

**P802.1ASdm/D1.2**

**August 18, 2023**

(Amendment to  
IEEE Std 802.1AS™-2020

as amended by

IEEE 802.1AS-2020/Cor 1-2021, P802.1ASdn, and P802.1ASdr)

## **Draft Standard for Local and metropolitan area networks—**

# **Timing and Synchronization for Time-Sensitive Applications**

## **Amendment: Hot Standby**

Sponsor

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of the  
IEEE Computer Society**

### **Time-Sensitive Networking Task Group of IEEE 802.1**

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This draft standard is an amendment. The scope of changes to the base standard is thus strictly limited, as detailed in the [PAR](#).

Information on participation in this project, and in the IEEE 802.1 Working Group can be found [here](#).

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## Project Authorization Request, Scope, Purpose, and Criteria for Standards Development (CSD)

The complete amendment PAR, as approved by IEEE NesCom 3rd June 2020, can be found at:

<https://development.standards.ieee.org/myproject-web/public/view.html#pardetail/8208>

**<<Editor's Note: The above date of approval of the amendment PAR and link will be updated when available. The text of the Scope of the Proposed changes and the Need for the Project reflect the updated PAR text approved by 802.1.>>**

The 'Scope of the Proposed changes' and the 'Need for the Project' specify the changes to be made by this amendment (see below).

### Scope of the Proposed changes:

This amendment specifies protocols, procedures, and managed objects for hot standby without use of the Best Master Clock Algorithm (BMCA), for time-aware systems, including:

- A function that transforms the synchronized times of two generalized Precision Time Protocol (gPTP) domains into one synchronized time for use by applications;
- A function that directs the synchronized time of one gPTP domain into a different gPTP domain; and
- Mechanisms that determine whether a gPTP domain has sufficient quality to be used for hot standby. This amendment specifies a Type-Length-Value (TLV) to enable 1  $\mu$ s time synchronization accuracy over 100 network hops. This amendment also addresses errors and omissions in the description of existing functionality.

### Need for the Project:

Hot standby is needed in some applications that use time synchronization (e.g., industrial automation and automotive in-vehicle networks). A TLV is needed by IEC/ IEEE 60802 Time-Sensitive Networking Profile for Industrial Automation to achieve the goal of 1  $\mu$ s time synchronization accuracy over 100 network hops while using existing silicon and low-cost crystal oscillators, i.e., not, for example, temperature compensated crystal oscillators.

### Maintenance items included in this amendment

This amendment includes 802.1 maintenance items 263, 278, 279, 280, 281, 283, 289, 292, 293, 294, 297, 307, 308, 327, 328, 331, 332, 334, and 335. These maintenance items are documented at <https://www.802-1.org/items?utf8=?&search=802.1AS&open=1&commit=Search>. The maintenance items do the following:

- a) Maintenance item 263 clarifies that a Signaling message can contain more than one message interval request TLV, but that the gPTP capable TLV is sent in a separate Signaling message.
- b) Maintenance item 278 makes it clear that the message interval request TLVs and state machines are optional (but the gPTP capable TLV and state machine is mandatory), and fixes inconsistencies between clause 5 and the PICS related to this.
- c) Maintenance item 279 fixes a typo, by changing externalPortConfiguration to externalPortConfigurationEnabled in one place.
- d) Maintenance item 280 updates some terminology in MIB descriptions.
- e) Maintenance item 281 fixes a typo, by changing Delay\_Req to Pdelay\_Req in one place.
- f) Maintenance item 283 changes "domain number" to "domainNumber."
- g) Maintenance item 289 removes several incorrect references to 11.2.13.
- h) Maintenance items 292 and 293 fix some capitalization.
- i) Maintenance item 294 corrects the description of the ptpTimescale flag in Table 10-9 for the case of domain 0.

- j) Maintenance item 297 makes editorial corrections to Table 11-6 and PICS item MINTA-6.
- k) Maintenance item 307 fixes a type by changing portEnabledOper to portOper in Figure and Figure .
- l) Maintenance item 308 fixes incorrect logic in the MDSyncSendSM state machine.
- m) Maintenance item 327 removes the undefined variable domainEnabled from the GtpCapableTransmit and GtpCapableReceive state machines.
- n) Maintenance item 328 corrects a comment that is incorrect in functions used by the GtpCapableIntervalSetting, AnnounceIntervalSetting, SyncIntervalSetting, and LinkDelayIntervalSetting state machines.
- o) Maintenance item 331 corrects the computation of meanLinkDelay for the case where CMLDS is used.
- p) Maintenance item 332 corrects the setting of the variable newInfo in the PortStateSettingExt state machine so that an Announce message is sent immediately when the port state is set to MasterPort.
- q) Maintenance item 334 corrects the error in the computation of clockSourcePhaseOffset in Figure 10-6.
- r) Maintenance item 335 clarifies the difference between currentTime and localTime used in 10.2.9.2.1.
- s) Maintenance item 360 adds ranges and defaults for syncReceiptTimeout and announceReceiptTimeout in 10.7.3.1 and 10.7.3.2, respectively, and fixes the range for gPtpCapableReceiptTimeout in 10.7.3.3.

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Draft Standard for Local and metropolitan area networks—Timing and Synchronization for Time-Sensitive Applications—  
Amendment: Hot Standby

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**Abstract:** This amendment to IEEE Std 802.1AS™-2020 specifies hot standby, and addresses errors and omissions in the description of existing functionality.

**Keywords:** hot standby, best master, frequency offset, Grandmaster Clock, Grandmaster PTP Instance, PTP End Instance, PTP Relay Instance, IEEE 802.1AS™, phase offset, synchronization, syntonization, time-aware system

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## Introduction

This introduction is not part of IEEE Std 802.1ASdm<sup>TM</sup>-20xx, IEEE Standard for Local and metropolitan area networks—Timing and Synchronization for Time-Sensitive Applications—Amendment: Hot Standby

The first edition of IEEE Std 802.1AS was published in 2011. A first corrigendum, IEEE Std 802.1AS<sup>TM</sup>-2011/Cor1-2013, provided technical and editorial corrections. A second corrigendum, IEEE Std 802.1AS<sup>TM</sup>-2011/Cor2-2015 provided additional technical and editorial corrections.

The second edition, IEEE Std 802.1AS-2020, added support for multiple gPTP domains, Common Mean Link Delay Service, external port configuration, and Fine Timing Measurement for 802.11 transport. Backward compatibility with IEEE Std 802.1AS-2011 was maintained. A corrigendum, IEEE Std 802.1AS<sup>TM</sup>-2020/Cor1-2021, provides technical and editorial corrections.

This amendment to IEEE Std 802.1AS<sup>TM</sup>-2020 specifies hot standby, and addresses errors and omissions in the description of existing functionality. Hot standby guards against the failure of a single Grandmaster PTP Instance or the failure of communication from that Grandmaster PTP Instance to a Clock Target.

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# Draft IEEE Standard for Local and metropolitan area networks— Timing and Synchronization for Time- Sensitive Applications Amendment: Hot Standby

[This amendment is based on IEEE Std 802.1AS™-2020 as modified by IEEE Std 802.AS-2020/Cor1, P802.1ASdn, and P802.1ASdr.]

NOTE—The editing instructions contained in this amendment define how to merge the material contained therein into the existing base standard and its amendments to form the comprehensive standard.

The editing instructions are shown in ***bold italic***. Four editing instructions are used: change, delete, insert, and replace. ***Change*** is used to make corrections in existing text or tables. The editing instruction specifies the location of the change and describes what is being changed by using ~~striketrough~~ (to remove old material) and underscore (to add new material). ***Delete*** removes existing material. ***Insert*** adds new material without disturbing the existing material. Deletions and insertions may require renumbering. If so, renumbering instructions are given in the editing instruction. ***Replace*** is used to make changes in figures or equations by removing the existing figure or equation and replacing it with a new one. Editing instructions, change markings, and this NOTE will not be carried over into future editions because the changes will be incorporated into the base standard.<sup>1</sup>

<sup>1</sup>Notes in text, tables, and figures are given for information only, and do not contain requirements needed to implement the standard.

### 3. Definitions

*Insert the following definitions in Clause 3, and renumber the definitions as appropriate:*

**3.4 default:** The value of an implementation parameter or capability selection, in the absence of an overriding profile standard requirement or explicit configuration.

**3.7 epoch:** The origin of a timescale.

*Change 3.14 as follows:*

**3.14 Grandmaster PTP Instance:** A PTP Instance containing the Grandmaster Clock.

NOTE—None of the PTP Ports of a Grandmaster PTP Instance is in the SlavePort state.

*Change 3.23 as follows:*

**3.23 PTP Link:** Within a domain, a network segment between two PTP Ports using the peer-to-peer delay mechanism of IEEE Std 802.1AS. The peer-to-peer delay mechanism is designed to measure the propagation time over such a link.

NOTE—A PTP Link between PTP Ports of PTP Instances is also a gPTP **C**ommunication **P**ath (see 3.11).

*Insert the following definitions in Clause 3, and renumber the definitions as appropriate:*

**3.24 profile standard:** A standard specifying the capabilities, options, and parameter values of one or more base standards for use in a specific target environment or application.

**3.28 sdoId:** The sdoId attribute is the primary mechanism for providing isolation of PTP Instances operating under a PTP Profile specified by one QSDO from PTP Instances operating under a PTP Profile specified by a different QSDO (see 7.1.3 and 7.1.4 of IEEE Std 1588-2019). The value of 1 for the majorSdoId is restricted to the PTP Profiles developed by the IEEE 802.1. Working Group (see 16.5 of IEEE Std 1588-2019).

**3.33 timescale:** A measure of elapsed time since an epoch.

## 5. Conformance

### *Change 5.4.1 as follows:*

#### 5.4.1 ~~Summary of~~ PTP Instance requirements

An implementation of a PTP Instance shall:

- a) Implement the generalized precision time protocol (gPTP) requirements specified in Clause 8.
- b) Support the requirements for time-synchronization state machines (10.1.2, 10.2.1, 10.2.2, 10.2.3, 10.2.4, 10.2.5, and 10.2.6).
- c) Support at least one PTP Port.
- d) On each supported PTP Port, implement the PortSyncSyncReceive state machine (10.2.8).
- e) Implement the ClockSlaveSync state machine (10.2.13).
- f) ~~Support the following best master clock algorithm (BMCA) requirements:~~
  - 1) ~~Implement the BMCA (10.3.2, 10.3.3, 10.3.4, 10.3.5, 10.3.6, 10.3.8, and 10.3.10).~~
  - 2) ~~For domain 0, implement specifications for externalPortConfigurationEnabled value of FALSE (10.3.1).~~
  - 3) ~~Implement the PortAnnounceReceive state machine (10.3.11).~~
  - 4) ~~Implement the PortAnnounceInformation state machine (10.3.12).~~
  - 5) ~~Implement the PortStateSelection state machine (10.3.13).~~
  - 6) ~~Have the BMCA as the default mode of operation, with externalPortConfiguration FALSE, on domain 0.~~
  - 7) ~~Implement at least one of the possibilities for externalPortConfigurationEnabled (i.e., FALSE, meaning the BMCA is used, and TRUE, meaning external port configuration is used) on domains other than domain 0.~~
- g) Implement the SiteSyncSync state machine (10.2.7).
- h) Implement the state machines related to signaling gPTP capability (~~10.4~~10.4.1, 10.4.2, 10.4.3).

NOTE 1—5.4.1 h) does not include the GptpCapableIntervalSetting state machine (10.4.4)

NOTE 2—The GptpCapableTransmit and GptpCapableReceive state machines are required; however, they can be disabled by setting the managed object gptpCapableStateMachinesEnabled (see 14.8.55) to FALSE.

- i) For receipt of all messages and for transmission of all messages except Announce (see 10.6.3) ~~and Signaling (see 10.6.4)~~, support the message requirements as specified in 10.5, 10.6, and 10.7.

NOTE 2—The support of Signaling messages is mandatory [5.4.1 i)] because the state machines related to the signaling of gPTP capability are mandatory [5.4.1 h)].

- j) Support the performance requirements in B.1 and B.2.4.

#### 5.4.2 PTP Instance options

### *Change 5.4.2 c) as follows:*

- c) Support the following for Grandmaster PTP Instance capability:
  - 1) Support the media-independent master capability specified in item b) of 5.4.2.
  - 2) Support the requirements for a grandmaster-capable PTP Instance (10.1.3).
  - 3) Implement the ClockMasterSyncSend state machine (10.2.9).

- 4) Implement the ClockMasterSyncOffset state machine (10.2.10).
- 5) Implement the ClockMasterReceive state machine (10.2.11).

**Delete 5.4.2 e), and renumber subsequent list items as appropriate.**

**Insert the following after 5.4.2 h), and renumber subsequent list items as appropriate:**

- i) Implement the GtpCapableIntervalSetting state machine.

**Change 5.4.2 k) as follows:**

- k) Support the use of a remote management protocol. A PTP Instance that claims to support remote management shall:
  - 1) State which remote management protocol standard(s) or specification(s) are supported (see A.19).
  - 2) State which standard(s) or specification(s) for managed object definitions and encodings are supported for use by the remote management protocol (see A.19).
  - 3) If the Simple Network Management Protocol (SNMP) is supported as a remote management protocol, support the managed object definitions specified as Structure of Management Information version 2 (SMIv2) Management Information Base (MIB) modules in Clause 15.
  - 4) If YANG is supported with a remote management protocol, support the YANG data model ieee802-dot1as-ptp in Clause 17.
  - 5) If YANG is supported with a remote management protocol, and if hot standby is supported, support the YANG data model ieee802-dot1as-ptp in Clause 17.

<<Editor's note: A few of the managed objects of the hot standby YANG module (ieee802-dot1as-hs) are not specific to hot standby; they are more general. These include, parentDS.gmpresent and portDS.gtpCapableStateMachinesEnabled (these were needed for 60802). Should these be mentioned in k)5) above, since they might be needed even if hot standby is not supported? And, if they are mentioned here, should they also be mentioned in RMGT-5 of A.19? Comments are requested.>>

**Insert the following after 5.4.2 l), and renumber subsequent list items as appropriate:**

- m) Implement hot standby as specified in Clause 18, 14.8, 9.3.3.4, 9.4.3.4, 9.5.3.3, and 9.6.2.6.
- n) Implement the Drift\_Tracking TLV as specified in 10.2.4.25, 10.2.4.26, 10.2.4.27, 11.2.14, 11.2.15, and 11.4.4.

**Insert the following new subclause after 5.4.2, and renumber subsequent subclauses as appropriate:**

### 5.4.3 PTP Instance defaults and recommendations

This standard identifies specific capabilities and parameter values as defaults. ~~Where not further qualified, by reference to the use of other capabilities or by profile standards, these defaults are intended to facilitate deployment in target environments and applications with time-aware system implementations conformant to prior revisions of this standard, including IEEE Std 802.1AS-2020 and IEEE Std 802.1AS-2011. If these defaults are not further qualified by reference to the use of other capabilities or by profile standards, the defaults are intended to facilitate deployment and to provide interoperability in target environments and applications with time-aware system implementations conformant to this standard, including the prior revisions of IEEE Std 802.1AS-2020 and IEEE Std 802.1AS-2011. Mandatory requirements that this standard does not identify as defaults, or otherwise qualify, are not subsetted or modified by profile standards.~~ Mandatory requirements that this standard does not identify as defaults, or call out in some other way, are not subsetted or modified by profile standards.

Changes to capability support or parameter defaults arising from conformance to a profile standard should be identified in the PICS for that profile standard. Changes to defaults by pre-configuration of parameter values by the supplier of a protocol implementation should be identified by Additional Information (see A.3.1) in a completed PICS for this standard.

An implementation of a PTP instance shall, by default, support the following best master clock algorithm (BMCA) requirements:

- a) Implement the BMCA (10.3.2, 10.3.3, 10.3.4, 10.3.5, 10.3.6, 10.3.8, and 10.3.10).
- b) For domain 0, implement specifications for externalPortConfigurationEnabled value of FALSE (10.3.1).
- c) Implement the PortAnnounceReceive state machine (10.3.11).
- d) Implement the PortAnnounceInformation state machine (10.3.12).
- e) Implement the PortStateSelection state machine (10.3.13).
- f) Have the BMCA as the default mode of operation, with externalPortConfigurationEnabled FALSE, on domain 0.
- g) Implement at least one of the possibilities for externalPortConfigurationEnabled (i.e., FALSE, meaning the BMCA is used, and TRUE, meaning external port configuration is used) on domains other than domain 0.

***Change 5.4.4 as follows:***

#### **5.4.4 PTP Relay Instance requirements**

An implementation of a PTP Relay Instance shall:

- a) Support more than one PTP Port.
- b) ~~Support the PTP Instance requirements specified in 5.4.~~
- c) Support the media-independent master capability specified in item b) of 5.4.2.

### **5.5 MAC-specific timing and synchronization methods for full-duplex IEEE 802.3 links**

***Change 5.5 as follows:***

An implementation of a time-aware system with IEEE 802.3 media access control (MAC) services to physical ports shall:

- a) Support full-duplex operation, as specified in 4.2 and Annex 4A of IEEE Std 802.3-2018.
- b) Support the requirements as specified in Clause 11.
- c) ~~Implement the SyncIntervalSetting state machine (10.3.18).~~
- d) Support two-step capability on receive as specified in 11.2.14.
- e) Support two-step capability on transmit as specified in 11.2.15.

An implementation of a PTP Instance with IEEE 802.3 MAC services to physical ports may:

- f) Support asymmetry measurement mode as specified in 10.3.12, 10.3.13, 10.3.16, 11.2.14, 11.2.15, 11.2.19, and 14.8.45.
- g) Support one-step capability on receive as specified in 11.2.14.

- h) Support one-step capability on transmit as specified in 11.2.15.
- i) Support the OneStepTxOperSetting state machine specified in 11.2.16.
- j) Support propagation delay averaging, as specified in 11.2.19.3.4.
- k) Implement the LinkDelayIntervalSetting state machine (11.2.21).



## 7. Time synchronization model for a packet network

### 7.2 Architecture of a time-aware network

#### *Change 7.2.1 as follows:*

##### 7.2.1 General

A time-aware network consists of a number of interconnected time-aware systems that support the gPTP defined within this standard. These time-aware systems can be any networking device, including, for example, bridges, routers, and end stations. A set of time-aware systems that are interconnected by gPTP-capable network elements is called a *gPTP network*. Each instance of gPTP that the time-aware systems support is in one *gPTP domain*, and the instances of gPTP are said to be part of that gPTP domain. A time-aware system can support, and therefore be part of, more than one gPTP domain. The entity of a single time-aware system that executes gPTP in one gPTP domain is called a *PTP Instance*. A time-aware system can contain multiple PTP Instances, which are each associated with a different gPTP domain. There are two types of PTP Instances:

- a) PTP End Instance, which, if not a Grandmaster PTP Instance, is a recipient of time information, and
- b) PTP Relay Instance, which, if not a Grandmaster PTP Instance, receives time information from the Grandmaster PTP Instance (perhaps indirectly through other PTP Relay Instances), applies corrections to compensate for delays in the local area network (LAN) and the PTP Relay Instance itself, and retransmits the corrected information.

This standard defines mechanisms for delay measurements using standard-based procedures for the following:

- c) IEEE 802.3 Ethernet using full-duplex point-to-point links (11)
- d) IEEE 802.3 EPON links (Clause 13)
- e) IEEE 802.11 wireless (Clause 12)
- f) Generic coordinated shared networks (CSNs, e.g., MoCA and G.hn) (Clause 16)

This standard specifies time distribution mechanisms that are tolerant to some faults if the network paths are redundant. One of these time distribution mechanisms updates the distribution path in reaction to changes, whereas the other one relies on redundant PTP instances on top of the redundant network paths.

##### 7.2.2 Time-aware network consisting of a single gPTP domain

#### *Change the second paragraph of 7.2.2 as follows:*

Any PTP Instance with clock sourcing capabilities can be a potential Grandmaster PTP Instance, and a selection method (the *best master clock algorithm*, or BMCA) ensures that all of the PTP Instances in ~~a~~the gPTP domain use the same Grandmaster PTP Instance.<sup>2</sup> The BMCA is largely identical to that used in IEEE Std 1588-2019, but somewhat simplified. In Figure 7-1 the BMCA process has resulted in the Grandmaster PTP Instance being on the network backbone. If, however, the access network fails, the systems on a local network automatically switch over to one of the potential Grandmaster PTP Instances on the local network that is as least as “good” as any other. For example, in Figure 7-2, the access network link has failed, and a potential Grandmaster PTP Instance that has a GNSS reference source has become the active Grandmaster PTP Instance. As a result, now two gPTP domains exist where there used to be one.

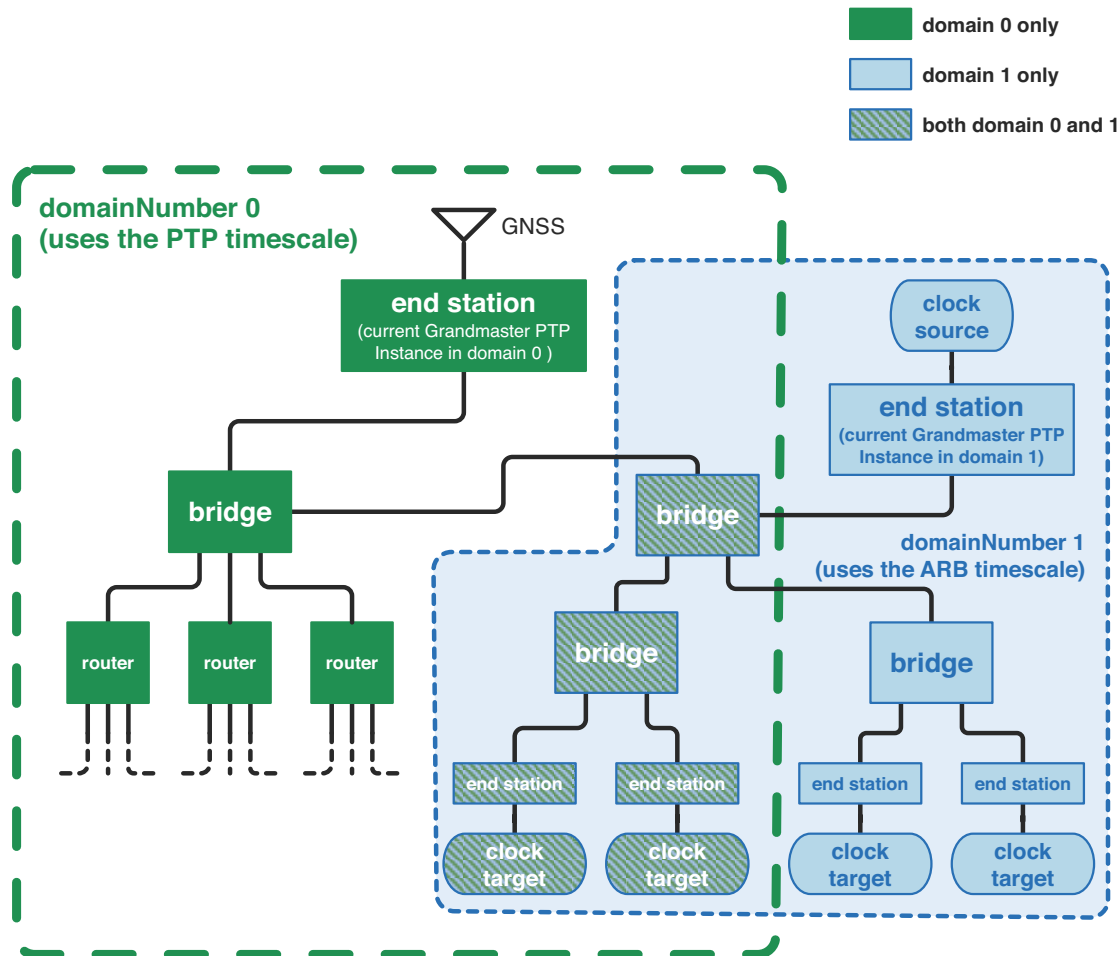
<sup>2</sup> There are, however, short periods during network reconfiguration when more than one Grandmaster PTP Instance might be active while the BMCA process is taking place.

~~Finally, note that when a time-aware system supports more than one domain, one of the domains supported must be domain 0 for backward compatibility with the 2011 edition of this standard, though domain 0 is not necessarily active in a time-aware system.~~

***Remove the last sentence of 7.2.2.***

### **7.2.3 Time-aware network consisting of multiple gPTP domains**

***Replace Figure 7-3 with the following:***



**NOTE 1**—All the “bridges” and “routers” in this figure are examples of time-aware systems that contain at least one PTP Relay Instance, and the end stations are time-aware systems that contain at least one PTP End Instance. The PTP Links in this figure can use any of the media specified in this standard.

**NOTE 2** —This figure differs from the 2020 edition, with Corrigendum 1 applied, of this standard in that “domain number” is replaced by “domainNumber” in two places, ~~and~~ the sentence indicating that the PTP links in the figure can use any of the media specified in this standard is added, and the text at the top of the figure is made a NOTE to this figure.

**Figure 7-3—Time-aware network example for multiple gPTP domains**

*Change 7.2.4 as follows:*

**7.2.4 ~~Time-aware networks with dundant Grandmaster PTP Instances and/or redundant paths~~Time-aware networks using BMCA**

**7.2.4.1 ~~General~~**

~~Redundancy has many levels of sophistication, performance, and cost. Therefore, the appropriate level and/or amount of redundancy required in a time-aware network can be very different for each application. Nonetheless, all solutions for redundancy consist of a detection component, a correction component, and an action component. The detection component detects that something is not working correctly. The correction component determines the appropriate corrective action. The action component performs the required action(s) to fix the detected problem.~~

**7.2.4.2 ~~Redundancy specified in this standard (BMCA)~~**

This standard provides a basic level of redundancy as follows:

- A detection component that triggers when no Sync or Announce messages are received for a defined period of time ~~the current Grandmaster PTP Instance stops working (i.e., loss of Sync messages and Announce messages for a period of time) or if the link to the Grandmaster PTP Instance goes down (i.e., immediate loss of Sync messages and Announce messages).~~
- A correction component that triggers the Best Master Clock Algorithm (BMCA) and the sending of Announce messages so that a new Grandmaster PTP Instance can be elected.
- An action component, where the winning Grandmaster PTP Instance starts sending Announce messages and Sync messages ~~and all the PTP Instances listen to this new Grandmaster PTP Instance.~~

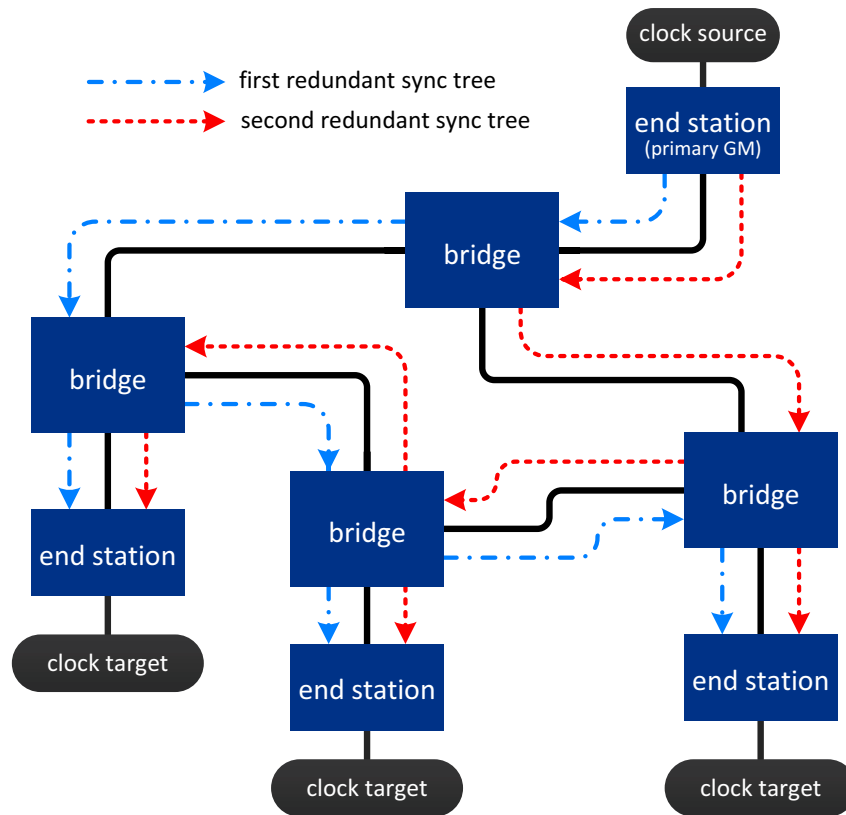
**7.2.4.3 ~~Redundacy not fully specified in this standard~~**

In addition to providing the basic level of redundancy, this standard provides the ability to support more sophisticated network configurations that provide additional levels of Grandmaster PTP Instance and clock path redundancy. Figure 7-4 ~~through and Figure 7-6~~ Figure 7-5 are examples of such networks that provide these additional levels of redundancy as a way to deal with these failures. The information necessary to implement and configure these network configurations is contained in this standard.

In order to take advantage of these failure correction configurations, new types of fault detection are required. The category of fault detection where a Grandmaster PTP Instance completely fails and stops sending clock information is supported as mentioned above.

Other types of faults involve instability of the Grandmaster Clock, such as time glitches, excess jitter, or wander, or various other impairments that could occur in the Grandmaster Clock. Techniques for identifying these types of failures, and the appropriate correction necessary, are not specified in this standard. However, if other techniques or standards are used for detection and correction of these types of failures, this standard provides the means to recover from these errors.

Figure 7-4 shows an example network realizing two redundant synchronization trees from a single GM, with each synchronization tree in a different gPTP domain (there are a total of two gPTP domains).



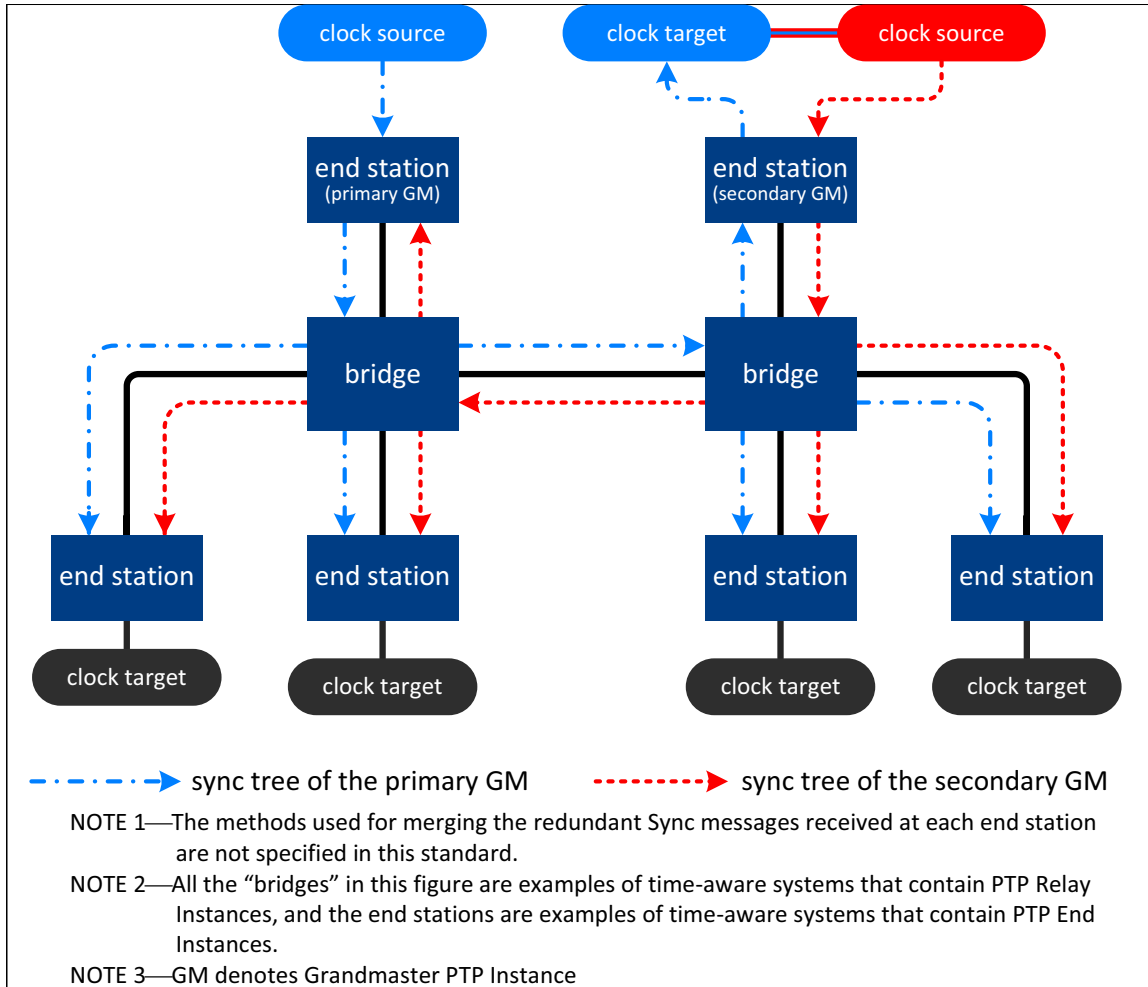
NOTE 1—The methods used for merging the redundant Sync messages received at each end station are not specified in this standard.

NOTE 2—All the “bridges” in this figure are examples of time-aware systems that contain PTP Relay Instances, and the end stations are examples of time-aware systems that contain PTP End Instances.

NOTE 3—GM denotes Grandmaster PTP Instance

**Figure 7-4—Time-aware network example for synchronization path redundancy, with one clock source providing time to two domains**

Figure 7-5 shows an example network with two redundant GMs, one as primary GM and the other as secondary GM, where each GM has one of the two redundant synchronization trees originating from it. This example supports hot-standby operating mode. In this mode, the secondary GM has to be synchronized to the primary GM, because it is part of the synchronization tree of the primary GM as shown in the figure.



**Figure 7-5—Time-aware network example for GM redundancy with one primary GM and one secondary GM, which are separated in two gPTP domains**

~~Figure 7-6 shows another example network, which involves ring topology, using the redundancy features of both Figure 7-4 and Figure 7-5.~~

~~For the techniques shown in the examples of Figure 7-4, Figure 7-5, and Figure 7-6, the detection component, correction component, and action component are not fully specified in this standard.~~

***Delete Figure 7-6.***

***Add the following new subclause after 7.2.4, and renumber figures as necessary:***

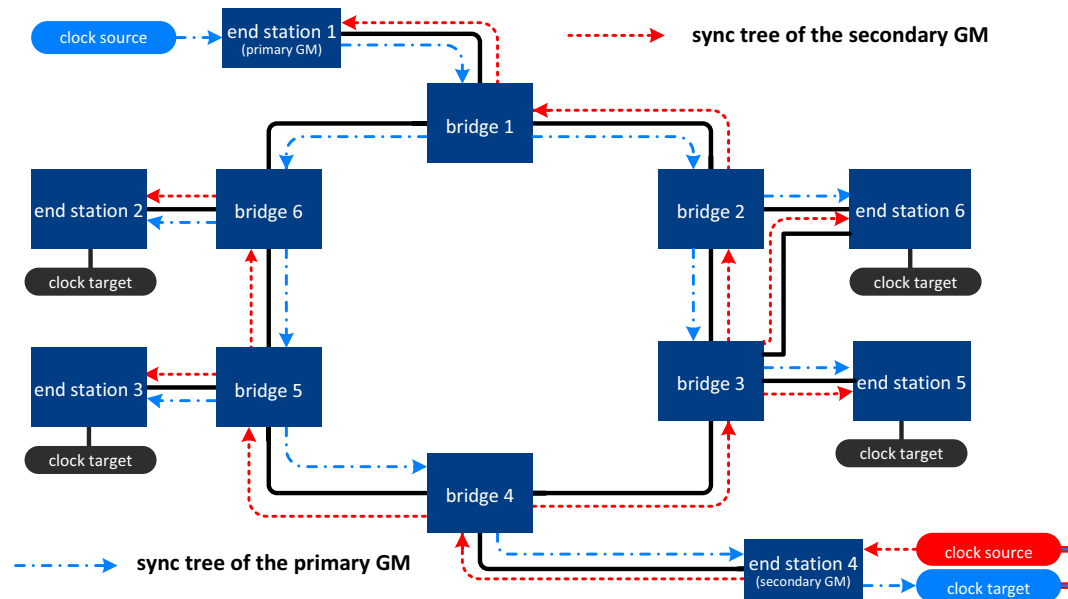
### **7.2.5 Time-aware network with hot standby**

Figure 7-6 shows an example of a time-aware network with hot standby (see Clause 18). The network has GM redundancy and path redundancy to all the bridges and end stations 1, 4, and 6 (path redundancy for end stations 2, 3, and 5 is possible, but is not illustrated to avoid clutter in the figure). In addition, for simplicity the optional split functionality (see 18.5.3.4) is assumed to be disabled.

End station 1 is the primary GM (see Clause 18), and its synchronization spanning tree is illustrated by the blue arrows. End station 4 is the secondary GM, and its synchronization spanning tree is illustrated by the red arrows. Except for the links to end stations 2, 3, and 5, the network can tolerate the failure of any single link or GM and the bridges and end stations remain synchronized by the same GM. For example, if the link between bridge 3 and end station 6 fails, end station 6 is still synchronized by the blue domain via bridge 2. If the link between bridge 5 and bridge 6 fails, bridge 6 is still synchronized by the blue domain and bridge 5 is still synchronized by the red domain. If the secondary GM fails, all the bridges and end stations are still synchronized by the primary GM.

If the primary GM fails, all the bridges and end stations are still synchronized by the secondary GM. In this case, if syncLocked is TRUE the primary PTP Instance of end station 4 stops receiving Sync messages and isSynced changes to FALSE; the hotStandbySystemState of end station 4 is then change to NOT\_REDUNDANT (see Clause 18). If syncLocked is FALSE, and assuming the bridges and end stations other than the GMs are not grandmaster capable, the time difference between the primary and secondary domain eventually exceeds the primarySecondaryOffsetThresh (see Clause 18) and the hotStandbySystemState of end station 4 changes to NOT\_REDUNDANT. In both syncLocked cases, the secondary GM no longer takes timing from the primary domain.

If the link between bridge 4 and bridge 5 fails, all the PTP Instances except for bridge 4 and end station 4 are then synchronized by the blue domain. Bridge 4 and end station 4 are synchronized by the red domain; in addition, bridges 1, 2, and 3, and end stations 1, 4, 5, and 6 receive timing on the red domain. Since end station 4 does not receive timing on the blue domain, its hotStandbySystemState changes to NOT\_REDUNDANT, and the red domain GM does not take timing from the blue domain.



NOTE 1—All the “bridges” in this figure are examples of time-aware systems that contain PTP Relay Instances, and the end stations are examples of time-aware systems that contain PTP End Instances.

NOTE 2—GM denotes Grandmaster PTP Instance

**Figure 7-6—Time-aware network example for hot standby with both GM and partial path redundancy**

## 7.4 PTP Instance architecture

*Change the first paragraph of 7.4 and the caption of Figure 7-8 as follows:*

The model of a PTP Instance and its interfaces to higher-layer applications are shown in Figure 7-7. The interfaces are those specified in Clause 9.



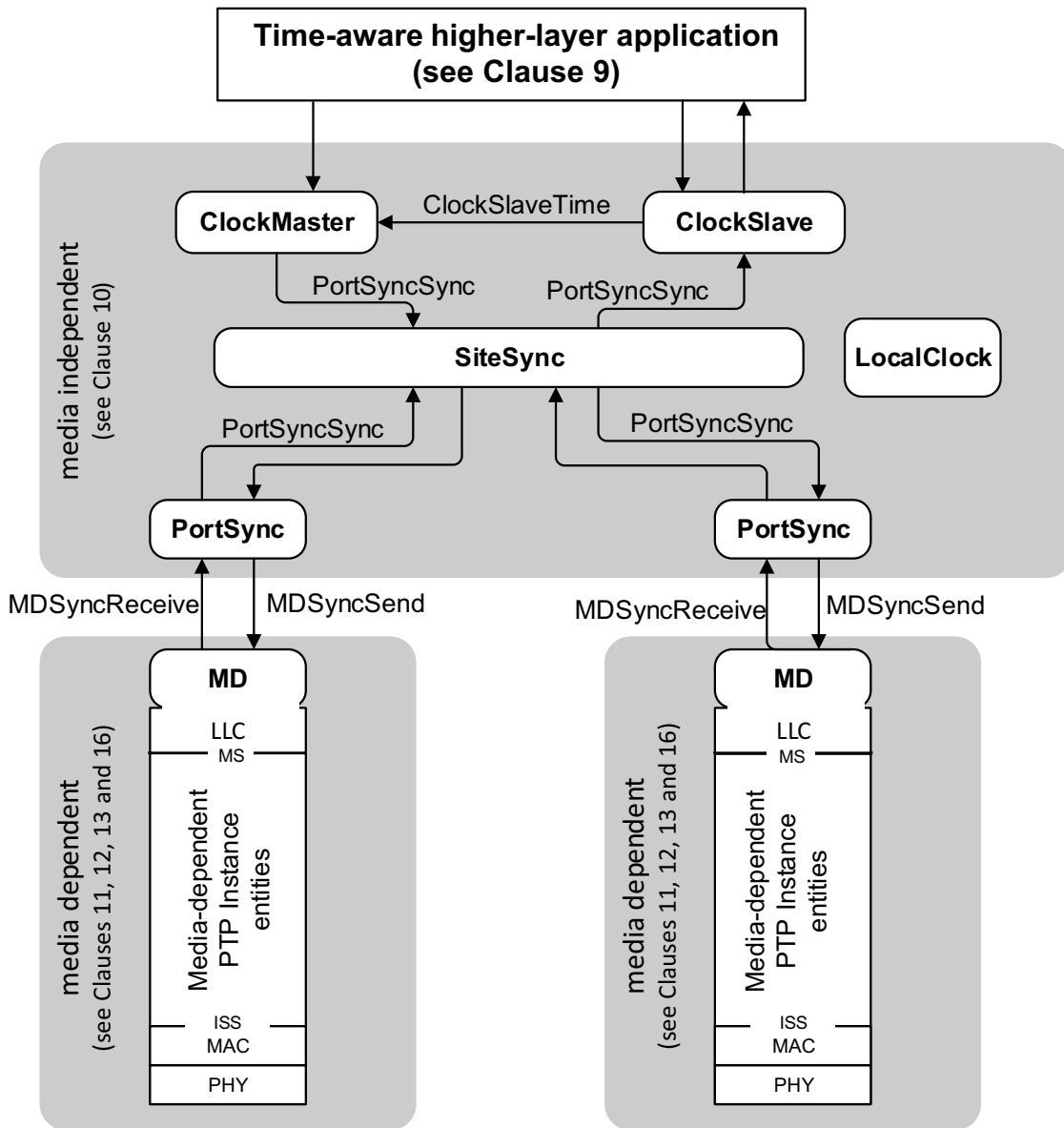


Figure 7-7—Model for PTP Instance and its interfaces to higher-layer applications

## 8. IEEE 802.1AS concepts and terminology

*Change 8.1 as follows:*

### 8.1 gPTP domain

A gPTP domain, hereafter called simply a *domain*, consists of one or more PTP Instances and links that meet the requirements of this standard and communicate with each other as defined by the IEEE 802.1AS protocol. A gPTP domain defines the scope of gPTP message communication, state, operations, data sets, and timescale.

A domain is identified by two attributes: ~~domain-number~~domainNumber and sdoId. The sdoId of a domain is a 12-bit unsigned integer. The sdoId is structured as a two-part attribute as follows:

- The most significant 4 bits are named the majorSdoId, and
- The least significant 8 bits are named the minorSdoId.

A time-aware system shall support one or more domains, each with a distinct ~~domain number~~domainNumber in the range 0 through 127. ~~A time-aware system shall support the domain whose domain-number is 0, and that domain-number shall not be changed to a nonzero value.~~ Unless otherwise specified in this standard, the operation of gPTP and the timescale in any given domain is independent of operation in any other domain.

The value of majorSdoId for a gPTP domain shall be 0x1. The value of minorSdoId for a gPTP domain shall be 0x00.

~~NOTE 1—The above requirements for majorSdoId and minorSdoId are for gPTP domains. The requirements for the Common Mean Link Delay Service (CMLDS) are given in 11.2.17.~~

Both the domainNumber and the sdoId are carried in the common header of all PTP messages (see 10.6.2.2).

NOTE 1—The above requirements for majorSdoId and minorSdoId are for gPTP domains. The requirements for the Common Mean Link Delay Service (CMLDS) are given in 11.2.17.

NOTE 2—In the 2011 edition of this standard, the attribute majorSdoId was named transportSpecific, and its value was specified as 0x1 in 10.5.2.2.1 of Corrigendum 1. The attribute minorSdoId did not exist in the 2011 edition, but its location in the common header was a reserved field, which was specified to be transmitted as 0 and ignored on receipt.

Unless otherwise stated, information in the remainder of this document is per domain.

NOTE 3—In steady state, all PTP Instances in a gPTP domain are traceable to a single Grandmaster PTP Instance.

Since the IEEE 802.1 Working Group has only one single unique sdoId value, the PTP Profile specified in the present standard is isolated from other PTP Profiles (see 16.5.2 of IEEE Std 1588-2019).

### 8.5 Ports

#### 8.5.2 Port identity

*Change 8.5.2.3 as follows:*

### 8.5.2.3 Port number

The portNumber values for the PTP Ports on a time-aware system shall be distinct in the range ~~1, 2, ..., 0x1~~ through 0xFFFE.

The portNumber value 0 is assigned to the interface between the ClockMaster and ClockSource entities (see 10.1 and Figure 10-1). The value 0xFFFF is reserved. The assignment of portNumber 0 to this interface helps to simplify the SiteSync and PortStateSelection state machines; this interface is not a PTP Port.

## 9. Application Interfaces

### 9.2 ClockSourceTime interface

#### 9.2.2 ClockSourceTime.invoke function parameters

*Change 9.2.2.1 as follows:*

##### 9.2.2.1 domainNumber (UInteger8)

This parameter is the ~~domain-number~~domainNumber of the gPTP domain to which this ClockSource entity is providing time.

*Change 9.2.2.3 as follows:*

##### 9.2.2.3 timeBaseIndicator (UInteger16)

The timeBaseIndicator is a ~~binary~~-value that is set by the ClockSource entity. The ClockSource entity changes the value whenever its time base changes. The ClockSource entity shall change the value of timeBaseIndicator if and only if there is a phase or frequency change.

NOTE—While the clock that supplies time to the ClockSource entity can be lost, ~~i.e., the PTP Instance can enter holdover~~, the ClockSource entity itself is not lost. The ClockSource entity ensures that timeBaseIndicator changes if the source of time is lost.

### 9.3 ClockTargetEventCapture interface

#### 9.3.2 ClockTargetEventCapture.invoke parameters

*Change 9.3.2.1 as follows:*

##### 9.3.2.1 domainNumber (UInteger8)

This parameter is the ~~domain-number~~domainNumber of the ClockSlave entity that is requested to provide the synchronized time of the signaled event.

*Change 9.3.3 as follows:*

#### 9.3.3 ClockTargetEventCapture.result parameters

```
ClockTargetEventCapture.result {  
    domainNumber,  
    slaveTimeCallback,  
    gmPresent,  
    isSynced,  
    grandmasterIdentity  
}
```

##### 9.3.3.1 domainNumber (UInteger8)

This parameter is the ~~domain-number~~domainNumber of the ClockSlave entity that is providing the synchronized time of the signaled event.

### 9.3.3.2 slaveTimeCallback (ExtendedTimestamp)

The value of slaveTimeCallback is the time, relative to the Grandmaster Clock, that the corresponding ClockTargetEventCapture.invoke function is invoked.

NOTE—The invocation of the ClockTargetEventCapture.invoke function and the detection of this invocation by the ClockSlave entity are simultaneous in this abstract interface.

### 9.3.3.3 gmPresent (Boolean)

The value of gmPresent is set equal to the value of the global variable gmPresent (see 10.2.4.13). This parameter indicates to the ClockTarget whether a Grandmaster PTP Instance is present.

### 9.3.3.4 isSynced (Boolean)

The value of the the variable isSynced (see 18.4.1.1). This parameter shall be present if the optional hot standby feature is implemented (see Clause 18).

### 9.3.3.5 grandmasterIdentity (ClockIdentity)

The value of grandmasterIdentity is the clockIdentity of the Grandmaster PTP Instance.

<<Editor's note: The data type for grandmasterIdentity was missing; it was inserted here and in other subclauses of Clause 9 where needed.>>

## 9.4 ClockTargetTriggerGenerate interface

### 9.4.2 ClockTargetTriggerGenerate.invoke parameters

*Change 9.4.2.1 as follows:*

#### 9.4.2.1 domainNumber (UInteger8)

This parameter is the ~~domain number~~domainNumber of the ClockSlave entity that is requested to signal an event at the specified time.

*Change 9.4.3 as follows:*

### 9.4.3 ClockTargetTriggerGenerate.result parameters

```
ClockTargetTriggerGenerate.result {  
    domainNumber,  
    errorCondition,  
    gmPresent,  
    isSynced,  
    grandmasterIdentity  
}
```

#### 9.4.3.1 domainNumber (UInteger8)

This parameter is the ~~domain number~~domainNumber of the ClockSlave entity that is triggering an event at the specified time.

