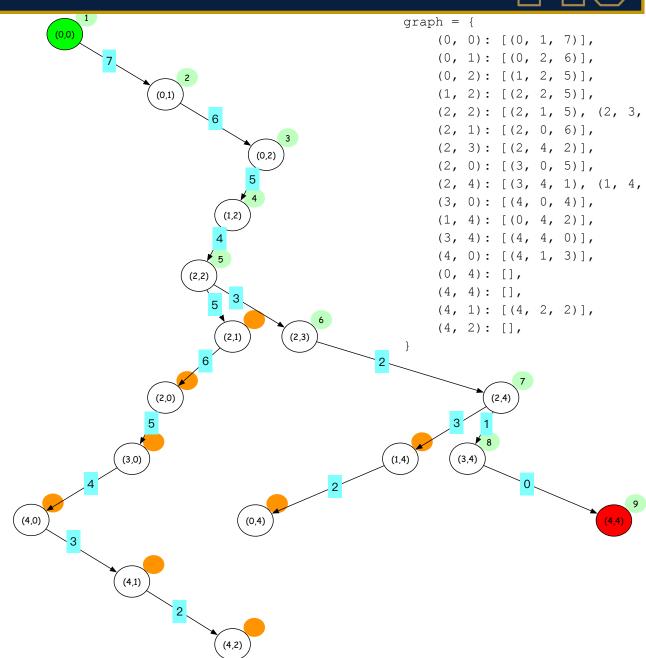


Greedy Best-First Search:

• In given weighted graph, just pick the best local option every time.

We use the Manhattan distance as the estimation to guide the greedy best-first search.

 Sometimes, it works well and find the best solution (lucky).



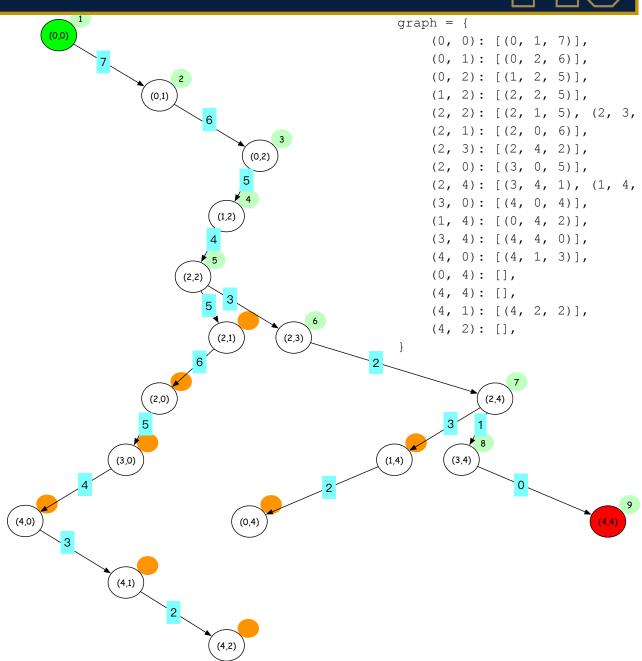


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- Sometimes, it works well and find the best solution (lucky).
- However, the result might be not optimal, and we have no chance to fix it.

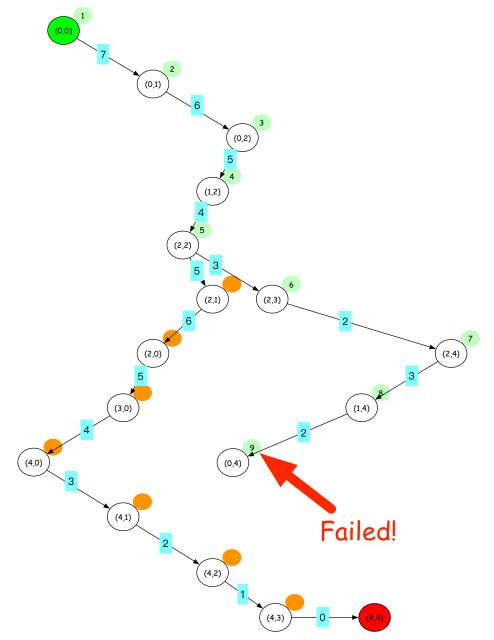




Greedy Best-First Search:

- In given weighted graph, just pick the best local option every time.
- However, the result might be not optimal, and we have no chance to fix it. Consider the example below.

