

Numerical Differential Equations

Template from Leon Lee @leon0241 on GitHub.

1 Preliminaries

Definition 1.1: Studying Differential Equations

Analytically - to find a solution in terms of y ; **qualitatively** - study the properties based on analytical and/or graphical techniques; **computationally** - find an approximation solution numerically.

Recall 1.2: ODE Terminology

Ordinary has one independent variable. **Autonomous** does not depend on the independent variable. Linear if in form

$$\frac{dy}{dt} = a(t)y + b(t)$$

where $a(t)$ and $b(t)$ do not need to be linear in t .

Example 1.3: Examples of (Non)-Linear Equations

$$(1) \frac{dy}{dx} = kx^2y - cx^3, \quad (2) x' = e^{-t\beta}x, \quad (3) \dot{u} = \cos(\omega t)u - \sin(\omega t).$$

All the above are linear equations with y , x and u being the dependent variables and parameters k, c, β, ω . Below is a *linear system*

$$\dot{x} = u, \quad \dot{u} = -\omega^2x \quad \text{for an unspecified independent var.}$$

The below are *non-linear* equations

$$\frac{du}{dt} = u^2, \quad \frac{dx}{dt} = kx^2t - cx^3$$

or (Duffing's oscillator) a *non-autonomous, non-linear* system

$$\frac{dx}{dt} = u, \quad \frac{du}{dt} = -\beta x - \alpha x^3 - \delta u - \gamma \cos(\omega t).$$

Recall 1.4: Lipschitz Continuity

A function f is Lipschitz continuous if there exists some constant $L > 0$ (the Lipschitz constant) such that

$$|f(t, u) - f(t, v)| \leq L|u - v|, \quad \text{for all } t \text{ and } u, v \in D.$$

Remark 1.5: Lipschitz Continuity: Easier

If f is continuous in t and continuously differentiable (that is, the derivative exists and is itself continuous) in y then

$$\frac{\partial f}{\partial y} \text{ is continuous} \implies L = \sup_D \left| \frac{\partial f}{\partial y} \right|.$$

Proof; by the mean value theorem. MVT states that if f is continuous and differentiable over some $[a, b]$ then there exists c such that

$$f'(c) = \frac{f(b) - f(a)}{b - a} \implies f(b) - f(a) = f'(c)(b - a).$$

It immediately follows how we can find L .

Theorem 1.6: Picard's Theorem (Existence/Uniqueness)

Let f be a cts func of its arguments in a region of the plane containing the rectangle

$$D = \{(t, y) \mid t_0 \leq t \leq T, \quad |y - y_0| \leq K\}$$

where T, K are some constants. Further suppose f is Lipschitz cts. Domain D is closed and bounded so there exists some

$$M_f = \max\{|f(t, u)| : (t, u) \in D\}.$$

If $M_f(T - t_0) \leq K$ then there exists a unique, continuously differentiable function $y(t)$ defined on $[t_0, T]$ such that

$$\frac{dy}{dt} = f(t, y), \quad y(t_0) = y_0.$$

Note. This can be condensed and/or rewritten, do it.

Theorem 1.7: Local Existence and Uniqueness

Suppose $f : \mathbb{R} \times \mathbb{R}^d \rightarrow \mathbb{R}^d$ is cts and has cts partial derivatives in w.r.t all variables in some neighbourhood of (t_0, y_0) . Then there some interval $I = (t_0 - \delta, t_0 + \delta)$ such that some y satisfies

$$\frac{dy}{dt} = f(t, y), \quad y(t_0) = y_0.$$

Recall 1.8: Slope Fields

Consider the ODE $y' = f(t, y)$. Any point $f(t_0, y_0)$ gives the slope of y ; do this for all points of the plane and draw, typically with arrow.

Recall 1.9: Taylor Series Expansion

The standard definition of a Taylor series for some function f centered about a is given

$$\sum_{n=0}^{\infty} \frac{f^{(n)}(a)}{n!} (x - a)^n.$$

Of course this requires f to be infinitely differentiable. If f is a most $N + 1$ times differentiable on some interval $[a, b]$, we have the Taylor remainder theorem (remark $x_0 \in (a, b)$ and τ between x_0 and x)

$$f(x) = \sum_{n=0}^N \frac{f^{(n)}(x_0)}{n!} (x - x_0)^n + \frac{f^{(N+1)}(\tau)}{(N + 1)!} (x - x_0)^{N+1}.$$

Example 1.10: Non-Autonomous Equations

2 Euler’s Method and Taylor Series Methods

Definition 2.1: Euler’s Method

y_{n+1} = y_n + hf(t_n, y_n), where y_i ≈ y(t_i).