**9.4.3.2 errorCondition (Boolean)**

A value of FALSE indicates that the ClockTargetTriggerGenerate.result function was invoked at the time, relative to the Grandmaster Clock, contained in the corresponding ClockTargetTriggerGenerate.invoke function. A value of TRUE indicates that the ClockTargetTriggerGenerate.result function could not be invoked at the synchronized time contained in the corresponding ClockTargetTriggerGenerate.invoke function.

NOTE—For example, the ClockTargetTriggerGenerate.result function is invoked with errorCondition = TRUE if the requested slaveTimeCallback is a time prior to the synchronized time when the corresponding ClockTargetTriggerGenerate.invoke function is invoked. As another example, the ClockTargetTriggerGenerate.result function is invoked with errorCondition = TRUE if a discontinuity in the synchronized time causes the requested slaveTimeCallback to be skipped over.

**9.4.3.3 gmPresent (Boolean)**

~~The value of gmPresent is set equal to the value of the global variable gmPresent (see 10.2.4.13). This parameter indicates to the ClockTarget whether a Grandmaster PTP Instance is present.~~ As specified in 9.3.3.3.

**9.4.3.4 isSynced (Boolean)**

As specified in 9.3.3.4.

**9.4.3.5 grandmasterIdentity (ClockIdentity)**

As specified in 9.3.3.5.

**9.5 ClockTargetClockGenerator interface****9.5.2 ClockTargetClockGenerator.invoke parameters**

*Change 9.5.2.1 as follows:*

**9.5.2.1 domainNumber (UInteger8)**

This parameter is the ~~domain-number~~ domainNumber of the ClockSlave entity that is requested to deliver a periodic clock signal.

*Change 9.5.3 as follows:*

**9.5.3 ClockTargetClockGenerator.result parameters**

```
ClockTargetClockGenerator.result {
    domainNumber,
    slaveTimeCallback,
    isSynced,
    grandmasterIdentity
}
```

**9.5.3.1 domainNumber (UInteger8)**

This parameter is the ~~domain-number~~ domainNumber of the ClockSlave entity that is delivering a periodic clock signal.

**9.5.3.2 slaveTimeCallback (ExtendedTimestamp)**

The value of slaveTimeCallback is the synchronized time of this event.

**9.5.3.3 isSynced Boolean**

As specified in 9.3.3.4.

**9.5.3.4 grandmasterIdentity (ClockIdentity)**

As specified in 9.3.3.5.

**9.6 ClockTargetPhaseDiscontinuity interface**

*Change 9.6.2 as follows:*

**9.6.2 ClockTargetPhaseDiscontinuity.result parameters**

```
ClockTargetPhaseDiscontinuity.result {
    domainNumber,
    gmIdentity,
    gmTimeBaseIndicator,
    lastGmPhaseChange,
    lastGmFreqChange,
    isSynced,
    grandmasterIdentity
}
```

**9.6.2.1 domainNumber (UInteger8)**

This parameter is the ~~domain-number~~domainNumber of the ClockSlave entity that is providing discontinuity information.

**9.6.2.2 gmIdentity (ClockIdentity)**

If gmPresent (see 10.2.4.13) is TRUE, the value of gmIdentity is the ClockIdentity of the current Grandmaster PTP Instance. If gmPresent is FALSE, the value of gmIdentity is 0x0.

**9.6.2.3 gmTimeBaseIndicator (UInteger16)**

The value of gmTimeBaseIndicator is the timeBaseIndicator of the current Grandmaster PTP Instance.

**9.6.2.4 lastGmPhaseChange (ScaledNs)**

The value of the global lastGmPhaseChange parameter (see 10.2.4.16) received from the Grandmaster PTP Instance.

**9.6.2.5 lastGmFreqChange (Float64)**

The value of lastGmFreqChange parameter (see 10.2.4.17) received from the Grandmaster PTP Instance.

**9.6.2.6 isSynced (Boolean)**As specified in 9.3.3.4.**9.6.2.7 grandmasterIdentity (ClockIdentity)**As specified in 9.3.3.5.



## 10. Media-independent layer specification

### 10.2 Time-synchronization state machines

**Change 10.2.1 and Figure 10-2 as follows:**

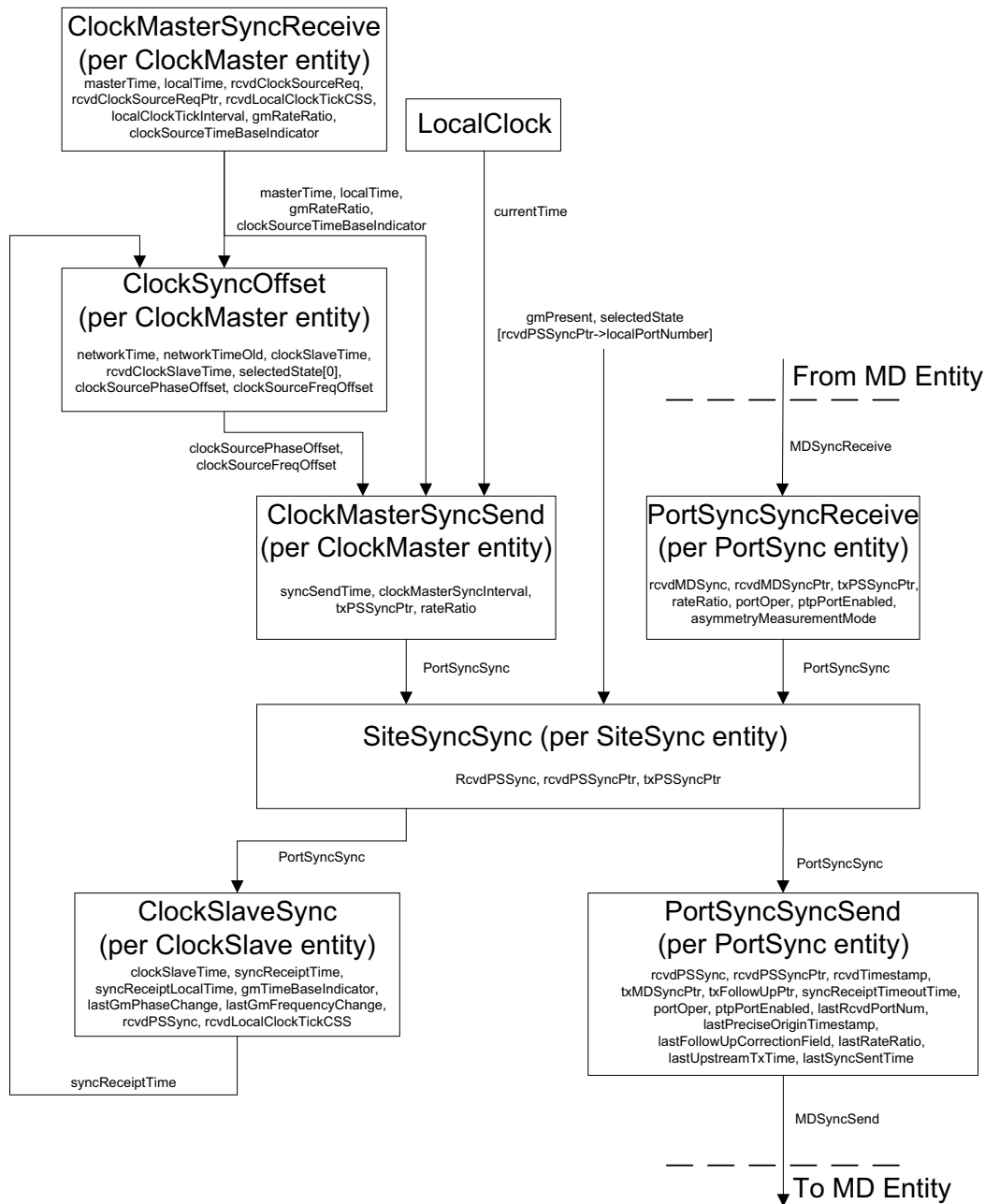
#### 10.2.1 Overview

The time-synchronization function in a PTP Instance is specified by a number of cooperating state machines. Figure 10-2 illustrates these state machines, their local variables, their interrelationships, and the global variables and structures used to communicate between them. The figure indicates the interaction between the state machines and the media-dependent layer and LocalClock entity.

The ClockMasterSyncReceive, ClockMasterSyncOffset, and ClockMasterSyncSend state machines are optional for PTP Instances that are not grandmaster-capable (see 8.6.2.1 and 10.1.3). These state machines may be present in a PTP Instance that is not grandmaster-capable; however, any information supplied by them, via the ClockMasterSyncSend state machine, to the SiteSyncSync state machine is not used by the SiteSyncSync state machine if the PTP Instance is not grandmaster-capable.

The media-independent layer state machines in Figure 10-2 are as follows:

- a) ClockMasterSyncReceive (one instance per PTP Instance): receives ClockSourceTime.invoke functions from the ClockSource entity and notifications of LocalClock entity ticks (see 10.2.4.18), updates masterTime, and provides masterTime to ClockMasterSyncOffset and ClockMasterSyncSend state machines.
- b) ClockMasterSyncOffset (one instance per PTP Instance): receives syncReceiptTime from the ClockSlave entity and masterTime from the ClockMasterSyncReceive state machine, computes phase offset and frequency offset between masterTime and syncReceiptTime if the PTP Instance is not the Grandmaster PTP Instance, and provides the frequency and phase offsets to the ClockMasterSyncSend state machine.
- c) ClockMasterSyncSend (one instance per PTP Instance): receives masterTime from the ClockMasterSyncReceive state machine, receives phase and frequency offset between masterTime and syncReceiptTime from the ClockMasterSyncOffset state machine, and provides masterTime (i.e., synchronized time) and the phase and frequency offset to the SiteSync entity using a PortSyncSync structure.
- d) PortSyncSyncReceive (one instance per PTP Instance, per PTP Port): receives time-synchronization information from the MD entity of the corresponding PTP Port, computes accumulated rateRatio, computes syncReceiptTimeoutTime, and sends the information to the SiteSync entity.
- e) SiteSyncSync (one instance per PTP Instance): receives time-synchronization information, accumulated rateRatio, and syncReceiptTimeoutTime from the PortSync entity of the current slave port or from the ClockMaster entity; and sends the information to the PortSync entities of all the ports and to the ClockSlave entity.
- f) PortSyncSyncSend (one instance per PTP Instance, per PTP Port): receives time-synchronization information from the SiteSync entity, requests that the MD entity of the corresponding PTP Port send a time-synchronization event message, receives the syncEventEgressTimestamp for this event message from the MD entity, uses the most recent time-synchronization information received from the SiteSync entity and the timestamp to compute time-synchronization information that will be sent by the MD entity in a general message (e.g., for full-duplex IEEE 802.3 media) or a subsequent event message (e.g., for IEEE 802.11 media), and sends this latter information to the MD entity.
- g) ClockSlaveSync (one instance per PTP Instance): receives time-synchronization information from the SiteSync entity; computes clockSlaveTime and syncReceiptTime; sets syncReceiptLocalTime,



NOTE 1—selectedState for each port and gmPresent are set by Port State Selection state machine (see 10.3.12)

NOTE 2—currentTime is a global variable that is always equal to the current time relative to the local oscillator

NOTE 3—application interfaces to higher layers are not shown

NOTE 4—The ClockMasterSyncReceive, ClockMasterSyncSend, and ClockSyncOffset state machines are optional for PTP Instances that are not grandmaster-capable.

**Figure 10-2—Time-synchronization state machines—overview and interrelationships**

GmTimeBaseIndicator, lastGmPhaseChange, and lastGmFreqChange; sends clockSlaveTime to the ClockMaster entity; and provides information to the ClockTarget entity (via the ClockTargetPhaseDiscontinuity interface; see 9.6) to enable that entity to determine if a phase or frequency discontinuity has occurred.

## 10.2.2 Data structures communicated between state machines

### 10.2.2.1 MDSyncSend

*Change 10.2.2.1.2 as follows:*

#### 10.2.2.1.2 domainNumber (UInteger8)

This parameter is the ~~domain-number~~domainNumber of the gPTP domain in which this structure is sent.

NOTE—The ~~domain-number~~domainNumber member is not essential because the state machines that send and receive this structure are per domain, and each state machine implicitly knows the number of the domain in which it operates.

### 10.2.2.2 MDSyncReceive

*Change 10.2.2.2.2 as follows:*

#### 10.2.2.2.2 domainNumber (UInteger8)

This parameter is the ~~domain-number~~domainNumber of the gPTP domain in which this structure is sent.

NOTE—The ~~domain-number~~domainNumber member is not essential because the state machines that send and receive this structure are per domain, and each state machine implicitly knows the number of the domain in which it operates.

### 10.2.2.3 PortSyncSync

*Change 10.2.2.3.2 as follows:*

#### 10.2.2.3.2 domainNumber (UInteger8)

This parameter is the ~~domain-number~~domainNumber of the gPTP domain in which this structure is sent.

NOTE—The ~~domain-number~~domainNumber member is not essential because the state machines that send and receive this structure are per domain, and each state machine implicitly knows the number of the domain in which it operates.

## 10.2.3 Overview of global variables used by time synchronization state machines

*Change 10.2.3 as follows:*

Subclauses 10.2.4 and 10.2.5 define global variables used by time synchronization state machines whose scopes are as follows:

- Per PTP Instance (i.e., per domain)
- Per PTP Instance, per PTP Port
- Instance used by the Common Mean Link Delay Service (CMLDS) (see 11.2.17) (i.e., variable is common across all ~~LinkPorts~~Link Ports)
- Instance used by CMLDS, per ~~LinkPort~~Link Port

Table 10-1 summarizes the scope of each global variable of 10.2.4 and 10.2.5.

**Table 10-1—Summary of scope of global variables used by  
time synchronization state machines (see 10.2.4 and 10.2.5)**

Variable name	Subclause of definition	Per PTP Instance (i.e., per domain)	Per PTP Instance, per PTP Port	Instance used by CMLDS (i.e., variable is common across all <del>LinkPort</del> Link k Ports)	Instance used by CMLDS, per <del>LinkPort</del> Link Port
BEGIN	10.2.4.1	Yes	No	Yes	No
clockMasterSyncInterval	10.2.4.2	Yes	No	No	No
clockSlaveTime	10.2.4.3	Yes	No	No	No
syncReceiptTime	10.2.4.4	Yes	No	No	No
syncReceiptLocalTime	10.2.4.5	Yes	No	No	No
clockSourceFreqOffset	10.2.4.6	Yes	No	No	No
clockSourcePhaseOffset	10.2.4.7	Yes	No	No	No
clockSourceTimeBaseIndicator	10.2.4.8	Yes	No	No	No
clockSourceTimeBaseIndicatorOld	10.2.4.9	Yes	No	No	No
clockSourceLastGmPhaseChange	10.2.4.10	Yes	No	No	No
clockSourceLastGmFreqChange	10.2.4.11	Yes	No	No	No
currentTime	10.2.4.12	Yes	No	No	No
gmPresent	10.2.4.13	Yes	No	No	No
gmRateRatio	10.2.4.14	Yes	No	No	No
gmTimeBaseIndicator	10.2.4.15	Yes	No	No	No
lastGmPhaseChange	10.2.4.16	Yes	No	No	No
lastGmFreqChange	10.2.4.17	Yes	No	No	No
localClockTickInterval	10.2.4.18	Yes	No	No	No
localTime	10.2.4.19	Yes	No	No	No
selectedState	10.2.4.20	Yes	No	No	No
masterTime	10.2.4.21	Yes	No	No	No
thisClock	10.2.4.22	Yes	No	Yes	No
parentLogSyncInterval	10.2.4.23	Yes	No	No	No
instanceEnable	10.2.4.24	Yes	No	No	No
syncReceiptTimeoutTime	10.2.4.25	Yes	No	No	No
asCapable	10.2.5.1	No	Yes	No	No
asymmetryMeasurementMode	10.2.5.2	No	Yes <sup>1</sup>	No	Yes
syncReceiptTimeoutTimeInterval	10.2.5.3	No	Yes	No	No
currentLogSyncInterval	10.2.5.4	No	Yes	No	No
initialLogSyncInterval	10.2.5.5	No	Yes	No	No
syncInterval	10.2.5.6	No	Yes	No	No
neighborRateRatio	10.2.5.7	No	Yes <sup>1</sup>	No	Yes

**Table 10-1—Summary of scope of global variables used by  
time synchronization state machines (see 10.2.4 and 10.2.5) (continued)**

Variable name	Subclause of definition	Per PTP Instance (i.e., per domain)	Per PTP Instance, per PTP Port	Instance used by CMLDS (i.e., variable is common across all <del>LinkPorts</del> <u>Link k Ports</u> )	Instance used by CMLDS, per <del>LinkPort</del> <u>Link Port</u>
meanLinkDelay	10.2.5.8	No	Yes <sup>1</sup>	No	Yes
delayAsymmetry	10.2.5.9	No	Yes <sup>1</sup>	No	Yes
computeNeighborRateRatio	10.2.5.10	No	Yes <sup>1</sup>	No	Yes
computeMeanLinkDelay	10.2.5.11	No	Yes <sup>1</sup>	No	Yes
portOper <sup>2</sup>	10.2.5.12	No	Yes	No	Yes
ptpPortEnabled	10.2.5.13	No	Yes	No	No
thisPort	10.2.5.14	No	Yes	No	Yes
syncLocked	10.2.5.15	No	Yes	No	No
neighborGptpCapable	10.2.5.16	No	Yes	No	No
syncSlowdown	10.2.5.17	No	Yes	No	No
oldSyncInterval	10.2.5.18	No	Yes	No	No
gPtpCapableMessageSlowdown	10.2.5.19	No	Yes	No	No
gPtpCapableMessageInterval	10.2.5.20	No	Yes	No	No
oldGptpCapableMessageInterval	10.2.5.21	No	Yes	No	No
currentLogGptpCapableMessageInterval	10.2.5.22	No	Yes	No	No
initialLogGptpCapableMessageInterval	10.2.5.23	No	Yes	No	No
<u>syncGrandmasterIdentity</u>	<u>10.2.4.25</u>	<u>Yes</u>	<u>No</u>	<u>No</u>	<u>No</u>
<u>syncStepsRemoved</u>	<u>10.2.4.26</u>	<u>Yes</u>	<u>No</u>	<u>No</u>	<u>No</u>
<u>driftTrackingTlvSupport</u>	<u>10.2.4.27</u>	<u>Yes</u>	<u>No</u>	<u>No</u>	<u>No</u>
<u>rcvdPSSyncCSS</u>	<u>10.2.4.28</u>	<u>Yes</u>	<u>No</u>	<u>No</u>	<u>No</u>
<u>rcvdLocalClockTickCSS</u>	<u>10.2.4.29</u>	<u>Yes</u>	<u>No</u>	<u>No</u>	<u>No</u>
<u>rateRatioDrift</u>	<u>10.2.4.30</u>	<u>Yes</u>	<u>No</u>	<u>No</u>	<u>No</u>

<sup>1</sup> The instance of this variable that is per PTP Instance, per PTP Port exists only for domain 0.<sup>2</sup> There is one instance of this variable per physical port, which is accessible by all PTP Ports and ~~LinkPorts~~Link Ports associated with the physical port.

**10.2.4 Per PTP Instance global variables**

***Insert 10.2.4.25, 10.2.4.26, 10.2.4.27, 10.2.4.28, 10.2.4.29, and 10.2.4.30, and renumber subsequent subclauses as necessary:***

**10.2.4.25 syncGrandmasterIdentity:** the clockIdentity carried in the syncGrandmasterIdentity field of the Drift\_Tracking TLV carried by the most recently received Sync message (twoStep flag FALSE) or Follow\_Up message (twoStep flag TRUE). If the received Sync or Follow\_Up message does not carry a Drift\_Tracking TLV, syncGrandmasterIdentity is set to the null value 0xFFFF FFFF FFFF FFFF (see the section “Unassigned and NULL EUI values” of the IEEE Registration Authority tutorial “Guidelines for Use of Extended Unique Identifier (EUI), Organizationally Unique Identifier (OUI), and Company ID (CID)” [B30]). The data type for syncGrandmasterIdentity is ClockIdentity.

**10.2.4.26 syncStepsRemoved:** the value of the syncStepsRemoved field of the Drift\_Tracking TLV carried by the most recently received Sync message (twoStep flag FALSE) or Follow\_Up message (twoStep flag TRUE). If the received Sync or Follow\_Up message does not carry a Drift\_Tracking TLV, syncStepsRemoved is set to 0xFFFF. The data type for syncStepsRemoved is UInteger16.

**10.2.4.27 driftTrackingTlvSupport:** An indicator of whether the PTP Instance supports the Drift\_Tracking TLV and the feature is enabled. The value is TRUE if the Drift\_Tracking TLV is supported and the managed object driftTrackingTlvSupportEnabled (see 14.9) is TRUE. The value is FALSE if the Drift\_Tracking TLV is not support, or if the TLV is supported and the managed object driftTrackingTlvSupportEnabled is FALSE. The data type for driftTrackingTlvSupport is Boolean.

NOTE—The Drift\_Tracking TLV is transported only on full-duplex, point-to-point links as specified in Clause 11. Even if driftTrackingTlvSupport is TRUE, the Drift\_Tracking TLV is not transported on links other than full-duplex, point-to-point links, and is not received by the PTP Ports at the other end of these links.

**10.2.4.28 rcvdPSSyncCSS:** A Boolean variable that is set to TRUE when a PortSyncSync structure is received from the SiteSyncSync state machine of the SiteSync entity by the ClockSlaveSync state machine. This variable is reset by the ClockSlaveSync state machine.

**10.2.4.29 rcvdLocalClockTickCSS:** A Boolean variable that is set to TRUE when the LocalClock entity updates its time. This variable is reset by the ClockSlaveSync state machine.

**10.2.4.30 rateRatioDrift:** the value of the rateRatioDrift field of the Drift\_Tracking TLV carried by the most recently received Sync message (twoStep flag FALSE) or Follow\_Up message (twoStep flag TRUE). If the received Sync or Follow\_Up message does not carry a Drift\_Tracking TLV, rateRatioDrift is set to 0xFFFF. The data type for rateRatioDrift is Integer32.

<<Editor’s note: Since rateRatioDrift is an internal global variable, it is possible to make it a Float64, while keeping the rateRatioDrift field as the scaled Integer32. This was done for rateRatio (e.g., the internal rateRatio variables of the various state machines have data type Float64, but the rateRatio field of the Follow\_Up TLV and the managed object defaultDS.cumulativeRateRatio are the scaled Integer32 variables. Should the internal rateRatioDrift variable be a Float64 that stores the actual rateRatio?>>

**10.2.5 Per port global variables**

***Change 10.2.5.9 as follows:***

**10.2.5.9 delayAsymmetry:** The asymmetry in the propagation delay on the link attached to this port. If propagation delay asymmetry is not modeled, then delayAsymmetry is zero. The data type for delayAsymmetry is ScaledNs. There is one instance of this variable for CMLDS (see 11.2.17), and there is also one instance of this variable for each domain ~~that uses the instance-specific peer-to-peer delay mechanism~~. The instance of this variable for CMLDS is relative to the local clock. The instance of this variable for ~~each~~ domain ~~that uses the instance-specific peer-to-peer delay mechanism~~ is relative to the grandmaster time base for that domain. The instance of delayAsymmetry for CMLDS is used where needed in all computations done by the CMLDS peer-to-peer delay mechanism. The instance of delayAsymmetry for a domain is used where needed in all computations done for that domain.

**<<Editor's note:** The resolution of comments#25 and #8 indicated that delayAsymmetry is domain specific, and therefore the instance of delayAsymmetry for CMLDS is not needed and should be deleted. However, on looking into this further, it was realized that while the computations done when sending and receiving Sync messages indeed do use delayAsymmetry for the respective domain irrespective of whether the CMLDS or transport-specific peer-to-peer delay mechanism is used, delayAsymmetry is needed in the CMLDS peer-to-peer delay computations because it is added and subtracted (see 11.2.17.2 and 11.2.19.3.4). Therefore, the instance of delayAsymmetry for CMLDS is retained. It is indicated above and in other relevant places which delayAsymmetry is used in the respective computations. Note also that the delayAsymmetry for a domain is referred to as "domain specific delayAsymmetry" rather than "transport-specific peer-to-peer delay mechanism delayAsymmetry" because this delayAsymmetry is used where needed in all domain-specific computations, i.e., not only in transport-specific peer delay computations.

Note also that the CMLDS delayAsymmetry is needed only for compatibility with IEEE Std 1588-2019 . If 1588 had not chosen to add and subtract delayAsymmetry (this would have been possible because the effect of adding and subtracting it cancels), CMLDS delayAsymmetry would not be needed.>>

**<<Editor's note:** References to delayAsymmetry in this standard do not explicitly indicate "domain-specific delayAsymmetry" or "CMLDS delayAsymmetry" because in all cases it is clear whether the respective computations are part of CMLDS or are domain-specific. For example, CMLDS computations are only done as part of the CMLDS peer-to-peer delay mechanism, and it is always indicated if this mechanism is used; all other computations are domain-specific. Comments are requested on whether this is sufficiently clear.>>

*Change 10.2.5.13 as follows:*

**10.2.5.13 ptpPortEnabled:** A Boolean that is administratively set to TRUE if time-synchronization is to be enabled on this PTP Port.

NOTE 1—It is expected that the value of ptpPortEnabled ~~will be~~ is set via the management interface (see 14.8.4). A physical port ~~PTP Port~~ can be enabled for data transport but not for synchronization transport.

NOTE 2—The variable ptpPortEnabled was named pttPortEnabled in the 2011 edition of this standard. Only the name of this variable has changed; the definition and function of this variable are the same as in the 2011 edition. The name change is reflected in many state machines.

## 10.2.8 PortSyncSyncReceive state machine

*Add the following definition to 10.2.8.1:*

### 10.2.8.1 State machine variables

**10.2.8.1.5 TEMP:** A temporary variable used to reduce clutter in the state diagram (see Figure 10-4). The data type for TEMP is Integer16.

**10.2.9 ClockMasterSyncSend state machine****10.2.9.2 State machine functions**

**10.2.9.2.1 setPSSyncCMSS (gmRateRatio):** Creates a PortSyncSync structure to be transmitted, and returns a pointer to this structure. The members are set as follows:

*Add the following NOTE 1 to 10.2.9.2.1, just after item c)2), and renumber subsequent NOTES as needed:*

NOTE 1—The definition of `currentTime` (see 10.2.4.12) and `localTime` (see 10.2.4.19), and the statement `localTime = currentTime` in the `RECEIVE_SOURCE_TIME` state of the `ClockMasterSyncReceive` state machine, imply that `currentTime` and `localTime` are always equal and the quantity  $\text{gmRateRatio} \times (\text{currentTime} - \text{localTime})$  is always zero. Both `localTime` and `masterTime` are updated by the `ClockMasterSyncReceive` state machine. The updates occur both when `sourceTime` is received from the `ClockSource` (indicated by the variable `rcvdClockSourceReq` being `TRUE`) and when the `LocalClock` entity itself updates by one tick (indicated by the variable `rcvdLocalClockTickCMSR` being `TRUE`). `masterTime` is either updated to the `sourceTime` (if the update occurs due to receipt of `sourceTime` from the `ClockSource`) or incremented by one `LocalClock` tick interval multiplied by `gmRateRatio`. `localTime` is updated by setting it equal to `currentTime`. The result is that `localTime` is always equal to the value of `currentTime` when `masterTime` was most recently updated. `localTime`, `currentTime`, and `masterTime` are per PTP Instance global variables. When they are updated, their values are known to all state machines. In setting (i) `preciseOriginTimestamp` to `masterTime` with fractional `ns` truncated, and `followUpCorrectionField` to (ii) the fractional `ns` part of `masterTime`, plus (iii)  $\text{gmRateRatio} \times (\text{currentTime} - \text{localTime})$ , the contributions (i) plus (ii) give the estimate of the `ClockSource` time when `masterTime` was most recently updated, because this latest computed `masterTime` includes any updates due to `LocalClock` ticks. The quantity (iii) is zero in the model here, because `currentTime` and `localTime` are equal in this model. However, in practice there could be situations where `currentTime` and `localTime` are not equal. For example, if the `ClockMasterSyncReceive` state machine is implemented as a different module than the `ClockMasterSyncSend` state machine, with `masterTime` and `localTime` being sent from the `ClockMasterSyncReceive` state machine to the `ClockMasterSyncSend` state machine, then `currentTime` and `localTime` will not be equal. In this case, `currentTime` in c)2) is the value of the `LocalClock` time when the `ClockMasterSyncSend` state machine receives `masterTime` and `localTime` from the `ClockMasterSyncReceive` state machine.

<<Editor's note: The following NOTES were proposed by comments #6 and #7 against D1.1, to replace NOTE 1 above. Comments are requested on whether to use the current text, the proposed text, or some modification of either or both texts:

NOTE 1—Both `localTime` (10.2.4.19) and `masterTime` (10.2.4.21) are updated by the `ClockMasterSyncReceive` state machine (10.2.11). The updates occur both when `sourceTime` (9.2.2.2) is received from the `ClockSource` (9.2) - indicated by the variable `rcvdClockSourceReq` (10.2.11.1.1) being `TRUE` - and when the `LocalClock` entity (10.1.2.1) itself updates by one tick - indicated by the variable `rcvdLocalClockTickCMSR` (10.2.11.1.3) being `TRUE`. `masterTime` is updated by `updateMasterTime()` (see 10.2.11.2.2). `localTime` is updated by setting it equal to `currentTime` (Figure 10-7). The result is that `localTime` is equal to the value of `currentTime` when `masterTime` was most recently updated. `localTime`, `currentTime`, and `masterTime` are per PTP Instance global variables. When they are updated, their values are known to all state machines.

NOTE 2—`currentTime` (10.2.4.19) and `localTime` (10.2.4.19) may not always be equal, despite the statement `localTime = currentTime` in the `RECEIVE_SOURCE_TIME` state of the `ClockMasterSyncReceive` state machine (10.2.11). For example, if the `ClockMasterSyncReceive` state machine is implemented as a different module than the `ClockMasterSyncSend` state machine (10.2.9), with `masterTime` (10.2.4.21) and `localTime` being sent from the `ClockMasterSyncReceive` state machine to the `ClockMasterSyncSend` state machine, then `currentTime` and `localTime` will not be equal. In this case, `currentTime` in c)2) is the value of the `LocalClock` entity (10.1.2.1) when the `ClockMasterSyncSend` state machine receives `masterTime` and `localTime` from the `ClockMasterSyncReceive` state machine.>>

*Change item m) of 10.2.9.2.1 as follows:*

- m) `domainNumber` is set equal to the ~~domain number~~ `domainNumber` of this gPTP domain.

*Change the first paragraph in 10.2.9.3 as follows:*



**10.2.9.3 State diagram**

The ClockMasterSyncSend state machine shall implement the function specified by the state diagram in Figure 10-5, the local variables specified in 10.2.9.1, the functions specified in 10.2.9.2, the structure specified in 10.2.2.3, and the relevant global variables and functions specified in 10.2.4 through 10.2.6. The state machine receives masterTime and clockSourceTimeBaseIndicator from the ClockMasterSyncReceive state machine, and phase and frequency offset between masterTime and syncReceiptTime from the ClockMasterSyncOffset state machine. It provides masterTime (i.e., synchronized time) and the phase and frequency offset to the SiteSync entity via a PortSyncSync structure.

*Change the heading of 10.2.10 as follows:*

**10.2.10 ClockMasterSyncOffset state machine**

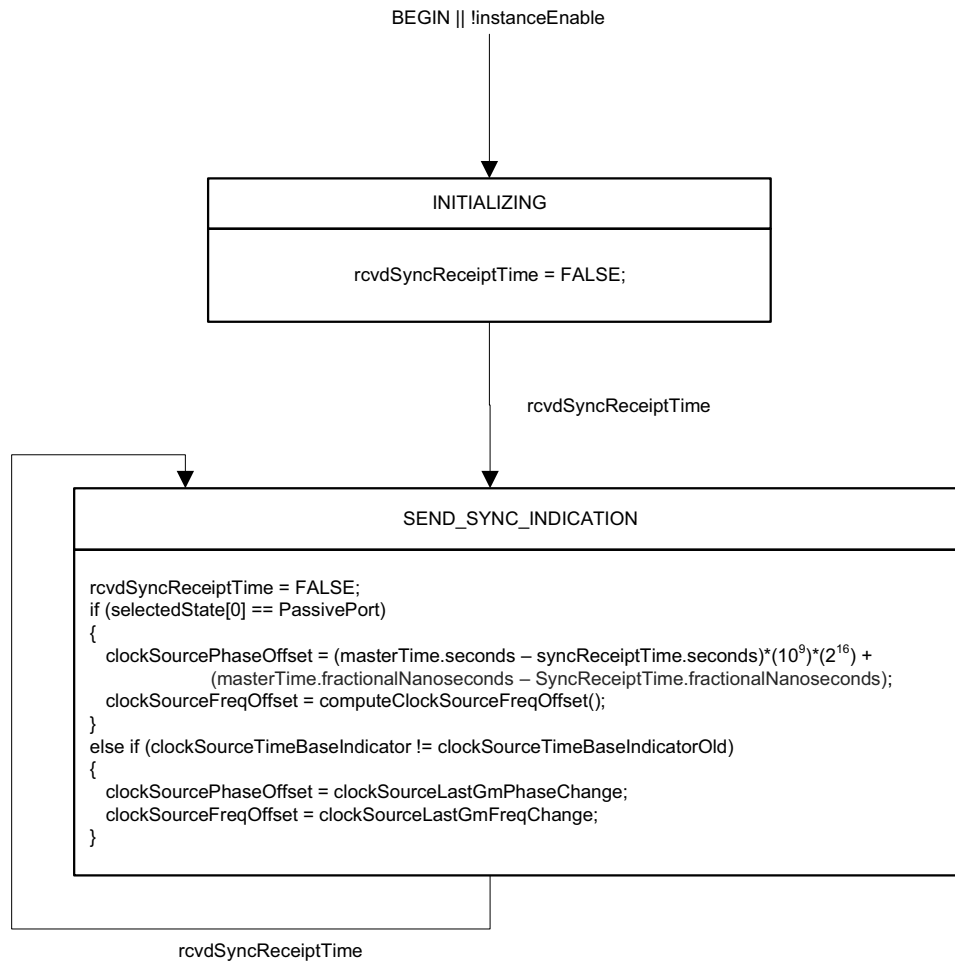
*Change 10.2.10.3 as follows:*

**10.2.10.3 State diagram**

The ClockMasterSyncOffset state machine shall implement the function specified by the state diagram in Figure , the local variable specified in 10.2.10.1, the function specified in 10.2.10.2, and the relevant global variables and functions specified in 10.2.4 through 10.2.6. The state machine receives syncReceiptTime from the ClockSlaveSync state machine and masterTime from the ClockMasterSyncReceive state machine. It computes clockSourcePhaseOffset and clockSourceFrequency offset if this PTP Instance is not currently the Grandmaster PTP Instance, i.e., if selectedState[0] is equal to PassivePort.

The ClockMasterSyncOffset state machine is optional for PTP Instances that are not grandmaster-capable (see 8.6.2.1, 10.1.3, and 10.2.1). This state machine may be present in a PTP Instance that is not grandmaster-capable; however, any information supplied by it, via the ClockMasterSyncSend state machine, to the SiteSyncSync state machine is not used by the SiteSyncSync state machine if the PTP Instance is not grandmaster-capable.

*Replace Figure 10-6 with the following:*



**NOTE**—This figure differs from the 2020 edition, with Corrigendum 1 applied, of this standard in that the statement `clockSourcePhaseOffset = masterTime.seconds – syncReceiptTime;` in the `SEND_SYNC_INDICATION` state is replaced by

$$\text{clockSourcePhaseOffset} = (\text{masterTime.seconds} - \text{syncReceiptTime.seconds}) \times (10^9) \times (2^{16}) + (\text{masterTime.fractionalNanoseconds} - \text{SyncReceiptTime.fractionalNanoseconds});$$

**Figure 10-6—ClockMasterSyncOffset state machine**

## 10.2.12 PortSyncSyncSend state machine

### 10.2.12.2 State machine functions

*Change item j) of 10.2.12.2.1 as follows:*

- j) domainNumber is set equal to the ~~domain-number~~domainNumber of this gPTP domain (see 8.1).

## 10.2.13 ClockSlaveSync state machine

*Change 10.2.13.1 as follows:*

### 10.2.13.1 State machine variables

The following variables ~~are~~is used in the state diagram in Figure 10-9 (in 10.2.13.3):

~~10.2.13.1.1 revdPSSyncCSS: A Boolean variable that notifies the current state machine when a PortSyncSync structure is received from the SiteSyncSync state machine of the SiteSync entity. This variable is reset by this state machine.~~

~~10.2.13.1.2 revdLocalClockTickCSS: A Boolean variable that notifies the current state machine when the LocalClock entity updates its time. This variable is reset by this state machine.~~

10.2.13.1.3 revdPSSyncPtrCSS: A pointer to the received PortSyncSync structure.

### 10.2.13.2 State diagram

*Change the first paragraph of 10.2.13.2 as follows:*

The ClockSlaveSync state machine shall implement the function specified by the state diagram in Figure 10-9, the local variables specified in 10.2.13.1, the functions specified in 10.2.13.2, and the relevant global variables and functions specified in 10.2.4 through 10.2.6. The state machine receives a PortSyncSync structure from the SiteSyncSync state machine. It computes syncReceiptTime and clockSlaveTime, and sets syncReceiptLocalTime (i.e., the time relative to the LocalClock entity corresponding to syncReceiptTime), GmTimeBaseIndicator, lastGmPhaseChange, and lastGmFreqChange. It provides clockSlaveTime to the ClockMasterSyncOffset state machine, and provides information to the ClockTarget entity (via the ClockTargetPhaseDiscontinuity interface; see 9.6) to enable that entity to determine if a phase or frequency discontinuity has occurred.

## 10.3 Best master clock selection, external port configuration, and announce interval setting state machines

### 10.3.1 Best master clock selection and external port configuration overview

*Change 10.3.1.3 as follows:*

### 10.3.1.3 External port configuration overview

In external port configuration (i.e., method b) of 10.3.1.1), an external entity determines the synchronization spanning tree and sets the PTP Port states accordingly. The method used by the external entity to determine the synchronization spanning tree is outside the scope of this standard. However, as with the BMCA, Announce messages are used to transport information on the time-synchronization spanning tree and Grandmaster PTP Instance time properties information from one PTP Instance to the next in the tree. The external entity sets the state of a PTP Port by setting the value of externalPortConfigurationPortDS.desiredState to the desired state.

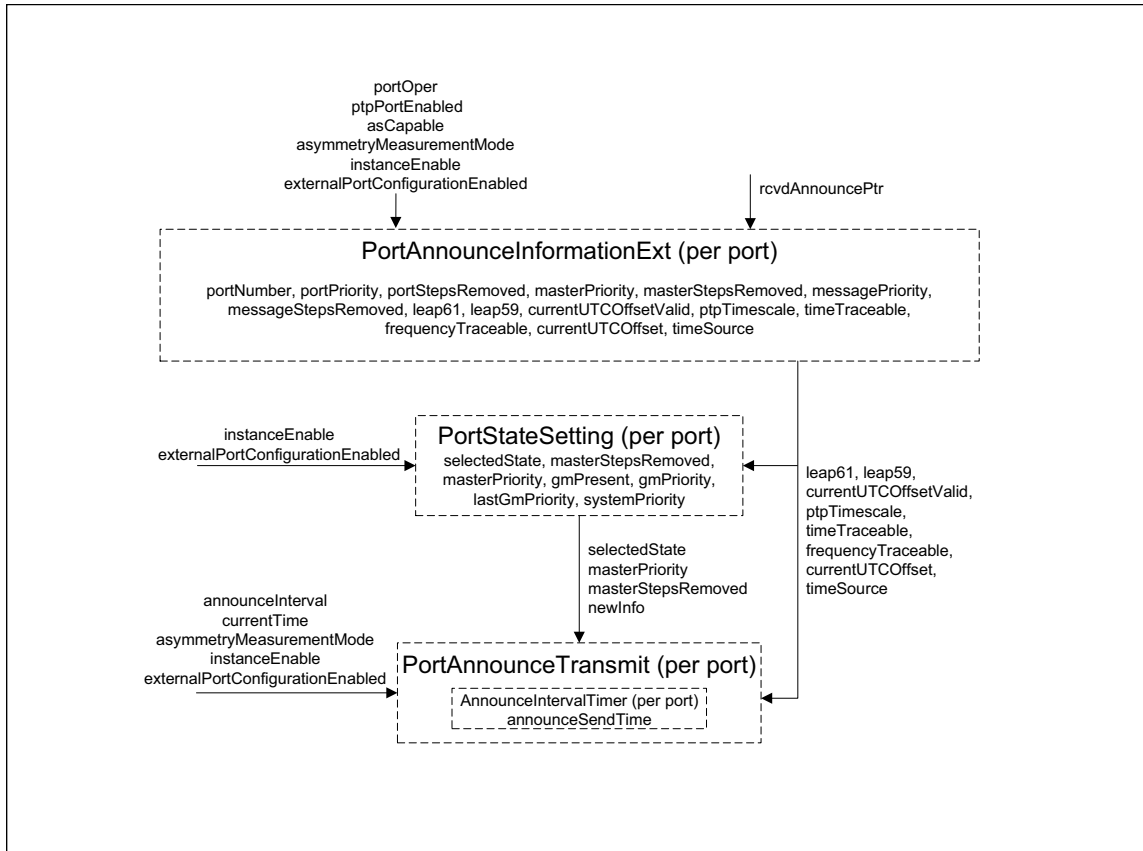
In the case of external port configuration, the time-synchronization spanning tree and the desired PTP Port states are controlled by an external entity, but the Grandmaster PTP Instance might change. If the slave port of a PTP Instance that is gmCapable (priority1 < 255, see 8.6.2.1) is no longer asCapable, the state of this PTP Port changes from SlavePort to DisabledPort (see 10.3.6.2) and the PTP Instance becomes Grandmaster for the time synchronization (sub-)tree where it is the root (10.3.15.2.2 d).3)). The original time-synchronization spanning tree can be split into disjunct subtrees with different, mutually unsynchronized, Grandmaster PTP Instances. If the port becomes asCapable again, its PTP Port state is again set to the desired state. If the slave port of a PTP Instance that is not gmCapable is no longer asCapable, the state of this PTP Port changes from SlavePort to DisabledPort (see 10.3.6.2). However, in this case gmPresent is set to FALSE (see 10.3.15.2.2 d) 3)), and the PTP Instance does not send Sync messages on any of its ports.

In external port configuration, there is no supervision of announce receipt timeout (see 10.7.3.2).

## 10.3.7 Overview of best master clock selection, external port configuration, and announce interval setting state machines

### 10.3.7.2 External port configuration state machines overview

*Replace Figure 10-12 with the following:*



This figure differs from the 2020 edition, with Corrigendum 1 applied, of this standard in that the variable newInfo is removed from the link between the PortAnnounceInformationExt and PortAnnounceTransmit blocks and added to the link between the PortStateSetting and PortAnnounceTransmit blocks.

**Figure 10-12—External port configuration state machines—overview and interrelationships**

### 10.3.8 Overview of global variables used by best master clock selection, external port configuration, and announce interval setting state machines

#### *Change 10.3.8 as follows:*

Subclauses and 10.3.10 define global variables used by best master clock selection, external port configuration, and announce interval setting state machines whose scopes are as follows:

- Per PTP Instance (i.e., per domain)
- Per PTP Instance, per PTP Port
- Instance used by CMLDS (see 11.2.17) (i.e., variable is common across all ~~Link Ports~~ Link Ports)
- Instance used by CMLDS, per ~~Link Port~~ Link Port

Table 10-3 summarizes the scope of each global variable of and 10.3.10.

#### *Change 10.3.9 as follows:*

**Table 10-3—Summary of scope of global variables used by  
best master clock selection, external port configuration, and announce interval setting  
state machines (see 10.3.9 and 10.3.10)**

Variable name	Subclause of definition	Per PTP Instance (i.e., per domain)	Per PTP Instance, per PTP Port	Instance used by CMLDS (i.e., variable is common across all Link Ports Link Ports)	Instance used by CMLDS, per Link Port Link Port
reselect	10.3.9.1	Yes	No	No	No
selected	10.3.9.2	Yes	No	No	No
masterStepsRemoved	10.3.9.3	Yes	No	No	No
leap61	10.3.9.4	Yes	No	No	No
leap59	10.3.9.5	Yes	No	No	No
currentUtcOffsetValid	10.3.9.6	Yes	No	No	No
ptpTimescale	10.3.9.7	Yes	No	No	No
timeTraceable	10.3.9.8	Yes	No	No	No
frequencyTraceable	10.3.9.9	Yes	No	No	No
currentUtcOffset	10.3.9.10	Yes	No	No	No
timeSource	10.3.9.11	Yes	No	No	No
sysLeap61	10.3.9.12	Yes	No	No	No
sysLeap59	10.3.9.13	Yes	No	No	No
sysCurrentUtcOffsetValid	10.3.9.14	Yes	No	No	No
sysPtpTimescale	10.3.9.15	Yes	No	No	No
sysTimeTraceable	10.3.9.16	Yes	No	No	No
sysFrequencyTraceable	10.3.9.17	Yes	No	No	No
sysCurrentUtcOffset	10.3.9.18	Yes	No	No	No
sysTimeSource	10.3.9.19	Yes	No	No	No
systemPriority	10.3.9.20	Yes	No	No	No
gmPriority	10.3.9.21	Yes	No	No	No
lastGmPriority	10.3.9.22	Yes	No	No	No
pathTrace	10.3.9.23	Yes	No	No	No
externalPortConfigurationEnabled	10.3.9.24	Yes	No	No	No
lastAnnouncePort	10.3.9.25	Yes	No	No	No
announceReceiptTimeoutTimeInterval	10.3.10.1	No	Yes	No	No
announceSlowdown	10.3.10.2	No	Yes	No	No
oldAnnounceInterval	10.3.10.3	No	Yes	No	No
infoIs	10.3.10.4	No	Yes	No	No
masterPriority	10.3.10.5	No	Yes	No	No

**Table 10-3—Summary of scope of global variables used by  
best master clock selection, external port configuration, and announce interval setting  
state machines (see 10.3.10) (continued)**

Variable name	Subclause of definition	Per PTP Instance (i.e., per domain)	Per PTP Instance, per PTP Port	Instance used by CMLDS (i.e., variable is common across all <del>Link Ports</del> Link Ports)	Instance used by CMLDS, per <del>Link Port</del> Link Port
currentLogAnnounceInterval	10.3.10.6	No	Yes	No	No
initialLogAnnounceInterval	10.3.10.7	No	Yes	No	No
announceInterval	10.3.10.8	No	Yes	No	No
messageStepsRemoved	10.3.10.9	No	Yes	No	No
newInfo	10.3.10.10	No	Yes	No	No
portPriority	10.3.10.11	No	Yes	No	No
portStepsRemoved	10.3.10.12	No	Yes	No	No
rcvdAnnouncePtr	10.3.10.13	No	Yes	No	No
rcvdMsg	10.3.10.14	No	Yes	No	No
updtInfo	10.3.10.15	No	Yes	No	No
annLeap61	10.3.10.16	No	Yes	No	No
annLeap59	10.3.10.17	No	Yes	No	No
annCurrentUtcOffsetValid	10.3.10.18	No	Yes	No	No
annPtpTimescale	10.3.10.19	No	Yes	No	No
annTimeTraceable	10.3.10.20	No	Yes	No	No
annFrequencyTraceable	10.3.10.21	No	Yes	No	No
annCurrentUtcOffset	10.3.10.22	No	Yes	No	No
annTimeSource	10.3.10.23	No	Yes	No	No
receivedPathTrace	10.3.10.24	No	Yes	No	No

### 10.3.9 Per PTP Instance global variables

**10.3.9.1 reselect:** A Boolean array of length numberPorts+1 (see 8.6.2.8). Setting reselect[j], where  $0 \leq j \leq \text{numberPorts}$ , to TRUE causes the STATE\_SELECTION block of the PortStateSelection state machine (see 10.3.13) to be re-entered, which in turn causes the PTP Port state of each PTP Port of the PTP Instance to be updated (via the function updtStatesTree(); see 10.3.13.2.4). This variable is used only by the BMCA, i.e., not by the ~~explicit port state~~ external port configuration option.

**10.3.9.2 selected:** A Boolean array of length numberPorts+1 (see 8.6.2.8). selected[j], where  $0 \leq j \leq \text{numberPorts}$ , is set to TRUE immediately after the PTP Port states of all the ports are updated. This value indicates to the PortAnnounceInformation state machine (see 10.3.12) that it can update the portPriorityVector and other variables for each PTP Port. This variable is used by both the BMCA and the ~~explicit port state~~ external port configuration option; however, its value does not impact the ~~explicit port state~~ external port configuration option (see the NOTE in 10.3.16.3).

NOTE—Array elements 0 of the `reselect` and `selected` arrays are not used, except that the function `clearReselectTree()` sets `reselect[0]` to `FALSE` when it sets the entire array to zero and the function `setSelectedTree()` sets `selected[0]` to `TRUE` when it sets the entire array to `TRUE`. This action is taken only for convenience, so that array element `j` can correspond to PTP Port `j`. Note also that, in contrast, `selectedState[0]` is ~~not~~ used (see 10.2.4.20)

**10.3.9.3 masterStepsRemoved:** The value of `stepsRemoved` for the PTP Instance, after the PTP Port states of all the ports have been updated (see 10.3.13.2.4 for details on the computation of `masterStepsRemoved`). The data type for `masterStepsRemoved` is `UInteger16`. This variable is used by both the BMCA and the ~~explicit-port-state~~external port configuration option.

