

CS061:

Machine Organization & Assembly Language

Lab 1

Agenda

1. Presentation:
 - a. Intro to Assembly Language
 - b. Basics of LC-3 Programming
 - c. Setting up LC-3 Tools
 - d. Lab Descriptions
2. Work Time / Questions / Demos

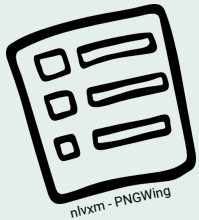
What is Assembly Language?

- **Assembly Language:** Human-readable representation of instructions that correspond to CPU instructions.
- Formal Definitions:
 - Oxford: “A low-level symbolic code converted by an assembler.”
 - Merriam Webster: “a programming language that consists of instructions that are mnemonic codes for corresponding machine language instructions.”

Let's Back Up A Little



- **CPU** (central processing unit): Responsible for executing a set of instructions that change the state of the system.
- CPU only understands limited set of instructions.
 - Think of it like a dog!
 - Dogs can't understand full human conversation, but understand commands like "sit" or "stay".
 - CPU can't understand a raw C++ program. It only understands "add x to y" or "set value at memory address x to value y", etc.
- Assembly language is the readable representation of the raw instructions the CPU understands.

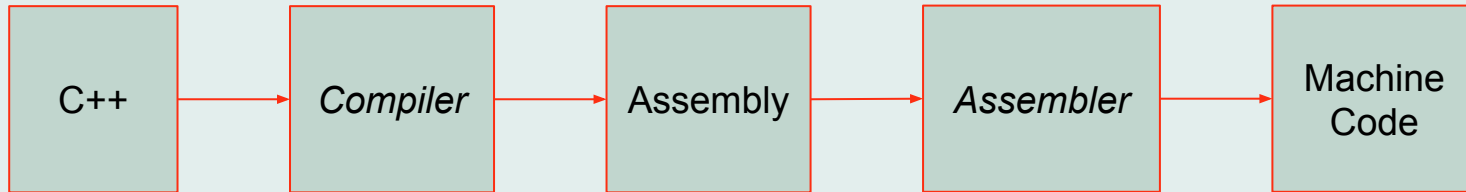


Assembler

- Computers can only understand 1s and 0s.
- **Assembler:** Converts assembly language (human-readable representation) to machine code (1s and 0s).

Running Hello World

- How do we get from a full C++ program to something that the CPU understands?
- Compiler takes C++ source code and converts it into assembly language (**if using gcc*).
- Then an assembler converts the assembly language to machine code!



But...Why?

- Why learn how to read/write assembly language?
- Careers:
 - Computer/Software Security
 - Reverse Engineering
 - Embedded Systems Developers
 - Compiler Writer
- Understanding how computers work under the hood.
 - OS Development
 - High Performance Computing
 - Software Testing - Static Analysis
 - Debugging

LC-3

- **LC-3** (Little Computing 3): Educational assembly language.
 - Easier to learn/understand than actual assembly language!
- **ISA** (Instruction-Set Architecture):
 - **15 instructions.**
 - Basic things like addition, bitwise NOT, bitwise AND, save value at memory, load value from memory, etc.
 - Still missing some basic things: no subtraction instruction or multiply instruction.
- **8 registers (16-bits)**
 - Registers are like global integer variables.
 - Can store/read some value at them.
 - *R0 - R7*
- LC-3 program consists of pseudo-ops, instructions, and data.

Show me code!!!

```
;-----  
; Name: Doe, Sam  
; Email: sdoe013@ucr.edu  
;  
; Lab: Lab 1, Ex 1  
; Section: 021  
; TA: Sanchit Goel, Westin Montano  
;  
; -----  
  
.ORIG x3000  
; Program Code Here  
  
AND R1, R1, #0  
  
LD R1, DEC_6  
  
DO_WHILE_LOOP  
    ADD R1, R1, #-1  
    BRp DO_WHILE_LOOP  
END_DO_WHILE_LOOP  
  
HALT  
; Local Data  
DEC_6    .FILL #6  
  
.END
```

What does it mean?

```
; -----  
; Name: Doe, Sam  
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;  
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    ADD R1, R1, #-1  
    BRp DO_WHILE_LOOP  
END_DO_WHILE_LOOP
```

```
HALT  
; Local Data  
DEC_6 .FILL #6  
  
.END
```

- Comments start with ‘;’
- Header section - Put this in all your exercises!
- **Pseudo-Op:** Not a CPU instruction, but an assembler directive.
 - Executes at assembly time!
- All pseudo-ops start with “.”
- “.ORIG” tells the assembler to put code/data starting at that memory address (x3000)
- Instruction that indicates the end of a program!

