

SybilLimit: **Secure** **Decentralized** **Reputation System**

Petar Maymounkov
petar@protocol.ai

SybilLimit: Secure Decentralized Reputation System

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- Based on SybilLimit (2010, Yu, et al.)
- Some simplifications, small improvements

Introduction

Context

In search of **robust** decentralized architectures:

- Open membership
- Surviving (malfunction, Sybil attacks, **Sybil infections**)
- Scalable
- Live (responsive)

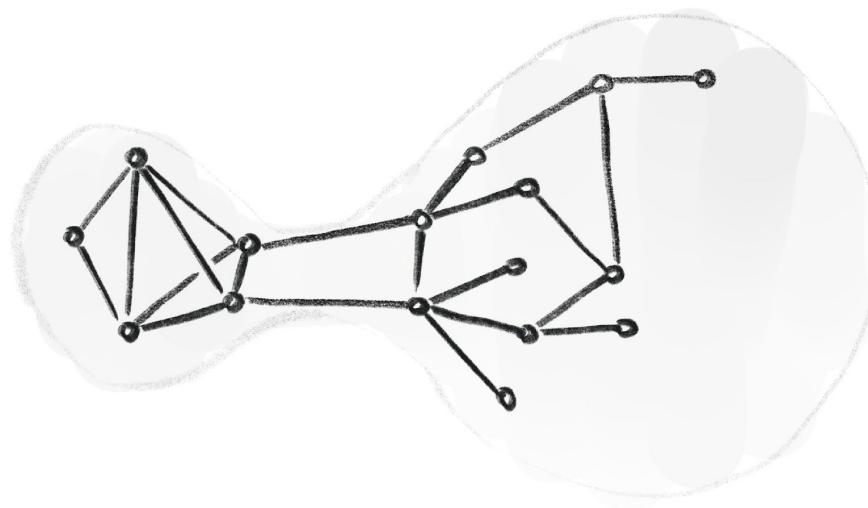
For broadcast (gossipsub), DHT (Kademlia) and in general.

Summary

- + Technical ideas are strong.
- Solution has a non-intuitive “user interface”.
- Solution not strong for consensus, which handles <25% infection.
- + I was able to reuse techniques to design a new provably robust decentralized system for broadcast and consensus.

Reputation system

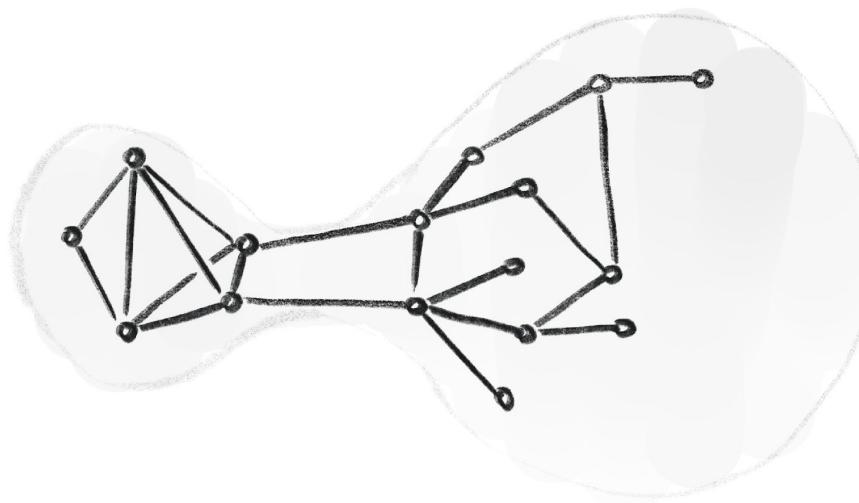
Motivated by the belief that “communities” are well-connected, with sparse bottlenecks between “communities”.



Reputation system

What is a “**reputation system**”?

- Add/remove edge/vertex
- Query: How well “**connected**” are two nodes?

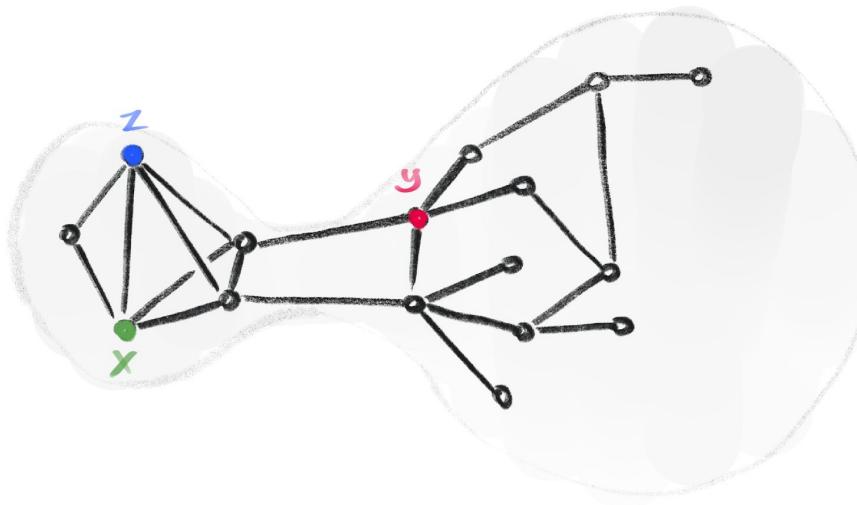


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$$\text{conn}(x, y) = \min_{\text{cut}} \frac{\text{edges across cut}}{\text{edges on side of } x}$$



Reputation system

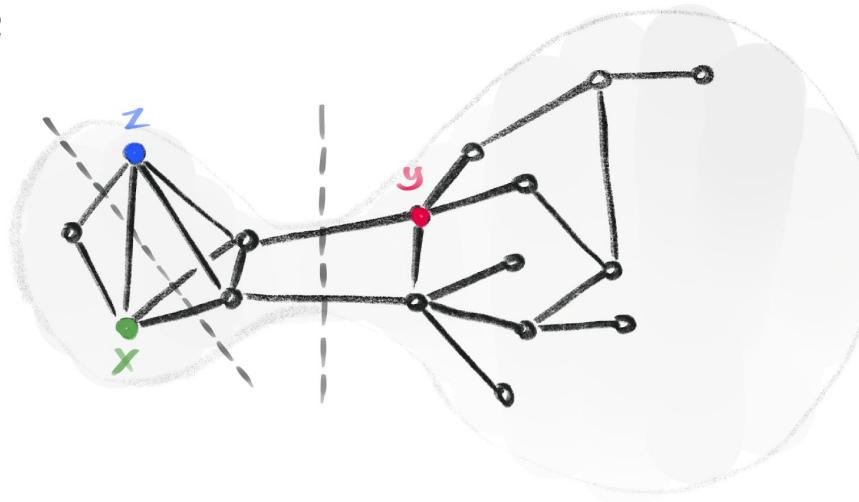
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$$\text{conn}(x, y) = \min_{\text{cut}} \frac{\text{edges across cut}}{\text{edges on side of } x}$$

$$\text{conn}(x, y) = \frac{2}{10} = 0.2$$

$$\text{conn}(x, z) = \frac{4}{5} = 0.8$$



Reputation system

What is a “**reputation system**”?

- Add/remove edge/vertex
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Diffusion = Distribution of the last vertex of a random walk.

Diffusions from well connected vertices **intersect** more than diffusions from badly connected ones.