**10.3.9.4 leap61:** A Boolean variable whose value is `TRUE` if the last minute of the current UTC day, relative to the current Grandmaster Clock, contains 61 s and `FALSE` if the last minute of the current UTC day does not contain 61 s. This variable is used by both the BMCA and the ~~explicit-port-state~~external port configuration option.

**10.3.9.5 leap59:** A Boolean variable whose value is `TRUE` if the last minute of the current UTC day, relative to the current Grandmaster Clock, contains 59 s and `FALSE` if the last minute of the current UTC day does not contain 59 s. This variable is used by both the BMCA and the ~~explicit-port-state~~external port configuration option.

**10.3.9.6 currentUtcOffsetValid:** A Boolean variable whose value is `TRUE` if `currentUtcOffset` (see 10.3.9.10), relative to the current Grandmaster Clock, is known to be correct and `FALSE` if `currentUtcOffset` is not known to be correct. This variable is used by both the BMCA and the ~~explicit-port-state~~external port configuration option.

**10.3.9.7 ptpTimescale:** A Boolean variable whose value is `TRUE` if the timescale of the current Grandmaster Clock is PTP (see 8.2.1) and `FALSE` if the timescale is ARB. This variable is used by both the BMCA and the ~~explicit-port-state~~external port configuration option.

**10.3.9.8 timeTraceable:** A Boolean variable whose value is `TRUE` if both `clockSlaveTime` [i.e., the synchronized time maintained at the slave (see 10.2.4.3)] and `currentUtcOffset` (see 10.3.9.10), relative to the current Grandmaster Clock, are traceable to a primary reference and `FALSE` if one or both are not traceable to a primary reference. This variable is used by both the BMCA and the ~~explicit-port-state~~external port configuration option.

**10.3.9.9 frequencyTraceable:** A Boolean variable whose value is `TRUE` if the frequency that determines `clockSlaveTime`, i.e., the frequency of the `LocalClockEntity` multiplied by the most recently computed `rateRatio` by the `PortSyncSyncReceive` state machine (see 10.2.8.1.4), is traceable to a primary reference and `FALSE` if this frequency is not traceable to a primary reference. This variable is used by both the BMCA and the ~~explicit-port-state~~external port configuration option.

**10.3.9.10 currentUtcOffset:** The difference between TAI time and UTC time, i.e., TAI time minus UTC time, in seconds, and relative to the current Grandmaster Clock, when known. Otherwise, the value has no meaning (see 10.3.9.6). The data type for `currentUtcOffset` is `Integer16`. This variable is used by both the BMCA and the ~~explicit-port-state~~external port configuration option.

NOTE—For example, 2006-01-01 00:00:00 UTC and 2006-01-01 00:00:33 TAI represent the same instant of time. At this time, `currentUtcOffset` was equal to 33 s.<sup>1</sup>

**10.3.9.11 timeSource:** The value of the `timeSource` attribute of the current Grandmaster PTP Instance. The data type for `timeSource` is `TimeSource` (see 8.6.2.7). This variable is used by both the BMCA and the ~~explicit-port-state~~external port configuration option.

<sup>1</sup>Note also that a leap second was not added at the end of the last UTC minute of 2005-12-31.



**10.3.9.12 sysLeap61:** A Boolean variable whose value is TRUE if the last minute of the current UTC day, relative to the ClockMaster entity of this PTP Instance, contains 61 s and FALSE if the last minute of the current UTC day does not contain 61 s. This variable is used by both the BMCA and the ~~explicit-port~~ stateexternal port configuration option.

**10.3.9.13 sysLeap59:** A Boolean variable whose value is TRUE if the last minute of the current UTC day, relative to the ClockMaster entity of this PTP Instance, contains 59 s and FALSE if the last minute of the current UTC day does not contain 59 s. This variable is used by both the BMCA and the ~~explicit-port~~ stateexternal port configuration option.

**10.3.9.14 sysCurrentUtcOffsetValid:** A Boolean variable whose value is TRUE if currentUtcOffset (see 10.3.9.10), relative to the ClockMaster entity of this PTP Instance, is known to be correct and FALSE if currentUtcOffset is not known to be correct. This variable is used by both the BMCA and the ~~explicit-port~~ stateexternal port configuration option.

**10.3.9.15 sysPtpTimescale:** A Boolean variable whose value is TRUE if the timescale of the ClockMaster entity of this PTP Instance is PTP (see 8.2.1) and FALSE if the timescale of the ClockMaster entity of this PTP Instance is ARB. This variable is used by both the BMCA and the ~~explicit-port~~ stateexternal port configuration option.

**10.3.9.16 sysTimeTraceable:** A Boolean variable whose value is TRUE if both masterTime [i.e., the time maintained by the ClockMaster entity of this PTP Instance (see 10.2.4.21)] and currentUtcOffset (see 10.3.9.10), relative to the ClockMaster entity of this PTP Instance, are traceable to a primary reference and FALSE if one or both are not traceable to a primary reference. This variable is used by both the BMCA and the ~~explicit-port~~ stateexternal port configuration option.

**10.3.9.17 sysFrequencyTraceable:** A Boolean variable whose value is TRUE if the frequency that determines masterTime of the ClockMaster entity of this PTP Instance, i.e., the frequency of the LocalClockEntity multiplied by the most recently computed gmRateRatio by the ClockMasterSyncReceive state machine (see 10.2.4.14 and 10.2.11), is traceable to a primary reference and FALSE if this frequency is not traceable to a primary reference. This variable is used by both the BMCA and the ~~explicit-port~~ stateexternal port configuration option.

**10.3.9.18 sysCurrentUtcOffset:** The difference between TAI time and UTC time, i.e., TAI time minus UTC time, in seconds, and relative to the ClockMaster entity of this PTP Instance, when known. Otherwise, the value has no meaning (see 10.3.9.14). The data type for sysCurrentUtcOffset is Integer16. This variable is used by both the BMCA and the ~~explicit-port~~ stateexternal port configuration option.

NOTE—See the NOTE in 10.3.9.10 for more detail on the sign convention.

**10.3.9.19 sysTimeSource:** The value of the timeSource attribute of the ClockMaster entity of this PTP Instance (see 8.6.2.7). The data type for sysTimeSource is TimeSource.

**10.3.9.20 systemPriority:** The systemPriority vector for this PTP Instance. The data type for systemPriority is UInteger224 (see 10.3.5).

**10.3.9.21 gmPriority:** The current gmPriorityVector for the PTP Instance. The data type for gmPriority is UInteger224 (see 10.3.5).

**10.3.9.22 lastGmPriority:** The previous gmPriorityVector for the PTP Instance, prior to the most recent invocation of the PortStateSelection state machine. The data type for lastGmPriority is UInteger224 (see 10.3.4). lastGmPriority is used only by the BMCA, i.e., not by the ~~explicit-port~~ stateexternal port configuration option.

**10.3.9.23 pathTrace:** An array that contains the clockIdentities of the successive PTP Instances that receive, process, and send Announce messages. The data type for pathTrace is ClockIdentity[N], where N is the number of PTP Instances, including the Grandmaster PTP Instance, that the Announce information has traversed. This variable is used by both the BMCA and the ~~explicit-port-state~~external port configuration option.

NOTE 1—N is equal to stepsRemoved+1 (see 10.6.3.2.6). The size of the pathTrace array can change after each reception of an Announce message, up to the maximum size for the respective medium. For example, the maximum value of N for a full-duplex IEEE 802.3 medium is 179. This is obtained from the fact that the number of PTP octets in an Announce message is  $68 + 8N$ , where N is the number of entries in the pathTrace array (see 10.6.3.1 and Table 10-11), and the maximum payload size for full-duplex IEEE 802.3 media is 1500 octets. Setting  $68 + 8N = 1500$ , and solving for N gives  $N = 179$ .

NOTE 2—The current behavior for the path trace feature is documented in 10.3.11.2.1 and 10.3.16.2.1 and is as follows:

- Item c) of 10.3.11.2.1, the description of the qualifyAnnounce() function of the PortAnnounceReceive state machine, indicates that if a path trace TLV is present and one of the elements of the pathSequence array field is equal to the clockIdentity of the clock where the TLV is being processed, the Announce message is not qualified.
- Item d) of 10.3.11.2.1 (qualifyAnnounce() function) indicates that if the Announce message is qualified and a path trace TLV is present, the pathSequence array of the TLV is copied to the pathTrace array (described in this subclause) and the clockIdentity of the PTP Instance that processes the Announce message is appended to the array. However, if a path trace TLV is not present, the path trace array is empty.
- Item f) of 10.3.16.2.1, the description of the txAnnounce() function of the PortAnnounceTransmit state machine, indicates that a path trace TLV is constructed and appended to an Announce message just before the Announce message is transmitted only if the pathTrace array is not empty and appending the TLV does not cause the media-dependent layer frame to exceed any respective maximum size. If appending the TLV does cause a respective maximum frame size to be exceeded or if the pathTrace array is empty, the TLV is not appended.
- As a result of the behaviors of the qualifyAnnounce() and txAnnounce() functions described in this note, the path trace feature is not used, i.e., a path trace TLV is not appended to an Announce message and the pathTrace array is empty, once appending a clockIdentity to the TLV would cause the frame carrying the Announce message to exceed its maximum size.

NOTE 3—Once the value of stepsRemoved of an Announce message reaches 255, the Announce message is not qualified [see item b) of 10.3.11.2.1].

**10.3.9.24 externalPortConfigurationEnabled:** A variable whose value indicates whether PTP Port states are externally configured or determined by the BMCA. The data type shall be Boolean. The value TRUE indicates that the PTP Port states are externally configured; the value FALSE indicates that the PTP Port states are determined by the BMCA. This variable is used by both the BMCA and the ~~explicit-port-state~~external port configuration option.

**10.3.9.25 lastAnnouncePort:** The PTP Port number of the PTP Port on which the most recent Announce message was received. This variable is used by the PortAnnounceInformationExt and PortStateSettingExt state machines for the ~~explicit-port-state~~external port configuration option. This variable is not used by the BMCA. The data type for this variable is UInteger16.

*Change 10.3.10 as follows:*

## 10.3.10 Per-port global variables

**10.3.10.1 announceReceiptTimeoutTimeInterval:** The time interval after which announce receipt timeout occurs if an Announce message has not been received during the interval. The value of announceReceiptTimeoutTimeInterval is equal to announceReceiptTimeout (see 10.7.3.2) multiplied by the announceInterval (see 10.3.10.8) for the PTP Port at the other end of the link to which this PTP Port is attached. The value of announceInterval for the PTP Port at the other end of the link is computed from logMessageInterval of the received Announce message (see 10.6.2.2.14). The data type for

announceReceiptTimeoutTimeInterval is UScaledNs. This variable is used only by the BMCA, i.e., not by the ~~explicit-port state~~external port configuration option.

**10.3.10.2 announceSlowdown:** A Boolean that is set to TRUE if the AnnounceIntervalSetting state machine (see Figure 10-19 in item 10.3.17.3) receives a TLV that requests a larger Announce message transmission interval (see 10.7.2.2) and FALSE otherwise. When announceSlowdown is set to TRUE, the PortAnnounceTransmit state machine (see Figure ) continues to send Announce messages at the old (i.e., faster) rate until a number of Announce messages equal to announceReceiptTimeout (see 10.7.3.2) have been sent, but with the logMessageInterval field of the PTP common header set equal to the new announce interval (i.e., corresponding to the slower rate). After announceReceiptTimeout Announce messages have been sent, subsequent Announce messages are sent at the new (i.e., slower) rate and with the logMessageInterval field of the PTP common header set to the new announce interval. This variable is used by both the BMCA and the ~~explicit-port state~~external port configuration option. When announceSlowdown is set to FALSE, the PortAnnounceTransmit state machine immediately sends Announce messages at the new (i.e., slower) rate.

NOTE—If a receiver of Announce messages requests a slower rate, the receiver ~~will~~ continues to use the upstream announceInterval value, which it obtains from the logMessageInterval field of received Announce messages, until it receives an Announce message where that value has changed. If, immediately after requesting a slower Announce message rate, up to announceReceiptTimeout minus one consecutive Announce messages sent to the receiver are lost, announce receipt timeout could occur if the sender had changed to the slower rate immediately. Delaying the slowing down of the sending rate of Announce messages for announceReceiptTimeout messages prevents announce receipt timeout from occurring until at least announceReceiptTimeout Announce messages have been lost. Note that networks with high packet loss can still experience announce receipt timeout under high-packet-loss conditions; however, the announce receipt timeout condition occurs only after at least announceReceiptTimeout Announce messages have been lost.

**10.3.10.3 oldAnnounceInterval:** The saved value of the previous announce interval, when a new announce interval is requested via a Signaling message that contains a message interval request TLV. The data type for oldAnnounceInterval is UScaledNs. This variable is used by both the BMCA and the ~~explicit-port state~~external port configuration option.

**10.3.10.4 infoIs:** An Enumeration2 that takes the values Received, Mine, Aged, or Disabled to indicate the origin and state of the PTP Port's time-synchronization spanning tree information:

- a) If infoIs is Received, the PTP Port has received current information (i.e., announce receipt timeout has not occurred and, if gmPresent is TRUE, sync receipt timeout also has not occurred) from the master PTP Instance for the attached gPTP communication path.
- b) If infoIs is Mine, information for the PTP Port has been derived from the SlavePort for the PTP Instance (with the addition of SlavePort stepsRemoved). This includes the possibility that the SlavePort is the PTP Port whose portNumber is 0, i.e., the PTP Instance is the root of the gPTP domain.
- c) If infoIs is Aged, announce receipt timeout or, when gmPresent is TRUE, sync receipt timeout has occurred.
- d) If portOper, ptpPortEnabled, and asCapable are not all TRUE, infoIs is Disabled.

The variable infoIs is used only by the BMCA, i.e., not by the ~~explicit-port state~~external port configuration option.

**10.3.10.5 masterPriority:** The masterPriorityVector for the PTP Port. The data type for masterPriority is UInteger224 (see 10.3.4). This variable is used by both the BMCA and the ~~explicit-port state~~external port configuration option.

**10.3.10.6 currentLogAnnounceInterval:** The current value of the logarithm to base 2 of the mean time interval, in seconds, between the sending of successive Announce messages (see 10.7.2.2). This value is set

in the `AnnounceIntervalSetting` state machine (see 10.3.17). The data type for `currentLogAnnounceInterval` is `Integer8`. This variable is used by both the BMCA and the ~~explicit-port-state~~external port configuration option.

**10.3.10.7 initialLogAnnounceInterval:** The initial value of the logarithm to base 2 of the mean time interval, in seconds, between the sending of successive Announce messages (see 10.7.2.2). The data type for `initialLogAnnounceInterval` is `Integer8`. This variable is used by both the BMCA and the ~~explicit-port-state~~external port configuration option.

**10.3.10.8 announceInterval:** A variable containing the mean Announce message transmission interval for the PTP Port. This value is set in the `AnnounceIntervalSetting` state machine (see 10.3.17). The data type for `announceInterval` is `UScaledNs`. This variable is used by both the BMCA and the ~~explicit-port-state~~external port configuration option.

**10.3.10.9 messageStepsRemoved:** The value of `stepsRemoved` contained in the received Announce information. The data type for `messageStepsRemoved` is `UInteger16`. This variable is used by both the BMCA and the ~~explicit-port-state~~external port configuration option.

**10.3.10.10 newInfo:** A Boolean variable that is set to cause a PTP Port to transmit Announce information; specifically, it is set when an announce interval has elapsed (see Figure ), PTP Port states have been updated, and `portPriority` and `portStepsRemoved` information has been updated with newly determined `masterPriority` and `masterStepsRemoved` information. This variable is used by both the BMCA and the ~~explicit-port-state~~external port configuration option.

**10.3.10.11 portPriority:** The `portPriorityVector` for the PTP Port. The data type for `portPriority` is `UInteger224` (see 10.3.4). This variable is used only by the BMCA, i.e., not by the ~~explicit-port-state~~external port configuration option.

**10.3.10.12 portStepsRemoved:** The value of `stepsRemoved` for the PTP Port. `portStepsRemoved` is set equal to `masterStepsRemoved` (see 10.3.9.3) after `masterStepsRemoved` is updated. The data type for `portStepsRemoved` is `UInteger16`. This variable is used by both the BMCA and the ~~explicit-port-state~~external port configuration option.

**10.3.10.13 rcvdAnnouncePtr:** A pointer to a structure that contains the fields of a received Announce message. This variable is used by both the BMCA and the explicit PTP Port state configuration option.

**10.3.10.14 rcvdMsg:** A Boolean variable that is `TRUE` if a received Announce message is qualified and `FALSE` if it is not qualified. This variable is used only by the BMCA, i.e., not by the ~~explicit-port-state~~external port configuration option.

**10.3.10.15 updtInfo:** A Boolean variable that is set to `TRUE` to indicate that the `PortAnnounceInformation` state machine (see 10.3.12) should copy the newly determined `masterPriority` and `masterStepsRemoved` to `portPriority` and `portStepsRemoved`, respectively. This variable is used by both the BMCA and the ~~explicit-port-state~~external port configuration option; however, its value does not impact the ~~explicit-port-state~~external port configuration option (see the NOTE in 10.3.16.3).

**10.3.10.16 annLeap61:** A global variable in which the `leap61` flag (see 10.6.2.2.8) of a received Announce message is saved. The data type for `annLeap61` is `Boolean`. This variable is used by both the BMCA and the ~~explicit-port-state~~external port configuration option.

**10.3.10.17 annLeap59:** A global variable in which the `leap59` flag (see 10.6.2.2.8) of a received Announce message is saved. The data type for `annLeap59` is `Boolean`. This variable is used by both the BMCA and the ~~explicit-port-state~~external port configuration option.

**10.3.10.18 annCurrentUtcOffsetValid:** A global variable in which the currentUtcOffsetValid flag (see 10.6.2.2.8) of a received Announce message is saved. The data type for annCurrentUtcOffsetValid is Boolean. This variable is used by both the BMCA and the ~~explicit port state~~external port configuration option.

**10.3.10.19 annPtpTimescale:** A global variable in which the ptpTimescale flag (see 10.6.2.2.8) of a received Announce message is saved. The data type for annPtpTimescale is Boolean. This variable is used by both the BMCA and the ~~explicit port state~~external port configuration option.

**10.3.10.20 annTimeTraceable:** A global variable in which the timeTraceable flag (see 10.6.2.2.8) of a received Announce message is saved. The data type for annTimeTraceable is Boolean. This variable is used by both the BMCA and the ~~explicit port state~~external port configuration option.

**10.3.10.21 annFrequencyTraceable:** A global variable in which the frequencyTraceable flag (see 10.6.2.2.8) of a received Announce message is saved. The data type for annFrequencyTraceable is Boolean. This variable is used by both the BMCA and the ~~explicit port state~~external port configuration option.

**10.3.10.22 annCurrentUtcOffset:** A global variable in which the currentUtcOffset field (see 10.6.3.2.1) of a received Announce message is saved. The data type for annCurrentUtcOffset is Integer16. This variable is used by both the BMCA and the ~~explicit port state~~external port configuration option.

**10.3.10.23 annTimeSource:** A global variable in which the timeSource field (see 10.6.3.2.1) of a received Announce message is saved. The data type for annTimeSource is TimeSource (see 8.6.2.7). This variable is used by both the BMCA and the ~~explicit port state~~external port configuration option.

**10.3.10.24 receivedPathTrace:** An array in which the pathSequence array field of the path trace TLV of the most recently received Announce message is saved. The data type for receivedPathTrace is clockIdentity[N], where N is the number of entries in the pathSequence array field.

### 10.3.11 PortAnnounceReceive state machine

*Change 10.3.11.3 as follows:*

#### 10.3.11.3 State diagram

The PortAnnounceReceive state machine shall implement the function specified by the state diagram in Figure 10-13, the local variable specified in 10.3.11.1, the function specified in 10.3.11.2, and the relevant global variables specified in 10.2.4, 10.2.5, , ~~and 10.3.10, and 11.2.13~~. The state machine is not used if externalPortConfigurationEnabled is TRUE. The state machine receives Announce information from the MD entity of the same PTP Port, determines if the Announce message is qualified, and if so, sets the rcvdMsg variable.

### 10.3.12 PortAnnounceInformation state machine

*Add the following definition to 10.3.12.1:*

#### 10.3.12.1 State machine variables

**10.3.12.1.4 TEMP:** A temporary variable used to reduce clutter in the state diagram (see Figure 10-14). The data type for TEMP is Integer16.

*Change 10.3.12.3 as follows:*



### 10.3.12.3 State diagram

The PortAnnounceInformation state machine shall implement the function specified by the state diagram in Figure 10-14, the local variables specified in 10.3.12.1, the functions specified in 10.3.12.2, and the relevant global variables specified in 10.2.4, 10.2.5, , and 10.3.10, and 11.2.13. This state machine is used only if externalPortConfigurationEnabled is FALSE (if this variable is TRUE, the PortAnnounceInformationExt state machine of 10.3.14.3 is used instead). The state machine receives new qualified Announce information from the PortAnnounceReceive state machine (see 10.3.11) of the same PTP Port and determines if the Announce information is better than the current best master information it knows. The state machine also updates the current best master information when it receives updated PTP Port state information from the PortStateSelection state machine (see 10.3.13) and when announce receipt timeout or, when gmPresent is TRUE, sync receipt timeout occurs.

### 10.3.13 PortStateSelection state machine

#### 10.3.13.2 State machine functions

*Change 10.3.13.2.4 as follows:*

**10.3.13.2.4 UpdtStatesTree():** Performs the following operations (see 10.3.4 and 10.3.5 for details on the priority vectors):

- a) Computes the gmPathPriorityVector for each PTP Port that has a portPriorityVector and for which neither announce receipt timeout nor, if gmPresent is TRUE, sync receipt timeout have occurred,
- b) Saves gmPriority (see 10.3.9.21) in lastGmPriority (see 10.3.9.22), computes the gmPriorityVector for the PTP Instance and saves it in gmPriority, chosen as the best of the set consisting of the systemPriorityVector (for this PTP Instance) and the gmPathPriorityVector for each PTP Port for which the clockIdentity of the master port is not equal to thisClock (see 10.2.4.22),
- c) Sets the per PTP Instance global variables leap61, leap59, currentUtcOffsetValid, ptpTimescale, timeTraceable, frequencyTraceable, currentUtcOffset, and timeSource as follows:
  - 1) If the gmPriorityVector was set to the gmPathPriorityVector of one of the ports, then leap61, leap59, currentUtcOffsetValid, ptpTimescale, timeTraceable, frequencyTraceable, currentUtcOffset, and timeSource are set to annLeap61, annLeap59, annCurrentUtcOffsetValid, annPtpTimescale, annTimeTraceable, annFrequencyTraceable, annCurrentUtcOffset, and annTimeSource, respectively, for that PTP Port.
  - 2) If the gmPriorityVector was set to the systemPriorityVector, then leap61, leap59, currentUtcOffsetValid, ptpTimescale, timeTraceable, frequencyTraceable, currentUtcOffset, and timeSource are set to sysLeap61, sysLeap59, sysCurrentUtcOffsetValid, sysPtpTimescale, sysTimeTraceable, sysFrequencyTraceable, sysCurrentUtcOffset, and sysTimeSource, respectively.
- d) Updates the timePropertiesDS members leap61, leap59, currentUtcOffsetValid, ptpTimescale, timeTraceable, frequencyTraceable, currentUtcOffset, and timeSource with the values of the corresponding global variables computed in item c) above.
- e) Updates the parentDS members grandmasterIdentity, grandmasterClockQuality.clockClass, grandmasterClockQuality.clockAccuracy, grandmasterClockQuality.offsetScaledLogVariance, grandmasterPriority1, grandmasterPriority2 with the clockIdentity, clockClass, clockAccuracy, offsetScaledLogVariance, priority1, and priority2 attributes, respectively, of the systemIdentity component of the gmPriority computed in item b) above.
- f) Updates the parentDS member parentPortIdentity with the sourcePortIdentity component of the gmPriority computed in item b) above.

- g) Computes the masterPriorityVector for each PTP Port.
- h) Computes masterStepsRemoved, which is equal to one of the following:
  - 1) messageStepsRemoved (see 10.3.10.9) for the PTP Port associated with the gmPriorityVector, incremented by 1, if the gmPriorityVector is not the systemPriorityVector, or
  - 2) 0 if the gmPriorityVector is the systemPriorityVector.
- i) Sets currentDS.stepsRemoved equal to masterStepsRemoved.
- j) Assigns the PTP Port state for PTP Port j, and sets selectedState[j] equal to this PTP Port state, as follows, for j = 1, 2, ..., numberPorts:
  - 1) If the PTP Port is disabled (infoIs == Disabled), then selectedState[j] is set to DisabledPort.
  - 2) If asymmetryMeasurementMode is TRUE, then selectedState[j] is set to PassivePort, and updtInfo is set to FALSE.
  - 3) If announce receipt timeout, or sync receipt timeout with gmPresent set to TRUE, has occurred (infoIs = Aged), then selectedState[j] is set to MasterPort, and updtInfo is set to TRUE.
  - 4) If the portPriorityVector was derived from another PTP Port on the PTP Instance or from the PTP Instance itself as the root (infoIs == Mine), then selectedState[j] is set to MasterPort. In addition, updtInfo is set to TRUE if the portPriorityVector differs from the masterPriorityVector or portStepsRemoved differs from masterStepsRemoved.
  - 5) If the portPriorityVector was received in an Announce message, announce receipt timeout, or sync receipt timeout with gmPresent TRUE, has not occurred (infoIs == Received), and the gmPriorityVector is now derived from the portPriorityVector, then selectedState[j] is set to SlavePort, and updtInfo is set to FALSE. The per port global variable receivedPathTrace, for this port, is copied to the per PTP Instance global array pathTrace, and, if it is not empty, thisClock is appended to pathTrace.
  - 6) If the portPriorityVector was received in an Announce message, announce receipt timeout, or sync receipt timeout with gmPresent TRUE, has not occurred (infoIs == Received), the gmPriorityVector is not now derived from the portPriorityVector, the masterPriorityVector is not better than the portPriorityVector, and the sourcePortIdentity component of the portPriorityVector *does not* reflect another PTP Port on the PTP Instance, then selectedState[j] is set to PassivePort, and updtInfo is set to FALSE.
  - 7) If the portPriorityVector was received in an Announce message, announce receipt timeout, or sync receipt timeout with gmPresent TRUE, has not occurred (infoIs == Received), the gmPriorityVector is not now derived from the portPriorityVector, the masterPriorityVector is not better than the portPriorityVector, and the sourcePortIdentity component of the portPriorityVector *does* reflect another PTP Port on the PTP Instance, then selectedState[j] is set to PassivePort, and updtInfo is set to FALSE.
  - 8) If the portPriorityVector was received in an Announce message, announce receipt timeout, or sync receipt timeout with gmPresent TRUE, has not occurred (infoIs == Received), the gmPriorityVector is not now derived from the portPriorityVector, and the masterPriorityVector is better than the portPriorityVector, then selectedState[j] is set to MasterPort, and updtInfo is set to TRUE.
- k) Updates gmPresent as follows:
  - 1) gmPresent is set to TRUE if the priority1 field of the rootSystemIdentity of the gmPriorityVector is less than 255.
  - 2) gmPresent is set to FALSE if the priority1 field of the rootSystemIdentity of the gmPriorityVector is equal to 255.
- l) Assigns the PTP Port state for PTP Port 0 (see 8.5.2.3), and sets selectedState[0] as follows:
  - 1) if selectedState[j] is set to SlavePort for any PTP Port with portNumber j, j = 1, 2, ..., numberPorts, selectedState[0] is set to PassivePort.

- 2) if selectedState[j] is *not* set to SlavePort for any PTP Port with portNumber j, j = 1, 2, ..., numberPorts, selectedState[0] is set to SlavePort.
- m) If the clockIdentity member of the systemIdentity (see 10.3.2) member of gmPriority (see 10.3.9.21) is equal to thisClock (see 10.2.4.22), i.e., if the current PTP Instance is the Grandmaster PTP Instance, the pathTrace array is set to contain the single element thisClock (see 10.2.4.22).

***Change 10.3.13.3 as follows:***

### **10.3.13.3 State diagram**

The PortStateSelection state machine shall implement the function specified by the state diagram in Figure 10-15, the functions specified in 10.3.13.1, and the relevant global variables specified in 10.2.4, 10.2.5, , *and* 10.3.10, ~~and 11.2.13~~. This state machine is used only if externalPortConfigurationEnabled is FALSE (if this variable is TRUE, the PortStateSettingExt state machine is used instead). The state machine updates the gmPathPriority vector for each PTP Port of the PTP Instance, the gmPriorityVector for the PTP Instance, and the masterPriorityVector for each PTP Port of the PTP Instance. The state machine determines the PTP Port state for each PTP Port and updates gmPresent.

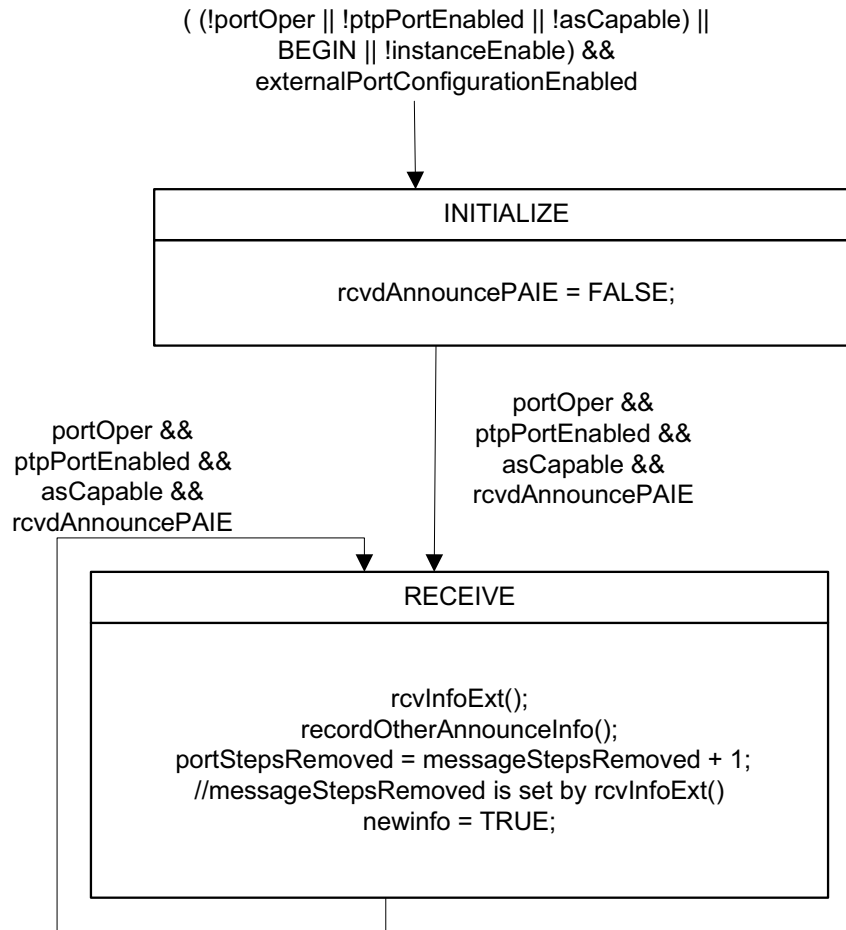
### **10.3.14 PortAnnounceInformationExt state machine**

***Change 10.3.14.3 as follows:***

### **10.3.14.3 State diagram**

The PortAnnounceInformationExt state machine shall implement the function specified by the state diagram in Figure 10-16, the local variables specified in 10.3.14.1, the functions specified in 10.3.14.2, and the relevant global variables specified in 10.2.4, 10.2.5, , *and* 10.3.10, ~~and 11.2.13~~. This state machine is used only if externalPortConfigurationEnabled is TRUE (if this variable is FALSE, the PortAnnounceInformation state machine of 10.3.12.3 is used instead). The state machine receives Announce information from the MD entity of the same PTP Port and saves the information..





**NOTE**—This figure differs from the 2020 edition, with Corrigendum 1 applied, of this standard in that the statement `newinfo = TRUE;` is added after the existing statements in the RECEIVE state.

**Figure 10-16—PortAnnounceInformationExt state machine**

### 10.3.15 PortStateSettingExt state machine

*Change 10.3.15.2.2 as follows:*

**10.3.15.2.2 updtPortState(j):** Performs the following operations for PTP Port *j* (see 10.3.4 and 10.3.5 for details on the priority vectors):

- a) Sets the per PTP Instance global variables `leap61`, `leap59`, `currentUtcOffsetValid`, `ptpTimescale`, `timeTraceable`, `frequencyTraceable`, `currentUtcOffset`, and `timeSource` as follows:
  - 1) If the PTP Port state of any PTP Port of this PTP Instance ~~other than PTP Port 0~~ (see 8.5.2.3) is `SlavePort`, then `leap61`, `leap59`, `currentUtcOffsetValid`, `ptpTimescale`, `timeTraceable`, `frequencyTraceable`, `currentUtcOffset`, and `timeSource` are set to `annLeap61`, `annLeap59`, `annCurrentUtcOffsetValid`, `annPtpTimescale`, `annTimeTraceable`, `annFrequencyTraceable`, `annCurrentUtcOffset`, and `annTimeSource`, respectively, for that PTP Port.

- 2) If no PTP Port of this PTP Instance ~~other than PTP Port 0 (see 8.5.2.3)~~ has the PTP Port state SlavePort, then leap61, leap59, currentUtcOffsetValid, ptpTimescale, timeTraceable, frequencyTraceable, currentUtcOffset, and timeSource are set to sysLeap61, sysLeap59, sysCurrentUtcOffsetValid, sysPtpTimescale, sysTimeTraceable, sysFrequencyTraceable, sysCurrentUtcOffset, and sysTimeSource, respectively.
- b) Update the timePropertiesDS members leap61, leap59, currentUtcOffsetValid, ptpTimescale, timeTraceable, frequencyTraceable, currentUtcOffset, and timeSource with the values of the corresponding global variables computed in item a) above.
- c) Computes masterStepsRemoved as follows:
  - 1) If the PTP Port state of any PTP Port of this PTP Instance ~~other than PTP Port 0 (see 8.5.2.3)~~ is SlavePort, then masterStepsRemoved is set equal to portStepsRemoved for that PTP Port.
  - 2) If no PTP Port of this PTP Instance ~~other than PTP Port 0 (see 8.5.2.3)~~ has the PTP Port state SlavePort, then masterStepsRemoved is set equal to 0.
- d) Sets currentDS.stepsRemoved equal to masterStepsRemoved.
- e) Assigns the PTP Port state for PTP Port j, and sets selectedState[j] equal to this PTP Port state, as follows:
  - 1) If disabledExt is TRUE, selectedState[j] is set to DisabledPort, else
  - 2) If asymmetryMeasurementMode is TRUE, selectedState[j] is set to PassivePort, else
  - 3) selectedState[j] is set to portStateInd. If portStateInd is equal to MasterPort, newInfo is set to TRUE.
- f) Updates gmPresent as follows:
  - 1) If the PTP Port state of any PTP Port of this PTP Instance ~~other than PTP Port 0 (see 8.5.2.3)~~ is SlavePort and the priority1 field of the rootSystemIdentity of the messagePriorityPAIE of the slave port is less than 255, gmPresent is set to TRUE, else
  - 2) If the PTP Port state of any PTP Port of this PTP Instance ~~other than PTP Port 0 (see 8.5.2.3)~~ is SlavePort and the priority1 field of the rootSystemIdentity of the messagePriorityPAIE of the slave PTP Port is equal to 255, gmPresent is set to FALSE, else
  - 3) If no PTP Port of this PTP Instance ~~other than PTP Port 0 (see 8.5.2.3)~~ has the PTP Port state SlavePort, gmPresent is set to TRUE if priority1 for this PTP Instance is less than 255 and FALSE if priority1 for this PTP Instance is equal to 255.
- g) Assigns the PTP Port state for PTP Port 0, and sets selectedState[0] as follows:
  - 1) If selectedState[j] is set to SlavePort, selectedState[0] is set to PassivePort.
  - 2) If selectedState[j] is *not* set to SlavePort and selectedState[k] is not equal to SlavePort for every k not equal to 0 or j, selectedState[0] is set to SlavePort.
- h) Computes the gmPriorityVector as follows:
  - 1) If selectedState[j] is set to SlavePort, the gmPriorityVector is set equal to messagePriorityPAIE for PTP Port j.
  - 2) If selectedState[j] is *not* set to SlavePort and selectedState[k] is not equal to SlavePort for every k not equal to 0 or j, the gmPriorityVector is set equal to the systemPriorityVector.
- i) Update the parentDS members grandmasterIdentity, grandmasterClockQuality.clockClass, grandmasterClockQuality.clockAccuracy, grandmasterClockQuality.offsetScaledLogVariance, grandmasterPriority1, grandmasterPriority2 with the clockIdentity, clockClass, clockAccuracy, offsetScaledLogVariance, priority1, and priority2 attributes, respectively, of the systemIdentity component of the gmPriorityVector computed in item i) above.
- j) Update the parentDS member parentPortIdentity with the sourcePortIdentity component of the gmPriorityVector computed in item i) above.
- k) Computes the masterPriorityVector for PTP Port j.

- l) If no PTP Port of this PTP Instance has the PTP Port state SlavePort, the pathTrace array is set to contain the single element thisClock (see 10.2.4.22).

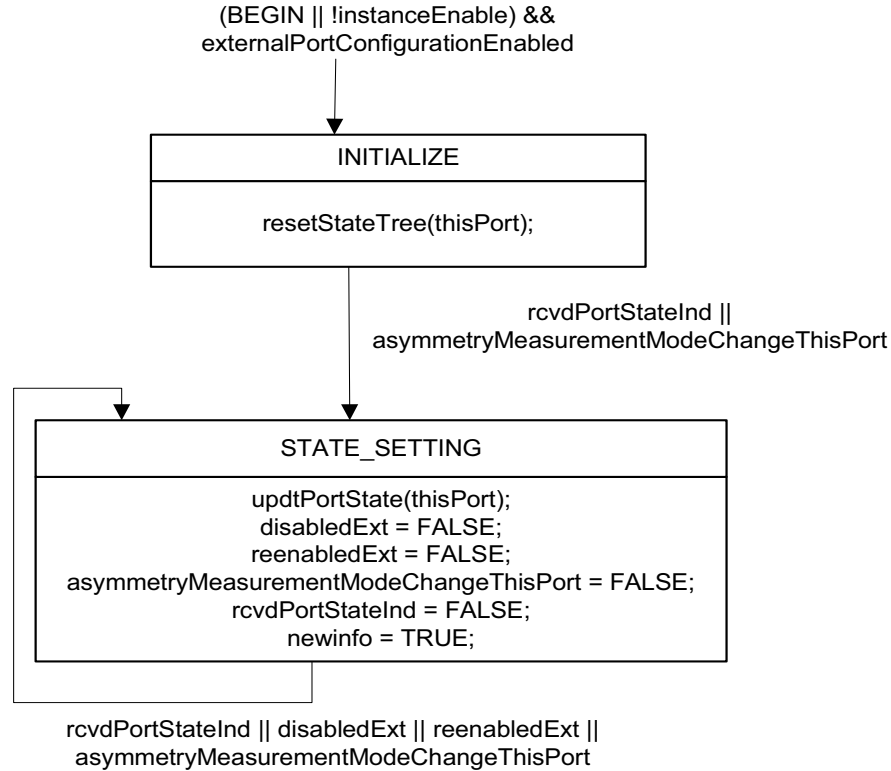
***Change 10.3.15.3 as follows:***

### 10.3.15.3 State diagram

The PortStateSettingExt state machine shall implement the function specified by the state diagram in Figure , the local variables specified in 10.3.15.1, the functions specified in 10.3.15.2, and the relevant global variables specified in 10.2.4, 10.2.5, , [and](#) 10.3.10, ~~and 11.2.13~~. This state machine is used only if externalPortConfigurationEnabled is TRUE (if this variable is FALSE, the PortStateSelection state machine of 10.3.13.3 is used instead). A separate instance of this state machine runs on each PTP Port (unlike the PortStateSelection state machine, for which a single instance runs in the PTP Instance and performs operations on all the ports).

The state machine updates the gmPriorityVector for the PTP Instance and the masterPriorityVector for each PTP Port of the PTP Instance. The state machine determines the PTP Port state for each PTP Port and updates gmPresent.

**NOTE 1**—It is possible to use the external port configuration mechanism to misconfigure the network, e.g., to produce a configuration where one or more PTP Instances have more than one slave port. Detecting and correcting misconfigurations is outside the scope of this standard.



**NOTE 2**—This figure differs from the 2020 edition, with Corrigendum 1 applied, of this standard in that the statement newinfo = TRUE; is added after the existing statements in the STATE\_SETTING state.

**Figure 10-17—PortStateSettingExt state machine**

### 10.3.16 PortAnnounceTransmit state machine

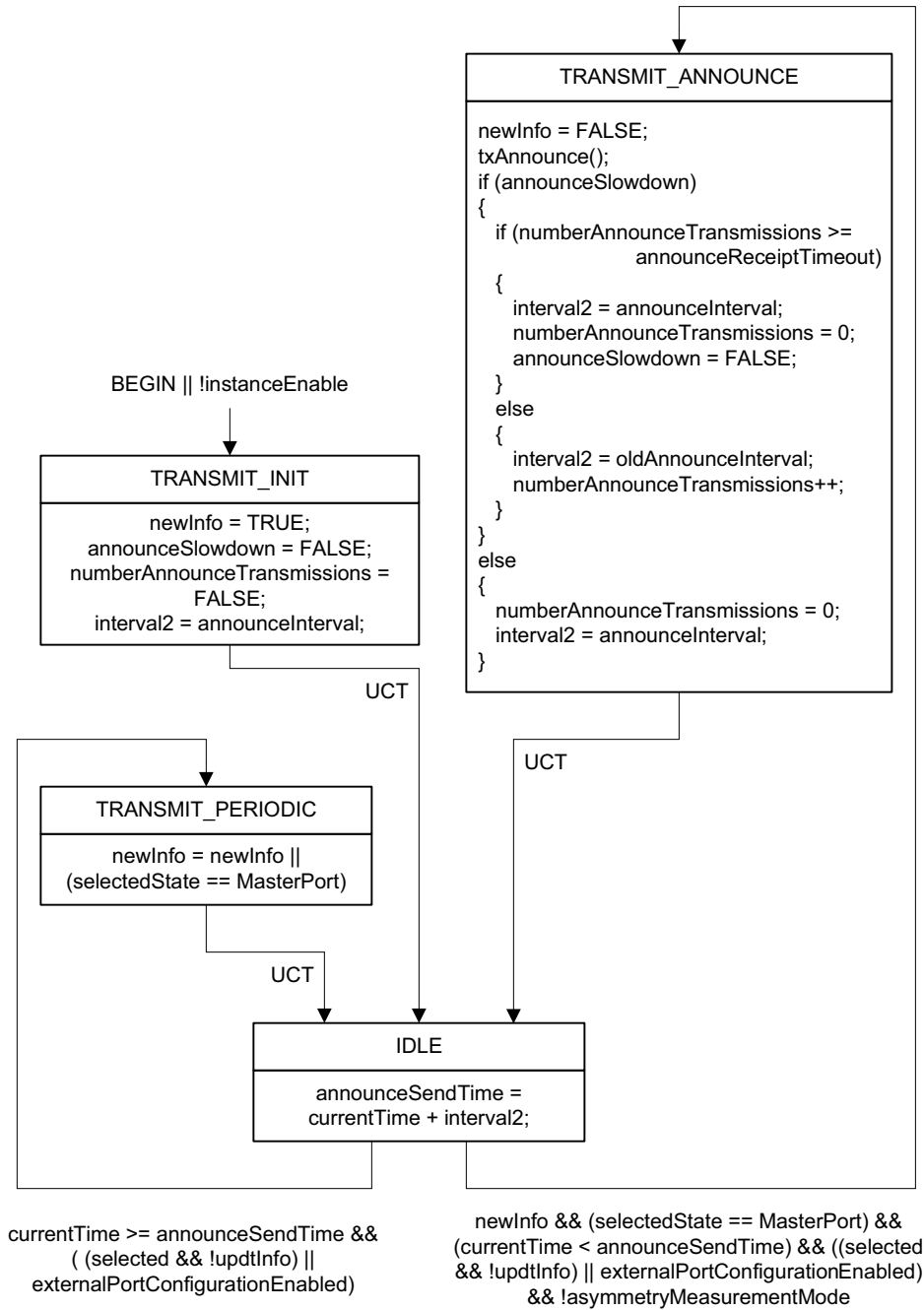
*Change 10.3.16.3 as follows:*

#### 10.3.16.3 State diagram

The PortAnnounceTransmit state machine shall implement the function specified by the state diagram in Figure 10-17, the local variables specified in 10.3.16.1, the functions specified in 10.3.16.2, and the relevant global variables specified in 10.2.4, 10.2.5, , and 10.3.10, ~~and 11.2.13~~. The state machine transmits Announce information to the MD entity when an announce interval has elapsed, PTP Port states have been updated, and portPriority and portStepsRemoved information has been updated with newly determined masterPriority and masterStepsRemoved information.

NOTE 1—When the external port configuration option is used (i.e., externalPortConfigurationEnabled is TRUE; see 10.3.9.24) the values of the variables updInfo and selected do not affect the operation of the PortAnnounceTransmit state machine because the term of the conditions in which they appear, i.e., (selected && !updInfo) || externalPortConfiguartionEnabled, evaluates to TRUE when externalPortConfigurationEnabled is TRUE.

*Replace Figure 10-18 with the following:*



**NOTE 2**—This figure differs from the 2020 edition, with Corrigendum 1 applied, of this standard in that the statement `interval2 = announceInterval` is added after the existing statements in the `TRANSMIT_INIT` state.

**Figure 10-18—PortAnnounceTransmit state machine**

### 10.3.17 AnnounceIntervalSetting state machine

*Add the following definition to 10.3.17.1:*

#### 10.3.17.1 State machine variables

**10.3.17.1.6 TEMP:** A temporary variable used to reduce clutter in the state diagram (see Figure ). The data type for TEMP is Integer16.

#### 10.3.17.2 State machine functions

*Change 10.3.17.2.2 as follows:*

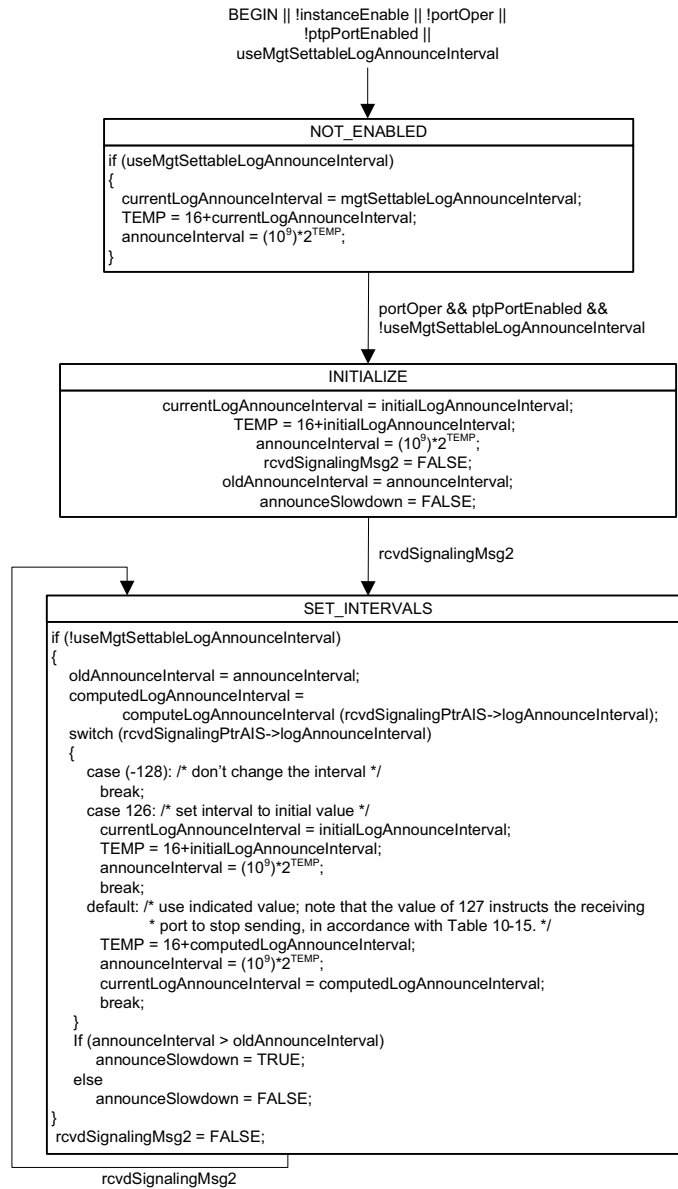
**10.3.17.2.2 computeLogAnnounceInterval (logRequestedAnnounceInterval):** An Integer8 function that computes and returns the logAnnounceInterval, based on the logRequestedAnnounceInterval. This function is defined as indicated below. It is defined here so that the detailed code that it invokes does not need to be placed into the state machine diagram.

```
Integer8 computeLogAnnounceInterval (logRequestedAnnounceInterval)
Integer8 logRequestedAnnounceInterval;
{
    Integer8 logSupportedAnnounceIntervalMax,
              logSupportedClosestLongerAnnounceInterval;
    if (isSupportedLogAnnounceInterval (logRequestedAnnounceInterval))
        // The requested Announce Interval is supported and returned
        return (logRequestedAnnounceInterval);
    else
    {
        if (logRequestedAnnounceInterval > logSupportedAnnounceIntervalMax)
            // Return the fastestslowest supported rate, even if faster than the
            requested rate
            return (logSupportedAnnounceIntervalMax);
        else
            // Return the fastest supported rate that is still slower than
            // the requested rate.
            return (logSupportedClosestLongerAnnounceInterval);
    }
}
```

### 10.3.17.3 State diagram

*Replace Figure 10-19 with the following:*





**NOTE**—This figure differs from the 2020 edition, with Corrigendum 1 applied, of this standard in that the variable `LogAnnounceInterval` at the end of line 5 of the `SET_INTERVALS` state is changed to `logAnnounceInterval` (i.e., the capitalization is corrected), and the statement `oldAnnounceInterval = announceInterval` is added immediately after the first open curly brace in the `SET_INTERVALS` state.