The algorithm takes in h, f, (t_0, y_0) and N. Remark N steps gives N + 1 time points t_0, t_1, ..., t_N = t_0 + Nh. Remark h = (b - a)/N.

Example 2.2: From Where?

Euler’s method is exactly a first order Taylor series method. Expand the Taylor series (with remainder, first order) about t_0 at t_1,

y(t_1) = y(t_0) + y'(t_0)(t_1 - t_0) + 1/2 y''(tau)(t_1 - t_0)^2.

where tau is between t_1 and t_0. We know nada of tau, but we approx the sol by last term die (h^2 -> 0). With t_1 = t_0 + h we are done.

Lemma 2.3: An Inequality

Sequence (v_n) of non-neg, for A > 1 and B > 0 constants.

v_{n+1} ≤ Av_n + B ⇒ v_n ≤ A^n v_0 + (A^n - 1)/(A - 1) B.

Theorem 2.4: Convergence of Euler’s Method

Consider y'(t) = f(t, y) with y(a) = y_0. Suppose unique, twice diff. sol y(t) on [a, b]. Further suppose f is cts everywhere with cts, bounded partial derivatives

|∂f/∂y| ≤ L, L > 0.

Then for some constant D, solution y_n at t_n satisfies

|e_n| = |y_n - y(t_n)| ≤ Dh where D = e^{(b-a)L} M / (2L) : |y''(t)| ≤ M.

It is stressed that D does not depend on N or h, but only on features of the problem and the solution. If h is not small enough, the error compounds and can become very large (recall examples).

Definition 2.5: The Big O (Landau) Notation

A quantity is (of order) O(h^p) if it decays at least as quickly as h^p when h is small enough. More formally, we write z = O(h^p) if there exist h_0, C ∈ ℝ_+ such that |z| ≤ Ch^p for all 0 < h ≤ h_0.

We are concerned with p. Say ... is of order p. Conventionally

z = O(h^{p+1}) ⇒ z = O(h^p).

Definition 2.6: Flow Maps

Idea “we look at a set of initial point and describe the journey to their endpoints”.

Φ_{t_0, h}(y_0) = y(t_0 + h; t_0, y_0)

that we read “given t_0, h”. Determining the flow map is equivalent to solving the ODE; idea being to replace the map by an approx. i.e

ḡ_{t, h}(y) = y + hf(t, y) is Euler’s method as a flow map approx.

We have, from outside the course (at this moment),

linear ODE y' = Ay ⇒ Φ_t(y) = e^{tA} y.

Recall 2.7: How to Differentiate

Suppose z = f(x, y) and x = g(t), y = h(t). Then, from SVCDE,

dz/dt = ∂z/∂x dx/dt + ∂z/∂y dy/dt.

So if y' = f(t, y) then

y'' = d(f(t, y))/dt = ∂f/∂t dt/dt + ∂f/∂y dy/dt = f_t + f_y · f.

Example 2.8: From Where? Remastered

(Second order Taylor series method.) Recall

y(t_0 + h) = y(t_0) + y'(t_0)h + 1/2 y''(t_0)h^2 + 1/6 y'''(t_0)h^3 + H.O.T.

Then a second order method instead drops the y''' term (consider Taylor remainder if wanted) and we know y' = f and y'' = f_t + f_y · f. So, by considering a typical ODE, we have

y_{n+1} ≈ y_n + hf(t_n, y_n) + 1/2 h^2 (f_t(t_n, y_n) + f_y(t_n, y_n) · f(t_n, y_n)).

Higher order Taylor series methods can of course be used. General case is a pain to express, but in practice is it simple (implicit diff).

Definition 2.9: One-Step Methods

Methods that approximate the solution through iteration of an approximate flow map (i.e Euler, Taylor) are one-step methods.

3 Convergence of One-Step Methods

Definition 3.1: Global Error

The *global error* after n steps is the diff. between approx and exact solution; $e_n := y_n - y(t_n)$. A good approx. will produce a global error that is small in norm at each step; i.e

$$\max_{n=0,\dots,N} \|e_n\| \leq \delta, \quad (\text{user specified typically } \delta = 10^{-4}).$$

Definition 3.2: Convergent Methods

For any given T , method *convergent* if

$$\lim_{h \rightarrow 0} \max_{n=0,\dots,N} \|e_n\| = 0.$$

Remark this is **worst case error**. As $h \rightarrow 0$ we have $N \rightarrow \infty$.

