

## 1 Preliminaries

For population  $N(t)$ , and max supported pop  $N^*$ . Let  $n = \frac{N}{N^*}$ .

**Logistic model:**

$$\frac{dn}{dt} = rn(1-n) \quad \left| \begin{array}{l} \text{General} \\ \text{Solution} \end{array} \right. : n(t) = \frac{ce^{rt}}{1+ce^{rt}}.$$

**Logistic model (with predation):**  $\frac{dn}{dt} = rn(1 - \frac{n}{s}) - \frac{n^2}{1-n^2}$

**Lotka-Volterra predator-prey model:**  $n$  = prey population,  $p$  = predator population:

$$\frac{dn}{dt} = rn(1-p), \quad \frac{dp}{dt} = p(1-n)$$

**Converting non-autonomous systems to autonomous:**

For non-autonomous system of  $d$  equations:

$$\frac{dy_1}{dt} = f_1(t, y_1, \dots, y_d),$$

$$\frac{dy_2}{dt} = f_2(t, y_1, \dots, y_d),$$

$\vdots$

$$\frac{dy_d}{dt} = f_d(t, y_1, \dots, y_d).$$

We can make it into an autonomous one by introducing the variable  $y_{d+1} \equiv t$  in place of  $t$ , (thus creating an autonomous system of  $d+1$  variables) and introduce new indep var  $s$  such that  $ds/dt = 1$ :

$$\frac{dy_1}{ds} = f_1(t, y_1, \dots, y_{d+1}),$$

$$\frac{dy_2}{ds} = f_2(t, y_1, \dots, y_{d+1}),$$

$\vdots$

$$\frac{dy_d}{ds} = f_d(t, y_1, \dots, y_{d+1}),$$

$$\frac{dy_{d+1}}{ds} = 1.$$

**Picard's Thm:** let the function  $f(\cdot, \cdot)$  be a continuous function of its arguments in a region of the plane containing the rectangle  $D = \{(t, y) : t_0 \leq t \leq T, |y - y_0| \leq K\}$ ,  $T, K > 0$ . Suppose  $\exists L$  (Lipschitz constant) such that  $|f(t, u) - f(t, v)| \leq L|u - v|$  whenever  $(t, u), (t, v) \in D$ . Since  $D$  closed & bounded,  $\exists M_f > 0$  such that  $M_f = \max\{|f(t, u)| : (t, u) \in D\}$ . Assume  $M_f(T - t_0) \leq K$ . Then there exists a unique continuously differentiable function  $t \mapsto y(t)$ , defined on  $[t_0, T]$ , such that  $\frac{dy}{dt} = f(t, y), y(t_0) = y_0$ .

*Note:*  $\frac{\partial f}{\partial y}$  being continuous  $\Rightarrow$  existence of a Lipschitz constant,

We can take  $L = \sup_D \left| \frac{\partial f}{\partial y} \right|$ .

**Thm 1.3.2 (Local Existence / Uniqueness of Solutions for ODE Systems):** Suppose  $f : \mathbb{R} \rightarrow \mathbb{R}^d \times \mathbb{R}^d$  is continuous and has continuous partial derivatives w.r.t. all components of the dependant var  $y$  in a neighbourhood of the point  $(t_0, y_0)$ . Then there is an interval  $I = (t_0 - \delta, t_0 + \delta)$  containing a unique function  $y$  (continuously differentiable on  $I$ ) that satisfies  $\frac{dy}{dt} = f(t, y), y(t_0) = y_0$ .

## 2 Euler's Method and Taylor Series Methods

**Euler's Method:** for finding approximate solutions to  $\frac{dy}{dt} = f(t, y)$  in time interval  $t \in [a, b]$ .

Let  $t_m = a + mh$ ,  $h \in \mathbb{R}$  small. Start with  $y_0 = y(t_0) = y(a)$  and use iteration function  $y_{n+1} = y_n + hf(t_n, y_n)$ .  $y_i \approx y(t_i)$ .

**Error:**  $e_n = |y_n - y(t_n)|$ .

**Lemma 2.6.1:** Suppose a given sequence of non-negative numbers  $(v_n)$  satisfies  $v_{n+1} \leq Av_n + B$ ,  $A > 1, B > 0$ . Then for

$$n = 0, 1, 2, \dots, \quad v_n \leq A^n v_0 + \frac{A^n - 1}{A - 1} B.$$

**Bound on error:**  $|e_n| \leq e^{(b-a)L} \frac{M}{2L} h = Dh$  where  $h$  is the timestep,  $M$  bounds  $|y''(t)|$ , and  $L$  bounds  $\left| \frac{\partial y}{\partial t} \right|$ .