**Figure 10-19—AnnounceIntervalSetting state machine**

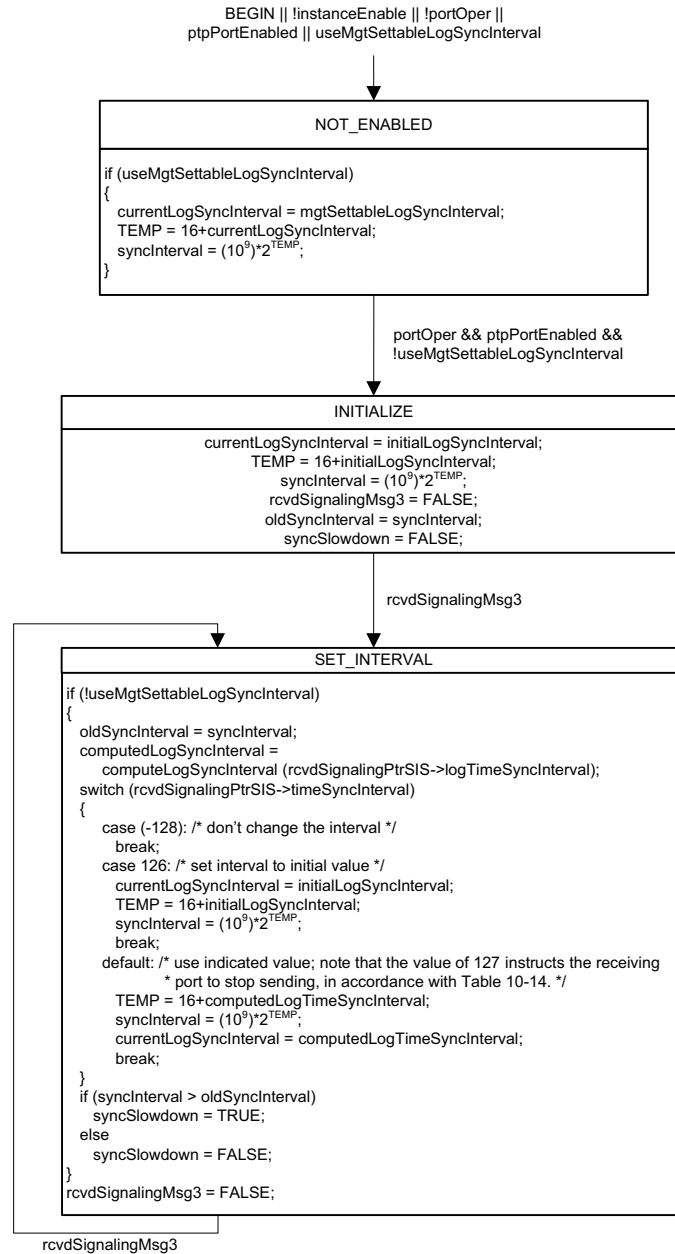
**10.3.18 SyncIntervalSetting state machine***Add the following definition to 10.3.18.1:***10.3.18.1 State machine variables****10.3.18.1.6 TEMP:** A temporary variable used to reduce clutter in the state diagram (see Figure ). The data type for TEMP is Integer16.**10.3.18.2 State machine functions***Change 10.3.18.2.2 as follows:***10.3.18.2.2 computeLogSyncInterval (logRequestedSyncInterval):** An Integer8 function that computes and returns the logSyncInterval, based on the logRequestedSyncInterval. This function is defined as indicated below. It is defined here so that the detailed code that it invokes does not need to be placed into the state machine diagram.

```

Integer8 computeLogSyncInterval (logRequestedSyncInterval)
Integer8 logRequestedSyncInterval;
{
    Integer8 logSupportedSyncIntervalMax, logSupportedClosestLongerSyncInterval;
    if (isSupportedLogSyncInterval (logRequestedSyncInterval))
        // The requested Sync Interval is supported and returned
        return (logRequestedSyncInterval);
    else
    {
        if (logRequestedSyncInterval > logSupportedSyncIntervalMax)
            // Return the fastedslowest supported rate, even if faster than the
            requested rate
            return (logSupportedSyncIntervalMax);
        else
            // Return the fastest supported rate that is still slower than
            // the requested rate.
            return (logSupportedClosestLongerSyncInterval);
    }
}

```

**10.3.18.3 State diagram***Replace Figure 10-20 with the following:*



**NOTE**—This figure differs from the 2020 edition, with Corrigendum 1 applied, of this standard in that the statement `oldSyncInterval = syncInterval` is added immediately after the first open curly brace in the SET\_INTERVALS state

**Figure 10-20—SyncIntervalSetting state machine**

**Delete the NOTE immediately after Figure 10-20.**

*Delete the NOTE immediately after Figure 10-20.*

## 10.4 State machines related to signaling gPTP capability

*Insert 10.4.1, and renumber subsequent subclauses as appropriate:*

### 10.4.1 Enabling and disabling the state machines

If the managed object `gtpCapableStateMachinesEnabled` (see 14.8.55) is TRUE, the `GtpCapableTransmit` and `GtpCapableReceive` state machines shall be enabled and shall function as specified in 10.4.2 and 10.4.3, respectively. The `GtpCapableIntervalSetting` state machine is enabled if it is implemented.

If the managed object `gtpCapableStateMachinesEnabled` is FALSE, the `GtpCapableTransmit` and `GtpCapableReceive` state machines shall be disabled. The `GtpCapableIntervalSetting` state machine shall be disabled if it is implemented.

If the managed object `gtpCapableStateMachinesEnabled` is FALSE, the global variable `neighborGtpCapable` for the port (see 10.2.5.16) shall be set to TRUE.

NOTE 1—The global variable `neighborGtpCapable` indicates to a PTP Port whether the PTP Port at the other end of the attached PTP Link is capable of invoking gPTP, and is used in the determination of `asCapable` (see 10.2.5.1, 11.2.2, 12.4, and 13.4). If `gtpCapableStateMachinesEnabled` is TRUE, the variable is set to TRUE or FALSE by the `GtpCapableTransmit` state machine. If `gtpCapableStateMachinesEnabled` is FALSE, the variable is automatically set to TRUE, i.e., it is assumed that the network is engineered such that the PTP Port at the other end of the attached link is capable of invoking gPTP. In this case, it is essential that the network be engineered to fulfill this condition; if `neighborGtpCapable` is TRUE and the PTP Port at the other end of the link is not capable of invoking gPTP, undesirable behavior can occur.

NOTE 2—If the `GtpCapableTransmit` and `GtpCapableReceive` state machines are disabled, the exchange of peer delay messages can occur immediately, i.e., it is not necessary to wait until gPTP capability is determined by the `GtpCapableTransmit` and `GtpCapableReceive` state machines.

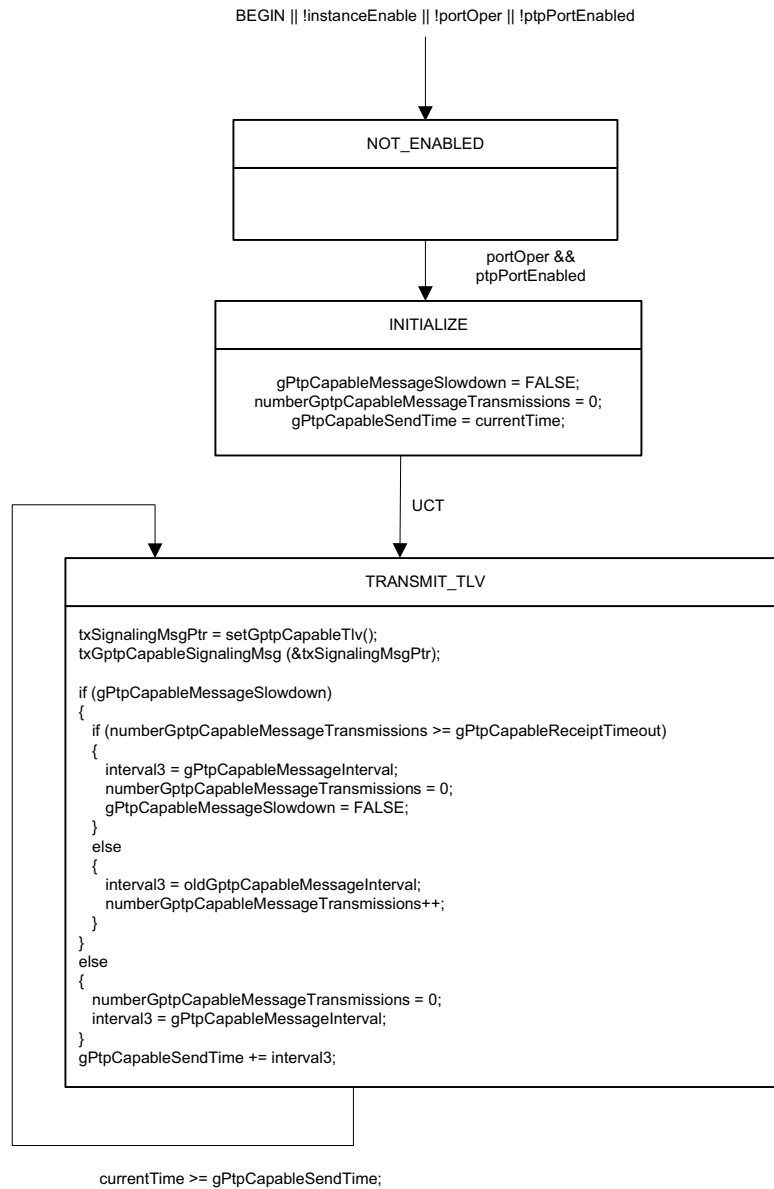
### 10.4.2 `GtpCapableTransmit` state machine

#### 10.4.2.1 State machine variables

*Change 10.4.2.1.1 as follows:*

**10.4.2.1.1 ~~intervalTimer~~gPtpCapableSendTime:** ~~A variable used to save the~~The time, relative to the  
LocalClock entity, at which~~when the gPTP-capable message interval timer is set (see Figure 10-21). A~~  
Signaling message containing a gPTP-capable TLV is next ~~sent when this timer expires~~. The data type for  
gPtpCapableSendTime~~intervalTimer~~ is UScaledNs.

***Replace Figure 10-21 with the following:***



**NOTE**—This figure differs from the 2020 edition, with Corrigendum 1 applied, of this standard in that the condition !domainEnabled is removed from the set of logical ORs at the ingress of the NOT\_ENABLED state and the variable intervalTimer is renamed gPtpCapableSendTime.

**Figure 10-21—GtpCapableTransmit state machine**

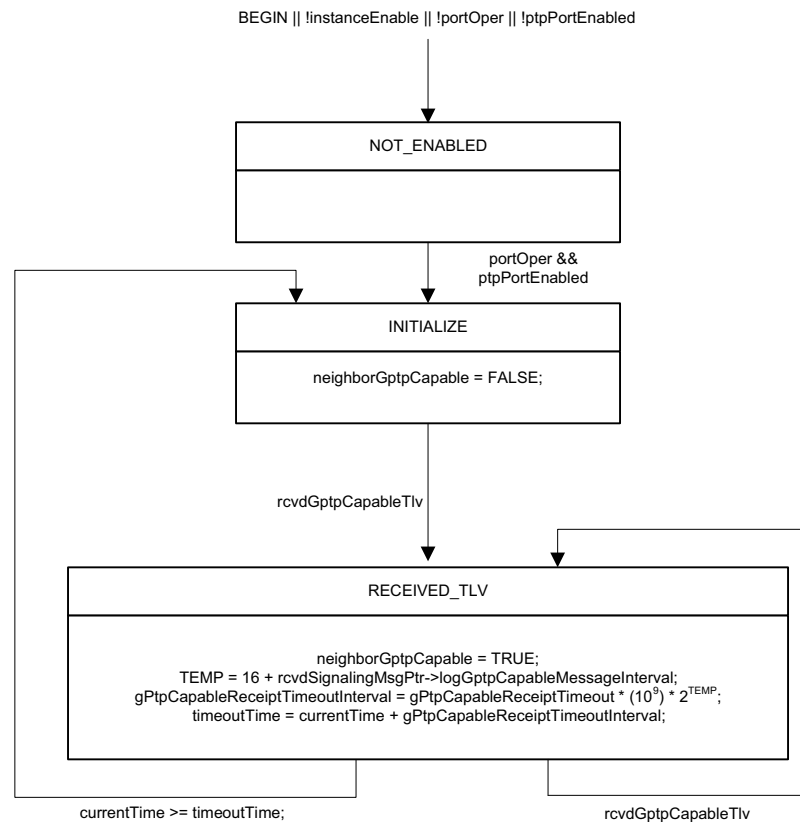
### 10.4.3 GptpCapableReceive state machine

*Add the following definition to 10.4.3.1:*

#### 10.4.3.1 State machine variables

**10.4.3.1.5 TEMP:** A temporary variable used to reduce clutter in the state diagram (see Figure ). The data type for TEMP is Integer16.

**Replace Figure 10-22 with the following:**



**NOTE**—This figure differs from the 2020 edition, with Corrigendum 1 applied, of this standard in that the condition !domainEnabled is removed from the set of logical ORs at the ingress of the NOT\_ENABLED state.

**Figure 10-22—GptpCapableReceive state machine**

**10.4.4 GtpCapableIntervalSetting state machine***Add the following definition to 10.4.4.1:***10.4.4.1 State machine variables****10.4.4.1.6 TEMP:** A temporary variable used to reduce clutter in the state diagram (see Figure ). The data type for TEMP is Integer16.**10.4.4.2 State machine functions***Change 10.4.4.2.2 as follows:***10.4.4.2.2 computeLogGtpCapableMessageInterval (logRequestedGtpCapableMessageInterval):**

An Integer8 function that computes and returns the logGtpCapableMessageInterval, based on the logRequestedGtpCapableMessageInterval. This function is defined as indicated below. It is defined here so that the detailed code that it invokes does not need to be placed into the state machine diagram.

```

Integer8 computeLogGtpCapableMessageInterval (logRequestedGtpCapableMessageInterval)
Integer8 logRequestedGtpCapableMessageInterval;
{
    Integer8 logSupportedGtpCapableMessageIntervalMax,
              logSupportedClosestLongerGtpCapableMessageInterval;
    if (isSupportedLogGtpCapableMessageInterval
        (logRequestedGtpCapableMessageInterval))
        // The requested gPTP-capable Message Interval is supported and returned
        return (logRequestedGtpCapableMessageInterval);
    else
    {
        if (logRequestedGtpCapableMessageInterval >
            logSupportedGtpCapableMessageIntervalMax)
            // Return the fastedslowest supported rate, even if faster than the
            requested rate
            return (logSupportedGtpCapableMessageIntervalMax);
        else

```

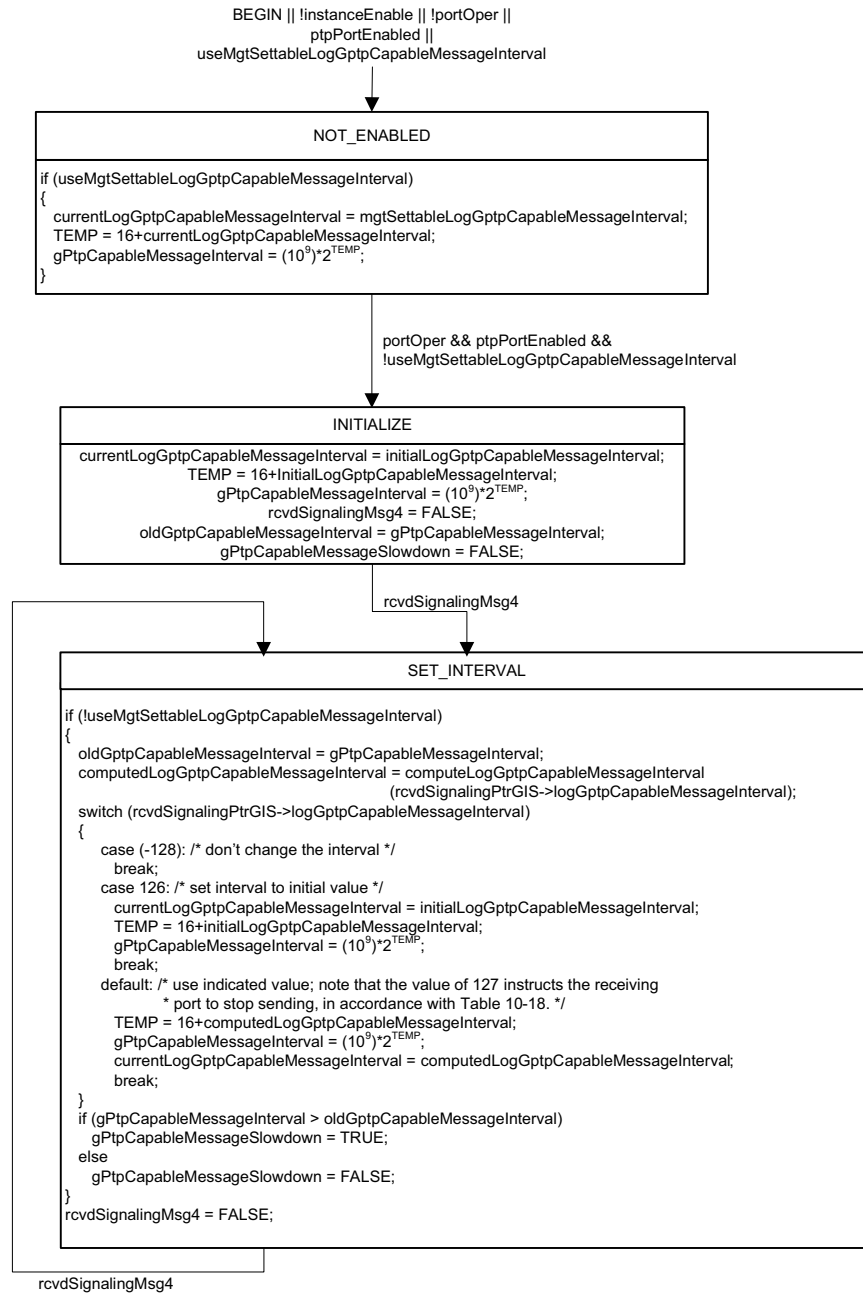


```
// Return the fastest supported rate that is still slower than  
// the requested rate.  
return (logSupportedClosestLongerGptpCapableMessageInterval);
```

```
    }  
}
```

#### 10.4.4.3 State diagram

*Replace Figure 10-23 with the following:*



**NOTE**—This figure differs from the 2020 edition, with Corrigendum 1 applied, of this standard in that the statement `oldGtpCapableMessageInterval = gPtpCapableMessageInterval` is added immediately after the first open curly brace in the SET\_INTERVALS state.

**Figure 10-23—GtpCapableIntervalSetting state machine**

*Delete the NOTE immediately after Figure 10-23.*

## 10.5 Message attributes

### 10.5.3 Addresses

*Add the following sentence, as a new paragraph, after the NOTE in 10.5.3.*

If the transport is full-duplex IEEE 802.3, all Announce and Signaling messages shall use the MAC address of the respective egress physical port as the source address.

*Change 10.5.5 as follows:*

### 10.5.5 Subtype

The subtype ~~of~~for the Announce and Signaling messages is indicated by the majorSdoId field (see 10.6.2.2.1).

~~NOTE—The subtype for all PTP messages is indicated by the majorSdoId field.~~

## 10.6 Message formats

### 10.6.2 Header

#### 10.6.2.2 Header field specifications

*Change 10.6.2.2.6 as follows:*

#### 10.6.2.2.6 domainNumber (UInteger8)

10.6.2 The value is the gPTP ~~domain-number~~domainNumber specified in 8.1.

#### 10.6.2.2.8 flags (Octet2)

*Change Table 10-9 as follows:*

#### 10.6.2.2.10 messageTypeSpecific (Octet4)

*Change Table 10-10 as follows:*

**Table 10-9—Values of flag bits**

Octet	Bit	Message types	Name	Value
0	0	All	alternateMasterFlag in Announce, Sync, Follow_Up, and Delay_Resp messages	Not used in this standard; transmitted as FALSE and ignored on reception
0	1	Sync, Pdelay_Resp	twoStepFlag	<p><i>For Sync messages:</i></p> <p>a) For a one-step transmitting PTP Port (see 11.1.3 and 11.2.13.9), the value is FALSE.</p> <p>b) For a two-step transmitting PTP Port, the value is TRUE.</p> <p><i>For Pdelay_Resp messages:</i></p> <p>The value is transmitted as TRUE and ignored on reception.</p>
0	2	All	unicastFlag	Not used in this standard; transmitted as FALSE and ignored on reception
0	3	All	Reserved	Not used by IEEE Std 1588-2019; reserved as FALSE and ignored on reception
0	4	All	Reserved	Not used by IEEE Std 1588-2019; reserved as FALSE and ignored on reception
0	5	All	PTP profileSpecific 1	Not used in this standard; transmitted as FALSE and ignored on reception
0	6	All	PTP profileSpecific 2	Not used in this standard; transmitted as FALSE and ignored on reception
0	7	All	Reserved	Not used in this standard; transmitted as FALSE and ignored on reception
1	0	Announce	leap61	The value of the global variable leap61 (see 10.3.9.4)
1	1	Announce	leap59	The value of the global variable leap59 (see 10.3.9.5)
1	2	Announce	currentUtcOffsetValid	The value of the global variable currentUtcOffsetValid (see 10.3.9.6)
1	3	Announce	ptpTimescale	<u>For domain 0, transmitted as TRUE and ignored on receipt. For domains other than domain 0, the value of the global variable ptpTimescale (see 10.3.9.7)</u>
1	4	Announce	timeTraceable	The value of the global variable timeTraceable (see 10.3.9.8)
1	5	Announce	frequencyTraceable	The value of the global variable frequencyTraceable (see 10.3.9.9)
1	6	All	Reserved	Not used by IEEE Std 1588-2019; reserved as FALSE and ignored on reception
1	7	All	Reserved	Not used in this standard; reserved as FALSE and ignored on reception

**Table 10-10—messageTypeSpecific semantics**

Value of messageType	Description
Follow_Up, Pdelay_Resp_Follow_Up, Announce, Signaling, Management	For the General message class, this field is reserved; it is transmitted as 0 and ignored on reception.
Sync, <del>Delay_Req</del> Pdelay_Req, Pdelay_Resp	For the Event message class, this field may be used for internal implementation as specified in this subclause.

#### 10.6.4 Signaling message

*Change 10.6.4.1 as follows:*

##### 10.6.4.1 General Signaling message specifications

The fields of the body of the Signaling message shall be as specified in .Table 10-13 and 10.6.4.2.

**Table 10-13—Signaling message fields**

Bits								Octets	Offset
7	6	5	4	3	2	1	0		
header (see 10.6.2)								34	0
targetPortIdentity								10	34
<del>message interval request TLV, gPTP-capable TLV, or</del> <del>gPTP-capable message interval request TLV</del> <u>one or more TLVs</u>								<del>16</del> <u>N</u>	44

*Change 10.6.4.2.2 as follows:*

##### 10.6.4.2.2 ~~Message interval request TLV or gPTP-capable TLV~~TLVs carried in one Signaling message

The Signaling message carries either: ~~the message interval request TLV, defined in 10.6.4.3, or the gPTP-capable TLV, defined in 10.6.4.4, but not both. If it is desired to send both TLVs, two Signaling messages must be sent.~~

- one or both interval request TLVs, defined in 10.6.4.3 and 10.6.4.5, or
- the gPTP-capable TLV, defined in 10.6.4.4.

*Change 10.6.4.5.6 as follows:*

##### 10.6.4.5.6 logGptpCapableMessageInterval (Integer8)

The value is the logarithm to base 2 of the mean time interval, desired by the PTP Port that sends this TLV, between successive Signaling messages that contain the gPTP-capable TLV (see 10.6.4.4), sent by the PTP

Port at the other end of the link. The format and allowed values of `logGptpCapableMessageInterval` are the same as the format and allowed values of `initialLogGptpCapableMessageInterval` (see 10.7.2.5).

The values 127, 126, and –128 are interpreted as defined in Table 10-14.

**Table 10-14—Interpretation of special values of `logGptpCapableMessageInterval`**

Value	Instruction to PTP Instance that receives this TLV
127	Instructs the PTP Port that receives this TLV to stop sending Signaling messages that contain the gPTP-capable TLV.
126	Instructs the PTP Port that receives this TLV to set <code>currentlogGptpCapableMessageInterval</code> to the value of <code>initialLogGptpCapableMessageInterval</code> (see 10.7.2.4).
–128	Instructs the PTP Port that receives this TLV not to change the mean time interval between successive Signaling messages that contain the gPTP-capable TLV.
All values in the ranges [–127, –25] and [25, 125] are reserved.	

When a Signaling message that contains this TLV is sent by a PTP Port, the value of `gPtpCapableReceiptTimeoutTimeInterval` for that PTP Port (see 10.3.10.1) shall be set equal to `gPtpCapableReceiptTimeout` (see 10.7.3.3) multiplied by the value of the interval, in seconds, reflected by `logGptpCapableMessageInterval`.

## 10.7 Protocol timing characterization

### 10.7.2 Message transmission intervals

#### *Change 10.7.2.2 as follows:*

##### 10.7.2.2 Announce message transmission interval

The logarithm to the base 2 of the announce interval (in seconds) is carried in the `logMessageInterval` field of the Announce message.

When `useMgtSettableLogAnnounceInterval` (see 14.8.14) is FALSE, the `initialLogAnnounceInterval` specifies the announce interval when the PTP Port is initialized and the value to which the announce interval is set when a message interval request TLV is received with the `logAnnounceInterval` field set to 126 (see the `AnnounceIntervalSetting` state machine in 10.3.17). The `currentLogAnnounceInterval` specifies the current value of the announce interval. The default value of `initialLogAnnounceInterval` is 0. Every PTP Port supports the value 127; the PTP Port does not send Announce messages when `currentLogAnnounceInterval` has this value (see 10.3.17). A PTP Port may support other values, except for the reserved values indicated in Table 10-17. A PTP Port ignores requests for unsupported values (see 10.3.17). The `initialLogAnnounceInterval` and `currentLogAnnounceInterval` are per-PTP Port attributes.

When `useMgtSettableLogAnnounceInterval` is TRUE, `currentLogAnnounceInterval` is set equal to `mgtSettableLogAnnounceInterval` (see 14.8.15), and `initialLogAnnounceInterval` is ignored.

Announce messages shall be transmitted such that the value of the arithmetic mean of the intervals, in seconds, between message transmissions is within  $\pm 30\%$  of  $2^{\text{currentLogAnnounceInterval}}$ . In addition, a PTP Port shall transmit Announce messages such that at least 90% of the inter-message intervals are within  $\pm 30\%$  of the value of  $2^{\text{currentLogAnnounceInterval}}$ . The interval between successive Announce messages should

not exceed twice the value of  $2^{\text{portDS} - \text{currentLogAnnounceInterval}}$  in order to prevent causing an announceReceiptTimeout event. The PortAnnounceTransmit state machine (see 10.3.16) is consistent with these requirements, i.e., the requirements here and the requirements of the PortAnnounceTransmit state machine can be met simultaneously.

NOTE 1—A minimum number of inter-message intervals is necessary in order to verify that a PTP Port meets these requirements. The arithmetic mean is the sum of the inter-message interval samples divided by the number of samples. For more detailed discussion of statistical analyses, see Papoulis [B25].

NOTE 2—If useMgtSettableLogAnnounceInterval is FALSE, the value of initialLogAnnounceInterval is the value of the mean time interval between successive Announce messages when the PTP Port is initialized. The value of the mean time interval between successive Announce messages can be changed, e.g., if the PTP Port receives a Signaling message that carries a message interval request TLV (see 10.6.4.3) and the current value is stored in currentLogAnnounceInterval. The value of the mean time interval between successive Announce messages can be reset to the initial value, e.g., by a message interval request TLV for which the value of the field logAnnounceInterval is 126 (see 10.6.4.3.8).

NOTE 3—A PTP Port that requests (using a Signaling message that contains a message interval request TLV; see 10.6.4 and 10.3.17) that the PTP Port at the other end of the attached link set its currentLogAnnounceInterval to a specific value can determine if the request was honored by examining the logMessageInterval field of subsequent received Announce messages.

### *Change 10.7.2.3 as follows:*

#### **10.7.2.3 Time-synchronization event message transmission interval**

The logarithm to the base 2 of the sync interval (in seconds) is carried in the logMessageInterval field of the time-synchronization messages.

When useMgtSettableLogSyncInterval (see 14.8.19) is FALSE, the initialLogSyncInterval specifies the sync interval when the PTP Port is initialized and the value to which the sync interval is set when a message interval request TLV is received with the logTimeSyncInterval field set to 126 (see the SyncIntervalSetting state machine in 10.3.18). The default value is media-dependent; the value is specified in the respective media-dependent clauses. The initialLogSyncInterval is a per-PTP Port attribute.

The currentLogSyncInterval specifies the current value of the sync interval and is a per-PTP Port attribute.

When useMgtSettableLogSyncInterval is TRUE, currentLogSyncInterval is set equal to mgtSettableLogSyncInterval (see 14.8.20), and initialLogSyncInterval is ignored.

When the value of syncLocked is FALSE, time-synchronization messages shall be transmitted such that the value of the arithmetic mean of the intervals, in seconds, between message transmissions is within  $\pm 30\%$  of  $2^{\text{currentLogSyncInterval}}$ . In addition, a PTP Port shall transmit time-synchronization messages such that at least 90% of the inter-message intervals are within  $\pm 30\%$  of the value of  $2^{\text{currentLogSyncInterval}}$ . The interval between successive time-synchronization messages should not exceed twice the value of  $2^{\text{portDS} - \text{currentLogSyncInterval}}$  in order to prevent causing a syncReceiptTimeout event. The PortSyncSyncSend state machine (see 10.2.12) is consistent with these requirements, i.e., the requirements here and the requirements of the PortSyncSyncSend state machine can be met simultaneously.

NOTE 1—A minimum number of inter-message intervals is necessary in order to verify that a PTP Port meets these requirements. The arithmetic mean is the sum of the inter-message interval samples divided by the number of samples. For more detailed discussion of statistical analyses, see Papoulis [B25].

NOTE 2—If useMgtSettableLogSyncInterval is FALSE the value of initialLogSyncInterval is the value of the sync interval when the PTP Port is initialized. The value of the sync interval can be changed, e.g., if the PTP Port receives a Signaling message that carries a message interval request TLV (see 10.6.4.3) and the current value is stored in currentLogSyncInterval. The value of the sync interval can be reset to the initial value, e.g., by a message interval request TLV for which the value of the field logTimeSyncInterval is 126 (see 10.6.4.3.7).

### 10.7.3 Timeouts

#### *Change 10.7.3.1 as follows:*

##### 10.7.3.1 syncReceiptTimeout

The value of this attribute tells a slave port the number of sync intervals to wait without receiving synchronization information, before assuming that the master is no longer transmitting synchronization information and that the BMCA needs to be run, if appropriate. The condition of the slave port not receiving synchronization information for syncReceiptTimeout sync intervals is known as *sync receipt timeout*.

The default value shall be 3. The range shall be 2 through 255. The syncReceiptTimeout is a per-PTP Port attribute.

#### *Change 10.7.3.2 as follows:*

##### 10.7.3.2 announceReceiptTimeout

The value of this attribute tells a slave port the number of announce intervals to wait without receiving an Announce message, before assuming that the master is no longer transmitting Announce messages, and that the BMCA needs to be run, if appropriate. The condition of the slave port not receiving an Announce message for announceReceiptTimeout announce intervals is known as *announce receipt timeout*.

The default value shall be 3. The range shall be 2 through 255. The announceReceiptTimeout is a per-PTP Port attribute.

#### *Change 10.7.3.3 as follows:*

##### 10.7.3.3 gPtpCapableReceiptTimeout

The value of this attribute tells a PTP Port the number of gPTP-capable message intervals to wait without receiving from its neighbor a Signaling message containing a gPTP-capable TLV, before determining that its neighbor is no longer invoking gPTP.

NOTE—A determination that its neighbor is no longer invoking gPTP ~~will~~ causes the PTP Port to set asCapable to FALSE.

The default value shall be 9. The range shall be 2 through 255.



## 11. Media-dependent layer specification for full-duplex point-to-point links

### 11.1 Overview

#### 11.1.1 Propagation delay measurement

<<Editor's note: The resolution of comment #11 against D1.1 introduced the terms "transport-specific peer-to-peer delay mechanism" to refer to peer delay that runs only on domain 0, and "CMLDS peer-to-peer delay mechanism" to refer to peer delay that runs as part of CMLDS. The changes to 11.1.1 below introduce this terminology.>>

*Change the first paragraph of 11.1.1 as follows:*

The measurement of propagation delay on a full-duplex point-to-point PTP Link using the peer-to-peer delay mechanism is illustrated in Figure 11-1. The instance of the peer-to-peer delay mechanism that runs only on domain 0 is referred to as the transport-specific peer-to-peer delay mechanism. The instance of the peer-to-peer delay mechanism that runs as part of Common Mean Link Delay Service (CMLDS, see 11.2.17) is referred to as the CMLDS peer-to-peer delay mechanism. The CMLDS peer-to-peer delay mechanism is the same as the peer-to-peer delay mechanism described in IEEE Std 1588-2019, specialized to a two-step PTP Port<sup>1</sup> and sending the requestReceiptTimestamp and the responseOriginTimestamp separately [see item c) 8) of 11.4.2 in IEEE Std 1588-2019]. The transport-specific peer-to-peer delay mechanism is mathematically equivalent to the CMLDS peer-to-peer delay mechanism, though its computations are organized differently. The description in this subclause focuses on the transport-specific peer-to-peer delay mechanism. In addition, when it is not important in this standard to distinguish the transport-specific CMLDS peer-to-peer delay mechanisms, the mechanism is simply referred to as the "peer-to-peer delay mechanism," or more briefly as the "peer delay mechanism."

The measurement is made by each port at the end of every full-duplex point-to-point PTP Link. Thus, both ports sharing a PTP Link will independently make the measurement, and both ports will know the propagation delay as a result. This allows the time-synchronization information described in 11.1.3 to be transported irrespective of the direction taken by a Sync message. The propagation delay measurement is made on ports otherwise blocked by non-PTP algorithms (e.g., Rapid Spanning Tree Protocol) used to eliminate cyclic topologies. This enables either no loss of synchronization or faster resynchronization, after a reconfiguration, because propagation delays are already known and do not have to be initially measured when the reconfiguration occurs.

*Change the second paragraph of 11.1.1 as follows:*

Since the propagation delay measurement is made using timestamps relative to the LocalClock entities at each port at the ends of the PTP Link and the resulting mean delay is expressed in the responder timebase (see 11.2.19.3.4), there is no need to measure the mean delay for the PTP Link in each domain because the mean delay is the same in each domain. In addition, the quantity neighborRateRatio (see 10.2.5.7) is the ratio of the responder to requester LocalClock frequency and is also the same in all domains. Therefore, the propagation delay, and neighborRateRatio measurement ~~is~~ made using the peer-to-peer delay mechanism (i.e., nrrPdelay, see 11.2.13.13) are domain-independent. Single instances of the respective state machines that cause these measurements to be made are invoked, rather than one instance per domain, and the results are available to all domains. The PTP messages used for the measurements (i.e., Pdelay\_Req, Pdelay\_Resp, and Pdelay\_Resp\_Follow\_Up; see 11.4.5 through 11.4.7) carry 0 in the domainNumber field, but this value is not used.

*Change the last paragraph of 11.1.1 as follows:*

<sup>1</sup>See 3.1.86 of IEEE Std 1588-2019 for the definition of a two-step PTP Port.

The rate ratio of the responder relative to the initiator is the quantity `neighborRateRatio` (see 10.2.5.7). When computed by the peer-to-peer delay mechanism, it is computed by the function `computePdelayRateRatio()` (see 11.2.19.3.3) of the `MDPdelayReq` state machine (see 11.2.19) using successive values of  $t_3$  and  $t_4$ . As indicated in the description of `computePdelayRateRatio()`, any scheme that uses this information is acceptable as long as the performance requirements of B.2.4 are met. One example scheme is given in NOTE 1 of 11.2.19.3.3.

## 11.2 State machines for MD entity specific to full-duplex point-to-point links

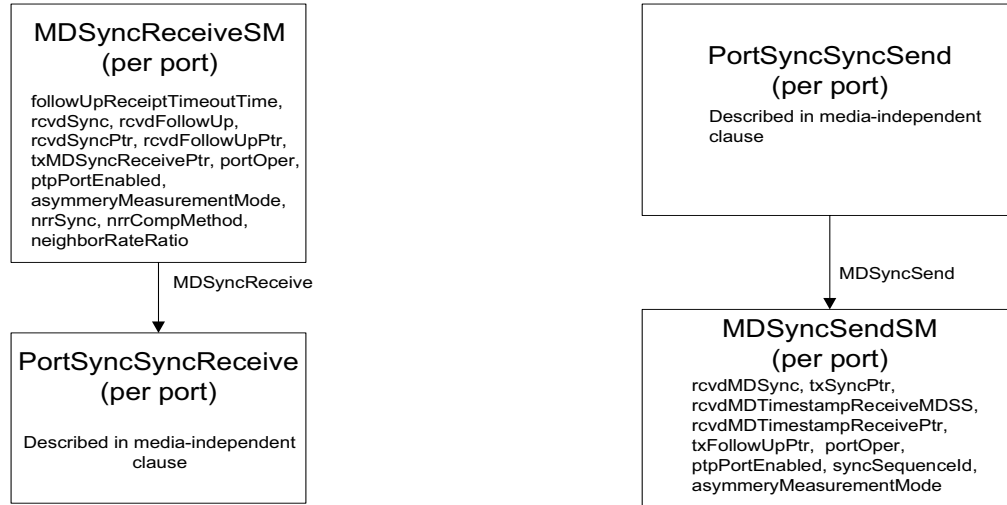
*Change 11.2.1 as follows:*

### 11.2.1 General

This subclause describes the media-dependent state machines for an MD entity, for full-duplex point-to-point links. The state machines are all per port because an instance of each is associated with an MD entity. The state machines are as follows:

- a) `MDSyncReceiveSM` (shown in Figure 10-2, and in more detail in Figure 11-4): receives `Sync` and, if the received information is two-step, `Follow_Up` messages, and sends the time-synchronization information carried in these messages to the `PortSync` entity of the same PTP Port. In addition, if the global variable `driftTrackingTlvSupport` is `TRUE`, this state machine computes the ratio of the frequency of the `LocalClock` entity in the PTP Instance at the other end of the attached PTP Link to the frequency of the `LocalClock` entity in this PTP Instance using the information contained in successive `Drift_Tracking` TLVs, and places the result in the global variable `nrrSync` (see 11.2.13.14). The global variable `neighborRateRatio` (see 10.2.5.7) is set equal to `nrrSync` or `nrrPdelay` (see 11.2.13.13) depending on whether the value of the global variable `nrrCompMethod` (see 11.2.13.15) is equal to `Sync` or `Pdelay`, respectively. There is one instance of this state machine per PTP Port, per domain.
- b) `MDSyncSendSM` (shown in Figure 10-2, and in more detail in Figure 11-4): receives an `MDSyncSend` structure from the `PortSync` entity of the same PTP Port; transmits a `Sync` message; uses the `syncEventEgressTimestamp`, corrected for `egressLatency`, and information contained in the `MDSyncSend` structure to compute information needed for the `Sync` message if the PTP Port is currently operating as a one-step PTP Port and for the corresponding `Follow_Up` message if the PTP Port is currently operating as a two-step PTP Port; and transmits the `Follow_Up` message if the PTP Port is two-step. There is one instance of this state machine per PTP Port, per domain.

*Replace Figure 11-4 with the following:*

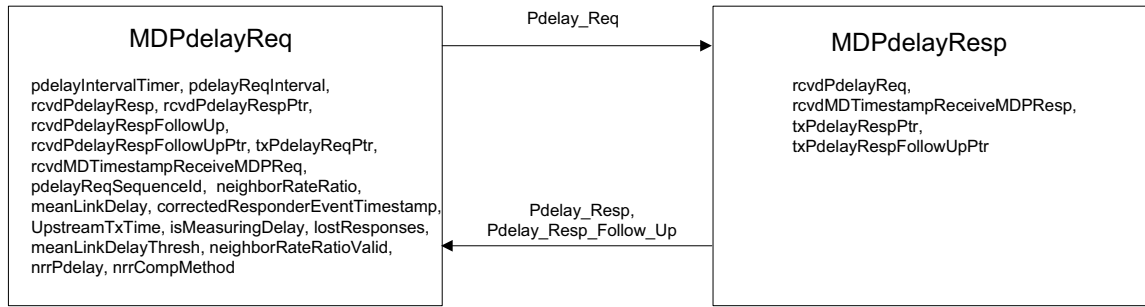


NOTE—This figure differs from the 2020 edition, with Corrigendum 1 applied, of this standard in that the the variables `nrrSync`, `nrrCompMethod`, and `neighborRateRatio` are added to the `MDSyncReceiveSM` block.

**Figure 11-4—Detail of MD entity time-synchronization state machines for full-duplex point-to-point links**

- c) `MDPdelayReq` (shown in Figure 11-5): transmits a `Pdelay_Req` message, receives `Pdelay_Resp` and `Pdelay_Resp_Follow_Up` messages corresponding to the transmitted `Pdelay_Req` message, uses the information contained in successive `Pdelay_Resp` and `Pdelay_Resp_Follow_Up` messages to compute the ratio of the frequency of the `LocalClock` entity in the PTP Instance at the other end of the attached PTP Link to the frequency of the `LocalClock` entity in this PTP Instance, places the result in the global variable `nrrPdelay` (see 11.2.13.13), and uses the information obtained from the message exchange and the computed frequency ratio to compute propagation delay on the attached PTP Link. There is one instance of this state machine for all the domains, per port.
- d) `MDPdelayResp` (shown in Figure 11-5): receives a `Pdelay_Req` message from the MD entity at the other end of the attached PTP Link, and responds with `Pdelay_Resp` and `Pdelay_Resp_Follow_Up` messages. There is one instance of this state machine for all the domains, per port.

**Replace Figure 11-5 with the following:**



NOTE—This figure differs from the 2020 edition, with Corrigendum 1 applied, of this standard in that the the variables nrrPdelay and nrrCompMethod are added to the MDPdelayReq block.

**Figure 11-5—Peer-to-peer delay mechanism state machines—  
overview and interrelationships**

- e) SyncIntervalSetting state machine (not shown): receives a Signaling message that contains a Message Interval Request TLV (see 10.6.4.3), and sets the global variables that give the duration of the mean intervals between successive Sync messages. There is one instance of this state machine per PTP Port, per domain.
- f) LinkDelayIntervalSetting state machine (not shown): receives a Signaling message that contains a Message Interval Request TLV (see 10.6.4.3), and sets the global variables that give the duration of the mean intervals between successive Pdelay\_Req messages. There is one instance of this state machine for all the domains, per port.

Figure 10-2, Figure 11-4, and Figure 11-5 are not themselves state machines, but illustrate the machines, their interrelationships, the principle variables and messages used to communicate between them, their local variables, and performance parameters. The figures do not show the service interface primitives between the media-dependent layer and the logical link control (LLC). Figure 11-5 is analogous to Figure 10-2; while Figure 10-2 applies to the general time-synchronization protocol, Figure 11-5 is limited to the peer-to-peer delay mechanism for measurement of propagation delay in full-duplex point-to-point links. Figure 11-4 shows greater detail of the MDSyncReceiveSM and MDSyncSendSM state machines for full-duplex point-to-point links than Figure 10-2.

The state machines described in 11.2 use some of the global per PTP Instance system variables specified in 10.2.4, the global per-port variables specified in 10.2.5, and the functions specified in 10.2.6.

### ***Change 11.2.2 as follows:***

#### **11.2.2 Determination of asCapable and asCapableAcrossDomains**

There is one instance of the global variable asCapable (see 10.2.5.1) per PTP Port, per domain. There is one instance of the global variable asCapableAcrossDomains (see 11.2.13.12), per port, that is common across, and accessible by, all the domains.

The per-PTP Port global variable asCapable (see 10.2.5.1) indicates whether the IEEE 802.1AS protocol is operating, in this domain, on the PTP Link attached to this PTP Port, and can provide the time-synchronization performance described in B.3. asCapable is used by the PortSync entity, which is media-independent; however, the determination of asCapable is media-dependent.

The per-port global variable `asCapableAcrossDomains` is set by the `MDPdelayReq` state machine (see 11.2.19 and Figure 11-9). For a port attached to a full-duplex point-to-point PTP Link, `asCapableAcrossDomains` shall be set to TRUE if and only if it is determined, via the transport-specific or CMLDS peer-to-peer delay mechanism, that the following conditions hold for the port:

- a) The port is exchanging peer delay messages with its neighbor,
- b) The measured delay does not exceed `meanLinkDelayThresh`,
- c) The port does not receive multiple `Pdelay_Resp` or `Pdelay_Resp_Follow_Up` messages in response to a single `Pdelay_Req` message, and
- d) The port does not receive a response from itself or another PTP Port of the same PTP Instance.

NOTE 1—If a PTP Instance implements only domain 0 and the `MDPdelayReq` and `MDPdelayResp` state machines are invoked on domain 0 (see 11.2.19), `asCapableAcrossDomains` is still set by the `MDPdelayReq` state machine. If only domain 0 is implemented, the state machines may implement the transport-specific peer-to-peer delay mechanism (see 11.2.17).

The default value of `meanLinkDelayThresh` shall be set as specified in Table 11-1.

**Table 11-1—Value of `meanLinkDelayThresh` for various links**

Link	Value of <code>meanLinkDelayThresh</code> (ns) (see NOTE)
100BASE-TX, 1000BASE-T	800 <sub>10</sub>
100BASE-FX, 1000BASE-X	FFFF FFFF FFFF FFFF FFFF FFFF <sub>16</sub>
NOTE—The actual propagation delay for 100BASE-TX and 1000BASE-T links is expected to be smaller than the above respective threshold. If the measured mean propagation delay (i.e., <code>meanLinkDelay</code> ; see 10.2.5.8) exceeds this threshold, it is assumed that this is due to the presence of equipment that does not implement gPTP. For 100BASE-FX and 1000BASE-X links, the actual propagation delay can be on the order of, or larger than, the delay produced by equipment that does not implement gPTP; therefore, such equipment cannot be detected by comparing measured propagation delay with a threshold. In this case, <code>meanLinkDelayThresh</code> is set to the largest possible value (i.e., all 1s).	

The per-PTP Port, per-domain global variable `asCapable` shall be set to TRUE if and only if the following conditions hold:

- e) The value of `asCapableAcrossDomains` is TRUE, and
- f) One of the following conditions holds:
  - 1) The value of `neighborGtpCapable` for this PTP Port is TRUE, or
  - 2) The value of `domainNumber` is zero, and the value of `sdoId` for peer delay messages received on this PTP Port is 0x100.

NOTE 2—Condition f) 2) ensures backward compatibility with the 2011 edition of this standard. A PTP Instance compliant with the current edition of this standard that is attached, via a full-duplex point-to-point PTP Link, to a PTP Instance compliant with the 2011 edition of this standard will not receive Signaling messages that contain the gPTP capable TLV and will not set `neighborGtpCapable` to TRUE. However, condition f) 2) ensures that `asCapable` for this PTP Port and domain (i.e., domain 0) will still be set to TRUE if condition e) holds because the peer delay messages received from the time-aware system compliant with the 2011 edition of this standard will have `sdoId` set to 0x100.

***Change 11.2.12 as follows:***

**11.2.12 Overview of MD entity global variables**

Subclause defines global variables used by the MD entity state machines whose scopes are as follows:

- Per PTP Instance (i.e., per domain)
- Per PTP Instance, per PTP Port
- Instances used by CMLDS (see 11.2.17) (i.e., variable is common across all ~~LinkPorts~~Link Ports)
- Instances used by CMLDS, per ~~LinkPort~~Link Port

Table 11-2 summarizes the scope of each global variable of .

