

Problem Set 0

Name: Christopher Jeff Sanchez

Problem 0-1.

$$\begin{aligned}
 A &= \left\{ i + \binom{5}{i} \mid i \in \mathbb{Z}, 0 \leq i \leq 4 \right\} \\
 &= \left\{ 0 + \binom{5}{0}, 1 + \binom{5}{1}, 2 + \binom{5}{2}, 3 + \binom{5}{3}, 4 + \binom{5}{4} \right\} \\
 &= \{1, 6, 12, 13, 8\}
 \end{aligned} \tag{1}$$

$$\begin{aligned}
 B &= \{3i \mid i \in \{1, 2, 3, 4, 5\}\} \\
 &= \{3, 6, 9, 12, 15\}
 \end{aligned} \tag{2}$$

(a) $A \cap B = \{6, 12\}$

(b) $|A \cup B| = 7$

(c) $|A - B| = 3$

Problem 0-2.

$$\begin{aligned}
 X &= \{\text{\# of heads in three coin flips}\} \\
 &= \{\text{HHH, HHT, HTH, THH, TTH, THT, HTT, TTT}\} \\
 &= \{3, 2, 2, 2, 1, 1, 1, 0\}
 \end{aligned} \tag{3}$$

$$\begin{aligned}
 Y &= \{\text{products of two six-sided dice}\} \\
 &= \{1*1, 1*2, 2*1, \dots, 6*5, 6*6\} \\
 &= \begin{matrix} & \begin{matrix} 1 & 2 & 3 & 4 & 5 & 6 \end{matrix} \\ \begin{matrix} 1 \\ 2 \\ 3 \\ 4 \\ 5 \\ 6 \end{matrix} & \begin{matrix} 1 & 2 & 3 & 4 & 5 & 6 \\ 2 & 4 & 6 & 8 & 10 & 12 \\ 3 & 6 & 9 & 12 & 15 & 18 \\ 4 & 8 & 12 & 16 & 20 & 24 \\ 5 & 10 & 15 & 20 & 25 & 30 \\ 6 & 12 & 18 & 24 & 30 & 36 \end{matrix} \end{matrix}
 \end{aligned} \tag{4}$$

(a) $E[X] = (3 \cdot 1 + 2 \cdot 3 + 1 \cdot 3)/8 = 12/8 = 1.5$

(b) $E[Y] = 426/36 = 11.\overline{833}$

(c) $E[\{X + Y\}] = 438/44 = 9.\overline{954}$

Problem 0-3.

$$A = 600/6 = 100 \quad (5)$$

$$B = 60 \bmod 42 = 17 \quad (6)$$

(a) $A \bmod 2 = 0, B \bmod 2 = 0, \therefore A \equiv B \pmod{2}$

(b) $A \bmod 3 = 1, B \bmod 3 = 0, \therefore A \not\equiv B \pmod{3}$

(c) $A \bmod 4 = 0, B \bmod 4 = 2, \therefore A \not\equiv B \pmod{4}$

Problem 0-4. Prove by induction that $\sum_{i=1}^n i^3 = [\frac{n(n+1)}{2}]^2$, for any $n \geq 1$.

Proof. Let $P(n) : \sum_{i=1}^n i^3 = [\frac{n(n+1)}{2}]^2$.

Base case, $P(1)$:

$$\sum_{i=1}^1 i^3 = 1 \quad (7)$$

$$\left[\frac{1(1+1)}{2} \right]^2 = 1 \quad (8)$$

Hence, base case is true. For the induction step, assuming $P(n)$ is true, we get $P(n+1)$:

$$\sum_{i=1}^{n+1} i^3 = (n+1)^3 + \sum_{i=1}^n i^3 \quad (9)$$

$$= (n+1)^3 + \left[\frac{n(n+1)}{2} \right]^2 \quad (10)$$

$$= \frac{4(n+1)^3 + n^2(n+1)^2}{4} = \frac{(4(n+1) + n^2)(n+1)^2}{4} \quad (11)$$

$$= \frac{(n^2 + 4n + 1)(n+1)^2}{4} = \frac{(n+2)^2(n+1)^2}{4} \quad (12)$$

$$= \left[\frac{(n+1)((n+1)+1)}{2} \right]^2 \quad (13)$$

■

Problem 0-5. Prove **by induction** that every connected undirected graph $G = (V, E)$ for which $|E| = |V| - 1$ is acyclic.

