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OS FINAL PROJECT REPORT (CHECKPOINT 4)

Introduction

In the beginning of project checkpoint, we are asked to implement the producer-consumer functions that produce alphabets in alphabetical order that loops for infinitely using the thread yielding that we have to code it ourselves with the help of the hints given by the project.

As the project progresses to the next checkpoint, we are asked to implement a scheduler which is called myTimer0Handler, and improve the testing file with the given instruction in the checkpoint assignment.

Third checkpoint starts to implement the semaphores to ensure the fairness of the code and prevent it from raising errors, especially the producer-consumer updates.

Going to the fourth checkpoint, the last change that is left to be made is adding another feature which is another producer that generates numbers ‘0’ to ‘9’ that loops for infinitely.

test3threads.c

The code starts off with the global initialising of mutex, full, empty, p1 (turn for Producer1), p2 (turn for Producer2) which is mainly for the semaphores that is to be used between the producer-consumer function. As for the other global variables like buffer (3-deep char), buffer head, and buffer tail, it is mainly used to keep up with the bounded buffer circular queue structure.

```
__data __at (0x20) int mutex;
__data __at (0x22) int full;
__data __at (0x24) int empty;
__data __at (0x28) int p1;
__data __at (0x29) int p2;

__data __at (0x3D) char buffer [3];
__data __at (0x30) char buffer_head;
__data __at (0x31) char buffer_tail;
```

The main function first creates the semaphores that are introduced in the global variables by filling the initial space using the SemaphoreCreate function. Then it continues with initialising the buffer data with the empty character. The buffer head and tail initialise as 0 that is to be updated later on. Lastly, the thread created for Producer1 and Producer2 which is used to generate for Consumer.

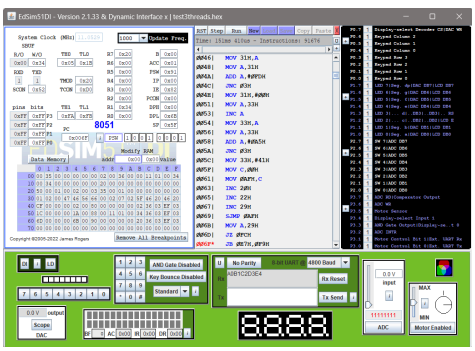
```
void main(void) {
    /*
        * @@@ [1 pt] initialize glob
        * @@@ [4 pt] set up Producer
        * Because both are infinite
        * in this function and no re
    */
    // create semaphore
    SemaphoreCreate(mutex, 1);
    SemaphoreCreate(full, 0);
    SemaphoreCreate(empty,3);
    SemaphoreCreate(p1, 1);
    SemaphoreCreate(p2, 0);

    //initialize buffer, head&tail
    buffer[0] = ' ';
    buffer[1] = ' ';
    buffer[2] = ' ';
    buffer_head = 0;
    buffer_tail = 0;

    ThreadCreate(Producer1);
    ThreadCreate(Producer2);
    Consumer();
}
```

The Producer1, Producer2, and Consumer function will be discussed deeper in the later part, but the main context used are the updates of buffer data using alphabets and numbers that are going to be fed to the consumer so that it can be output later on.

UART Fair Version



The Producer and Consumer uses the semaphores to maintain its fairness (mutex & empty) whereas the Producer1 and Producer2 maintain it by relying on the turns (Producer1 & Producer2) which is to be switch upon finishing its respective task so that alphabets and digits alternately generated for the consumer to send it to the SBUF. Through this, we can see the result on the right: ‘A0B1C2…’ which shows that the fairness that is designed in the code is working.

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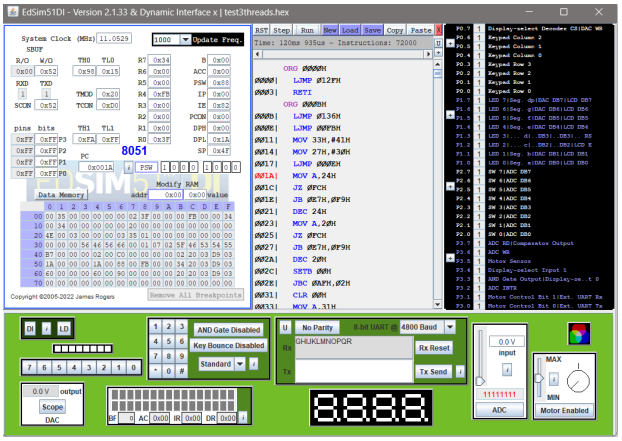
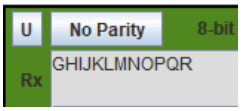
UART Unfair Version



After removing the fairness using turns (Producer1 & Producer2) semaphores, the result would appear in 2 different formats where first one will only print out alphabets and second would only print out digits. This is the problem that can be prevented using additional semaphores to keep the generated data in track.