Remark 3.3: Notational Remarks

We use $y(t+h) = \Phi_h(y(t))$ i.e let's look at the flow map of y from t to $t+h$. Determining the flow map is equivalent in many ways to solving the ODE. We are given a *one step numerical method* that we describe with $\Psi_h(y)$; an approximation of the flow map.

Definition 3.4: Local Error

The *local error* of a one-step num. method is the difference between the flow map and its approximation i.e.

$$le(y, h) = \Psi_h(y) - \Phi_h(y).$$

Definition 3.5: Consistent (at Order p) Methods

Suppose some method satisfies

$$\|le(y, h)\| \leq Ch^{p+1}$$

for $p \geq 1$ and some constant C depending on y and its derivatives. Then the method is *consistent* at order p with $p \geq 1$.

Definition 3.6: Stable Methods

Suppose method satisfies an h -dependent Lipschitz condition on D (below). Then the method is *stable*.

$$\|\Psi_h(u) - \Psi_h(v)\| \leq (1 + h\bar{L})\|u - v\|, \quad \forall u, v \in D.$$

Theorem 3.7: Convergence of One-Step Methods

Consider some ODE and some one-step numerical method Ψ_h that is both *stable* and *consistent* c Then the global error satisfies

$$\max_{n=0,\dots,N} \|e_n\| = \mathcal{O}(h^p)$$

i.e any stable and consistent one step method will converge.

Construction of More One-Step Methods

Remark 3.8: The Idea

$$\begin{aligned} \frac{dy}{dt} &= f(y), \quad y(t_0) = y_0; \\ \int_t^{t+h} \frac{dy}{dt} &= \int_t^{t+h} f(y) \implies y(t+h) - y(t) = \int_t^{t+h} f(y). \end{aligned}$$

How do we approximate the RHS?

Example 3.9: Euler's Method

$$\int_t^{t+h} f(y) \approx hf(y(t)) \implies y(t+h) - y(t) \approx hf(y(t));$$

i.e Euler's method; $y_{n+1} - y_n = hf(y_n)$, an explicit method.

Example 3.10: Trapezoidal Rule

$$\begin{aligned} \int_t^{t+h} f(y(\tau))d\tau &\approx \frac{h}{2}(f(y(t)) + f(y(t+h))) \\ \implies y_{n+1} &= y_n + \frac{h}{2}(f(y_n) + f(y_{n+1})), \quad \text{an implicit method.} \end{aligned}$$

Definition 3.11: Polynomial Interpolation

Let \mathbb{P}_s denote the space of real polynomials degree $\leq s$. Set of s *abscissa* c_i , correspond. data g_i , there exists a unique polynomial $P(x) \in \mathbb{P}_{s-1}$ through them. The *Lagrange* interpolating polynomials

$$\ell_i(x) = \prod_{j=1, j \neq i}^s \frac{x - c_j}{c_i - c_j}, \quad P(x) = \sum_{i=1}^s g_i \ell_i(x).$$

Definition 3.12: Quadrature Rules

A method to approx. a definite integral of one independent variable is a *numerical quadrature rule*. A quadrature rule has order p if it exactly integrates any polynomial in \mathbb{P}_{p-1} . We have $p \geq s$ always.

Theorem 3.13: Quadrature Formula

$$\begin{aligned} b_i &= \int_0^1 \ell_i(x)dx, \\ \int_{t_0}^{t_0+h} g(t)dt &= \int_0^1 h \cdot g(t_0 + hx)dx \approx h \sum_{i=1}^s b_i g(t_0 + hc_i). \end{aligned}$$

Definition 3.14: Collocation Methods

We construct a polynomial that passes through y_0 and agrees with the ODE at s nodes; then the numerical solution is the value of the polynomial at y_1 . Below are necessary conditions;

$$\begin{aligned} u(t_0) &= y_0, \\ u'(t_0 + c_i h) &= f(u(t_0 + c_i h)), \quad i = 1, \dots, s. \end{aligned}$$

Example 3.15: Construction of Collocation Methods

We construct u' . Consider F_i to be the values of the (as yet undetermined) polynomial at the nodes i.e $F_i := u'(t_0 + c_i h)$. Then

$$u'(t) = \sum_{i=1}^s F_i \ell_i \left(\frac{t - t_0}{h} \right)$$

and by integrating over $[t_0, t_0 + c_i h]$ with a change of variables $x = (t - t_0)/h$ we obtain

$$u(t_0 + c_i h) = y_0 + h \sum_{j=1}^s F_j \int_0^{c_i} \ell_j(x)dx, \quad i = 1, \dots, s.$$