How do we math?

; Program Code Here

Destination Source Operand

AND R1, R1, #0

LD R1, DEC_6

DO_WHILE_LOOP

ADD R1, R1, #-1

BRp DO_WHILE_LOOP

END_DO_WHILE_LOOP

- General Instruction Format:
- *Instruction, Destination Register, Source Register, Source Register 2 / Operand*
- AND instruction performs bitwise AND between source register (R1) and operand (#0) then stores it in destination register (R1).
- $R1 \leftarrow R1 \& 0$ (*Register Transfer Notation*)
- Sets R1 to 0
- Subtracts R1 by 1.
- ADD adds source register (R1) with operand (#-1) and stores value in destination register (R1).
- $R1 \leftarrow R1 - 1$

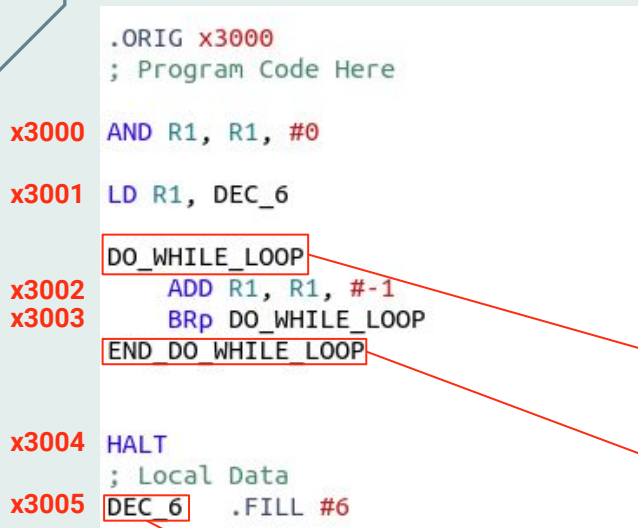
Code is Memory

```
.ORIG x3000
; Program Code Here

x3000 AND R1, R1, #0
x3001 LD R1, DEC_6

DO_WHILE_LOOP
x3002 ADD R1, R1, #-1
x3003 BRp DO_WHILE_LOOP
END DO WHILE LOOP

x3004 HALT
; Local Data
x3005 DEC_6 .FILL #6
```

The diagram shows a block of assembly code. Several labels and instructions are highlighted with red rectangular boxes: 'DO_WHILE_LOOP' at line x3002, 'END DO WHILE LOOP' at line x3003, 'DEC_6' at line x3005, and the instruction 'ADD R1, R1, #-1' at line x3002. Red arrows originate from these boxes and point towards the right-hand list of bullet points, specifically to the 'Labels' section and its sub-points.

- Think of memory as a contiguous block of values with an address representing location of value.
- Each instruction takes up one memory space.
- For reference, every memory space holds a 16-bit value.
- **Labels:** Alias for a memory address!
 - DO_WHILE_LOOP refers to the address x3002 which corresponds to the ADD instruction.
 - END_DO_WHILE_LOOP refers to the address x3004.

- DEC_6 refers to the address at x3005.
- The .FILL pseudo-op sets the value at address x3005 to 6.

Code is Memory

Memory Address	Value (Hex)	Value (Decimal)	Instruction/Pseudo-Op
▶ x3000	x5260	21088	<i>AND R1, R1, #0</i>
▶ x3001	x2203	8707	<i>LD R1, DEC_6</i>
▶ x3002	x127F	4735	<i>ADD R1, R1, #-1</i>
▶ x3003	x03FE	1022	<i>BRp DO_WHILE_LOOP</i>
▶ x3004	xF025	-4059	<i>HALT</i>
▶ x3005	x0006	6	<i>DEC_6 .FILL #6</i>

Movin' Values

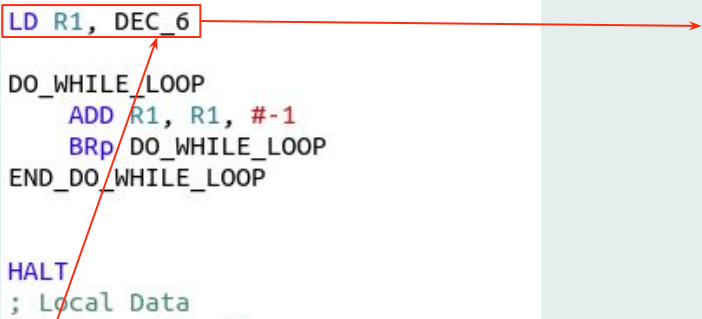
```
.ORIG x3000
; Program Code Here

x3000 AND R1, R1, #0

x3001 LD R1, DEC_6

DO_WHILE_LOOP
x3002     ADD R1, R1, #-1
x3003     BRp DO_WHILE_LOOP
END_DO_WHILE_LOOP

x3004 HALT
; Local Data
x3005 DEC_6 .FILL #6
```