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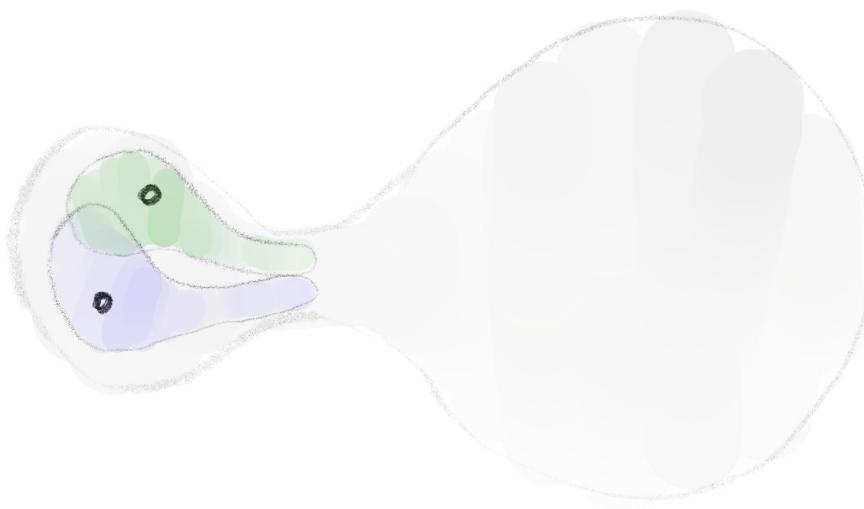
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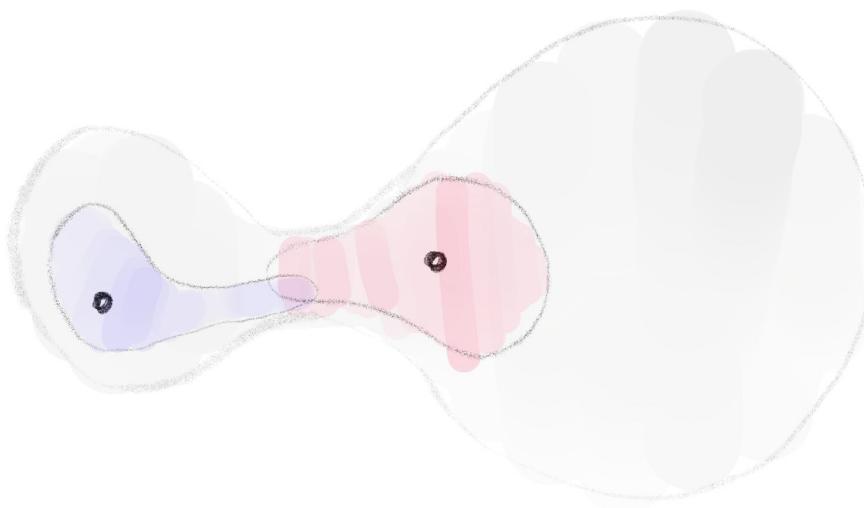
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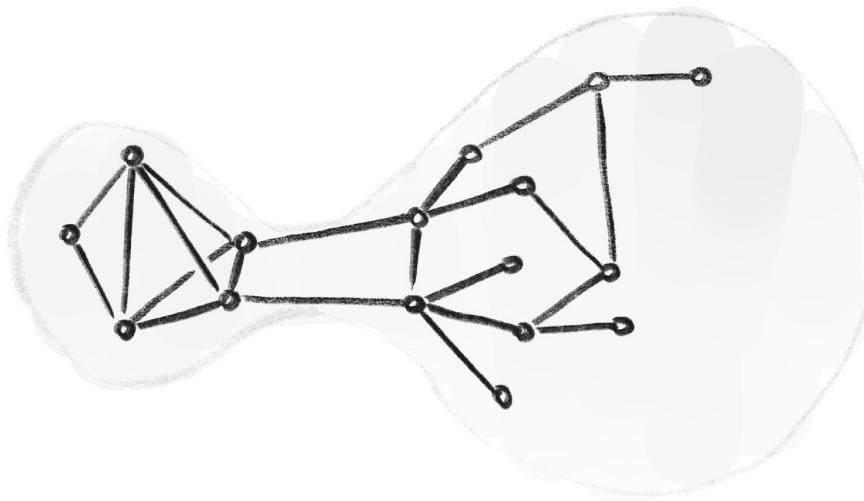
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Decentralized reputation system

What is a “**decentralized reputation system**”?

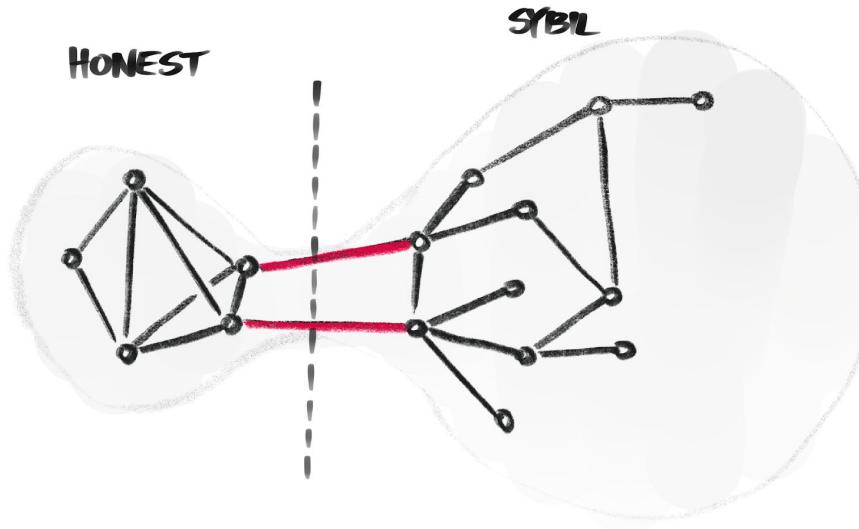
- Vertices are participants (compute nodes)
- Participants can add/remove **only incident** edges
- Query: Any two participant can compute their connectedness



Secure decentralized reputation system

What is a “**secure decentralized reputation system**”?

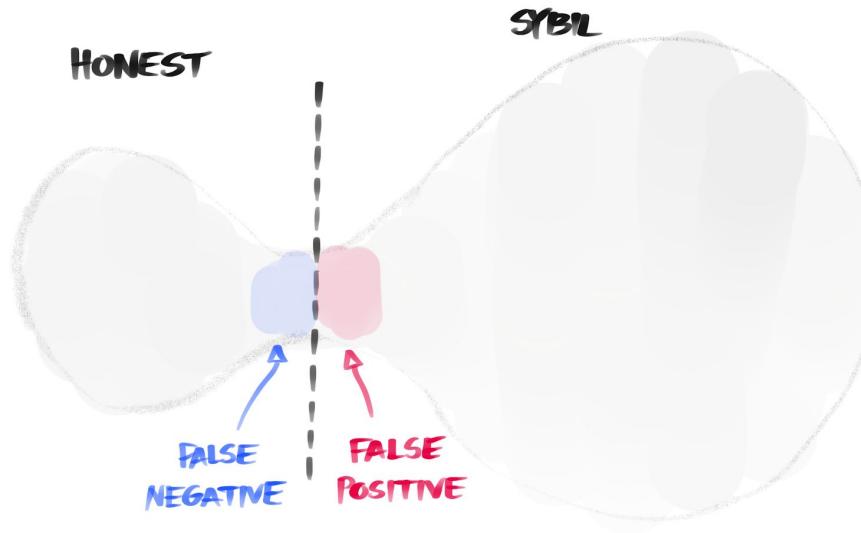
- Honest participants + unlimited amount of **Sybil** participants
- Honest “verifier” can check whether a “suspect” is well-connected
 - Query: *Is my connectedness to you below a threshold (e.g. %5)? Yes/no?*



Secure decentralized reputation system

What is a “**secure decentralized reputation system**”?