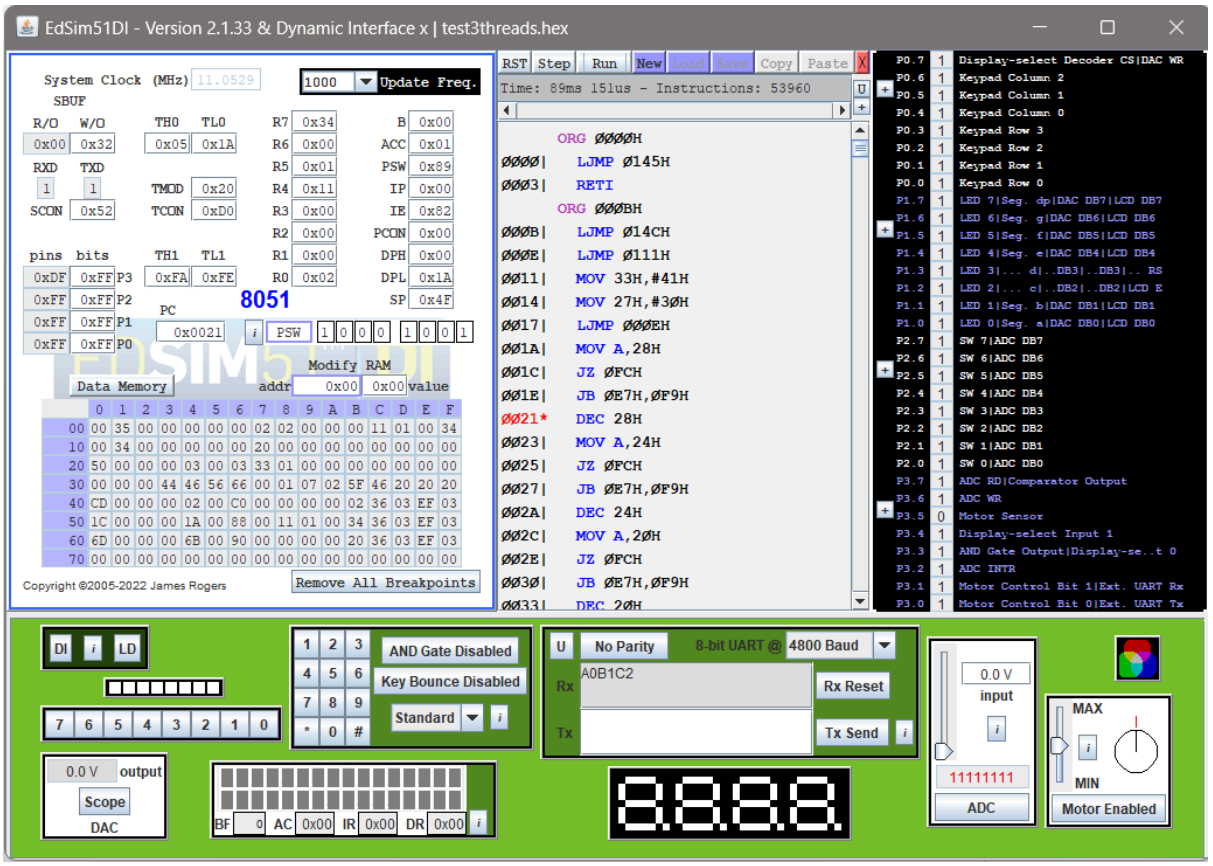


Typescript

```
Christian owen@LAPTOP-TL3OUSE0 /cygdrive/d/University/Semester5/OS/Checkpoints/C
4
$ make clean
rm *.hex *.ihx *.lnk *.lst *.map *.mem *.rel *.rst *.sym
rm: cannot remove '*.ihx': No such file or directory
rm: cannot remove '*.lnk': No such file or directory
make: *** [Makefile:25: clean] Error 1

Christian owen@LAPTOP-TL3OUSE0 /cygdrive/d/University/Semester5/OS/Checkpoints/C
4
$ make
sdcc -c  test3threads.c
test3threads.c:104: warning 158: overflow in implicit constant conversion
sdcc -c  preemptive.c
preemptive.c:195: warning 85: in function ThreadCreate unreferenced function arg
ument : 'fp'
sdcc  -o test3threads.hex test3threads.rel preemptive.rel
```