0	0	0	1	0
1	1	0	1	0
0	0	0	0	0
0	1	1	1	1
0	0	0	0	0



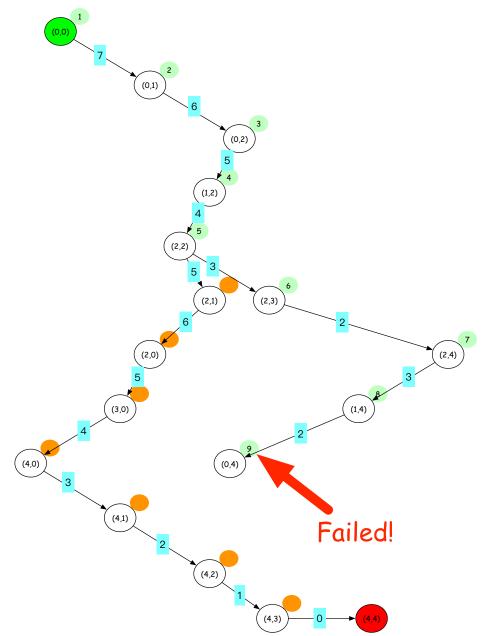


Greedy Best-First Search:

- In given weighted graph, just pick the best local option every time.
- However, the result might be not optimal, and we have no chance to fix it. Consider the example below.

0	0	0	1	0
1	1	0	1	0
0	0	0	0	0
0	1	1	1	1
0	0	0	0	0

Why?



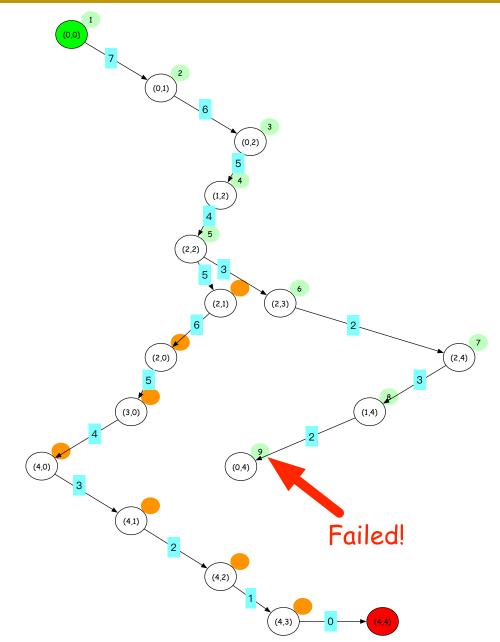


Greedy Best-First Search:

- In given weighted graph, just pick the best local option every time.
- However, the result might be not optimal, and we have no chance to fix it. Consider the example below.

0	0	_	1	0	
1	1	0	1	0	
0	0	0	0	0	
0	1	1	1	1	
0	0	0	0	0	

The core idea of greedy best-first search is to always pursue the path with the smallest estimated future cost. However, the future is inherently uncertain, as the estimation of future costs may be inaccurate. In such cases, greedy best-first search not only risks failing to find the optimal solution but may even fail entirely.





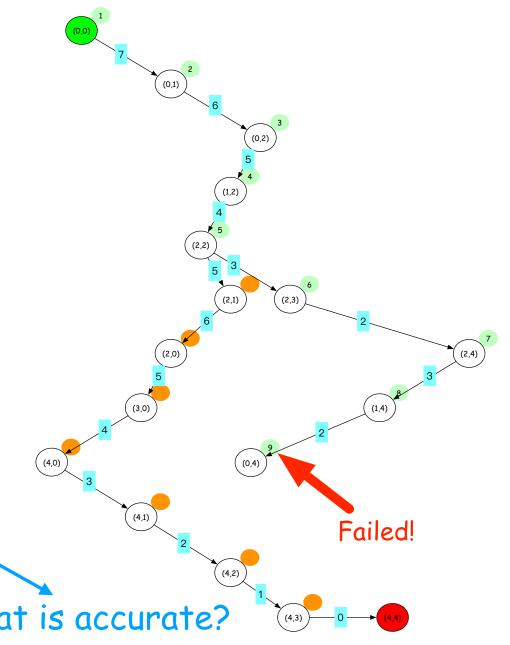
Greedy Best-First Search:

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0	0	0	1	0
1	1	0	1	0
0	0	0	0	0
0	1	1	1	1
0	0	0	0	0

Why?

