



# CS120: Computer Networks

## **Lecture 5. ACK**

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# The ACK Mechanism

- For ACK Transmitter: an acknowledgement (ACK for short) is a small control frame that a protocol sends back to its peer saying that it has received the earlier frame
- For ACK Receiver: the receipt of an acknowledgement indicates to the sender of the original frame that its frame was successfully delivered.



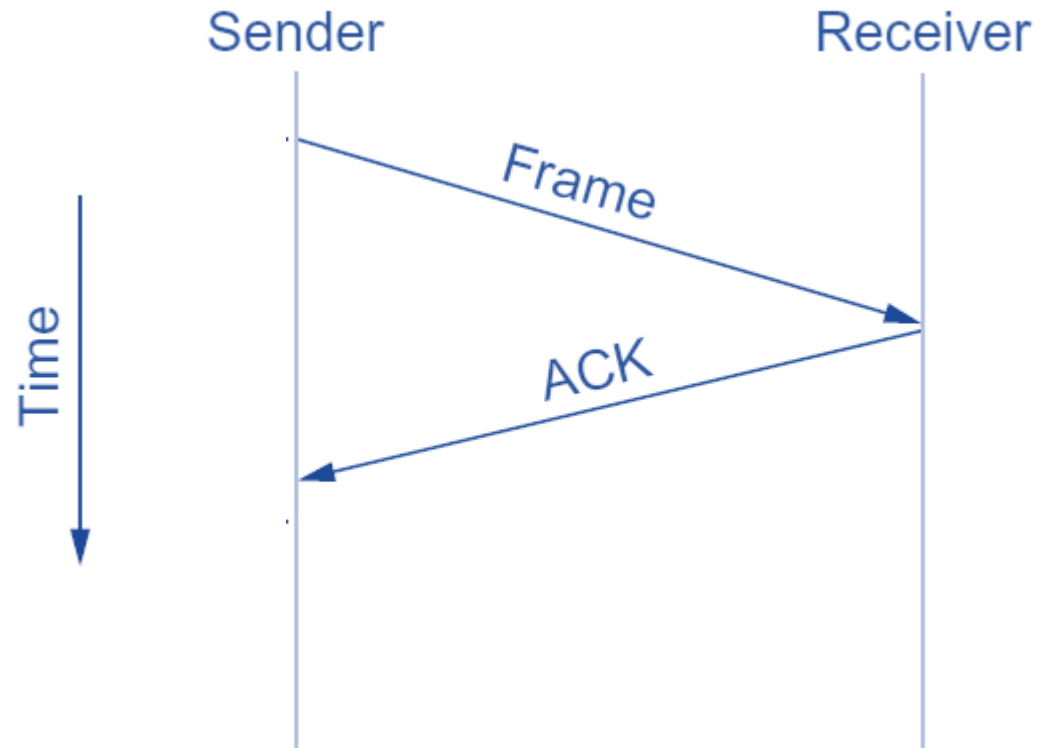
Have you heard that ?

# ACK Schemes

- Stop-and-Wait
- Sliding Window

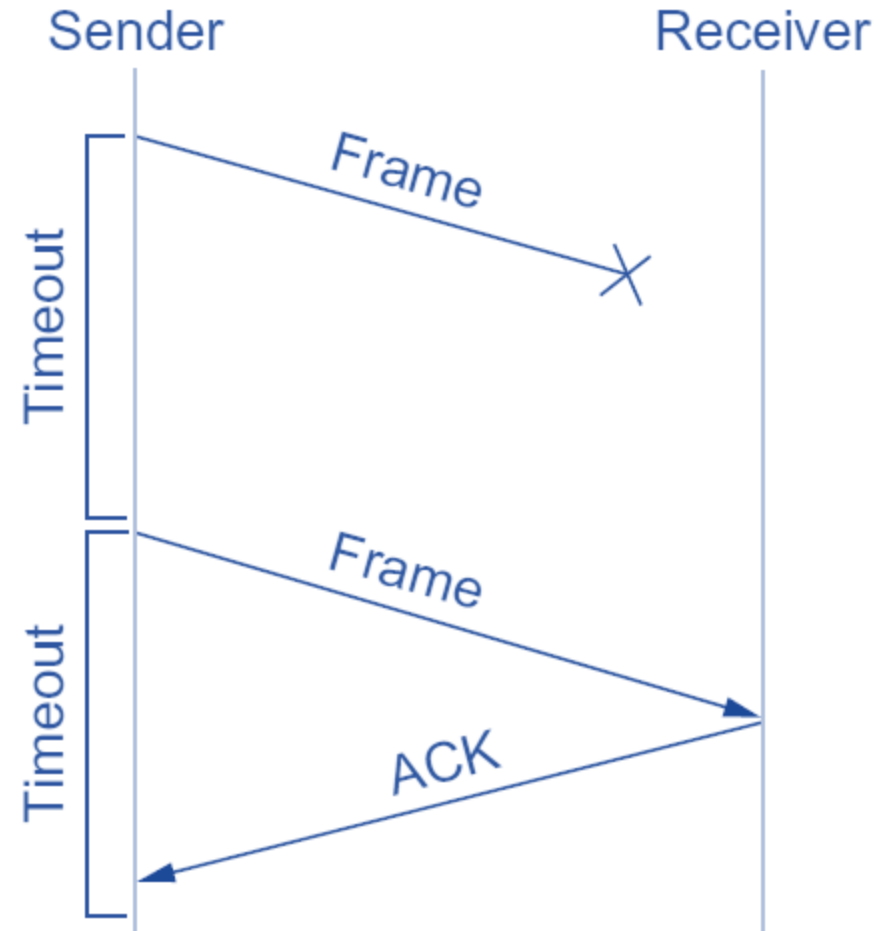
# Stop-and-Wait

- Case 0: (understanding the timeline)



