

Optimal block execution of trades

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Problem statement, first stab

Traditionally, OTC service by brokers to institutional clients.

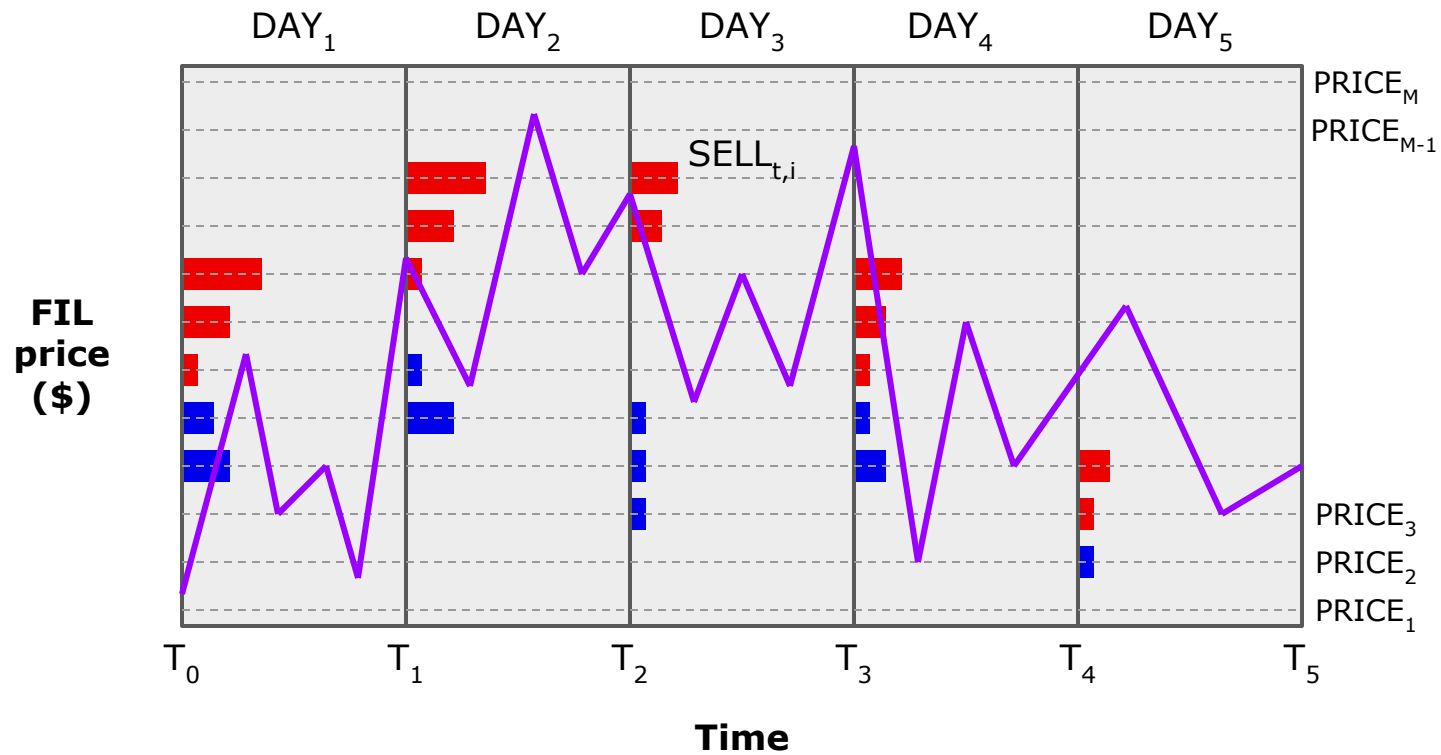
Electronic version, around 2005.

Now crypto OTC services.

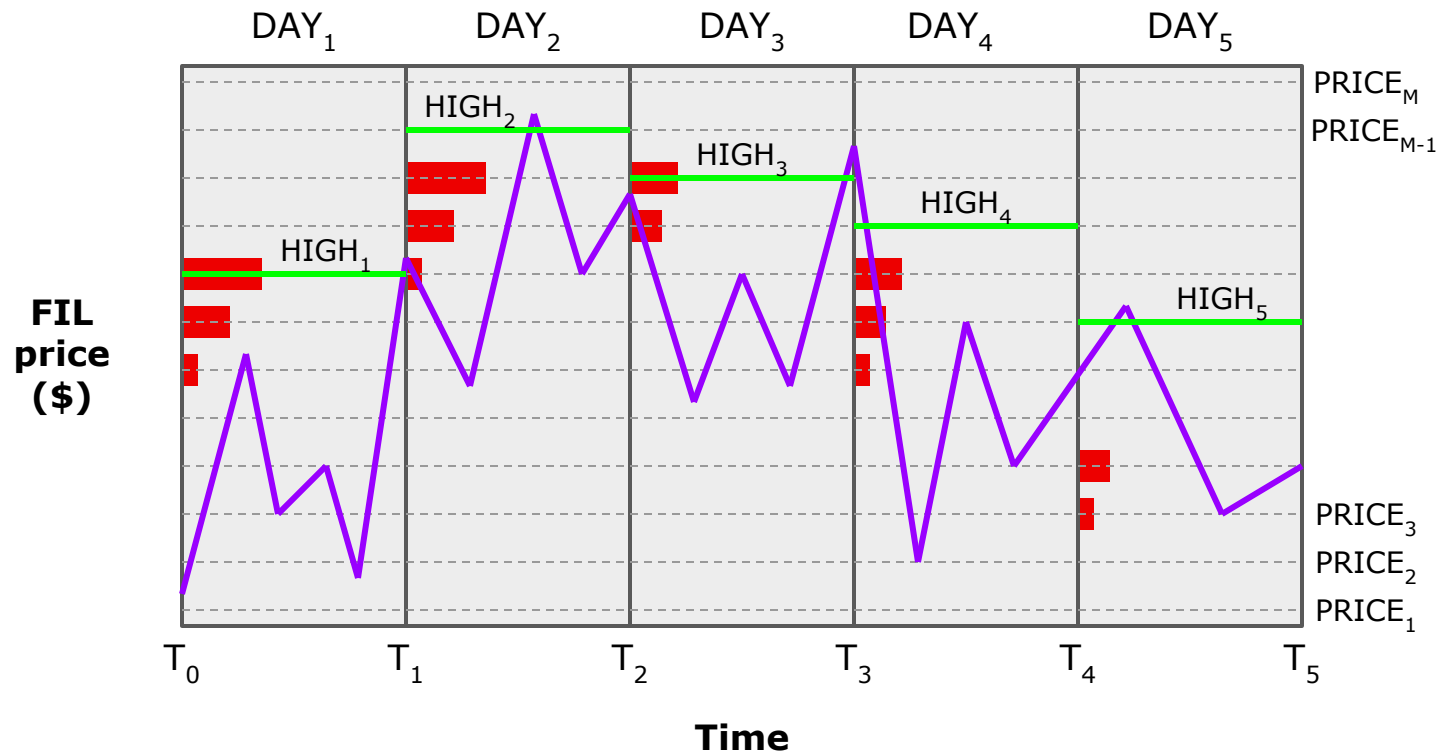
**Sell 1 FIL before a deadline for as much USD as possible,
given a (probabilistic) model of FIL/USD.**