The core idea of greedy best-first search is to always pursue the path with the smallest estimated future cost. However, the future is inherently uncertain, as the estimation of future costs may be inaccurate. In such cases, greedy best-first search not only risks failing to find the optimal solution but may even fail entirely. What is accurate?

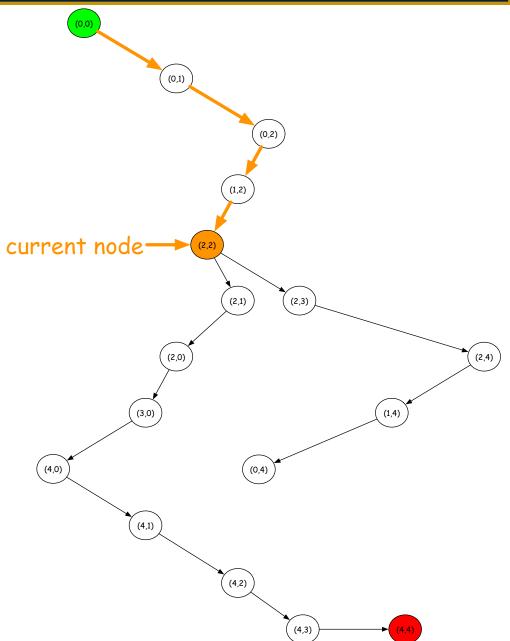




Dijkstra's Algorithm:

 The most accurate information we can obtain is the recorded events that have already occurred, specifically the cost incurred up to the current node.

0	0	0	1	0
1	1	0	1	0
0	0	0	0	0
0	1	1	1	1
0	0	0	0	0





Dijkstra's Algorithm:

 The most accurate information we can obtain is the recorded events that have already occurred, specifically the cost incurred up to the current node

· Weighted graph.

						1.
o†	ain	is t	he			(0,1)
re	ed,					1
	rrer	nt n	ode			(0,2)
						(1,2)
	0	0	0	1	0	
	1	1	0	1	0	(2,2)
	0	0	0	0	0	
	0	1	1	1	1	$(2,1) \qquad (2,3)$
	0	0	0	0	0	(2,4)
						(3,0)
						1
					(4,0	(0,4)
						(4.1)
						1
						(4,2)
						$ \begin{array}{c} 1 \\ (4.3) \\ \end{array} $
						(1,0)



Dijkstra's Algorithm:

- The most accurate information we can obtain is the recorded events that have already occurred, specifically the cost incurred up to the current node.
- · Weighted graph.
- The algorithm uses a priority queue (min-heap) to always expand the node with the smallest current distance.

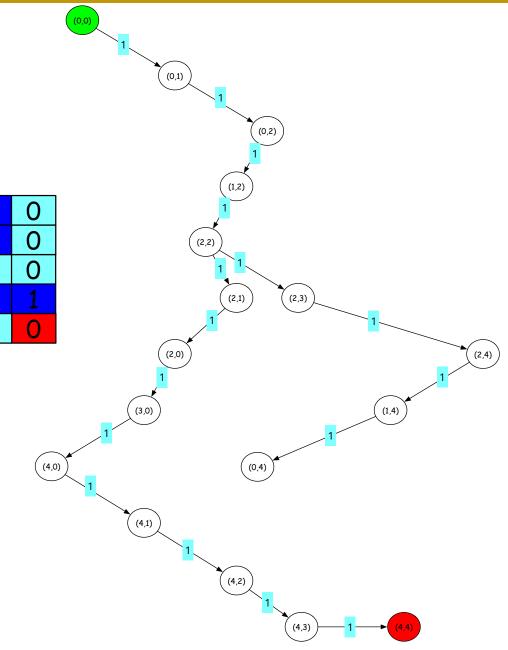
When the current node is (2, 2)

(2, 1, 5)	(2, 3, 5)		

Assume we pop (2, 1)