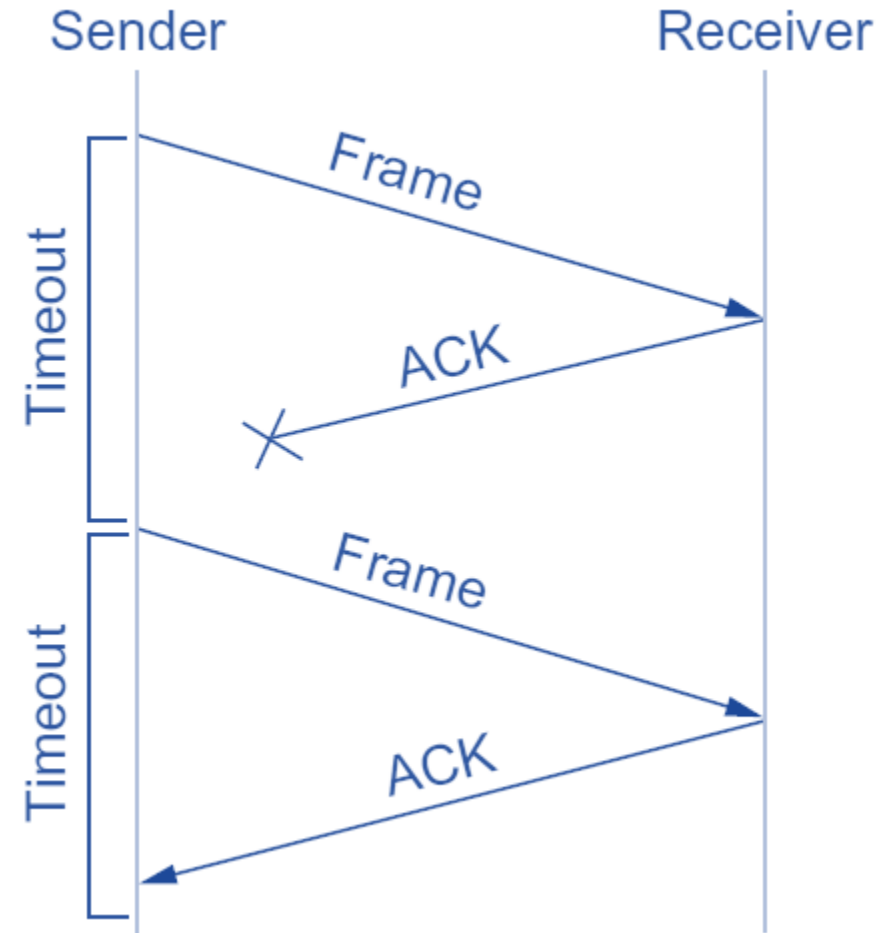
# Stop-and-Wait

- Case 1: Frame Loss
  - Sender time out
  - Sender retransmits



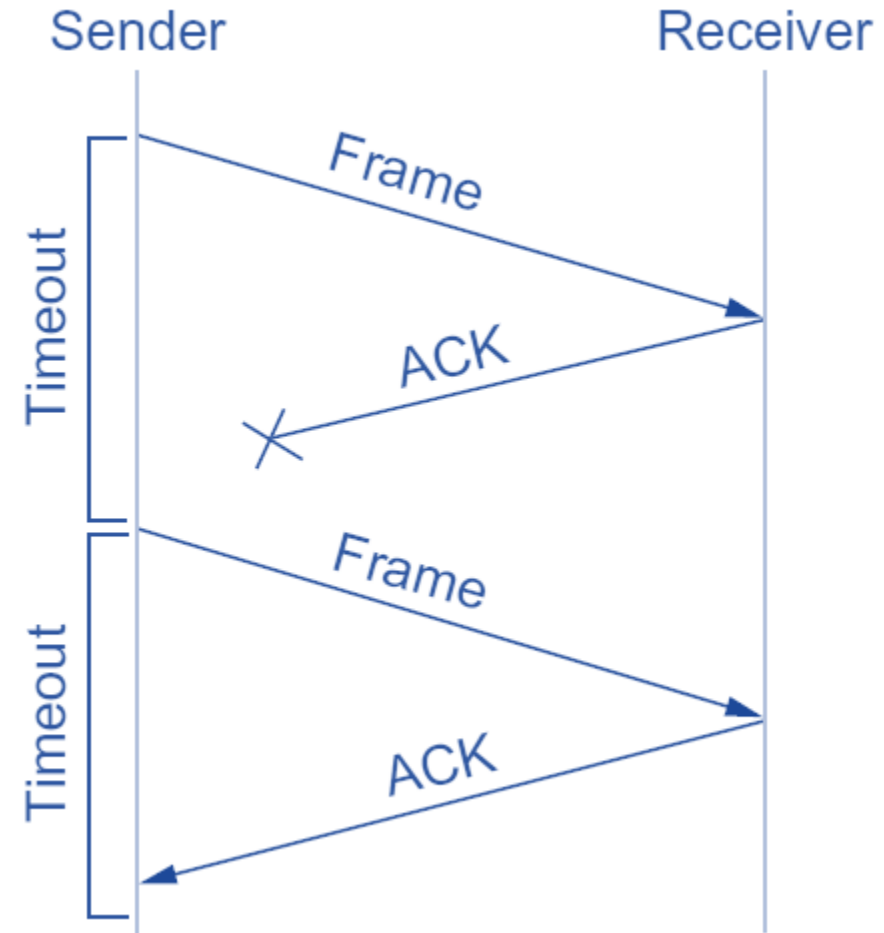
# Stop-and-Wait

- Case 2: ACK Loss
  - Sender time out
  - Sender retransmits



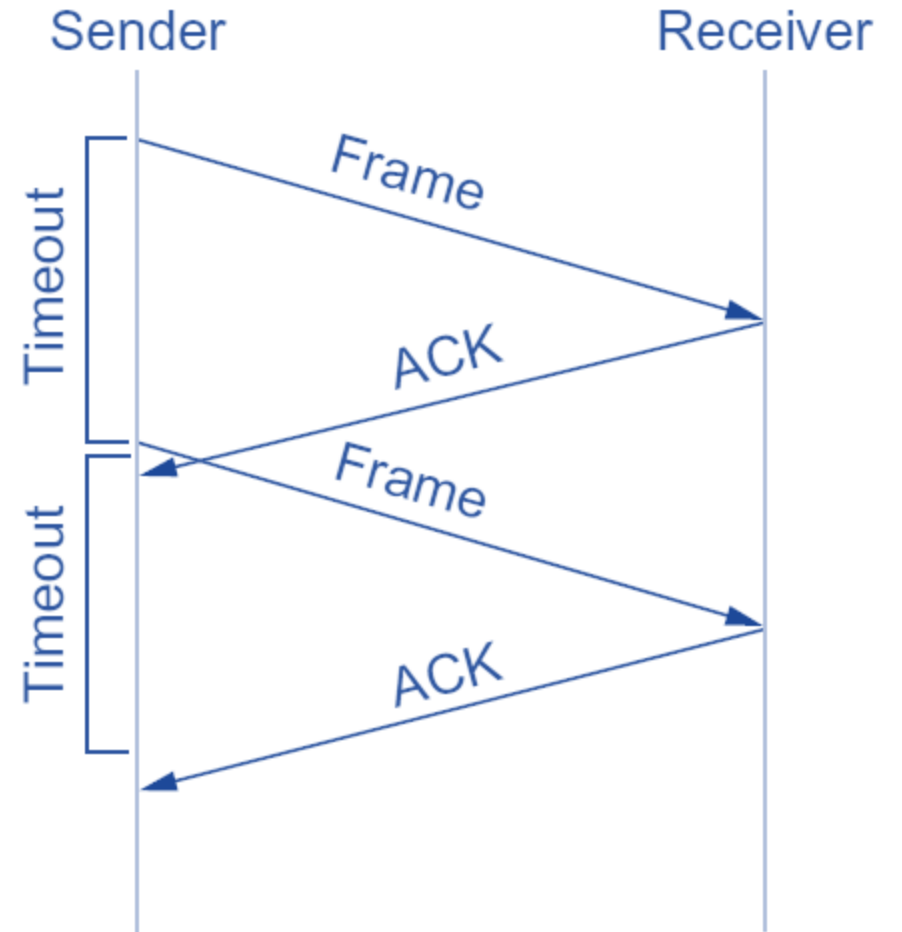
# Stop-and-Wait

- Case 2: ACK Loss
  - Sender time out
  - Sender retransmits
- Duplicated Frames