Running Producer1:



Below is a semaphore wait implementation for producer1 that decreases the turn (producer1), empty, and mutex.

```
ØØ1A|   MOV A,28H
ØØ1C|   JZ ØFCH
ØØ1E|   JB ØE7H,ØF9H
ØØ21*   DEC 28H
ØØ23|   MOV A,24H
ØØ25|   JZ ØFCH
ØØ27|   JB ØE7H,ØF9H
ØØ2A|   DEC 24H
ØØ2C|   MOV A,2ØH
ØØ2E|   JZ ØFCH
ØØ3Ø|   JB ØE7H,ØF9H
ØØ33|   DEC 2ØH
```

Below is a semaphore signal implementation for producer 1 that increases the turn (producer2), mutex, and full.

```
ØØ63|   INC 2ØH
ØØ65|   INC 22H
ØØ67|   INC 29H
```

Explanation:
This part of my test3threads file is the producer who updates the buffer data using loops alphabets from ‘A’ to ‘Z’ back to ‘A’ and so on while making sure that the buffer tail is properly updated so that Producer2 and Consumer can be updated by the changes. The semaphore is also used here to keep in track with the switching system between Producer1, Producer2 and Consumer so that it will not raise an error.

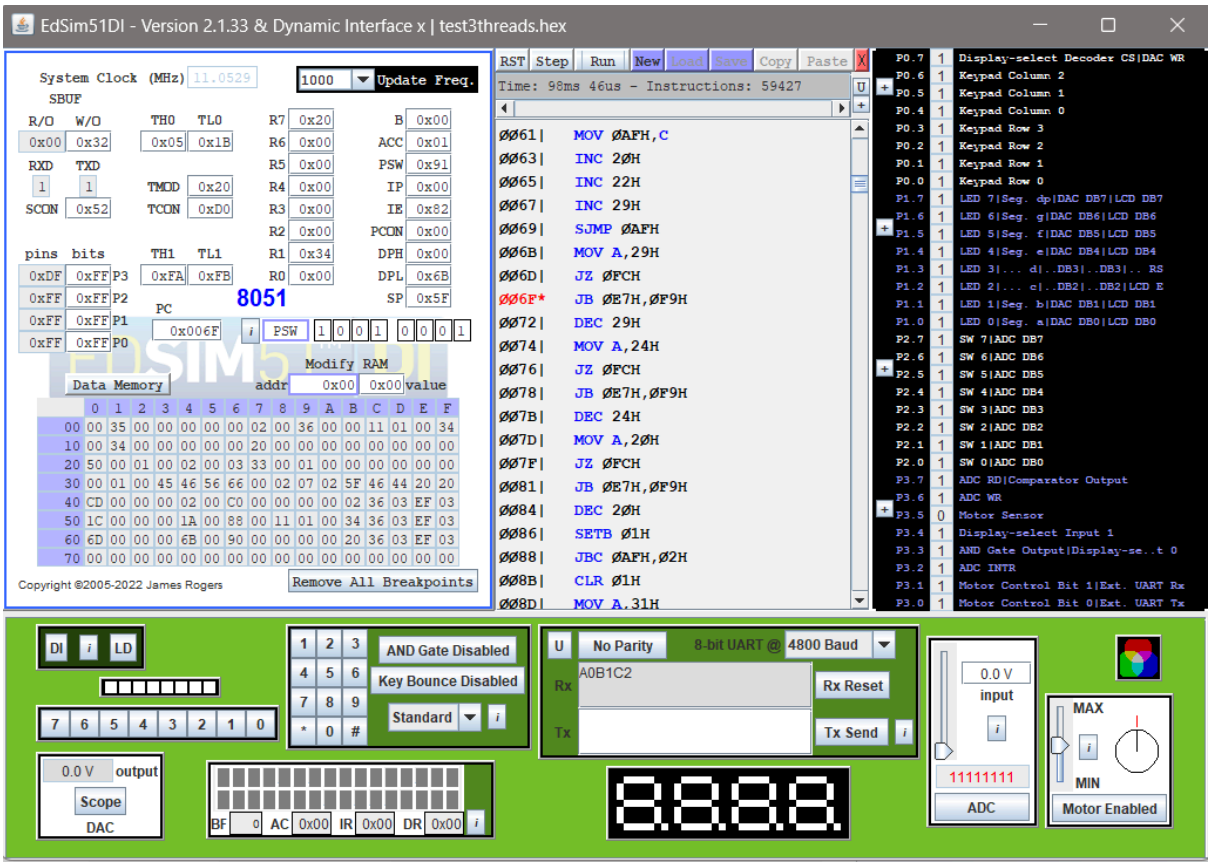
```
void Producer1(void) { // A to Z
    /*
        * @@@ [2 pt]
        * initialize producer data structure, and then ent
        * an infinite loop (does not return)
    */
    static __data __at (0x33) char buffer_next = 'A';

    while (1) {
        /*
            * @@@ [6 pt]
            * wait for the buffer to be available,
            * and then write the new data into the buffer
        */
        SemaphoreWait(p1);
        SemaphoreWait(empty);
        SemaphoreWait(mutex);

        __critical{
            buffer[buffer_tail] = buffer_next;
            buffer_tail++;
            if (buffer_tail > 2) buffer_tail = 0;
            buffer_next++;
            if(buffer_next > 'Z') buffer_next = 'A';
        }

        SemaphoreSignal(mutex);
        SemaphoreSignal(full);
        SemaphoreSignal(p2);
    }
}
```