- Honest participants + unlimited amount of **Sybil** participants
- Honest “verifier” can check whether a “suspect” is well-connected
 - Query: *Is my connectedness to you below a threshold (e.g. %5)? Yes/no?*
- Size of the bottleneck affects the correctness of the answer



Experimental results

Vague sense of numbers:

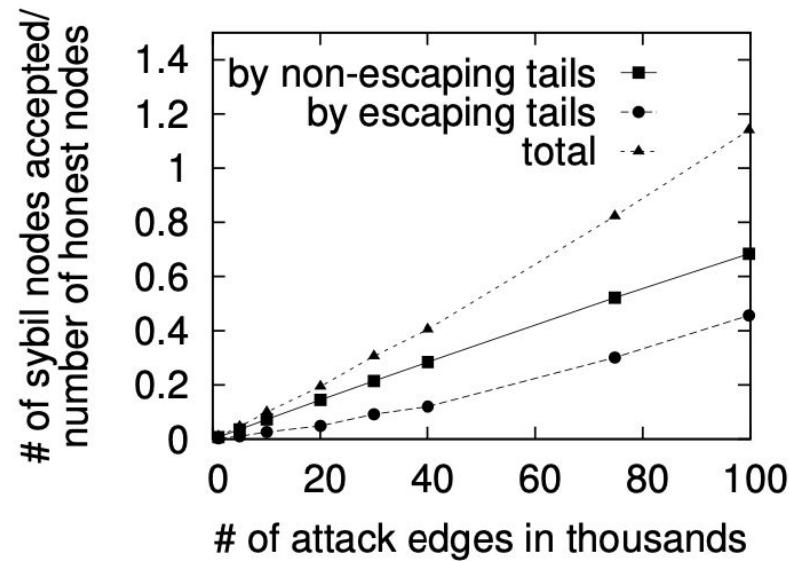
{ attack edges = 7% of honest nodes
infection = accepted sybils / accepted honest = 80% }

Consensus breaks
after 25% infection

Caveats:

Topology-dependent. Using random walk diffusion (PageRank walk is better).

Data set	Friendster
Data set source	[33]
Date crawled	Nov-Dec 2005
# nodes	932,512
# undirected edges	7,835,974
w used in SybilLimit	10
r used in SybilLimit	8,000



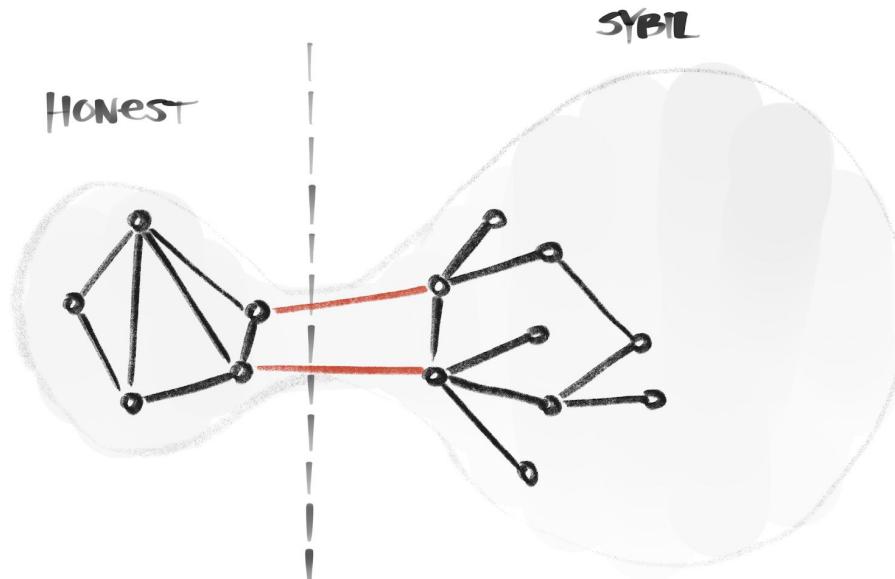
Applications?

Applications vary based on what you model as a vertex/edge.

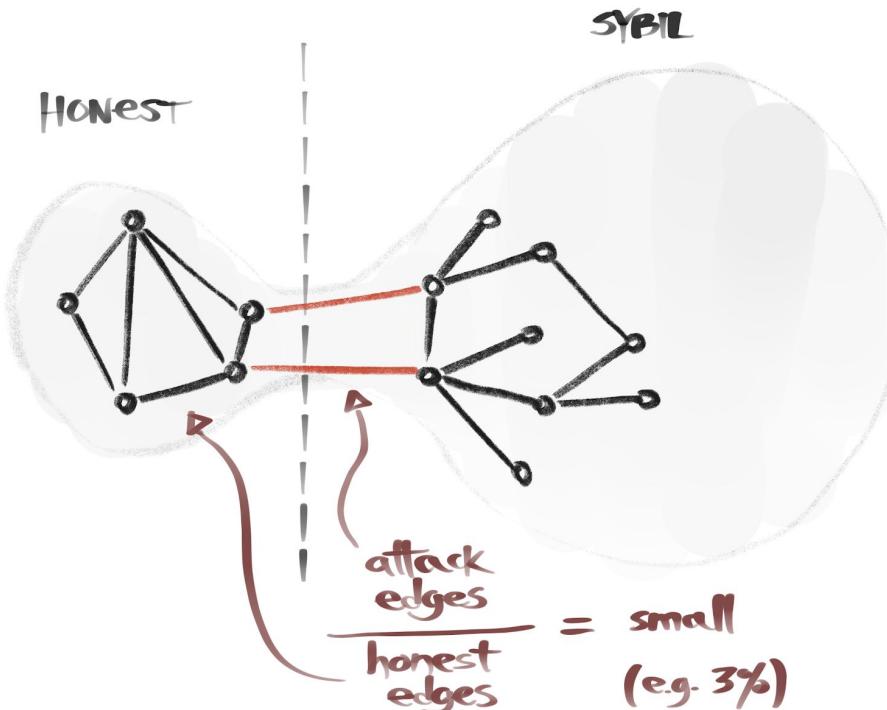
-
- Social / subjective:**
 - Social network: edge = subjective trust (based on manual user input)
 - Socio-economic / subjective:**
 - Club with membership fees: edge = monetary transaction
 - Historical / objective:**
 - Lineage relationships: “Alice joined the network via Bob and Chandra”
 - Introspective / objective:**
 - App-related history: “Alice performed good work for Bob in past epoch.”
- User interface
not intuitive?