  - Solution: frame number



# Stop-and-Wait

- Case 3: ACK Late
  - Sender time out
  - Sender retransmits



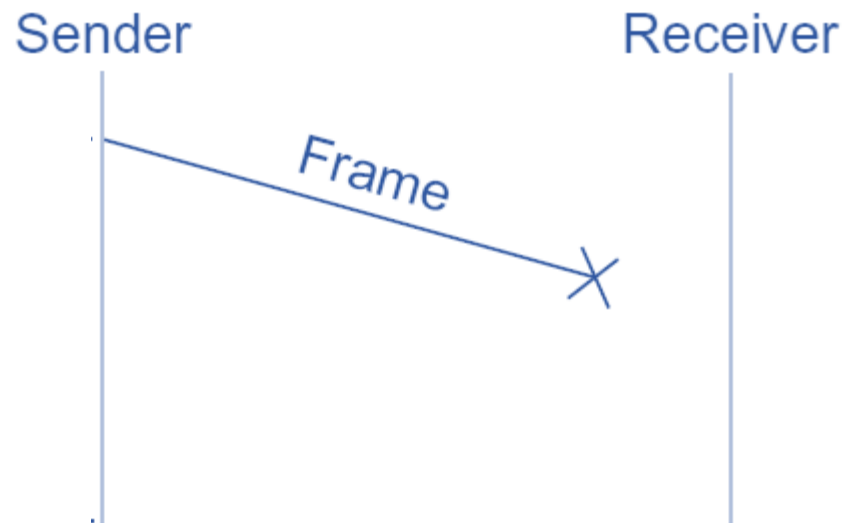


# Demo: Stop-and-Wait

- [https://www2.tkn.tu-berlin.de/teaching/rn/animations/gbn\\_sr/](https://www2.tkn.tu-berlin.de/teaching/rn/animations/gbn_sr/)

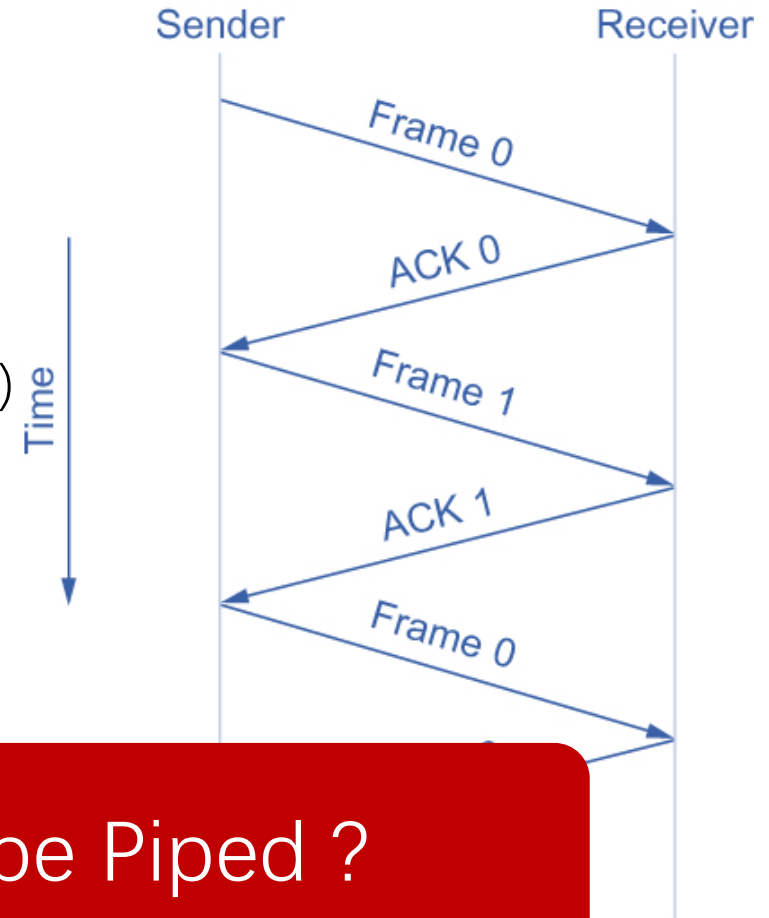
# How about NACK ?