(2, 3, 5) (2, 0, 6)			
---------------------	--	--	--

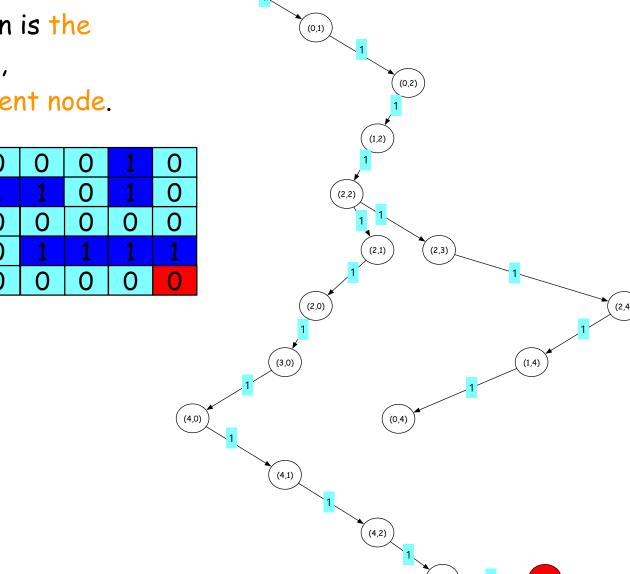
We will pop out (2, 3) then
(It's like an insurance for exit wrong paths)





Dijkstra's Algorithm:

- The most accurate information we can obtain is the recorded events that have already occurred, specifically the cost incurred up to the current node.
- · Weighted graph.
- The algorithm uses a priority queue (min-heap) to always expand the node with the smallest current distance.
- The process ends when the target node is reached or all reachable nodes are processed.
 - if found, the best (Optimality).
 - if reached all, no result.

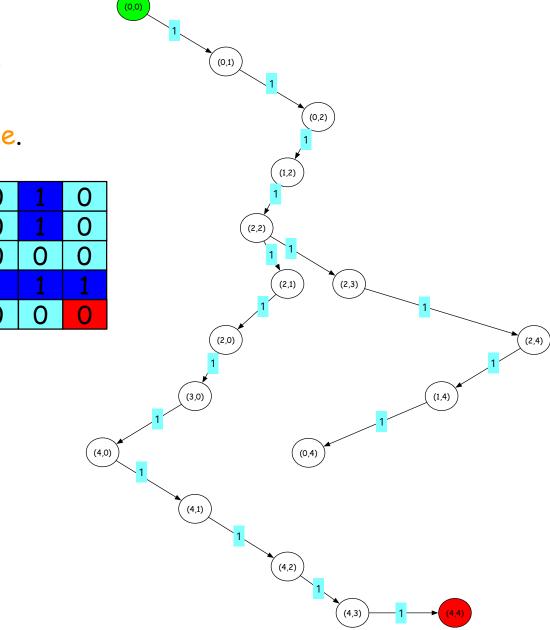




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Past vs. Future g(n) vs. h(n)

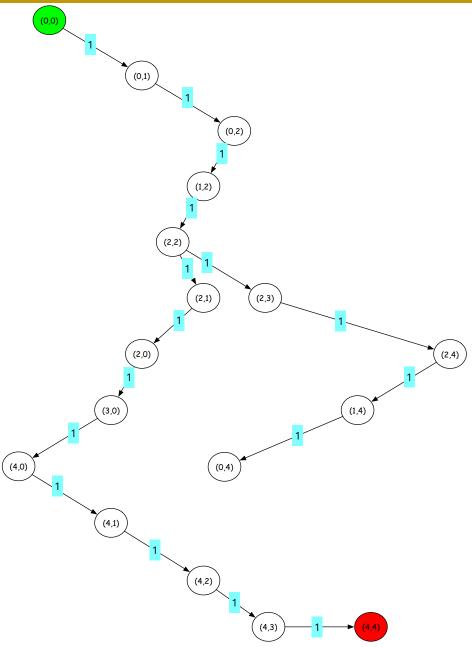




Dijkstra's Algorithm:

 The most accurate information we can obtain is the recorded events that have already occurred, specifically the cost incurred up to the current node.