Running Producer2:



Below is a semaphore wait implementation for producer2 that decreases the turn (producer2), empty, and mutex.

```
ØØ6B|   MOV A,29H
ØØ6D|   JZ ØFCH
ØØ6F*   JB ØE7H,ØF9H
ØØ72|   DEC 29H
ØØ74|   MOV A,24H
ØØ76|   JZ ØFCH
ØØ78|   JB ØE7H,ØF9H
ØØ7B|   DEC 24H
ØØ7D|   MOV A,2ØH
ØØ7F|   JZ ØFCH
ØØ81|   JB ØE7H,ØF9H
ØØ84|   DEC 2ØH
```

Below is a semaphore signal implementation for producer2 that increases the turn (producer1), mutex, and full.

```
ØØB4|   INC 2ØH
ØØB6|   INC 22H
ØØB8|   INC 28H
```

Explanation:
The other part of the Producer function in the file is the Producer2 that has the same functionality as Producer1 but in this case, the buffer data is updated with loop numbers from ‘0’ to ‘9’ back to ‘0’ and so on. The semaphore that is implemented here is also the same as Producer1, except the semaphore for turns will be flipped to let Producer1 update the buffer data in the next turn.

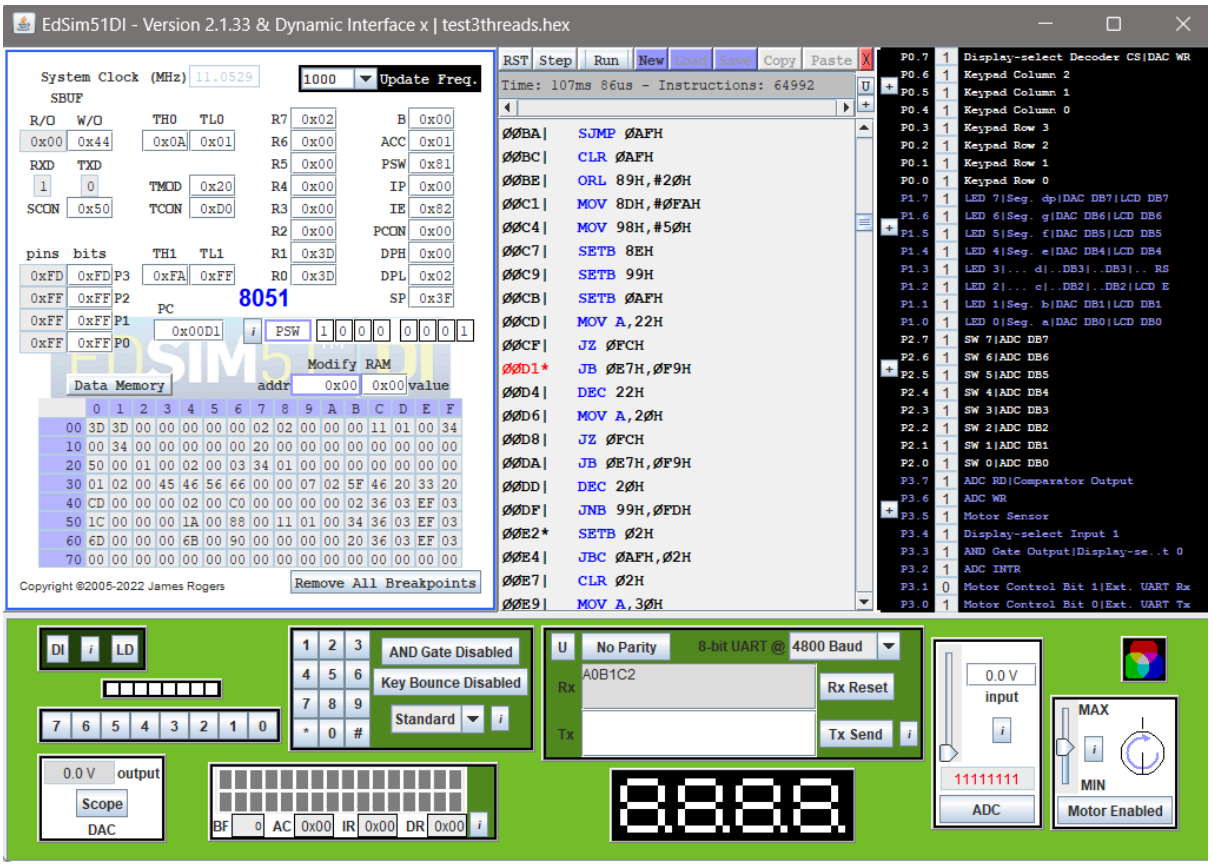
```
void Producer2(void) { // 0 to 9
    /*
        * @@@ [2 pt]
        * initialize producer data structure, and then e
        * an infinite loop (does not return)
    */
    static __data __at (0x27) char buffer_next2 = '0';

    while (1) {
        /*
            * @@@ [6 pt]
            * wait for the buffer to be available,
            * and then write the new data into the buffe
        */
        SemaphoreWait(p2);
        SemaphoreWait(empty);
        SemaphoreWait(mutex);

        __critical{
            buffer[buffer_tail] = buffer_next2;
            buffer_tail++;
            if (buffer_tail > 2) buffer_tail = 0;
            buffer_next2++;
            if(buffer_next2 > '9') buffer_next2 = '0';