- Negative ACK
  - Receiver sends NACK to indicate frame loss through sequence number
  - If frame loss is after sender's idle
    - The receiver has no way to notice the loss



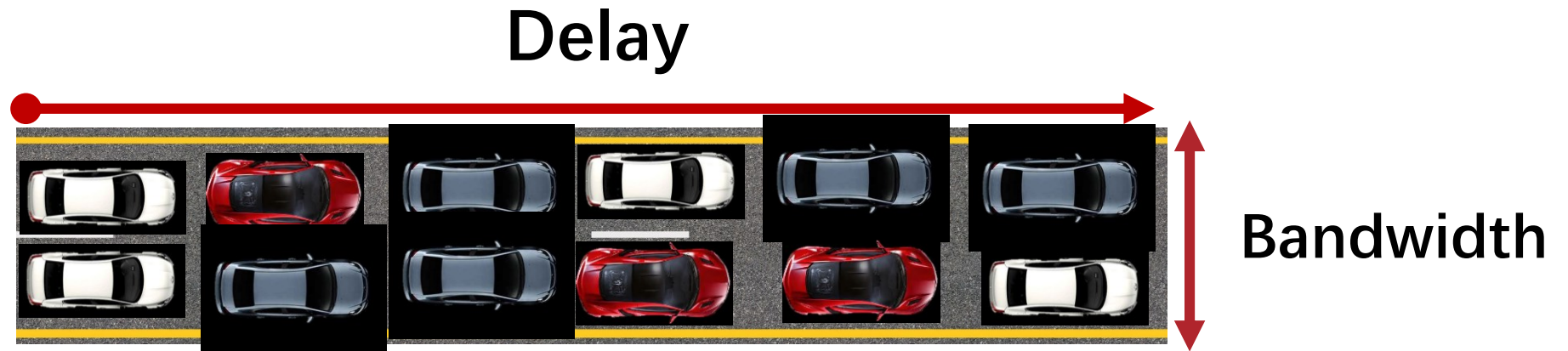
# Stop-and-Wait

- Efficiency Problem
  - 1.5Mbps bandwidth
  - 45ms RTT
  - 1KB frame
    - Effective Rate =  $1024 \cdot 8 / (1024 \cdot 8 / 1.5\text{Mbps} + 45\text{ms})$   
about 160kbps
- Solution
  - Pipeline



How Many Packets Can be Piped ?

# Delay $\times$ Bandwidth



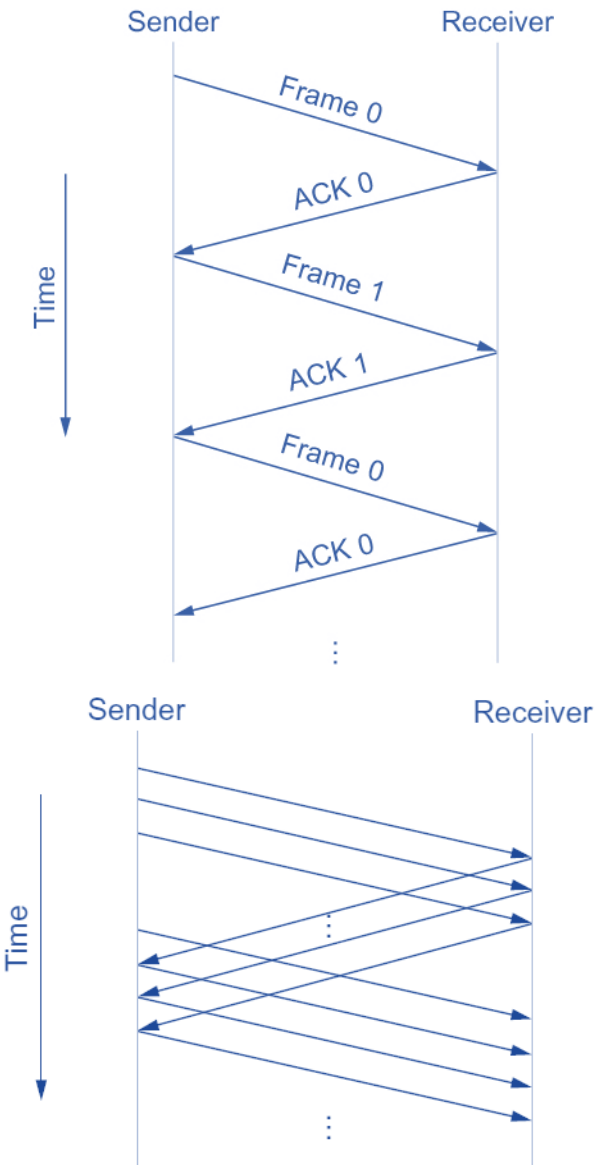
# Delay $\times$ Bandwidth

- Use to judge whether the link has been fully utilized

Table 1.1 Sample Delay $\times$ Bandwidth Products				
Link type	Bandwidth (typical)	One-way distance (typical)	Round-trip delay	RTT $\times$ Bandwidth
Dial-up	56 kbps	10 km	87 $\mu$ s	5 bits
Wireless LAN	54 Mbps	50 m	0.33 $\mu$ s	18 bits
Satellite	45 Mbps	35,000 km	230 ms	10 Mb
Cross-country fiber	10 Gbps	4,000 km	40 ms	400 Mb