Recall:

- A graph is connected if each pair of vertices has at least one connecting edge.
- An undirected graph simply means that all edges has no direction component.
- An acyclic graph is a graph that contains no cycles. When traversing the graph from vertex to vertex, no same vertex shall be visited twice.
- $|E| = |V| - 1$ means that there is one less edge as there are vertices.

Proof. Using induction on the number of vertices, $n = |V|$, in a graph $G = (V, E)$. Define $P(n)$ as: "Given that G is a connected undirected graph, if $|E| = |V| - 1$, then G is acyclic."

Base case, $P(1)$: $|V| = 1$, hence $|E| = 0$. A graph with only one vertex can't create a cycle since a cycle requires a nonempty sequence of edges. Thus the base case holds.

For the induction step, we consider a "shrink-down, build-up" approach: ¹ Consider an $(n+1)$ vertex graph G' that is connected, undirected, and has $|V| = n + 1$, and $|E| = n$. We know that any connected graph has a spanning tree, and a leaf of that spanning tree can be removed such that we are left with a graph G that is still connected, undirected, and has $|V| = n$, and $|E| = n - 1$. If we assume $P(n)$, then G is acyclic. Now we see if we can get G' from G by adding a vertex.

Case 1: Connecting the vertex v_{n+1} to an edge in G . This is not possible since all edges in G are already connected to 2 vertices (since it is a connected graph with one less edge than vertices).

Case 2: Connecting the edge e_n to two vertices already in G . This is not possible since if $e_n = \{v_{n-1}, v_n\}$, then by case 1, there won't be any edge for v_{n+1} and it won't be a part of the connected graph G .

Case 3: Connecting the new vertex and edge together with a vertex in G . That is, $e_n = \{v_n, v_{n+1}\}$. Since G is said to be acyclic, we know that v_n doesn't belong to any prior cycle. For v_n and v_{n+1} to create a new cycle, there should be a path from v_n to v_{n+1} and back to v_n . Since v_{n+1} is only connected to e_n , then there is no path from v_{n+1} back to v_n without repeating the edge. Therefore, the two vertices don't create a new cycle.

By considering the three cases, we see that G remains acyclic. ■

¹This is to avoid the "Build-up Error," where one assumes that every size $n+1$ graph with some property can be built-up from a size n graph with the same property. This assumption is not true by default. Note that the property in question is that of the antecedent, which is necessary to assume $P(n)$.

Problem 0-6. Submit your implementation to `alg.mit.edu`.

```
1 def count_long_subarray(A):
2     '''
3     Input: A      | Python Tuple of positive integers
4     Output: count | number of longest increasing subarrays of A
5     '''
6     count = 0
7     prev_a = 0
8     current_subarray_size = 0
9     max_subarray_size = 0
10
11    for a in A:
12        print("a =", a)
13
14        if a > prev_a: # while increasing
15            # keep track of length of increasing-subarray
16            current_subarray_size += 1
17            print("subarray_size =", current_subarray_size)
18            prev_a = a
19
20            if current_subarray_size > max_subarray_size:
21                max_subarray_size = current_subarray_size
22                count = 1
23                print(
24                    "update max_subarray_size =",
25                    max_subarray_size,
26                    "reset count =",
27                    count,
28                )
29            elif current_subarray_size == max_subarray_size:
30                count += 1
31                print("increment count =", count)
32
33            continue
34
35            # if not increasing anymore, end of subarray
36            prev_a = a # the update in the 1st if-case won't be executed, update here
37            current_subarray_size = 1 # 1 because current_subarray = [ prev_a ]
38            print("reset current_subarray_size =", 1)
39
40    print("COUNT =", count)
41    return count
```