**Table 11-2 — Summary of scope of global variables used by time synchronization state machines (see 10.2.4 and 10.2.5)**

Variable name	Subclause of definition	Per PTP Instance (i.e., per domain)	Per PTP Instance, per PTP Port	Instance used by CMLDS (i.e., variable is common across all <del>LinkPorts</del> Link Ports)	Instance used by CMLDS, per <del>LinkPort</del> Link Port
currentLogPdelayReqInterval	11.2.13.1	No	Yes <sup>1</sup>	No	Yes
initialLogPdelayReqInterval	11.2.13.2	No	Yes <sup>1</sup>	No	Yes
pdelayReqInterval	11.2.13.3	No	Yes <sup>1</sup>	No	Yes
allowedLostResponses	11.2.13.4	No	Yes <sup>1</sup>	No	Yes
allowedFaults	11.2.13.5	No	Yes <sup>1</sup>	No	Yes
isMeasuringDelay	11.2.13.6	No	Yes <sup>1</sup>	No	Yes
meanLinkDelayThresh	11.2.13.7	No	Yes <sup>1</sup>	No	Yes
syncSequenceId	11.2.13.8	No	Yes	No	No
oneStepReceive	11.2.13.9	No	Yes	No	No
oneStepTransmit	11.2.13.10	No	Yes	No	No
oneStepTxOper	11.2.13.11	No	Yes	No	No
asCapableAcrossDomains	11.2.13.12	No	No	No	Yes
nrrPdelay	11.2.13.13	No	Yes	No	No
nrrSync	11.2.13.14	No	Yes	No	No
nrrCompMethod	11.2.13.15	No	Yes	No	No

<sup>1</sup> The instance of this variable that is per PTP Instance, per PTP Port exists only for domain 0.

**11.2.13 MD entity global variables**

*Change 11.2.13.5 as follows:*

**11.2.13.5** allowedFaults: The number of faults [i.e., instances where meanLinkDelay (see 10.2.5.8) exceeds meanLinkDelayThresh (see 11.2.13.7) and/or the computation of ~~neighborRateRatio~~nrrPdelay (see

11.2.13.13 is invalid (see 11.2.19.2.10)] above which asCapableAcrossDomains is set to FALSE, i.e., the port is considered not capable of interoperating with its neighbor via the IEEE 802.1AS protocol (see 10.2.5.1). The data type for allowedFaults is UInteger8. The default value and range of allowedFaults is given in 11.5.4. There is one instance of this variable for all the domains (per port). The variable is accessible by all the domains.

***Add the following definitions to 11.2.13:***

**11.2.13.13 nrrPdelay:** The neighborRateRatio (see 10.2.5.7) computed using the transport-specific or CMLDS peer-to-peer delay mechanism (see 11.2.19). The data type for nrrPdelay is Float64.

**11.2.13.14 nrrSync:** The neighborRateRatio (see 10.2.5.7) computed using the Drift\_Tracking TLV carried by successive Sync or, if two-step, Follow\_Up messages (see 11.2.14). The data type for nrrSync is Float64.

**11.2.13.15 nrrCompMethod:** An Enumeration1 that takes on the values Sync and Pdelay to indicate the source of the value of neighborRateRatio (see 10.2.5.7):

- a) If the value is Sync and driftTrackingTlvSupport (see 10.2.4.27) is TRUE, neighborRateRatio is populated with the value of nrrSync.
- b) If the value is Pdelay or if driftTrackingTlvSupport (10.2.4.27) is FALSE, neighborRateRatio is populated with the value of nrrPdelay.

**11.2.14 MDSyncReceiveSM state machine**

***Add the following definition to 11.2.14.1:***

**11.2.14.1 State machine variables**

**11.2.14.1.2 rcvdSync:** A Boolean variable that ~~notifies the current state machine~~ is set to TRUE when a Sync message is received. This variable is reset to FALSE by the current state machine.

**11.2.14.1.8 TEMP:** A temporary variable used to reduce clutter in the state diagram (see Figure ). The data type for TEMP is Integer16.

**11.2.14.2 State machine functions**

***Insert 11.2.14.2.1 f), g), and h), and renumber subsequent list items as appropriate.***

- f) If the value of the global variable driftTrackingTlvSupport (see 10.2.4.27) is TRUE and the most recently received Sync message (if twoStepFlag is FALSE) or Follow\_Up message (if twoStepFlag is TRUE) contains a Drift\_Tracking TLV:
  - 1) syncGrandmasterIdentity (see 10.2.4.25) is set equal to the syncGrandmasterIdentity of the Drift\_Tracking TLV.
  - 2) If the value of syncStepsRemoved of the Drift\_Tracking TLV is not equal to 0xFFFF, syncStepsRemoved (see 10.2.4.26) is set equal to the value of syncStepsRemoved of the Drift\_Tracking TLV plus one. If the value of syncStepsRemoved of the Drift\_Tracking TLV is equal to 0xFFFF, syncStepsRemoved (see 10.2.4.26) is set equal to 0xFFFF.
  - 3) nrrSync (see 11.2.13.14) is computed using the following information conveyed by successive Sync and, if twoStepFlag is TRUE, Follow\_Up messages:
    - i) the syncEventIngressTimestamp values for the respective Sync messages.
    - ii) The correctedResponderEventTimestamp values, whose data type is UScaledNs, obtained by adding (A) the seconds field of the syncEgressTimestamp of the Drift\_Tracking TLV, multiplied by  $10^9$ , (B) the nanoseconds field of the syncEgressTimestampTimestamp of



- the Drift\_Tracking TLV, and (C) delayAsymmetry (see 10.2.5.9) if and only if this state machine is invoked by CMLDS.
- iii) the rateRatioDrift field of the Drift\_Tracking TLV.

NOTE 1—If delayAsymmetry does not change during the time interval over which nrrSync is computed, it is not necessary to add it in f3)ii)(B) because in that case it will be canceled when computing the difference between earlier and later correctedResponderEventTimestamps.

NOTE 2—The methodology in f3) for computing nrrSync follows the methodology described in 11.2.19.3.3 for computing nrrPdelay using information contained in peer delay messages. See the NOTES and description in 11.2.19.3.3 for more detail.

NOTE 3—The use of the rateRatioDrift, e.g., in compensating for frequency drift when computing rateRatio (see the RECEIVED\_SYNC state of the PortSyncSyncReceive state machine of Figure 10-4), is implementation-specific or profile standard-specific.

- g) If the value of the global variable driftTrackingTlvSupport (see 10.2.4.27) is TRUE and the most recently received Sync message (if twoStepFlag is FALSE) or Follow\_Up message (if twoStepFlag is TRUE) does not contain a Drift\_Tracking TLV:
- 1) syncGrandmasterIdentity (see 10.2.4.25) is set equal to 0xFFFF FFFF FFFF FFFF.
  - 2) syncStepsRemoved (see 10.2.4.26) is set equal to 0xFFFF.
- h) the value of neighborRateRatio (see 10.2.5.7) is set equal to the value of nrrSync (see 11.2.13.14) or nrrPdelay (see 11.2.13.13) based on the value of nrrCompMethod (see 11.2.13.15).

<<Editor's note: In D1.1, the above items f) and g) were inserted after the existing item f). In D1.2 they are inserted before the existing item f), along with the new item h), so that in the case where the Drift\_Tracking TLV is used, neighborRateRatio will be computed before it is used to update upstreamTxTime in the existing item f).>>

#### *Change 11.2.14.2.1 f) as follows:*

- f) upstreamTxTime is set equal to the syncEventIngressTimestamp for the most recently received Sync message (see 11.4.3), minus the mean propagation time on the PTP Link attached to this PTP Port (meanLinkDelay; see 10.2.5.8) divided by neighborRateRatio (see 10.2.5.7), ~~and, if and only if the state machine is invoked by the instance-specific peer-to-peer delay mechanism,~~ minus delayAsymmetry (see 10.2.5.9) for this PTP Port divided by the sum of rateRatio [see item e) in this subclause] and the quantity neighborRateRatio – 1 [see item e) in this subclause]. The syncEventIngressTimestamp is equal to the timestamp value measured relative to the timestamp measurement plane, minus any ingressLatency (see 8.4.3). The upstreamTxTime can be written as follows:

~~State machine invoked by instance-specific peer-to-peer delay mechanism:~~

$$\text{upstreamTxTime} = \text{syncEventIngressTimestamp} - (\text{meanLinkDelay}/\text{neighborRateRatio}) - (\text{delayAsymmetry}/(\text{rateRatio} + \text{neighborRateRatio} - 1))$$

~~State machine invoked by CMLDS:~~

$$\text{upstreamTxTime} = \text{syncEventIngressTimestamp} - (\text{meanLinkDelay}/\text{neighborRateRatio})$$

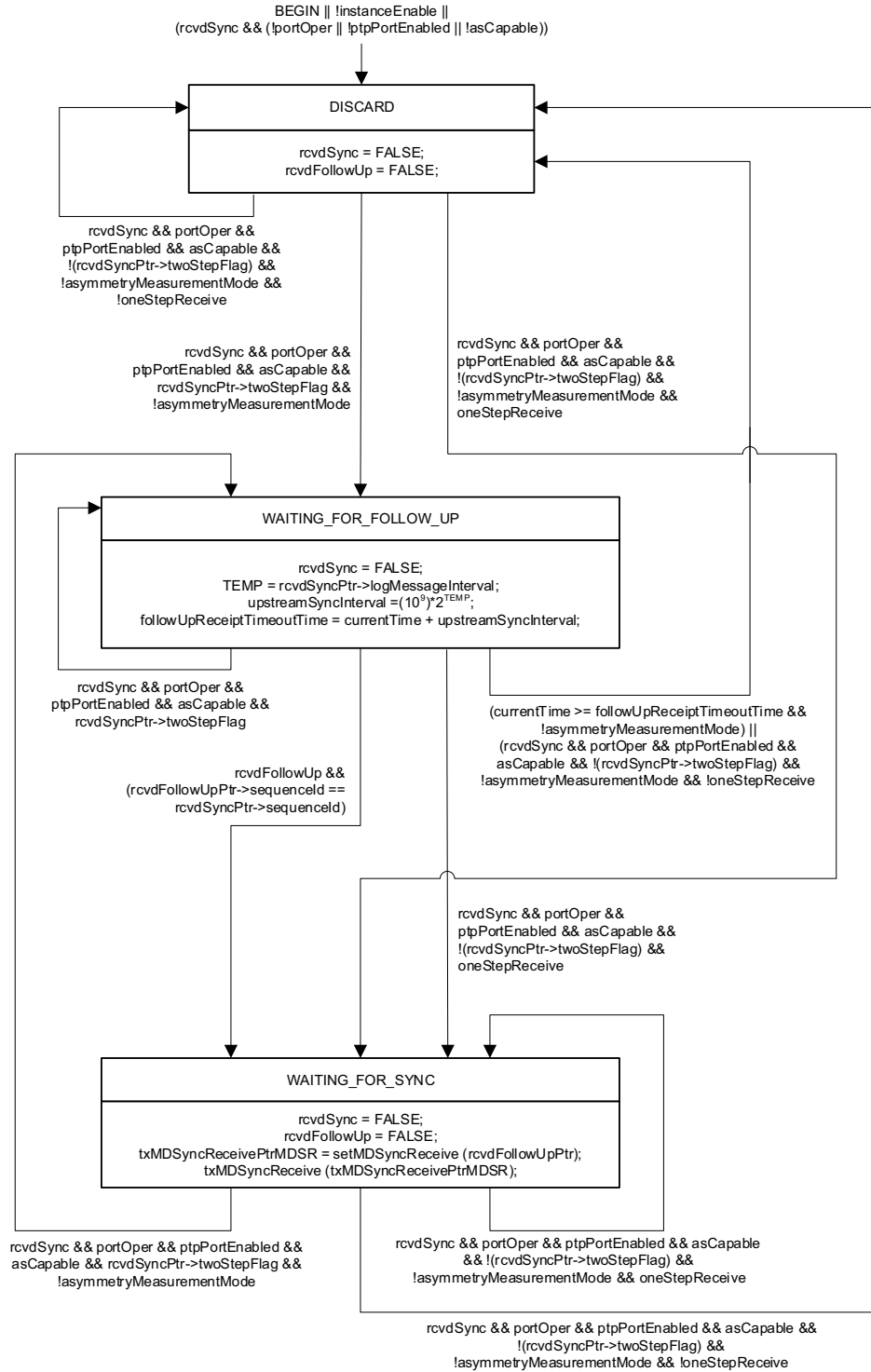
NOTE 1—The mean propagation time is divided by neighborRateRatio to convert it from the time base of the PTP Instance at the other end of the attached PTP Link to the time base of the current PTP Instance. If the ~~instance~~transport-specific peer-to-peer delay mechanism is used (i.e., portDS.delayMechanism is P2P), delayAsymmetry is divided by rateRatio to convert it from the time base of the Grandmaster Clock to the time base of the current PTP Instance. The first quotient is then subtracted from syncEventIngressTimestamp, and the second quotient is subtracted from syncEventIngressTimestamp if the ~~instance~~transport-specific peer-to-peer delay mechanism is used. The syncEventIngressTimestamp is measured relative to the time base of the current PTP Instance. See 11.2.17.2 for more detail.



NOTE 2—The difference between the mean propagation time in the Grandmaster Clock time base, the time base of the PTP Instance at the other end of the PTP Link, and the time base of the current PTP Instance is usually negligible. The same is true of any delayAsymmetry (see NOTE 2 of 11.2.19.3.4).

### 11.2.14.3 State diagram

*Replace Figure 11-6 with the following:*



**NOTE**—This figure differs from the 2020 edition, with Corrigendum 1 applied, of this standard in that the variable portEnabledOper at the entry to the DISCARD state is changed to portOper.

**Figure 11-6—MDSyncReceiveSM state machine**

## 11.2.15 MDSyncSendSM state machine

*Insert 11.2.15.2.3 g) and h), and renumber subsequent list items as appropriate.*

- g) If the value of the global variable driftTrackingTlvSupport (see 10.2.4.27) is TRUE and the PTP Instance is the Grandmaster PTP Instance, a Drift\_Tracking TLV is attached to the transmitted Follow\_Up message with the syncGrandmasterIdentity, syncStepsRemoved, and syncEgressTimestamp fields set as follows:
  - 1) syncGrandmasterIdentity is set equal to the defaultDS.clockIdentity.
  - 2) syncStepsRemoved is set equal to 0.
  - 3) syncEgressTimestamp is set equal to the syncEventEgressTimestamp of the associated Sync message.
  - 4) rateRatioDrift (see 11.4.4.4.9) is the measured estimate of the rate of change per second of the frequency of the Grandmaster Clock relative to the frequency of the local clock at the Grandmaster PTP Instance. If the Grandmaster Clock and the Local Clock at the Grandmaster PTP Instance are the same, rateRatioDrift is zero.
- h) If the value of the global variable driftTrackingTlvSupport (see 10.2.4.27) is TRUE and the PTP Instance is not the Grandmaster PTP Instance, a Drift\_Tracking TLV is attached to the transmitted Follow\_Up message with the syncGrandmasterIdentity, syncStepsRemoved, and syncEgressTimestamp fields set as follows:
  - 1) syncGrandmasterIdentity is set equal to the value of the global variable syncGrandmasterIdentity (see 10.2.4.25).
  - 2) syncStepsRemoved is set equal to the value of the global variable syncStepsRemoved (see 10.2.4.26).
  - 3) syncEgressTimestamp is set equal to the syncEventEgressTimestamp of the associated Sync message.
  - 4) rateRatioDrift (see 11.4.4.4.9) is the measured estimate of the rate of change per second of the frequency of the Grandmaster Clock relative to the frequency of the local clock of the PTP Instance that sends the message.

*Change 11.2.15.2.3 i) as follows:*

- i) ~~scaledLastGmFreqChange~~lastGmFreqChange is set equal to the ~~scaledLastGmFreqChange~~lastGmFreqChange member of the most recently received MDSyncSend structure (see 10.2.2.1 and 11.2.11), multiplied by  $2^{41}$ .

*Insert 11.2.15.2.5 g) and h), and renumber subsequent list items as appropriate.*

- g) If the value of the global variable driftTrackingTlvSupport (see 10.2.4.27) is TRUE and the PTP Instance is the Grandmaster PTP Instance, a Drift\_Tracking TLV is attached to the transmitted Sync message with the syncGrandmasterIdentity, syncStepsRemoved, and syncEgressTimestamp fields set as follows:
  - 1) syncGrandmasterIdentity is set equal to the defaultDS.clockIdentity.
  - 2) syncStepsRemoved is set equal to 0.
  - 3) syncEgressTimestamp is set equal to the syncEventEgressTimestamp of the associated Sync message.
  - 4) rateRatioDrift (see 11.4.4.4.9) is the measured estimate of the rate of change per second of the frequency of the Grandmaster Clock relative to the frequency of the local clock at the Grandmaster PTP Instance. If the Grandmaster Clock and the Local Clock at the Grandmaster PTP Instance are the same, rateRatioDrift is zero.
- h) If the value of the global variable driftTrackingTlvSupport (see 10.2.4.27) is TRUE and the PTP Instance is not the Grandmaster PTP Instance, a Drift\_Tracking TLV is attached to the transmitted

Sync message with the syncGrandmasterIdentity, syncStepsRemoved, and syncEgressTimestamp fields set as follows:

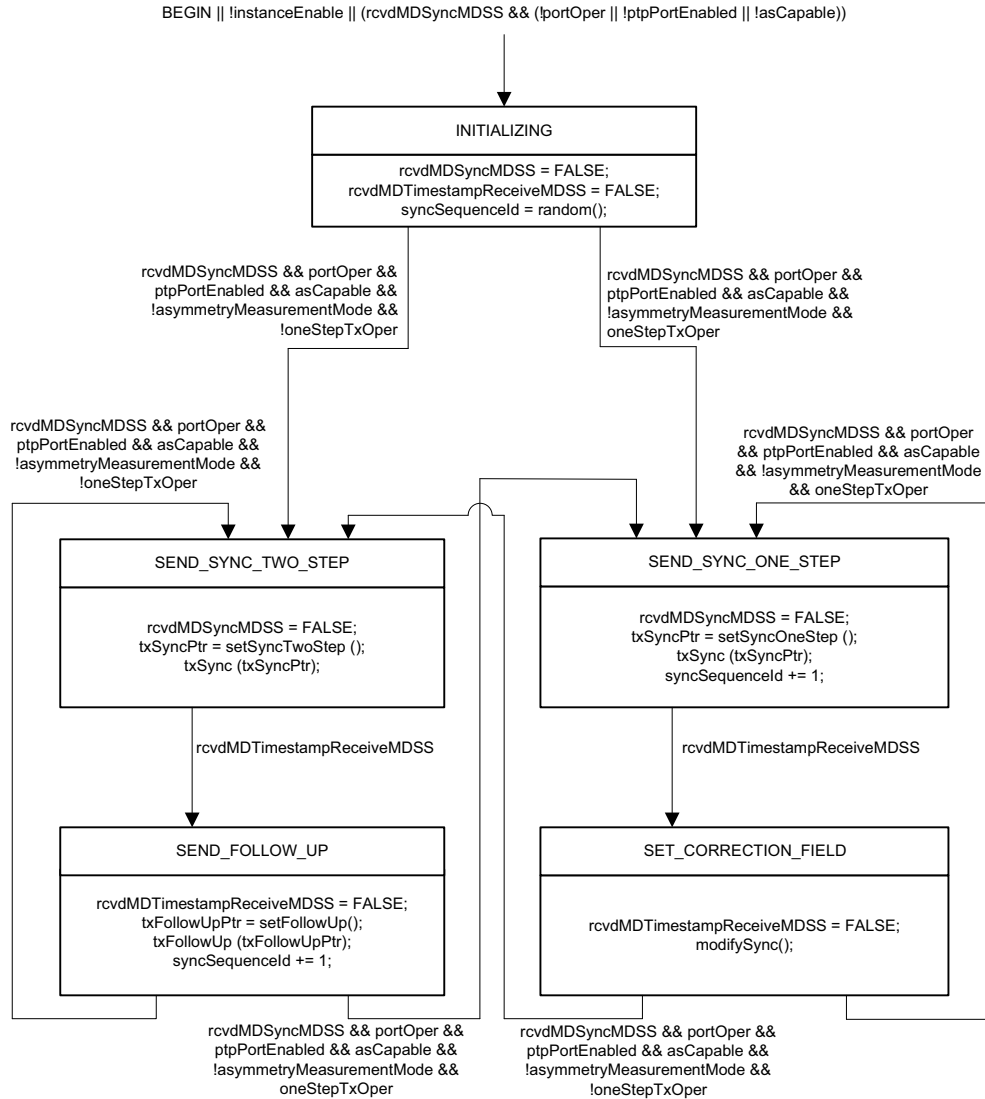
- 1) syncGrandmasterIdentity is set equal to the value of the global variable syncGrandmasterIdentity (see 10.2.4.25).
- 2) syncStepsRemoved is set equal to the value of the global variable syncStepsRemoved (see 10.2.4.26) .
- 3) syncEgressTimestamp is set equal to the syncEventEgressTimestamp of the associated Sync message.
- 4) rateRatioDrift (see 11.4.4.4.9) is the measured estimate of the rate of change per second of the frequency of the Grandmaster Clock relative to the frequency of the local clock of the PTP Instance that sends the message.

***Change 11.2.15.2.5 i) as follows:***

- i) ~~sealedLastGmFreqChange~~<sup>scaledLastGmFreqChange</sup> is set equal to the ~~sealedLastGmFreqChange~~<sup>scaledLastGmFreqChange</sup> member of the most recently received MDSyncSend structure (see 10.2.2.1 and 11.2.11), multiplied by  $2^{41}$ .

### **11.2.15.3 State diagram**

***Replace Figure 11-7 with the following:***



**NOTE**—This figure differs from the 2020 edition, with Corrigendum 1 applied, of this standard in that (a) the variable portEnabledOper is changed to portOper on the links between the SEND\_FOLLOW\_UP and SEND\_SYNC\_TWO\_STEP states, between the SEND\_FOLLOW\_UP and SEND\_SYNC\_ONE\_STEP states, and between the SET\_CORRECTION\_FIELD and SEND\_SYNC\_TWO\_STEP states; and (b) the statement ‘syncSequenceId += ‘ in the last line of the SEND\_SYNC\_TWO\_STEP state is moved to the last line of the SEND\_FOLLOW\_UP state.

**Figure 11-7—MDSyncSendSM state machine**

## 11.2.17 Common Mean Link Delay Service (CMLDS)

*Change 11.2.17.1 as follows:*

### 11.2.17.1 General

Each ~~port~~PTP Port or LinkPort of a time-aware system invokes a single instance of the MDPdelayReq state machine (see 11.2.19) and the MDPdelayResp state machine (see 11.2.20). If the time-aware system implements more than one domain, these two state machines shall provide a Common Mean Link Delay Service (CMLDS), as described in this subclause, that measures mean propagation delay on the PTP Link attached to the port and the neighbor rate ratio for the port (i.e., the ratio of the frequency of the LocalClock entity of the time-aware system at the other end of the PTP Link attached to this port, to the frequency of the LocalClock entity of this time-aware system). The CMLDS makes the mean propagation delay and neighbor rate ratio available to all active domains. If the time-aware system implements one domain, and if ~~(the~~ domainNumber of this domain is 0; ~~(see 8.1), these two state machines may also provide the CMLDS; however, if they~~the state machines do not also provide the CMLDS (i.e., if only the PTP Instance ~~transport-specific peer delay mechanism is provided), they shall be invoked on domain 0. In other words, if the domain-number~~domainNumber is not 0, portDS.delayMechanism (see Table 14-8 in 14.8.5) must not be P2P. If CMLDS is used, the LocalClock entity for CMLDS and the LocalClock entity for each PTP Instance shall be the same LocalClock.

NOTE 1—In the above ~~sentence~~, the ~~condition that~~case where the time-aware system implements only one domain implicitly assumes that IEEE 802.1AS is the only PTP profile present on the respective port of the time-aware system, i.e., no other PTP profiles are implemented on that port. If other PTP profiles that use the CMLDS are present on the port, the CMLDS must be provided.

In accordance with IEEE Std 1588-2019, the term *Link Port* refers to a port of the CMLDS. A PTP Port for which portDS.delayMechanism is COMMON\_P2P uses the CMLDS provided by the Link Port whose cmlDsLinkPortDS.portIdentity.portNumber (see 14.16.2) is equal to the commonServicesPortDS.cmlDsLinkPortPortNumber (see 14.14.2) for this PTP Port.

The value of majorSdoId for the CMLDS shall be 0x2. The value of minorSdoId for the Common Mean Link Delay Service shall be 0x00. As a result, the value of sdoId for the Common Mean Link Delay Service is 0x200.

NOTE 2—The above requirements for majorSdoId and minorSdoId are for the CMLDS. The requirements for gPTP domains, including ~~instance~~transport-specific peer delay messages, are given in 8.1.

NOTE 3—The requirements of this subclause for CMLDS are consistent with the requirements of IEEE Std 1588-2019.

If a PTP Port that invokes the CMLDS receives a Pdelay\_Req message with majorSdoId value of 0x1, minorSdoId value of 0x00, and domainNumber value of 0, the PTP Port shall respond with PTP ~~Instance~~transport-specific peer delay messages (i.e., the Pdelay\_Resp and Pdelay\_Resp\_Follow\_Up corresponding to this Pdelay\_Req) using the ~~instance~~transport-specific peer-to-peer delay mechanism. These instance-specific messages have majorSdoId value of 0x1, minorSdoId value of 0x00, and domainNumber value of 0.

NOTE 3—The above requirement ensures:

- a) Backward compatibility with time-aware systems that comply with the 2011 version of this standard if domain 0 is implemented, and
- b) Compatibility with time-aware systems that implement only one domain and invoke the MDPdelayReq and MDPdelayResp state machines on domain 0.

NOTE 4—In general, a port can receive:

- a) Peer delay messages of the CMLDS,
- b) PTP ~~Instance~~transport-specific peer delay messages of domain 0 (with sdoId of 0x100), and
- c) If there are other PTP profiles on the neighbor port that use ~~instance~~transport-specific peer delay, peer delay messages of those profiles.

The port responds to the messages of type a) if it invokes CMLDS, the messages of type b) if it invokes gPTP domain 0, and the messages of type c) if it invokes the respective other PTP profiles.

The CMLDS shall be enabled on a Link Port if the value of portDS.delayMechanism (see 14.8.5) is COMMON\_P2P for at least one PTP Port that is enabled (i.e., for which portOper and ptpPortEnabled are both TRUE) and corresponds to the same physical port as the Link Port (see 14.1). The value of cmlDsLinkPortEnabled is TRUE if the CMLDS is enabled on the Link Port and FALSE if the CMLDS is not enabled on the Link Port.

***Change 11.2.17.2 as follows:***

#### **11.2.17.2 Differences between ~~instance~~transport-specific peer-to-peer delay mechanism and CMLDS peer-to-peer delay mechanism in computations of ~~mean-link-delay~~meanLinkDelay and ~~effect-of-delayAsymmetry~~nrrPdelay**

The MDPdelayReq state machine (see 11.2.19), ~~and~~ MDPdelay\_Resp state machine (see 11.2.20), ~~and~~ ~~MDSyncReceiveSM state machine (see 11.2.14)~~ perform various computations of ~~mean-link-delay~~meanLinkDelay and ~~of the effect of delayAsymmetry~~nrrPdelay differently, depending on whether the respective computations are done using the ~~instance~~transport-specific peer-to-peer delay mechanism or using CMLDS. The resulting values of meanLinkDelay and nrrPdelay are the same for both transport-specific peer delay and CMLDS peer delay; however, the organization of the computations for transport-specific peer delay and CMLDS peer delay are different. Some of the computations are done at the Initiator and some of the computations are done at the Responder. It is necessary for the Initiator and the Responder to both perform the transport-specific peer delay computations or both perform the CMLDS peer delay computations in order to obtain the correct results for meanLinkDelay and nrrPdelay. The differences are ~~described as follows:~~

- a) Both ~~Instance~~transport-specific peer delay and CMLDS peer delay computes ~~mean-link-delay~~meanLinkDelay averaged over the two directions, and adds delayAsymmetry separately when computing upstreamTxTime by the setMDSyncReceiveMDSR() function of the MDSyncReceive state machine. However, in the computation of meanLinkDelay, CMLDS subtracts delayAsymmetry

from the correctionField when sending the pdelayReq message at the Initiator (see Figure 11-1) and adds delayAsymmetry back when computing the quantity  $t_3$  in the function computePropTime() (at the Initiator), while transport-specific peer delay does not subtract and add back delayAsymmetry, while CMLDS corrects the computed mean link delay for delayAsymmetry and therefore does not need to add it separately (see 11.2.14.2.1).

- b) ~~Instance~~Transport-specific peer delay sets the correctionField of a transmitted Pdelay\_Req message to 0, while CMLDS peer delay sets it to  $-\text{delayAsymmetry}$  (see 11.2.19.3.1).
- c) ~~Instance~~Transport-specific peer delay sets the correctionField of Pdelay\_Resp equal to the fractional nanoseconds portion of the pdelayReqEventIngressTimestamp of the corresponding Pdelay\_Req, while CMLDS peer delay sets the correctionField of Pdelay\_Resp equal to minus the fractional nanoseconds portion of the pdelayReqEventIngressTimestamp of the corresponding Pdelay\_Req (see 11.2.20.3.1).
- d) ~~Instance~~Transport-specific peer delay sets the correctionField of Pdelay\_Resp\_Follow\_Up equal to the fractional nanoseconds portion of the pdelayRespEventEgressTimestamp, while CMLDS peer delay sets the correctionField of Pdelay\_Resp\_Follow\_Up equal to the sum of the correctionField of the corresponding Pdelay\_Req and the fractional nanoseconds portion of the pdelayRespEventEgressTimestamp (see 11.2.20.3.3).
- e) When computing the quantity  $t_2$  mean link delay for the meanLinkDelay computation [i.e., the quantity  $D$  in Equation (11-5), see 11.2.19.3.4], the correctionField of the Pdelay\_Resp message, divided by  $2^{16}$ , is added ~~when computing the quantity  $t_2$  if instance~~transport-specific peer delay is used, while it is subtracted if CMLDS peer delay is used (see 11.2.19.3.4), and the quantity delayAsymmetry, divided by  $2^{16}$ , is added if CMLDS peer delay is used.
- f) When computing ~~neighborRateRatio~~nrrPdelay, the computation of the correctedResponderEventTimestamp must be corrected for delayAsymmetry if, and only if, CMLDS peer delay is used. The reason for this correction is that, with CMLDS peer delay, delayAsymmetry was subtracted from the Pdelay\_Req correctionField, and then the Pdelay\_Resp\_Follow\_Up correctionField was set equal to the sum of the Pdelay\_Req correctionField and the fractional nanoseconds portion of the PdelayRespEventEgressTimestamp, while with ~~instance~~transport-specific peer delay, the correctionField of Pdelay\_Req was set equal to 0 (see 11.2.19.3.1, 11.2.19.3.3, and 11.2.20.3.3).

The computations in this standard for the ~~instance~~transport-specific peer-to-peer delay mechanism are the same as in IEEE Std 802.1AS-2011, for backward compatibility. However, the computations in this standard for CMLDS peer delay ~~must need to~~ be consistent with IEEE Std 1588-2019 because CMLDS peer delay can be used by other PTP profiles, in addition to the PTP profile included in IEEE Std 802.1AS, that might be present in a gPTP node. Therefore, the computations for the ~~instance~~transport-specific peer-to-peer delay mechanism and CMLDS peer delay are different (i.e., are organized differently), even though they produce the same results.

## 11.2.19 MDPdelayReq state machine

### 11.2.19.2 State machine variables

*Change 11.2.19.2.10 as follows:*

**11.2.19.2.10 ~~neighborRateRatio~~nrrPdelayValid:** A Boolean variable that indicates whether the function computePdelayRateRatio() (see 11.2.19.3.3) successfully computed ~~neighborRateRatio~~nrrPdelay (see 11.2.13.13).

*Change 11.2.19.2.12 as follows:*

**11.2.19.2.12 portEnabled0:** A Boolean variable whose value is equal to ptpPortEnabled (see 10.2.5.13) if this state machine is invoked by the ~~instance~~transport-specific peer-to-peer delay mechanism and is equal to



cmldsLinkPortEnabled (see 11.2.18.1) if this state machine is invoked by the CMLDS peer-to-peer delay mechanism.

### 11.2.19.3 State machine functions

#### *Change 11.2.19.3.1 as follows:*

**11.2.19.3.1 setPdelayReq():** Creates a structure containing the parameters (see 11.4) of a Pdelay\_Req message to be transmitted, and returns a pointer, txPdelayReqPtr (see 11.2.19.2.6), to this structure. The parameters are set as follows:

- a) sourcePortIdentity is set equal to the port identity of the port corresponding to this MD entity (see 8.5.2).
- b) sequenceId is set equal to pdelayReqSequenceId (see 11.2.19.2.8).
- c) correctionField is set to
  - 1) 0 if this state machine is invoked by the ~~instance~~transport-specific peer-to-peer delay mechanism, and
  - 2) –delayAsymmetry (i.e., the negative of delayAsymmetry) if this state machine is invoked by the CMLDS.
- d) The remaining parameters are set as specified in 11.4.2 and 11.4.5.

#### *Change 11.2.19.3.3 as follows:*

**11.2.19.3.3 computePdelayRateRatio():** Computes ~~neighborRateRatio~~nrrPdelay (see 11.2.13.13) using the following information conveyed by successive Pdelay\_Resp and Pdelay\_Resp\_Follow\_Up messages:

- a) The pdelayRespEventIngressTimestamp (see 11.3.2.1) values for the respective Pdelay\_Resp messages
- b) The correctedResponderEventTimestamp values, whose data type is UScaledNs, obtained by adding the following fields of the received Pdelay\_Resp\_Follow\_Up message:
  - 1) The seconds field of the responseOriginTimestamp field, multiplied by 10<sup>9</sup>
  - 2) The nanoseconds field of the responseOriginTimestamp ~~parameter~~field
  - 3) The correctionField, divided by 2<sup>16</sup>
  - 4) delayAsymmetry (see 10.2.5.9), if and only if this state machine is invoked by CMLDS

NOTE 1—If delayAsymmetry does not change during the time interval over which ~~neighborRateRatio~~nrrPdelay is computed, it is not necessary to ~~subtract~~add it if this state machine is invoked by CMLDS because in that case it will be canceled when computing the difference between earlier and later correctedResponderEventTimestamps.

Any scheme that uses the preceding information, along with any other information conveyed by the successive Pdelay\_Resp and Pdelay\_Resp\_Follow\_Up messages, to compute ~~neighborRateRatio~~nrrPdelay is acceptable as long as the performance requirements specified in B.2.4 are met. If ~~neighborRateRatio~~nrrPdelay is successfully computed, the Boolean ~~neighborRateRatio~~nrrPdelayValid (see 11.2.19.2.10) is set to TRUE. If ~~neighborRateRatio~~nrrPdelay is not successfully computed (e.g., if the MD entity has not yet exchanged a sufficient number of peer delay messages with its peer), the Boolean ~~neighborRateRatioValid~~nrrPdelayValid is set to FALSE.

NOTE 2—As one example, ~~neighborRateRatio~~nrrPdelay can be estimated as the ratio of the elapsed time of the LocalClock entity of the time-aware system at the other end of the PTP Link attached to this port, to the elapsed time of the LocalClock entity of this time-aware system. This ratio can be computed for the time interval between a set of received Pdelay\_Resp and Pdelay\_Resp\_Follow Up messages and a second set of received Pdelay\_Resp and Pdelay\_Resp\_Follow Up messages some number of Pdelay\_Req message transmission intervals later, i.e.,

$$\frac{\text{correctedResponderEventTimestamp}_N - \text{correctedResponderEventTimestamp}_0}{\text{pdelayRespEventIngressTimestamp}_N - \text{pdelayRespEventIngressTimestamp}_0}$$

where  $N$  is the number of Pdelay\_Req message transmission intervals separating the first set of received Pdelay\_Resp and Pdelay\_Resp\_Follow\_Up messages and the second set, and the successive sets of received Pdelay\_Resp and Pdelay\_Resp\_Follow\_Up messages are indexed from 0 to  $N$  with the first set indexed 0.

NOTE 3—This function must account for non-receipt of Pdelay\_Resp and/or Pdelay\_Resp\_Follow\_Up for a Pdelay\_Req message and also for receipt of multiple Pdelay\_Resp messages within one Pdelay\_Req message transmission interval.

#### ***Change 11.2.19.3.4 as follows:***

**11.2.19.3.4 computePropTime():** Computes the mean propagation delay on the PTP Link attached to this MD entity,  $D$ , and returns this value.  $D$  is given by Equation (11-5).

$$D = \frac{r \cdot (t_4 - t_1) - (t_3 - t_2)}{2} \quad (11-5)$$

where

- $t_4$  is pdelayRespEventIngressTimestamp (see 11.3.2.1) for the Pdelay\_Resp message received in response to the Pdelay\_Req message sent by the MD entity, in nanoseconds; the pdelayRespEventIngressTimestamp is equal to the timestamp value measured relative to the timestamp measurement plane, minus any ingressLatency (see 8.4.3)
- $t_1$  is pdelayReqEventEgressTimestamp (see 11.3.2.1) for the Pdelay\_Req message sent by the P2PPort entity, in nanoseconds
- $t_2$  is the sum of (1) the ns field of the requestReceiptTimestamp, (2) the seconds field of the requestReceiptTimestamp multiplied by  $10^9$ , ~~and~~ (3) the correctionField multiplied by  $s$  (see 11.2.19.2.13) and then divided by  $2^{16}$  (i.e., the correctionField is expressed in nanoseconds plus fractional nanoseconds), of the Pdelay\_Resp message received in response to the Pdelay\_Req message sent by the MD entity, and (4) -delayAsymmetry divided by  $2^{16}$  if this state machine is invoked by CMLDS (i.e., delayAsymmetry is subtracted in the case of CMLDS).
- $t_3$  is the sum of (1) the ns field of the responseOriginTimestamp, (2) the seconds field of the responseOriginTimestamp multiplied by  $10^9$ , and (3) the correctionField divided by  $2^{16}$  (i.e., the correctionField is expressed in nanoseconds plus fractional nanoseconds), of the Pdelay\_Resp\_Follow\_Up message received in response to the Pdelay\_Req message sent by the MD entity.
- $r$  is the current value of ~~neighborRateRatio~~ nrrPdelay for this MD entity (see 10.2.5.7)

When CMLDS is used, Eq. (11-5) can be rewritten as shown in Eq. (11-6) (see NOTE 3 below) using the definitions of  $t_2$  and  $t_3$  above for the CMLDS case.

Propagation delay averaging may be performed, as described in 11.1.2 by Equation (11-2), Equation (11-3), and Equation (11-4). In this case, the successive values of propagation delay computed using Equation (11-5) are input to either Equation (11-2) or Equation (11-4), and the computed average propagation delay is returned by this function.

NOTE 1—Equation (11-5) defines  $D$  as the mean propagation delay relative to the time base of the time-aware system at the other end of the attached PTP Link. It is divided by ~~neighborRateRatio~~ nrrPdelay (see 10.2.5.7) to convert it to the time base of the current time-aware system when subtracting from syncEventIngressTimestamp to compute upstreamTxTime [see item f) in 11.2.14.2.1].

NOTE 2—The difference between mean propagation delay relative to the Grandmaster Clock time base and relative to the time base of the time-aware system at the other end of the attached PTP Link is usually negligible. To see this, note that the former can be obtained from the latter by multiplying the latter by the ratio of the Grandmaster Clock frequency to the frequency of the LocalClock entity of the time-aware system at the other end of the PTP Link attached to this port. This ratio differs from 1 by 200 ppm or less. For example, for a worst-case frequency offset of the LocalClock entity of the time-aware system at the other end of the PTP Link, relative to the Grandmaster Clock, of 200 ppm, and a measured propagation time of 100 ns, the difference in  $D$  relative to the two time bases is 20 ps.

NOTE 3—In IEEE Std 1588-2019, the computation of ~~Equation (11-5)~~ meanLinkDelay is organized differently from the organization used for instance-specific peer delay in the present standard. Using the definitions of  $t_2$  and  $t_3$  above, Equation (11-5) can be rewritten as shown in Equation (11-6).

$$D = [r \cdot (t_4 - t_1) - (\text{responseOriginTimestamp} - \text{requestReceiptTimestamp}) + (\text{correctionField of Pdelay\_Resp} - (\text{correctionField of Pdelay\_Resp\_Follow\_Up}))] / 2 \quad (11-6)$$

$$D = [r \cdot (t_4 - t_1) - (\text{responseOriginTimestamp} - \text{requestReceiptTimestamp}) + s \cdot (\text{correctionField of Pdelay\_Resp}) - (\text{correctionField of Pdelay\_Resp\_Follow\_Up}) - \text{delayAsymmetry}] / 2 \quad (11-6)$$

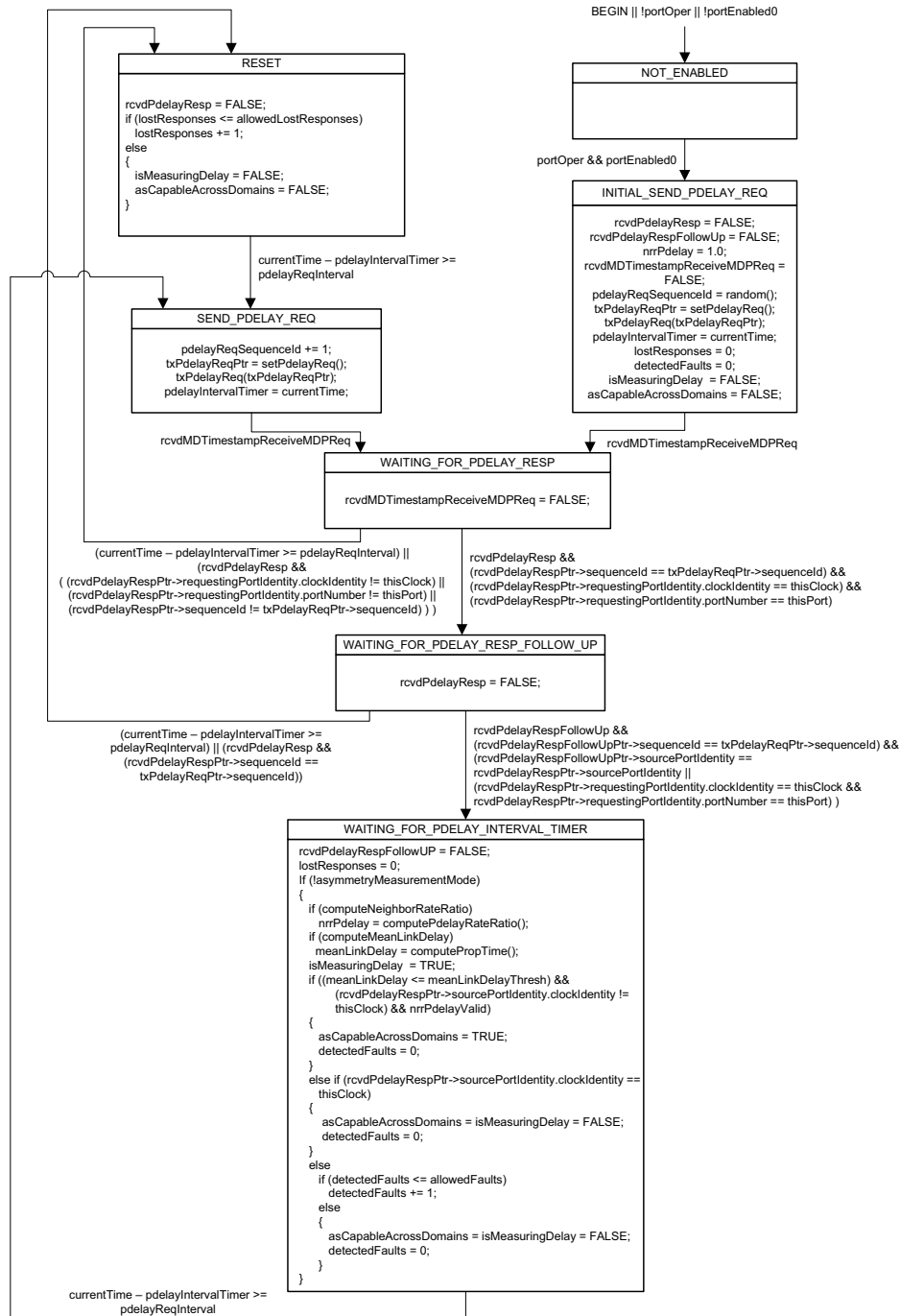
where each term is expressed in units of nanoseconds as described in the definitions of  $t_1$ ,  $t_2$ ,  $t_3$ , and  $t_4$  above. In IEEE Std 1588-2019, the fractional nanoseconds portion of  $t_2$  is subtracted from the correctionField of Pdelay\_Resp, rather than added as in ~~this the present standard for instance-specific peer delay [see 11.2.20.3.1 d)1]~~; however, the correctionField of Pdelay\_Resp is then subtracted in Equation (11-6) rather than added [in Eq. (11-6) for instance-specific peer delay, where  $s=1$ ], and the two minus signs ~~cancel each other~~ result in the same sign for the fractional nanoseconds portion of  $t_2$  and the seconds and nanoseconds portion of  $t_2$  (requestReceiptTimestamp). The computations of  $D$  in this standard and IEEE Std 1588-2019 are mathematically equivalent. The organization of the computation with CMLDS must be used in the present standard for interoperability with ~~used in~~ IEEE Std 1588-2019 (see 11.2.17.2). This organization must be used with CMLDS in the present standard, for the case where CMLDS is used, is consistent with the organization in ~~for interoperability with~~ IEEE Std 1588-2019 (see 11.2.17.2).

#### 11.2.19.4 State diagram

**Change NOTE 1 of 11.2.19.4 as follows:**

NOTE 1—The ratio of the frequency of the LocalClock entity of the time-aware system at the other end of the attached PTP Link to the frequency of the LocalClock entity of the current time-aware system, ~~neighborRateRatio~~ nrPdelay, retains its most recent value when a Pdelay\_Resp and/or Pdelay\_Resp\_Follow\_Up message is lost.

**Replace Figure 11-9 with the following:**



NOTE—This figure differs from the 2020 edition, with Corrigendum 1 applied, of this standard in that (a) the condition “else if (rcvdPdelayRespPtr->sourcePortIdentity.clockIdentity != thisClock)” in the state WAITING\_FOR\_PDELAY\_INTERVAL\_TIMER is changed to “else if (rcvdPdelayRespPtr->sourcePortIdentity.clockIdentity == thisClock)” (i.e., “!=” is changed to “==”), and (b) the variables neighborRateRatio and neighborRateRatioValid are replaced by nrrPdelay and nrrPdelayValid, respectively.

Figure 11-9—MDPdelayReq state machine

*Delete the NOTE immediately after Figure 11-9.*

## 11.2.20 MDPdelayResp state machine

### 11.2.20.2 State machine variables

*Change 11.2.20.2.5 as follows:*

**11.2.20.2.5 portEnabled1:** A Boolean variable whose value is equal to ptpPortEnabled (see 10.2.5.13) if this state machine is invoked by the ~~instance~~transport-specific peer-to-peer delay mechanism and is equal to cmlDsLinkPortEnabled (see 11.2.18.1) if this state machine is invoked by the CMLDS peer delay mechanism.

### 11.2.20.3 State machine functions

*Change 11.2.20.3.1 as follows:*

**11.2.20.3.1 setPdelayResp():** Creates a structure containing the parameters (see 11.4) of a Pdelay\_Resp message to be transmitted, and returns a pointer, txPdelayRespPtr (see 11.2.20.2.3), to this structure. The parameters are set as follows:

- a) sourcePortIdentity is set equal to the port identity of the port corresponding to this MD entity (see 8.5.2).
- b) sequenceId is set equal to the sequenceId field of the corresponding Pdelay\_Req message.
- c) requestReceiptTimestamp is set equal to the pdelayReqEventIngressTimestamp (see 11.3.2) of the corresponding Pdelay\_Req message, with any fractional nanoseconds portion truncated.
- d) correctionField is set equal to the following:
  - 1) The fractional nanoseconds portion of the pdelayReqEventIngressTimestamp of the corresponding Pdelay\_Req message if this state machine is invoked by the ~~instance~~transport-specific peer-to-peer delay mechanism and

- 2) Minus the fractional nanoseconds portion of the `pdelayReqEventIngressTimestamp` of the corresponding `Pdelay_Req` message if this state machine is invoked by CMLDS peer-to-peer delay mechanism.
- e) `requestingPortIdentity` is set equal to the `sourcePortIdentity` field of the corresponding `Pdelay_Req` message.
- f) The remaining parameters are set as specified in 11.4.2 and 11.4.6.

### ***Change 11.2.20.3.3 as follows:***

**11.2.20.3.3 `setPdelayRespFollowUp()`:** Creates a structure containing the parameters (see 11.4) of a `Pdelay_Resp_Follow_Up` message to be transmitted, and returns a pointer, `txPdelayRespFollowUpPtr` (see 11.2.20.2.4), to this structure. The parameters are set as follows:

- a) `sourcePortIdentity` is set equal to the port identity of the port corresponding to this MD entity (see 8.5.2).
- b) `sequenceId` is set equal to the `sequenceId` field of the corresponding `Pdelay_Req` message.
- c) `responseOriginTimestamp` is set equal to the `pdelayRespEventEgressTimestamp` (see 11.3.2) of the corresponding `Pdelay_Resp` message, with any fractional nanoseconds truncated.
- d) `correctionField` is set equal to the following:
  - 1) The fractional nanoseconds portion of the `pdelayRespEventEgressTimestamp` of the corresponding `Pdelay_Resp` message if this state machine is invoked by the ~~instance~~transport-specific peer-to-peer delay mechanism and
  - 2) The sum of the `correctionField` of the corresponding `Pdelay_Req` message and the fractional nanoseconds portion of the `pdelayRespEventEgressTimestamp` of the corresponding `Pdelay_Resp` message if this state machine is invoked by CMLDS peer-to-peer delay mechanism.
- e) `requestingPortIdentity` is set equal to the `sourcePortIdentity` field of the corresponding `Pdelay_Req` message.
- f) The remaining parameters are set as specified in 11.4.2 and 11.4.6.

