

CS120: Computer Networks

Lecture 4. Framing and Error Detection

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Quantifying Error

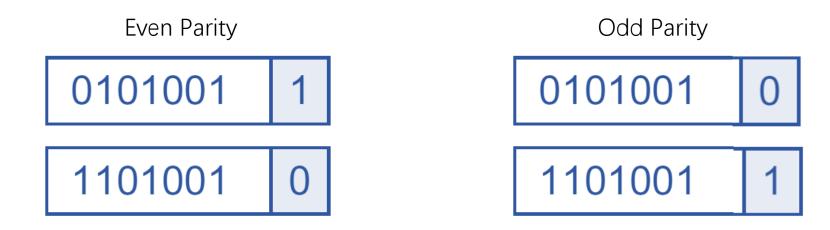
- Shannon capacity gives a throughput upper bound.
 - Actual throughput << channel capacity
 - => underutilized
 - Actual throughput >= channel capacity
 - => error
- Bit error rate (BER): error bits/transmitted bits
 - 10⁻¹, 10⁻², 10⁻³,···
- Trade off
 - High rate -> low reliability

Handling Error

- Error Detection
 - Parity Check
 - Checksum
 - Cyclic Redundancy Check (CRC)
- Error Correction
 - Hamming code
 - Others: convolutional code, fountain code, etc.

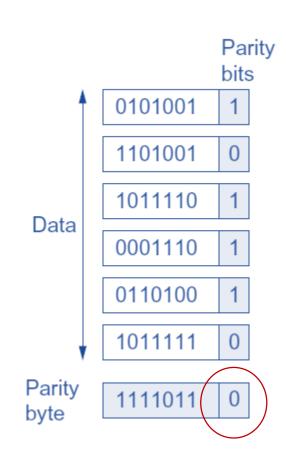
Parity Check

• Method: adding one extra bit to a 7-bit code to balance the number of 1s in the byte.



2D Parity Check

- Add one byte to check the "columns"
 - 1-bit error
 - 2-bit error
 - 3-bit error
 - 4-bit error
 - not all



Checksum

Method: add all the bytes up use ones' complement arithmetic;
 then take ones' complement of the result.

Ip header

Ones' Complement

Bits	Unsigned	Signed	Ones' Complement
1111	15	-1	-0
1110	14	-2	-1
1101	13	-3	-2
1100	12	-4	-3
1011	11	-5	-4
1010	10	-6	-5
1001	9	-7	-6
1000	8	-8	-7
0000	0	0	+0

Ones' Complement

- Calculation Examples
 - Signed -4 + (-2)
 - 1100+1110=11010
 - ignore carry
 - 1010 (i.e., -6)
 - Ones' Complement -4 + (-2)
 - 1101+1011=11000
 - shift and add carry
 - 1000+1=1001 (i.e., -6)

Checksum

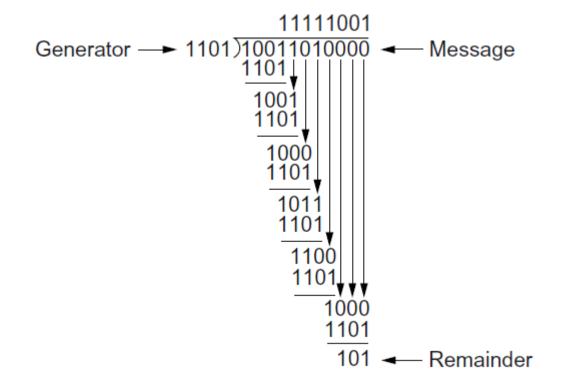
Method: add all the bytes up use ones' complement arithmetic;
 then take ones' complement of the result.