(Market Making is block execution on both sides of the market.)

Model of a trade execution



If only selling, HIGH PRICE between syncs determines limit order matches.



Assumptions and objective

Assume given $PROB_i$: Probability that the next day highest price is $PRICE_i$.

Let USD_t = dollars earned from selling FIL 1 between T_t and T_{LAST} using strategy.

Objective:

the unknowns
(the strategy)



maximize $E[USD_0]$ over the choices of $SELL_{i,t}$

such that $SELL_{i,t} \geq 0$ for all i, t
and $\sum_i SELL_{i,t} \leq 1$ for all t

Solution

$$\text{USD}_{\text{LAST}} = 0$$

$$\text{USD}_t = \text{"matched orders between } T_t \text{ and } T_{t+1}\text{"} + \text{"unsold FIL"} \cdot \text{USD}_{t+1}$$

$$= \text{"orders with } \text{SELL}_{i,t} \leq \text{HIGH}_t\text{"} + \text{"unsold FIL"} \cdot \text{USD}_{t+1}$$

$$= \sum_{p:\text{price}} [\text{HIGH}_t = \text{PRICE}_p] \cdot \left(\sum_{i \leq p} \text{PRICE}_i \cdot \text{SELL}_{t,i} + (1 - \sum_{i \leq p} \text{SELL}_{t,i}) \cdot \text{USD}_{t+1} \right)$$

indicator variable:
1 if condition met,
0 otherwise

proceeds from
matched sell
orders

remaining FIL
(not matched)

$$E[\text{USD}_t] = \sum_{p:\text{price}} \text{PROB}_p \cdot \left(\sum_{i \leq p} \text{PRICE}_i \cdot \text{SELL}_{t,i} + (1 - \sum_{i \leq p} \text{SELL}_{t,i}) \cdot E[\text{USD}_{t+1}] \right)$$

maximize $E[\text{USD}_t]$ over $\text{SELL}_{t,i}$ using linear programming

Implementation intermission

github.com/petar/BlockExecExpMax.jl

Discussion and improvements

Caveat: Maximizing expectation, $E[\text{USD}_0]$, can produce bi-modal (all-or-nothing) outcomes.

Intuitively, we want to maximize expectation **and** minimize variance.

However, minimizing variance contradicts maximizing expectation.

So, we must embed our preferences in a single objective function.

Sharpe ratio: $E[\text{USD}_0] / \text{StdDev}[\text{USD}_0]$

Weighted: $E[\text{USD}_0] - 3 * \text{StdDev}[\text{USD}_0]$

Optimization challenges

Maximizing Sharpe (or any variance-based objectives) is **not** linear programming.

$\text{maximize}_x \text{ Sharpe}$ **and** $\text{maximize}_x \text{ Sharpe}^2$ **find the same x.**

$$\begin{aligned}\text{Sharpe}^2 &= E[\text{USD}_0]^2 / \text{StdDev}[\text{USD}_0]^2 \\ &= E[\text{USD}_0]^2 / \text{Var}[\text{USD}_0] \\ &= E[\text{USD}_0]^2 / (E[\text{USD}_0^2] - E[\text{USD}_0]^2)\end{aligned}$$

Therefore, the Sharpe maximization is **rational** and **quadratic** (in the unknowns $\text{SELL}_{t,i}$).

Possible approaches

Linear/quadratic/conic/semidefinite programming find global optima, but are harder to work with: objective needs to be rewritten cleverly.

Alternatively, one could find an optimum for the objective (e.g. Sharpe) using gradient methods like **(Stochastic) Gradient Descent**. Global optimum is not always found, but no need to rewrite the problem objective.

Takeaway

Every “control” problem can be somehow written as a mathematical optimization problem:

maximize $f(x)$, subject to $a(x) \geq 0$, $b(x) \geq 0$, etc.

With a little effort you can find a decent solution to any optimization problem.

(Use Julia as a one-stop-shop for math.)