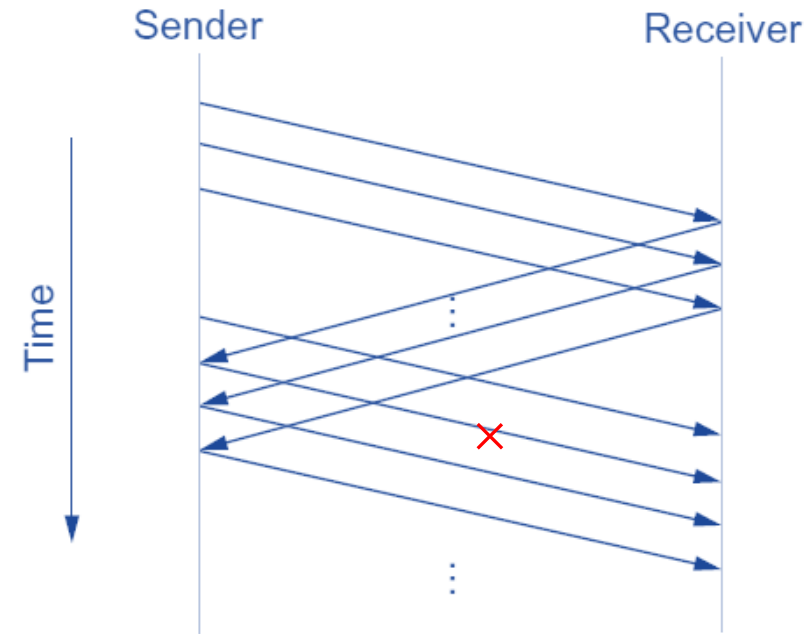
# Delay $\times$ Bandwidth

- Efficiency Problem
  - 1.5Mbps bandwidth
  - 45ms RTT
  - 1KB Frame
    - Effective Rate = 160kbps
- Solution
  - Pipeline
  - Full pipe situation:
    - $1.5\text{Mbps} \times 45\text{ms} / 1\text{KB} = 8$  frames in flight



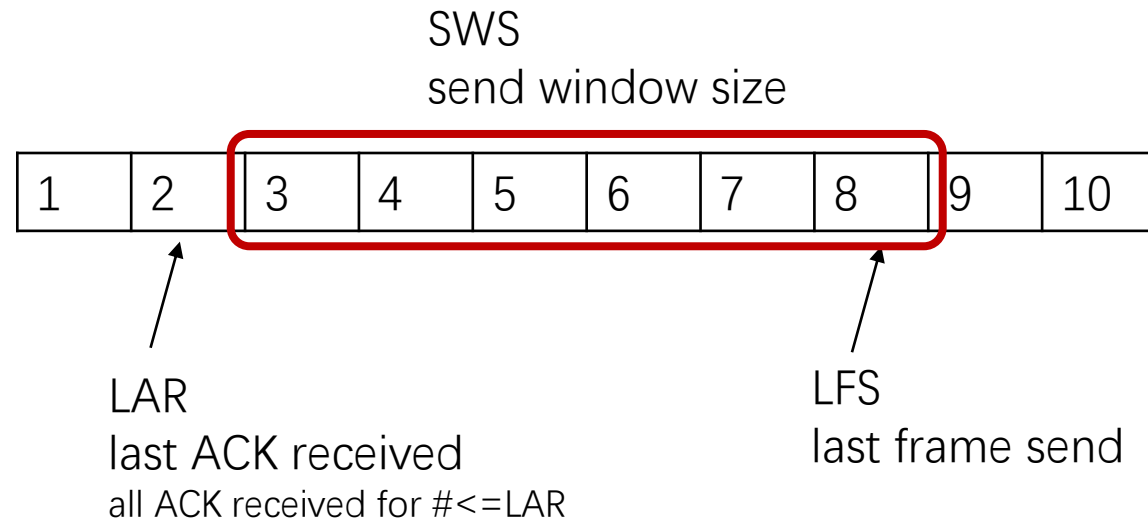
# Sliding Window – pipelined transmitting

- Sender Buffer
  - Retransmit
- Receiver Buffer
  - Handle out-of-order frames

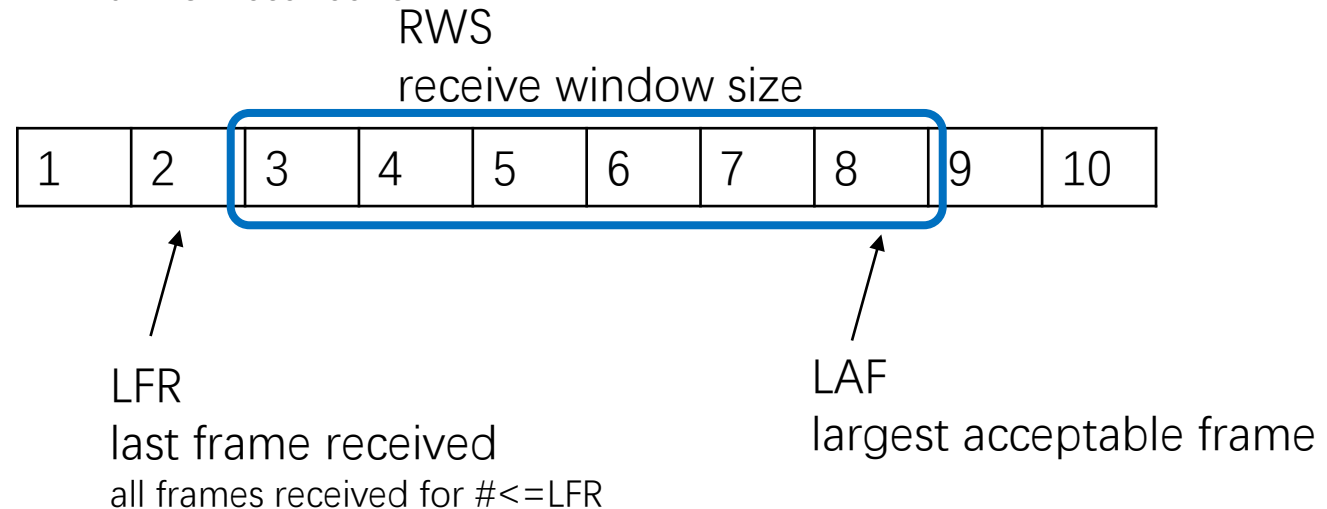


# Sliding Window

Sender Buffer:



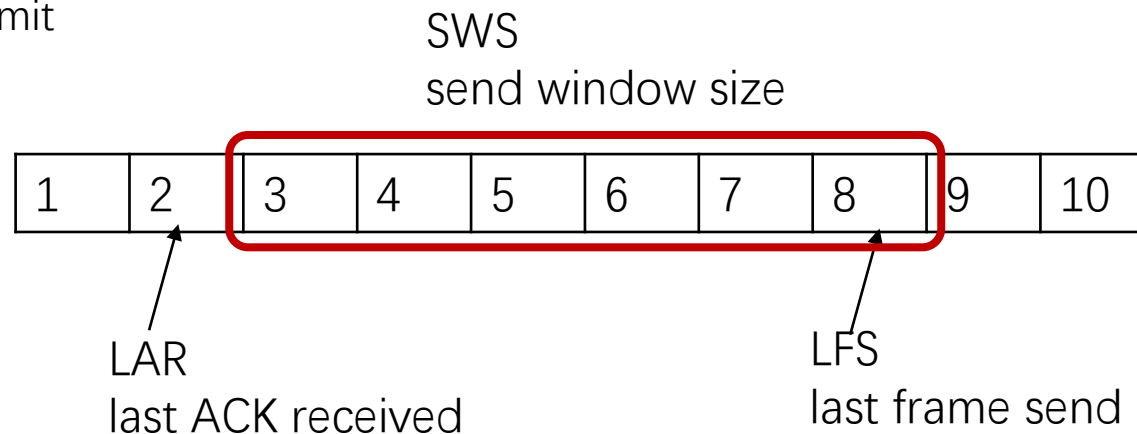
Receiver Buffer:





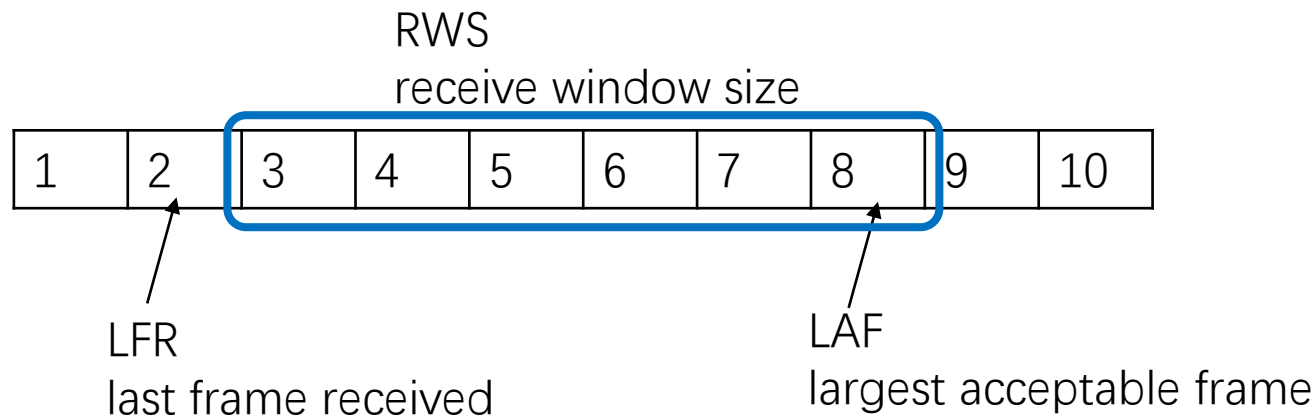
# Sliding Window

- Sender Protocol
  - Always maintain  $LFS - LAR \leq SWS$
  - When an ACK with sequence number #SeqNum arrives
    - If  $\#SeqNum \leq LAR$  or  $\#SeqNum > LFS$ 
      - No action
    - If  $LFR < \#SeqNum \leq LAF$ 
      - Move LAR to #SeqNum, increase LFS to send new packet
  - Associate a timer with each frame sender transmits
    - If timeout
      - Retransmit



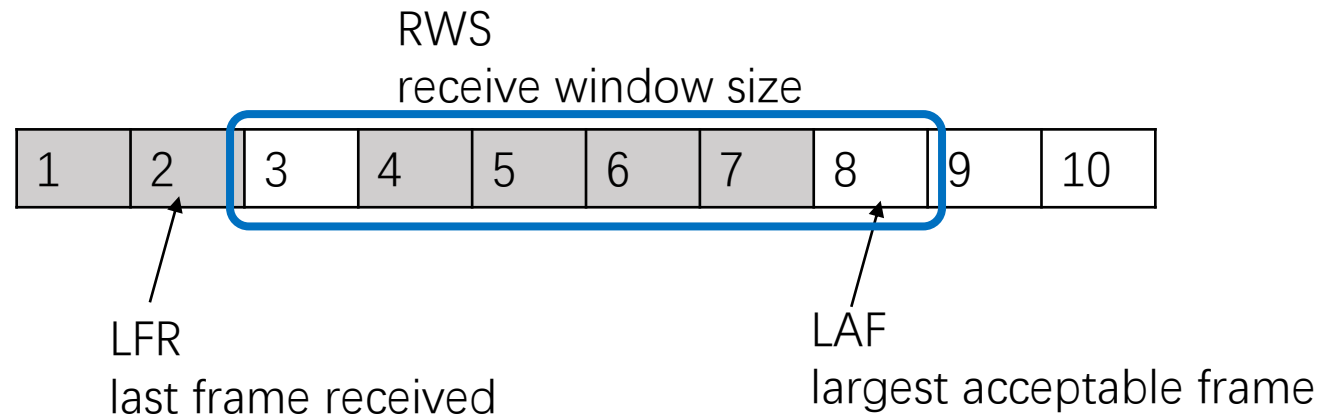
# Sliding Window

- Receiver Protocol
  - Always maintain  $LAF - LFR \leq RWS$
  - When a frame with sequence number  $\#SeqNum$  arrives
    - If  $\#SeqNum \leq LFR$  or  $\#SeqNum > LAF$ 
      - Discard frame, send accumulative ACK.
    - If  $LFR < \#SeqNum \leq LAF$ 
      - Accept frame, send accumulative ACK, modify LFT and LAF.