0x4500+0x012b+0x190a+0x0000+0xff11+0x0a14+0x4508+0xe000+0x00fb=0x28e5d =>0x8e5f=>0x71a0

Checksum

- Method: add all the bytes up use ones' complement arithmetic;
 then take ones' complement of the result.
 - Fast calculation
 - Weak protection

Cyclic Redundancy Check (CRC)



CRC Performance

- CRC v.s. Checksum
 - CRC protects more bits
 - CRC takes more time/resource to calculate
- CRC v.s. Hash
 - CRC does not protect data integration
 - Hash functions are more complex than CRC
- CRC Polynomials
 - http://users.ece.cmu.edu/~koopman/crc/

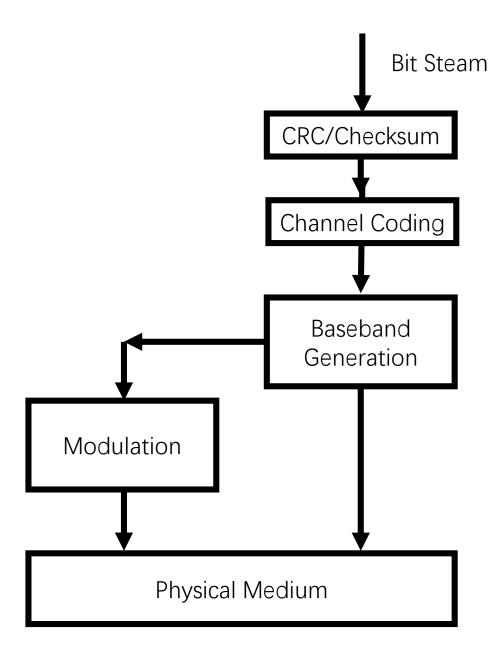
Table 2.3	Table 2.3 Common CRC Polynomials										
CRC	C(x)										
CRC-8	$x^8 + x^2 + x^1 + 1$										
CRC-10	$x^{10} + x^9 + x^5 + x^4 + x^1 + 1$										
CRC-12	$x^{12} + x^{11} + x^3 + x^2 + x + 1$										
CRC-16	$x^{16} + x^{15} + x^2 + 1$										
CRC-CCITT	$x^{16} + x^{12} + x^5 + 1$										
CRC-32	$x^{32} + x^{26} + x^{23} + x^{22} + x^{16} + x^{12} + x^{11}$										
	$+x^{10} + x^8 + x^7 + x^5 + x^4 + x^2 + x + 1$										

CRC Size	CDC Dal	mamial	HD	Hamming weights for number of bits corrupted:									
(bits)	CRC Poly	momiai	пи	1 bit	2 bits	3 bits	4 bits	5 bits	6 bits				
16	CCITT-16	0x8810	4	0	0	0	84	0	2 430				
16	[Baicheva0	0] 0xC86C	6	0	0	0	0	0	2 191				
15	CAN	0x62CC	6	0	0	0	0	0	4 314				
12	CRC-12	0xC07	4	0	0	0	575	0	28809				
12		0x8F8	5	0	0	0	0	1 452	13 258				
8	DARC-8	0x9C	2	0	66	0	2 039	13 122	124 248				
8	CRC-8	0xEA	4	0	0	0	2 984	0	253 084				
7	CRC-7	0x48	3	0	0	216	2 690	27 051	226 856				
7		0x5B	4	0	0	0	5 589	0	451 125				