## **11.2.21 `LinkDelayIntervalSetting` state machine**

### **11.2.21.2 State machine variables**

#### ***Change 11.2.21.2.3 as follows:***

**11.2.21.2.3 `portEnabled3`:** A Boolean variable whose value is equal to `ptpPortEnabled` (see 10.2.5.13) if this state machine is invoked by the ~~instance~~transport-specific peer-to-peer delay mechanism and is equal to `cmldsLinkPortEnabled` (see 11.2.18.1) if this state machine is invoked by the CMLDS peer-to-peer delay mechanism.

#### ***Add the following definition to 11.2.21.2:***

### **11.2.21.2 State machine variables**

**11.2.21.2.7 TEMP:** A temporary variable used to reduce clutter in the state diagram (see Figure 11-11). The data type for TEMP is `Integer16`.

### 11.2.21.3 State machine functions

*Change 11.2.21.3.2 as follows:*

**11.2.21.3.2 computeLogPdelayReqInterval (logRequestedPdelayReqInterval):** An Integer8 function that computes and returns the logPdelayReqInterval, based on the logRequestedPdelayReqInterval. This function is defined as indicated below. It is defined here so that the detailed code that it invokes does not need to be placed into the state machine diagram.

```
Integer8 computeLogPdelayReqInterval (logRequestedPdelayReqInterval)
Integer8 logRequestedPdelayReqInterval;
{
    Integer8 logSupportedPdelayReqIntervalMax,
              logSupportedClosestLongerPdelayReqInterval;
    if (isSupportedLogPdelayReqInterval (logRequestedPdelayReqInterval))
        // The requested Pdelay_Req Interval is supported and returned
        return (logRequestedPdelayReqInterval);
    else
    {
        if (logRequestedPdelayReqInterval > logSupportedPdelayReqIntervalMax)
            // Return the fastedslowest supported rate, even if faster than the
            requested rate
            return (logSupportedPdelayReqIntervalMax);
        else
            // Return the fastest supported rate that is still slower than
            // the requested rate.
            return (logSupportedClosestLongerPdelayReqInterval);
    }
}
```

*Change the first paragraph of 11.2.21.4 as follows:*

### 11.2.21.4 State diagram

The LinkDelayIntervalSetting state machine shall implement the function specified by the state diagram in Figure 11-10, the local variables specified in 11.2.21.2, the functions specified in 11.2.21.3, the messages specified in 10.6 and 11.4, the relevant global variables specified in 10.2.5, , and 11.2.18, the relevant managed objects specified in 14.8 and 14.14, and the relevant timing attributes specified in 10.7 and 11.5. This state machine is responsible for setting the global variables that give the duration of the mean interval between successive Pdelay\_Req messages and also the global variables that control whether meanLinkDelay and ~~neighborRateRatio~~nrrPdelay are computed, both at initialization and in response to the receipt of a Signaling message that contains a Message Interval Request TLV (see 10.6.4.3).

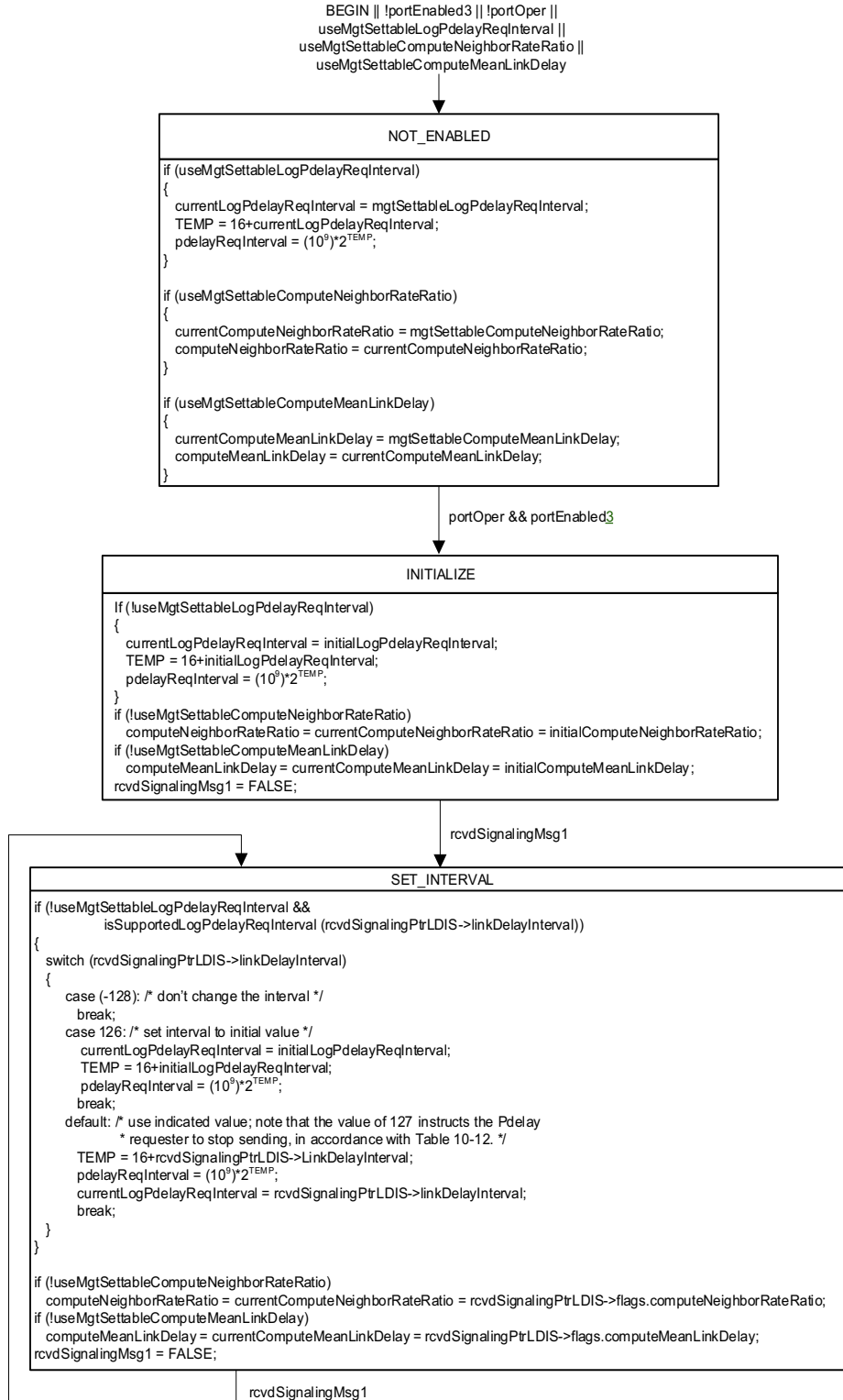


Figure 11-10—LinkDelayIntervalSetting state machine



## 11.3 Message attributes

### 11.3.6 Subtype

The subtype ~~of~~<sup>for</sup> the Sync, Follow\_Up, Pdelay\_Req, Pdelay\_Resp, and Pdelay\_Resp\_Follow\_Up messages is indicated by the majorSdoId field (see 10.6.2.2.1).

~~NOTE—The subtype for all PTP messages is indicated by the majorSdoId field.~~

## 11.4 Message formats

### 11.4.2 Header

#### 11.4.2.6 correctionField (Integer64)

*Change Table 11-6 as follows:*

**Table 11-6—Value of correctionField**

Message type	Value
Follow_Up, Sync (sent by a one-step PTP Port; see 11.1.3 and 11.2.13.9)	Corrections for fractional nanoseconds (see 10.2.9 and Figure 10-5), difference between preciseOriginTimestamp field (if sent by a two-step PTP Port) or originTimestamp field (if sent by a one-step PTP Port) and current synchronized time (see 11.2.15.2.3 and Figure ), and asymmetry <del>corrections</del> (see 8.3, 11.2.14.2.1, and 11.2.15.2.3; <del>the</del> <sup>the</sup> quantity delayAsymmetry is used in the computation of upstreamTxTime in 11.2.14.2.1, and upstreamTxTime is used in computing an addition to the correctionField in 11.2.15.2.3).
Pdelay_Resp, Pdelay_Resp_Follow_Up	Corrections for fractional nanoseconds (see Figure 11-9 and Figure 11-10) if the message is sent by the <del>instance</del> <sup>transport</sup> -specific peer-to-peer delay mechanism; or  For Pdelay_Resp, minus the corrections for fractional nanoseconds (see 11.2.20.3.1, Figure 11-9, and Figure 11-10), and for Pdelay_Resp_Follow_Up, the sum of the correctionField of the corresponding Pdelay_Req message and the fractional nanoseconds portion of the pdelayRespEventEgressTimestamp of the corresponding Pdelay_Resp message, if this state machine is invoked by CMLDS <u>peer-to-peer delay mechanism</u> .
Sync (sent by a two-step PTP Port), Pdelay_Req, Announce, Signaling	The value is 0 (see 10.6.2.2.9) if the message is sent by the <del>instance</del> <sup>transport</sup> -specific peer-to-peer delay mechanism, or  The value is 0 for Sync (sent by a two-step PTP Port), Announce, and Signaling, and <del>delayAsymmetry</del> for Pdelay_Req (see 11.2.19.3.1), if the message is sent by CMLDS <u>peer-to-peer delay mechanism</u> .
NOTE—IEEE Std 1588-2019 describes asymmetry corrections for the Pdelay_Req and Pdelay_Resp messages. However, the peer-to-peer delay mechanism computes the mean propagation delay. Here, where the gPTP communication path is a full-duplex point-to-point PTP Link, these corrections cancel in the mean propagation delay computation and therefore are not needed.	

### 11.4.3 Sync message

*Change 11.4.3.1 as follows:*

#### 11.4.3.1 General Sync message specifications

If the twoStep flag of the PTP common header (see Table 10-9) of the Sync message is TRUE, the fields of the Sync message shall be as specified in Table 11-8. If the twoStep flag of the PTP common header of the Sync message is FALSE, the fields of the Sync message shall be as specified in Table 11-9 and 11.4.3.2. Carrying the Drift\_Tracking TLV is optional.

**Table 11-8—Sync message fields if twoStep flag is TRUE**

Bits								Octets	Offset
7	6	5	4	3	2	1	0		
header (see 11.4.2)								34	0
reserved								10	34

**Table 11-9—Sync message fields if twoStep flag is FALSE**

Bits								Octets	Offset
7	6	5	4	3	2	1	0		
header (see 11.4.2)								34	0
originTimestamp								10	34
Follow_Up information TLV								32	44
<u>Drift_Tracking TLV (optional)</u>								<u>30</u>	<u>76</u>

#### 11.4.3.2 Sync message field specifications if twoStep flag is FALSE

*Insert 11.4.3.2.2, and renumber subsequent subclauses as appropriate:*

##### 11.4.3.2.2 Drift\_Tracking TLV

The Sync message may carry the Drift\_Tracking TLV, defined in 11.4.4.4.

### 11.4.4 Follow\_Up message

*Change 11.4.4.1 as follows:*

#### 11.4.4.1 General Follow\_Up message specifications

The fields of the Follow\_Up message shall be as specified in Table 11-10 and 11.4.4.2. Carrying the Drift\_Tracking TLV is optional.

Table 11-10—Follow\_Up message fields

Bits								Octets	Offset
7	6	5	4	3	2	1	0		
header (see 11.4.2)								34	0
preciseOriginTimestamp								10	34
Follow_Up information TLV								32	44
Drift_Tracking TLV (optional)								30	76

11.4.4.2 Follow\_Up message field specifications

Insert 11.4.4.2.3, and renumber subsequent subclauses as appropriate:

11.4.4.2.3 Drift\_Tracking TLV

The Follow\_Up message may carry the Drift\_Tracking TLV, defined in 11.4.4.4.

Insert 11.4.4.4, and renumber subsequent subclauses and as appropriate:

11.4.4.4 Drift\_Tracking TLV definition

11.4.4.4.1 General

The fields of the Drift\_Tracking TLV shall be as specified in Table 11-12 and in 11.4.4.4.2 through 11.4.4.4.9. This TLV is a standard organization extension TLV for the Sync or Follow\_Up message, as specified in 14.3 of IEEE Std 1588-2019.

Insert Table 11-12, and renumber subsequent tables as appropriate:

Table 11-12—Drift\_Tracking TLV

Bits								Octets	Offset
7	6	5	4	3	2	1	0		
tlvType								2	0
lengthField								2	2
organizationId								3	4
organizationSubType								3	7
syncEgressTimestamp								12	10
syncGrandmasterIdentity								8	22
syncStepsRemoved								2	30
rateRatioDrift								4	32

**11.4.4.4.2 tlvType (Enumeration16)**

The value of the tlvType is 0x3.

NOTE—This value indicates the TLV is a vendor and standard-organization extension TLV, as specified in 14.3.2.1 and Table 52 of IEEE Std 1588-2019. The tlvType is specified in that standard as ORGANIZATION\_EXTENSION with a value of 0x3.

**11.4.4.4.3 lengthField (UInteger16)**

The value of lengthField is 26.

**11.4.4.4.4 organizationId (Octet3)**

The value of organizationId is 00-80-C2.

**11.4.4.4.5 organizationSubType (Enumeration24)**

The value of organizationSubType is 6.

**11.4.4.4.6 syncEgressTimestamp (ExtendedTimestamp)**

The value of the syncEgressTimestamp field is the timestamp, based on the local clock, when the Sync message was sent by that PTP Relay Instance~~is the seconds and nanoseconds portion of the associated Sync message~~ (see 11.4.3.2).

**11.4.4.4.7 syncGrandmasterIdentity (ClockIdentity)**

The value is the value of the clockIdentity component of the rootSystemIdentity of the gmPriorityVector (see 10.3.5) of the PTP Instance that transmits the Sync message.

**11.4.4.4.8 syncStepsRemoved (UInteger16)**

The value is the value of syncMasterStepsRemoved (see 10.3.9.3) for the PTP Instance that transmits the Sync message.

**11.4.4.4.9 rateRatioDrift (Integer32)**

The value of rateRatioDrift is equal to  $(RR_{drift} - 1.0) \times (2^{41})$ , truncated to the next smaller signed integer, where  $RR_{drift}$  is the measured estimate of the rate of change per second of the ratio of the frequency of the Grandmaster Clock to the frequency of the Local Clock entity in the PTP Instance that sends the message.

NOTE—The above scaling allows the representation of rates of change of fractional frequency offset in the range  $[-(2^{-10} - 2^{-41}), 2^{-10} - 2^{-41}] \text{ s}^{-1}$ , with granularity of  $2^{-41}$ . This range is approximately  $[-9.766 \times 10^{-4}, 9.766 \times 10^{-4}] \text{ s}^{-1}$ .

**11.5 Protocol timing characterization****11.5.2 Message transmission intervals****11.5.2.2 Pdelay\_Req message transmission interval**

*Change 11.5.2.2 as follows:*

When useMgtSettableLogPdelayReqInterval (see 14.16.12) is FALSE, the initialLogPdelayReqInterval specifies the following:

- a) The mean time interval between successive Pdelay\_Req messages sent over a PTP Link when the port is initialized, and
- b) The value to which the mean time interval between successive Pdelay\_Req messages is set when a message interval request TLV is received with the logLinkDelayIntervalField set to 126 (see 11.2.21).

The currentLogPdelayReqInterval specifies the current value of the mean time interval between successive Pdelay\_Req messages. The default value of initialLogPdelayReqInterval is 0. Every port supports the value 127; the port does not send Pdelay\_Req messages when currentLogPdelayReqInterval has this value (see 11.2.21). A port may support other values, except for the reserved values indicated in Table 10-15. A port shall ignore requests for unsupported values (see 11.2.21). The initialLogPdelayReqInterval and currentLogPdelayReqInterval are per-port attributes.

When useMgtSettableLogPdelayReqInterval is TRUE, currentLogPdelayReqInterval is set equal to mgtSettableLogPdelayReqInterval (see 14.16.13), and initialLogPdelayReqInterval is ignored.

Pdelay\_Req messages shall be transmitted such that the value of the arithmetic mean of the intervals, in seconds, between Pdelay\_Req message transmissions is not less than the value of  $0.9 \times 2^{\text{currentLogPdelayReqInterval}}$ .

NOTE 1—A minimum number of inter-message intervals is necessary to verify that a PTP Port meets these requirements. The arithmetic mean is the sum of the inter-message interval samples divided by the number of samples. For more detailed discussion of statistical analyses, see Papoulis [B25].

NOTE 2—If useMgtSettableLogPdelayReqInterval is FALSE, the value of initialLogPdelayReqInterval is the value of the mean time interval between successive Pdelay\_Req messages when the port is initialized. The value of the mean time interval between successive Pdelay\_Req messages can be changed, e.g., if the port receives a Signaling message that carries a message interval request TLV (see 10.6.4.3) and the current value is stored in currentLogPdelayReqInterval. The value of the mean time interval between successive Pdelay\_Req messages can be reset to the initial value, e.g., by a message interval request TLV for which the value of the field logLinkDelayInterval is 126 (see 10.6.4.3.6).

NOTE 3—A port that requests (using a Signaling message that contains a message interval request TLV; see 10.6.4 and 11.2.21) that the port at the other end of the attached PTP Link set its currentLogPdelayReqInterval to a specific value can determine if the request was honored by examining the logMessageInterval field of subsequent received Pdelay\_Req messages.

NOTE 4—The MDPdelayReq state machine ensures that the times between transmission of successive Pdelay\_Req messages, in seconds, are not smaller than  $2^{\text{currentLogPdelayReqInterval}}$ . This is consistent with the requirements of the current subclause and IEEE Std 1588-2019, which requires that the ~~logarithm to the base 2 of the mean~~ value of the arithmetic mean of the intervals, in seconds, between Pdelay\_Req message transmissions is not less than the value of  $0.9 \times 2^{\text{currentLogPdelayReqInterval}}$  ~~smaller than the interval computed from the value of the portDS.logMinPdelayReqInterval member of the data set of the transmitting PTP Instance.~~ The sending of Pdelay\_Req messages is governed by the LocalClock and not the synchronized time (i.e., the estimate of the Grandmaster Clock time). Since the LocalClock frequency can be slightly larger than the Grandmaster Clock frequency (e.g., by 100 ppm, which is the specified frequency accuracy of the LocalClock; see B.1.1), it is possible for the time intervals between successive Pdelay\_Req messages to be slightly less than  $2^{\text{currentLogPdelayReqInterval}}$  when measured relative to the synchronized time. The factor 0.9 allow a margin of 10% for the arithmetic mean of the successive intervals between Pdelay\_Req messages.

***Change 11.5.4 as follows:***

#### 11.5.4 allowedFaults

The variable allowedFaults (see 11.2.13.5) is the number of faults above which asCapableAcrossDomains is set to FALSE, i.e., the port is considered not capable of interoperating with its neighbor via the IEEE 802.1AS protocol (see 10.2.5.1). In this context, the term *faults* refers to instances where

- a) The computed mean propagation delay, i.e., meanLinkDelay (see 10.2.5.8), exceeds the threshold, meanLinkDelayThresh (see 11.2.13.7) and/or
- b) The computation of ~~neighborRateRatio~~nrrPdelay is invalid (see 11.2.19.2.10).

The default value of allowedFaults shall be 9. The range shall be 1 through 255.

NOTE—The above description of allowedFaults uses the variable asCapableAcrossDomains (see 11.2.13.12). When only one domain is active, asCapableAcrossDomains is equivalent to the variable asCapable.

#### *Change 11.6 as follows:*

#### 11.6 Control of computation of neighborRateRatio

If the variable driftTrackingTlvSupport (see 10.2.4.27) is TRUE and nrrCompMethod (see 11.2.13.15) is equal to Pdelay, the value of neighborRateRatio (see 10.2.5.7) is set equal to nrrSync (see 11.2.13.13).

If the variable driftTrackingTlvSupport (see 10.2.4.27) is TRUE and nrrCompMethod (see 11.2.13.15) is equal to Pdelay, the value of neighborRateRatio (see 10.2.5.7) is set equal to nrrPdelay (see 11.2.13.13). In this case, the computation of nrrPdelay is controlled as described below. The description below applies only to the computation of nrrPdelay, and not to the computation of nrrSync.

The variable computeNeighborRateRatio (see 10.2.5.10) indicates whether ~~neighborRateRatio~~nrrPdelay is to be computed by this port when the peer-to-peer delay mechanism is invoked.

When useMgtSettableComputeNeighborRateRatio (see 14.16.16) is FALSE, computeNeighborRateRatio is initialized to the value of initialComputeNeighborRateRatio.

The currentComputeNeighborRateRatio specifies the current value of computeNeighborRateRatio. The default value of initialComputeNeighborRateRatio is TRUE. The initialComputeNeighborRateRatio and currentComputeNeighborRateRatio are per-port attributes.

When useMgtSettableComputeNeighborRateRatio is TRUE, currentComputeNeighborRateRatio is set equal to mgtSettableComputeNeighborRateRatio (see 14.16.17), and initialComputeNeighborRateRatio is ignored.

NOTE—If useMgtSettableComputeNeighborRateRatio is FALSE, the value of initialComputeNeighborRateRatio determines whether ~~neighborRateRatio~~nrrPdelay is computed by the peer delay mechanism when the port is initialized. The value of computeNeighborRateRatio can be changed, e.g., if the port receives a Signaling message that carries a message interval request TLV (see 10.6.4.3) and the current value is stored in currentComputeNeighborRateRatio.

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## 13. Media-dependent layer specification for interface to IEEE 802.3 Ethernet passive optical network link

### 13.3 Message format

#### 13.3.1 TIMESYNC message

##### 13.3.1.2 TIMESYNC message field specifications

*Change 13.3.1.2.15 as follows:*

##### 13.3.1.2.15 domainNumber (UInteger8)

This field is specified as the gPTP ~~domain-number~~domainNumber (see 8.1).



## 14. Timing and synchronization management

### 14.1 General

*Change 14.1.1 as follows:*

#### 14.1.1 Data set hierarchy

This clause defines the set of managed objects, and their functionality, that allow administrative configuration of clock parameters and timing and synchronization protocols.

Management data models typically represent data for the physical device (i.e., time-aware system). The specifications for discovery, management address, and security for the physical device are typically covered by standards of the management mechanism, which are outside the scope of this standard. For the management information model of this standard, the scope of work is the data contained within a time-aware system. From a management perspective, the time-aware system contains a list of one or more PTP Instances. Each entry in the list is a set of managed data sets for the respective PTP Instance.

Conformance for each managed object is optional. This standard operates correctly using default values; therefore, management is not essential. Since the management mechanism is not limited to remote protocols (e.g., SNMP, NETCONF), management can use a local mechanism with a simple interface (e.g., DIP switches). Therefore, each product can determine the support of managed objects as appropriate for its management mechanism.

The following hierarchy summarizes the managed data sets within a gPTP Node:

- a) instanceList[]
  - 1) defaultDS
  - 2) currentDS
  - 3) parentDS
  - 4) timePropertiesDS
  - 5) pathTraceDS
  - 6) acceptableMasterTableDS
  - 7) ptpInstanceSyncDS
  - 8) driftTrackingDS
  - 9) portList[]
    - i) portDS
    - ii) descriptionPortDS
    - iii) portStatisticsDS
    - iv) acceptableMasterPortDS
    - v) externalPortConfigurationPortDS
    - vi) asymmetryMeasurementModeDS
    - vii) commonServicesPortDS
- b) commonServices
  - 1) commonMeanLinkDelayService
    - i) cmlDsDefaultDS
    - ii) cmlDsLinkPortList[]
      - cmlDsLinkPortDS
      - cmlDsLinkPortStatisticsDS
      - cmlDsAsymmetryMeasurementModeDS

- 2) hotStandbyService
  - i) hotStandbySystemList[]
    - hotStandbySystemDS
    - hotStandbySystemDescriptionDS
- 3) Future common services can follow.

The instanceList is indexed using a number that is unique per PTP Instance within the time-aware system, applicable to the management context only (i.e., not used in PTP messages). The domainNumber of the PTP Instance must not be used as the index to instanceList since it is possible for a time-aware system to contain multiple PTP Instances using the same domainNumber. The portList is indexed using a number that is unique per logical port (i.e., PTP Port) in the PTP Instance (see 8.5.1). Since the portNumber of a logical port can have any value in the range 1, 2, 3, ..., 0xFFFFE (see 8.5.2.3), the portList index and portNumber values for a logical port ~~will~~are not necessarily ~~be~~ the same. PTP Instances and logical ports may be created or deleted dynamically in implementations that support dynamic create/delete of devices. Unless otherwise indicated, the data sets and managed objects under the instanceList[] are maintained separately for each PTP Instance supported by the time-aware system.

Following the instanceList[] and all the data sets of each instanceList[] member is an overall structure for common services. That structure contains one sub-structure for each common service. At present there ~~is~~are ~~only one~~two common services, namely the Common Mean Link Delay Service (CMLDS) and the Hot Standby Service. ~~, and the~~The corresponding sub-structures ~~is~~are the commonMeanLinkDelayService structure and the hotStandbyService structure, respectively. The item “~~f~~Future common services can follow” is a placeholder for any common services that might be defined in the future. The commonMeanLinkDelayService structure and the hotStandbyService structure contains the data sets and lists that are needed by the Common Mean Link Delay Service and the Hot Standby Service, respectively.

The commonMeanLinkDelayService structure contains the cmlldsLinkPortList, which is a list of CMLDS logical ports, i.e., Link Ports (see 11.2.17), of the time-aware system that ~~will~~runs the common service. The CMLDS must be implemented (i.e., a CMLDS executable must be present) on every physical port for which there is a PTP Port of a PTP Instance that can use the CMLDS (i.e., where portDS.delayMechanism of that PTP Instance can have the value COMMON\_P2P). Therefore, the cmlldsLinkPortList[] must include Link Ports that correspond to all such physical ports. As is the case for the portList of a PTP Instance, the cmlldsLinkPortList is indexed using a number that is unique per Link Port that invokes the CMLDS (see 8.5.1). Since the portNumber of a logical port (i.e., PTP Port or CMLDS Link Port) can have any value in the range 1, 2, 3, ..., 0xFFFFE (see 8.5.2.3), the cmlldsLinkPortList index and cmlldsLinkPortDS.portIdentity.portNumber values for a logical port of the Common Mean Link Delay Service ~~will~~are not necessarily ~~be~~ the same. CMLDS Link Ports may be created or deleted dynamically in implementations that support dynamic create/delete of devices.

The Common Mean Link Delay Service Data Sets are not maintained separately for each PTP Instance. Rather, a single copy of the commonServices.cmlldsDefaultDS is maintained for the time-aware system, and a single copy of each data set under the cmlldsLinkPortList[] is maintained per Link Port of the time-aware system.

A PTP Instance can use the commonServicesPortDS to determine which Link Port it must use when it obtains information provided by the Common Mean Link Delay Service (see 14.14).

The hotStandbyService structure contains the hotStandbySystemList, which is a list of instances of the Hot Standby Service. Since a Hot Standby Service Instance is associated with two PTP Instances, but a time-aware system can have more than two PTP Instances, there can, in general, be multiple instances of the Hot Standby Service (i.e., one Hot Standby Service Instance associated with each set of two distinct PTP Instances). A hotStandbySystemDS and a hotStandbySystemDescriptionDS are maintained for each instance of the Hot Standby Service.

Each instance of the Hot Standby Service (i.e., each hotStandbySystemList member) is associated with two PTP Instances, i.e., the primary PTP Instance and the secondary PTP Instance (see 18.1 and 18.2). The indices of these PTP Instances in the instanceList are contained in the members primaryPtpInstanceIndex and secondaryPtpInstanceIndex, respectively, of the hotStandbySystemDS for this instance of the Hot Standby Service.

NOTE—This hierarchy is intended to support a wide variety of time-aware system implementations. Examples include the following:

- a) A time-aware system containing four PTP Relay Instances, each of which use the same physical ports, but different domainNumber values.
- b) A time-aware system containing four PTP Relay Instances and two hotStandbySystem entities, with two PTP Instances associated with one of the hotStandbySystem entities and the other two PTP Instances associated with the other hotStandbySystem entity.
- bc) A time-aware system that represents a chassis with slots for switch/router cards, where each switch/router card is represented as a PTP Instance using distinct physical ports and all PTP Instances can use the same domainNumber.

#### 14.1.2 Data set descriptions

*Insert 14.1.2 g), and renumber the list as appropriate:*

- g) The PTP Instance Synchronization Parameter Data Set (ptpInstanceSyncDS in 14.1.1; see Table 14-6), which represents time synchronization status information for the PTP Instance.

*Insert 14.1.2 h), and renumber the list as appropriate.*

- h) The Drift Tracking Parameter Data Set (driftTrackingDS in 14.1.1; see Table xxx), which contains a managed object used to enable or disable the Drift\_Tracking TLV.

*Insert 14.1.2 r), and renumber the list as appropriate:*

- r) The Hot Standby Service Hot Standby System Parameter Data Set (hotStandbySystemDS in 14.1.1; see Table 14-20), which describes the attributes of the respective instance of the Hot Standby Service.

*Insert 14.1.2 s), and renumber the list as appropriate:*

- s) The Hot Standby Service Hot Standby System Description Parameter Data Set (hotStandbySystemDescriptionDS, see Table 14-21), which represents descriptive information for the instance of the Hot Standby Service.

#### 14.2 Default Parameter Data Set (defaultDS)

*Change 14.2.16 as follows:*

##### 14.2.16 domainNumber

The value is the ~~domain number~~ domainNumber of the gPTP domain for this instance of gPTP supported by the time-aware system (see 8.1).

NOTE—The PTP Instance for which domainNumber is 0 has constraints applied to it, e.g., timescale (see 8.2.1).

#### 14.3 Current Parameter Data Set (currentDS)

*Change 14.3.3 as follows:*

**14.3.3 offsetFromMaster**

The value is an implementation-specific or profile standard-specific representation of the current value of the time difference between a slave and the Grandmaster Clock, as computed by the slave, and as specified in 10.2.10. The value is computed by an algorithm that is implementation-specific or profile standard-specific. The data type shall be TimeInterval. The default value is implementation specific.

NOTE—For example, the inputs to this implementation-specific algorithm could be the successive values of clockSourcePhaseOffset (see 10.2.4.7) of the ClockMasterSyncOffset state machine (see 10.2.10 and Figure 10-6).

**14.4 Parent Parameter Data Set (parentDS)**

*Insert 14.4.8, and renumber subsequent subclauses as appropriate:*

**14.4.8 gmPresent**

The value is the value of the per PTP Instance global variable gmPresent (see 10.2.4.13).

*Change Table 14-3 as follows:*

**Table 14-3—parentDS table**

Name	Data type	Operations supported <sup>1</sup>	References
parentPortIdentity	PortIdentity (see 6.4.3.7)	R	14.4.2
cumulativeRateRatio	Integer32	R	14.4.3
grandmasterIdentity	ClockIdentity	R	14.4.4
grandmasterClockQuality.clock Class	UInteger8	R	14.4.5.2; 7.6.2.5 of IEEE Std 1588-2019
grandmasterClockQuality.clock Accuracy	Enumeration8	R	14.4.5.3; 7.6.2.6 of IEEE Std 1588-2019
grandmasterClockQuality.offset ScaledLogVariance	UInteger16	R	14.4.5.4
grandmasterPriority1	UInteger8	R	14.4.6
grandmasterPriority2	UInteger8	R	14.4.7
<u>gmPresent</u>	<u>Boolean</u>	<u>R</u>	<u>14.4.8</u>

<sup>1</sup> R = Read only access; RW = Read/write access.

*Insert 14.8 and Table 14-7, and renumber subsequent subclauses and tables as appropriate:*

## 14.8 PTP Instance Synchronization Parameter Data Set (ptpInstanceSyncDS)

### 14.8.1 General

The ptpInstanceSyncDS describes the synchronization status of the PTP Instance. The ptpInstanceSyncDS shall be implemented if the optional hot standby feature (see Clause 18) is implemented and may be implemented otherwise.

### 14.8.2 isSynced

The value of the global variable isSynced (see 18.4.1.1).

### 14.8.3 offsetFromMasterMax

The value is the threshold for offsetFromMaster (see 18.4.1.2), below which the PTP Instance is considered to be synchronized. For values less than or equal to zero, the PTP Instance is considered synchronized if and only if offsetFromMaster is zero. For values greater than zero, the PTP Instance is considered synchronized if and only if the equation

$$-\text{offsetFromMasterThreshold} \leq \text{offsetFromMaster} \leq \text{offsetFromMasterThreshold}$$

is satisfied.

### 14.8.4 rxSyncCountSlavePThresh

The value of rxSyncCountSlavePThresh is the threshold for rxSyncCountSlaveP (see 18.4.1.4), above which the PTP Instance is considered to be synchronized.

### 14.8.5 offsetMaxExceededCountThresh

The value of offsetMaxExceededCountThresh (see 18.4.1.7) is the threshold for the number of consecutive exceedances of offsetFromMasterMax (see 17.5.1.3) by offsetFromMaster (see 17.5.1.2), at which the PTP Instance is no longer in the SYNCED state (see 17.5.3.3).

### 14.8.6 offsetMaxMetCountThresh

The value of offsetMaxMetCountThresh (18.4.1.9) is the threshold for the number of consecutive occurrences of offsetFromMaster (see 17.5.1.2) being within offsetFromMasterMax (see 17.5.1.3), at which the PTP Instance is placed in the SYNCED state (see 17.5.3.3).

### 14.8.7 ptpInstanceSyncDS table

There is one ptpInstanceSyncDS table per PTP Instance, as detailed in Table 14-6.

*Insert 14.9 and Table Table 14-7, and renumber subsequent subclauses as appropriate.*

## 14.9 Drift Tracking Parameter Data Set (driftTrackingDS)

### 14.9.1 General

The driftTrackingDS contains a managed object that is used to enable or disable the optional Drift\_Tracking TLV.

**Table 14-6—ptpInstanceSyncDS table**

Name	Data type	Operations supported <sup>1</sup>	References
isSynced	Boolean	R	14.8.2
offsetFromMasterMax	TimeInterval	RW	14.8.3
rxSyncCountSlavePThresh	UInteger32	RW	14.8.4
offsetMaxExceededCountThresh	UInteger32	RW	14.8.5
offsetMaxMetCountThresh	UInteger32	RW	14.8.6

<sup>1</sup> R = Read only access; RW = Read/write access.

#### 14.9.2 driftTrackingTlvSupport

The value of driftTrackingTlvSupport indicates whether the Drift\_Tracking TLV is enabled or disabled. If the value is TRUE, the TLV is enabled, i.e., the global variable driftTrackingTlvSupport (see 10.2.4.27) is set to TRUE. If the value is FALSE, the TLV is disabled, i.e., the global variable driftTrackingTlvSupport is set to FALSE.

#### 14.9.3 driftTrackingDS table

There is one driftTrackingDS table per PTP Instance, as detailed in Table 14-6.

**Table 14-7—driftTrackingDS table**

Name	Data type	Operations supported <sup>1</sup>	References
driftTrackingTlvSupport	Boolean	RW	14.9.2

<sup>1</sup> R = Read only access; RW = Read/write access.

### 14.8 Port Parameter Data Set (portDS)

*Change 14.8.5 as follows:*

#### 14.8.5 delayMechanism

The value of delayMechanism indicates the mechanism for measuring mean propagation delay and neighbor rate ratio on the link attached to this PTP Port and is taken from the enumeration in ~~Table 14-8~~ Table 14-8. If the ~~domain-number~~ domainNumber is not 0, portDS.delay mechanism must not be P2P (see 11.2.17).

**Table 14-8—delayMechanism enumeration**

Delay mechanism	Value	Specification
P2P	02	The PTP Port uses the <a href="#">transport-specific</a> peer-to-peer delay mechanism
COMMON_P2P	03	The PTP Port uses the CMLDS <a href="#">peer-to-peer delay mechanism</a>
SPECIAL	04	The PTP Port uses a transport that has a native time transfer mechanism and, therefore, does not use the peer-to-peer delay mechanism (e.g., IEEE 802.11, IEEE 802.3 EPON)
	All other values reserved	
NOTE—The enumeration values are consistent with Table 21 in IEEE Std 1588-2019.		

*Insert 14.8.55, and renumber subsequent subclauses as appropriate:*

#### **14.8.55 gptpCapableStateMachinesEnabled**

A Boolean that is used to enable or disable the GptpCapableTransmit, GptpCapableReceive, and GptpCapableIntervalSetting state machines. If the value is TRUE, the GptpCapableTransmit and GptpCapableReceive state machines are enabled, and the GptpCapableIntervalSetting state machine is enabled if it is implemented. If the value is FALSE, the GptpCapableTransmit and GptpCapableReceive state machines are disabled, and the GptpCapableIntervalSetting state machine is disabled if it is implemented.. The default value is TRUE.

#### **14.8.56 nrrSync**

The value of the global variable nrrSync (see 11.2.13.14).

The value is an estimate of the ratio of the frequency of the LocalClock entity of the time-aware system at the other end of the link attached to this Link Port, to the frequency of the LocalClock entity of this time-aware system (see 10.2.5.7). nrrSync is expressed as the fractional frequency offset stored in the global variable nrrSync (see 11.2.13.13) multiplied by  $2^{41}$ , i.e., the quantity  $(\text{nrrSync} - 1.0)(2^{41})$ .

#### **14.8.57 nrrCompMethod**

The value of the global variable nrrCompMethod (see 11.2.13.15).

#### **14.8.55 portDS table**

*Insert the following items after the final item of Table 14-10:*

.

#### **14.10 Port Parameter Statistics Data Set (portStatisticsDS)**

*Change the heading of 14.10.9 as follows:*

**Table 14-10—portDS table**

Name	Data type	Operations supported <sup>1</sup>	References
gtpCapableStateMachinesEnabled	Boolean	RW	14.8.55
nrrSync	Integer32	R	14.8.56
nrrCompMethod	Enumeration2	RW	14.8.57

<sup>1</sup> R = Read only access; RW = Read/write access.

#### 14.10.9 rxP~~TP~~<sup>tp</sup>PacketDiscardCount

*Insert 14.10.20, and renumber subsequent subclauses as appropriate:*

#### 14.10.20 rxSyncCountSlaveP

This counter increments ~~every~~<sup>whenever</sup> time synchronization information is received on ~~the~~<sup>a</sup> PTP Port ~~whose~~<sup>when its</sup> port state is SlavePort. The receipt of time synchronization information is denoted by one of the following events:

- A transition to TRUE from FALSE of the rcvdSync variable of the MDSyncReceiveSM state machine (see 11.2.14.1.2 and Figure 11-6) when in the DISCARD, WAITING\_FOR\_SYNC, or WAITING\_FOR\_FOLLOW\_UP states; or
- rcvdIndication transitions to TRUE (see Figure 12-7).

This counter ~~is initialized to zero, and~~ resets to zero when the port state is not SlavePort.

*Change Table 14-12 as follows:*

**Table 14-12—portStatisticsDS table**

Name	Data type	Operations supported <sup>1</sup>	References
rxSyncCount	UInteger32	R	14.10.2
rxOneStepSyncCount	UInteger32	R	14.10.3
rxFollowUpCount	UInteger32	R	14.10.4
rxPdelayRequestCount	UInteger32	R	14.10.5
rxPdelayResponseCount	UInteger32	R	14.10.6
rxPdelayResponseFollowUpCount	UInteger32	R	14.10.7
rxAnnounceCount	UInteger32	R	14.10.8
rxP <del>TP</del> <sup>tp</sup> PacketDiscardCount	UInteger32	R	14.10.9
syncReceiptTimeoutCount	UInteger32	R	14.10.10
announceReceiptTimeoutCount	UInteger32	R	14.10.11
pdelayAllowedLostResponsesExceededCount	UInteger32	R	14.10.12



**Table 14-12—portStatisticsDS table**

Name	Data type	Operations supported <sup>1</sup>	References
txSyncCount	UInteger32	R	14.10.13
txOneStepSyncCount	UInteger32	R	14.10.14
txFollowUpCount	UInteger32	R	14.10.15
txPdelayRequestCount	UInteger32	R	14.10.16
txPdelayResponseCount	UInteger32	R	14.10.17
txPdelayResponseFollowUpCount	UInteger32	R	14.10.18
txAnnounceCount	UInteger32	R	14.10.19
<u>rxSyncCountSlaveP</u>	<u>UInteger32</u>	<u>R</u>	<u>14.10.20</u>

<sup>1</sup>R= Read only access.

### 14.13 Asymmetry Measurement Mode Parameter Data Set (asymmetryMeasurementModeDS)

*Change 14.13.3 as follows:*

#### 14.13.3 asymmetryMeasurementModeDS table

There is one asymmetryMeasurementModeDS table for the single PTP Instance whose ~~domain number~~domainNumber is 0, per PTP Port, as detailed in Table 14-15. This data set is used only when there is a single gPTP domain and CMLDS is not used.

**Table 14-15—asymmetryMeasurementModeDS table**

Name	Data type	Operations supported <sup>1</sup>	References
asymmetryMeasurementMode	Boolean	RW	14.13.2

<sup>1</sup> R = Read only access; RW = Read/write access.

### 14.16 Common Mean Link Delay Service Link Port Parameter Data Set (cmldsLinkPortDS)

*Change 14.16.9 as follows:*

#### 14.16.9 ~~neighborRateRatio~~nrrPdelay

The value is an estimate of the ratio of the frequency of the LocalClock entity of the time-aware system at the other end of the link attached to this Link Port, to the frequency of the LocalClock entity of this time-aware system (see 10.2.5.7). ~~neighborRateRatio~~nrrPdelay is expressed as the fractional frequency offset stored in the global variable nrrPdelay (see 11.2.13.13) multiplied by 2<sup>41</sup>, i.e., the quantity (~~neighborRateRatio~~nrrPdelay – 1.0)(2<sup>41</sup>).

**14.16.27 cmldsLinkPortDS table***Change the eighth row of 14-18 (not including the header row) as follows:***Table 14-18—cmldsLinkPortDS table**

Name	Data type	Operations supported <sup>1</sup>	References
<del>neighborRateRatio</del> nrrPdelay	Integer32	R	14.16.9

<sup>1</sup> R = Read only access; RW = Read/write access.**14.17 Common Mean Link Delay Service Link Port Parameter Statistics Data Set (cmldsLinkPortStatisticsDS)***Change the heading of 14.17.5 as follows***14.17.5 rxP~~TP~~tpPacketDiscardCount***Change Table 14-19 as follows:.***Table 14-19—cmldsLinkPortStatisticsDS table**

Name	Data type	Operations supported <sup>1</sup>	References
rxPdelayRequestCount	UInteger32	R	14.17.2
rxPdelayResponseCount	UInteger32	R	14.17.3
rxPdelayResponseFollowUpCount	UInteger32	R	14.17.4
rxP <del>TP</del> tpPacketDiscardCount	UInteger32	R	14.17.5
pdelayAllowedLostResponsesExceededCount	UInteger32	R	14.17.6
txPdelayRequestCount	UInteger32	R	14.17.7
txPdelayResponseCount	UInteger32	R	14.17.8
txPdelayResponseFollowUpCount	UInteger32	R	14.17.9

<sup>1</sup> R= Read only access.*Insert 14.19, and renumber any subsequent subclauses as appropriate:*

**14.19 Hot Standby System Parameter Data Set (hotStandbySystemDS)****14.19.1 General**

The hotStandbySystemDS describes the attributes of the respective instance of the Hot Standby Service.

**14.19.2 primaryPtpInstanceIndex**

The value of primaryPtpInstanceIndex is the index (see 14.1.1) of the primary PTP Instance associated with this hotStandbySystem instance.

**14.19.3 secondaryPtpInstanceIndex.**

The value of secondaryPtpInstanceIndex is the index (see 14.1.1) of the secondary PTP Instance associated with this hotStandbySystem instance.

**14.19.4 hotStandbySystemEnable**

The value is the hotStandbySystemEnable attribute of the HotStandbySystem entity (see 18.5.1.2).

**14.19.5 hotStandbySystemState**

The value of hotStandbySystemState is the state of the hotStandbySystem, i.e., the value of the global variable hotStandbySystemState (see 18.5.1.1).

**14.19.6 hotStandbySystemSplitFunctionality**

If the value is TRUE, the optional split functionality (see 18.5.3.4) is used. If the value is FALSE, the optional split functionality is not used.

**14.19.7 primarySecondaryOffset**

The absolute value of the difference between the clockSlaveTimes (see 10.2.4.3) of the primary and secondary PTP Instances.

**14.19.8 primarySecondaryOffsetThresh**

The threshold for hotStandbySystemDS.primarySecondaryOffset (see 14.19.7), above which the hotStandbySystemState transitions from REDUNDANT to NOT\_REDUNDANT, or does not transition from NOT\_REDUNDANT or OUT\_OF\_SYNC to REDUNDANT even if other conditions for these transitions are satisfied.

**14.19.9 hotStandbySystemLogSyncTimeThresh**

The value of hotStandbySystemLogSyncTimeThresh is the logarithm to base 2 of the time interval, in seconds, after which the hotStandbySystem transitions from the OUT\_OF\_SYNC state to either the NOT\_REDUNDANT or REDUNDANT state, or from the NOT\_REDUNDANT to the REDUNDANT state, if all other conditions for the respective transition are met. The value -128 means that the transition time is zero, i.e., the transition occurs immediately.

<<Editor's note: The following corrections were made in the above paragraph, relative to the resolutions of comments #32 and #97 of D1.1: (a) The sentence indicating that the value -128 means that the transition time is zero, which was previously present in 14.9.6 and 14.9.7 of D1.1 (these were deleted) is retained, (b) the statement on the data type was removed, as the style used in Clause 14 is

to have the data type in the respective table (it is in Table 14-20 below), and (c) the data type is Integer8, not UScaledNs (this was a typo, as the resolution referred this to the variable notRedundantTransitionTime (the only occurrence of notRedundantTransitionTime was in 17.6.2.4 of D1.1, and this was changed to transitionTime by the resolution of comment #100.>>

#### 14.19.10 hotStandbySystemDS table

There is one hotStandbyDS table per time-aware system, as detailed in Table 14-20.

**Table 14-20—hotStandbySystemDS table**

Name	Data type	Operations supported <sup>1</sup>	References
primaryPtpInstanceIndex	UInteger32	RW	14.19.2
secondaryPtpInstanceIndex	UInteger32	RW	14.19.3
hotStandbySystemEnable	Boolean	RW	14.19.4
hotStandbySystemState	Enumeration8	R	14.19.5
hotStandbySystemSplitFunctionality	Boolean	RW	14.19.6
primarySecondaryOffset	ScaledNs	R	14.19.7
primarySecondaryOffsetThresh	ScaledNs	RW	14.19.8
hotStandbySystemLogSyncTimeThresh	Integer8	RW	14.19.9

<sup>1</sup> R = Read only access; RW = Read/write access.

*Insert 14.20, and renumber any subsequent subclauses as appropriate:*

1       **14.20 Hot Standby System Description Parameter Data Set**  
2       **(hotStandbySystemDescriptionDS)**

3  
4       **14.20.1 General**

5  
6       The hotStandbySystemDescriptionDS contains descriptive information for the respective instance of the Hot  
7       Standby Service.

8  
9       **14.20.2 userDescription**

10  
11       The user description is a character string whose maximum length is 128.

12  
13       **14.20.3 hotStandbySystemDescriptionDS table**

14  
15       There is one hotStandbySystemDescriptionDS table per hotStandbySystem instance, as detailed in Table 14-  
16       21.

17  
18  
19               **Table 14-21—hotStandbySystemDescriptionDS table**

20

Name	Data type	Operations supported <sup>1</sup>	References
userDescription	Octet128	RW	14.20.2

21  
22  
23  
24  
25       <sup>1</sup> RW= Read/write access.

26  
27  
28       <Editor’s Note: Contributions are requested to supply members of this data set.>>

## 16. Media-dependent layer specification for CSN

### 16.4 Path delay measurement over a CSN backbone

#### 16.4.3 Path delay measurement between CSN nodes

##### 16.4.3.2 Path delay measurement without network clock reference

*Change the first paragraph of 16.4.3.2 as follows:*

Each CSN node has a free-running local clock. The path delay measurement uses the peer-to-peer delay mechanism protocol, messages, and state machines described in Clause 11 for full-duplex point-to-point links, as illustrated by Figure 16-5. The criteria of 11.2.17 for determining whether the peer-to-peer delay mechanism is ~~instance-the transport-specific~~ peer-to-peer delay mechanism or ~~is provided by the~~ CMLDS peer-to-peer delay mechanism apply here.