**Thm 2.6.1** Consider the IVP  $\frac{dy}{dt} = f(t, y), y(a) = y_0$ . Suppose  $\exists$  unique, twice-differentiable function  $f$ , continuous everywhere with continuous, bounded partial derivative  $\left| \frac{\partial f}{\partial y} \right| < L$  with  $L > 0$ . Then for  $n = 0, 1, \dots, N$ , and some  $D > 0$ , the solution  $y_n$  given by Euler's method at  $t_n$  satisfies  $e_n = |y_n - y(t_n)| \leq Dh$  where  $h = (b - a)/N, t_n = a + hn$ .

**Big O (Landau) notation:** If method is  $O(h^p)$  then the error decays at least as quickly as  $h^p$  (for small  $h$ ).

$$(\exists h_0, C > 0) : (\forall 0 < h < h_0), |z| \leq Ch^p \Rightarrow z = O(h^p)$$

Say  $z$  is of order  $p$ .

**The Flow Map:** Consider IVP  $\frac{dy}{dt} = f(t, y), y(a) = y_0$  with unique solution in  $t \in [a, b]$ . Starting at arbitrary  $t_0 \in [a, b], y_0$  we may see where  $y(t)$  ends up, denoted  $y(t; t_0, y_0)$ .

Fix  $t_0$  and look at the **flow map**  $\Phi_{t_0, h}(y_0) = y(t_0 + h; t_0, y_0)$ . (actually a family of maps parameterised by  $h$ ).

**Numerical methods approximate flow maps:** Euler's method approximates flow map with  $\hat{\Phi}_{t, h}(y) = y + hf(t, y)$ .

**One-step methods:** approximate the solution through the iteration of an approximated flow map.

**Constructing Taylor series methods:**

Start with Taylor series:

$$y(t_0 + h) = y(t_0) + y'(t_0)h + \frac{1}{2}y''(t_0)h^2 + \frac{1}{6}y'''(t_0)h^3 + \dots$$

Also  $y'(t) = f(t, y) \Rightarrow$

$$y'' = \frac{d}{dt} f(t, y) = \frac{\partial}{\partial t} f(t, y) \frac{dt}{dt} + \frac{\partial}{\partial y} f(t, y) \frac{dy}{dt}$$

$$= \frac{\partial}{\partial t} f(t, y) + \frac{\partial}{\partial y} f(t, y) y' = f_t + f_y f. \quad (\text{By chain rule})$$

$$\Phi_{t, h}(y) = y + hf(t, y) + \frac{1}{2}h^2(f_t(t, y) + f_y(t, y)f(t, y)) + \frac{1}{6}y'''h^3 + \dots$$

Which we can truncate to get the 2nd order Taylor series method

$$\hat{\Phi}_{t, h}(y) = y + hf(t, y) + \frac{1}{2}h^2(f_t(t, y) + f_y(t, y)f(t, y)).$$

## 3 Convergence of One-Step Methods

**Def 3.1.2 (Convergence):** A method is said to be **convergent** iff for any  $T$ ,

$$\lim_{h \rightarrow 0} \max_{n=0, 1, \dots, N} \|e_n\| = 0.$$

**Def 3.2.1 (Local Error):** The **local error** of a one-step method is the difference between the flow map  $\Phi_h$  and it's discrete approximation  $\Psi_h$

$$le(y, h) = \Psi_h(y) - \Phi_h(y).$$

It measures how much error is introduced in a single timestep of size  $h$ .

**Def 3.2.2 (Consistency):** Suppose the local error for our method satisfies

$$\|le(y, h)\| \leq Ch^{p+1}$$

where  $C$  is a constant that depends on  $y(t)$  and it's derivatives, and  $p \geq 1$ . Then the method is **consistent** at order  $p$ .

**Def 3.2.3 (Stability):** Suppose that a method satisfies an  $h$ -independent Lipschitz condition on  $D$  (spatial domain)

$$\|\Psi_h(u) - \Psi_h(v)\| \leq (1 + h\hat{L})\|u - v\| \quad \forall u, v \in D.$$

Then the method is **stable**. *Note  $\hat{L}$  need not be the same Lipschitz constant as for the vector field.*

**Thm 3.2.1 (Convergence of One-Step Methods):** Given a differential equation and a one-step method  $\Psi_h$  which is **consistent** and **stable**. Then the method is **convergent**.

**Interpolating Polynomials:** Given  $s$  distinct *abscissa points*  $c_0, \dots, c_s$  and *data points*  $g_0, \dots, g_s$ , there exists a unique interpolating polynomial  $P(x) \in \mathbb{P}_{s-1}$  passing through all points  $(c_i, g_i)$ .