- Can we put a value at a memory address into a register, and vice-versa?
- **LD (Load Direct)** takes a value from memory and puts in the destination register (R1).
 - Operand (DEC_6) is a label.
 - DEC_6 refers to address x3005.
 - Value at x3005 is 6.
 - $R1 \leftarrow \text{Mem}[\text{DEC_6}]$
 - R1 will have 6 after instruction executes.
- 3 different types of loading/storing in LC-3.

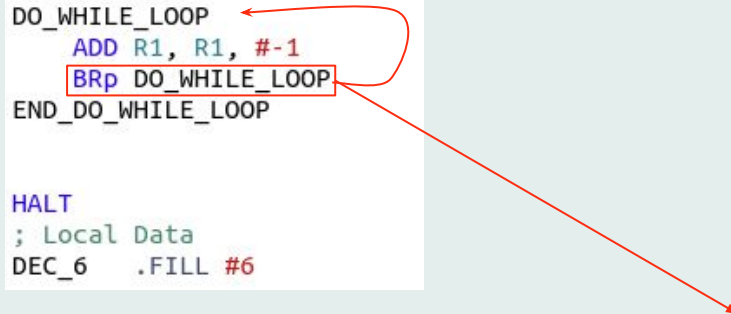
On One Condition

```
.ORIG x3000
; Program Code Here

x3000 AND R1, R1, #0
x3001 LD R1, DEC_6

DO_WHILE_LOOP
x3002     ADD R1, R1, #-1
x3003     BRp DO_WHILE_LOOP
END_DO_WHILE_LOOP

x3004 HALT
; Local Data
x3005 DEC_6 .FILL #6
```



- Sometimes we need to repeat a section of code (loops) or only run a section of code if something happens (conditions).
- BR - **Branch** jumps to a line of code based on a condition.
 - Condition specified by p (positive), n (negative), z (zero).
 - Condition set by any instruction that modifies registers.
 - The line to jump to is determined by the label operand.
- BRp will check if result of addition above is > 0 .
 - If it is, jumps to address corresponding to DO_WHILE_LOOP (x3002).
 - If not, continues to next line.

From the Top

```
.ORIG x3000  
; Program Code Here
```

```
AND R1, R1, #0
```

$R1 \leftarrow R1 \& 0$

= Set R1 to 0

```
LD R1, DEC_6
```

$R1 \leftarrow \text{Mem}[\text{DEC_6}]$

= Load R1 with 6

```
DO_WHILE_LOOP
```

```
    ADD R1, R1, #-1
```

$R1 \leftarrow R1 - 1$

= Subtract R1 by 1

```
    BRp DO_WHILE_LOOP
```

Branch

= Jump to DO_WHILE_LOOP if $R1 > 0$

```
END_DO_WHILE_LOOP
```

Repeat 5 times ($R1 = 6$, then $R1 = 5$, ... $R1 = 0$)

```
HALT
```

End Program

```
; Local Data
```

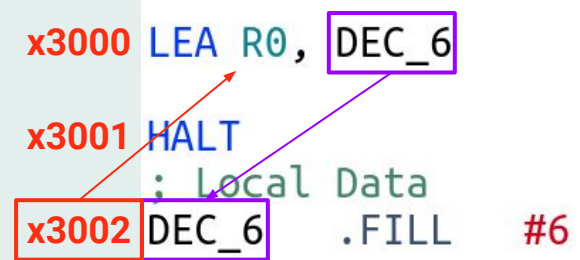
```
DEC_6    .FILL #6
```

This program does nothing meaningful 😊

```
.END
```


Load Effective Address (LEA)

- Sometimes need the memory address of a label in a register.
 - E.g. loading/storing, PUTS, etc
- **LEA Rx, LABEL**
 - Gets the memory address of the label and stores it in Rx.
 - E.g. LEA R0, DEC 6 : Gets the memory address of label DEC_6 (x3002) and stores it in R0.
 - R0 <- DEC_6.
 - R0 <- x3002