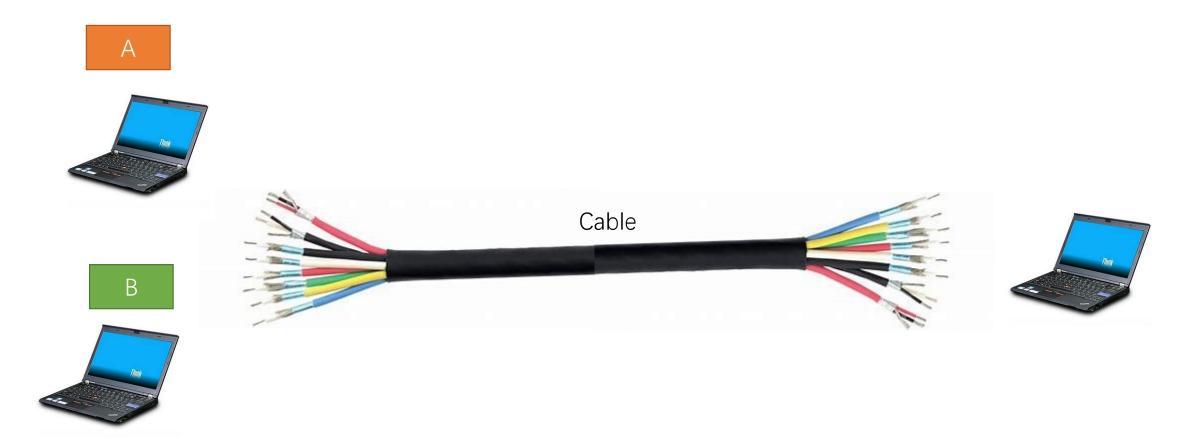
Handling Errors

- Error Detection
 - Parity Check
 - Checksum
 - Cyclic Redundancy Check (CRC)
- Error Correction
 - Hamming code
 - Others: convolutional code, fountain code, etc.
 - Trade off: efficiency and reliability
 - Retransmission is expensive
 - Storage, satellite, etc.
 - Errors are probable
 - Wi-Fi, cellular (channel is unstable, interference)