High-level plan

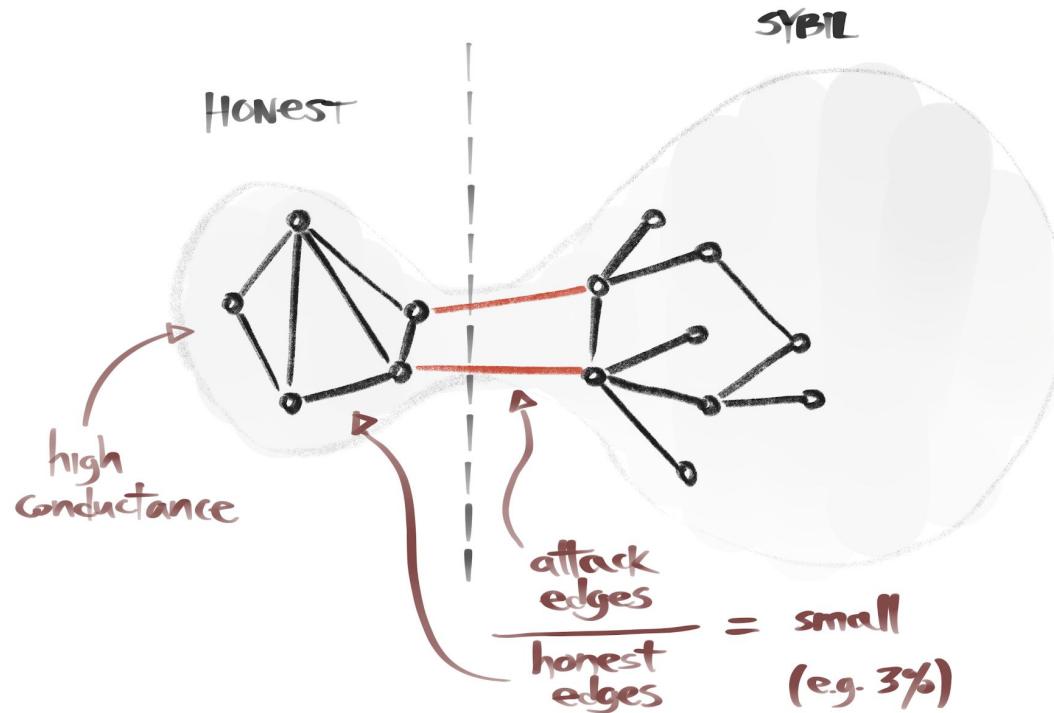
Model and assumptions



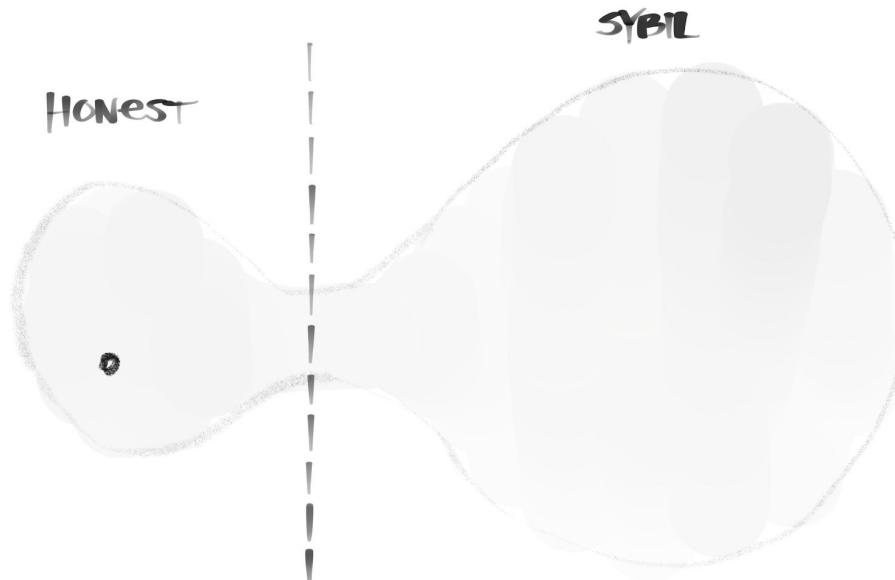
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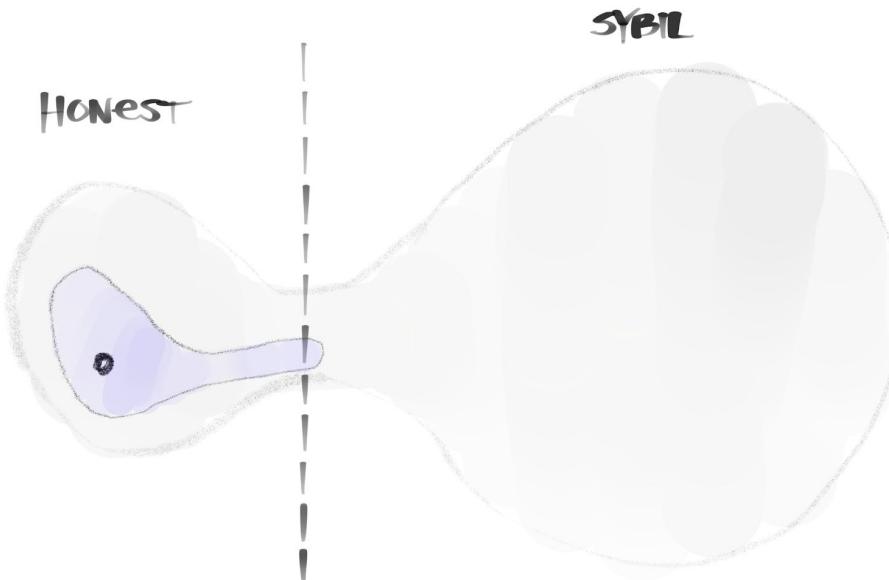
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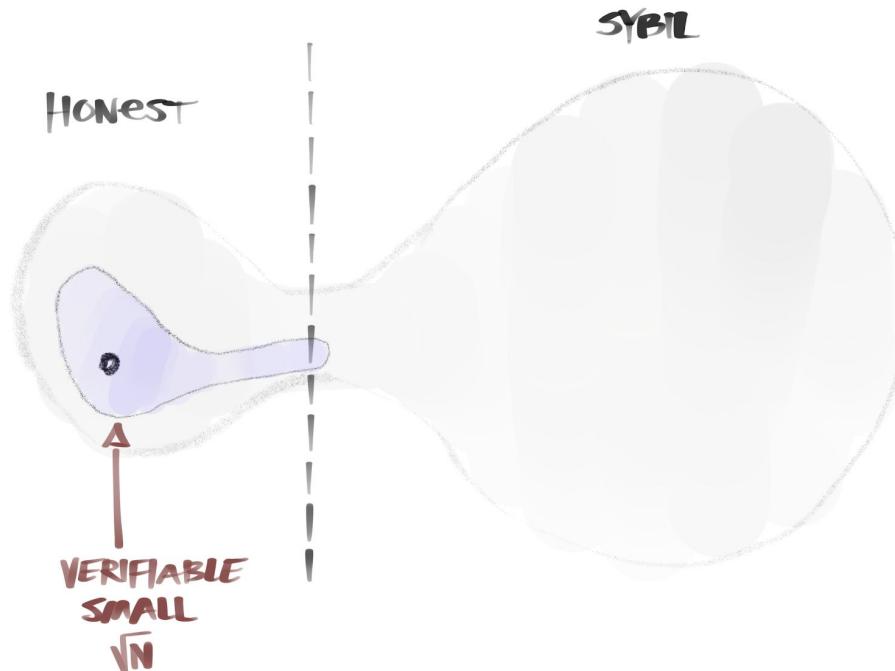
Proof idea: spectral view