By defining a_{ij} and b_i , we substitute this result into our required collocation conditions and achieve the coupled non-linear system (remark - it is explicit).

$$\begin{aligned} a_{ij} &:= \int_0^{c_i} \ell_j dx, \quad b_i := \int_0^1 \ell_i(x)dx, \quad i, j = 1, \dots, s; \\ F_i &= f \left(y_0 + h \sum_{j=1}^s a_{ij} F_j \right), \quad y_{n+1} = y_n + h \sum_{i=1}^s b_i F_i. \end{aligned}$$

Lemma 3.16: Collocation: A Continuous Approximation

We obtain a cts approx. of the sol $u(t)$ on each interval $[t_n, t_{n+1}]$.

Remark 3.17: Best Accuracy?

The Gauss-Legendre collocation methods uses nodes c_1 for

$$\begin{aligned} c_1 &= \frac{1}{2} \\ c_1 &= \frac{1}{2} - \frac{\sqrt{3}}{6}, c_2 = \frac{1}{2} + \frac{\sqrt{3}}{6}, \\ c_1 &= \frac{1}{2} - \frac{\sqrt{15}}{10}, c_2 = \frac{1}{2}, c_3 = \frac{1}{2} + \frac{\sqrt{15}}{10}. \end{aligned}$$

4 Runge-Kutta Methods

Definition 4.1: Runge-Kutta Methods

$$Y_i = y_n + h \sum_{j=1}^s a_{ij} f(Y_j), \quad i = 1, \dots, s$$

$$y_{n+1} = y_n + h \sum_{i=1}^s b_i f(Y_i).$$

Here s num. stages, b_i weights, a_{ij} internal coefficients. This the *same formula* as collocation method but w/o restriction on $a/b/c$. It generates a discrete flow-map approximation

$$\Psi_h(y) = y + h \sum_{i=1}^s b_i f(Y_i(y, h)).$$

Remark 4.2: Internal Consistency

We have $c_i = \sum_{j=1}^s a_{ij}$ (assumed in the notes; not necessarily true).

Definition 4.3: Butcher Tables

$$\begin{array}{c|c} C & A \\ \hline & b^T \end{array} \quad \text{i.e.} \quad \begin{array}{c|ccc} c_1 & a_{11} & \cdots & a_{1s} \\ \vdots & \vdots & & \vdots \\ c_s & a_{s1} & \cdots & a_{ss} \\ \hline & b_1 & \cdots & b_s \end{array}$$

If the matrix $A = (a_{ij})$ is strictly lower triangular then the method is *explicit*. Otherwise it is implicit.

Example 4.4: Examples Using the Butcher Table

The implicit midpoint rule

$$Y_1 = y_n + \frac{h}{2} f(Y_1), \quad y_{n+1} = y_n + h f(Y_1),$$

$$\text{i.e. } y_{n+1} = y_n + h f\left(\frac{y_{n+1} + y_n}{2}\right).$$

So in this case we have $s = 1$, $a_{11} = \frac{1}{2}$ and $b_i = 1$ and also $c_i = \frac{1}{2}$

$$\begin{array}{c|c} 1/2 & 1/2 \\ \hline & 1 \end{array} \quad \text{clearly implicit (see formula).}$$

The **four-stage** Runge-Kutta Method has

$$\begin{cases} Y_1 = y_n, \\ Y_2 = y_n + \frac{h}{2} f(Y_1), \\ Y_3 = y_n + \frac{h}{2} f(Y_2), \\ Y_4 = y_n + h f(Y_3), \end{cases} : y_{n+1} = h \left(\frac{1}{6} f(Y_1) + \frac{1}{3} f(Y_2) + \frac{1}{3} f(Y_3) + \frac{1}{6} f(Y_4) \right).$$

$$\begin{array}{c|ccc} 0 & & & \\ 1/2 & 1/2 & & \\ 1/2 & 0 & 1/2 & \\ 1 & 0 & 0 & 1 \\ \hline & 1/6 & 1/3 & 1/3 & 1/6 \end{array} \quad \text{consider } Y_i = y(t_n + c_i h).$$

Recall 4.5: Continuously Differentiable and Semisimple

Notation $C^{(n)}$ means function has a cts n -th derivative. Any eigenvalue λ is semisimple if the size of its largest Jordan block is 1.