Stringz

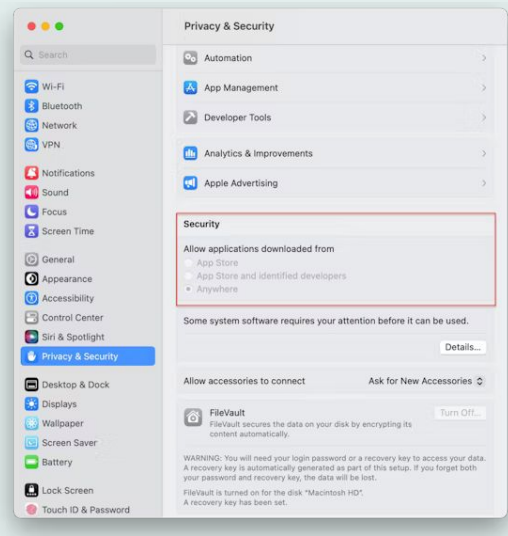
- **“.STRINGZ”**
 - *Pseudo-op* to tell assembler to create a string starting at a memory address.
 - Each character in the string *occupies 1 memory space*.
 - End of a string is denoted by 0 (\0 - the null character).
- **“PUTS”** (Print)
 - MSG address is the start of a string.
 - PUTS prints out the string *starting* at the address specified by R0.
 - PUTS will *always* use R0 as the starting address.

```
x3000 LEA R0, MSG
x3001 PUTS
x3002 HALT
      ; Local Data
x3003 MSG .STRINGZ "WOW! "
```

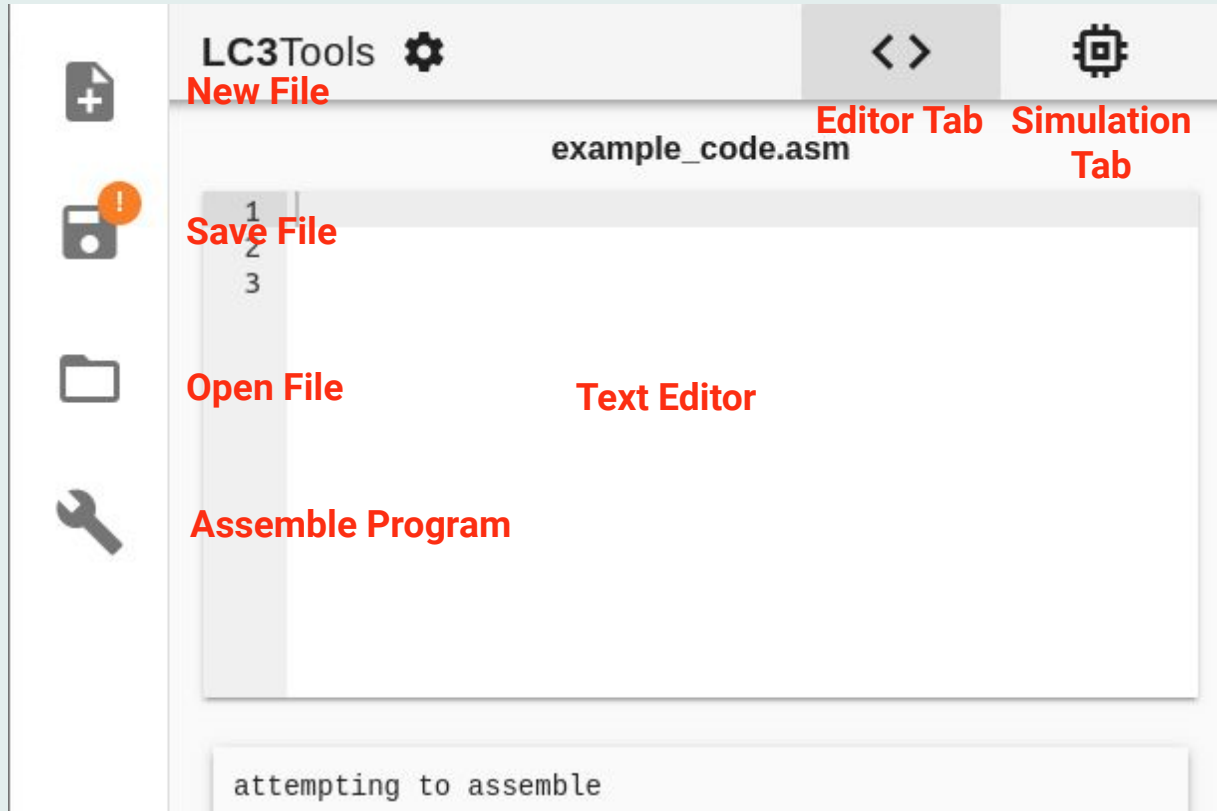
Address	x3003	x3004	x3005	x3006	x3007
Value	'W'	'O'	'W'	'!'	0

Your Turn!