        }

        SemaphoreSignal(mutex);
        SemaphoreSignal(full);
        SemaphoreSignal(p1);
    }
}
```

Running Consumer:



Below is a semaphore wait implementation for Consumer that decreases the empty and mutex semaphore.

```
ØØCD|    MOV A,22H
ØØCF|    JZ ØFCH
ØØD1*    JB ØE7H,ØF9H
ØØD4|    DEC 22H
ØØD6|    MOV A,2ØH
ØØD8|    JZ ØFCH
ØØDA|    JB ØE7H,ØF9H
ØØDD|    DEC 2ØH
```

Below is a semaphore signal implementation for producer2 that increases the mutex and full semaphore.

```
Ø1ØB|    INC 2ØH
Ø1ØD|    INC 24H
```

Explanation:
Just like the Producer1 and Producer2, the Consumer relies on the semaphore mutex and full to wait for its turn to take the buffer data into SBUF. It means that it needs to wait till the full semaphore increases till 3 to load the head of buffer data into SBUF. After that, the buffer head will be updated.

```c
void Consumer(void) {
    /* @@@ [2 pt] initialize Tx for polling */
    EA = 0;
        TMOD |= 0x20;
        TH1 = -6;
        SCON = 0x50;
        TR1 = 1;
        TI = 1;
    EA = 1;
    while (1) {
        /*
            * @@@ [2 pt] wait for new data from p
            * @@@ [6 pt] write data to serial por
            * poll for Tx to finish writing (TI),
            * then clear the flag
        */
        SemaphoreWait(full);
        SemaphoreWait(mutex);

        while (!TI);
        __critical{
            SBUF = buffer[buffer_head];
            buffer[buffer_head] = ' ';
            buffer_head++;
            if(buffer_head > 2) buffer_head = 0;
            TI = 0;
        }

        SemaphoreSignal(mutex);
        SemaphoreSignal(empty);
    }
}
```