Definition 4.6: The Jacobian ... And More

the Jacobian $f' = \left(\frac{\partial f_i}{\partial y_j} \right)$, $i \leq i, j \leq d$, notation $M = (m_{ij})$.

$$f'' = \left(\frac{\partial^2 f_i}{\partial y_j \partial y_k} \right) : \text{operator } \sum_{j=1}^d \sum_{k=1}^d \frac{\partial^2 f_i}{\partial y_j \partial y_k} a_j b_k \text{ maps } a, b \in \mathbb{R}^d \rightarrow \mathbb{R}.$$

There is a long spiel but we achieve, in summary,

$$\begin{aligned} y' &= f, \\ y'' &= f' f, \\ y''' &= f''(f, f) + f' f' f; \end{aligned} \quad y(t), f(y).$$

Recall Taylor series expansion

$$y(t+h) = y(t) + h y'(t) + \frac{1}{2} h^2 y''(t) + \frac{1}{6} h^3 y'''(t) + \mathcal{O}(h^4)$$

and the corresponding flow map

$$\Phi_h(y) = y + h f + \frac{h^2}{2} f' f + \frac{h^3}{6} (f''(f, f) + f' f' f) + \mathcal{O}(h^4).$$

The notes gives how to calculate the local error for;

$$\text{Euler's method: } le(y, h) = \frac{h^2}{2} f' f + \mathcal{O}(h^3),$$

$$\text{Trapezium rule: } le(y, h) = -\frac{1}{12} (f''(f, f) + f' f' f) h^3 + \mathcal{O}(h^4).$$

Theorem 4.7: Convergence of RK Methods

Respectively for $p = 1, 2, 3$ we require the following conditions;

$$\begin{aligned} \sum_{i=1}^s b_i &= 1, \\ \sum_{i=1}^s b_i &= 1, \quad \sum_{i=1}^s b_i c_i = \frac{1}{2}, \\ \sum_{i=1}^s b_i &= 1, \quad \sum_{i=1}^s b_i c_i = \frac{1}{2}, \quad \sum_{i=1}^s b_i c_i^2 = \frac{1}{3}, \quad \sum_{i=1}^s \sum_{j=1}^s b_i a_{ij} c_j = \frac{1}{6}. \end{aligned}$$

Stability of Runge-Kutta Methods

Remark 4.8: The Idea

We look at the local sol. behaviour in the neighbourhood of fixed points; assume ODE autonomous, often about origin. Key takeaway; any linear ODE system and its' RK approx. can be decomposed into scalar linear problems involving the eigenvalues as coeffs.

Definition 4.9: Fixed Points of an ODE

Point y^* if $f(y^*) = 0$; that is, $y(t) \equiv y^*$ when $y(0) = y^*$.

Definition 4.10: (Spurious) Fixed Points

Fixed points of Φ_t identical to ODE; \mathcal{F} . We have fixed points of num. method \mathcal{F}_h depend on h ; $\mathcal{F}_h = \{y \in \mathbb{R}^d : \Psi_h(y) = y\}$. In general (for RK methods) $\mathcal{F} \subseteq \mathcal{F}_h$, for Euler, implicit mp we have $\mathcal{F} = \mathcal{F}_h$.

Spurious point: $y^* \in \mathcal{F}_h$ but $y^* \notin \mathcal{F}$. How do we identify? As $h \rightarrow 0$, all spurious tend to infinity; so test multiple h .

Definition 4.11: Formal Definitions for Stability

- (i) y^* stable if for all $\epsilon > 0$, exists $\delta > 0$ s.t $\|y_0 - y^*\| < \delta \implies \|y(t; y_0) - y^*\| < \epsilon, \quad t > 0$;
- (ii) asymptotically stable if stable and exists $\gamma > 0$ s.t $\lim_{t \rightarrow \infty} \|y(t; y_0) - y^*\| = 0$.

This can be analogously translated for a general map $\Psi : \mathbb{R}^d \rightarrow \mathbb{R}^d$ by replacing y_0 with $y^n(y_0)$ where a fixed point is $y^* = \Psi(y^*)$.

Theorem 4.12: The Linearisation Theorem

Consider $\mathbb{R}^d : \frac{dy}{dt} = By + F(y)$ where $y(0) = y_0$; B is a constant matrix F is C^1 in a neighbourhood of y_0 with $F(0) = F'(0) = 0$ where F' is the Jacobian of F . If B 's evals all negative real part then $y = 0$ is an asymptotically stable fixed point. If any positive, y^* unstable.