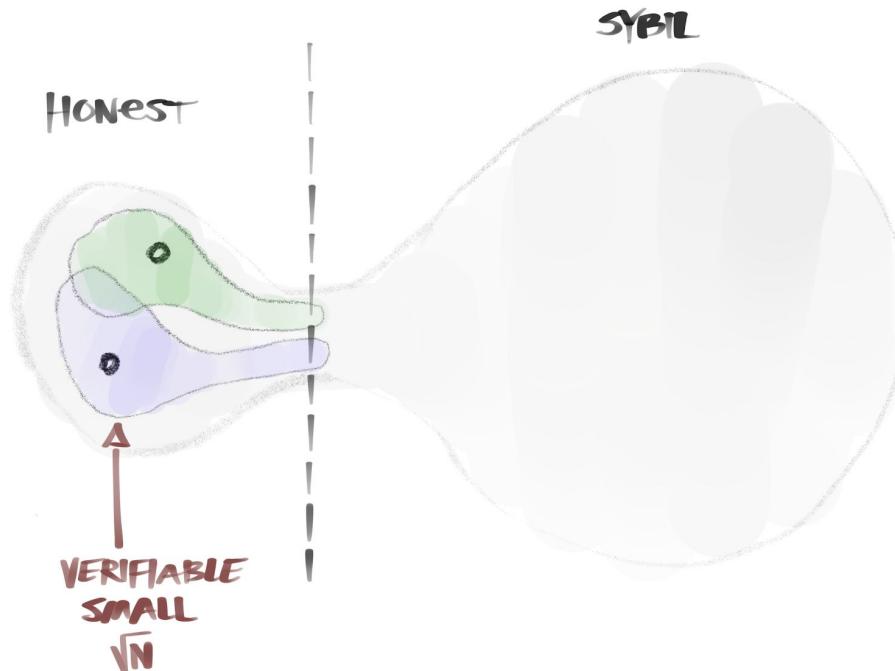
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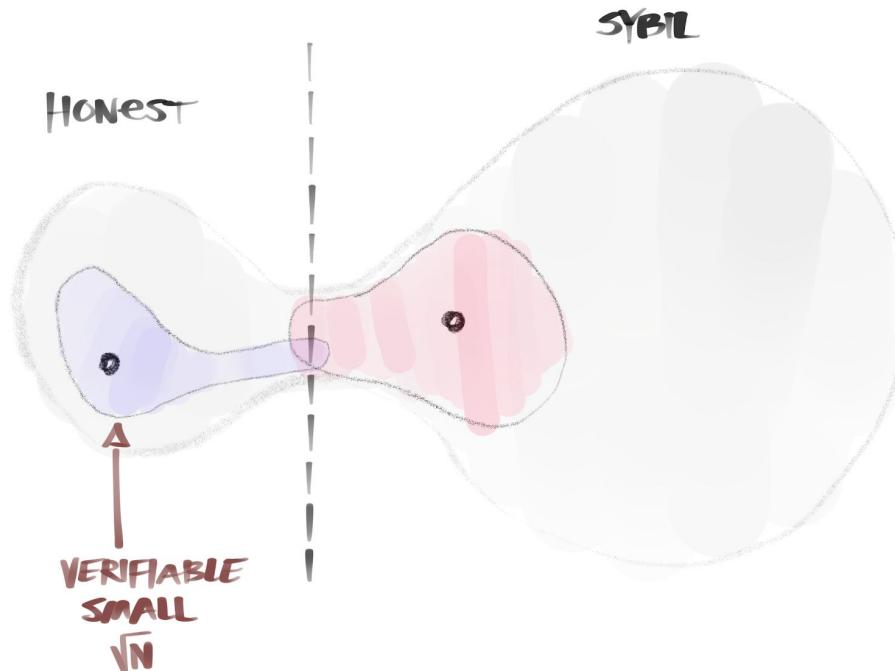
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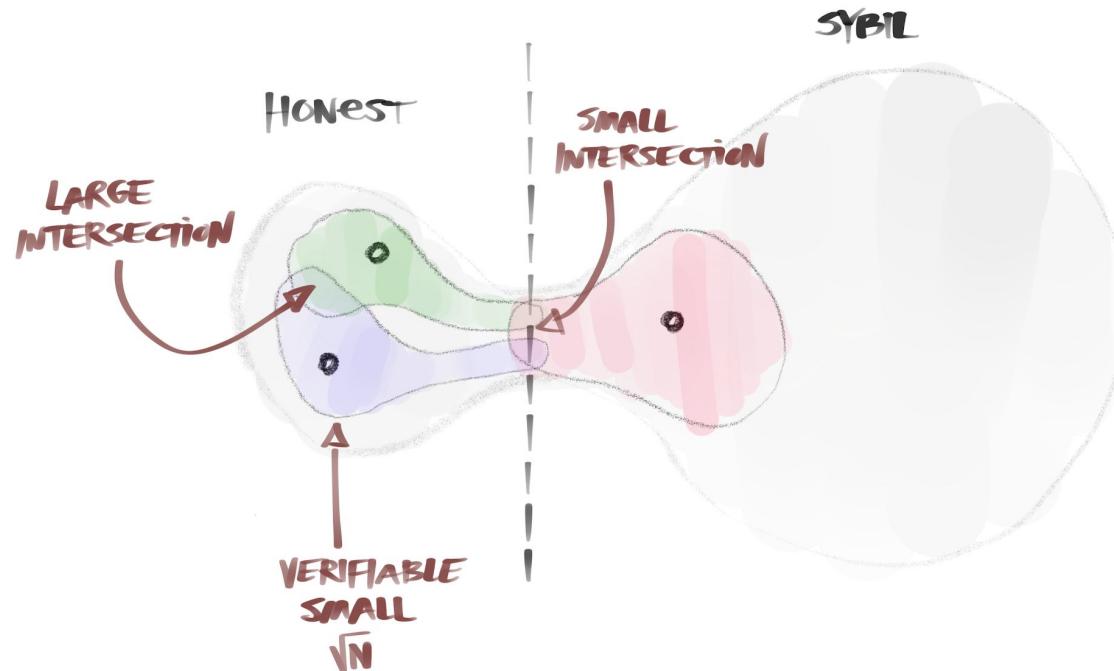
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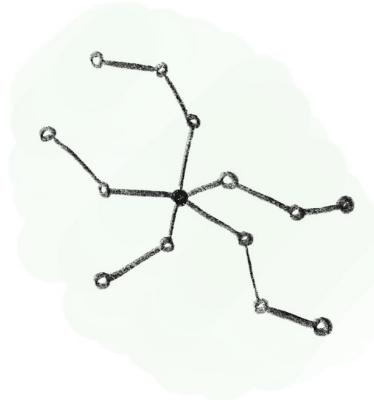
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Protocol and algorithms

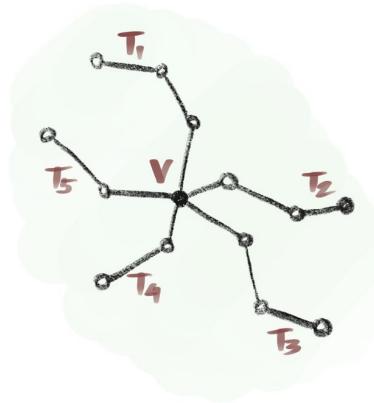
Verification algorithm

\sqrt{N} WALKS
 $\log N$ LENGTH



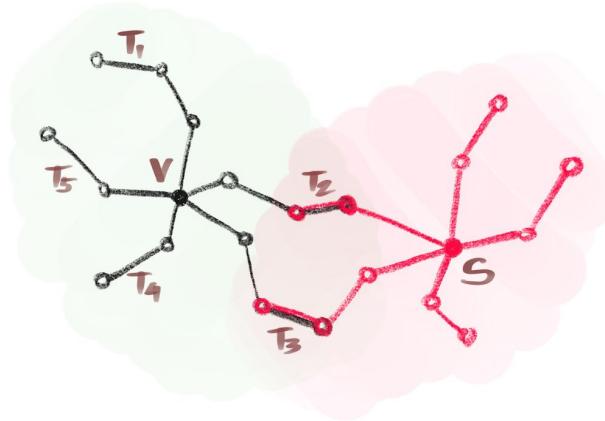
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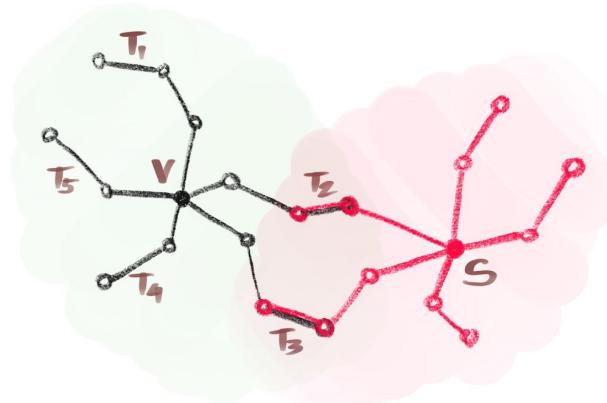
\sqrt{N} WALKS
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INTERSECTION
CONDITION

