## $\frac{\text{Quiz-4}}{\text{Data Structures}}$

## $\underline{\mathbf{Set}\ \mathbf{A}}$

Time Allowed: 60 minutes

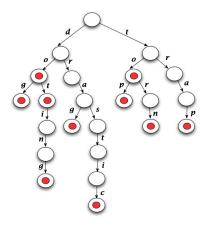
 $18 th\ April$ 

INSTI	RUCTIONS
1. T	This paper contains Multiple choice questions.
2. N	farking Scheme
	Marking scheme format : (correct, unattempted, incorrect)
	MCQs (+4, 0, -1)
	True/False(+4, 0, -2)
	Fill-ups(+4, 0, 0)
3. A	nswers have to be written in the space provided besides the question.
	o:
E tl	all possible strings of length at most $k$ ( $k > 0$ ) having only lowercase anglish alphabet are inserted in a Trie, what is the total number of nodes in the Trie?

2. Consider the following Trie :

Lets denote the ending node of strings "drag" and "drastic" by A and B respectively. Which of the following strings when inserted in this Trie end on the node represented by LCA(A, B)?

- (a) "dra"
- (b) "dr"
- (c) "d"
- (d) "drag"



3. Number of distinct strings in the Trie is always equal to the number of leaf nodes. State True / False

Ans : \_\_\_\_\_

4. Inserting a string in a Trie is O(height of the Trie), where height of the Trie is maximum distance from Root node to any of the leaf nodes before insertion. State True / False.

Ans : \_\_\_\_\_

5. Number of nodes at depth d ( depth of root node == 0 ) is equal to the number of distinct characters at d-th position in all the strings ( having length  $\geq d$  ) that have been inserted in the Trie. State True / False.

Ans : \_\_\_\_\_

6. Given a full binary tree of height H, and a node A at depth d ( depth of root node ==0 ) , find out how many unordered pairs of nodes u, v ( such that

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u \stackrel{!}{=} v ) exist such that LCA(u, v) == A ? Ans : ____
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7. Consider a tree, and its two nodes A and B , at depths d1 and d2 respectively ( depth of root node == 0 ). Let LCA(A, B) = L and depth(L) = d3. Let  $S = \{ LCA(u, B) \mid u \text{ is an ancestor of } A \}$ . Assume a node is also an ancestor of itself. What is the cardinality of S?

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(a) d3 + 1

(b) d1 + d2 - d3 + 1

(c) d1 + 1

(d) d2 + 1
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8. Consider the following pseudo-code of a function F that takes two arguments : nodes u, v that belong to a tree.

```
F(u, v):
    Initialize : l = 0, r = depth(v)
    while l <= r :
        d = (l + r) / 2
        A = d-th ancestor of v
        if A is an ancestor of u :
            return A
        else :
        l = mid + 1
        return u</pre>
```

Mark all the correct options:

- (a) F returns LCA(u, v) in some cases.
- (b) F returns LCA(u, v) in all cases.
- (c) F returns some ancestor of LCA(u, v) in all cases.
- (d) F returns ancestor of u in all cases.

9. If the first k nodes on the path from Root to node u are same as the first k nodes on the path from Root to node v, in a given Tree, the k-th node in any of these paths (from root to u or v) is always the LCA of u and v. State True / False

Ans : \_\_\_\_\_

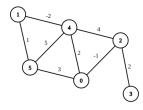
10. Given a tree, and two nodes u and v belonging to it. Any node A present in the path from u to v which satisfies the property, that both u and v are present in the subtree of A is LCA of u and v. State True / False.

Ans : \_\_\_\_\_

- 11. Consider a tree having N nodes and N 1 edges (  $N \ge 2$  ). Let us define a function F(u, v, R) which returns the LCA(u, v) when the tree is rooted at R. Given two nodes A and B ( A != B ), now consider the set S = { F(A, B, R) | R is any node in the tree }. What is the cardinality of S?
  - (a) N
  - (b) 1
  - (c) N/2
  - (d) distance(A, B) + 1
- 12. Given a rooted tree, consider nodes A, B, C, D. Let E = LCA(A, B) and F = LCA(C, D) and G = LCA(E, F). Assume nodes A, B, C, D, E, F, G are all distinct.