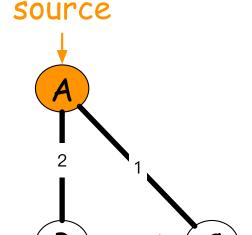
```
Func Dijkstra (graph, source):
    Output: Shortest distances (dist) from source to all
nodes and predecessors (prev)
    for each node v in graph:
        dist[v] = \infty
    dist[source] = 0
    Create a min-priority queue Q
   Add source into the O
    While Q is not empty:
        dist so far, u = Q.pop()
        for each neighbor v of u:
            if dist[u] + weight(u, v) < dist[v]:</pre>
                dist[v] = dist[u] + weight(u, v)
                push v into the Q
    return dist and prev
```





Dijkstra's Algorithm:

- The most accurate information we can obtain is the recorded events that have already occurred, specifically the cost incurred up to the current node.
- (Non-negative) Weighted graph.
 - Greedy Assumption:
 - Dijkstra's algorithm works by always expanding the node with the smallest current distance.
 - Once a node's shortest distance is finalized, the algorithm assumes that no shorter path to this node will ever be found.
 - This assumption is true only if all edge weights are non-negative.



Initiaization

Dis(A)=0, Dis(B)=inf, Dis(C)=inf

Choose A:

Dis(A)=0, Dis(B)=2, Dis(C)=1

Choose C: (do nothing)

Dis(A)=0, Dis(B)=2, Dis(C)=1

Choose B:

$$2 + -3 = -1 < Dis(C)$$



Dijkstra's Algorithm:

- The most accurate information we can obtain is the recorded events that have already occurred, specifically the cost incurred up to the current node.
- Time complexity: $O((|V| + |E|)\log |V|)$
 - When Using a Binary Min-Heap
 - Extracting the minimum node takes $O(\log |V|)$
 - At most V nodes in the heap, $O(|V| \log |V|)$
 - Updating nodes takes $O(\log |V|)$
 - At most relexing |E| edges, $O(|E|\log|V|)$

Relaxation is the process of checking whether a shorter path to a node can be found by passing through another node, and if so, updating the shortest known distance.

```
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    for each node v in graph:
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    dist[source] = 0
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                dist[v] = dist[u] + weight(u, v)
                push v into the Q
    return dist and prev
```



A* search: Heuristic = PAST & FUTURE

• For the current node n, e.g., n = (0,4)

(0,0)	(0,1)	(0,2)	(0,3)	(0,4)
(1,0)				(1,4)
(2,0)	(2,1)	(2,2)	(2,3)	(2,4)
(3,0)				(3,4)
(4,0)	(4,1)	(4,2)	(4,3)	(4,4)

- cost function g(n)
 - The cost from the start node to the current node n. It's how far (how much cost) you've already traveled.
 - g((0,4)) = 4
- heuristic function h(n)
 - It estimates the cost from the current node n to the goal node.
 - It's an educated guess of how much farther it is to the goal.
 - h((0,4)) = 4
- The total cost function f(n)
 - the sum of g(n) and h(n), i.e., f(n) = g(n) + h(n)
 - f((0,4)) = g((0,4)) + h((0,4)) = 4 + 4 = 8

 A^* can balance between paths that are already known (using g(n)) and paths that are likely to be shorter (using h(n)).