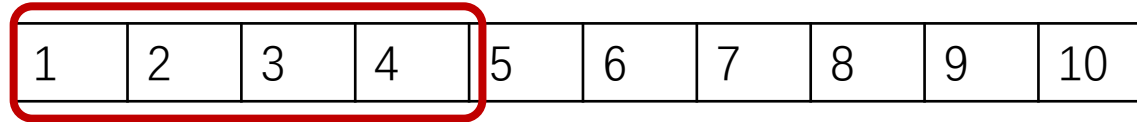
# Accumulative Ack

- If frame #3 is received
  - Ack #7, move LFR to 7, move LAF to 13
- If frame #8 is received
  - Ack #2



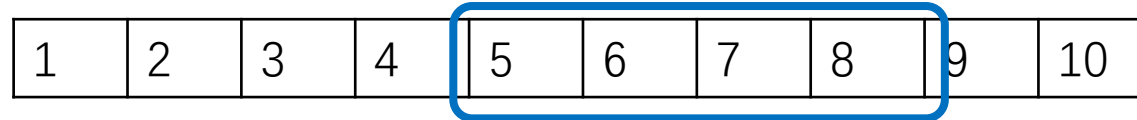
# ACK Loss

Sender



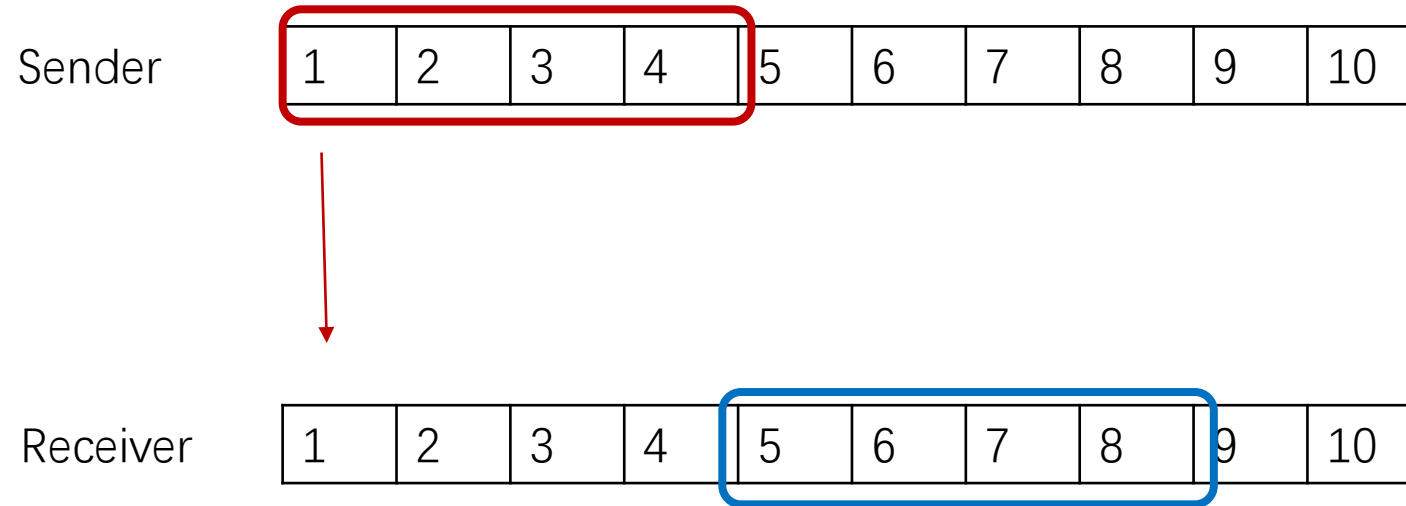
ack for 1,2,3,4, loss

Receiver

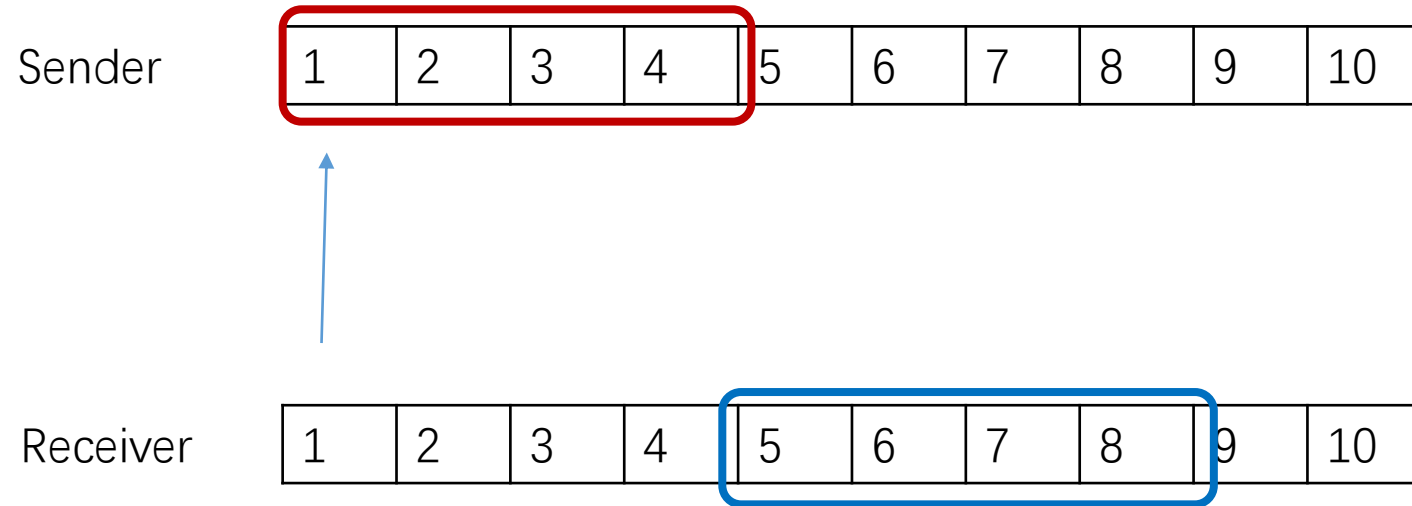


# ACK Loss

packet 1 time out, resend packet 1



# ACK Loss

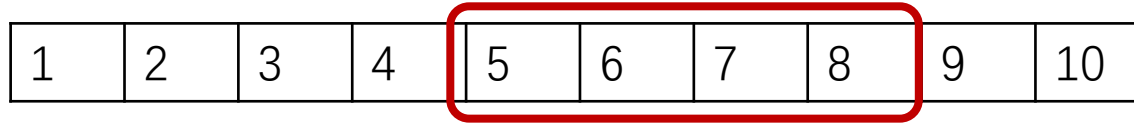


packet 1 rejected, send ack 4

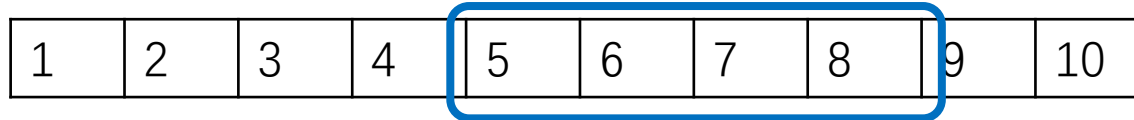
# ACK Loss

ack 4 received, slide window to 5

Sender



Receiver



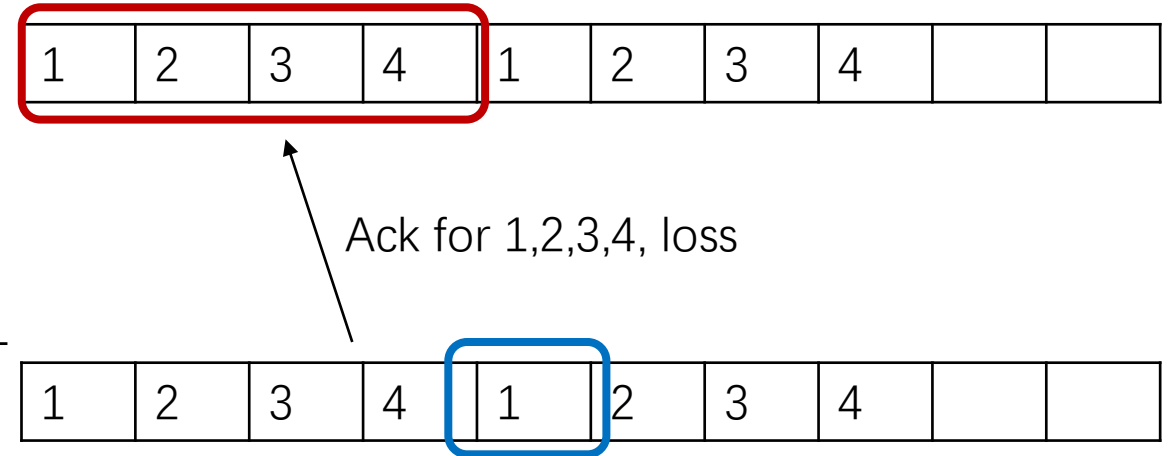
# Sliding Window

- Determine Window Size
  - Send Window Size: Pipeline depth
    - $\text{Delay} \times \text{Bandwidth}$