 ${\it Mark}$  all correct options :

- (a) LCA(A, C) == G
- (b) LCA(B, D) == G
- (c) LCA(A, C) == E
- (d) LCA(A, G) == G
- 13. Consider the following graph :



Mark all the correct statements for the given graph:

- (a) Dijkstra's Algorithm can be used to find the single source shortest path from source vertex 0.
- (b) Bellman Ford Algorithm can be used to find the single source shortest path from source vertex 0.
- (c) Floyd Warshall Algorithm can be used to find all pair shortest path.
- (d) None of these.
- 14. For a directed graph having V vertices and having no negative cycles, if instead of running the outer loop of Bellman Ford Algorithm for V 1 iterations, it is run for V iterations, the distances computed by the algorithm may be wrong. State True / False.

Ans : \_\_\_\_\_

15. DAG-BELLMAN-FORD(V, E, s):

- 1. TOPOLOGICALLYSORT(V, E)
- 2. INITIALIZE(V, E)
- 3. for i in 1..|V|-1:
- 4. for each vertex u in V taken in topological order :
- 5. for each edge originating at u : (u, v) in E :
- 6. RELAX(u, v)
- 7. if no edge was relaxed:
- 8. break

For an input graph which is a DAG, to the above function, how many times does the loop in line number 3 run?

Mark all the correct options:

(a) 1

- (b) 2
- (c) |V| 1
- (d) Can't say, depends on the input graph.
- 16. Consider a connected undirected graph G(V, E) having all edge weights greater than 0. Dijkstra's algorithm is applied from source vertex s. Let d[i] store the shortest distance from s to vertex i. Let  $E' = \{ (u, v) \mid d[u] == d[v] + w(u, v) \text{ or } d[v] == d[u] + w(u, v) \}$ , such that (u, v) denotes the edge joining vertices u and v, and w(u, v) denote the edge weight of edge (u, v). Graph represented by G'(V, E') will always form a tree. State True / False. Ans: \_\_\_\_\_
- 17. What is the worst case space complexity of Bellman Ford algorithm on an input Graph G(V, E)?
  - (a)  $\theta(|V|)$
  - (b)  $\theta(|V| + |E|)$
  - (c) O(|V|)
  - (d) O(|V| + |E|)

Note: We consider only the auxiliary space used for calculating the space complexity of an algorithm.

- 18. On applying Floyd Warshall Algorithm on a directed Graph G(V, E), it is found that distance(u, u) < 0 for some vertex  $u \in V$ . Mark all the correct options:
  - (a) G has at least one edge with negative edge weight.
  - (b) G has a negative cycle.
  - (c) All edges going out of u have negative edge weights.
  - (d) There exists an edge going out of u having negative edge weight.
- 19. Consider a connected undirected weighted graph G(V, E) with distinct edge weights, such that |V|, |E| > 3. Let S1, S2 be two non-empty mutually exclusive subsets of V, such that  $S1 \cup S2 == V$ .

Let  $E'=\{\ (u,\,v)\mid (u,\,v)\in E\ and\ (\ (\ u\in S1\ and\ v\ in\in S2\ )\ or\ (\ u\in S2\ and\ v\in S1\ )\ )\ \}$ 

The claim that edge e = (u, v) such that  $e \in E'$  and weight(e) is minimum, will belong to the Minimum Spanning Tree of G. State True / False.

Ans : \_\_\_\_\_

- 20. Given an undirected graph G(V,E), if you start the DFS from a node  $n\in V$ , what is the worst complexity of the algorithm ? Mark all the correct options :
  - (a) O(|V|)
  - (b) O(|V| + |E|)
  - (c) O(|V| \* |E|)
  - (d) O(|E|)