Lemma 4.13: Linearisation: The Sequel

Suppose $\frac{dy}{dt} = f(y)$ and f is C^2 . If evals of $J = f'(y^*)$ (where f' is Jacobian of f) all strictly lie in the left complex half-plane then y^* is asymptotically stable. If any one in the right, y^* is unstable.

Definition 4.14: Spectral Radius

Denoted $\rho(K)$ of matrix K ; "the smallest circle centered at origin enclosing all eigenvalues of K ". In practice; $\max \|\lambda_i\|$ for $\lambda_i \in \sigma(K)$.

Theorem 4.15: Linear Iteration

Consider $z_n = \|K^n y_0\|$. Then

- (i) $z_n \rightarrow 0$ for all y_0 iff $\rho(K) < 1$;
- (ii) $z_n \rightarrow \infty$ for some y_0 iff $\rho(K) > 1$;
- (iii) if $\rho(K) = 1$ and all eigenvalues on the unit circle are semisimple then $\{z_n\}$ remains bounded as $n \rightarrow \infty$.

Theorem 4.16: General (Smooth) Maps

Given a smooth (C^2) map Ψ , we consider the fixed point y^* for the iteration $y_{n+1} = \Psi(y_n)$. It

- (i) is stable if $\rho(\Psi'(y^*)) < 1$,
- (ii) is unstable if $\rho(\Psi'(y^*)) > 1$,
- (iii) and we cannot say if we have inequality with 1.

5 Linear Stability of Numerical Methods

Remark 5.1: The Idea

We consider linearised $\frac{dy}{dt} = By$ where B is a matrix with a basis of eigenvectors (i.e can be written as $\text{diag}(\lambda_i) = \Lambda$).

Theorem 5.2: Stability Function of RK Methods

All Runge-Kutta methods applied to $y' = \lambda y$ can be written as $y_{n+1} = R(h\lambda)y_n$ where $R(\mu = h\lambda) = P(\mu)/Q(\mu)$, $R :=$ “stab func”.

Theorem 5.3: Stability of The Origin

Consider $y' = By$ where B is $d \times d$ and has a basis of eigenvectors. Consider applying some RK method. Then the RK method has a stable (asymptotically stable) fixed point at the origin if and only if the same method has a stable (asymptotically) fixed point at the origin to $\frac{dy}{dt} = \lambda_i y$ for all λ_i .

Lemma 5.4: Diagonalisable with Stability Function

Consider $\frac{dy}{dt} = By$ with diagonalisable B . Consider an RK method with stab. func R . The origin is stable for this RK method at step-size h if and only if $\hat{R} = |R(\mu)| < 1$ for all $u = h\lambda_i$. Unstable if > 1 .

Example 5.5: Examples of Stability Functions

Recall we are requiring $\hat{R}(\mu) = |R(\mu)| < 1$ where $\mu = \lambda h$ (?).

- (i) Euler’s method $\hat{R}(\mu) = |1 + \mu|$; i.e a disc of radius 1 centered at the point -1 .
- (ii) Trapezium rule;

$$\hat{R}(\mu) = \left| \frac{1 + \mu/2}{1 - \mu/2} \right| < 1$$

i.e when $\mu/2$ is closer to -1 than 1 i.e the entire left complex half-plane (and thus is A-stable).

- (iii) The implicit Euler $z_{n+1} = z_n + h\lambda z_{n+1}$ has $\hat{R}(\mu) = |1 - \mu|^{-1}$ i.e the exterior of the disk of radius 1 centered at 1 (so included entire left half-plane and is thus A-stable).

What about RK4? We have

$$R(\mu) = 1 + \mu + \frac{1}{2}\mu^2 + \frac{1}{6}\mu^3 + \frac{1}{24}\mu^4$$

and agrees with the Taylor series expansion of $\exp(h\lambda)$ through to fourth order. It is tricky to solve.

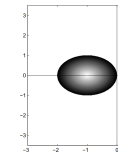


Figure 1: RK1

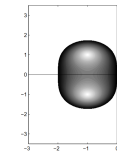


Figure 2: RK2

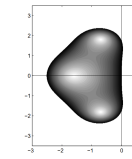


Figure 3: RK3

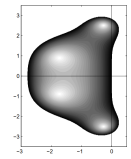


Figure 4: RK4

Definition 5.6: A-Stable Methods

If the stability region includes the entire left hand plane the method is *A-stable*. On linear systems, if the origin is stable for the ODE then it is also stable for the numerical method, regardless of stepsize.