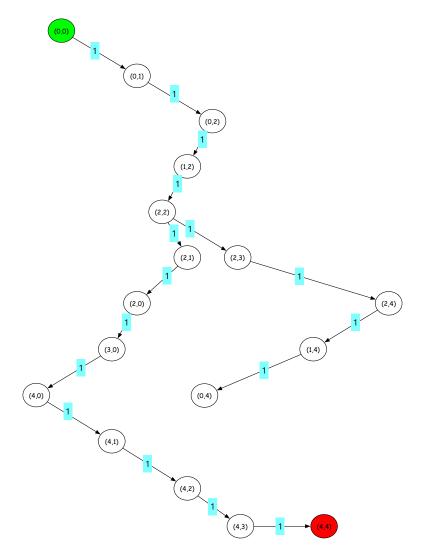


- How A* Works:
 - 1. Start at the initial node
 - 1. e.g., n = (0,0)
 - 2. The initial cost g((0,0)) = 0, as we haven't moved yet.
 - 3. f((0,0)) = g((0,0)) + h((0,0)) = h((0,0))
 - 2. Explore neighboring nodes
 - 1. For each neighbor, e.g., n = (0,1) and n = (1,0)
 - 1. g(n): The cost to reach the current node from the initial node.
 - 2. h(n): An estimate of how far this neighbor is from the goal node.
 - 3. Update the total cost
 - 1. For each neighboring node, compute f(n).
 - 2. put f(n) into a priority queue (a list of nodes to explore, ordered by their f(n) values).
 - 4. Choose the node with the smallest f(n) from the queue and repeat the process.
 - 5. Stop when you reach the goal node.
 - 1. No need to traverse all nodes. The first time A^* reaches the goal, it will have found the shortest path because of how f(n) is calculated.



A* search:

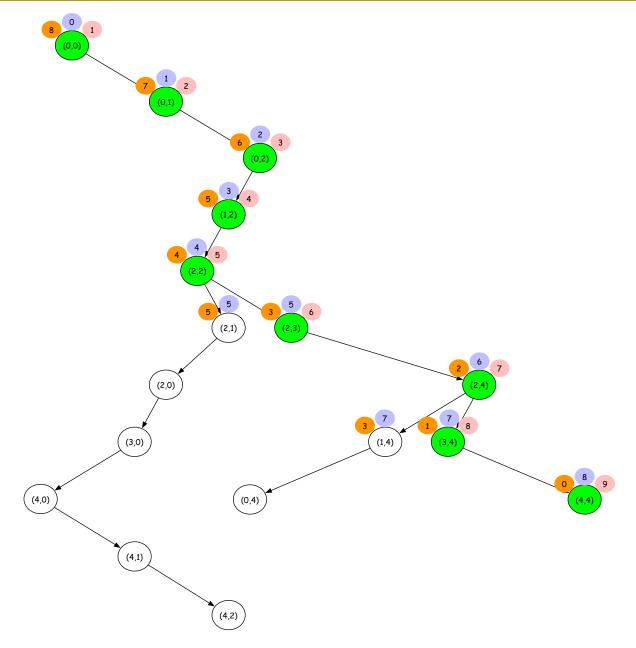
Heuristic = PAST & FUTURE



```
A Star Search (Grid, Start, Goal):
    Initialize an empty priority queue Q
   Initialize a VISITED to keep track of visited nodes
    Initialize a dictionary "g score" where g score[node] = ∞ for all nodes, except
g score[Start] = 0
    Initialize a dictionary "f score" where f score[node] = ∞ for all nodes, except
f score[Start] = h(Start, Goal)
    Initialize a came from to track the path
    Add Start to Q with priority f score[Start]
    while Q is not empty:
        current = Q.pop() // Get the node in Q with the lowest f score
        if current == Goal:
            return // Found the shortest path
        Mark current as visited
        for each neighbor of current in the 4 possible directions:
            if neighbor is not walkable or neighbor is in "visited":
                continue
            tentative g score = g score[current] + 1 // Each move costs 1
            if tentative g score < g score[neighbor]:</pre>
                Edit came from // This path to neighbor is better, so record it
                g score[neighbor] = tentative g score
                f score[neighbor] = g score[neighbor] + h(neighbor, Goal) // f(n)
= q(n) + h(n)
                if neighbor is not in Q:
                    Add neighbor to Q with priority f score[neighbor]
   return -1 // No path found
```