- Download LC3 Tools from <https://github.com/chiragsakhuja/lc3tools/releases/tag/v2.0.2>
 - The Lab 1 manual has the link.
 - Download LC3Tools-2.0.2.exe for Windows
 - Download LC3Tools-2.0.2.dmg for Mac
 - Download LC3Tools-2.0.2.AppImage for Linux
- On Mac, make sure to allow opening apps from anywhere.
 - To enable this, go to System Settings -> Privacy & Security -> Security.
 - Select the Anywhere option.





Using LC-3 Tools



Using LC-3 Simulation

Play
Reload Object File
Step Over
Step In
Step Over
Reset

LC3Tools  < > 

Registers

R0	x0000	0
R1	x0000	0
R2	x0000	0
R3	x0000	0
R4	x0000	0
R5	x0000	0
R6	x3000	12288
R7	x0000	0
PSR	x0002	2 CC: Z
PC	x0200	512
MCR	x0000	0

Memory

0	▶	x3000	x0000	0
1	▶	x3001	x0000	0
2	▶	x3002	x0000	0
3	▶	x3003	x0000	0
4	▶	x3004	x0000	0
5	▶	x3005	x0000	0
6	▶	x3006	x0000	0
7	▶	x3007	x0000	0
8	▶	x3008	x0000	0
9	▶	x3009	x0000	0
A	▶	x300A	x0000	0
B	▶	x300B	x0000	0
C	▶	x300C	x0000	0
D	▶	x300D	x0000	0
E	▶	x300E	x0000	0
F	▶	x300F	x0000	0
10	▶	x3010	x0000	0
11	▶	x3011	x0000	0
12	▶	x3012	x0000	0
13	▶	x3013	x0000	0
14	▶	x3014	x0000	0
15	▶	x3015	x0000	0
16	▶	x3016	x0000	0
17	▶	x3017	x0000	0
18	▶	x3018	x0000	0
19	▶	x3019	x0000	0
1A	▶	x301A	x0000	0
1B	▶	x301B	x0000	0
1C	▶	x301C	x0000	0
1D	▶	x301D	x0000	0
1E	▶	x301E	x0000	0
1F	▶	x301F	x0000	0

Console (click to focus)

Jump To Location

PC

← ← → →

Exercise 0

- Print “Hello World” to the Console (in the Simulation Tab).
- **“LEA”** (Load Effective Address)
 - LEA R0, MSG_TO_PRINT
 - Remember: label is alias for memory address.
 - LEA puts the memory address that the label (MSG_TO_PRINT) refers to into destination register (R0).
 - E.g. suppose MSG_TO_PRINT referred to address x3006 then after instruction runs, R0 = x3006.
- **“.STRINGZ”**
 - Pseudo-op to tell assembler to create a string starting at a memory address.
 - Each character in the string occupies 1 memory space.
 - End of a string is denoted by 0 (\0 - the null character).
- **“PUTS”** (Print)
 - MSG_TO_PRINT address is the start of a string.
 - PUTS prints out the string starting at the address specified by R0.
 - PUTS will *always* use R0 as the starting address.

Exercise 1

- Create a program that multiplies R1 by R2.
- Code given to you.
- Refer to previous slides to see how the assembly instructions work.

Demos

- When you complete all exercises in a lab, you must demo the lab to a TA or grader to receive credit.
- TAs & Graders will ask questions about your code and how it works.
- If you do not answer a question, you can always come back and re-demo (without penalty).
- Must demo before the next lab session to receive full credit for the lab!
 - May demo during any TA or Grader's office hours (TBA) but not the professor's.
 - **If demo before Friday, then will receive 1 extra point for the lab.**
 - If demo previous lab during the next lab session, there will be 3 point penalty.
 - 7 for demoing, 3 points for attendance = 10 points for full credit (11 is extra credit).
- To sign up for a demo, use the demo spreadsheet link that will be provided.
 - Net ID is the first part of your email. e.g. sd0e03 is the net ID if sd0e03@ucr.edu is the email.
 - Demos are usually taken in order.
 - **Please only sign up once you have completed all the exercises and fully understand your code.**