  - Receive Window Size: Flow control

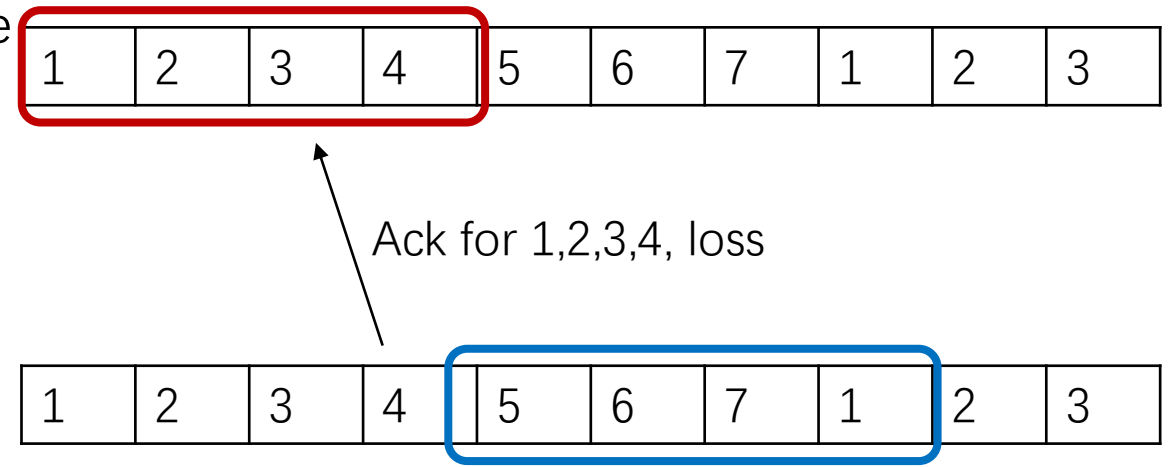


# Sliding Window

- Determine SeqNum Range
  - if Receive Window Size == 1
    - $\text{MaxSeqNum} \geq \text{Send Window Size} + 1$



- if Send Window Size == Receive Window Size
  - $\text{MaxSeqNum} \geq 2 * \text{Send Window Size}$



# Demo

- Sliding Window code in TCP  
`/net/ipv4/`
- Change Sliding Window Scheme
  - Show current congestion control scheme  
`cat /proc/sys/net/ipv4/tcp_congestion_control`
  - Show/change available congestion control scheme  
`sysctl net.ipv4.tcp_available_congestion_control[=XX]`
- [https://www2.tkn.tu-berlin.de/teaching/rn/animations/gbn\\_sr/](https://www2.tkn.tu-berlin.de/teaching/rn/animations/gbn_sr/)

# Reference

- Textbook 1.5.2
- Textbook 2.5