**Lagrange Polynomials:** For a set of abscissae  $c_0, \dots, c_s$ , the Lagrange polynomials  $\ell_i$ ,  $i = 1, \dots, s$  are defined by

$$\ell_i(x) = \prod_{\substack{j=1 \\ j \neq i}}^s \frac{x - c_j}{c_i - c_j}.$$

The Lagrange polynomial  $\ell_i$  is the interpolating polynomial through the data  $g_j = \begin{cases} 1 & \text{if } j = i \\ 0 & \text{if } j \neq i \end{cases}$ .  $\{\ell_i\}$  form a basis for  $\mathbb{P}_{s-1}$ , and any polynomial  $Q(x)$  has the simple form

$$Q(x) = \sum_{i=1}^s Q(c_i) \ell_i(x) = \sum_{i=1}^s g_i \ell_i(x).$$

**Numerical Quadrature:** Given a smooth function  $g(x) : \mathbb{R} \rightarrow \mathbb{R}$ , and  $s$  **quadrature points**  $0 \leq c_1 < \dots < c_s \leq 1$  we can estimate  $\int_0^1 g(x) dx$  by integrating the corresponding interpolating polynomial  $P(x) \in \mathbb{P}_{s-1}$ . Define the weights

$$b_i = \int_0^1 \ell_i(x) dx.$$

Then our approximate integral is

$$\int_0^1 g(x) dx \approx \int P(x) dx = \sum_{i=1}^s g(c_i) \int_0^1 \ell_i(x) dx = \sum_{i=1}^s b_i g(c_i)$$

Therefore for interval  $[t_0, t_0 + h]$  we have

$$\int_{t_0}^{t_0+h} g(x) dx \approx \int_{t_0}^{t_0+h} P(x) dx = \sum_{i=1}^s b_i g(t_0 + hc_i).$$

A quadrature rule has **order**  $p$  if it integrates any polynomial  $\in \mathbb{P}_{p-1}$  exactly. We always have  $p \geq s$ , and for optimal choice of  $c_i$  we have  $p = 2s$ .

**One-Step Collocation:** Given an ODE we wish to construct the **collocation polynomial**  $u(t) \in \mathbb{R}^d$  that satisfies

$$u(t_0) = y_0 u'(t_0 + c_i h) = f(u(t_0 + c_i h)).$$

In particular, it agrees with our solut at  $t_0$ , and it's derivative matches that of the solution at each  $c_1, \dots, c_s$ . We can use such a polynomial to approximate a numerical solution to our ODE by decomposing it's derivative  $u'$  into Lagrange polynomial components, and then integrating over  $[t_0, t_0 + h]$  to get  $u(t_0 + h) = u(t_1)$ .

Let  $F_i = u'(t_0 + hc_i)$  be the value of the derivative of the polynomial at node  $c_i$ . Then

$$F_i = f(y_0 + h \sum_{j=1}^s a_{ij} F_j), \quad (\text{A})$$

$$y_{n+1} = y_n + h \sum_{i=1}^s b_i F_i. \quad (\text{B})$$

where

$$a_{ij} = \int_0^{c_i} \ell_j(x) dx,$$

$$b_i = \int_0^1 \ell_i(x) dx.$$

First solve the  $sd$ -dimensional system of non-linear equations given by (A), and plug into (B).

**Rem 3.6.1 (Continuous Approximations):** Collocation provides a continuous approximation  $u(t)$  of the solution  $y(t)$

on each interval  $[t_n, t_{n+1}]$ .

**Rem 3.6.2 (Optimal Node Placement):** For optimal order of accuracy, use **Gauss-Legendre** collocation methods. This means placing nodes at roots of shifted Legendre polynomials. For  $s = 1, 2, 3$ , the optimal nodes are

$$c_1 = \frac{1}{2}, \quad p = 2,$$

$$c_1 = \frac{1}{2} - \frac{\sqrt{3}}{6}, \quad c_2 = \frac{1}{2} + \frac{\sqrt{3}}{6}, \quad p = 4,$$

$$c_1 = \frac{1}{2} - \frac{\sqrt{15}}{10}, \quad c_2 = \frac{1}{2}, \quad c_3 = \frac{1}{2} + \frac{\sqrt{15}}{10}, \quad p = 6.$$

**Runge-Kutta Methods:** are a generalisation of collocation methods, since they don't have coefficients that rely on integrals of Lagrange polynomials, they can have any coefficients. A **Runge-Kutta** method is any method of form

$$Y_i = y_n + h \sum_{j=1}^s a_{ij} f(Y_j), \quad i = 1, \dots, s,$$

$$y_{n+1} = y_n + h \sum_{i=1}^s b_i f(Y_i).$$

where  $s$  is the **number of stages**,  $b_i$  are the **weights**, and  $a_{ij}$  are the **internal coefficients**. Such a method generates the discrete flow-map

$$\Psi_h(y) = y + h \sum_{i=1}^s b_i f(Y_i(y, h)).$$