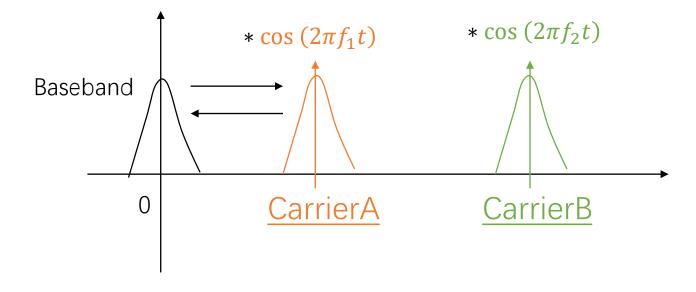
By Now



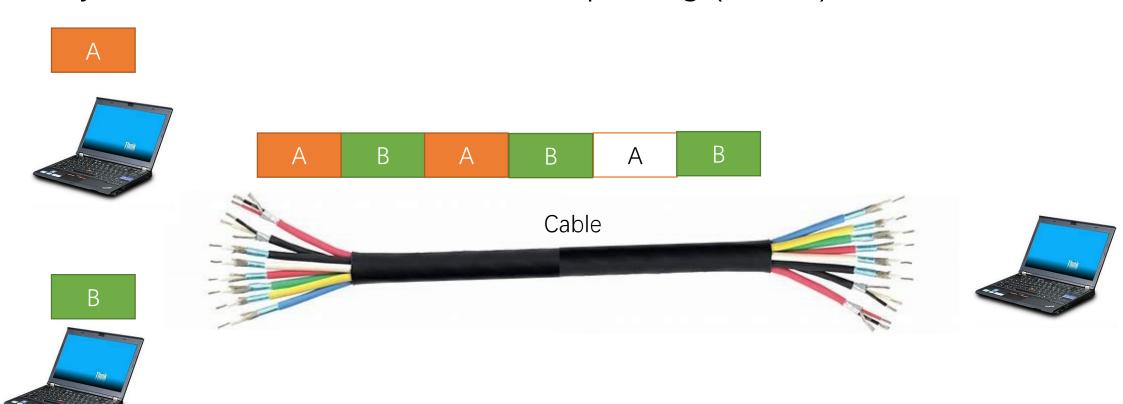
The Multiplexing Problem



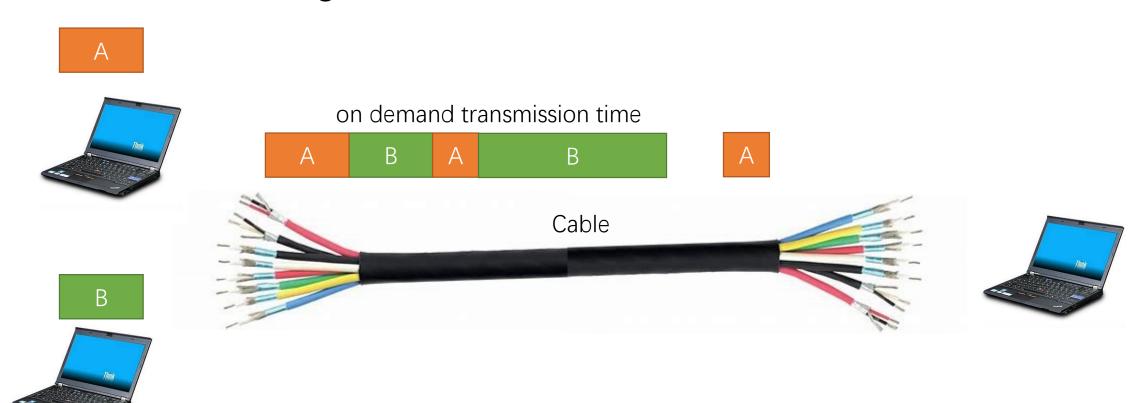
Frequency-division multiplexing (FDM)



• Synchronous Time-division Multiplexing (STDM)



Packet Switching



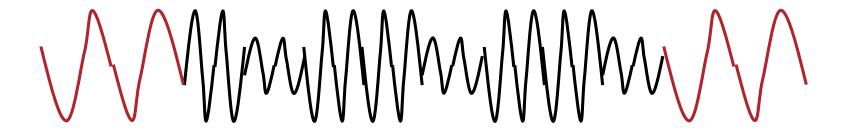
- Fixed Multiplexing
 - STDM, FDM (e.g. telecommunication network, WAN, etc.)
 - Predictable performance
 - Expensive
- Statistical Multiplexing
 - Packet Switching (e.g. Internet, etc.)
 - Performance is random
 - Cheap

Framing

- Why?
 - One transmitter should not occupy the communication resource forever
- Functionality
 - Help receiver recognize the start and the end of the transmission
- Framing Design
 - Frame detection
 - Frame termination

Framing - Detection

- Sync signal: find accurate start through special pattern (e.g., Preamble, Pilot, etc.)
 - Add a special pattern before or after each frame
 - Correlate the special pattern ···



Framing - Termination

- Sentinel-Based Approach
 - Use special patterns (sentinel characters) to indicate the start and end of the frame
 - Similar patterns may exist in payload
 - Solution: character stuffing
 - e.g., disp ('abc''') in matlab shows abc'; ' is the escape character



Example: BISYNC Frame

Framing - Termination

- Counting-based Approach
 - Use special patterns to indicate the start of the frame
 - Use number to indicate the length of the frame
 - Count may contain error
 - Solution: limit the maximum frame length

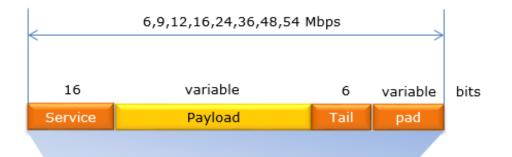


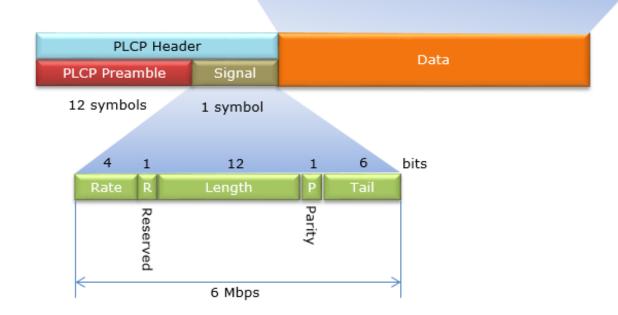
Example: DDCMP Frame

• Ethernet

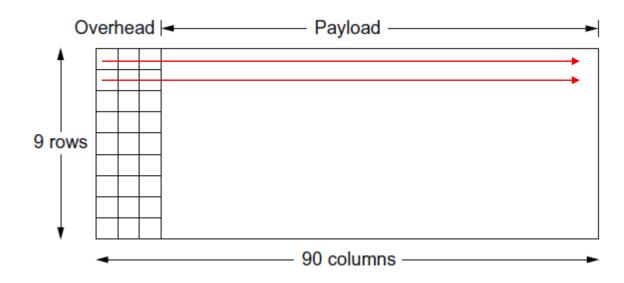
Address Address		Preamble	SFD	Destination MAC Address	Source MAC Address	EtherType	Payload	4	4	FCS	
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• Wi-Fi