*Change 16.4.3.3 as follows:*

##### 16.4.3.3 Intrinsic CSN path delay measurement

Some CSN technologies feature a native mechanism that provides a path delay measurement with accuracy similar to the accuracy the peer delay protocol provides. For these CSNs, the path delay can be provided using the native measurement method rather than using the Pdelay protocol defined in 11.2.19 and 11.2.20. Such a situation is described in more detail as follows. The CSN MD entity populates the following per-PTP Port and MD-entity global variables (described respectively in 10.2.5 and 11.2.13) as indicated:

- asCapable (10.2.5.1) is set to TRUE.
- neighborRateRatio (10.2.5.7) is set to the value provided by the native CSN measurement.
- meanLinkDelay (10.2.5.8) is set to the value provided by the native CSN measurement.
- computeNeighborRateRatio (10.2.5.10) is set to FALSE.
- computeMeanLinkDelay (10.2.5.11) is set to FALSE.
- isMeasuringDelay (11.2.13.6) is set to TRUE to indicate that the CSN MD entity is measuring path delay (in this case, using its internal mechanism).
- domainNumber (8.1) is set to the ~~domain number~~ domainNumber of this gPTP domain.

## 16.5 Synchronization messages

### 16.5.3 Synchronization message propagation on a CSN with network reference clock

#### 16.5.3.2 CSN ingress node

##### 16.5.3.2.2 CSN TLV

*Change 16.5.3.2.2.10 as follows:*

##### 16.5.3.2.2.10 domainNumber (UInteger8)

This parameter is the ~~domain number~~ domainNumber of this gPTP domain.

## 17. YANG Data Model

<<Editor's note: Clause 17, and the included YANG tree and YANG module, must be updated to reflect changes and additions to managed objects due to comments against draft D1.1 (some, but not all, of these changes and additions are included here.>

### 17.1 YANG framework

Clause 14 specifies the information model for management of this standard. The data model for a specific management mechanism is derived from the information model. Since YANG-based protocols are an example of a management mechanism, the YANG data model of this clause is derived from Clause 14.

NOTE 1 - The MIB modules specified in Clause 15 were also derived from Clause 14. Consequently, the capabilities and structure of the YANG data models are aligned with those represented by the MIB. However the YANG data model has not been derived from the MIB, and there has been no attempt to include data or modeling constructs that might appear in the MIB but not in the information model.

<<Editor's note: It has been the understanding of the editor that the 802.1AS-2020 MIB is not intended to be updated to reflect the new material in 802.1ASdm. If this is correct, the above NOTE 1 needs to be updated to reflect this, as in this case the capabilities of the YANG data models related to the new material of 802.1ASdm are not represented by the MIB. However, the editor recalls discussion at one point that indicated that the MIB might actually be updated to also reflect the 802.1ASdm material. In this case, the above NOTE 1 need not be changed; also, in this case maintenance item 280, related to the MIB, must be included.>>

The information model in Clause 14 is organized as a hierarchy of data sets. Each data set contains one or more related members (items of data that can be read or written). In the context of YANG, each data set is represented as a YANG “container”, and each member is represented as a YANG “leaf”.

*Change 17.1.1 as follows:*

#### 17.1.1 Relationship to the IEEE Std 1588 data model

The YANG data models specified in this standard are based on, and augment, those specified in IEEE Std 1588. In particular the `ieee802-dot1as-ptp.yang` module imports the `ieee1588-ptp` module as a whole, augmenting that module as necessary to meet the requirements of this standard. In addition, the `ieee802-dot1as-hs.yang` module imports the `ieee1588-ptp` and `ieee802-dot1as-ptp` modules as a whole, augmenting those modules as necessary to meet the requirements of this standard. ~~This~~<sup>These</sup> imports makes existing and new IEEE Std 1588 YANG capabilities not specifically addressed by the present standard available to its implementors without delay, without the need to revise or amend IEEE Std 802.1AS.

Some of the data sets in Clause 14 (e.g., `defaultDS`) are derived from IEEE Std 1588, and some of the data sets are unique to IEEE Std 802.1AS (i.e., not derived from IEEE Std 1588). For each data set in Clause 14 that is derived from IEEE Std 1588, a portion of the members are derived from IEEE Std 1588, and the remaining members are unique to IEEE Std 802.1AS. For the members that are derived from IEEE Std 1588, the specifications in both standards are analogous (i.e., same name, data type, semantics, etc).

The YANG data model for IEEE Std 1588-2019 is published as amendment IEEE Std 1588e. The YANG module of IEEE Std 1588e (`ieee1588-ptp.yang`) contains the hierarchy (tree) of data sets and their members.

The YANG modules of this clause (`ieee802-dot1as-ptp.yang` and `ieee802-dot1as-hs.yang`) ~~use~~<sup>uses</sup> the YANG “import” statement to import the YANG module of IEEE Std 1588e. This effectively uses the IEEE Std 1588

YANG tree as the foundation of the IEEE Std 802.1AS YANG tree. By importing the tree and its data set containers, all members from Clause 14 that are derived from IEEE Std 1588 are also imported.

The core of the YANG modules for IEEE Std 802.1AS consists of YANG “augment” statements, used to add members to the tree that are unique for IEEE Std 802.1AS.

NOTE 2 - IETF RFC8575 [B48] is the standard YANG data model for IEEE Std 1588-2008. The YANG data model of IEEE Std 1588e is effectively a newer version of RFC8575. Therefore, the YANG module of RFC8575 is not imported by the YANG module of this clause.

*Change 17.2 as follows:*

## 17.2 IEEE 802.1AS YANG models

This clause uses a UML-like representation to provide an overview of the hierarchy of the IEEE Std 802.1AS YANG data model.

A representation of the management model is provided in Figures 17-1 through 17-4. The purpose of the diagram is to express the model design in a concise manner. The structure of the representation shows the name of the object followed by a list of properties for the object. The properties indicate their type and accessibility. It should be noted that the representation is meant to express simplified semantics for the properties. It is not meant to provide the specific datatype used to encode the object in either MIB or YANG. In the representation, a box with a white background represents information that comes from sources outside of this IEEE standard. A box with a gray background represents objects that are defined by this IEEE standard.

NOTE 1 - OMG UML 2.5 [B49] conventions together with C++ language constructs are used in this clause as a representation to convey model structure and relationships.

NOTE 2 - This standard specifies YANG for Clause 14 of this standard. There are optional features in the YANG module of IEEE Std 1588 that are not specified in Clause 14, and therefore not shown in the figures of this subclause. If optional IEEE Std 1588 YANG features are implemented, conformance is specified by IEEE Std 1588.

For all figures, Clause 14 data that is imported from the `ieee1588-tp.yang` module is shown in white, and Clause 14 data in augments of `ieee802-dot1as-tp.yang` is shown in gray.

Figure 17-1 provides an overview of the IEEE Std 802.1AS YANG tree. The top level instance-list provides the list of one or more PTP Instances, each with data sets. For each PTP Instance, `port-ds-list` provides the list of one or more PTP Ports, each with data sets. The common-services apply to all PTP Instances, including the Common Mean Link Delay Service (`cmlds`).

Figure 17-2 provides detail for the data sets of each PTP Instance, including each data set member.

Figure 17-3 provides detail for the data sets of each PTP Port, including each data set member.

NOTE 2 - 14.8.4 specifies `ptpPortEnabled` (`ptp-port-enabled`), which is provided in YANG as the semantically equivalent node in `ieee1588-tp` named `port-enable` (in `port-ds` of Figure 17-3). 14.8.15 specifies `mgtSettableLogAnnounceInterval` (`mgt-settable-log-announce-interval`), which is provided in YANG as the semantically equivalent node in `ieee1588-tp` named `log-announce-interval` (in `port-ds` of Figure 17-3). 14.8.20 specifies `mgtSettableLogSyncInterval` (`mgt-settable-log-sync-interval`), which is provided in YANG as the semantically equivalent node in `ieee1588-tp` named `log-sync-interval` (in `port-ds` of Figure 17-3).



Figure 17-4 provides detail for the common services, including each data set member. The Common Mean Link Delay Service (cmls) has a data sets for the service itself (e.g., default-ds), and data sets for each PTP Link Port. The Hot Standby Service has data sets for each HotStandbySystem.

NOTE 3 - 14.16.9 specifies neighborRateRatio (neighbor-rate-ratio), which is provided in YANG as the semantically equivalent node in ieee1588-ptp named scaled-neighbor-rate-ratio (in link-port-ds of Figure 17-4).

***Replace Figure 17-1 with the following:***

***Replace Figure 17-2 with the following:***

***Replace Figure 17-3 with the following:***

***Replace Figure 17-4 with the following:***

## 17.3 Structure of YANG models

The YANG model specified by this standard uses the YANG modules summarized in Table 17-1.<sup>1</sup>

In the YANG module definitions, if any discrepancy between the “description” text and the corresponding definition in any other part of this standard occur, the definitions outside this clause (Clause 17) take precedence.

***Change Table 17-1 as follows.:***

## 17.5 YANG schema tree definitions

***Insert 17.5.2 and renumber subsequent subclauses as appropriate:***

### 17.5.2 Tree diagram for ieee802-dot1as-ptp.yang

```
module: ieee802-dot1as-hs

  augment /ptp:ptp/ptp:instances/ptp:instance:
    +--rw ptp-instance-sync-ds
    |   +--ro ptp-instance-state?                ptp-instance-state-type
    |   +--rw offset-from-master-threshold?      ptp:time-interval
```

<sup>1</sup>An amendment’s designation is often used to refer to functionality in an IEEE standard after the amendment has been incorporated in a revision of the standard, even if the functionality has been revised. The amendment that added each YANG module is identified to help locate the relevant provisions of this standard.

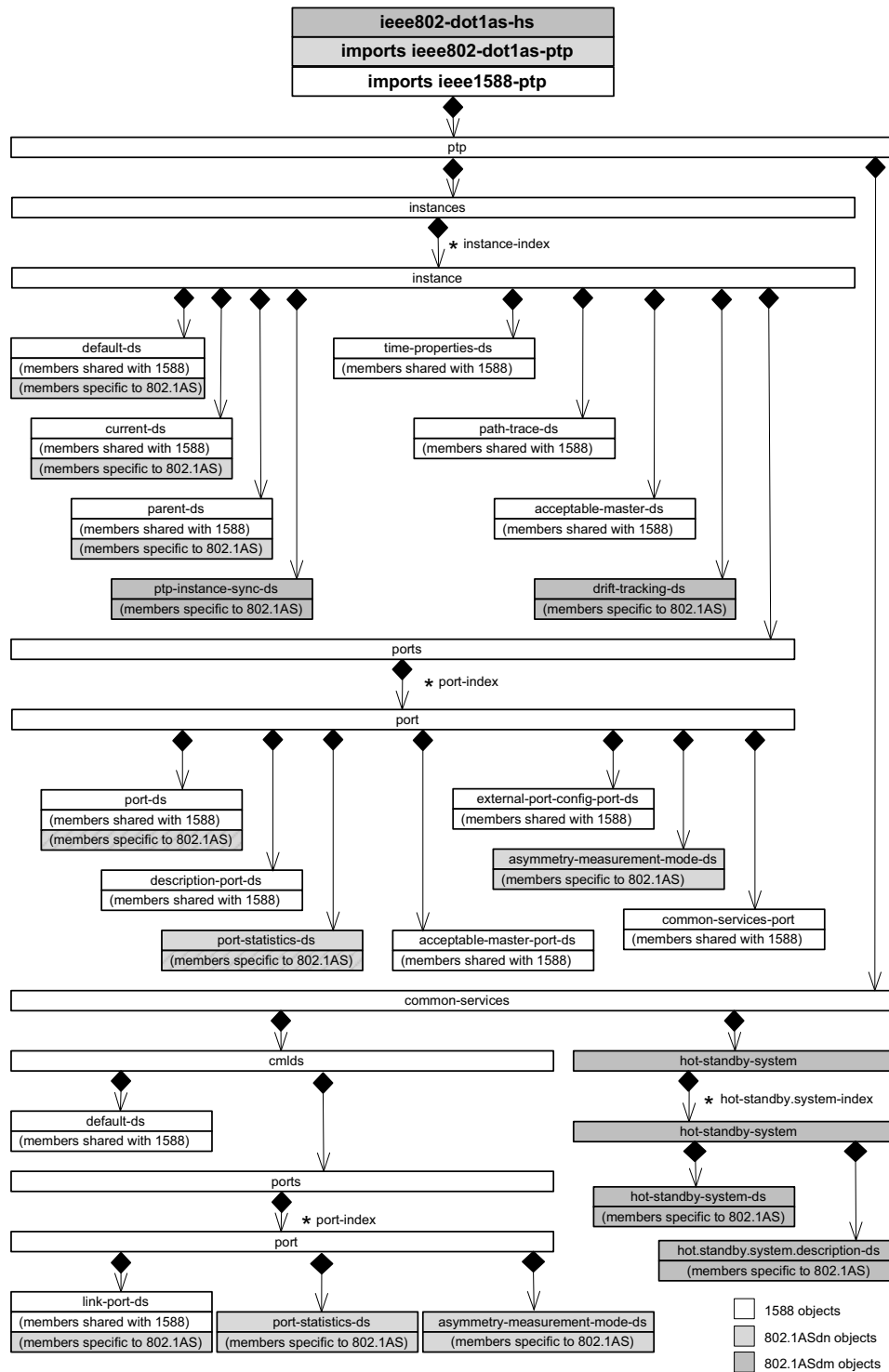


Figure 17-1—Overview of YANG tree

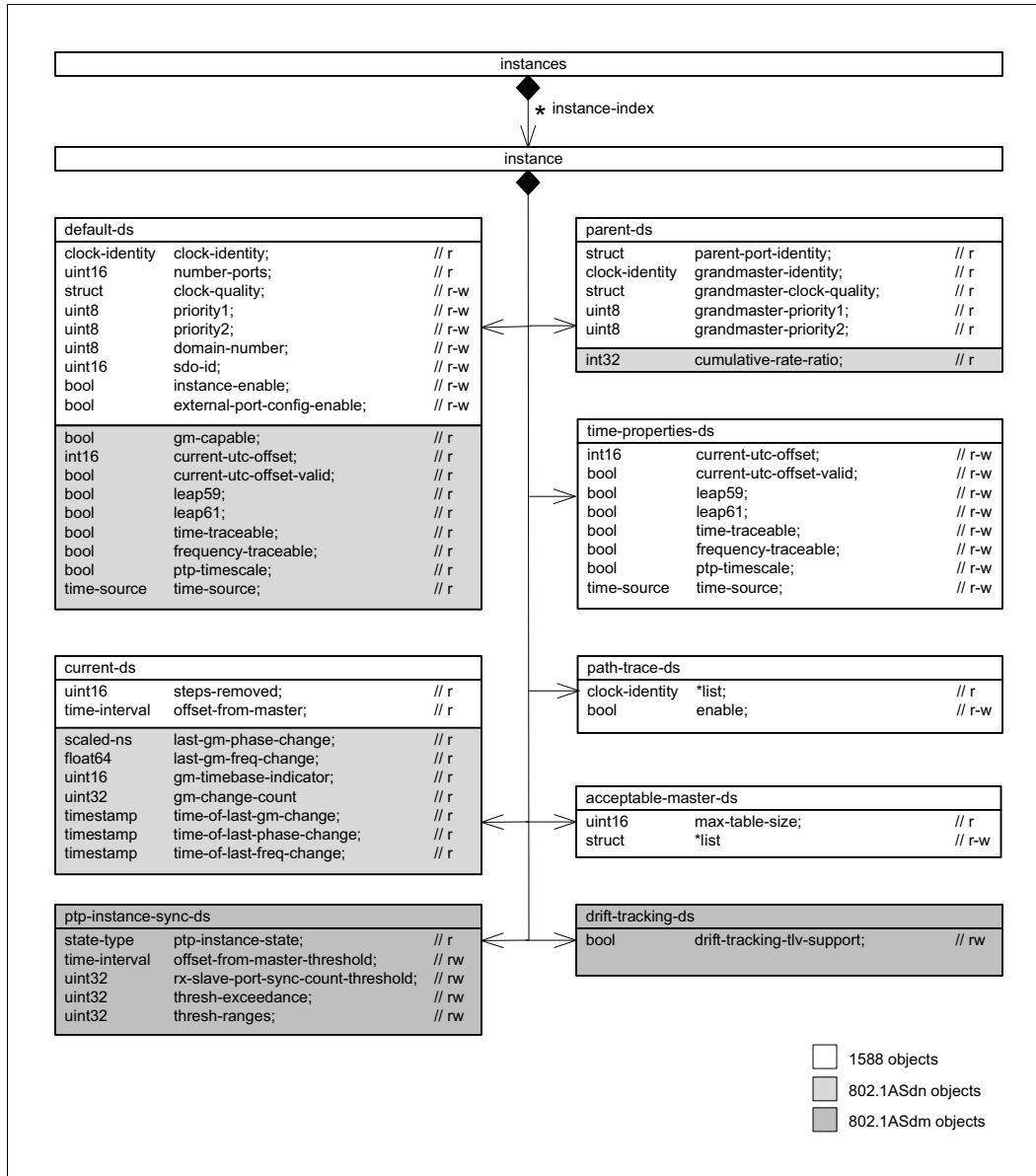


Figure 17-2—PTP Instance detail

```

|   +--rw rx_slave-port-sync-count-threshold?   uint32
|   +--rw thresh-exceedance?                     uint32
|   +--rw thresh-in-ranges?                       uint32
+--rw drift-tracking-ds
    +--rw drift-tracking-tlv-support?   boolean
augment /ptp:ptp/ptp:instances/ptp:instance/ptp:ports/ptp:port/ptp:port-ds:
    +--rw gptp-capable-state-machines-enabled?   boolean
augment /ptp:ptp/ptp:instances/ptp:instance/ptp:ports/ptp:port/dot1as-
ptp:port-statistics-ds:
    +--ro rx-sync-count-slave-p?   uint32
augment /ptp:ptp/ptp:common-services:

```

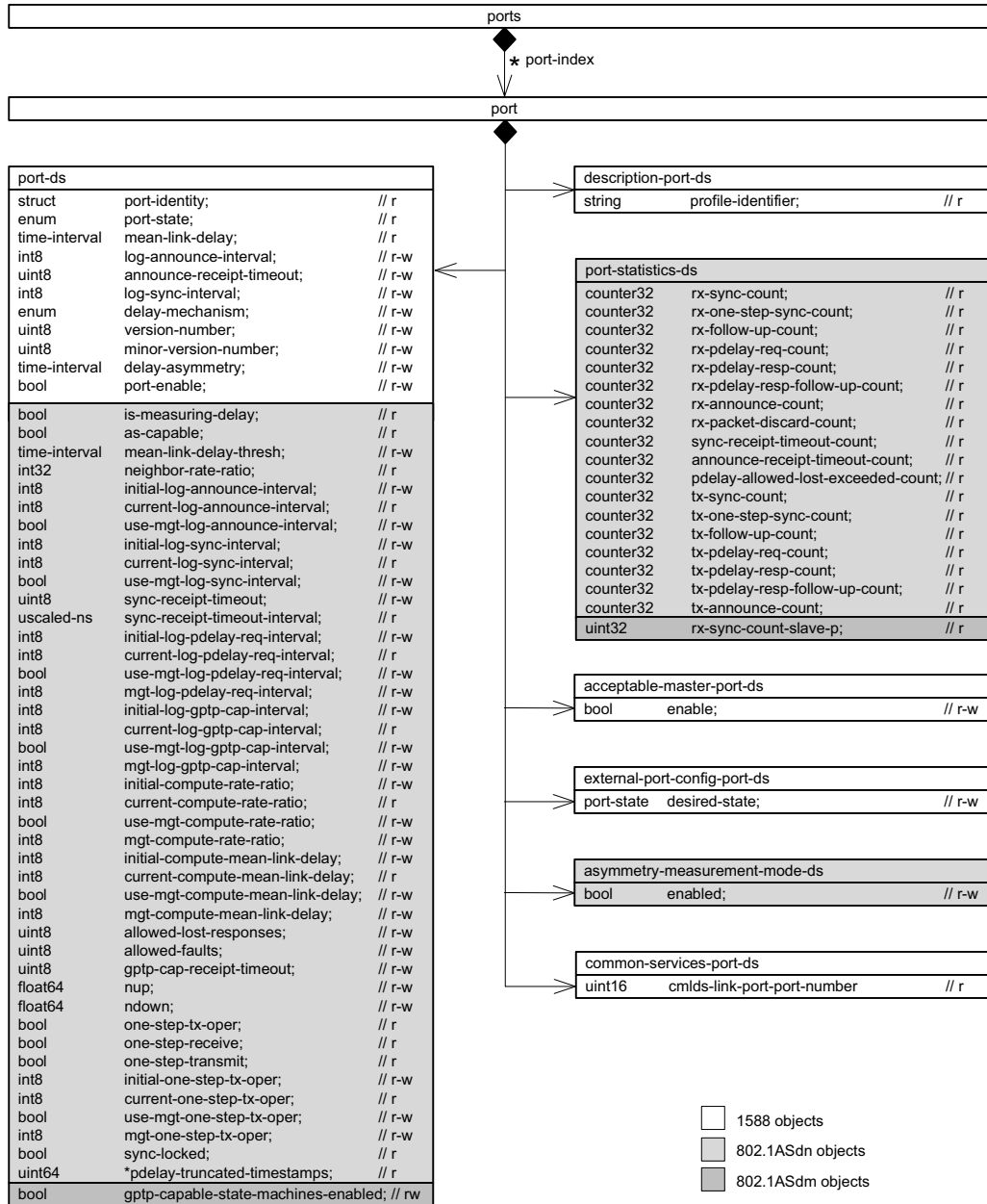


Figure 17-3—PTP Port detail

```

+--rw hot-standby-system {hot-standby}?
  +--rw hot-standby-sytem* [hot-standby-system-index]
    +--rw hot-standby-system-index          uint8
    +--rw hot-standby-system-ds
      | +--rw primary-ptp-instance-index?    uint32
      | +--rw secondary-ptp-instance-index?  uint32

```

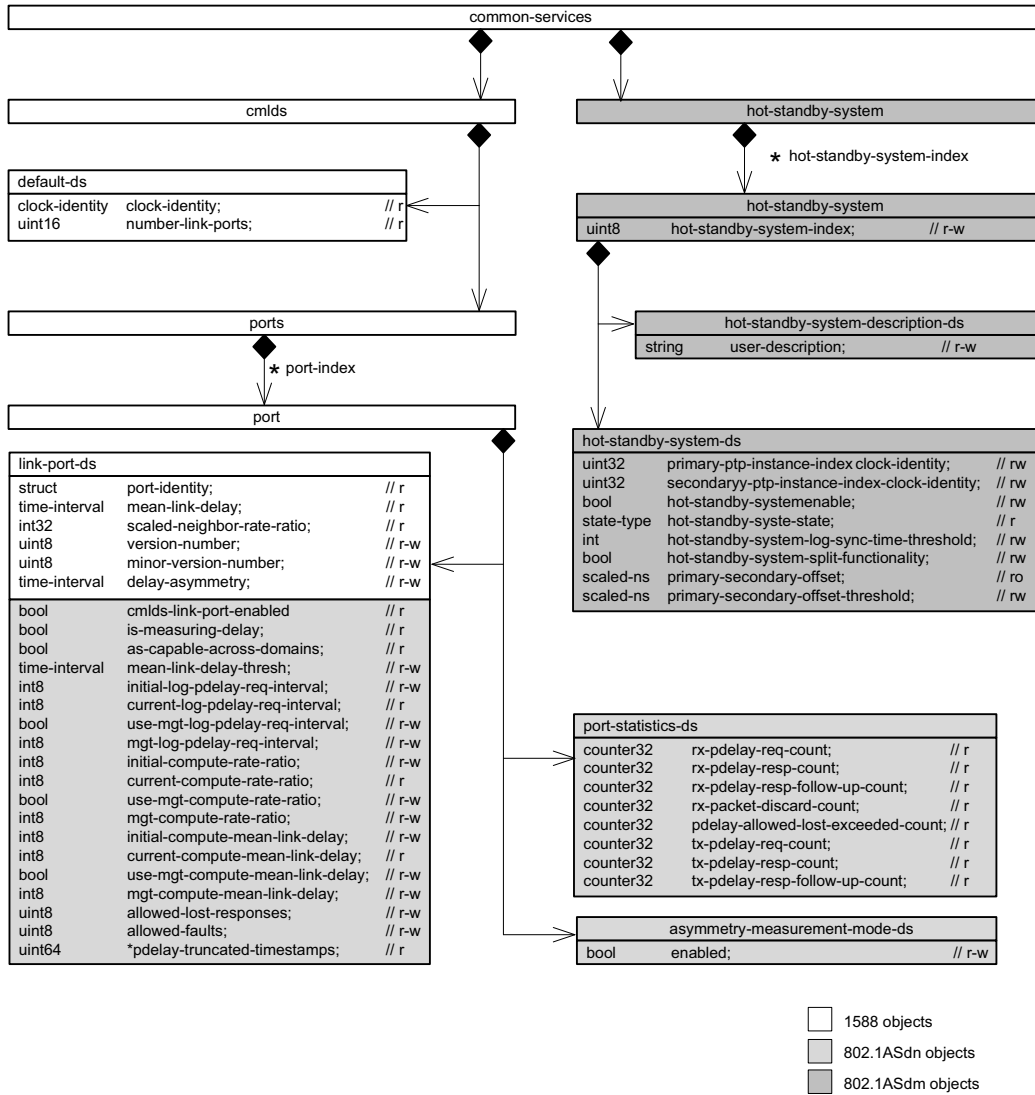


Figure 17-4—Common services detail

```

| +--rw hot-standby-system-enable?          boolean
| +--ro hot-standby-system-state?          hot-standby-
system-state-type
| +--rw hot-standby-system-log-sync-time-threshold?  int8
| +--rw hot-standby-system-split-functionality?      boolean
| +--ro primary-secondary-offset?                dot1as-ptp:scaled-
ns
| +--rw primary-secondary-offset-threshold?        dot1as-
ptp:scaled-ns
+--rw hot-standby-system-description-ds

```

**Table 17-1—Summary of the YANG modules**

Module	Managed functionality	YANG specification notes
ietf-yang-types	Type definitions	IETF RFC 6991 - Common YANG Data Types.
ieee1588-ntp	Clause 14	IEEE Std 1588e - MIB and YANG Data Models. IEEE Std 802.1ASdn imports this YANG module as its foundational tree, including a subset of members from Clause 14.
ieee802-dot1as-ntp	Clause 14	IEEE Std 802.1ASdn - YANG Data Model. The YANG module of this clause uses YANG augments to add members from Clause 14 that are unique to IEEE Std 802.1AS.
<a href="#">ieee802-dot1as-hs</a>	<a href="#">Clause 14</a>	<a href="#">IEEE Std 802.1ASdm - YANG Data Model.</a> <a href="#">The YANG module of this clause uses YANG augments to add members from Clause 14 that are unique to IEEE Std 802.1ASdm.</a>

---rw user-description? string

## 17.6 YANG modules<sup>1 2</sup>

*Insert 17.6.2 and renumber subsequent subclauses as appropriate:*

### 17.6.2 Module [ieee802-dot1as-hs.yang](#)

```

module ieee802-dot1as-hs {
  yang-version 1.1;
  namespace "urn:ieee:std:802.1AS:yang:ieee802-dot1as-hs";
  prefix dot1as-hs;

  import ieee1588-ntp {
    prefix ntp;
  }
  import ieee802-dot1as-ntp {
    prefix dot1as-ntp;
  }

  organization
    "IEEE 802.1 Working Group";
  contact
    "WG-URL: http://ieee802.org/1/
    WG-EMail: stds-802-1-1@ieee.org

    Contact: IEEE 802.1 Working Group Chair
    Postal: C/O IEEE 802.1 Working Group
    IEEE Standards Association
    445 Hoes Lane
    Piscataway, NJ 08854
    USA

```

<sup>1</sup>Copyright release for YANG modules: Users of this standard may freely reproduce the YANG modules contained in this subclause so that they can be used for their intended purpose.

<sup>2</sup>An ASCII version of the YANG modules are attached to the PDF version of this standard, and can be obtained by Web browser from the IEEE 802.1 Website at <https://1.ieee802.org/yang-modules/>.

```

1      E-mail: stds-802-1-chairs@ieee.org";
2
3      description
4          "Management objects that control hot standby systems as specified in
5          IEEE Std 802.1ASdm-d1.1.
6
7          Copyright (C) IEEE (2023).
8          This version of this YANG module is part of IEEE Std 802.1ASdm;
9          see the standard itself for full legal notices.";
10
11      revision 2023-07-12 {
12          description
13              "Published as part of IEEE Std 802.1ASdm-XXXX.
14              Initial version.";
15          reference
16              "IEEE Std 802.1ASdm-XXXX - YANG Data Model";
17      }
18
19      // The year (XXXX) will be replaced during publication.
20      // This is the 1st balloted draft D0.1
21      // of the YANG module for amendment IEEE P802.1ASdm.
22
23      feature hot-standby {
24          description
25              "This feature indicates that the device supports the hot-standby
26              functionality.";
27      }
28
29      typedef ptp-instance-state-type {
30          type enumeration {
31              enum not-capable {
32                  value 0;
33                  description
34                      "For all enabled PTP ports, asCapable is FALSE, i.e., the neighbor is
35                      not
36                      exchanging the messages that are required for conformance to this
37                      standard.";
38              }
39              enum synced {
40                  value 1;
41                  description
42                      "Time is synchronized to the requirements of this standard (see 17.5.3
43                      and 17.5.4), or to
44                      the requirements specified by a profile standard.";
45              }
46              enum not-synced {
47                  value 2;
48                  description
49                      "On the PTP Port in SlavePort state, the port failed to receive time
50                      synchronization
51                      event messages or is not synchronized according to the requirements
52                      of this standard (see 17.5.3 and
53                      17.5.4), or to the requirements specified by a profile standard.";
54              }
55              enum initializing {
56                  value 3;
57                  description
58                      "Time synchronization does not conform to the requirements of this
59                      standard.";
60              }
61          }
62      }

```

```

1      }
2    }
3  }
4
5  typedef hot-standby-system-state-type {
6    type enumeration {
7      enum init {
8        value 0;
9        description
10         "Initialization after the HotStandbySystem powers on and is enabled. In
11         this state, the system is
12         waiting for both PTP Instances to synchronize.";
13      }
14      enum redundant {
15        value 1;
16        description
17         "Both PTP Instances are synchronized according to the requirements of
18         the
19         respective application or profile standard (see 3.24). Time
20         synchronization is redundant.";
21      }
22      enum not-redundant {
23        value 2;
24        description
25         "One PTP Instance is synchronized, and the other PTP Instance is faulted
26         (not
27         synchronized). Time synchronization continues to meet the requirements
28         of the respective
29         application or profile standard (see 3.24). Time synchronization is not
30         redundant.";
31      }
32      enum out-of-sync {
33        value 3;
34        description
35         "The HotStandbySystem is adjusting phase/frequency of its local time
36         using the
37         data stored while the system was in the REDUNDANT or NOT_REDUNDANT
38         state, but the local
39         time will eventually drift relative to other time-aware systems. During
40         OUT_OF_SYNC state, time
41         synchronization might not meet the requirements of the respective
42         application or profile standard
43         (see 3.24).";
44      }
45    }
46  }
47
48  augment "/ptp:ptp/ptp:instances/ptp:instance" {
49    description
50     "Augment IEEE Std 1588 instanceList.";
51    container ptp-instance-sync-ds {
52      description
53       "The ptpInstanceSyncDS describes the synchronization status of the PTP
54       Instance.";
55      reference
56       "14.8 of IEEE Std 802.1ASdm-2023";
57      leaf ptp-instance-state {
58        type ptp-instance-state-type;
59        config false;

```



```

1      description
2          " State of the PTP Instance with respect to requirements of the
3      respective
4          application or profile standard";
5      reference
6          "14.8.2 of IEEE Std 802.1ASdm-2023";
7      }
8      leaf offset-from-master-threshold {
9          type ptp:time-interval;
10         config true;
11         description
12             "The value is the threshold for offsetFromMaster (see 17.5.1.2), below
13             which the PTP Instance is considered
14             to be synchronized.";
15         reference
16             "14.8.3";
17     }
18     leaf rx_slave-port-sync-count-threshold {
19         type uint32;
20         config true;
21         description
22             "The value of rxSlavePortSyncCountThreshold is the threshold for
23             rxSyncCountSlaveP (see 17.5.1.4), above
24             which the PTP Instance is considered to be synchronized.";
25         reference
26             "14.8.4";
27     }
28     leaf thresh-exceedance {
29         type uint32;
30         config true;
31         description
32             "The value of threshExceedance (see 17.5.1.7) is the threshold for the
33             number of consecutive exceedances of
34             offsetFromMasterThreshold (see 17.5.1.3) by offsetFromMaster (see
35             17.5.1.2), at which the PTP Instance is
36             no longer in the SYNCED state (see 17.5.3.3).";
37         reference
38             "14.8.5";
39     }
40     leaf thresh-in-ranges {
41         type uint32;
42         config true;
43         description
44             "The value of threshInRanges (17.5.1.9) is the threshold for the number
45             of consecutive occurrences of
46             offsetFromMaster (see 17.5.1.2) being within offsetFromMasterThreshold
47             (see 17.5.1.3), at which the PTP
48             Instance is placed in the SYNCED state (see 17.5.3.3).";
49         reference
50             "14.8.6";
51     }
52     }
53     container drift-tracking-ds {
54         description
55             "The driftTrackingDS contains a managed object that is used to enable or
56             disable the optional Drift_Tracking
57             TLV.";
58         reference
59             "14.9 of IEEE Std 802.1ASdm-2023";

```

```

1      leaf drift-tracking-tlv-support {
2          type boolean;
3          config true;
4          description
5              "The value of driftTrackingTlvSupport indicates whether the
6              Drift_Tracking TLV is enabled or disabled.";
7          reference
8              "14.9.2 of IEEE Std 802.1ASdm-2023";
9      }
10 }
11
12 augment "/ptp:ptp/ptp:instances/ptp:instance/ptp:ports/ptp:port/ptp:port-ds" {
13     description
14         "Augment IEEE Std 1588 commonServices.";
15     leaf gptp-capable-state-machines-enabled {
16         type boolean;
17         config true;
18         description
19             "A Boolean that is used to enable or disable the GptpCapableTransmit,
20             GptpCapableReceive, and
21             GptpCapableIntervalSetting state machines.";
22         reference
23             "14.8.55 of IEEE Std 802.1ASdm-2023";
24     }
25
26     augment "/ptp:ptp/ptp:instances/ptp:instance/ptp:ports/ptp:port/dotlas-
27     ptp:port-statistics-ds" {
28         description
29             "Augment IEEE Std 802.1AS PortStatisticsDS.";
30         leaf rx-sync-count-slave-p {
31             type uint32;
32             config false;
33             description
34                 "This counter increments whenever time synchronization information is
35                 received on a PTP Port
36                 when its port state is SlavePort.";
37             reference
38                 "14.10.20 of IEEE Std 802.1ASdm-2023";
39         }
40
41         augment "/ptp:ptp/ptp:common-services" {
42             description
43                 "Augment IEEE Std 1588 commonServices.
44
45                 14.6.9 of IEEE Std 802.1AS-2020 specifies neighborRateRatio
46                 (neighbor-rate-ratio), which is provided in YANG as the
47                 semantically equivalent node in ieee1588-ptp named
48                 scaled-neighbor-rate-ratio (in link-port-ds).";
49             container hot-standby-system {
50                 if-feature "hot-standby";
51                 description
52                     "The hotStandbyService structure contains the hotStandbySystemList, which
53                     is a list of instances of the Hot
54                     Standby Service.";
55                 reference
56                     "14.1.1.1 of IEEE Std 802.1ASdm-2023";

```

```

1      list hot-standby-sytem {
2          description
3              "List of instances of the Hot Standby Service";
4          reference
5              "14.1.1.1 of IEEE Std 802.1ASdm-2023";
6          key "hot-standby-system-index";
7          leaf hot-standby-system-index {
8              type uint8;
9              config true;
10         }
11         container hot-standby-system-ds {
12             description
13                 "The hotStandbySystemDS describes the attributes of the respective
14 instance of the Hot Standby Service.";
15             reference
16                 "14.19 of IEEE Std 802.1ASdm-2023";
17             leaf primary-ptp-instance-index {
18                 type uint32;
19                 config true;
20                 description
21                     "The value of primaryPtpInstanceIndex is the index (see 14.1.1) of
22 the primary PTP Instance associated with
23 this hotStandbySystem instance.";
24                 reference
25                     "14.19.2 of IEEE Std 802.1ASdm-2023";
26             }
27             leaf secondary-ptp-instance-index {
28                 type uint32;
29                 config true;
30                 description
31                     "The value of secondaryPtpInstanceIndex is the index (see 14.1.1)
32 of the secondaryPTP Instance associated
33 with this hotStandbySystem instance.";
34                 reference
35                     "14.19.3 of IEEE Std 802.1ASdm-2023";
36             }
37             leaf hot-standby-system-enable {
38                 type boolean;
39                 config true;
40                 description
41                     "The value is the hotStandbySystemEnable attribute of the
42 HotStandbySystem entity (see 17.6.1.2).";
43                 reference
44                     "14.19.4 of IEEE Std 802.1ASdm-2023";
45             }
46             leaf hot-standby-system-state {
47                 type hot-standby-system-state-type;
48                 config false;
49                 description
50                     "The value of hotStandbySystemState is the state of the
51 hotStandbySystem, i.e., the value of the global
52 variable hotStandbySystemState (see 17.6.1.1).";
53                 reference
54                     "14.19.5 of IEEE Std 802.1ASdm-2023";
55             }
56             leaf hot-standby-system-log-sync-time-threshold {
57                 type int8;
58                 config true;
59                 description

```

```

1         "The value of hotStandbySystemLogSyncTimeThreshold1 is the logarithm
2         to base 2 of the time interval, in
3         seconds, after which a hotStandbySystem in the INIT state
4         transitions from the INIT state to the
5         NOT_REDUNDANT state if the conditions to transition to the
6         REDUNDANT state are not met (see
7         17.6.4). The value -128 means that the transition time is zero,
8         i.e., the transition occurs immediately.";
9         reference
10        "14.19.6 of IEEE Std 802.1ASdm-2023";
11    }
12    leaf hot-standby-system-split-functionality {
13        type boolean;
14        config true;
15        description
16        "If the value is TRUE, the optional split functionality (see
17        17.6.3.4) is used. If the value is FALSE, the
18        optional split functionality is not used.";
19        reference
20        "14.19.8 of IEEE Std 802.1ASdm-2023";
21    }
22    leaf primary-secondary-offset {
23        type dotlas-ptp:scaled-ns;
24        config false;
25        description
26        "The absolute value of the difference between the clockSlaveTimes
27        (see 10.2.4.3) of the primary and
28        secondary PTP Instances.";
29        reference
30        "14.19.9 of IEEE Std 802.1ASdm-2023";
31    }
32    leaf primary-secondary-offset-threshold {
33        type dotlas-ptp:scaled-ns;
34        config true;
35        description
36        "The threshold for hotStandbySystemDS.primarySecondaryOffset (see
37        14.19.9), above which the
38        hotStandbySystemState transitions from REDUNDANT to NOT_REDUNDANT,
39        or does not transition from
40        NOT_REDUNDANT or OUT_OF_SYNC to REDUNDANT even if other conditions
41        for these transitions
42        are satisfied.";
43        reference
44        "14.19.10 of IEEE Std 802.1ASdm-2023";
45    }
46    }
47    container hot-standby-system-description-ds {
48        description
49        "The hotStandbySystemDescriptionDS contains descriptive information
50        for the respective instance of the Hot Standby Service.";
51        reference
52        "14.20 of IEEE Std 802.1ASdm-2023";
53        leaf user-description {
54            type string {
55                length "0..128";
56            }
57            config true;
58            description
59            "Configurable description of the hot standby system.";
60        }
61    }

```

```
1      reference
2      "14.20.3 of IEEE Std 802.1ASdm-2023";
3    }
4  }
5 }
6 }
7 }
8 }
9
10
11
12
13
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```

**Insert the following new Clause 18:**

## 18. Hot Standby

### 18.1 General

Hot standby includes:

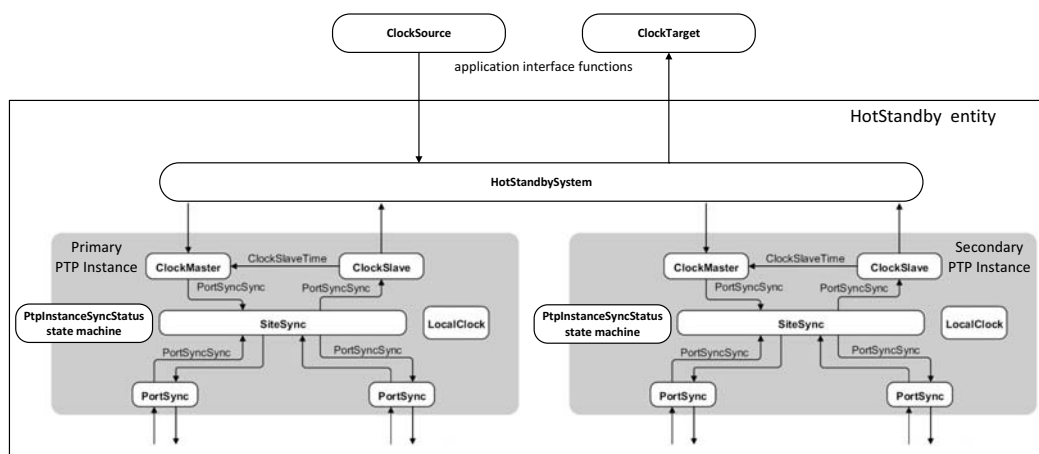
- A function that transforms the synchronized times of two generalized Precision Time Protocol (gPTP) domains into one synchronized time for use by applications;
- A function that directs the synchronized time of one gPTP domain into a different gPTP domain; and
- Mechanisms that determine whether a gPTP domain has sufficient quality to be used for hot standby.

For time synchronization using hot standby, two distinct domains are statically configured in the network and the Best Master Clock Algorithm (BMCA) is disabled for these domains. When hot standby is used for redundancy, a time-aware system operates two PTP Instances simultaneously, each in its own domain. One of the domains is considered the primary domain, and the other domain is considered the secondary domain. A time-aware system that has a primary domain and a secondary domain available uses the primary domain via the associated PTP instance. If the primary domain becomes unavailable (e.g., due to temporary or permanent failure of a physical link) and the secondary domain is still available to the time-aware system, the time aware system begins using the secondary domain.

Examples of hot standby are shown in 7.2.4 of this standard.

### 18.2 Overview

Figure 18-1 provides a model of hot standby for time synchronization.



**Figure 18-1—Model for hot standby entity in a time-aware system, and its interfaces to higher-layer applications**

Each PTP Instance has a corresponding `PtpInstanceSyncStatus` state machine (see 18.4). The `PtpInstanceSyncStatus` state machine monitors its PTP Instance to determine whether it is synchronized according to the requirements of the respective application or profile standard.

There is one `HotStandbySystem` entity, which interacts with the primary and secondary PTP Instances via the `HotStandbySystem` state machine (see 17.6) in order to provide a single value of time to the application (`ClockTarget`). This single value of time is based on the redundant values of time provided by either the primary PTP Instance, or the secondary PTP Instance, or both. There is one `HotStandbySystem` for both PTP Instances that use hot standby. The primary and secondary PTP Instances either both use the PTP timescale or both use the ARB timescale.

The `ClockTarget` entity represents the application that uses the synchronized time. The `ClockTarget` and its application interfaces are analogous to the `ClockTarget` specified in Clause 9. In the case of PTP Relay Instances, the `ClockTarget` is present if it is needed by the application.

The `ClockSource` entity is used to transfer the source of time when at least one of the PTP Instances is grandmaster-capable. The `ClockSource` and its application interfaces are analogous to the `ClockSource` specified in Clause 9. The `ClockSource` might be present if the PTP Instance is capable of becoming a Grandmaster PTP Instance.

There is one `LocalClock` entity for each PTP Instance. Therefore, Figure 18-1 shows two `LocalClock` entities.

NOTE—The `LocalClock` entities can be traceable to the same oscillator or to different oscillators.

### 18.3 PTP Instance configuration

PTP Instances that support hot standby are configured as follows:

- a) The two PTP Instances are enabled using `domainNumbers` as specified by the respective application or profile standard;
- b) For both PTP Instances, `externalPortConfigurationEnabled` is set to `TRUE`;
- c) If a PTP Instance (primary or secondary) is grandmaster, `externalPortConfigurationPortDS.desiredState` is configured to `MasterPort` or `PassivePort` for all PTP Ports (`portNumber` 1 and higher); otherwise, `externalPortConfiguration.desiredState` is configured to `SlavePort` for one PTP Port, and is configured to `MasterPort` or `PassivePort` for other PTP Ports; and
- d) Each PTP Instance shall have a corresponding `PtpInstanceSyncStatus` state machine (see 18.4).

### 18.4 `PtpInstanceSyncStatus` state machine

This state machine operates within the PTP Instance, to determine whether it is synchronized according to the requirements of the respective application or profile standard (see 3.24).

#### 18.4.1 `PtpInstanceSyncStatus` state machine global variables

**18.4.1.1 `isSynced`:** The synchronization status of the PTP Instance. The variable is a Boolean, and is `TRUE` if synchronization is sufficient according to the operation of the `PtpInstanceSyncStatus` state machine and `FALSE` otherwise. An application can set the `offsetFromMasterMax` (see 18.4.1.3),

rxSyncCountSlavePThresh (18.4.1.5), offsetMaxExceededCountThresh (18.4.1.7), and offsetMaxMetCountThresh (18.4.1.9) according to its requirements.

**18.4.1.2 offsetFromMaster:** The value of the managed object currentDS.offsetFromMaster (see 14.3.3).

**18.4.1.3 offsetFromMasterMax:** The value of the managed object hotstandbyDS.offsetFromMasterMax (see 14.8.3).

**18.4.1.4 rxSyncCountSlaveP:** The value of the managed object portStatisticsDS.rxSyncCountSlaveP (see 14.10.20).

**18.4.1.5 rxSyncCountSlavePThresh:** The value of the managed object hotstandbyDS.rxSyncCountSlavePThresh (see 14.8.4).

**18.4.1.6 offsetMaxExceededCount:** The current number of consecutive exceedances of offsetFromMasterMax (see 18.4.1.3) by offsetFromMaster (see 18.4.1.2). The data type of offsetMaxExceededCount is UInteger32.

**18.4.1.7 offsetMaxExceededCountThresh:** The threshold (see 14.8.5) for the number of consecutive exceedances of offsetFromMasterMax (see 18.4.1.3) by offsetFromMaster (see 18.4.1.2), at which the PTP Instance is no longer in the SYNCED state (see 18.4.3.3).

**18.4.1.8 offsetMaxMetCount:** The current number of consecutive occurrences of offsetFromMaster (see 18.4.1.2) being within offsetFromMasterMax (see 18.4.1.3). The data type of offsetMaxMetCount is UInteger32.

**18.4.1.9 offsetMaxMetCountThresh:** The threshold (see 14.8.6) for the number of consecutive occurrences of offsetFromMaster (see 18.4.1.2) being within offsetFromMasterMax (see 18.4.1.3), at which the PTP Instance is placed in the SYNCED state (see 18.4.3.3).

## 18.4.2 PtpInstanceSyncStatus state machine local variables

**18.4.2.1 lastRxSyncCountSlaveP:** Holds the last value of rxSyncCountSlaveP (see 18.4.1.4) that the state machine read for the SlavePort, as part of monitoring for a sync event message timeout.

## 18.4.3 PtpInstanceSyncStatus state machine functions

**18.4.3.1 isGtpCapable():** This function returns a boolean value that is TRUE when ptpPortEnabled (see 10.2.5.13) and asCapable (see 10.2.5.1) are TRUE for at least one PTP Port (i.e., port for which portNumber is not zero).

```
isGtpCapable()
{
    if (!instanceEnable)
        return (FALSE);
    for (int i = 1; i <= numberPorts; i++)
        // see 8.6.2.8 for numberPorts
        {
            if (portIsCapable (i))
                return (TRUE);
        }
    return (FALSE)
}
```



**18.4.3.2 isGm():** This function determines if the PTP Instance is a Grandmaster PTP Instance (TRUE) or not (FALSE) by searching the selectedState array (see 10.2.4.20) for absence or presence of at least one value equal to SlavePort (see Table 10-2).