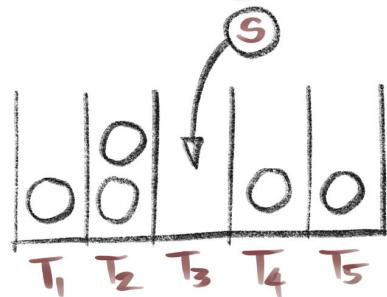
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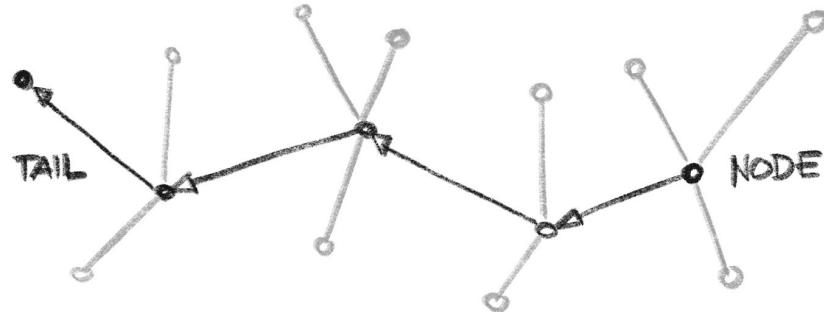
INTERSECTION
CONDITION

max counter – min counter ≤ 2

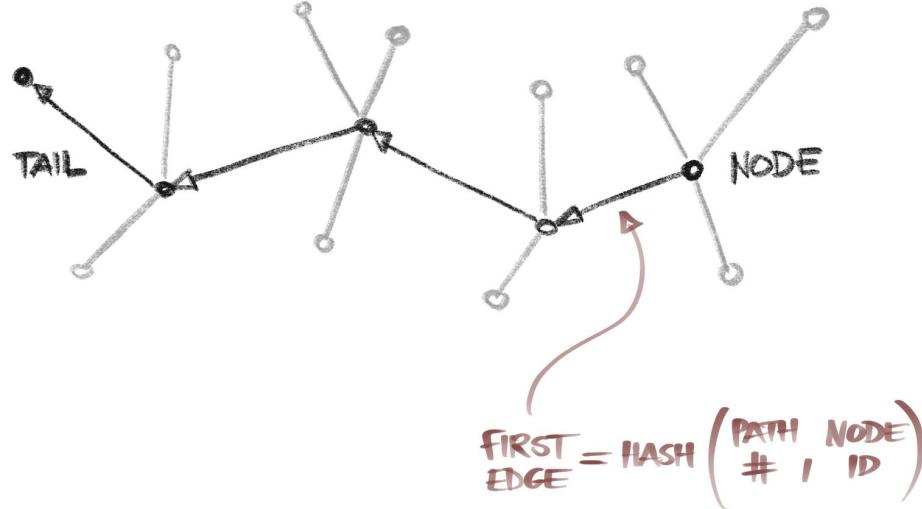


BALANCE
CONDITION

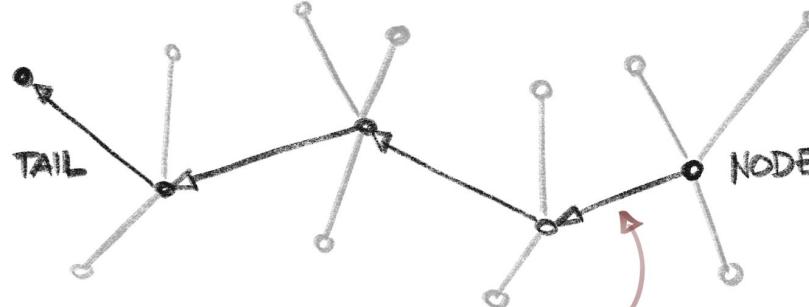
Verifiable paths



Verifiable paths



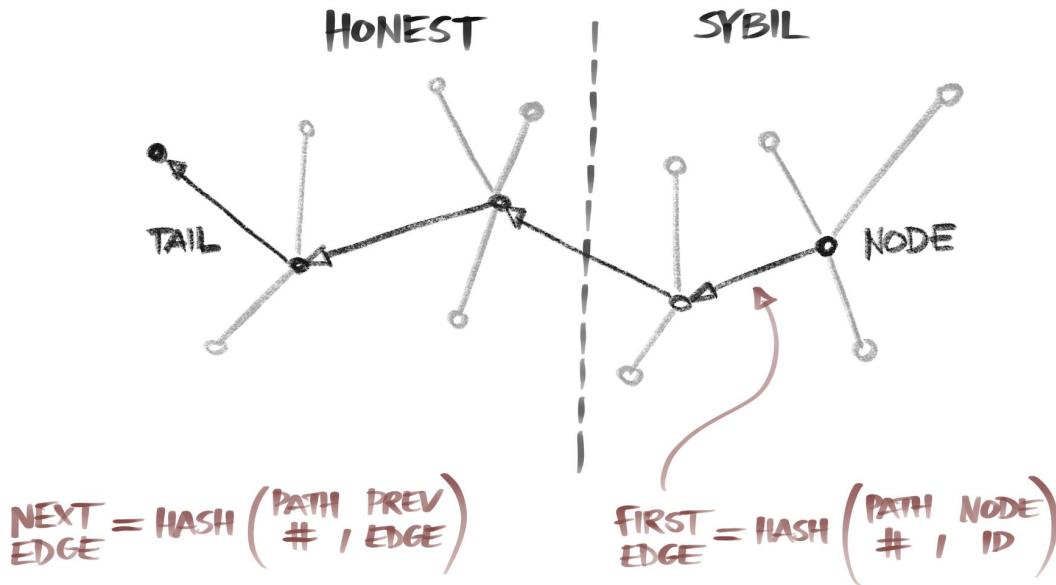
Verifiable paths



NEXT
EDGE = HASH $\left(\begin{matrix} \text{PATH}, \text{PREV} \\ \#, \text{EDGE} \end{matrix} \right)$

FIRST
EDGE = HASH $\left(\begin{matrix} \text{PATH}, \text{NODE} \\ \#, \text{ID} \end{matrix} \right)$

Verifiable paths



SYBILS CANNOT TAMPER
THE PATH
IN THE HONEST REGION

Proof of security

Theorem

MOST HONEST VERIFIERS:

- ACCEPT MOST HONEST SUSPECTS
- ACCEPT FEW SYBIL SUSPECTS

Theorem

$$\geq (1-\epsilon)N$$

MOST HONEST VERIFIERS:

$$\geq (1-\epsilon)N$$

- ACCEPT MOST HONEST SUSPECTS

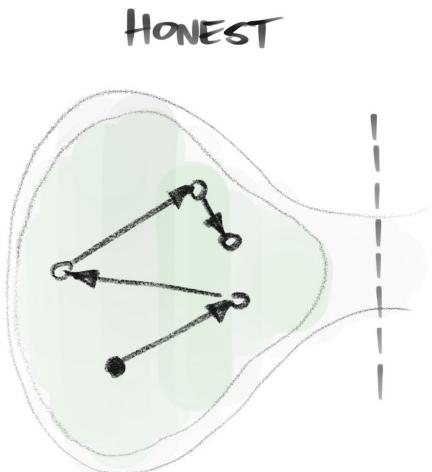
$$\leq \epsilon N$$

- ACCEPT FEW SYBIL SUSPECTS

N = number of honest nodes

$$\epsilon \approx \frac{\text{attack edges}}{\text{honest edges}}$$

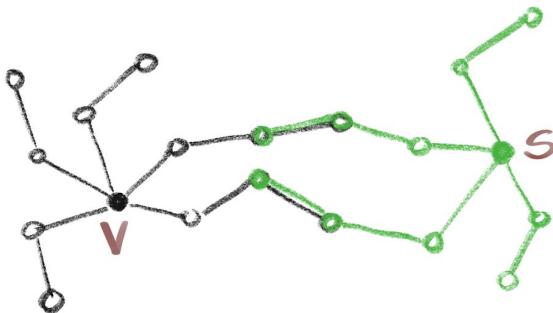
First step



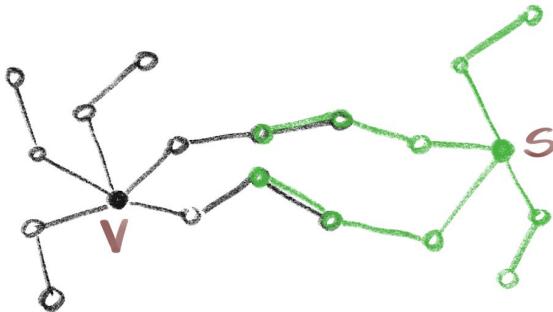
$\geq (1-\epsilon)N$
FOR MOST HONEST NODES,
THE TAIL OF A RANDOM WALK
IS
UNIFORM ACROSS
THE
HONEST EDGES
WITH HIGH PROBABILITY

$$\geq 1 - \frac{\text{attack edges} \cdot \log N}{N}$$

Accepting an honest suspect

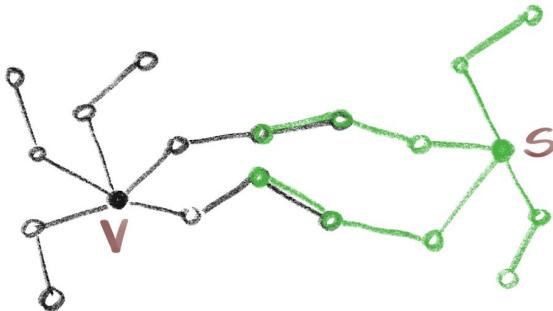


Accepting an honest suspect



DO \sqrt{N} RANDOM EDGES
INTERSECT
 \sqrt{N} RANDOM EDGES?

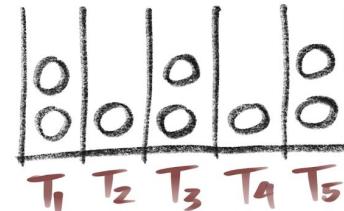
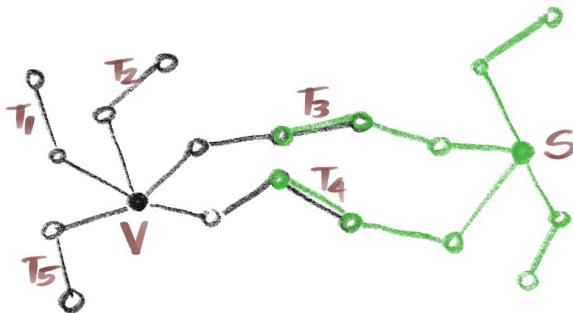
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INTERSECT
 \sqrt{N} RANDOM EDGES?

YES!
(BIRTHDAY PARADOX)

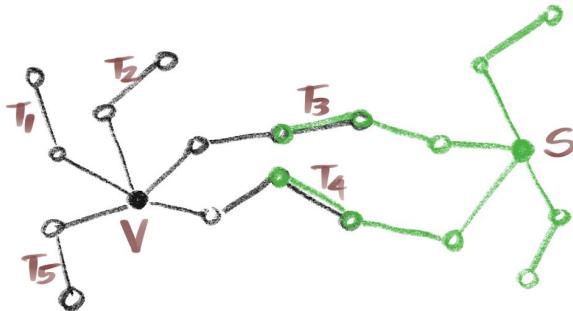
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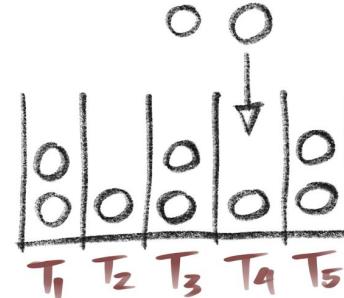
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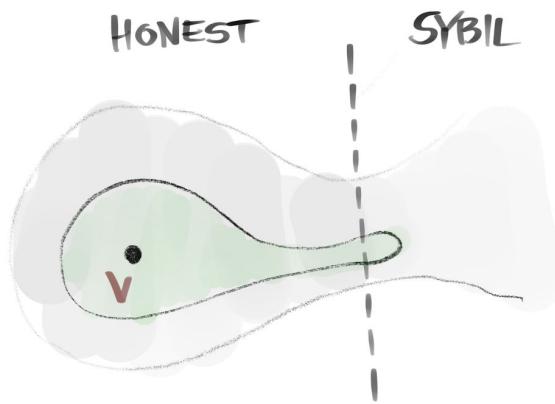
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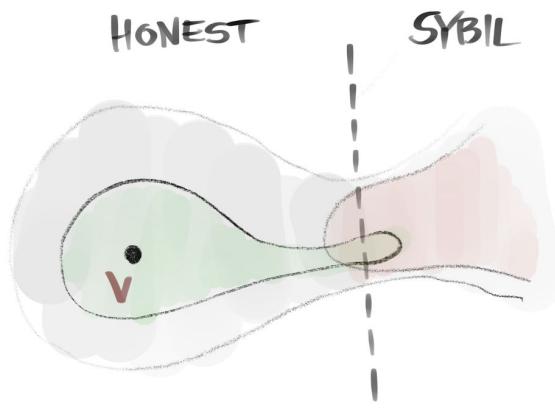


IF THE INTERSECTION
HAS
TWO OR MORE TAILS
THERE IS ONE TAIL
THAT CAN BE INSERTED
WITHOUT VIOLATING
BALANCE
(POWER OF TWO CHOICES)