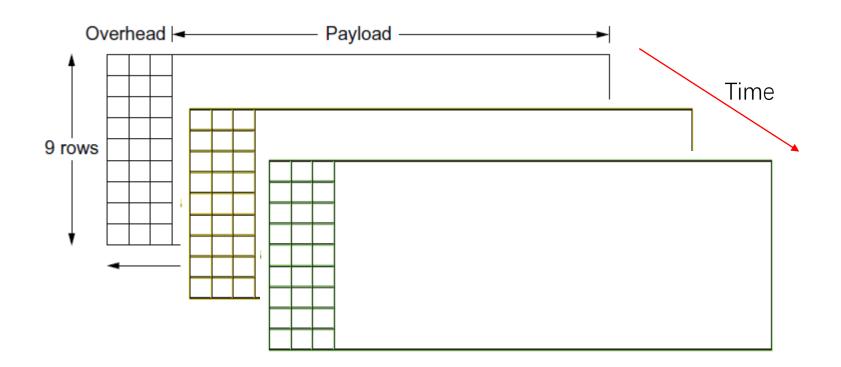




 Synchronous Optical Network (SONET)/Synchronous Digital Hierarchy (SDH)



 Synchronous Optical Network (SONET)/Synchronous Digital Hierarchy (SDH)



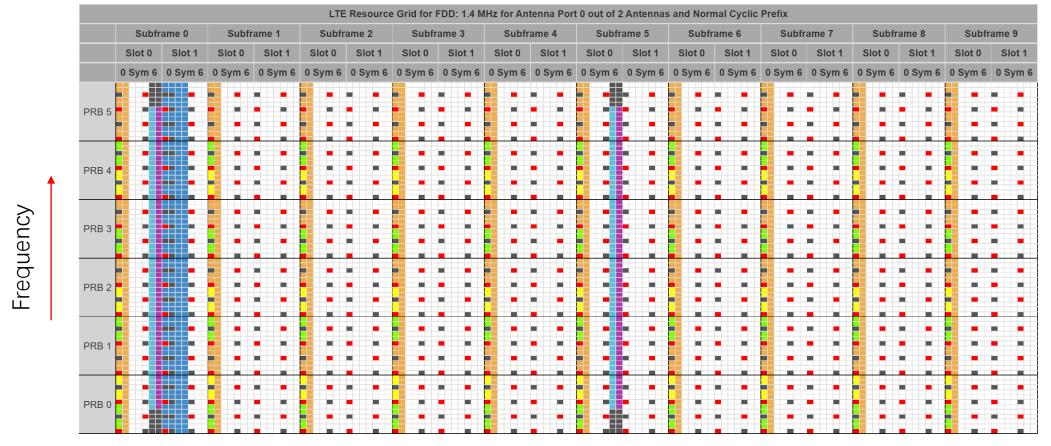
• 4G LTE FDD



LTE Resource Grid for FDD: 1.4 MHz for Antenna Port 0 out of 2 Antennas and Normal Cyclic Prefix																				
	Subframe 0		Subframe 1		Subframe 2		Subframe 3		Subframe 4		Subframe 5		Subframe 6		Subframe 7		Subframe 8		Subframe 9	
	Slot 0	Slot 1	Slot 0	Slot																
	0 Sym 6	0 Sym 6	0 Sym 6	0 Syr																

• 4G LTE FDD





source: https://dhagle.in/LTE.php

Reference

- Textbook 1.2.3
- Textbook 2.3
- Textbook 2.4