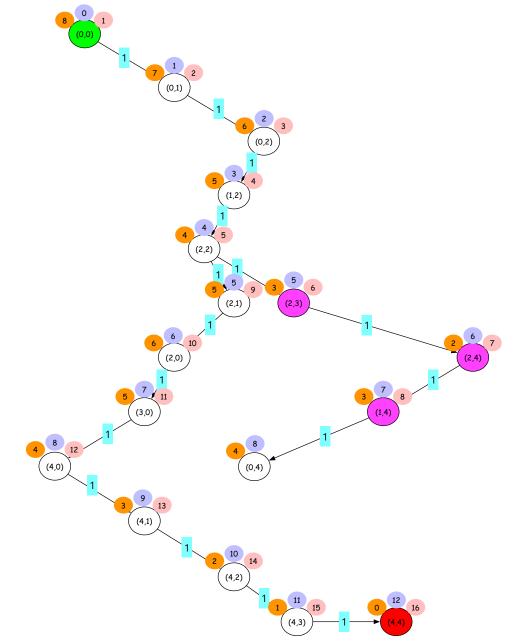


0	0	0	1	0
1	1	0	1	0
0	0	0	0	0
0	1	1	1	0
0	0	0	1	0





0	0	0	1	0
1	1	0	1	0
0	0	0	0	0
0	1	1	1	1
0	0	0	0	0

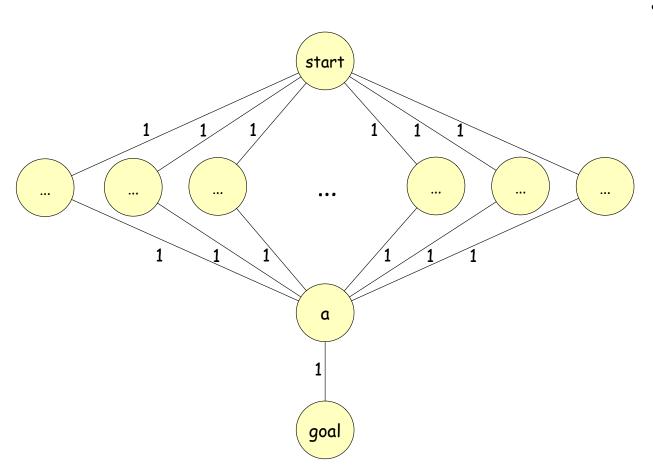




- A* search
 - Pros
 - Efficient Search: A* focuses on the most promising nodes first by leveraging the actual cost (g(n)) and heuristic estimate (h(n)); No need to traverse all nodes.
 - Guaranteed Optimality: A* guarantees the Shortest Path as long as the heuristic is admissible and consistent.
 - Customizable with Heuristics
 - Manhattan distance
 - Euclidean distance
 - Chebyshev distance
 - etc

- A* search
 - Cons
 - Memory-Intensive
 - Stores ALL explored nodes in memory. This
 can become a significant problem in large
 search spaces as memory usage grows
 rapidly.
 - Depends Heavily on the Heuristic
 - A poorly chosen or weak heuristic can cause A* to behave inefficiently, exploring many unnecessary nodes.
 - If the heuristic is NOT admissible or inconsistent, A* may not find the optimal solution.





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Summary



- Uninformed Search
 - · Search without guidance, i.e., blind search
 - Breadth-first Search (BFS)
 - Depth-first Search (DFS)
- Informed Search
 - Search with guidance, i.e., heuristic search
 - common heuristics
 - Manhattan Distance
 - Chebyshev Distance
 - Eunlidean Distance
 - Greedy Best-first Search (leverage the estimation regarding future)
 - Dijkstra's Algorithm (leverage the collected info regarding past)
 - A* search (leverage both past and future)

0	0	0	1	0
1	1	0	1	0
0	0	0	0	0
0	1	1	1	1
0	0	0	0	0

- · Define
 - node
 - edge
- Construct plain graph/tree
- · Construct weighted graph/tree

Homework: Eight Puzzle Game

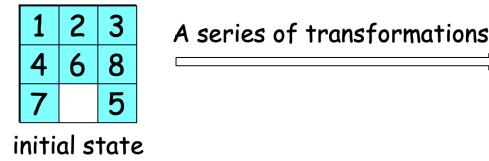


The 8-puzzle is a sliding puzzle that consists of a 3×3 grid containing 8 numbered tiles (1 to 8) and one empty space.





The goal is to rearrange the tiles from an arbitrary initial configuration into a specified goal configuration by sliding the tiles into the empty space.



goal state

Homework: Eight Puzzle Game







- 1. Can we transfer this problm into a general graph search problem? (node & edge)
 - 1. If it's okay, provide the graph; if not, provide the analysis to show why.
- 2. Slove this problem using dynamic construction of the search space.
 - 1. Implement BFS
 - 2. Implement DFS
 - 3. Compare the performance and provide your understanding.
- 3. Analyze how to build a weighted tree for this problem?
 - 1. Implement greedy best-first search, report your result.
 - 2. Implement dijkstra's algorithm, report your result.
 - 3. Implement A* search, report your result.