```

isGm()
{
    for (int i = 1; i <= numberPorts; i++)
        // see 8.6.2.8 for numberPorts
        {
            if (selectedState[i] == SlavePort)
                return (FALSE);
        }
    return (TRUE);
}

```

**18.4.3.3 portIsCapable(j):** This function returns a boolean value that is TRUE if PTP Port j is enabled and asCapable is TRUE, and FALSE otherwise.

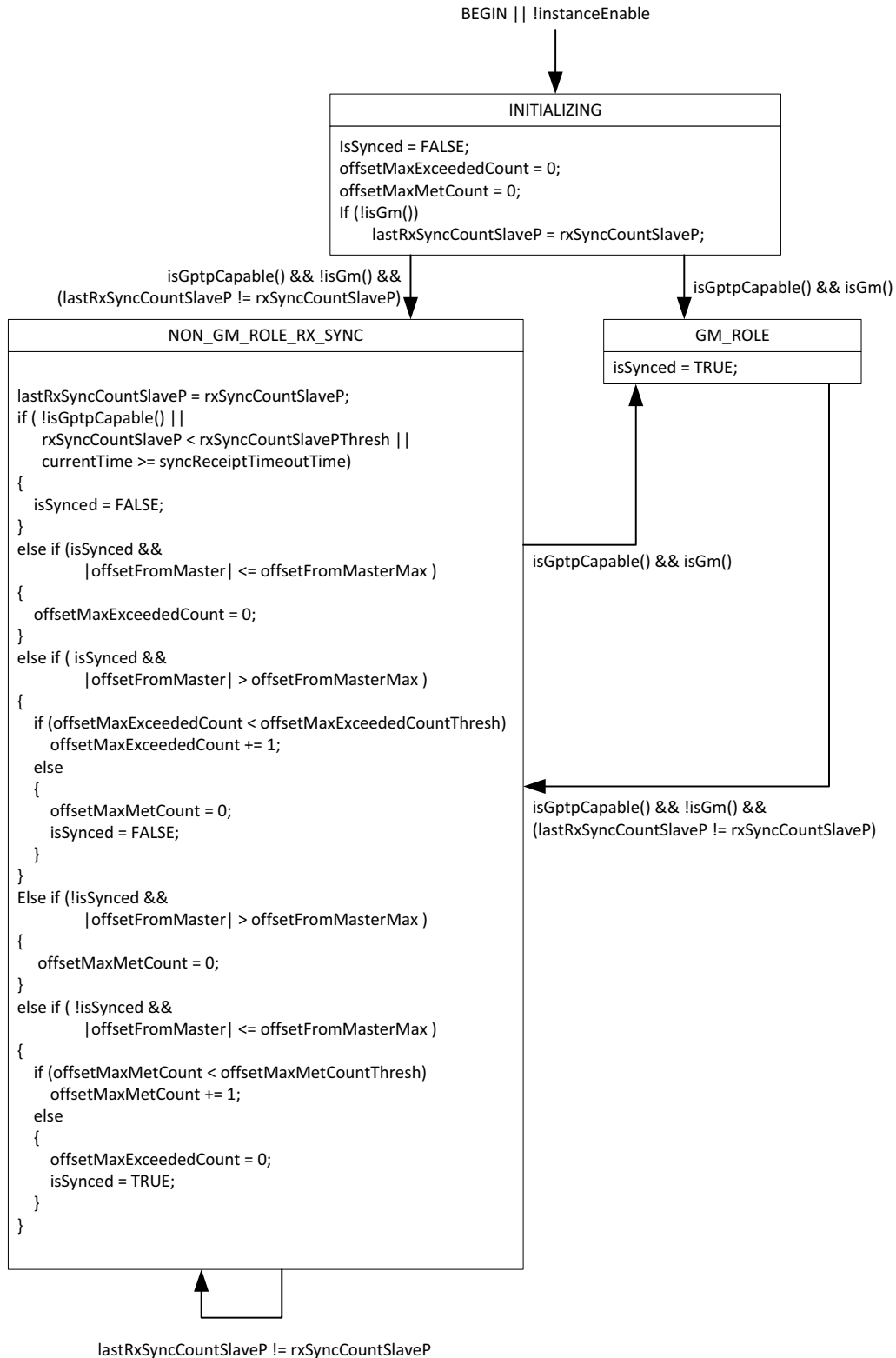
```

portIsCapable (j)
{
    if (ptpPortEnabled && asCapable)
        return (TRUE);
    else
        return (FALSE);
}

```

#### 18.4.4 State diagram

The PtpInstanceSyncStatus state machine shall implement the function specified by the state diagram in Figure 18-2, the variables specified in 18.4.1 and 18.4.2, and the functions specified in 18.4.3.



## 18.5 HotStandbySystem state machine

This state machine interacts with the primary and secondary PTP Instances in order to provide a single redundant time to the application.

### 18.5.1 HotStandbySystem state machine global variables

**18.5.1.1 hotStandbySystemState:** State of the HotStandbySystem. The variable is an enumeration that takes one of the following values:

- a) INIT: Initialization after the HotStandbySystem powers on and is enabled. In this state, the system is waiting for both PTP Instances to synchronize.
- b) REDUNDANT: Both PTP Instances are synchronized according to the requirements of the respective application or profile standard (see 3.24). Time synchronization is redundant.

NOTE—A PTP Instance is implicitly considered to be synchronized if it is a Grandmaster PTP Instance (see 18.4.3).

- c) NOT\_REDUNDANT: One PTP Instance is synchronized, and the other PTP Instance is faulted (not synchronized). Time synchronization continues to meet the requirements of the respective application or profile standard (see 3.24). Time synchronization is not redundant.
- d) OUT\_OF\_SYNC: The HotStandbySystem may adjust phase/frequency of its local time using the data stored while the system was in the REDUNDANT or NOT\_REDUNDANT state, but the local time will eventually drift relative to other time-aware systems. During OUT\_OF\_SYNC state, time synchronization might not meet the requirements of the respective application or profile standard (see 3.24).

**18.5.1.2 hotStandbySystemEnable:** A Boolean variable used to enable/disable the HotStandbySystem.

**18.5.1.3 hotStandbySystemLogSyncTimeThresh:** The value of the managed object hotStandbySystemLogSyncTimeThreshold (see 14.19.6) of the hotStandbySystemDS. The data type for hotStandbySystemLogSyncTimeThreshold is Integer8.

**18.5.1.4 primarySecondaryOffset:** The value of the managed object hotStandbySystemDS.primarySecondaryOffset (see 14.19.7), which is the absolute value of the difference between the clockSlaveTime variables (see 10.2.4.3) of the primary and secondary PTP Instances.

### 18.5.2 HotStandbySystem state machine local variables

**18.5.2.1 primary:** Reference to the data sets of the PTP Instance configured to use the primary domainNumber as specified by the respective application or profile standard (see 3.24). This references an element of the instanceList specified in 14.1.

**18.5.2.2 secondary:** Reference to the data sets of the PTP Instance configured to use the secondary domainNumber as specified by the respective application or profile standard (see 3.24). This references an element of the instanceList specified in 14.1.

**18.5.2.3 TEMP:** A temporary variable used to reduce clutter in the state diagram (see Figure 18-3). The data type for TEMP is Integer16.

**18.5.2.4 transitionTime:** The value of currentTime (see 10.2.4.12) after which the hotStandbySystem transitions from the OUT\_OF\_SYNC state to either the NOT-REDUNDANT or REDUNDANT state, or from the NOT\_REDUNDANT to the REDUNDANT state, if all other conditions for the respective transition is met. The data type for transitionTime is UScaledNs.

**18.5.2.5 primarySecondaryOffsetThresh:** The value of the managed object `hotStandbySystemDS.primarySecondaryOffsetThresh`, which is the threshold for `primarySecondaryOffset` (see 18.5.1.4), above which the `hotStandbySystemState` transitions from `REDUNDANT` to `NOT_REDUNDANT`, or stays in `NOT_REDUNDANT` or `OUT_OF_SYNC` even if other conditions for transitioning to `REDUNDANT` are satisfied.

### 18.5.3 HotStandbySystem requirements

#### 18.5.3.1 Primary grandmaster

When the primary PTP Instance is a Grandmaster PTP Instance (i.e., no PTP Port in `SlavePort` state), the `HotStandbySystem` shall transfer phase, frequency, `clockSourceTimeBaseIndicator` (see 10.2.4.8), `clockSourceLastGmPhaseChange` (see 10.2.4.10), and `clockSourceLastGmFreqChange` (see 10.2.4.11) from the `ClockSource` to the `ClockMaster` of the primary PTP Instance (see Figure 18-1). If no external source of time is implemented, the `ClockSource` is equivalent to the `LocalClock` of this PTP Instance.

When the primary PTP Instance is grandmaster, there is no requirement for the `HotStandbySystem` to receive time from the `ClockSlave` of the secondary PTP Instance.

#### 18.5.3.2 Secondary grandmaster

##### 18.5.3.2.1 HotStandbySystem in REDUNDANT state

When the secondary PTP Instance is a Grandmaster PTP Instance (i.e., no PTP Port in `SlavePort` state), and the `hotStandbySystemState` is `REDUNDANT`, the `HotStandbySystem` shall transfer phase, frequency, `clockSourceTimeBaseIndicator` (see 10.2.4.8), `clockSourceLastGmPhaseChange` (see 10.2.4.10), and `clockSourceLastGmFreqChange` (see 10.2.4.11) from the `ClockSlave` of the primary PTP Instance to the `ClockMaster` of the secondary PTP Instance (see Figure 18-1). By using phase from the primary PTP Instance, the secondary grandmaster can maintain continuity in the event of a fault in the primary grandmaster.

##### 18.5.3.2.2 HotStandbySystem in NOT\_REDUNDANT state

When the secondary PTP Instance is a Grandmaster PTP Instance (i.e., no PTP Port in `SlavePort` state), and the `hotStandbySystemState` is `NOT_REDUNDANT`, the `HotStandbySystem` shall transfer phase, frequency, `clockSourceTimeBaseIndicator` (see 10.2.4.8), `clockSourceLastGmPhaseChange` (see 10.2.4.10), and `clockSourceLastGmFreqChange` (see 10.2.4.11) from the `ClockSource` to the `ClockMaster` of the secondary PTP Instance (see Figure 18-1). If no external source of time is implemented, the `ClockSource` is equivalent to the `LocalClock` of this PTP Instance.

In this situation, the secondary PTP Instance is considered synchronized because it is a Grandmaster PTP Instance (see 18.4.4 and Figure 18-2). Since the `hotStandbySystemState` is `NOT_REDUNDANT`, this means that either the primary PTP Instance is not synchronized or `primarySecondaryOffset` exceeds `primarySecondaryOffsetThresh`. However, in the former case `primarySecondaryOffset` can exceed `primarySecondaryOffsetThresh` because in this case the primary and secondary PTP Instances receive timing from different sources. The result is that the `HotStandbySystem` remains in the `NOT_REDUNDANT` state, even if the primary PTP Instance subsequently becomes synchronized again by the primary Grandmaster PTP Instance, because `primarySecondaryOffset` continues to exceed `primarySecondaryOffsetThresh`. This situation would not occur if at the secondary PTP Instance grandmaster it were the primary PTP Instance that was synchronized and the secondary PTP Instance that was not synchronized; in this case, the secondary Grandmaster PTP Instance `ClockMaster` would receive timing from the primary PTP Instance `ClockSlave` (see 18.5.3.2.1) and the `hotStandbySystemState` would change to `REDUNDANT`.

The situation described in the previous paragraph is avoided if the optional split functionality (see 18.5.3.4) is used. In this case, timing is transferred directly from the secondary SiteSync entity to the primary SiteSync entity. When this happens, primarySecondaryOffset changes and, since now the primary PTP Instance is synchronized to the secondary PTP Instance, is within primarySecondaryOffsetThresh. The HotStandbySystem state changes to REDUNDANT and the secondary Grandmaster PTP Instance begins receiving timing from the primary PTP Instance. However, even if the optional split functionality is not used, an application can choose to re-initialize the secondary PTP Instance after the primary PTP Instance becomes synchronized. If this is done, the secondary Grandmaster PTP Instance receives timing from the primary PTP Instance ClockSlave after initialization and hotStandbySystemState is REDUNDANT.

<<Editor's note: It seems that an explicit, normative statement is needed that on initialization of the secondary Grandmaster PTP Instance, if the primary PTP Instance is synchronized the secondary receives timing from the primary and the hotStandbySystemState becomes REDUNDANT. It is not clear that this follows from Figure 17-3 and the text in 17.5 or its subclauses. Comments are requested. >>

### 18.5.3.2.3 Transition of hotStandbySystemState from REDUNDANT to NOT\_REDUNDANT

The secondary grandmaster shall conform to the time synchronization requirements of the respective application or profile standard (see 3.24) during a transition of HotStandbySystemState from REDUNDANT to NOT\_REDUNDANT due to a fault in its primary PTP Instance (i.e., primary.isSynced == FALSE).

NOTE—The secondary grandmaster is responsible for maintaining continuity (no "jump" in time) when a fault occurs in the primary grandmaster.

NOTE—The secondary domain global variables lastGmPhaseChange and lastGmFrequencyChange are propagated throughout the secondary domain, and can be used to compensate for any phase or frequency jump that occurs when the HotStandbySystem state becomes NOT\_REDUNDANT.

### 18.5.3.3 Slave HotStandbySystem

#### 18.5.3.3.1 General

A HotStandbySystem is said to be a Slave HotStandbySystem if and only if neither PTP Instance is a Grandmaster PTP Instance.

#### 18.5.3.3.2 HotStandbySystem in REDUNDANT state

When the HotStandbySystem is a slave (i.e., neither PTP Instance is a grandmaster), and hotStandbySystemState is REDUNDANT, the HotStandbySystem shall transfer phase, frequency, gmTimeBaseIndicator (see 10.2.4.15), lastGmPhaseChange (see 10.2.4.16), and lastGmFrequencyChange (see 10.2.4.17) from the ClockSlave of the primary PTP Instance to the ClockTarget (application) if a ClockTarget is present.

#### 18.5.3.3.3 HotStandbySystem in NOT\_REDUNDANT state

##### 18.5.3.3.3.1 Transfer of synchronized time to the ClockTarget

When the HotStandbySystem is a slave, and the HotStandbySystem is in the NOT\_REDUNDANT state, the HotStandbySystem shall transfer phase, frequency, gmTimeBaseIndicator (see 10.2.4.15), lastGmPhaseChange (see 10.2.4.16), and lastGmFrequencyChange (see 10.2.4.17) from the ClockSlave of the synchronized PTP Instance (i.e. isSynced == TRUE) to the ClockTarget if a ClockTarget is present.

#### 18.5.3.4 Split functionality

The HotStandbySystem may provide an optional split functionality. If the managed object hotStandbySystemDS.hotstandbySystemSplitFunctionality is set to TRUE, the split functionality is invoked. If the managed object hotStandbySystemDS.hotstandbySystemSplitFunctionality is set to FALSE, the split functionality is not invoked. The split functionality shall provide an interworking function (IWF) that transfers time synchronization information from the primary PTP Instance to the secondary PTP Instance when the secondary PTP Instance is in the NOT\_SYNCED state and the primary PTP Instance is in the SYNCED state, or from the secondary PTP Instance to the primary PTP Instance when the primary PTP Instance is in the NOT\_SYNCED state and the secondary PTP Instance is in the SYNCED state. The IWF provides the most recently received PortSyncSync structure at the slave port of the synchronized PTP Instance (isSynced == TRUE, see 18.4.1.1) to the unsynchronized PTP Instance (isSynced == FALSE) SiteSyncSync entity, in order to generate a PortSyncSync structure that is used by the master port of the unsynchronized PTP Instance, as follows:

- a) The domainNumber is changed from the synchronized PTP Instance domainNumber to the unsynchronized PTP Instance domainNumber;
  - b) localPortNumber is changed to the portNumber of the unsynchronized PTP Instance slave port; and
  - c) All other members of the synchronized PTP Instance PortSyncSync structure are provided to the unsynchronized PTP Instance SiteSync entity unchanged.
  - d) The flags leap61, leap59, currentUtcOffsetValid, ptpTimescale, timeTraceable, and frequencyTraceable, and the fields currentUtcOffset and timeSource of received Announce messages on the slave port of the unsynchronized PTP Instance are ignored. The IWF provides the values of these corresponding members of the timePropertiesDS of the synchronized PTP Instance to the unsynchronized PTP Instance. These values are then copied to:
    - i) the respective members of the timePropertiesDS of the unsynchronized PTP Instance;
    - ii) the unsynchronized PTP Instance global variables leap61, leap59, currentUtcOffsetValid, ptpTimescale, timeTraceable, and frequencyTraceable, and the fields currentUtcOffset and timeSource, respectively; and
    - iii) the unsynchronized PTP Instance global variables annLeap61, annLeap59, annCurrentUtcOffsetValid, annPtpTimescale, annTimeTraceable, and annFrequencyTraceable, and the fields annCurrentUtcOffset and annTimeSource, respectively.
  - e) The messagePriorityVector (see 10.3.4 and 10.3.5) information of received Announce messages on the slave port of the unsynchronized PTP Instance is ignored. The IWF provides the values of the corresponding members of the synchronized PTP Instance parentDS, i.e., grandmasterIdentity, grandmasterClockQuality.clockClass, grandmasterClockQuality.clockAccuracy, grandmasterClockQuality.offsetScaledLogVariance, grandmasterPriority1, grandmasterPriority2, which are copied to:
    - i) the corresponding members of the parentDS of the unsynchronized PTP Instance; and
    - ii) the gmPriorityVector of the unsynchronized PTP Instance.
- The IWF also provides the value of the stepsRemoved member of the synchronized PTP Instance currentDS, which is copied to:
- iii) the stepsRemoved member of the currentDS of the unsynchronized PTP Instance; and
  - iv) the gmPriorityVector of the unsynchronized PTP Instance.

NOTE 1—With the above, the secondary/primary PTP Instance state machines operate as though the time synchronization information had been received from the secondary/primary PTP Instance slave port. The SiteSync entity of the secondary/primary PTP Instance transfers the timing information to the PortSync entity of each master port of the secondary/primary PTP Instance. Each PortSyncSync state machine computes rateRatio, which now is relative to the primary/secondary PTP Instance GM. Each MDSyncSend state machine computes the fields of transmitted Sync and, in

the two-step case, Follow\_Up messages. The copied syncReceiptTimeout time is less than currentTime because sync receipt timeout has not occurred at the primary/secondary PTP Instance.

NOTE 2—The split functionality is used to transfer time synchronization information from the PTP Instance that is in the SYNCED state to the PTP Instance that is in the NOT\_SYNCED state. It is not meant to cover the case where the primary and secondary PTP Instances are both in the NOT\_SYNCED state.

### 18.5.3.5 Both PTP Instances not in the SYNCED state

When the primary and secondary PTP Instances are both not in the SYNCED state, the HotStandbySystemState is OUT\_OF\_SYNC and the HotStandbySystem shall transfer phase, frequency, gmTimeBaseIndicator (see 10.2.4.15), lastGmPhaseChange (see 10.2.4.16), and lastGmFrequencyChange (see 10.2.4.17) from the ClockSlave of the PTP Instance that was in the SYNCED state, before the HotStandbySystem transitioned to the OUT\_OF\_SYNC state, to the ClockTarget if a ClockTarget is present.

When the HotStandbySystem state is OUT\_OF\_SYNC, time synchronization performance is not required to meet the respective application or profile standard (see 3.24) requirements. Nevertheless, in order to mitigate LocalClock frequency drift, the PTP Instance should adjust the phase/frequency of its LocalClock using the data stored while the PTP Instance was in the SYNCED state and the HotStandbySystem was in the REDUNDANT or NOT\_REDUNDANT state.

### 18.5.4 State diagram

The HotStandbySystem state machine shall implement the function specified by the state diagram in Figure 17-3, the variables specified in 18.5.1 and 18.5.2, and the state requirements specified in 18.5.3.

## 18.6 PrimarySecondaryOffset state machine

This state machine updates the value of the global variable primarySecondaryOffset (see 18.5.1.4) whenever a new value of clockSlaveTime (see 10.2.4.3) is computed for either the primary or secondary PTP Instance. A new value of clockSlaveTime is computed when the global variable rcvdPSSyncCSS (see 10.2.4.28) or the global variable rcvdLocalClockTickCSS (see 10.2.4.29) is set to TRUE.

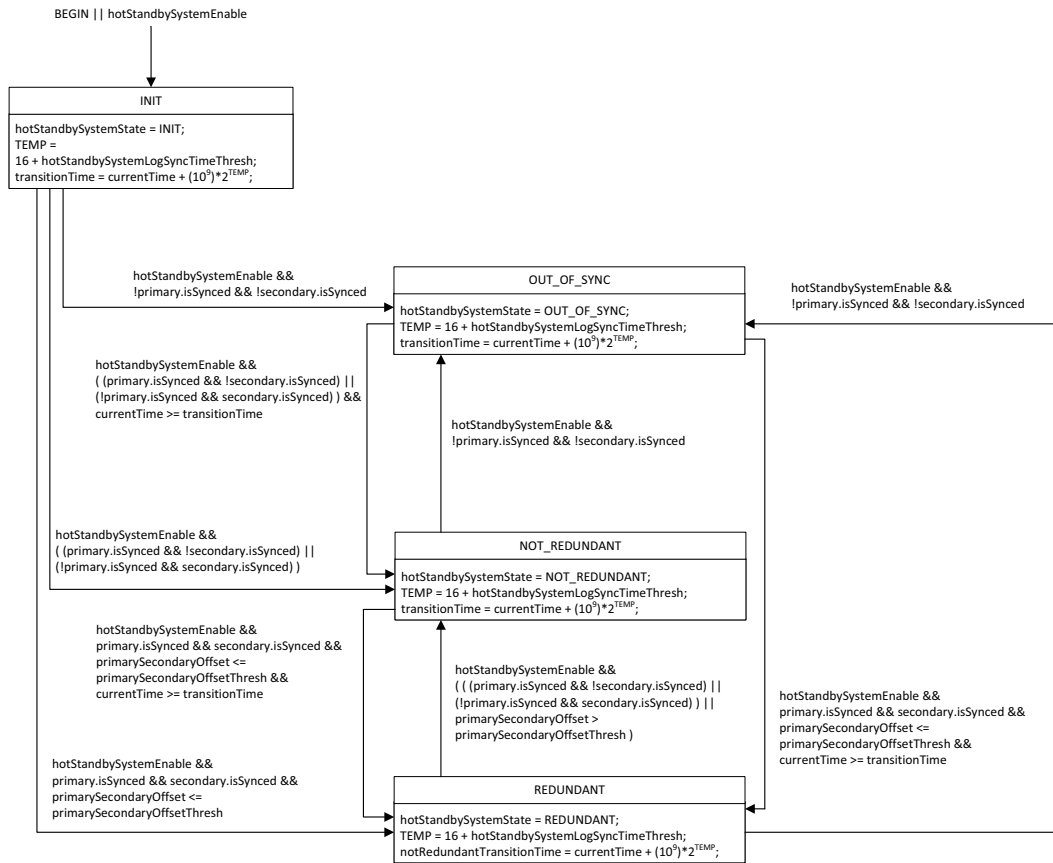
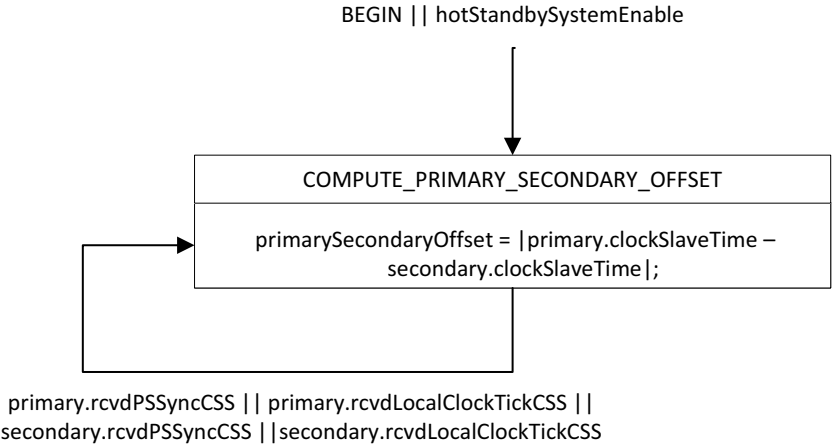


Figure 18-3—HotStandbySystem State Diagram

### 18.6.1 State diagram

The PrimarySecondaryOffset state machine shall implement the function specified by the state diagram in Figure 18-4..





**Figure 18-4—PrimarySecondaryOffset State Diagram**

## Annex A

(normative)

### Protocol Implementation Conformance Statement (PICS) proforma<sup>3</sup>

#### A.5 Major capabilities

*Change the table of A.5 as follows:*

Item	Feature	Status	References	Support
<del>DOM0</del>	<del>Does the time-aware system support a PTP Instance with domain number 0, in accordance with the requirements of 8.1?</del>	<del>M</del>	<del>item a) of 5.4.2, 8.1</del>	<del>Yes [ ]</del>
DOMADD	Does the time-aware system support one or more PTP Instances with <del>domain number</del> domainNumber in the range 1 to 127?	O	item f) of 5.4.2, 8.1	Yes [ ]    No [ ]
SIG	Does the PTP Instance transmit Signaling messages?	<del>O</del> <u>M</u>	<del>10.2.13, item e) of 5.4.2, item i) of 5.4.1, 10.4,</del> 10.6.4, A.8	Yes [ ] <del>No [ ]</del>
APPL	Does the PTP Instance support one or more of the application interfaces?	O	item i) of 5.4.2, Clause 9, A.20	Yes [ ]    No [ ]
<u>HOTSTDBY</u>	<u>Does the time-aware system support hot standby as specified in Clause 17?</u>	<u>O</u>	<u>Clause 17, 14.8, item l) of 5.4.2, 9.3.3.4, 9.4.3.4, 9.5.3.3, 9.6.2.6</u>	<u>Yes [ ]</u> <u>No [ ]</u>
<u>DRFTRK</u>	<u>Does the PTP Instance support the Drift Tracking TLV as specified in 10.2.4.25, 10.2.4.26, 10.2.4.27, 11.2.14, 11.2.15, and 11.4.4?</u>	<u>O</u>	<u>item n) of 5.4.2, 10.2.4.25, 10.2.4.26, 10.2.4.27, 11.2.14, 11.2.15, and 11.4.4</u>	<u>Yes [ ]</u> <u>No [ ]</u>

<sup>3</sup> Copyright release for PICS proformas: Users of this standard may freely reproduce the PICS proforma in this annex so that it can be used for its intended purpose and may further publish the completed PICS.

**A.7 Minimal time-aware system***Change the table of A.7 as follows:*

Item	Feature	Status	References	Support
MINTA-6	Is the <del>domain-number</del> domainNumber for all transmitted messages in the range 0 through 127, in compliance with the requirements of 8.1?	M	8.1	Yes [ ]
MINTA-8	Is the <del>domain-number</del> domainNumber for at least one of the gPTP domains supported by the time-aware system, in compliance with the requirements of 8.1?	M	8.1	Yes [ ]
MINTA-10	If path delay asymmetry is modeled by this PTP Instance, does it comply with the requirements of 8.3?	O	8.3	Yes [ ] No [ ] N/A [ ]
MINTA-15	Does the PTP Instance support the state machines related to signaling gPTP capability?	M	item h) of 5.4.1, 10.4.2	Yes [ ]
MINTA-16	For receive of all messages and for transmit of all messages except Announce <del>and Signaling</del> , does the PTP Instance support the message requirements?	M	item i) of 5.4.1, 10.5, 10.6, 10.7	Yes [ ]

*Insert the following items after MINTA-20 in the table of A.7:*

Item	Feature	Status	References	Support
MINTA-21	Do the PTP Ports of this PTP Instance implement the functionality of the GtpCapableIntervalSetting state machine in compliance with the requirements of 10.4.4 and Figure 10-23?	MINTA-15:O	10.4.4, item i) of 5.4.2, item i) of 5.4.1	Yes [ ] No [ ]
MINTA-22	Does the PTP Instance implement the PtpInstanceSyncStatus state machine of 18.4 and the PtpInstanceStatusDS of 14.8?	¬HOTST DBY:O	18.4, 14.8	Yes [ ] No [ ]

## A.8 Signaling

*Change SIG-7 of the table of A.8 as follows:*

Item	Feature	Status	References	Support
SIG-7	Does the message interval request TLV for Signaling messages comply with the requirements in 10.6.4.3?	<del>SIG:M</del> BMC-23:M, MIMSTR-16:M, MDFDPP-6:M	10.6.4.3	Yes [ ]

## A.9 BBest master clock

*Change BMC-11 of the table of A.9 as follows:*

Item	Feature	Status	References	Support
BMC-11	Is the PTP Port number equal to 1 in compliance with the requirements of 8.5.2.3?	<del>BRDG:M</del>	8.5.2.3	Yes [ ]    N/A [ ]

*Insert the following item after BMC-22 in the table of A.9:*

Item	Feature	Status	References	Support
BMC-23	Do the MasterPorts of this PTP Instance implement the functionality of the AnnounceIntervalSetting state machine in compliance with the requirements of 10.3.17 and Figure 10-19?	BMC:O	10.3.17, item h) of 5.4.2, item i) of 5.4.1	Yes [ ]    No [ ]

## A.10 Grandmaster-capable PTP Instance

*Change GMCAP-2 of the table of A.10 as follows:*

Item	Feature	Status	References	Support
GMCAP-2	Does the PTP Instance implement the functionality specified by the Clock <del>Master</del> SyncOffset state machine in compliance with the requirements of 10.2.10 and Figure 10-6?	GMCAP:M	10.2.10	Yes [ ]

**A.11 Media-independent master***Insert the following item after MIMSTR-15 in the table of A.11:*

Item	Feature	Status	References	Support
MIMSTR-16	Do the MasterPorts of this PTP Instance implement the functionality of the SyncIntervalSetting state machine in compliance with the requirements of 10.3.18 and Figure 10-20?	MIMSTR:O	10.3.18, item b)3) of 5.4.2, item i) of 5.4.1	Yes [ ] No [ ]

**A.13 Media-dependent, full-duplex point-to-point link***Delete MDFDPP-5 from the table of A.13, and renumber subsequent items as appropriate**Change MDFDPP-6 in the table of A.13 as follows:*

Item	Feature	Status	References	Support
MDFDPP-6	Does this port implement the functionality of the LinkDelayIntervalSetting state machine in compliance with the requirements of 11.2.21 and Figure 11-11?	MDFDPP: <del>M</del> <u>Q</u>	11.2.21, <u>item i) of 5.5, item i) of 5.4.1</u>	Yes [ ] <u>No [ ]</u>
<u>MDFDPP-36</u>	<u>Does this port support two-step capability on receive?</u>	<u>MDFDPP:M</u>	<u>11.2.14, item d) of 5.5</u>	<u>Yes [ ]</u> <u>No [ ]</u>
<u>MDFDPP-37</u>	<u>Does this port support two-step capability on transmit?</u>	<u>MDFDPP:M</u>	<u>11.2.15, item e) of 5.5</u>	<u>Yes [ ]</u> <u>No [ ]</u>

***Change the table of A.19 as follows:***

## **A.19 Remote management**

Item	Feature	Status	References	Support
	If item RMGT is not supported, mark N/A.			N/A [ ]
RMGT-1	What management protocol standard(s) or specification(s) are supported?	RMGT:M	item k) 1) of 5.4.2	
RMGT-2	What standard(s) or specification(s) for managed object definitions and encodings are supported for use by the remote management protocol?	RMGT:M	item k) 2) of 5.4.2	
RMGT-3	If the Simple Network Management Protocol (SNMP) is listed in RMGT-2, is the IEEE 8021-AS-MIB module fully supported (per its MODULE-COMPLIANCE)?	RMGT:O	item k) 3) of 5.4.2, Clause 15	Yes [ ] No [ ] <a href="#">N/A [ ]</a>
RMGT-4	If a remote management protocol that supports YANG is listed in RMGT-2, is the YANG data model ieee802-dot1as-ptp of Clause 17 supported?	RMGT:O	item k) 4) of 5.4.2, Clause 17	Yes [ ] No [ ]
RMGT-5	If a remote management protocol that supports YANG is listed in RMGT-2, and if hot standby is supported, is the YANG data model ieee802-dot1as-hs of Clause 17 supported?	RMGT:O	item k) 5) of 5.4.2, Clause 17, Clause 18	Yes [ ] No [ ] N/A [ ]

## Annex F

(informative)

### PTP profile included in this standard

#### F.3 PTP attribute values

*Change F.3 a) as follows:*

- a) A domain whose ~~domain-number~~domainNumber is 0 is present. A domain whose ~~domain~~  
~~number~~domainNumber is in the range 1 through 127 can be present (see 8.1).

## Annex G

(informative)

### The asymmetry compensation measurement procedure based on line-swapping

#### G.3 Measurement procedure

*Change the second paragraph of G.3 as follows:*

The criteria of 11.2.17 for determining whether the peer-to-peer delay mechanism is ~~instance~~ the transport-specific peer-to-peer delay mechanism or ~~is provided by~~ the CMLDS peer-to-peer delay mechanism apply here.

*Change item f) as follows:*

- f) The NMS reads and saves the multiple sets of (t3', t4') from ACC2 through the MIB. Then the NMS can compute the delay asymmetry, in units of time, as  $\frac{[(t4' - t4) \times \text{neighborRateRatio} - (t3' - t3)]}{2}$ . The NMS can use multiple sets of (t3, t4, t3', t4') to compute average values to get a more accurate result. The averaging method can be decided as required and is outside the scope of gPTP. If ACC1 and ACC2 are frequency synchronized, then neighborRateRatio is 1, and the delay asymmetry is  $\frac{[(t4' - t3') - (t4 - t3)]}{2}$ .



## Annex H

(informative)

## Bibliography

*Insert the following bibliography reference after the last bibliography reference:*

[B30] “Guidelines for Use of Extended Unique Identifier (EUI), Organizationally Unique Identifier (OUI), and Company ID (CID),” IEEE Registration Authority (<https://standards.ieee.org/products-programs/regauth/>).

## Annex Z

### COMMENTARY

**<<Editor's Note: This is a temporary Annex intended to record issues/resolutions thereof as the project proceeds. It also documents the revision history. It will be removed prior to Sponsor ballot, and should be ignored for the purpose of TG/WG ballot.>>**

#### Z.1 Revision history

##### Z.1.1 Revision 0.1

This is the initial version, prepared by the editor. It includes, in Clause 17, Annex Z.1 of IEC/IEEE 60802/D.1.2, which is the description of the hot standby feature and is taken from material previously contributed by Rodney Cummings. The Annex Z.1 material is modified as described in Z.1.1.1(a)(5) below. This initial version also includes maintenance items 263, 279, 281, 289, 292, 293, 294, 297, 307, and 308.

Clause 17 is a brand new clause. In Clause 17, track changes are used to show edits relative to Annex Z of 60802/D1.2, which is where this text was taken from. This is for the convenience of the reader, so that the reader can see what changes the editor has made. Various changes were necessary because the text in Annex Z of 60802 was written from the standpoint that this text would be part of 60802, not 802.1AS. It is intended that these track changes will be removed prior to SA ballot.

This initial draft is incomplete. It is planned to complete it before the initial TG ballot.

##### Z.1.2 Revision 0.2

This is a minor revision of the initial version, prepared by the editor. As with D0.1, it includes, in Clause 17, Annex Z.1 of IEC/IEEE 60802/D.1.2, which is the description of the hot standby feature and is taken from material previously contributed by Rodney Cummings. The Annex Z.1 material is modified as described in Z.1.2.1(a)(5) below. This initial version also includes maintenance items 263, 279, 281, 289, 292, 293, 294, 297, 307, and 308. However, the track changes of Clause 17 that were in D0.1 have been accepted, and several other minor edits were made (e.g., inserting the title line of Clause 11, replacing reference to “IEEE STD 802.1AS-2020” by “this standard”).

##### Z.1.2.1 Material needed for revision 0.2:

- a) Description and incorporation of the split functionality, and other relevant material, described in the following references:
  - 1) Feng Chen, *Segment Protection*, IEEE 802.1 presentation, March 2014 (available at <https://www.ieee802.org/1/files/public/docs2014/as-chen-segment-protection-0314-v01.pdf>).
  - 2) Rodney Cummings, *802.1AS Hot Standby Amendment: Scope Discussion*, Revision 1, IEEE 802.1 presentation, January 2020 (available at <https://www.ieee802.org/1/files/public/docs2020/new-cummings-as-hot-standby-scope-0120-v02.pdf>).
  - 3) Mike Potts, *IEEE 802.1AS Addendum for “Hot Standby” with multiple domain definition*, IEEE 802.1 presentation, 18 February 2020 (available at <https://www.ieee802.org/1/files/public/docs2020/dm-potts-as-hot-standby-multiple-domains-0220-v01.pdf>).
  - 4) Rodney Cummings, *802.1AS Hot Standby Amendment: 4 Domains*, Revision 3, IEEE 802.1 presentation, March 2020 (available at <https://www.ieee802.org/1/files/public/docs2020/dm-cummings-as-4-domains-0320-v04.pdf>).

- 5) Geoffrey M. Garner, *Initial Plans and Assumptions for 802.1ASdm*, IEEE 802.1 presentation, November 2, 2020 (available at <https://www.ieee802.org/1/files/public/docs2020/dm-garner-plans-1120-v00.pdf>).
- 6) Max Turner, *802.1ASdm - Redundancy*, IEEE 802.1 presentation, July 2021 (available at <https://www.ieee802.org/1/files/public/docs2021/dm-turner-ASredundancy-0721-v02.pdf>).
- b) Definition and specification of the NOT\_QUALITY state, along with any related functions.
- c) Addressing of all Editor's Notes.
- d) The editor must review the full Clause for completeness, consistency, and correctness.
- e) The following maintenance items must still be added: 278, 280, 283. Details are as follows:
  - 1) Discussion is needed for maintenance item #278. This maintenance item is to make the PICS and Clause 5 conformance statements consistent with respect to the interval setting state machines. These state machines were mandatory in 802.1AS-2011, but in 802.1AS-2020 the PICS and clause 5 are inconsistent on whether they are mandatory or optional. In initial discussion of the maintenance item, several participants indicated that the state machines are not needed by some profile standards (see 3.24) and applications, and that those applications and profiles that don't need them should not be required to include them, as the cost will be greater. The feature apparently is mainly needed for AVB. It was suggested that the feature could be optional in 802.1AS and mandatory in 802.1BA. However, during the presentation of *Insights and observations on TSN applied across ecosystems* on Monday, March 8, 2021 (available at <https://www.ieee802.org/1/files/public/docs2021/new-Kaltheuner-et-al-TSN-across-ecosystems-0321-v01.pdf>), it was indicated that if a feature that was previously mandatory, and needed by AVB, is now made optional, vendors may choose not to provide the feature in the future because the AVB market is smaller than, e.g., Industrial and Automotive. It also was mentioned, in the 802.1 maintenance meeting on Tuesday, March 9, 2021, that there is some chance that 802.1BA could be withdrawn if an editor for a revision is not found quickly because 802.1BA is approaching the deadline for re-affirmation. In any case, discussion is needed on how to proceed on this maintenance item.
  - 2) Maintenance item 280 is to change "time-aware Bridge" to "PTP Relay Instance" and "time-aware end station" to "PTP End Instance" in a MIB description field on p.363 of 802.1AS-2020.
  - 3) Maintenance item 283 is the following:  
 "IEEE Std 802.1AS-2020 uses both "domain number" and "domainNumber". The former is the concept, while the latter is the attribute, parameter, or name of the respective field. There are 38 instances of the former, and 66 instances of the latter. In contrast, IEEE Std 1588-2019 uses "domainNumber" almost exclusively; there is one instance of "domain numbers" and 106 instances of "domainNumber". It appears that at least some of the instances of "domain number" in 802.1AS-2020 should be "domainNumber. It would be useful to go through the document, check the usage of "domain number" and "domainNumber", and change one to the other where necessary."

The editor has reviewed the instances of "domain number" and "domainNumber." In the opinion of the editor, the instances of "domainNumber" are correct and need not be changed. Also, it is the editor's opinion that all the instances of "domain number" can be changed to "domainNumber" because all these uses pertain to the attribute "domainNumber" rather than the general concept "domain number." Can this change be made by a single editing instruction, or must there be a separate editing instruction (and change indicated by track changes) for each clause or subclause where "domain number" is used? If a single editing instruction is sufficient, where should it go? Note that two of the instances of "domain number" are in the MIB, on pp. 305 and 320 (in description fields). If a single editing instruction is used, how are the changes in the MIB handled (because, if the MIB is changed, the MIB revision and date also must change,

as well as the MIB file that will be attached). Note also that Maintenance item 280 also changes the MIB.

### Z.1.3 Revision 0.3

This revision version, prepared by the editor, addressess all comments against D0.2, except the following:

- a) Comment 95: Improvements to the formatting of Figure 18-1, Figure 18-2, and Figure 18-3, e.g., fixing the font size and making each figure fill the page as much as possible, are deferred to the next draft. However, all technical changes to the figures, required by comments against D0.2, are made.
- b) Comment 49: Updates to the MIB, and Clause 15, are deferred to working group ballot, when updates to managed objects will have been made. This is done so that all MIB updates can be made at once. Note that changing “domain number” to “domainNumber” in two places in the MIB also are deferred. However, “domain number” is changed to “domainNumber” in all other places, in accordance with comment 50.
- c) Comment 148: In deferring MIB changes, the part of comment 148 that requires MIB changes, i.e., addressing maintenance item 280 and the part of 283 that affects the MIB, is deferred.

Except for the maintenance items indicated above, all maintenance items listed in the frontmatter are included.

#### Z.1.3.1 Requests for contributions after D0.3

Contributions are requested on the following items:

- a) Contributions are needed on what new managed objects related to hot standby should be defined (see the editor’s note at the end of 18.5.1.3).
- b) Contributions are needed on how the split functionality (see Z.1.2.1) or equivalent should operate, and how it should be incorporated into the draft (see also the editor’s note at the beginning of 17.6).
- c) It has been suggested that whether a PTP Instance is synchronized also depends on the GM attributes, e.g., clockClass, clockAccuracy, offsetScaledLogVariance, etc. Contributions are requested to further discuss the details of this and make detailed proposals (see the editor’s note at the end of 18.4.3.3).

### Z.1.4 Revision 0.4

This revision version, prepared by the editor, addressess all comments against D0.3. The following comments against Revision D0.2 are still unaddressed:

- a) Comment 95: Improvements to the formatting of Figure 18-1, Figure 18-2, and Figure 18-3, e.g., fixing the font size and making each figure fill the page as much as possible, are deferred to the next draft. However, all technical changes to the figures, required by comments against D0.2, are made.
- b) Comment 49: Updates to the MIB, and Clause 15, are deferred to working group ballot, when updates to managed objects will have been made. This is done so that all MIB updates can be made at once. Note that changing “domain number” to “domainNumber” in two places in the MIB also are deferred. However, “domain number” is changed to “domainNumber” in all other places, in accordance with comment 50.
- c) Comment 148: In deferring MIB changes, the part of comment 148 that requires MIB changes, i.e., addressing maintenance item 280 and the part of 283 that affects the MIB, is deferred.

Except for the maintenance items indicated above, all maintenance items listed in the frontmatter are included.

#### Z.1.4.1 Requests for contributions after D0.4

Contributions are requested on the following items:

- a) Contributions are needed on what additional new managed objects related to hot standby should be defined (see the editor's note at the end of 18.5.1.3).
- b) Contributions are requested on the split functionality (see 18.5.3.4)
- c) The split functionality and, in general, the transfer of information from the primary and to the secondary PTP Instance do not address the transfer of time properties information (e.g., `offsetFromMaster`, `leap59`, `leap61`, etc.), nor of `ClockQuality` information). Should this be considered? (For that matter, the transfer of time properties and `ClockQuality` information from the `ClockSource` to the `ClockMaster`, either in Clause 17 or Clause 9, is not considered.) See the editor's note at the end of 18.5.3.4.
- d) It has been suggested that whether a PTP Instance is synchronized also depends on the GM attributes, e.g., `clockClass`, `clockAccuracy`, `offsetScaledLogVariance`, etc. Contributions are requested to further discuss the details of this and make detailed proposals (see the editor's note at the end of 18.4.3.3).

#### Z.1.5 Revision 0.5

This revision version, prepared by the editor, is the same as D0.4, except that track changes in Clause 17 have been accepted, the present subclause (Annex 1.5) has been added, and the draft number has been updated. The track changes were included in D0.4 for the convenience of readers. Since Clause 17 is a new clause, the only editing instruction for Clause 17 is the instruction at the very beginning to add the clause. Therefore, there was no danger of track changes in Clause 17 being confused with amendment editing instructions. However, D0.5 will be balloted (D0.4 was not balloted), and the practice is to not include track changes in balloted drafts; instead, changed text is indicated by change bars in the left margin. D0.5 does include change bars (as did D0.4, including Clause 17).

#### Z.1.6 Revision 0.6

This revision, prepared by the editor, addresses all comments against the current and previous revisions except for:

- a) Comment 95 against D0.2: Improvements to the formatting of Figure 18-1, Figure 18-2, and Figure 18-3, e.g., fixing the font size and making each figure fill the page as much as possible, are deferred to the next draft. However, all technical changes to the figures, required by comments against D0.2, are made.
- b) All comments related to the MIB. It was decided at the July 11, 2020 TSN TG meeting that revised MIB will not be included in P802.1ASdm. After repeated requests for a volunteer MIB editor, to which there were no responses, it was decided that an updated MIB will not be necessary in view of the fact that a YANG module will be included. As a result, Clause 15 has been removed from the draft.
- c) Comment 56 against D0.5. This comment indicates that the split functionality of 17.6.3.3.2.2 effectively duplicates the functionality of 17.6.3.2 Secondary grandmaster - both ensure (by different internal means) that the master-ports of both the primary and secondary PTP Instance communicate effectively the same time synchronization information to downstream slave ports.

Note that it is true that both the split functionality (17.6.3.3.2.2) and the Secondary GM section (17.6.3.2) both transfer timing from the primary to secondary PTP Instance. However, the two cases

are different. The split functionality is used when the secondary PTP Instance is not GM, and the transfer of the PortSyncSync structure, with translation, from primary to secondary allows the secondary domain state machines to work exactly as they do now in 802.1AS-2020. In contrast, in 17.6.3.2 the secondary is GM and, consequently, is expecting the information to come from the ClockMaster entity; that is why in this case the transfer occurs via the primary ClockSlave and secondary ClockMaster.

The currently formulation was chosen so that the existing 802.1AS-2020 state machines could be used as is, without modification. If a different formulation is chosen (e.g., a formulation that does not use the split functionality, does not use the secondary GM, or something entirely different, it needs to be determined exactly what changes are needed for the existing 802.1AS-2020 state machines.

More discussion of this topic is needed. Comments and/or contributions are requested.

### Z.1.7 Revision 0.7

This revision version, prepared by the editor, is the same as D0.6, except that track changes in Clause 17 have been accepted, the present subclause (Z.1.7) has been added, and the draft number has been updated. The track changes were included in D0.6 for the convenience of readers. Since Clause 17 is a new clause, the only editing instruction for Clause 17 is the instruction at the very beginning to add the clause. Therefore, there was no danger of track changes in Clause 17 being confused with amendment editing instructions. However, D0.7 will be balloted (D0.6 was not balloted), and the practice is to not include track changes in balloted drafts; instead, changed text is indicated by change bars in the left margin. D0.7 does include change bars (as did D0.6, including Clause 17).

### Z.1.8 Revision 0.8

This revision, prepared by the editor, addresses all comments against the current and previous revisions except for items a), b), and c) in Z.1.6 above. Note that item b) is no longer relevant as it has been decided not to include a MIB, and item c) is superseded by TG ballot comments against 0.7 related to the split functionality.

This revision has all changes to Clause 17 shown via track changes. This was possible because Clause 17 is a brand new clause and has no editing instructions other than to add the clause. This was done for the convenience of readers. Revision 0.8 will not be balloted; instead, a subsequent revision will be balloted.

All comments against D0.7 have been addressed, except for comment #69, which relates to adding a cross-reference to the cover page so that the title of the standard appears in the PDF bookmarks. This comment will be addressed in the next draft. Note that the resolutions to comments 3, 27, and 92 resulted in major changes to the HotStandbySystem state machine, and the resolution to comments 13, 51, 53, 54, and 55 resulted in major changes to the PtpInstanceSyncStatus state machine and addition of the ptp.

### Z.1.9 Revision 1.0

This revision version, prepared by the editor, is the same as D0.8, except that track changes in Clause 17 have been accepted, the present subclause (Z.1.9) has been added, and the draft number has been updated. The track changes were included in D0.8 for the convenience of readers. Since Clause 17 is a new clause, the only editing instruction for Clause 17 is the instruction at the very beginning to add the clause. Therefore, there was no danger of track changes in Clause 17 being confused with amendment editing instructions. However, D1.0 will be balloted (D0.8 was not balloted), and the practice is to not include track

changes in balloted drafts; instead, changed text is indicated by change bars in the left margin. D1.0 does include change bars (as did D0.8, including Clause 17).

### Z.1.10 Revision 1.1

This version, prepared by the editor, incorporates all comments against D1.0, but with the following exceptions:

- a) For comments 1 and 141, the approval date of and the link to the modified PAR cannot be inserted until the PAR approved by the IEEE SASB and placed on the IEEE SA web site. As of the date of completion of this draft, the information provided to the editor indicates that the modified PAR has been approved by NesCom but is awaiting approval by the SASB via letter ballot.
- b) Comments 3 and 144 ask that 802.1ASdr be incorporated into the draft. This will be done after P802.1ASdr completes SA ballot. In addition, comments 58 and 59 requests changes that depend on P802.1ASdr; these changes will be made after P802.1ASdr completes SA ballot.
- c) Comments 47 and 48 ask that permission be obtained for two definitions that were copied from IEEE Std 1588-2019. It is the understanding of the editor that respective forms related to this must be submitted to the IEEE SA by the 802.1 Chair. As of the date of completion of this draft, it is expected that the forms will be submitted. The editor has not yet been informed what additions, if any, are needed for the draft; however, any needed additions will be made when the editor is informed of this.

In addition, comment 95 against D0.2, Improvements to the formatting of Figure 18-1, Figure 18-2, and Figure 18-3, e.g., fixing the font size and making each figure fill the page as much as possible, is deferred to the next draft. Finally, a YANG module is not yet available; however, it is expected that this will be available, and will be incorporated in the draft along with any other material needed related to P802.1ASdn.

### Z.1.11 Revision 1.2

This version, prepared by the editor, incorporates all comments against D1.1, with the exception that the YANG Clause and module for hot standby does not reflect all changes and additions to managed objects due to comments against D1.1. In addition, the items in Z.1.10 above are still to be addressed.