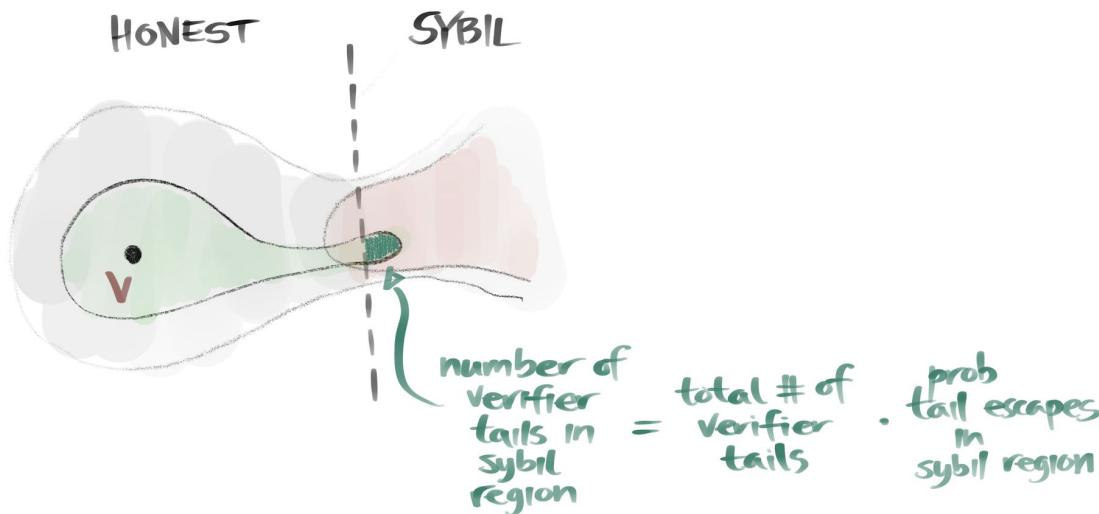
Rejecting a Sybil suspect



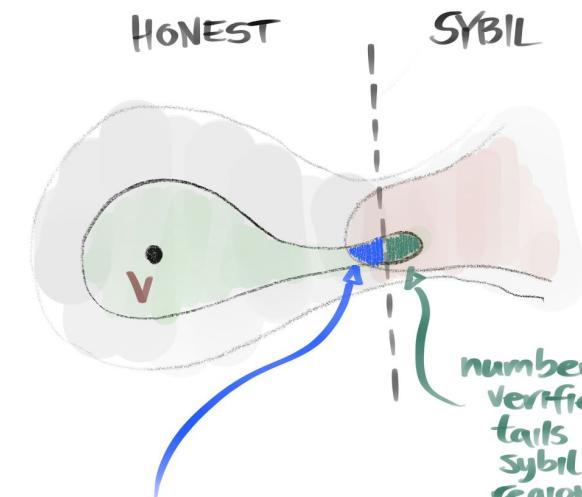
Rejecting a Sybil suspect



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Rejecting a Sybil suspect



number of
verifier tails
in
honest region
reachable
by
sybils

=

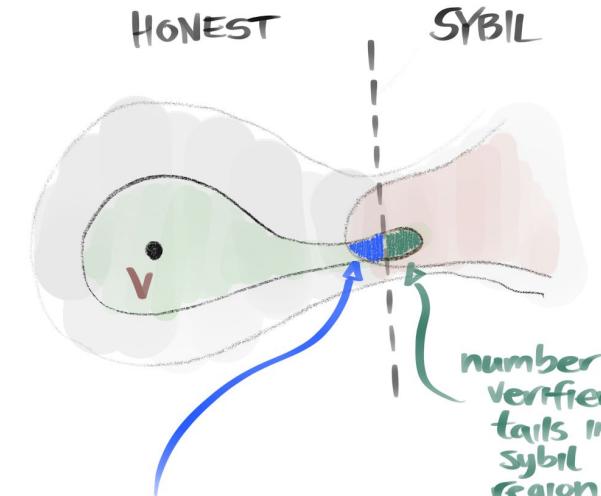
total number of
honest tails
reachable by
sybils

• prob
Verifier
picks
a tail
• total number
of
Verifier
tails

number of
Verifier
tails in
sybil
region

= total # of
Verifier
tails • prob
tail escapes
in
sybil region

Rejecting a Sybil suspect



number of verifier tails in honest region reachable by sybils

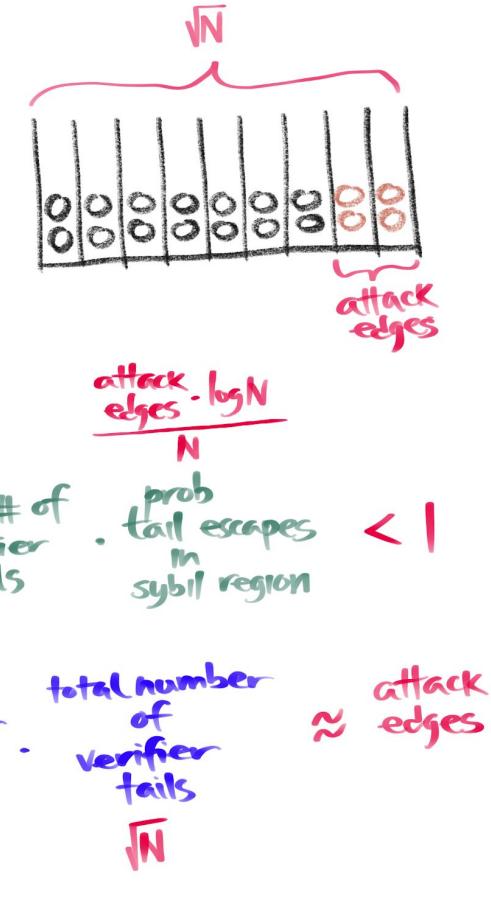
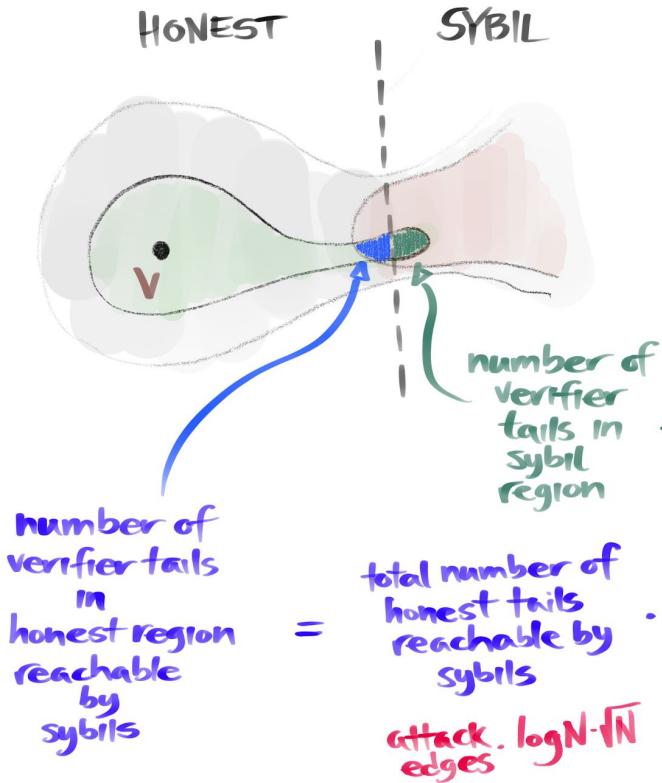
$$=$$

total number of honest tails reachable by sybils

attack. $\log N \cdot \sqrt{N}$ edges

$$\frac{\sqrt{N}}{\text{total # of verifier tails}} = \frac{\text{prob verifier picks a tail}}{\frac{1}{N}} \cdot \frac{\text{total number of verifier tails}}{\frac{\text{attack edges} \cdot \log N}{N}} \cdot \frac{\text{prob tail escapes in sybil region}}{\frac{N}{\text{attack edges}}} < 1$$

Rejecting a Sybil suspect



The end

Thank you!
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